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Handling of Overlapping IPv6 Fragments

Abstract

The fragmentation and reassembly algorithm specified in the base IPv6 specification allows fragments to overlap. This document demonstrates the security issues associated with allowing overlapping fragments and updates the IPv6 specification to explicitly forbid overlapping fragments.

Status of This Memo

This document specifies an Internet standards track protocol for the Internet community, and requests discussion and suggestions for improvements. Please refer to the current edition of the "Internet Official Protocol Standards" (STD 1) for the standardization state and status of this protocol. Distribution of this memo is unlimited.

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1. Introduction

Fragmentation is used in IPv6 when the IPv6 packet will not fit inside the path MTU to its destination. When fragmentation is performed, an IPv6 node uses a fragment header, as specified in Section 4.5 of the IPv6 base specification [RFC2460], to break down the datagram into smaller fragments that will fit in the path MTU. The destination node receives these fragments and reassembles them. The algorithm specified for fragmentation in [RFC2460] does not prevent the fragments from overlapping, and this can lead to some security issues with firewalls [RFC4942]. This document explores the issues that can be caused by overlapping fragments and updates the IPv6 specification to explicitly forbid overlapping fragments.

1.1. Conventions Used in This Document

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

2. Overlapping Fragments

Commonly used firewalls use the algorithm specified in [RFC1858] to weed out malicious packets that try to overwrite parts of the transport-layer header in order to bypass inbound connection checks. [RFC1858] prevents an overlapping fragment attack on an upper-layer protocol (in this case, TCP) by recommending that packets with a fragment offset of 1 be dropped. While this works well for IPv4 fragments, it will not work for IPv6 fragments. This is because the fragmentable part of the IPv6 packet can contain extension headers before the TCP header, making this check less effective.

3. The Attack

This attack describes how a malicious node can bypass a firewall using overlapping fragments. Consider a sufficiently large IPv6 packet that needs to be fragmented.



Figure 1: Large IPv6 Packet

This packet is split into several fragments by the sender so that the packet can fit inside the path MTU. Let's say the packet is split into two fragments.

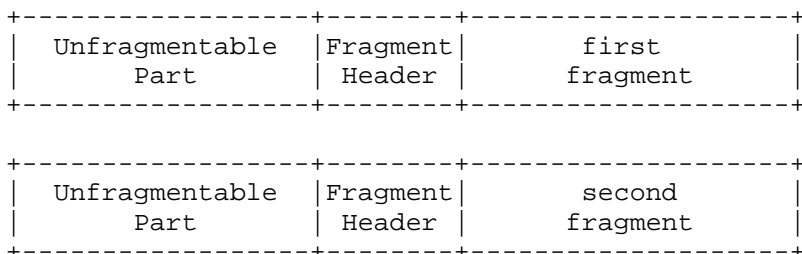


Figure 2: Fragmented IPv6 Packet

Consider the first fragment. Let's say it contains a destination options header (DOH) 80 octets long and is followed by a TCP header.

```

+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+<==FH
|NextHdr=DOH(60)|   Reserved   |   FragmentOffset = 0   |Res|1|
+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+
|                                     Identification=aaaabbbb                                     |
+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+<==DOH
|NextHdr=TCP(6) | HdrExtLen = 9 |                                     |
+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+
|                                     |                                     |
|                                     |                                     |
|                                     |                                     |
|                                     |                                     |
|                                     |                                     |
+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+<==TCP
|          Source Port          |          Destination Port          |
+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+
|          Sequence Number          |
+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+
|          Acknowledgment Number          |
+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+
| Offset| Reserved |U|A|P|R|S|F|          Window          |
+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+

```

Figure 3: First Fragment

The TCP header has the following values of the flags: S(YN)=1 and A(CK)=1. This may make an inspecting stateful firewall think that it is a response packet for a connection request initiated from the trusted side of the firewall. Hence, it will allow the fragment to pass. It will also allow the following fragments with the same Fragment Identification value in the fragment header to pass through.

A malicious node can form a second fragment with a TCP header that changes the flags and sets S(YN)=1 and A(CK)=0. This can change the packet on the receiving end to consider the packet as a connection request instead of a response. By doing this, the malicious node has bypassed the firewall's access control to initiate a connection request to a node protected by a firewall.

```
+-----+-----+-----+-----+-----+-----+-----+-----+<==FH  
|NextHdr=DOH(60)|   Reserved   |   FragmentOffset = 10   |Res|0|  
+-----+-----+-----+-----+-----+-----+-----+-----+  
|               Identification=aaaabbbb                |  
+-----+-----+-----+-----+-----+-----+-----+-----+<==TCP  
|      Source Port          |      Destination Port       |  
+-----+-----+-----+-----+-----+-----+-----+-----+  
|              Sequence Number                          |  
+-----+-----+-----+-----+-----+-----+-----+-----+  
|            Acknowledgment Number                      |  
+-----+-----+-----+-----+-----+-----+-----+-----+  
| Offset|Reserved|U|R|S|F|                               Window    |  
+-----+-----+-----+-----+-----+-----+-----+-----+
```

Figure 4: Second Fragment

Note that this attack is much more serious in IPv6 than in IPv4. In IPv4, the overlapping part of the TCP header does not include the source and destination ports. In IPv6, the attack can easily work to replace the source or destination port with an overlapping fragment.

4. Node Behavior

IPv6 nodes transmitting datagrams that need to be fragmented MUST NOT create overlapping fragments. When reassembling an IPv6 datagram, if one or more its constituent fragments is determined to be an overlapping fragment, the entire datagram (and any constituent fragments, including those not yet received) MUST be silently discarded.

Nodes MAY also provide mechanisms to track the reception of such packets, for instance, by implementing counters or alarms relating to these events.

5. Security Considerations

This document discusses an attack that can be used to bypass IPv6 firewalls using overlapping fragments. It recommends disallowing overlapping fragments in order to prevent this attack.

6. Acknowledgements

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7. References

7.1. Normative References

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- [RFC2460] Deering, S. and R. Hinden, "Internet Protocol, Version 6 (IPv6) Specification", RFC 2460, December 1998.

7.2. Informative References

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