

Network Working Group
Request for Comments: 4981
Category: Informational

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September 2007

Survey of Research towards Robust Peer-to-Peer Networks: Search Methods

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Abstract

The pace of research on peer-to-peer (P2P) networking in the last five years warrants a critical survey. P2P has the makings of a disruptive technology -- it can aggregate enormous storage and processing resources while minimizing entry and scaling costs.

Failures are common amongst massive numbers of distributed peers, though the impact of individual failures may be less than in conventional architectures. Thus, the key to realizing P2P's potential in applications other than casual file sharing is robustness.

P2P search methods are first couched within an overall P2P taxonomy. P2P indexes for simple key lookup are assessed, including those based on Plaxton trees, rings, tori, butterflies, de Bruijn graphs, and skip graphs. Similarly, P2P indexes for keyword lookup, information retrieval and data management are explored. Finally, early efforts to optimize range, multi-attribute, join, and aggregation queries over P2P indexes are reviewed. Insofar as they are available in the primary literature, robustness mechanisms and metrics are highlighted throughout. However, the low-level mechanisms that most affect robustness are not well isolated in the literature. Recommendations are given for future research.

Table of Contents

1. Introduction	3
1.1. Related Disciplines	6
1.2. Structured and Unstructured Routing	7
1.3. Indexes and Queries	9
2. Index Types	10
2.1. Local Index (Gnutella)	10
2.2. Central Index (Napster)	12
2.3. Distributed Index (Freenet)	13
3. Semantic Free Index	15
3.1. Origins	15
3.1.1. Plaxton, Rajaraman, and Richa (PRR)	15
3.1.2. Consistent Hashing	16
3.1.3. Scalable Distributed Data Structures (LH*)	16
3.2. Dependability	17
3.2.1. Static Dependability	17
3.2.2. Dynamic Dependability	18
3.2.3. Ephemeral or Stable Nodes -- $O(\log n)$ or $O(1)$ Hops	19
3.2.4. Simulation and Proof	20
3.3. Latency	21
3.3.1. Hop Count and the $O(1)$ -Hop DHTs	21
3.3.2. Proximity and the $O(\log n)$ -Hop DHTs	22
3.4. Multicasting	23
3.4.1. Multicasting vs. Broadcasting	23
3.4.2. Motivation for DHT-based Multicasting	23
3.4.3. Design Issues	24
3.5. Routing Geometries	25
3.5.1. Plaxton Trees (Pastry, Tapestry)	25
3.5.2. Rings (Chord, DKS)	27
3.5.3. Tori (CAN)	28
3.5.4. Butterflies (Viceroy)	29
3.5.5. de Bruijn (D2B, Koorde, Distance Halving, ODRI)	30
3.5.6. Skip Graphs	32
4. Semantic Index	33
4.1. Keyword Lookup	34
4.1.1. Gnutella Enhancements	36
4.1.2. Partition-by-Document, Partition-by-Keyword	38
4.1.3. Partial Search, Exhaustive Search	39
4.2. Information Retrieval	39
4.2.1. Vector Model (PlanetP, FASD, eSearch)	41
4.2.2. Latent Semantic Indexing (pSearch)	43
4.2.3. Small Worlds	43
5. Queries	44
5.1. Range Queries	45
5.2. Multi-Attribute Queries	48
5.3. Join Queries	50

5.4. Aggregation Queries	50
6. Security Considerations	52
7. Conclusions	52
8. Acknowledgments	53
9. References	54
9.1. Informative References	54

1. Introduction

Peer-to-peer (P2P) networks are those that exhibit three characteristics: self-organization, symmetric communication, and distributed control [1]. A self-organizing P2P network "automatically adapts to the arrival, departure and failure of nodes" [2]. Communication is symmetric in that peers act as both clients and servers. It has no centralized directory or control point. USENET servers and BGP peers have these traits [3] but the emphasis here is on the flurry of research since 2000. Leading examples include Gnutella [4], Freenet [5], Pastry [2], Tapestry [6], Chord [7], the Content Addressable Network (CAN) [8], pSearch [9], and Edutella [10]. Some have suggested that peers are inherently unreliable [11]. Others have assumed well-connected, stable peers [12].

This critical survey of P2P academic literature is warranted, given the intensity of recent research. At the time of writing, one research database lists over 5,800 P2P publications [13]. One vendor surveyed P2P products and deployments [14]. There is also a tutorial survey of leading P2P systems [15]. DePaoli and Mariani recently reviewed the dependability of some early P2P systems at a high level [16]. The need for a critical survey was flagged in the peer-to-peer research group of the Internet Research Task Force (IRTF) [17].

P2P is potentially a disruptive technology with numerous applications, but this potential will not be realized unless it is demonstrated to be robust. A massively distributed search technique may yield numerous practical benefits for applications [18]. A P2P system has potential to be more dependable than architectures relying on a small number of centralized servers. It has potential to evolve better from small configurations -- the capital outlays for high performance servers can be reduced and spread over time if a P2P assembly of general purpose nodes is used. A similar argument motivated the deployment of distributed databases -- one thousand, off-the-shelf PC processors are more powerful and much less expensive than a large mainframe computer [19]. Storage and processing can be aggregated to achieve massive scale. Wasteful partitioning between servers or clusters can be avoided. As Gedik and Liu put it, if P2P is to find its way into applications other than casual file sharing, then reliability needs to be addressed [20].

The taxonomy of Figure 1 divides the entire body of P2P research literature along four lines: search, storage, security, and applications. This survey concentrates on search aspects. A P2P search network consists of an underlying index (Sections 2 to 4) and queries that propagate over that index (Section 5).

Search [18, 21-29]

Semantic-Free Indexes [2, 6, 7, 30-52]

- Plaxton Trees

- Rings

- Tori

- Butterflies

- de Bruijn Graphs

- Skip Graphs

Semantic Indexes [4, 53-71]

- Keyword Lookup

- Peer Information Retrieval

- Peer Data Management

Queries [20, 22, 23, 25, 32, 38, 41, 56, 72-100]

- Range Queries

- Multi-Attribute Queries

- Join Queries

- Aggregation Queries

- Continuous Queries

- Recursive Queries

- Adaptive Queries

Storage

Consistency & Replication [101-112]

- Eventual consistency

- Trade-offs

Distribution [39, 42, 90, 92, 113-131]

- Epidemics, Bloom Filters

Fault Tolerance [40, 105, 132-139]

- Erasures Coding

- Byzantine Agreement

Locality [24, 43, 47, 140-160]

Load Balancing [37, 86, 100, 107, 151, 161-171]

- Security
 - Character [172-182]
 - Identity
 - Reputation and Trust
 - Incentives
 - Goals [25, 27, 71, 183-197]
 - Availability
 - Authenticity
 - Anonymity
 - Access Control
 - Fair Trading
- Applications [1, 198-200]
 - Memory [32, 90, 142, 201-222]
 - File Systems
 - Web
 - Content Delivery Networks
 - Directories
 - Service Discovery
 - Publish / Subscribe ...
- Intelligence [223-228]
 - GRID
 - Security...
- Communication [12, 92, 119, 229-247]
 - Multicasting
 - Streaming Media
 - Mobility
 - Sensors...

Figure 1: Classification of P2P Research Literature

This survey is concerned with two questions. The first, "How do P2P search networks work?" This foundation is important given the pace and breadth of P2P research in the last five years. In Section 2, we classify indexes as local, centralized and distributed. Since distributed indexes are becoming dominant, they are given closer attention in Sections 3 and 4. Section 3 compares distributed P2P indexes for simple key lookup; in particular, their origins (Section 3.1), dependability (Section 3.2), latency (Section 3.3), and their support for multicast (Section 3.4). It classifies those indexes according to their routing geometry (Section 3.5) -- Plaxton trees, rings, tori, butterflies, de Bruijn graphs and skip graphs. Section 4 reviews distributed P2P indexes supporting keyword lookup (Section 4.1) and information retrieval (Section 4.2). Section 5 probes the embryonic research on P2P queries; in particular, range queries (Section 5.1), multi-attribute queries (Section 5.2), join queries (Section 5.3), and aggregation queries (Section 5.4).

The second question, "How robust are P2P search networks?" Insofar as it is available in the research literature, we tease out the robustness mechanisms and metrics throughout Sections 2 to 5. Unfortunately, robustness is often more sensitive to low-level design choices than it is to the broad P2P index structure, yet these underlying design choices are seldom isolated in the primary literature [248]. Furthermore, there has been little consensus on P2P robustness metrics (Section 3.2). Section 8 gives recommendations to address these important gaps.

1.1. Related Disciplines

Peer-to-peer research draws upon numerous distributed systems disciplines. Networking researchers will recognize familiar issues of naming, routing, and congestion control. P2P designs need to address routing and security issues across network region boundaries [152]. Networking research has traditionally been host-centric. The Web's Universal Resource Identifiers are naturally tied to specific hosts, making object mobility a challenge [216].

P2P work is data-centric [249]. P2P systems for dynamic object location and routing have borrowed heavily from the distributed systems corpus. Some have used replication, erasure codes, and Byzantine agreement [111]. Others have used epidemics for durable peer group communication [39].

Similarly, P2P research is set to benefit from database research [250]. Database researchers will recognize the need to reapply Codd's principle of physical data independence, that is, to decouple data indexes from the applications that use the data [23]. It was the invention of appropriate indexing mechanisms and query optimizations that enabled data independence. Database indexes like B+ trees have an analog in P2P's distributed hash tables (DHTs). Wide-area, P2P query optimization is a ripe, but challenging, area for innovation.

More flexible distribution of objects comes with increased security risks. There are opportunities for security researchers to deliver new methods for availability, file authenticity, anonymity, and access control [25]. Proactive and reactive mechanisms are needed to deal with large numbers of autonomous, distributed peers. To build robust systems from cooperating but self-interested peers, issues of identity, reputation, trust, and incentives need to be tackled. Although it is beyond the scope of this paper, robustness against malicious attacks also ought to be addressed [195].

Possibly the largest portion of P2P research has majored on basic routing structures [18], where research on algorithms comes to the

fore. Should the overlay be "structured" or "unstructured"? Are the two approaches competing or complementary? Comparisons of the "structured" approaches (hypercubes, rings, toroids, butterflies, de Bruijn, and skip graphs) have weighed the amount of routing state per peer and the number of links per peer against overlay hop counts. While "unstructured" overlays initially used blind flooding and random walks, overheads usually trigger some structure, for example, super-peers and clusters.

P2P applications rely on cooperation between these disciplines. Applications have included file sharing, directories, content delivery networks, email, distributed computation, publish-subscribe middleware, multicasting, and distributed authentication. Which applications will be suited to which structures? Are there adaptable mechanisms that can decouple applications from the underlying data structures? What are the criteria for selection of applications amenable to a P2P design [1]?

Robustness is emphasized throughout the survey. We are particularly interested in two aspects. The first, dependability, was a leading design goal for the original Internet [251]. It deserves the same status in P2P. The measures of dependability are well established: reliability, a measure of the mean-time-to-failure (MTTF); availability, a measure of both the MTTF and the mean-time-to-repair (MTTR); maintainability; and safety [252]. The second aspect is the ability to accommodate variation in outcome, which one could call adaptability. Its measures have yet to be defined. In the context of the Internet, it was only recently acknowledged as a first-class requirement [253]. In P2P, it means planning for the tussles over resources and identity. It means handling different kinds of queries and accommodating changeable application requirements with minimal intervention. It means "organic scaling" [22], whereby the system grows gracefully, without a priori data center costs or architectural breakpoints.

In the following section, we discuss one notable omission from the taxonomy of P2P networking in Figure 1 -- routing.

1.2. Structured and Unstructured Routing

P2P routing algorithms have been classified as "structured" or "unstructured". Peers in unstructured overlay networks join by connecting to any existing peers [254]. In structured overlays, the identifier of the joining peer determines the set of peers that it connects to [254]. Early instantiations of Gnutella were unstructured -- keyword queries were flooded widely [255]. Napster [256] had decentralized content and a centralized index, so it only partially satisfies the distributed control criteria for P2P systems.

Early structured algorithms included Plaxton, Rajaraman and Richa (PRR) [30], Pastry [2], Tapestry [31], Chord [7], and the Content Addressable Network [8]. Mishchke and Stiller recently classified P2P systems by the presence or absence of structure in routing tables and network topology [257].

Some have cast unstructured and structured algorithms as competing alternatives. Unstructured approaches have been called "first generation", implicitly inferior to the "second generation" structured algorithms [2, 31]. When generic key lookups are required, these structured, key-based routing schemes can guarantee location of a target within a bounded number of hops [23]. The broadcasting unstructured approaches, however, may have large routing costs, or fail to find available content [22]. Despite the apparent advantages of structured P2P, several research groups are still pursuing unstructured P2P.

There have been two main criticisms of structured systems [61]. The first relates to peer transience, which in turn, affects robustness. Chawathe, et al. opined that highly transient peers are not well supported by DHTs [61]. P2P systems often exhibit "churn", with peers continually arriving and departing. One objection to concerns about highly transient peers is that many applications use peers in well-connected parts of the network. The Tapestry authors analyzed the impact of churn in a network of 1000 nodes [31]. Others opined that it is possible to maintain a robust DHT at relatively low cost [258]. Very few papers have quantitatively compared the resilience of structured systems. Loguinov, Kumar, et al. claimed that there were only two such works [24, 36].

The second criticism of structured systems is that they do not support keyword searches and complex queries as well as unstructured systems. Given the current file-sharing deployments, keyword searches seem more important than exact-match key searches in the short term. Paraphrased, "most queries are for hay, not needles" [61].

More recently, some have justifiably seen unstructured and structured proposals as complementary, and have devised hybrid models [259]. Their starting point was the observation that unstructured flooding or random walks are inefficient for data that is not highly replicated across the P2P network. Structured graphs can find keys efficiently, irrespective of replication. Castro, et al. proposed Structella, a hybrid of Gnutella built on top of Pastry [259]. Another design used structured search for rare items and unstructured search for massively replicated items [54].

However, the "structured versus unstructured routing" taxonomy is becoming less useful, for two reasons. Firstly, most "unstructured" proposals have evolved and incorporated structure. Consider the classic "unstructured" system, Gnutella [4]. For scalability, its peers are either ultrapeers or leaf nodes. This hierarchy is augmented with a query routing protocol whereby ultrapeers receive a hashed summary of the resource names available at leaf nodes. Between ultrapeers, simple query broadcast is still used, though methods to reduce the query load here have been considered [260]. Secondly, there are emerging schema-based P2P designs [59], with super-node hierarchies and structure within documents. These are quite distinct from the structured DHT proposals.

1.3. Indexes and Queries

Given that most, if not all, P2P designs today assume some structure, a more instructive taxonomy would describe the structure. In this survey, we use a database taxonomy in lieu of the networking taxonomy, as suggested by Hellerstein, Cooper, and Garcia-Molina [23, 261]. The structure is determined by the type of index (Sections 2, 3, and 4). Queries feature in lieu of routing (Section 5). The DHT algorithms implement a "semantic-free index" [216]. They are oblivious of whether keys represent document titles, meta-data, or text. Gnutella-like and schema-based proposals have a "semantic index".

Index engineering is at the heart of P2P search methods. It captures a broad range of P2P issues, as demonstrated by the Search/Index Links model [261]. As Manber put it, "the most important of the tools for information retrieval is the index -- a collection of terms with pointers to places where information about documents can be found" [262]. Sen and Wang noted that a "P2P network" usually consists of connections between hosts for application-layer signaling, rather than for the data transfer itself [263]. Similarly, we concentrate on the "signaled" indexes and queries.

Our focus here is the dependability and adaptability of the search network. Static dependability is a measure of how well queries route around failures in a network that is normally fault-free. Dynamic dependability gives an indication of query success when nodes and data are continually joining and leaving the P2P system. An adaptable index accommodates change in the data and query distribution. It enables data independence, in that it facilitates changes to the data layout without requiring changes to the applications that use the data [23]. An adaptable P2P system can support rich queries for a wide range of applications. Some applications benefit from simple, semantic-free key lookups [264]. Others require more complex, Structured Query Language (SQL)-like

queries to find documents with multiple keywords, or to aggregate or join query results from distributed relations [22].

2. Index Types

A P2P index can be local, centralized, or distributed. With a local index, a peer only keeps the references to its own data, and does not receive references for data at other nodes. The very early Gnutella design epitomized the local index (Section 2.1). In a centralized index, a single server keeps references to data on many peers. The classic example is Napster (Section 2.2). With distributed indexes, pointers towards the target reside at several nodes. One very early example is Freenet (Section 2.3). Distributed indexes are used in most P2P designs nowadays -- they dominate this survey.

P2P indexes can also be classified as non-forwarding and forwarding. When queries are guided by a non-forwarding index, they jump to the node containing the target data in a single hop. There have been semantic and semantic-free one-hop schemes [138, 265, 266]. Where scalability to a massive number of peers is required, these schemes have been extended to two hops [267, 268]. More common are the forwarding P2Ps, where the number of hops varies with the total number of peers, often logarithmically. The related trade-offs between routing state, lookup latency, update bandwidth, and peer churn are critical to total system dependability.

2.1. Local Index (Gnutella)

P2Ps with a purely local data index are becoming rare. In such designs, peers flood queries widely and only index their own content. They enable rich queries - the search is not limited to a simple key lookup. However, they also generate a large volume of query traffic with no guarantee that a match will be found, even if it does exist on the network. For example, to find potential peers on the early instantiations of Gnutella, 'ping' messages were broadcast over the P2P network and the 'pong' responses were used to build the node index. Then, small 'query' messages, each with a list of keywords, are broadcast to peers that respond with matching filenames [4].

There have been numerous attempts to improve the scalability of local-index P2P networks. Gnutella uses fixed time-to-live (TTL) rings, where the query's TTL is set less than 7-10 hops [4]. Small TTLs reduce the network traffic and the load on peers, but also reduce the chances of a successful query hit. One paper reported, perhaps a little too bluntly, that the fixed "TTL-based mechanism does not work" [67]. To address this TTL selection problem, they proposed an expanding ring, known elsewhere as iterative deepening [29]. It uses successively larger TTL counters until there is a

match. The flooding, ring, and expanding ring methods all increase network load with duplicated query messages. A random walk, whereby an unduplicated query wanders about the network, does indeed reduce the network load but massively increases the search latency. One solution is to replicate the query k times at each peer. Called random k -walkers, this technique can be coupled with TTL limits, or periodic checks with the query originator, to cap the query load [67]. Adamic, Lukose, et al. suggested that the random walk searches be directed to nodes with a higher degree, that is, with larger numbers of inter-peer connections [269]. They assumed that higher-degree peers are also capable of higher query throughputs. However, without some balancing design rule, such peers would be swamped with the entire P2P signaling traffic. In addition to the above approaches, there is the 'directed breadth-first' algorithm [29]. It forwards queries within a subset of peers selected according to heuristics on previous performance, like the number of successful query results. Another algorithm, called probabilistic flooding, has been modeled using percolation theory [270].

Several measurement studies have investigated locally indexed P2Ps. Jovanovic noted Gnutella's power law behaviour [70]. Sen and Wang compared the performance of Gnutella, Fasttrack [271], and Direct Connect [263, 272, 273]. At the time, only Gnutella used local data indexes. All three schemes now use distributed data indexes, with hierarchy in the form of Ultrapeers (Gnutella), Super-Nodes (FastTrack), and Hubs (Direct Connect). It was found that a very small percentage of peers have a very high degree and that the total system dependability is at the mercy of such peers. While peer up-time and bandwidth were heavy-tailed, they did not fit well with the Zipf distribution. Fortunately for Internet Service Providers, measures aggregated by IP prefix and Autonomous System (AS) were more stable than for individual IP addresses. A study of University of Washington traffic found that Gnutella and Kazaa together contributed 43% of the university's total TCP traffic [274]. They also reported a heavy-tailed distribution, with 600 external peers (out of 281,026) delivering 26% of Kazaa bytes to internal peers. Furthermore, objects retrieved from the P2P network were typically three orders of magnitude larger than Web objects -- 300 objects contributed to almost half the total outbound Kazaa bandwidth. Others reported Gnutella's topology mismatch, whereby only 2-5% of P2P connections link peers in the same Autonomous System (AS), despite over 40% of peers being in the top 10 ASs [65]. Together these studies underscore the significance of multimedia sharing applications. They motivate interesting caching and locality solutions to the topology mismatch problem.

These same studies bear out one main dependability lesson: total system dependability may be sensitive to the dependability of high-

degree peers. The designers of Scamp translated this observation to the design heuristic, "have the degree of each node be of nearly equal size" [153]. They analyzed a system of N peers, with mean degree $c \cdot \log(n)$, where link failures occur independently with probability e . If $d > 0$ is fixed and $c > (1+d)/(-\log(e))$, then the probability of graph disconnection goes to zero as $N \rightarrow \infty$. Otherwise, if $c < (1+d)/(-\log(e))$, then the probability of disconnection goes to one as $N \rightarrow \infty$. They presented a localizer, which finds approximate minima to a global function of peer degree and arbitrary link costs using only local information. The Scamp overlay construction algorithms could support any of the flooding and walking routing schemes above, or other epidemic and multicasting schemes for that matter. Resilience to high churn rates was identified for future study.

2.2. Central Index (Napster)

Centralized schemes like Napster [256] are significant because they were the first to demonstrate the P2P scalability that comes from separating the data index from the data itself. Ultimately, 36 million Napster users lost their service not because of technical failure, but because the single administration was vulnerable to the legal challenges of record companies [275].

There has since been little research on P2P systems with central data indexes. Such systems have also been called 'hybrid' since the index is centralized but the data is distributed. Yang and Garcia-Molina devised a four-way classification of hybrid systems [276]: unchained servers, where users whose index is on one server do not see other servers' indexes; chained servers, where the server that receives a query forwards it to a list of servers if it does not own the index itself; full replication, where all centralized servers keep a complete index of all available metadata; and hashing, where keywords are hashed to the server where the associated inverted list is kept. The unchained architecture was used by Napster, but it has the disadvantage that users do not see all indexed data in the system. Strictly speaking, the other three options illustrate the distributed data index, not the central index. The chained architecture was recommended as the optimum for the music-swapping application at the time. The methods by which clients update the central index were classified as batch or incremental, with the optimum determined by the query-to-login ratio. Measurements were derived from a clone of Napster called OpenNap[277]. Another study of live Napster data reported wide variation in the availability of peers, a general unwillingness to share files (20-40% of peers share few or no files), and a common understatement of available bandwidth so as to discourage other peers from sharing one's link [202].

Influenced by Napster's early demise, the P2P research community may have prematurely turned its back on centralized architectures. Chawathe, Ratnasamy, et al. opined that Google and Yahoo demonstrate the viability of a centralized index. They argued that "the real barriers to Napster-like designs are not technical but legal and financial" [61]. Even this view may be a little too harsh on the centralized architectures -- it implies that they always have an up-front capital hurdle that is steeper than for distributed architectures. The closer one looks at scalable 'centralized' architectures, the less the distinction with 'distributed' architectures seems to matter. For example, it is clear that Google's designers consider Google a distributed, not centralized, file system [278]. Google demonstrates the scale and performance possible on commodity hardware, but still has a centralized master that is critical to the operation of each Google cluster. Time may prove that the value of emerging P2P networks, regardless of the centralized-versus-distributed classification, is that they smooth the capital outlays and remove the single points of failure across the spectra of scale and geographic distribution.

2.3. Distributed Index (Freenet)

An important early P2P proposal for a distributed index was Freenet [5, 71, 279]. While its primary emphasis was the anonymity of peers, it did introduce a novel indexing scheme. Files are identified by low-level "content-hash" keys and by "secure signed-subspace" keys, which ensure that only a file owner can write to a file while anyone can read from it. To find a file, the requesting peer first checks its local table for the node with keys closest to the target. When that node receives the query, it too checks for either a match or another node with keys close to the target. Eventually, the query either finds the target or exceeds time-to-live (TTL) limits. The query response traverses the successful query path in reverse, depositing a new routing table entry (the requested key and the data holder) at each peer. The insert message similarly steps towards the target node, updating routing table entries as it goes, and finally stores the file there. Whereas early versions of Gnutella used breadth-first flooding, Freenet uses a more economic depth-first search [280].

An initial assessment has been done of Freenet's robustness. It was shown that in a network of 1000 nodes, the median query path length stayed under 20 hops for a failure of 30% of nodes. While the Freenet designers considered this as evidence that the system is "surprisingly robust against quite large failures" [71], the same datapoint may well be outside meaningful operating bounds. How many applications are useful when the first quartile of queries have path lengths of several hundred hops in a network of only 1000 nodes, per

Figure 4 of [71]? To date, there has been no analysis of Freenet's dynamic robustness. For example, how does it perform when nodes are continually arriving and departing?

There have been both criticisms and extensions of the early Freenet work. Gnutella proponents acknowledged the merit in Freenet's avoidance of query broadcasting [281]. However, they are critical on two counts: the exact file name is needed to construct a query; and exactly one match is returned for each query. P2P designs using DHTs, per Section 3, share similar characteristics -- a precise query yields a precise response. The similarity is not surprising since Freenet also uses a hash function to generate keys. However, the query routing used in the DHTs has firmer theoretical foundations. Another difference with DHTs is that Freenet will take time, when a new node joins the network, to build an index that facilitates efficient query routing. By the inventor's own admission, this is damaging for a user's first impressions [282]. It was proposed to download a copy of routing tables from seed nodes at startup, even though the new node might be far from the seed node. Freenet's slow startup motivated Mache, Gilbert, et al. to amend the overlay after failed requests and to place additional index entries on successful requests -- they claim almost an order of magnitude reduction in average query path length [280]. Clarke also highlighted the lack of locality or bandwidth information available for efficient query routing decisions [282]. He proposed that each node gather response times, connection times, and proportion of successful requests for each entry in the query routing table. When searching for a key that is not in its own routing table, it was proposed to estimate response times from the routing metrics for the nearest known keys and consequently choose the node that can retrieve the data fastest. The response time heuristic assumed that nodes close in the key space have similar response times. This assumption stemmed from early deployment observations that Freenet peers seemed to specialize in parts of the keyspace -- it has not been justified analytically. Kronfol drew attention to Freenet's inability to do keyword searches [283]. He suggested that peers cache lists of weighted keywords in order to route queries to documents, using Term Frequency Inverse Document Frequency (TFIDF) measures and inverted indexes (Section 4.2.1). With these methods, a peer can route queries for simple keyword lists or more complicated conjunctions and disjunctions of keywords. Robustness analysis and simulation of Kronfol's proposal remain open.

The vast majority of P2P proposals in following sections rely on a distributed index.

3. Semantic Free Index

Many of today's distributed network indexes are semantic. The semantic index is human-readable. For example, it might associate information with other keywords, a document, a database key, or even an administrative domain. It makes it easy to associate objects with particular network providers, companies, or organizations, as evidenced in the Domain Name System (DNS). However, it can also trigger legal tussles and frustrate content replication and migration [216].

Distributed Hash Tables (DHTs) have been proposed to provide semantic-free, data-centric references. DHTs enable one to find an object's persistent key in a very large, changing set of hosts. They are typically designed for [23]:

- a) low degree. If each node keeps routing information for only a small number of other nodes, the impact of high node arrival and departure rates is contained;
- b) low hop count. The hops and delay introduced by the extra indirection are minimized;
- c) greedy routing. Nodes independently calculate a short path to the target. At each hop, the query moves closer to the target; and
- d) robustness. A path to the target can be found even when links or nodes fail.

3.1. Origins

To understand the origins of recent DHTs, one needs to look to three contributions from the 1990s. The first two -- Plaxton, Rajaraman, and Richa (PRR) [30] and Consistent Hashing [49] -- were published within one month of each other. The third, the Scalable Distributed Data Structure (SDDS) [52], was curiously ignored in significant structured P2P designs despite having some similar goals [2, 6, 7]. It has been briefly referenced in other P2P papers [46, 284-287].

3.1.1. Plaxton, Rajaraman, and Richa (PRR)

PRR is the most recent of the three. It influenced the designs of Pastry [2], Tapestry [6], and Chord [7]. The value of PRR is that it can locate objects using fixed-length routing tables [6]. Objects and nodes are assigned a semantic-free address, for example a 160-bit key. Every node is effectively the root of a spanning tree. A message routes toward an object by matching longer address suffixes, until it encounters either the object's root node or another node

with a 'nearby' copy. It can route around link and node failure by matching nodes with a related suffix. The scheme has several disadvantages [6]: global knowledge is needed to construct the overlay; an object's root node is a single point of failure; nodes cannot be inserted and deleted; and there is no mechanism for queries to avoid congestion hot spots.

3.1.2. Consistent Hashing

Consistent Hashing [288] strongly influenced the designs of Chord [7] and Koorde [37]. Karger, et al. introduced Consistent Hashing in the context of the Web-caching problem [49]. Web servers could conceivably use standard hashing to place objects across a network of caches. Clients could use the approach to find the objects. For normal hashing, most object references would be moved when caches are added or deleted. On the other hand, Consistent Hashing is "smooth" -- when caches are added or deleted, the minimum number of object references move so as to maintain load balancing. Consistent Hashing also ensures that the total number of caches responsible for a particular object is limited. Whereas Litwin's Linear Hashing (LH*) scheme requires 'buckets' to be added one at a time in sequence [50], Consistent Hashing allows them to be added in any order [49]. There is an open Consistent Hashing problem pertaining to the fraction of items moved when a node is inserted [165]. Extended Consistent Hashing was recently proposed to randomize queries over the spread of caches to significantly reduce the load variance [289]. Interestingly, Karger [49] referred to an older DHT algorithm by Devine that used "a novel autonomous location discovery algorithm that learns the buckets' locations instead of using a centralized directory" [51].

3.1.3. Scalable Distributed Data Structures (LH*)

In turn, Devine's primary point of reference was Litwin's work on SDDSs and the associated LH* algorithm [52]. An SDDS satisfies three design requirements: files grow to new servers only when existing servers are well loaded; there is no centralized directory; and the basic operations like insert, search, and split never require atomic updates to multiple clients. Honicky and Miller suggested the first requirement could be considered a limitation since expansion to new servers is not under administrative control [286]. Litwin recently noted numerous similarities and differences between LH* and Chord [290]. He found that both implement key search. Although LH* refers to clients and servers, nodes can operate as peers in both. Chord 'splits' nodes when a new node is inserted, while LH* schedules 'splits' to avoid overload. Chord requests travel $O(\log n)$ hops, while LH* client requests need, at most, two hops to find the target. Chord stores a small number of 'fingers' at each node. LH* servers

store $N/2$ to N addresses while LH^* clients store 1 to N addresses. This trade-off between hop count and the size of the index affects system robustness, and bears striking similarity to recent one- and two-hop P2P schemes in Section 2. The arrival and departure of LH^* clients does not disrupt LH^* server metadata at all. Given the size of the index, the arrival and departure of LH^* servers are likely to cause more churn than that of Chord nodes. Unlike Chord, LH^* has a single point of failure, the split coordinator. It can be replicated. Alternatively, it can be removed in later LH^* variants, though details have not been progressed for lack of practical need [290].

3.2. Dependability

We make four overall observations about their dependability. Dependability metrics fall into two categories: static dependability, a measure of performance before recovery mechanisms take over; and dynamic dependability, for the most likely case in massive networks where there is continual failure and recovery ("churn").

3.2.1. Static Dependability

Observation A: Static dependability comparisons show that no $O(\log n)$ DHT geometry is significantly more dependable than the other $O(\log n)$ geometries.

Gummadi, et al. compared the tree, hypercube, butterfly, ring, XOR, and hybrid geometries. In such geometries, nodes generally know about $O(\log n)$ neighbors and route to a destination in $O(\log n)$ hops, where N is the number of nodes in the overlay. Gummadi, et al. asked "Why not the ring?" They concluded that only the ring and XOR geometries permit flexible choice of both neighbors and alternative routes [24]. Loguinov, et al. added the de Bruijn graph to their comparison [36]. They concluded that the classical analyses, for example the probability that a particular node becomes disconnected, yield no major differences between the resilience of Chord, CAN, and de Bruijn graphs. Using bisection width (the minimum edge count between two equal partitions) and path overlap (the likelihood that backup paths will encounter the same failed nodes or links as the primary path), they argued for the superior resilience of the de Bruijn graph. In short, ring, XOR, and de Bruijn graphs all permit flexible choice of alternative paths, but only in de Bruijn are the alternate paths independent of each other [36].

3.2.2. Dynamic Dependability

Observation B: Dynamic dependability comparisons show that DHT dependability is sensitive to the underlying topology maintenance algorithms.

Li, et al. give the best comparison to date of several leading DHTs during churn [291]. They relate the disparate configuration parameters of Tapestry, Chord, Kademlia, Kelips, and OneHop to fundamental design choices. For each of these DHTs, they plotted the optimal performance in terms of lookup latency (milliseconds) and fraction of failed lookups. The results led to several important insights about the underlying algorithms, for example: increasing routing table size is more cost-effective than increasing the rate of periodic stabilization; learning about new nodes during the lookup process sometimes eliminates the need for stabilization; and parallel lookups reduce latency due to timeouts more effectively than faster stabilization. Similarly, Zhuang, et al. compared keep-alive algorithms for DHT failure detection [292]. Such algorithmic comparisons can significantly improve the dependability of DHT designs.

In Figure 2, we propose a taxonomy for the topology maintenance algorithms that influence dependability. The algorithms can be classified by how nodes join and leave, how they first detect failures, how they share information about topology updates, and how they react when they receive information about topology updates.

Normal Updates

Joins (passive; active) [293]

Leaves (passive; active) [293]

Fault Detection [292]

Maintenance

Proactive (periodic or keep-alive probes)

Reactive (correction-on-use, correction-on-failure) [294]

Report

Negative (all dead nodes, nodes recently failed)

Positive (all live nodes; nodes recently recovered) [292]

Topology Sharing: yes/ no [292]

Multicast Tree (explicit, implicit) [267, 295]

Gossip (timeouts; number of contacts) [39]

Corrective Action

Routing

Rerouting actions

(reroute once; route in parallel [291]; reject)

Routing timeouts

(TCP-style, virtual coordinates) [296]

Topology

Update action (evict/ replace/ tag node)

Update timeliness (immediate, periodic[296], delayed [297])

Figure 2: Topology Maintenance in Distributed Hash Tables

3.2.3. Ephemeral or Stable Nodes -- $O(\log n)$ or $O(1)$ Hops

Observation C: Most DHTs use $O(\log n)$ geometries to suit ephemeral nodes. The $O(1)$ hop DHTs suit stable nodes and deserve more research attention.

Most of the DHTs in Section 3.5 assume that nodes are ephemeral, with expected lifetimes of one to two hours. Therefore, they mostly use an $O(\log n)$ geometry. The common assumption is that maintenance of full routing tables in the $O(1)$ hop DHTs will consume excessive bandwidth when nodes are continually joining and leaving. The corollary is that, when they run on stable infrastructure servers [298], most of the DHTs in Section 3.5 are less than optimal -- lookups take many more hops than necessary, wasting latency and bandwidth budgets. The $O(1)$ hop DHTs suit stable deployments and high lookup rates. For a churning 1024-node network, Li, et al. concluded that OneHop is superior to Chord, Tapestry, Kademlia, and Kelips in terms of latency and lookup success rate [291]. For a 3000-node network, they concluded that "OneHop is only preferable to Chord when the deployment scenario allows a communication cost

greater than 20 bytes per node per second" [291]. This apparent limitation needs to be put in context. They assumed that each node issues only one lookup every 10 minutes and has a lifetime of only 60 minutes. It seems reasonable to expect that in some deployments, nodes will have a lifetime of weeks or more, a maintenance bandwidth of tens of kilobits per second, and a load of hundreds of lookups per second. $O(1)$ hop DHTs are superior in such situations. OneHop can scale at least to many tens of thousands of nodes [267]. The recent $O(1)$ hop designs [267, 295] are vastly outnumbered by the $O(\log n)$ DHTs in Section 3.5. Research on the algorithms of Figure 2 will also yield improvements in the dependability of the $O(1)$ hop DHTs.

3.2.4. Simulation and Proof

Observation D: Although not yet a mature science, the study of DHT dependability is helped by recent simulation and formal development tools.

While there are recent reference architectures [294, 298], much of the DHT literature in Section 3.5 does not lend itself to repeatable, comparative studies. The best comparative work to date [291] relies on the Peer-to-Peer Simulator (P2PSIM) [299]. At the time of writing, it supports more DHT geometries than any other simulator. As the study of DHTs matures, we can expect to see the simulation emphasis shift from geometric comparison to a comparison of the algorithms of Figure 2.

P2P correctness proofs generally rely on less-than-complete formal specifications of system invariants and events [7, 45, 300]. Li and Plaxton expressed concern that "when many joins and leaves happen concurrently, it is not clear whether the neighbor tables will remain in a 'good' state" [47]. While acknowledging that guaranteeing consistency in a failure-prone network is impossible, Lynch, Malkhi, et al. sketched amendments to the Chord algorithm to guarantee atomicity [301]. More recently, Gilbert, Lynch, et al. gave a new algorithm for atomic read/write memory in a churning distributed network, suggesting it to be a good match for P2P [302]. Lynch and Stoica show in an enhancement to Chord that lookups are provably correct when there is a limited rate of joins and failures [303]. Fault Tolerant Active Rings is a protocol for active joins and leaves that was formally specified and proven using B-method tools [304]. A good starting point for a formal DHT development would be the numerous informal API specifications [22, 305, 306]. Such work could be informed by other efforts to formally specify routing invariants [307, 308].

3.3. Latency

The key metrics for DHT latency are:

- 1) Shortest-Path Distance and Diameter. In graph theory, the shortest-path distance is the minimum number of edges in any path between two vertices of the graph. Diameter is the largest of all shortest-path distances in a graph [309]. Networking synonyms for distance on a DHT are "hop count" and "lookup length".
- 2) Latency and Latency Stretch. Two types of latency are relevant here -- network-layer latency and overlay latency. Network-layer latency has been referred to as "proximity" or "locality" [24]. Stretch is the cost of an overlay path between two nodes, divided by the cost of the direct network path between those nodes [310]. Latency stretch is also known as the "relative delay penalty" [311].

3.3.1. Hop Count and the $O(1)$ -Hop DHTs

Hop count gives an approximate indication of path latency. $O(1)$ -hop DHTs have path latencies lower than the $O(\log n)$ -hop DHTs [291]. This significant advantage is often overlooked on account of concern about the messaging costs to maintain large routing tables (Section 3.2.3). Such concern is justified when the mean node lifetime is only a few hours and the mean lookup interval per node is more than a few seconds (the classic profile of a P2P file-sharing node). However, for a large, practical operating range (node lifetimes of days or more, lookup rates of over tens of lookups per second per node, up to ~100,000 nodes), the total messaging cost in $O(1)$ hop DHTs is lower than in $O(\log n)$ DHTs [312]. Lookups and routing table maintenance contribute to the total messaging cost. If a deployment fits this operating range, then $O(1)$ -hop DHTs will give lower path latencies and lower total messaging costs. An additional merit of the $O(1)$ -hop DHTs is that they yield lower lookup failure rates than their $O(\log N)$ -hop counterparts [291].

Low hop count can be achieved in two ways: each node has a large $O(N)$ index of nodes; or the object references can be replicated on many nodes. Beehive [313], Kelips [39], LAND [310], and Tulip [314] are examples of the latter category. Beehive achieves $O(1)$ hops on average and $O(\log n)$ hops in the worst case, by proactive replication of popular objects. Kelips replicates the 'file index'. It incurs $O(\sqrt{N})$ storage costs for both the node index and the file index. LAND uses $O(\log n)$ reference pointers for each stored object and an $O(\log n)$ index to achieve a worst-case $1+e$ stretch, where $0 < e$. The Kelips-like Tulip [314] requires 2 hops per lookup. Each node

maintains $2\sqrt{N}\log(N)$ links to other nodes and objects are replicated on $O(\sqrt{N})$ nodes.

The DHTs with a large $O(N)$ node index can be divided into two groups: those for which the index is always $O(N)$; and those for which the index opportunistically ranges from $O(\log n)$ to $O(N)$. Linear Hashing (LH*) servers [52], OneHop [267], and lh-Calot [295] fall into the former category. EpiChord [315] and Accordion [316] are examples of the latter.

3.3.2. Proximity and the $O(\log n)$ -Hop DHTs

If one chooses not to use single-hop DHTs, hop count is a weak indicator of end-to-end path latency. Some hops may incur large delays because of intercontinental or satellite links. Consequently, numerous DHT designs minimize path latency by considering the proximity of nodes. Gummadi, et al. classified the proximity methods as follows [24]:

- 1) Proximity Neighbor Selection (PNS). The nodes in the routing table are chosen based on the latency of the direct hop to those nodes. The latency may be explicitly measured [317], or it may be estimated using one of several synthetic coordinate systems [150, 154, 318]. As a lower bound on PNS performance, Dabek, et al. showed that lookups on $O(\log n)$ DHTs take at least 1.5 times the average roundtrip time of the underlying network [154].
- 2) Proximity Route Selection (PRS). At lookup time, the choice of the next-hop node relies on the latency of the direct hop to that node. PRS is less effective than PNS, though it may complement it [24]. Some of the routing geometries in Section 3.5 do not support PNS and/or PRS [24].
- 3) Proximity Identifier Selection (PIS). Node identifiers indicate geographic position. PIS frustrates load balancing, increases the risk of correlated failures, and is not often used [24].

The proximity study by Gummadi, et al. assumed recursive routing, though they suggested that PNS would also be superior to PRS with iterative routing [24]. Dabek, et al. found that recursive lookups take 0.6 times as long as iterative lookups [150].

Beyond the explicit use of proximity information, redundancy can help to avoid slow paths and servers. One may increase the number of replicas [150], use parallel lookups [291, 316], use alternate routes on failure [150], or use multiple gateway nodes to enter the DHT [317].

3.4. Multicasting

3.4.1. Multicasting vs. Broadcasting

"Multicasting" here means sending a message to a subset of an overlay's nodes. Nodes explicitly join and leave this subset, called a "multicast group". "Broadcasting" here is a special case of multicasting in which a message is sent to all nodes in the overlay. Broadcasting relies on overlay membership messages -- it does not need extra group membership messaging. Castro, et al. said multicasting on structured overlays is either "flooding" (one overlay per group) or "tree-based" (one tree per group) [319]. These are synonyms for broadcasting and multicasting respectively.

The first DHT-based designs for multicasting were CAN multicast [320], Scribe [241], Bayeux [242], and i3 [231]. They were based on CAN [8], Pastry [2], Tapestry [31], and Chord [7] respectively. El-Ansary, et al. devised the first DHT-based broadcasting scheme [321]. It was based on Chord.

Multicast trees can be constructed using reverse-path forwarding or forward-path forwarding. Scribe uses reverse-path forwarding [241]. Bayeux uses forward-path forwarding [242]. Borg, a multicast design based on Pastry, uses a combination of forward-path and reverse-path forwarding to minimize latency [237].

3.4.2. Motivation for DHT-based Multicasting

Multicasting complements DHT search capability. DHTs naturally support exact match queries. With multicasting, they can support more complex queries. Multicasting also enables the dissemination and collection of global information.

Consider, for example, aggregation queries like minimum, maximum, count, sum, and average (Section 5.4). A node at the root of a dissemination tree might multicast such a query [322]. The leaf nodes return local results towards the root node. Successive parents aggregate the result so that eventually the root node can compute the global result. Such queries may help to monitor the capacity and health of the overlay itself.

Why bother with structured overlays for multicasting? In Section 2.1, we saw that Gnutella can multicast complex queries without them [4]. Castro, et al. posed the question, "Should we build Gnutella on a structured overlay?" [259]. While acknowledging that their study was preliminary, they did conclude that "we see no reason to build Gnutella on top of an unstructured overlay" [259]. The supposedly high maintenance costs of structured overlays were outweighed by

query cost savings. The structured overlay ensured that nodes were only visited once during a complex query. It also helped to accurately limit the total number of nodes visited. Pai, et al. acknowledged that multicast trees based on structured overlays contribute to simple routing rules, low delay and low delay variation [323]. However, they opted for unstructured, gossip-based multicasting for reliability reasons: data loss near the tree root affects all subtended nodes; interior node failures must be repaired quickly; interior nodes are obliged to disseminate more than their fair share of traffic, giving leaf nodes a "free ride". The most promising research direction is to improve on the Bimodal Multicasting approach [324]. It combines the bandwidth efficiency and low latency of structured, best-effort multicasting trees with the reliability of unstructured gossip protocols.

3.4.3. Design Issues

None of the early structured overlay multicast designs addressed all of the following issues [325]:

- 1) Heterogeneous Node Capacity. Nodes differ in their processing, memory, and network capacity. Multicast throughput is largely determined by the node with smallest throughput [325]. To limit the multicasting load on a node, one might cap its out-degree. If the same node receives further join requests, it refers them to its children ("pushdown") [240]. Bharambe, et al. explored several pushdown strategies but found them inadequate to deal with heterogeneity [326]. They concluded that the heterogeneity issue remains open, and should be addressed before deploying DHTs for high-bandwidth multicasting applications. Independently, Zhang et al. partially tackled heterogeneity by allowing nodes in their CAM-Chord and CAM-Koorde designs to vary out-degree according to the node's capacity [325]. However, they made no mention of the "pushdown" issue -- they did not describe topology maintenance when the out-degree limit is reached.
- 2) Reliability (Dynamic Membership). If a multicast tree is to be resilient, it must survive dynamic membership. There are several ways to deal with dynamic membership: ensure that the root node of the multicasting tree does not handle all requests to join or leave the multicast group [242]; use multiple interior-node-disjoint trees to avoid single points of failure in tree structures [322]; and split the root node into several replicas and partition members across them [241]. For example, Bayeux requires the root node to track all group membership changes whereas Scribe does not [241]. CAN-multicast uses a single, well-known host to bootstrap the join operations [320]. The earliest DHT-based broadcasting work by El-Ansary, et al. did not

address the issue of dynamic membership [321]. Ghodsi, et al. addressed it in a subsequent paper, though, giving two broadcast algorithms that accommodate routing table inconsistencies [327]. One algorithm achieves a more optimal multicasting network at the expense of greater correction overhead. Splitstream, based on Scribe and Pastry, redundantly striped content across multiple interior-node-disjoint multicast trees -- if one interior node fails, then only one stripe is lost [240].

- 3) Large Any-Source Multicast Groups. Any group member should be allowed to send multicast messages. The group should scale to a very large number of hosts. CAN-based multicast was the first application-level multicast scheme to scale to groups of several thousands of nodes without restricting the service model to a single source [320]. Bayeux scales to large groups but has a single root node for each multicast group. It supports the any-source model only by having the root node operate as a reflector for multiple senders [242].

3.5. Routing Geometries

In Sections 3.5.1 to 3.5.6, we introduce the main geometries for simple key lookup and survey their robustness mechanisms.

3.5.1. Plaxton Trees (Pastry, Tapestry)

Work began in March 2000 on a structured, fault-tolerant, wide-area Dynamic Object Location and Routing (DOLR) system called Tapestry [6, 155]. While DHTs fix replica locations, a DOLR API enables applications to control object placement [31]. Tapestry's basic location and routing scheme follows Plaxton, Rajaraman, and Richa (PRR) [30], but it remedies PRR's robustness shortcomings described in Section 3.1. Whereas each object has one root node in PRR, Tapestry uses several to avoid a single point of failure. Unlike PRR, it allows nodes to be inserted and deleted. Whereas PRR required a total ordering of nodes, Tapestry uses 'surrogate routing' to incrementally choose root nodes. The PRR algorithm does not address congestion, but Tapestry can put object copies close to nodes generating high query loads. PRR nodes only know of the nearest replica, whereas Tapestry nodes enable selection from a set of replicas (for example, to retrieve the most up to date). To detect routing faults, Tapestry uses TCP timeouts and UDP heartbeats for detection, sequential secondary neighbours for rerouting, and a 'second chance' window so that recovery can occur without the overhead of a full node insertion. Tapestry's dependability has been measured on a testbed of about 100 machines and on simulations of

about 1000 nodes. Successful routing rates and maintenance bandwidths were measured during instantaneous failures and ongoing churn [31].

Pastry, like Tapestry, uses Plaxton-like prefix routing [2]. As in Tapestry, Pastry nodes maintain $O(\log n)$ neighbours and route to a target in $O(\log n)$ hops. Pastry differs from Tapestry only in the method by which it handles network locality and replication [2]. Each Pastry node maintains a 'leaf set' and a 'routing table'. The leaf set contains $1/2$ node IDs on either side of the local node ID in the node ID space. The routing table, in row r , column c , points to the node ID with the same r -digit prefix as the local node, but with an $r+1$ digit of c . A Pastry node periodically probes leaf set and routing table nodes, with periodicity of Tls and Trt and a timeout $Tout$. Mahajan, Castry, et al. analyzed the reliability versus maintenance cost trade-offs in terms of the parameters l , Tls , Trt , and $Tout$ [328]. They concluded that earlier concerns about excessive maintenance cost in a churning P2P network were unfounded, but suggested follow-up work for a wider range of reliability targets, maintenance costs, and probe periods. Rhea Geels, et al. concluded that existing DHTs fail at high churn rates [329]. Building on a Pastry implementation from Rice University, they found that most lookups fail to complete when there is excessive churn. They conjectured that short-lived nodes often leave the network with lookups that have not yet timed out, but no evidence was provided to confirm the theory. They identified three design issues that affect DHT performance under churn: reactive versus periodic recovery of peers; lookup timeouts; and choice of nearby neighbours. Since reactive recovery was found to add traffic to already congested links, the authors used periodic recovery in their design. For lookup timeouts, they advocated an exponentially weighted moving average of each neighbour's response time, over alternative fixed timeout or 'virtual coordinate' schemes. For selection of nearby neighbours, they found that 'global sampling' was more effective than simply sampling a 'neighbour's neighbours' or 'inverse neighbours'. Castro, Costa, et al. have refuted the suggestion that DHTs cannot cope with high churn rates [330]. By implementing methods for continuous detection and repair, their MSPastry implementation achieved shorter routing paths and a maintenance overhead of less than half a message per second per node.

There have been more recent proposals based on these early Plaxton-like schemes. Kademlia uses a bit-wise exclusive or (XOR) metric for the 'distance' between 160-bit node identifiers [45]. Each node keeps a list of contact nodes for each section of the node space that is between 2^i and $2^{(i+1)}$ from itself ($0 \leq i < 160$). Longer-lived nodes are deliberately given preference on this list -- it has been found in Gnutella that the longer a node has been active, the more likely

it is to remain active. Like Kademlia, Willow uses the XOR metric [32]. It implements a Tree Maintenance Protocol to 'zipper' together broken segments of a tree. Where other schemes use DHT routing to inefficiently add new peers, Willow can merge disjoint or broken trees in $O(\log n)$ parallel operations.

3.5.2. Rings (Chord, DKS)

Chord is the prototypical DHT ring, so we first sketch its operation. Chord maps nodes and keys to an identifier ring [7, 34]. Chord supports one main operation: find a node with the given key. It uses Consistent Hashing (Section 3.1) to minimize disruption of keys when nodes join and leave the network. However, Chord peers need only track $O(\log n)$ other peers, not all peers as in the original consistent hashing proposal [49]. It enables concurrent node insertions and deletions, improving on PRR. Compared to Pastry, it has a simpler join protocol. Each Chord peer tracks its predecessor, a list of successors, and a finger table. Using the finger table, each hop is at least half the remaining distance around the ring to the target node, giving an average lookup hop count of $(1/2)\log n(\text{base } 2)$. Each Chord node runs a periodic stabilization routine that updates predecessor and successor pointers to cater to newly added nodes. All successors of a given node need to fail for the ring to fail. Although a node departure could be treated the same as a failure, a departing Chord node first notifies the predecessor and successors, so as to improve performance.

In their definitive paper, Chord's inventors critiqued its dependability under churn [34]. They provided proofs on the behaviour of the Chord network when nodes in a stable network fail, stressing that such proofs are inadequate in the general case of a perpetually churning network. An earlier paper had posed the question, "For lookups to be successful during churn, how regularly do the Chord stabilization routines need to run?" [331]. Stoica, Morris, et al. modeled a range of node join/departure rates and stabilization periods for a Chord network of 1000 nodes. They measured the number of timeouts (caused by a finger pointing to a departed node) and lookup failures (caused by nodes that temporarily point to the wrong successor during churn). They also modeled the 'lookup stretch', the ratio of the Chord lookup time to optimal lookup time on the underlying network. They demonstrated the latency advantage of recursive lookups over iterative lookups, but there remains room for delay reduction. For further work, the authors proposed to improve resilience to network partitions, using a small set of known nodes or 'remembered' random nodes. To reduce the number of messages per lookup, they suggested an increase in the size of each step around the ring, accomplished via a larger number of fingers at each node. Much of the paper assumed independent, equally

likely node failures. Analysis of correlated node failures, caused by massive site or backbone failures, will be more important in some deployments. The paper did not attempt to recommend a fixed optimal stabilization rate. Liben-Nowell, Balakrishnan, et al. had suggested that optimum stabilization rate might evolve according to measurements of peers' behaviour [331] -- such a mechanism has yet to be devised.

Alima, El-Ansary, et al. considered the communication costs of Chord's stabilization routines, referred to as 'active correction', to be excessive [332]. Two other robustness issues also motivated their Distributed K-ary Search (DKS) design, which is similar to Chord. Firstly, the total system should evolve for an optimum balance between the number of peers, the lookup hop count, and the size of the routing table. Secondly, lookups should be reliable -- P2P algorithms should be able to guarantee a successful lookup for key/value pairs that have been inserted into the system. A similar lookup-correctness issue was raised elsewhere by one of Chord's authors; "Is it possible to augment the data structure to work even when nodes (and their associated finger lists) just disappear?" [333] Alima, El-Ansary, et al. asserted that P2Ps using active correction, like Chord, Pastry, and Tapestry, are unable to give such a guarantee. They propose an alternate 'correction-on-use' scheme, whereby expired routing entries are corrected by information piggybacking lookups and insertions. A prerequisite is that lookup and insertion rates are significantly higher than node arrival, departure, and failure rates. Correct lookups are guaranteed in the presence of simultaneous node arrivals or up to f concurrent node departures, where f is configurable.

3.5.3. Tori (CAN)

Ratnasamy, Francis, et al. developed the Content-Addressable Network (CAN), another early DHT widely referenced alongside Tapestry, Pastry, and Chord [8, 334]. It is arranged as a virtual d -dimensional Cartesian coordinate space on a d -torus. Each node is responsible for a zone in this coordinate space. The designers used a heuristic thought to be important for large, churning P2P networks: keep the number of neighbours independent of system size. Consequently, its design differs significantly from Pastry, Tapestry, and Chord. Whereas they have $O(\log n)$ neighbours per node and $O(\log n)$ hops per lookup, CAN has $O(d)$ neighbours and $O(dn^{1/d})$ hop count. When CAN's system-wide parameter d is set to $\log(n)$, CAN converges to their profile. If the number of nodes grows, a major rearrangement of the CAN network may be required [151]. The CAN designers considered building on PRR, but opted for the simple, low-state-per-node CAN algorithm instead. They had reasoned that a PRR-based design would not perform well under churn, given node

departures and arrivals would affect a logarithmic number of nodes [8].

There have been preliminary assessments of CAN's resilience. When a node leaves the CAN in an orderly fashion, it passes its own Virtual ID (VID), its neighbours' VIDs and IP addresses, and its key/value pairs to a takeover node. If a node leaves abruptly, its neighbours send recovery messages towards the designated takeover node. CAN ensures the recovery messages reach the takeover node, even if nodes die simultaneously, by maintaining a VID chain with Chord's stabilization algorithm. Some initial 'proof of concept' resilience simulations were run using the Network Simulator (NS) [335] for up to a few hundred nodes. Average hop counts and lookup failure probabilities were plotted against the total number of nodes for various node failure rates [8]. The CAN team documented several open research questions pertaining to state/hop count trade-offs, resilience, load, locality, and heterogeneous peers [44, 334].

3.5.4. Butterflies (Viceroy)

Viceroy approximates a butterfly network [46]. It generally has constant degree like CAN. Like Chord, Tapestry, and Pastry, it has logarithmic diameter. It improves on these systems, inasmuch as its diameter is better than CAN and its degree is better than Chord, Tapestry, and Pastry. As with most DHTs, it utilizes Consistent Hashing. When a peer joins the Viceroy network, it takes a random but permanent 'identity' and selects its 'level' within the network. Each peer maintains general ring pointers ('predecessor' and 'successor'), level ring pointers ('nextonlevel' and 'prevonlevel'), and butterfly pointers ('left', 'right', and 'up'). When a peer departs, it normally passes its key pairs to a successor, and notifies other peers to find a replacement peer.

The Viceroy paper scoped out the issue of robustness. It explicitly assumed that peers do not fail [46]. It assumed that join and leave operations do not overlap, so as to avoid the complication of concurrency mechanisms like locking. Kaashoek and Karger were somewhat critical of Viceroy's complexity [37]. They also pointed to its fault-tolerance blind spot. Li and Plaxton suggested that such constant-degree algorithms deserve further consideration [47]. They offered several pros and cons. The limited degree may increase the risk of a network partition, or inhibit use of local neighbours (for the simple reason that there are less of them). On the other hand, it may be easier to reason about the correctness of fixed-degree networks. One of the Viceroy authors has since proposed constant-degree peers in a two-tier, locality-aware DHT [310] -- the lower degree maintained by each lower-tier peer purportedly improves network adaptability. Another Viceroy author has since explored an

alternative bounded-degree graph for P2P, namely the de Bruijn graph [336].

3.5.5. de Bruijn (D2B, Koorde, Distance Halving, ODRI)

De Bruijn graphs have had numerous refinements since their inception [337, 338]. Schlumberger was the first to use them for networking [339]. Two research teams independently devised the 'generalized' de Bruijn graph that accommodates a flexible number of nodes in the system [340, 341]. Rowley and Bose studied fault-tolerant rings overlaid on the de Bruijn graph [342]. Lee, Liu, et al. devised a two-level de Bruijn hierarchy, whereby clusters of local nodes are interconnected by a second-tier ring [343].

Many of the algorithms discussed previously are 'greedy' in that each time a query is forwarded, it moves closer to the destination. Unfortunately, greedy algorithms are generally suboptimal -- for a given degree, the routing distance is longer than necessary [344]. Unlike these earlier P2P designs, de Bruijn graphs of degree k achieve an asymptotically optimal diameter $\log n$, where n is the number of nodes in the system and k can be varied to improve resilience. If there are $O(\log n)$ neighbours per node, the de Bruijn hop count is $O(\log n / \log \log n)$. To illustrate de Bruijn's practical advantage, consider a network with one million nodes of degree 20: Chord has a diameter of 20, while de Bruijn has a diameter of 5 [36]. In 2003, there were a quick succession of de Bruijn proposals -- D2B [345], Koorde [37], Distance Halving [132, 336], and the Optimal Diameter Routing Infrastructure (ODRI) [36].

Fraigniaud and Gauron began the D2B design by laying out an informal problem statement: keys should be evenly distributed; lookup latency should be small; traffic load should be evenly distributed; updates of routing tables and redistribution of keys should be fast when nodes join or leave the network. They defined a node's "congestion" to be the probability that a lookup will traverse it. Apart from its optimal de Bruijn diameter, they highlighted D2B's merits: a constant expected update time when nodes join and leave ($O(\log n)$ with high probability (w.h.p.)); the expected node congestion is $O((\log n)/n)$ ($O((\log n)^2/n)$ w.h.p.) [345]. D2B's resilience was discussed only in passing.

Koorde extends Chord to attain the optimal de Bruijn degree/diameter trade-off above [37]. Unlike D2B, Koorde does not constrain the selection of node identifiers. Also unlike D2B, it caters to concurrent joins, by extension of Chord's functionality. Kaashoek and Karger investigated Koorde's resilience to a rather harsh failure scenario: "in order for a network to stay connected when all nodes fail with probability of $1/2$, some nodes must have degree

$\text{omega}(\log n)$ " [37]. They sketched a mechanism to increase Koorde's degree for this more stringent fault tolerance, losing de Bruijn's constant degree advantage. Similarly, to achieve a constant-factor load balance, Koorde would have to sacrifice its degree optimality. They suggested that the ability to trade the degree, and hence the maintenance overhead, against the expected hop count may be important for churning systems. They also identified an open problem: find a load-balanced, degree optimal DHT. Datta, Girdzijauskas, et al. showed that for arbitrary key distributions, de Bruijn graphs fail to meet the dual goals of load balancing and search efficiency [346]. They posed the question, "(Is there) a constant routing table sized DHT which meets the conflicting goals of storage load balancing and search efficiency for an arbitrary and changing key distribution?"

Distance Halving was also inspired by de Bruijn [336] and shares its optimal diameter. Naor and Wieder argued for a two-step "continuous-discrete" approach for its design. The correctness of its algorithms is proven in a continuous setting. The algorithms are then mapped to a discrete space. The source x and target y are points on the continuous interval $[0,1)$. Data items are hashed to this same interval. $\langle \text{str} \rangle$ is a string that determines how messages leave any point on the ring: if bit t of the string is 0, the left leg is taken; if it is 1, the right leg is taken. $\langle \text{str} \rangle$ increases by one bit each hop, giving a sequence by which to step around the ring. A lookup has two phases. In the first, the lookup message containing the source, target, and the random string hops toward the midpoint of the source and target. On each hop, the distance between $\langle \text{str} \rangle(x)$ and $\langle \text{str} \rangle(y)$ is halved, by virtue of the specific 'left' and 'right' functions. In the second phase, the message steps 'backward' from the midpoint to the target, removing the last bit in $\langle \text{str} \rangle$ at each hop. 'Join' and 'leave' algorithms were outlined but there was no consideration of recovery times or message load on churn. Using the Distance Halving properties, the authors devised a caching scheme to relieve congestion in a large P2P network. They have also modified the algorithm to be more robust in the presence of random faults [132].

Solid comparisons of DHT resilience are scarce, but Loguinov, Kumar, et al. give just that in their ODRI paper [36]. They compare Chord, CAN, and de Bruijn in terms of routing performance, graph expansion and clustering. At the outset, they give the optimal diameter (the maximum hop count between any two nodes in the graph) and average hop count for graphs of fixed degree. De Bruijn graphs converge to both optima, and outperform Chord and CAN on both counts. These optima impact both delay and aggregate lookup load. They present two clustering measures (edge expansion and node expansion), which are interesting for resilience. Unfortunately, after decades of de Bruijn research, they have no exact solution. De Bruijn was shown to

be superior in terms of path overlap - "de Bruijn automatically selects backup paths that do not overlap with the best shortest path or with each other" [36].

3.5.6. Skip Graphs

Skip Graphs have been pursued by two research camps [38, 41]. They augment the earlier Skip Lists [347, 348]. Unlike earlier balanced trees, the Skip List is probabilistic -- its insert and delete operations do not require tree rearrangements and so are faster by a constant factor. The Skip List consists of layers of ordered linked lists. All nodes participate in the bottom layer 0 list. Some of these nodes participate in the layer 1 list with some fixed probability. A subset of layer 1 nodes participate in the layer 2 list, and so on. A lookup can proceed quickly through the list by traversing the sparse upper layers until it is close to, or at, the target. Unfortunately, nodes in the upper layers of a Skip List are potential hot spots and single points of failure. Unlike Skip Lists, Skip Graphs provide multiple lists at each level for redundancy, and every node participates in one of the lists at each level.

Each node in a Skip Graph has $\theta(\log n)$ neighbours on average, like some of the preceding DHTs. The Skip Graph's primary edge over the DHTs is its support for prefix and proximity search. DHTs hash objects to a random point in the graph. Consequently, they give no guarantees over where the data is stored. Nor do they guarantee that the path to the data will stay within the one administration as far as possible [38]. Skip graphs, on the other hand, provide for location-sensitive name searches. For example, to find the document `docname` on the node `user.company.com`, the Skip Graph might step through its ordered lists for the prefix `com.company.user` [38]. Alternatively, to find an object with a numeric identifier, an algorithm might search the lowest layer of the Skip Graph for the first digit, the next layer for the next digit, in the same vein until all digits are resolved. Being ordered, Skip Graphs also facilitate range searches. In each of these examples, the Skip Graph can be arranged such that the path to the target, as far as possible, stays within an administrative boundary. If one administration is detached from the rest of the Skip Graph, routing can continue within each of the partitions. Mechanisms have been devised to merge disconnected segments [157], though at this stage, segments are re-merged one at a time. A parallel merge algorithm has been flagged for future work.

The advantages of Skip Graphs come at a cost. To be able to provide range queries and data placement flexibility, Skip Graph nodes require many more pointers than their DHT counterparts. An increased number of pointers implies increased maintenance traffic. Another

shortcoming of at least one of the early proposals was that no algorithm was given to assign keys to machines. Consequently, there are no guarantees on system-wide load balancing or on the distance between adjacent keys [100]. Aspnes, Kirsch, et al. have recently devised a scheme to reduce the inter-machine pointer count from $O(m \log m)$, where m is the number of data elements, to $O(n \log n)$, where n is the number of nodes [100]. They proposed a two-layer scheme -- one layer for the Skip Graph itself and the second 'bucket layer'. Each machine is responsible for a number of buckets and each bucket elects a representative key. Nodes locally adjust their load. They accept additional keys if they are below their threshold or disperse keys to nearby nodes if they are above threshold. There appear to be numerous open issues: simulations have been done but analysis is outstanding; mechanisms are required to handle the arrival and departure of nodes; there were only brief hints as to how to handle nodes with different capacities.

4. Semantic Index

Semantic indexes capture object relationships. While the semantic-free methods (DHTs) have firmer theoretic foundations and guarantee that a key can be found if it exists, they do not capture the relationships between the document name and its content or metadata on their own. Semantic P2P designs do. However, since their design is often driven by heuristics, they may not guarantee that scarce items will be found.

So what might the semantically indexed P2Ps add to an already crowded field of distributed information architectures? At one extreme, there are the distributed relational database management systems (RDBMSs), with their strong consistency guarantees [284]. They provide strong data independence, the flexibility of SQL queries, and strong transactional semantics -- Atomicity, Consistency, Isolation and Durability (ACID) [349]. They guarantee that the query response is complete -- all matching results are returned. The price is performance. They scale to perhaps 1000 nodes, as evidenced in Mariposa [350, 351], or require query caching front ends to constrain the load [284]. Database research has "arguably been cornered into traditional, high-end, transactional applications" [72]. Then there are distributed file systems, like the Network File System (NFS) or the Serverless Network File Systems (xFS), with little data independence, low-level file retrieval interfaces, and varied consistency [284]. Today's eclectic mix of Content Distribution Networks (CDNs) generally deload primary servers by redirecting Web requests to a nearby replica. Some intercept the HTTP requests at the DNS level and then use consistent hashing to find a replica [23]. Since this same consistent hashing was a forerunner to the DHT

approaches above, CDNs are generally constrained to the same simple key lookups.

The opportunity for semantically indexed P2Ps, then, is to provide:

- a) graduated data independence, consistency, and query flexibility, and
- b) probabilistically complete query responses, across
- c) very large numbers of low-cost, geographically distributed, dynamic nodes.

4.1. Keyword Lookup

P2P keyword lookup is best understood by considering the structure of the underlying index and the algorithms by which queries are routed over that index. Figure 3 summarizes the following paragraphs by classifying the keyword query algorithms, index structures, and metrics. The research has largely focused on scalability, not dependability. There have been very few studies that quantify the impact of network churn. One exception is the work by Chawathe, et al. on the Gia system [61]. Gia's combination of algorithms from Figure 3 (receiver-based flow control, biased random walk, and one-hop replication) gave 2-4 orders of magnitude improvement in query success rates in churning networks.

QUERY

Query routing

Flooding: Peers only index local files so queries must propagate widely [4]

Policy-based: Choice of the next hop node: random; most/least recently used; most files shared; most results [265, 352]

Random walks: Parallel [67] or biased random walks [61, 66]

Query forwarding

Iterative: Nodes perform iterative unicast searches of ultrapeers, until the desired number of results is achieved. See Gnutella UDP Extension for Scalable Searches (GUESS) [265, 353]

Recursive

Query flow control

Receiver-controlled: Receivers grant query tokens to senders, so as to avoid overload [61]

Reactive: sender throttles queries when it notices receivers are discarding packets [61, 66]

Dynamic Time To Live: In the Dynamic Query Protocol, the sender adjusts the time-to-live on each iteration based on the number of results received, the number of connections left, and the number of nodes already theoretically reached by the search [354]

INDEX

Distribution

Compression: Leaf nodes periodically send ultrapeers compressed query routing tables, as in the Query Routing Protocol [260]

One hop replication: Nodes maintain an index of content on their nearest neighbors [61, 352]

Partitioning

By document [210]

By keyword: Use an inverted list to find a matching document, either locally or at another peer [21]. Partition by keyword sets [355]

By document and keyword: Also called Multi-Level Partitioning [21]

METRIC

Query load: Queries per second per node/link [65, 265]

Degree: The number of links per node [66, 352]. Early P2P networks approximated power-law networks, where the number of nodes with L links is proportional to L^{-k} , where k is a constant [65]

Query delay: Reported in terms of time and hop count [61, 66]

Query success rate: The "Collapse Point" is the per-node query rate at which the query success rate drops below 90% [61]. See also [61, 265, 352].

Figure 3: Keyword Lookup in P2P Systems

4.1.1.1. Gnutella Enhancements

Perhaps the most widely referenced P2P system for simple keyword match is Gnutella [4]. Gnutella queries contain a string of keywords. Gnutella peers answer when they have files whose names contain all the keywords. As discussed in Section 2.1, early versions of Gnutella did not forward the document index. Queries were flooded and peers searched their own local indexes for filename matches. An early review highlighted numerous areas for improvement [65]. It was estimated that the query traffic alone from 50,000 early-generation Gnutella nodes would amount to 1.7% of the total U.S. Internet backbone traffic at December 2000 levels. It was speculated that high-degree Gnutella nodes would impede dependability. An unnecessarily high percentage of Gnutella traffic crossed Autonomous System (AS) boundaries -- a locality mechanism may have found suitable nearby peers.

Fortunately, there have since been numerous enhancements within the Gnutella Developer Forum. At the time of writing, it has been reported that Gnutella has almost 350,000 unique hosts, of which nearly 90,000 accept incoming connections [356]. One of the main improvements is that an index of filename keywords, called the Query Routing Table (QRT), can now be forwarded from 'leaf peers' to its 'ultrapeers' [260]. Ultrapeers can then ensure that the leaves only receive queries for which they have a match, dramatically reducing the query traffic at the leaves. Ultrapeers can have connections to many leaf nodes (~10-100) and a small number of other ultrapeers (<10) [260]. Originally, a leaf node's QRT was not forwarded by the parent ultrapeer to other ultrapeers. More recently, there has been a proposal to distribute aggregated QRTs amongst ultrapeers [357]. To further limit traffic, QRTs are compressed by hashing, according to the Query Routing Protocol (QRP) specification [281]. This same specification claims QRP may reduce Gnutella traffic by orders of magnitude, but cautions that simulation is required before mass deployment. A known shortcoming of QRP was that the extent of query propagation was independent of the popularity of the search terms. The Dynamic Query Protocol addressed this [358]. It required leaf nodes to send single queries to high-degree ultrapeers that adjust the queries' time-to-live (TTL) bounds according to the number of received query results. An earlier proposal, called the Gnutella UDP Extension for Scalable Searches (GUESS) [353], similarly aimed to reduce the number of queries for widely distributed files. GUESS reuses the non-forwarding idea (Section 2). A GUESS peer repeatedly queries single ultrapeers with a TTL of 1, with a small timeout on each query to limit load. It chooses the number of iterations and selects ultrapeers so as to satisfy its search needs. For adaptability, a small number of experimental Gnutella nodes have

implemented eXtensible Markup Language (XML) schemas for richer queries [359, 360]. None of the above Gnutella proposals explicitly assess robustness.

The broader research community has recently been leveraging aspects of the Gnutella design. Lv, Ratnasamy, et al. exposed one assumption implicit in some of the early DHT work -- that designs "such as Gnutella are inherently not scalable, and therefore should be abandoned" [66]. They argued that by making better use of the more powerful peers, Gnutella's scalability issues could be alleviated. Instead of its flooding mechanism, they used random walks. Their preliminary design to bias random walks towards high capacity nodes did not go as far as the ultrapeer proposals in that the indexes did not move to the high-capacity nodes. Chawathe, Ratnasamy, et al. chose to extend the Gnutella design with their Gia system, in response to the perceived shortcomings of DHTs in Section 1.2 [61]. Compared to the early Gnutella designs, they incorporated several novel features. They devise a topology adaptation algorithm so that most peers are attached to high-degree peers. They use a random walk search algorithm, in lieu of flooding, and bias the query load towards higher-degree peers. For 'one-hop replication', they require all nodes to keep pointers to content on adjacent peers. To implement a receiver-controlled token-based flow control, a peer must have a token from its neighbouring peer before it sends a query to it. Chawathe, Ratnasamy, et al. show by simulations that the combination of these features provides a scalability improvement of three to five orders of magnitude over Gnutella "while retaining significant robustness". The main robustness metrics they used were the 'collapse point' query rate (the per-node query rate at which the successful query rate falls below 90%) and the average hop count immediately prior to collapse. Their comparison with Gnutella did not take into account the Gnutella enhancements above -- this was left as future work. Castro, Costa, and Rowstron argued that if Gnutella were built on top of a structured overlay, then both the query and overlay maintenance traffic could be reduced [259]. Yang, Vinograd, et al. explore various policies for peer selection in the GUESS protocol, since the issue is left open in the original proposal [265]. For example, the peer initiating the query could choose peers that have been "most recently used" or that have the "most files shared". Various policy pitfalls are identified. For example, good peers could be overloaded, victims of their own success. Alternatively, malicious peers could encourage the querying peer to try inactive peers. They conclude that a "most results" policy gives the best balance of robustness and efficiency. Like Castro, Costa, and Rowstron, they concentrated on the static network scenario. Cholvi, Felber, et al. very briefly describe how similar "least recently used" and "most often used" heuristics can be used by a peer to select peer 'acquaintances' [352]. They were motivated by the

congestion associated with Gnutella's TTL-limited flooding. Recognizing that the busiest peers can quickly become overloaded central hubs for the entire network, they limit the number of acquaintances for any given peer to 25. They sketch a mechanism to decrement a query's TTL multiple times when it traverses "interested peers". In summary, these Gnutella-related investigations are characterized by a bias for high-degree peers and very short directed query paths, a disdain for flooding, and concern about excessive load on the 'better' peers. Generally, the robustness analysis for dynamic networks (content updates and node arrivals/departures) remains open.

4.1.2. Partition-by-Document, Partition-by-Keyword

One aspect of P2P keyword search systems has received particular attention: should the index be partitioned by document or by keyword? The issue affects scalability. To be partitioned by document, each node has a local index of documents for which it is responsible. Gnutella is a prime example. Queries are generally flooded in systems partitioned by document. On the other hand, a peer may assume responsibility for a set of keywords. The peer uses an inverted list to find a matching document, either locally or at another peer. If the query contains several keywords, inverted lists may need to be retrieved from several different peers to find the intersection [21]. The initial assessment by Li, Loo, et al. was that the partition-by-document approach was superior [210]. For one scenario of a full-text Web search, they estimated the communications costs to be about six times higher than the feasible budget. However, wanting to exploit prior work on inverted list intersection, they studied the partition-by-keyword strategy. They proposed several optimizations that put the communication costs for a partition-by-keyword system within an order of magnitude of feasibility. There had been a couple of prior papers that suggested partitioned-by-keyword designs incorporate DHTs to map keywords to peers [355, 361]. In Gnawali's Keyword-set Search System (KSS), the index is partitioned by sets of keywords [355]. Terpstra, Behnel, et al. point out that by keeping keyword pairs or triples, the number of lists per document in KSS is squared or tripled [362]. Shi, Guangwen, et al. interpreted the approximations of Li, Loo, et al. to mean that neither approach is feasible on its own [21]. Their Multi-Level Partitioning (MLP) scheme incorporates both partitioning approaches. They arrange nodes into a group hierarchy, with all nodes in the single 'level 0' group, and with the same nodes subdivided into k logical subgroups on 'level 1'. The subgroups are again divided, level by level, until level l. The inverted index is partitioned by document between groups and by keyword within groups. MLP avoids the query flooding normally associated with systems partitioned by document, since a small number of nodes in each group

process the query. It reduces the bandwidth overheads associated with inverted list intersection in systems partitioned solely by keyword, since groups can calculate the intersection independently over the documents for which they are responsible. MLP was overlaid on SkipNet, per Section 3.5.6 [38]. Some initial analyses of communications costs and query latencies were provided.

4.1.3. Partial Search, Exhaustive Search

Much of the research above addresses partial keyword search. Daswani, et al. highlighted the open problem of efficient, comprehensive keyword search [25]. How can exhaustive searches be achieved without flooding queries to every peer in the network? Terpstra, Behnel et al. couched the keyword search problem in rendezvous terms: dynamic keyword queries need to 'meet' with static document lists [362]. Their Bitzipper scheme is partitioned by document. They improved on full flooding by putting document metadata on $2\sqrt{n}$ nodes and forwarding queries through only $6\sqrt{n}$ nodes. They reported that Bitzipper nodes need only 1/166th of the bandwidth of full-flooding Gnutella nodes for an exhaustive search. An initial comparison of query load was given. There was little consideration of either static or dynamic resilience; that is, of nodes failing, of documents continually changing, or of nodes continually joining and leaving the network.

4.2. Information Retrieval

The field of Information Retrieval (IR) has matured considerably since its inception in the 1950s [363]. A taxonomy for IR models has been formalized [262]. It consists of four elements: a representation of documents in a collection; a representation of user queries; a framework describing relationships between document representations and queries; and a ranking function that quantifies an ordering amongst documents for a particular query. Three main issues motivate current IR research -- information relevance, query response time, and user interaction with IR systems. The dominant IR trends for searching large text collections are also threefold [262]. The size of collections is increasing dramatically. More complicated search mechanisms are being found to exploit document structure, to accommodate heterogeneous document collections, and to deal with document errors. Compression is in favour -- it may be quicker to search compact text or retrieve it from external devices. In a distributed IR system, query processing has four parts. Firstly, particular collections are targeted for the search. Secondly, queries are sent to the targeted collections. Queries are then evaluated at the individual collections. Finally, results from the collections are collated.

So how do P2P networks differ from distributed IR systems? Bawa, Manku, et al. presented four differences [62]. They suggested that a P2P network is typically larger, with tens or hundreds of thousands of nodes. It is usually more dynamic, with node lifetimes measured in hours. They suggested that a P2P network is usually homogeneous, with a common resource description language. It lacks the centralized "mediators" found in many IR systems that assume responsibility for selecting collections, for rewriting queries, and for merging ranked results. These distinctions are generally aligned with the peer characteristics in Section 1. One might add that P2P nodes display more symmetry -- peers are often both information consumers and producers. Daswani, Garcia-Molina, et al. pointed out that, while there are IR techniques for ranked keyword search at moderate scale, research is required so that ranking mechanisms are efficient at the larger scale targeted by P2P designs [25]. Joseph and Hoshiai surveyed several P2P systems using metadata techniques from the IR toolkit [60]. They described an assortment of IR techniques and P2P systems, including various metadata formats, retrieval models, bloom filters, DHTs, and trust issues.

In the ensuing paragraphs, we survey P2P work that has incorporated information retrieval models, particularly the Vector Model and the Latent Semantic Indexing Model. We omit the P2P work based on Bayesian models. Some have pointed to such work [60], but made no explicit mention of the model [364]. One early paper on P2P content-based image retrieval also leveraged the Bayesian model [365]. For the former two models, we briefly describe the design, then try to highlight robustness aspects. On robustness, we are again stymied for lack of prior work. Indeed, a search across all proceedings of the Annual ACM Conference on Research and Development in Information Retrieval for the words "reliable", "available", "dependable", or "adaptable" did not return any results at the time of writing. In contrast, a standard text on distributed database management systems [366] contains a whole chapter on reliability. IR research concentrates on performance measures. Common performance measures include recall, the fraction of the relevant documents that has been retrieved and precision, the fraction of the retrieved documents that is relevant [262]. Ideally, an IR system would have high recall and high precision. Unfortunately techniques favouring one often disadvantage the other [363].

4.2.1. Vector Model (PlanetP, FASD, eSearch)

The vector model [367] represents both documents and queries as term vectors, where a term could be a word or a phrase. If a document or query has a term, the weight of the corresponding dimension of the vector is non-zero. The similarity of the document and query vectors gives an indication of how well a document matches a particular query.

The weighting calculation is critical across the retrieval models. Amongst the numerous proposals for the probabilistic and vector models, there are some commonly recurring weighting factors [363]. One is term frequency. The more a term is repeated in a document, the more important the term is. Another is inverse document frequency. Terms common to many documents give less information about the content of a document. Then there is document length. Larger documents can bias term frequencies, so weightings are sometimes normalized against document length. The expression "TFIDF weighting" refers to the collection of weighting calculations that incorporate term frequency and inverse document frequency, not just to one. Two weighting calculations have been particularly dominant -- Okapi [368] and pivoted normalization [369]. A distributed version of Google's Pagerank algorithm has also been devised for a P2P environment [370]. It allows incremental, ongoing Pagerank calculations while documents are inserted and deleted.

A couple of early P2P systems leveraged the vector model. Building on the vector model, PlanetP divided the ranking problem into two steps [215]. In the first, peers are ranked for the probability that they have matching documents. In the second, higher-priority peers are contacted and the matching documents are ranked. An Inverse Peer Frequency, analogous to the Inverse Document Frequency, is used to rank relevant peers. To further constrain the query traffic, PlanetP contacts only the first group of m peers to retrieve a relevant set of documents. In this way, it repeatedly contacts groups of m peers until the top k document rankings are stable. While the PlanetP designers first quantified recall and precision, they also considered reliability. Each PlanetP peer has a global index with a list of all other peers, their IP addresses, and their Bloom filters. This large volume of shared information needs to be maintained. Klampanos and Jose saw this as PlanetP's primary shortcoming [371]. Each Bloom filter summarized the set of terms in the local index of each peer. The time to propagate changes, be they new documents or peer arrivals/departures, was studied by simulation for up to 1000 peers. The reported propagation times were in the hundreds of seconds. Design workarounds were required for PlanetP to be viable across slower dial-up modem connections. For future work, the authors were

considering some sort of hierarchy to scale to larger numbers of peers.

A second early system using the vector model is the Fault-tolerant, Adaptive, Scalable Distributed (FASD) search engine [283], which extended the Freenet design (Section 2.3) for richer queries. The original Freenet design could find a document based on a globally unique identifier. Kronfol's design added the ability to search, for example, for documents about "apples AND oranges NOT bananas". It uses a TFIDF weighting scheme to build a document's term vector. Each peer calculates the similarity of the query vector and local documents and forwards the query to the best downstream peer. Once the best downstream peer returns a result, the second-best peer is tried, and so on. Simulations with 1000 nodes gave an indication of the query path lengths in various situations -- when routing queries in a network with constant rates of node and document insertion, when bootstrapping the network in a "worst-case" ring topology, or when failing randomly and specifically selected peers. Kronfol claimed excellent average-case performance -- less than 20 hops to retrieve the same top n results as a centralized search engine. There were, however, numerous cases where the worst-case path length was several hundred hops in a network of only 1000 nodes.

In parallel, there have been some P2P designs based on the vector model from the University of Rochester -- pSearch [9, 372] and eSearch [373]. The early pSearch paper suggested a couple of retrieval models, one of which was the Vector Space Model, to search only the nodes likely to have matching documents. To obtain approximate global statistics for the TFIDF calculation, a spanning tree was constructed across a subset of the peers. For the m top terms, the term-to-document index was inserted into a Content-Addressable Network [334]. A variant that mapped terms to document clusters was also suggested. eSearch is a hybrid of the partition-by-document and partition-by-term approaches (Section 4.1.2) eSearch nodes are primarily partitioned by term. Each is responsible for the inverted lists for some top terms. For each document in the inverted list, the node stores the complete term list. To reduce the size of the index, the complete term lists for a document are only kept on nodes that are responsible for top terms in the document. eSearch uses the Okapi term weighting to select top terms. It relies on the Chord DHT [34] to associate terms with nodes storing the inverted lists. It also uses automatic query expansion. This takes the significant terms from the top document matches and automatically adds them to the user's query to find additional relevant documents. The eSearch performance was quantified in terms of search precision, the number of retrieved documents, and various load-balancing metrics. Compared to the more common proposals for partitioning by

keywords, eSearch consumed 6.8 times the storage space to achieve faster search times.

4.2.2. Latent Semantic Indexing (pSearch)

Another retrieval model used in P2P proposals is Latent Semantic Indexing (LSI) [374]. Its key idea is to map both the document and query vectors to a concept space with lower dimensions. The starting point is a $t \times N$ weighting matrix, where t is the total number of indexed terms, N is the total number of documents, and the matrix elements could be TFIDF rankings. Using singular value decomposition, this matrix is reduced to a smaller number of dimensions, while retaining the more significant term-to-document mappings. Baeza-Yates and Ribeiro-Neto suggested that LSI's value is a novel theoretic framework, but that its practical performance advantage for real document collections had yet to be proven [262]. pSearch incorporated LSI [9]. By placing the indices for semantically similar documents close in the network, Tang, Xu, et al. touted significant bandwidth savings relative to the early full-flooding variant of Gnutella [372]. They plotted the number of nodes visited by a query. They also explored the trade-off with accuracy, the percentage match between the documents returned by the distributed pSearch algorithm and those from a centralized LSI baseline. In a more recent update to the pSearch work, Tang, Dwarkadas, et al. summarized LSI's shortcomings [375]. Firstly, for large document collections, its retrieval quality is inherently inferior to Okapi. Secondly, singular value decomposition consumes excessive memory and computation time. Consequently, the authors used Okapi for searching while retaining LSI for indexing. With Okapi, they selected the next node to be searched and selected documents on searched nodes. With LSI, they ensured that similar documents are clustered near each other, thereby optimizing the network search costs. When retrieving a small number of top documents, the precision of LSI+Okapi approached that of Okapi. However, if retrieving a large number of documents, the LSI+Okapi precision is inferior. The authors want to improve this in future work.

4.2.3. Small Worlds

The "small world" concept originally described how people are interconnected by short chains of acquaintances [376]. Kleinberg was struck by the algorithmic lesson of the small world, namely "that individuals using local information are collectively very effective at constructing short paths between two points in a social network" [377]. Small world networks have a small diameter and a large clustering coefficient (a large number of connections amongst relevant nodes) [378].

The small world idea has had a limited impact on peer-to-peer algorithms. It has influenced only a few unstructured [62, 378-380] and structured [344, 381] algorithms. The most promising work on "small worlds" in P2P networks are those concerned with the information retrieval metrics, precision and recall [62, 378, 380].

5. Queries

Database research suggests directions for P2P research. Hellerstein observed that, while work on fast P2P indexes is well underway, P2P query optimization remains a promising topic for future research [23]. Kossman reviewed the state of the art of distributed query processing, highlighting areas for future research: simulation and query optimization for networks of tens of thousands of servers and millions of clients; non-relational data types (e.g., XML, text, and images); and partial query responses since on the Internet, "failure is the rule rather than the exception" [19]. A primary motivation for the P2P system, PIER, was to scale from the largest database systems of a few hundred nodes to an Internet environment in which there are over 160 million nodes [22]. Litwin and Sahri have also considered ways to combine distributed hashing, more specifically the Scalable Distributed Data Structures, with SQL databases, claiming to be first to implement scalable distributed database partitioning [382]. Motivated by the lack of transparent distribution in current distributed databases, they measure query execution times for Microsoft SQL servers aggregated by means of an SDDS layer. One of their starting assumptions was that it is too challenging to change the SQL query optimizer.

Database research also suggests the approach to P2P research. Researchers of database query optimization were divided between those looking for optimal solutions in special cases and those using heuristics to answer all queries [383]. Gribble, et al. cast query optimization in terms of the data placement problem, which is to "distribute data and work so the full query workload is answered with lowest cost under the existing bandwidth and resource constraints" [250]. They pointed out that even the static version of this problem is NP-complete in P2P networks. Consequently, research on massive, dynamic P2P networks will likely progress using both strategies of early database research - heuristics and special-case optimizations.

If P2P networks are going to be adaptable, if they are to support a wide range of applications, then they need to accommodate many query types [72]. Up to this point, we have reviewed queries for keys (Section 3) and keywords (Sections 4.1. and 4.2). Unfortunately, a major shortcoming of the DHTs in Section 3.5 is that they primarily support exact-match, single-key queries. Skip Graphs support range and prefix queries, but not aggregation queries. Here we probe below

the language syntax to identify the open research issues associated with more expressive P2P queries [25]. Triantafillou and Pitoura observed the disparate P2P designs for different types of queries and so outlined a unifying framework [76]. To classify queries, they considered the number of relations (single or multiple), the number of attributes (single or multiple), and the type of query operator. They described numerous operators: equality, range, join, and "special functions". The latter referred to aggregation (like sum, count, average, minimum, and maximum), grouping and ordering. The following sections approximately fit their taxonomy -- range queries, multi-attribute queries, join queries and aggregation queries. There has been some initial P2P work on other query types -- continuous queries [20, 22, 73], recursive queries [22, 74], and adaptive queries [23, 75]. For these, we defer to the primary references.

5.1. Range Queries

The support of efficient range predicates in P2P networks was identified as an important open research issue by Huebsch, et al. [22]. Range partitioning has been important in parallel databases to improve performance, so that a transaction commonly needs data from only one disk or node [22]. One type of range search, longest prefix match, is important because of its prevalence in routing schemes for voice and data networks alike. In other applications, users may pose broad, inexact queries, even though they require only a small number of responses. Consequently, techniques to locate similar ranges are also important [77]. Various proposals for range searches over P2P networks are summarized in Figure 4. Since the Scalable Distributed Data Structure (SDDS) has been an important influence on contemporary Distributed Hash Tables (DHTs) [49-51], we also include ongoing work on SDDS range searches.

PEER-TO-PEER (P2P)
Locality Sensitive Hashing (Chord) [77]
Prefix Hash Trees (unspecified DHT) [78, 79]
Space Filling Curves (CAN) [80]
Space Filling Curves (Chord) [81]
Quadrees (Chord) [82]
Skip Graphs [38, 41, 83, 100]
Mercury [84]
P-Grid [85, 86]

SCALABLE DISTRIBUTED DATA STRUCTURES (SDDS)
RP* [87, 88]

Figure 4: Solutions for Range Queries on P2P and SDDS Indexes

The papers on P2P range search can be divided into those that rely on an underlying DHT (the first five entries in Figure 4) and those that do not (the subsequent three entries). Bharambe, Agrawal, et al. argued that DHTs are inherently ill-suited to range queries [84]. The very feature that makes for their good load balancing properties, randomized hash functions, works against range queries. One possible solution would be to hash ranges, but this can require a priori partitioning. If the partitions are too large, partitions risk overload. If they are too small, there may be too many hops.

Despite these potential shortcomings, there have been several range query proposals based on DHTs. If hashing ranges to nodes, it is entirely possible that overlapping ranges map to different nodes. Gupta, Agrawal, et al. rely on locality sensitive hashing to ensure that, with high probability, similar ranges are mapped to the same node [77]. They propose one particular family of locality sensitive hash functions, called min-wise independent permutations. The number of partitions per node and the path length were plotted against the total numbers of peers in the system. For a network with 1000 nodes, the hop count distribution was very similar to that of the exact-matching Chord scheme. Was it load-balanced? For the same network with 50,000 partitions, there were over two orders of magnitude variation in the number of partitions at each node (first and ninety-ninth percentiles). The Prefix Hash Tree is a trie in which prefixes are hashed onto any DHT. The preliminary analysis suggests efficient doubly logarithmic lookup, balanced load, and fault resilience [78, 79]. Andrzejak and Xu were perhaps the first to propose a mapping from ranges to DHTs [80]. They use one particular Space Filling Curve, the Hilbert curve, over a Content Addressable Network (CAN) construction (Section 3.5.3). They maintain two properties: nearby ranges map to nearby CAN zones; if a range is split into two sub-ranges, then the zones of the sub-ranges partition the zone of the primary range. They plot path length and load proxy measures (the total number of messages and nodes visited) for three algorithms to propagate range queries: brute force, controlled flooding, and directed controlled flooding. Schmidt and Parashar also advocated Space Filling Curves to achieve range queries over a DHT [81]. However, they point out that, while Andrzejak and Xu use an inverse Space Filling Curve to map a one-dimensional space to d-dimensional zones, they map a d-dimensional space back to a one-dimensional index. Such a construction gives the ability to search across multiple attributes (Section 5.2). Tanin, Harwood, et al. suggested quadrees over Chord [82], and gave preliminary simulation results for query response times.

Because DHTs are naturally constrained to exact-match, single-key queries, researchers have considered other P2P indexes for range searches. Several were based on Skip Graphs [38, 41], which, unlike

the DHTs, do not necessitate randomizing hash functions and are therefore capable of range searches. Unfortunately, they are not load balanced [83]. For example, in SkipNet [48], hashing was added to balance the load -- the Skip Graph could support range searches or load balancing, but not both. One solution for load-balancing relies on an increased number of 'virtual' servers [168] but, in their search for a system that can both search for ranges and balance loads, Bharambe, Agrawal, et al. rejected the idea [84]. The virtual servers work assumed load imbalance stems from hashing; that is, by skewed data insertions and deletions. In some situations, the imbalance is triggered by a skewed query load. In such circumstances, additional virtual servers can increase the number of routing hops and increase the number of pointers that a Skip Graph needs to maintain. Ganesan, Bawa, et al. devised an alternate method to balance load [83]. They proposed two Skip Graphs, one to index the data itself and the other to track load at each node in the system. Each node is able to determine the load on its neighbours and the most (least) loaded nodes in the system. They devise two algorithms: NBRADJUST balances load on neighbouring nodes; using REORDER, empty nodes can take over some of the tuples on heavily loaded nodes. Their simulations focus on skewed storage load, rather than on skewed query loads, but they surmise that the same approach could be used for the latter.

Other proposals for range queries avoid both the DHT and the Skip Graph. Bharambe, Agrawal, et al. distinguish their Mercury design by its support for multi-attribute range queries and its explicit load balancing [84]. In Mercury, nodes are grouped into routing hubs, each of which is responsible for various query attributes. While it does not use hashing, Mercury is loosely similar to the DHT approaches: nodes within hubs are arranged into rings, like Chord [34]; for efficient routing within hubs, k long-distance links are used, like Symphony [381]. Range lookups require $O(((\log n)^2)/k)$ hops. Random sampling is used to estimate the average load on nodes and to find the parts of the overlay that are lightly loaded. Whereas Symphony assumed that nodes are responsible for ranges of approximately equal size, Mercury's random sampling can determine the location of the start of the range, even for non-uniform ranges [84]. P-Grid [42] does provide for range queries, by virtue of the key ordering in its tree structures. Ganesan, Bawa, et al. critiqued its capabilities [83]: P-Grid assumes fixed-capacity nodes; there was no formal characterization of imbalance ratios or balancing costs; every P-Grid periodically contacts other nodes for load information.

The work on Scalable Distributed Data Structures (SDDSs) has progressed in parallel with P2P work and has addressed range queries. Like the DHTs above, the early SDDS Linear Hashing (LH*) schemes were not order-preserving [52]. To facilitate range queries, Litwin,

Niemat, et al. devised a Range Partitioning variant, RP* [87]. There are options to dispense with the index, to add indexes to clients, and to add them to servers. In the variant without an index, every query is issued via multicasting. The other variants also use some multicasting. The initial RP* paper suggested scalability to thousands of sites, but a more recent RP* simulation was capped at 140 servers [88]. In that work, Tsangou, Ndiaye, et al. investigated TCP and UDP mechanisms by which servers could return range query results to clients. The primary metrics were search and response times. Amongst the commercial parallel database management systems, they reported that the largest seems only to scale to 32 servers (SQL Server 2000). For future work, they planned to explore aggregation of query results, rather than establishing a connection between the client and every single server with a response.

All in all, it seems there are numerous open research questions on P2P range queries. How realistic is the maintenance of global load statistics considering the scale and dynamism of P2P networks? Simulations at larger scales are required. Proposals should take into account both the storage load (insert and delete messages) and the query load (lookup messages). Simplifying assumptions need to be attacked. For example, how well do the above solutions work in networks with heterogeneous nodes, where the maximum message loads and index sizes are node-dependent?

5.2. Multi-Attribute Queries

There has been some work on multi-attribute P2P queries. As late as September 2003, it was suggested that there has not been an efficient solution [76].

Again, an early significant work on multi-attribute queries over aggregated commodity nodes germinated amongst SDDSs. k-RP* [89] uses the multi-dimensional binary search tree (or k-d tree, where k indicates the number of dimensions of the search index) [384]. It builds on the RP* work from the previous section and inherits their capabilities for range search and partial match. Like the other SDDSs, k-RP* indexes can fit into RAM for very fast lookup. For future work, Litwin and Neimat suggested a) a formal analysis of the range search termination algorithm and the k-d paging algorithm, b) a comparison with other multi-attribute data structures (quad-trees and R-trees) and c) exploration of query processing, concurrency control, and transaction management for k-RP* files [89]. On the latter point, others have considered transactions to be inconsequential to the core problem of supporting more complex queries in P2P networks [72].

In architecting their secure wide-area Service Discovery Service (SDS), Hodes, Czerwinski, et al. considered three possible designs for multi-criteria search -- Centralization, Mapping and Flooding [90]. These correlate to the index classifications of Section 2 -- Central, Distributed, and Local. They discounted the centralized, Napster-like index for its risk of a single point of failure. They considered the hash-based mappings of Section 3, but concluded that it would not be possible to adequately partition data. A document satisfying many criteria would be wastefully stored in many partitions. They rejected full flooding for its lack of scalability. Instead, they devised a query filtering technique, reminiscent of Gnutella's query routing protocol (Section 4.1). Nodes push proactive summaries of their data rather than waiting for a query. Summaries are aggregated and stored throughout a server hierarchy, to guide subsequent queries. Some initial prototype measurements were provided for total load on the system, but not for load distribution. They put several issues forward for future work. The indexing needs to be flexible to change according to query and storage workloads. A mesh topology might improve on their hierarchic topology since query misses would not propagate to root servers. The choice is analogous to BGP meshes and DNS trees.

More recently, Cai, Frank, et al. devised the Multi-Attribute Addressable Network (MAAN) [91]. They built on Chord to provide both multi-attribute and range queries, claiming to be the first to service both query types in a structured P2P system. Each MAAN node has $O(\log n)$ neighbours, where N is the number of nodes. MAAN multi-attribute range queries require $O(\log n + N \cdot S_{\min})$ hops, where S_{\min} is the minimum range selectivity across all attributes. Selectivity is the ratio of the query range to the entire identifier range. The paper assumed that a locality preserving hash function would ensure balanced load. Per Section 5.1, the arguments by Bharambe, Agrawal, et al. have highlighted the shortcomings of this assumption [84]. MAAN required that the schema must be fixed and known in advance -- adaptable schemas were recommended for subsequent attention. The authors also acknowledged that there is a selectivity breakpoint at which full flooding becomes more efficient than their scheme. This begs for a query resolution algorithm that adapts to the profile of queries. Cai and Frank followed up with RDFPeers [55]. They differentiate their work from other RDF proposals by a) guaranteeing to find query results if they exist and b) removing the requirement of prior definition of a fixed schema. They hashed <subject, predicate, object> triples onto the MAAN and reported routing hop metrics for their implementation. Load imbalance across nodes was reduced to less than one order of magnitude, but the specific measure was the number of triples stored per node - skewed query loads were not considered. They plan to improve load balancing with the virtual servers of Section 5.1 [168].

5.3. Join Queries

Two research teams have done some initial work on P2P join operations. Harren, Hellerstein, et al. initially described a three-layer architecture -- storage, DHT and query processing. They implemented the join operation by modifying an existing Content Addressable Network (CAN) simulator, reporting "significant hot-spots in all dimensions: storage, processing, and routing" [72]. They progressed their design more recently in the context of PIER, a distributed query engine based on CAN [22, 385]. They implemented two equi-join algorithms. In their design, a key is constructed from the "namespace" and the "resource ID". There is a namespace for each relation and the resource ID is the primary key for base tuples in that relation. Queries are multicast to all nodes in the two namespaces (relations) to be joined. Their first algorithm is a DHT version of the symmetric hash join. Each node in the two namespaces finds the relevant tuples and hashes them to a new query namespace. The resource ID in the new namespace is the concatenation of join attributes. In the second algorithm, called "fetch matches", one of the relations is already hashed on the join attributes. Each node in the second namespace finds tuples matching the query and retrieves the corresponding tuples from the first relation. They leveraged two other techniques, namely the symmetric semi-join rewrite and the Bloom filter rewrite, to reduce the high bandwidth overheads of the symmetric hash join. For an overlay of 10,000 nodes, they simulated the delay to retrieve tuples and the aggregate network bandwidth for these four schemes. The initial prototype was on a cluster of 64 PCs, but it has more recently been expanded to PlanetLab.

Triantafillou and Pitoura considered multicasting to large numbers of peers to be inefficient [76]. They therefore allocated a limited number of special peers, called range guards. The domain of the join attributes was divided, one partition per range guard. Join queries were sent only to range guards, where the query was executed. Efficient selection of range guards and a quantitative evaluation of their proposal were left for future work.

5.4. Aggregation Queries

Aggregation queries invariably rely on tree-structures to combine results from a large number of nodes. Examples of aggregation queries are Count, Sum, Maximum, Minimum, Average, Median, and Top-K [92, 386, 387]. Figure 5 summarizes the tree and query characteristics that affect dependability.

Tree type: Doesn't use DHT [92], use internal DHT trees [95], use independent trees on top of DHTs
 Tree repair: Periodic [93], exceptional [32]
 Tree count: One per key, one per overlay [56]
 Tree flexibility: Static [92], dynamic

 Query interface: install, update, probe [98]
 Query distribution: multicast [98], gossip [92]
 Query applications: leader election, voting, resource location, object placement and error recovery [98, 388]
 Query semantics
 Consistency: Best-effort, eventual [92], snapshot / interval / single-site validity [99]
 Timeliness [388]
 Lifetime: Continuous [97, 99], single-shot
 No. attributes: Single, multiple
 Query types: Count, sum, maximum, minimum, average, median, top k [92, 386, 387]

Figure 5: Aggregation Trees and Queries in P2P Networks

Key: Astrolabe [92]; Cone [93]; Distributed Approximative System Information Service (DASIS) [95]; Scalable Distributed Information

Management System (SDIMS) [98]; Self-Organized Metadata Overlay (SOMO) [56]; Wildfire [99]; Willow [32]; Newscast [97]

The fundamental design choices for aggregation trees relate to how the overlay uses DHTs, how it repairs itself when there are failures, how many aggregation trees there are, and whether the tree is static or dynamic (Figure 5). Astrolabe is one of the most influential P2P designs included in Figure 5, yet it makes no use of DHTs [92]. Other designs make use of the internal trees of Plaxton-like DHTs. Others build independent tree structures on top of DHTs. Most of the designs repair the aggregation tree with periodic mechanisms similar to those used in the DHTs themselves. Willow is an exception [32]. It uses a Tree Maintenance Protocol to "zip" disjoint aggregation trees together when there are major failures. Yalagandula and Dahlin found reconfigurations at the aggregation layer to be costly, suggesting more research on techniques to reduce the cost and frequency of such reconfigurations [98]. Many of the designs use multiple aggregation trees, each rooted at the DHT node responsible for the aggregation attribute. On the other hand, the Self-Organized Metadata Overlay [56] uses a single tree and is vulnerable to a single point of failure at its root.

At the time of writing, researchers have just begun exploring the performance of queries in the presence of churn. Most designs are for best-effort queries. Bawa, et al. devised a better consistency model, called Single-Site Validity [99] to qualify the accuracy of results when there is churn. Its price was a five-fold increase in the message load, when compared to an efficient but best-effort Spanning Tree. Gossip mechanisms are resilient to churn, but they delay aggregation results and incur high message cost for aggregation attributes with small read-to-write ratios.

6. Security Considerations

An initial list of references to research on P2P security is given in Figure 1, Section 1. This document addresses P2P search. P2P storage, security, and applications are recommended for further investigation in Section 8.

7. Conclusions

Research on peer-to-peer networks can be divided into four categories -- search, storage, security and applications. This critical survey has focused on search methods. While P2P networks have been classified by the existence of an index (structured or unstructured) or the location of the index (local, centralized, and distributed), this survey has shown that most have evolved to have some structure, whether it is indexes at superpeers or indexes defined by DHT algorithms. As for location, the distributed index is most common. The survey has characterized indexes as semantic and semantic-free. It has also critiqued P2P work on major query types. While much of it addresses work from 2000 or later, we have traced important building blocks from the 1990s.

The initial motivation in this survey was to answer the question, "How robust are P2P search networks?" The question is key to the deployment of P2P technology. Balakrishnan, Kaashoek, et al. argued that the P2P architecture is appealing: the startup and growth barriers are low; they can aggregate enormous storage and processing resources; "the decentralized and distributed nature of P2P systems gives them the potential to be robust to faults or intentional attacks" [18]. If P2P is to be a disruptive technology in applications other than casual file sharing, then robustness needs to be practically verified [20].

The best comparative research on P2P dependability has been done in the context of Distributed Hash Tables (DHTs) [291]. The entire body of DHT research can be distilled to four main observations about dependability (Section 3.2). Firstly, static dependability comparisons show that no $O(\log n)$ DHT geometry is significantly more

dependable than the other $O(\log n)$ geometries. Secondly, dynamic dependability comparisons show that DHT dependability is sensitive to the underlying topology maintenance algorithms (Figure 2). Thirdly, most DHTs use $O(\log n)$ geometries to suit ephemeral nodes, whereas the $O(1)$ hop DHTs suit stable nodes - they deserve more research attention. Fourthly, although not yet a mature science, the study of DHT dependability is helped by recent simulation tools that support multiple DHTs [299].

We make the following four suggestions for future P2P research:

- 1) Complete the companion P2P surveys for storage, security, and applications. A rough outline has been suggested in Figure 1, along with references. The need for such surveys was highlighted within the peer-to-peer research group of the Internet Research Task Force (IRTF) [17].
- 2) P2P indexes are maturing. P2P queries are embryonic. Work on more expressive queries over P2P indexes started to gain momentum in 2003, but remains fraught with efficiency and load issues.
- 3) Isolate the low-level mechanisms affecting robustness. There is limited value in comparing robustness of DHT geometries (like rings versus de Bruijn graphs), when robustness is highly sensitive to underlying topology maintenance algorithms (Figure 2).
- 4) Build consensus on robustness metrics and their acceptable ranges. This paper has teased out numerous measures that impinge on robustness, for example, the median query path length for a failure of $x\%$ of nodes, bisection width, path overlap, the number of alternatives available for the next hop, lookup latency, average live bandwidth (bytes/node/sec), successful routing rates, the number of timeouts (caused by a finger pointing to a departed node), lookup failure rates (caused by nodes that temporarily point to the wrong successor during churn), and clustering measures (edge expansion and node expansion). Application-level robustness metrics need to drive a consistent assessment of the underlying search mechanics.

8. Acknowledgments

This document was adapted from a paper in Elsevier's Computer Networks:

J. Risson & T. Moors, Survey of Research towards Robust Peer-to-Peer Networks: Search Methods, Computer Networks 51(7)2007.

We thank Bill Yeager, Ali Ghodsi, and several anonymous reviewers for thorough comments that significantly improved the quality of earlier versions of this document.

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