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## Diameter Base Protocol

### Status of this Memo

This document specifies an Internet standards track protocol for the Internet community, and requests discussion and suggestions for improvements. Please refer to the current edition of the "Internet Official Protocol Standards" (STD 1) for the standardization state and status of this protocol. Distribution of this memo is unlimited.

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### Abstract

The Diameter base protocol is intended to provide an Authentication, Authorization and Accounting (AAA) framework for applications such as network access or IP mobility. Diameter is also intended to work in both local Authentication, Authorization & Accounting and roaming situations. This document specifies the message format, transport, error reporting, accounting and security services to be used by all Diameter applications. The Diameter base application needs to be supported by all Diameter implementations.

### Conventions Used In This Document

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14, RFC 2119 [KEYWORD].

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## 1. Introduction

Authentication, Authorization and Accounting (AAA) protocols such as TACACS [TACACS] and RADIUS [RADIUS] were initially deployed to provide dial-up PPP [PPP] and terminal server access. Over time, with the growth of the Internet and the introduction of new access technologies, including wireless, DSL, Mobile IP and Ethernet, routers and network access servers (NAS) have increased in complexity and density, putting new demands on AAA protocols.

Network access requirements for AAA protocols are summarized in [AAAREQ]. These include:

### Failover

[RADIUS] does not define failover mechanisms, and as a result, failover behavior differs between implementations. In order to provide well defined failover behavior, Diameter supports application-layer acknowledgements, and defines failover algorithms and the associated state machine. This is described in Section 5.5 and [AAATRANS].

### Transmission-level security

[RADIUS] defines an application-layer authentication and integrity scheme that is required only for use with Response packets. While [RADEXT] defines an additional authentication and integrity mechanism, use is only required during Extensible Authentication Protocol (EAP) sessions. While attribute-hiding is supported, [RADIUS] does not provide support for per-packet confidentiality. In accounting, [RADACCT] assumes that replay protection is provided by the backend billing server, rather than within the protocol itself.

While [RFC3162] defines the use of IPsec with RADIUS, support for IPsec is not required. Since within [IKE] authentication occurs only within Phase 1 prior to the establishment of IPsec SAs in Phase 2, it is typically not possible to define separate trust or authorization schemes for each application. This limits the usefulness of IPsec in inter-domain AAA applications (such as roaming) where it may be desirable to define a distinct certificate hierarchy for use in a AAA deployment. In order to provide universal support for transmission-level security, and enable both intra- and inter-domain AAA deployments, IPsec support is mandatory in Diameter, and TLS support is optional. Security is discussed in Section 13.

#### Reliable transport

RADIUS runs over UDP, and does not define retransmission behavior; as a result, reliability varies between implementations. As described in [ACCMGMT], this is a major issue in accounting, where packet loss may translate directly into revenue loss. In order to provide well defined transport behavior, Diameter runs over reliable transport mechanisms (TCP, SCTP) as defined in [AAATrans].

#### Agent support

[RADIUS] does not provide for explicit support for agents, including Proxies, Redirects and Relays. Since the expected behavior is not defined, it varies between implementations. Diameter defines agent behavior explicitly; this is described in Section 2.8.

#### Server-initiated messages

While RADIUS server-initiated messages are defined in [DYNAUTH], support is optional. This makes it difficult to implement features such as unsolicited disconnect or reauthentication/reauthorization on demand across a heterogeneous deployment. Support for server-initiated messages is mandatory in Diameter, and is described in Section 8.

#### Auditability

RADIUS does not define data-object security mechanisms, and as a result, untrusted proxies may modify attributes or even packet headers without being detected. Combined with lack of support for capabilities negotiation, this makes it very difficult to determine what occurred in the event of a dispute. While implementation of data object security is not mandatory within Diameter, these capabilities are supported, and are described in [AAACMS].

#### Transition support

While Diameter does not share a common protocol data unit (PDU) with RADIUS, considerable effort has been expended in enabling backward compatibility with RADIUS, so that the two protocols may be deployed in the same network. Initially, it is expected that Diameter will be deployed within new network devices, as well as within gateways enabling communication between legacy RADIUS devices and Diameter agents. This capability, described in [NASREQ], enables Diameter support to be added to legacy networks, by addition of a gateway or server speaking both RADIUS and Diameter.

In addition to addressing the above requirements, Diameter also provides support for the following:

#### Capability negotiation

RADIUS does not support error messages, capability negotiation, or a mandatory/non-mandatory flag for attributes. Since RADIUS clients and servers are not aware of each other's capabilities, they may not be able to successfully negotiate a mutually acceptable service, or in some cases, even be aware of what service has been implemented. Diameter includes support for error handling (Section 7), capability negotiation (Section 5.3), and mandatory/non-mandatory attribute-value pairs (AVPs) (Section 4.1).

#### Peer discovery and configuration

RADIUS implementations typically require that the name or address of servers or clients be manually configured, along with the corresponding shared secrets. This results in a large administrative burden, and creates the temptation to reuse the RADIUS shared secret, which can result in major security vulnerabilities if the Request Authenticator is not globally and temporally unique as required in [RADIUS]. Through DNS, Diameter enables dynamic discovery of peers. Derivation of dynamic session keys is enabled via transmission-level security.

#### Roaming support

The ROAMOPS WG provided a survey of roaming implementations [ROAMREV], detailed roaming requirements [ROAMCRIT], defined the Network Access Identifier (NAI) [NAI], and documented existing implementations (and imitations) of RADIUS-based roaming [PROXYCHAIN]. In order to improve scalability, [PROXYCHAIN] introduced the concept of proxy chaining via an intermediate server, facilitating roaming between providers. However, since RADIUS does not provide explicit support for proxies, and lacks auditability and transmission-level security features, RADIUS-based roaming is vulnerable to attack from external parties as well as susceptible to fraud perpetrated by the roaming partners themselves. As a result, it is not suitable for wide-scale deployment on the Internet [PROXYCHAIN]. By providing explicit support for inter-domain roaming and message routing (Sections 2.7 and 6), auditability [AAACMS], and transmission-layer security (Section 13) features, Diameter addresses these limitations and provides for secure and scalable roaming.

In the decade since AAA protocols were first introduced, the capabilities of Network Access Server (NAS) devices have increased substantially. As a result, while Diameter is a considerably more sophisticated protocol than RADIUS, it remains feasible to implement



within embedded devices, given improvements in processor speeds and the widespread availability of embedded IPsec and TLS implementations.

### 1.1. Diameter Protocol

The Diameter base protocol provides the following facilities:

- Delivery of AVPs (attribute value pairs)
- Capabilities negotiation
- Error notification
- Extensibility, through addition of new commands and AVPs (required in [AAAREQ]).
- Basic services necessary for applications, such as handling of user sessions or accounting

All data delivered by the protocol is in the form of an AVP. Some of these AVP values are used by the Diameter protocol itself, while others deliver data associated with particular applications that employ Diameter. AVPs may be added arbitrarily to Diameter messages, so long as the required AVPs are included and AVPs that are explicitly excluded are not included. AVPs are used by the base Diameter protocol to support the following required features:

- Transporting of user authentication information, for the purposes of enabling the Diameter server to authenticate the user.
- Transporting of service specific authorization information, between client and servers, allowing the peers to decide whether a user's access request should be granted.
- Exchanging resource usage information, which MAY be used for accounting purposes, capacity planning, etc.
- Relaying, proxying and redirecting of Diameter messages through a server hierarchy.

The Diameter base protocol provides the minimum requirements needed for a AAA protocol, as required by [AAAREQ]. The base protocol may be used by itself for accounting purposes only, or it may be used with a Diameter application, such as Mobile IPv4 [DIAMMIP], or network access [NASREQ]. It is also possible for the base protocol to be extended for use in new applications, via the addition of new commands or AVPs. At this time the focus of Diameter is network access and accounting applications. A truly generic AAA protocol used by many applications might provide functionality not provided by Diameter. Therefore, it is imperative that the designers of new applications understand their requirements before using Diameter.

See Section 2.4 for more information on Diameter applications.

Any node can initiate a request. In that sense, Diameter is a peer-to-peer protocol. In this document, a Diameter Client is a device at the edge of the network that performs access control, such as a Network Access Server (NAS) or a Foreign Agent (FA). A Diameter client generates Diameter messages to request authentication, authorization, and accounting services for the user. A Diameter agent is a node that does not authenticate and/or authorize messages locally; agents include proxies, redirects and relay agents. A Diameter server performs authentication and/or authorization of the user. A Diameter node MAY act as an agent for certain requests while acting as a server for others.

The Diameter protocol also supports server-initiated messages, such as a request to abort service to a particular user.

#### 1.1.1. Description of the Document Set

Currently, the Diameter specification consists of a base specification (this document), Transport Profile [AAATrans] and applications: Mobile IPv4 [DIAMMIP], and NASREQ [NASREQ].

The Transport Profile document [AAATrans] discusses transport layer issues that arise with AAA protocols and recommendations on how to overcome these issues. This document also defines the Diameter failover algorithm and state machine.

The Mobile IPv4 [DIAMMIP] application defines a Diameter application that allows a Diameter server to perform AAA functions for Mobile IPv4 services to a mobile node.

The NASREQ [NASREQ] application defines a Diameter Application that allows a Diameter server to be used in a PPP/SLIP Dial-Up and Terminal Server Access environment. Consideration was given for servers that need to perform protocol conversion between Diameter and RADIUS.

In summary, this document defines the base protocol specification for AAA, which includes support for accounting. The Mobile IPv4 and the NASREQ documents describe applications that use this base specification for Authentication, Authorization and Accounting.

## 1.2. Approach to Extensibility

The Diameter protocol is designed to be extensible, using several mechanisms, including:

- Defining new AVP values
- Creating new AVPs
- Creating new authentication/authorization applications
- Creating new accounting applications
- Application authentication procedures

Reuse of existing AVP values, AVPs and Diameter applications are strongly recommended. Reuse simplifies standardization and implementation and avoids potential interoperability issues. It is expected that command codes are reused; new command codes can only be created by IETF Consensus (see Section 11.2.1).

### 1.2.1. Defining New AVP Values

New applications should attempt to reuse AVPs defined in existing applications when possible, as opposed to creating new AVPs. For AVPs of type Enumerated, an application may require a new value to communicate some service-specific information.

In order to allocate a new AVP value, a request **MUST** be sent to IANA [IANA], along with an explanation of the new AVP value. IANA considerations for Diameter are discussed in Section 11.

### 1.2.2. Creating New AVPs

When no existing AVP can be used, a new AVP should be created. The new AVP being defined **MUST** use one of the data types listed in Section 4.2.

In the event that a logical grouping of AVPs is necessary, and multiple "groups" are possible in a given command, it is recommended that a Grouped AVP be used (see Section 4.4).

In order to create a new AVP, a request **MUST** be sent to IANA, with a specification for the AVP. The request **MUST** include the commands that would make use of the AVP.

### 1.2.3. Creating New Authentication Applications

Every Diameter application specification **MUST** have an IANA assigned Application Identifier (see Section 2.4) or a vendor specific Application Identifier.

Should a new Diameter usage scenario find itself unable to fit within an existing application without requiring major changes to the specification, it may be desirable to create a new Diameter application. Major changes to an application include:

- Adding new AVPs to the command, which have the "M" bit set.
- Requiring a command that has a different number of round trips to satisfy a request (e.g., application foo has a command that requires one round trip, but new application bar has a command that requires two round trips to complete).
- Adding support for an authentication method requiring definition of new AVPs for use with the application. Since a new EAP authentication method can be supported within Diameter without requiring new AVPs, addition of EAP methods does not require the creation of a new authentication application.

Creation of a new application should be viewed as a last resort. An implementation MAY add arbitrary non-mandatory AVPs to any command defined in an application, including vendor-specific AVPs without needing to define a new application. Please refer to Section 11.1.1 for details.

In order to justify allocation of a new application identifier, Diameter applications MUST define one Command Code, or add new mandatory AVPs to the ABNF.

The expected AVPs MUST be defined in an ABNF [ABNF] grammar (see Section 3.2). If the Diameter application has accounting requirements, it MUST also specify the AVPs that are to be present in the Diameter Accounting messages (see Section 9.3). However, just because a new authentication application id is required, does not imply that a new accounting application id is required.

When possible, a new Diameter application SHOULD reuse existing Diameter AVPs, in order to avoid defining multiple AVPs that carry similar information.

#### 1.2.4. Creating New Accounting Applications

There are services that only require Diameter accounting. Such services need to define the AVPs carried in the Accounting-Request (ACR)/ Accounting-Answer (ACA) messages, but do not need to define new command codes. An implementation MAY add arbitrary non-mandatory AVPs (AVPs with the "M" bit not set) to any command defined in an

application, including vendor-specific AVPs, without needing to define a new accounting application. Please refer to Section 11.1.1 for details.

Application Identifiers are still required for Diameter capability exchange. Every Diameter accounting application specification **MUST** have an IANA assigned Application Identifier (see Section 2.4) or a vendor specific Application Identifier.

Every Diameter implementation **MUST** support accounting. Basic accounting support is sufficient to handle any application that uses the ACR/ACA commands defined in this document, as long as no new mandatory AVPs are added. A mandatory AVP is defined as one which has the "M" bit set when sent within an accounting command, regardless of whether it is required or optional within the ABNF for the accounting application.

The creation of a new accounting application should be viewed as a last resort and **MUST NOT** be used unless a new command or additional mechanisms (e.g., application defined state machine) is defined within the application, or new mandatory AVPs are added to the ABNF.

Within an accounting command, setting the "M" bit implies that a backend server (e.g., billing server) or the accounting server itself **MUST** understand the AVP in order to compute a correct bill. If the AVP is not relevant to the billing process, when the AVP is included within an accounting command, it **MUST NOT** have the "M" bit set, even if the "M" bit is set when the same AVP is used within other Diameter commands (i.e., authentication/authorization commands).

A DIAMETER base accounting implementation **MUST** be configurable to advertise supported accounting applications in order to prevent the accounting server from accepting accounting requests for unbillable services. The combination of the home domain and the accounting application Id can be used in order to route the request to the appropriate accounting server.

When possible, a new Diameter accounting application **SHOULD** attempt to reuse existing AVPs, in order to avoid defining multiple AVPs that carry similar information.

If the base accounting is used without any mandatory AVPs, new commands or additional mechanisms (e.g., application defined state machine), then the base protocol defined standard accounting application Id (Section 2.4) **MUST** be used in ACR/ACA commands.

### 1.2.5. Application Authentication Procedures

When possible, applications SHOULD be designed such that new authentication methods MAY be added without requiring changes to the application. This MAY require that new AVP values be assigned to represent the new authentication transform, or any other scheme that produces similar results. When possible, authentication frameworks, such as Extensible Authentication Protocol [EAP], SHOULD be used.

### 1.3. Terminology

#### AAA

Authentication, Authorization and Accounting.

#### Accounting

The act of collecting information on resource usage for the purpose of capacity planning, auditing, billing or cost allocation.

#### Accounting Record

An accounting record represents a summary of the resource consumption of a user over the entire session. Accounting servers creating the accounting record may do so by processing interim accounting events or accounting events from several devices serving the same user.

#### Authentication

The act of verifying the identity of an entity (subject).

#### Authorization

The act of determining whether a requesting entity (subject) will be allowed access to a resource (object).

#### AVP

The Diameter protocol consists of a header followed by one or more Attribute-Value-Pairs (AVPs). An AVP includes a header and is used to encapsulate protocol-specific data (e.g., routing information) as well as authentication, authorization or accounting information.

#### Broker

A broker is a business term commonly used in AAA infrastructures. A broker is either a relay, proxy or redirect agent, and MAY be operated by roaming consortiums. Depending on the business model, a broker may either choose to deploy relay agents or proxy agents.

#### Diameter Agent

A Diameter Agent is a Diameter node that provides either relay, proxy, redirect or translation services.

#### Diameter Client

A Diameter Client is a device at the edge of the network that performs access control. An example of a Diameter client is a Network Access Server (NAS) or a Foreign Agent (FA).

#### Diameter Node

A Diameter node is a host process that implements the Diameter protocol, and acts either as a Client, Agent or Server.

#### Diameter Peer

A Diameter Peer is a Diameter Node to which a given Diameter Node has a direct transport connection.

#### Diameter Security Exchange

A Diameter Security Exchange is a process through which two Diameter nodes establish end-to-end security.

#### Diameter Server

A Diameter Server is one that handles authentication, authorization and accounting requests for a particular realm. By its very nature, a Diameter Server MUST support Diameter applications in addition to the base protocol.

#### Downstream

Downstream is used to identify the direction of a particular Diameter message from the home server towards the access device.

#### End-to-End Security

TLS and IPsec provide hop-by-hop security, or security across a transport connection. When relays or proxy are involved, this hop-by-hop security does not protect the entire Diameter user session. End-to-end security is security between two Diameter nodes, possibly communicating through Diameter Agents. This security protects the entire Diameter communications path from the originating Diameter node to the terminating Diameter node.

#### Home Realm

A Home Realm is the administrative domain with which the user maintains an account relationship.

#### Home Server

See Diameter Server.

#### Interim accounting

An interim accounting message provides a snapshot of usage during a user's session. It is typically implemented in order to provide for partial accounting of a user's session in the case of a device reboot or other network problem prevents the reception of a session summary message or session record.

#### Local Realm

A local realm is the administrative domain providing services to a user. An administrative domain MAY act as a local realm for certain users, while being a home realm for others.

#### Multi-session

A multi-session represents a logical linking of several sessions. Multi-sessions are tracked by using the Acct-Multi-Session-Id. An example of a multi-session would be a Multi-link PPP bundle. Each leg of the bundle would be a session while the entire bundle would be a multi-session.

#### Network Access Identifier

The Network Access Identifier, or NAI [NAI], is used in the Diameter protocol to extract a user's identity and realm. The identity is used to identify the user during authentication and/or authorization, while the realm is used for message routing purposes.

#### Proxy Agent or Proxy

In addition to forwarding requests and responses, proxies make policy decisions relating to resource usage and provisioning. This is typically accomplished by tracking the state of NAS devices. While proxies typically do not respond to client Requests prior to receiving a Response from the server, they may originate Reject messages in cases where policies are violated. As a result, proxies need to understand the semantics of the messages passing through them, and may not support all Diameter applications.

#### Realm

The string in the NAI that immediately follows the '@' character. NAI realm names are required to be unique, and are piggybacked on the administration of the DNS namespace. Diameter makes use of the realm, also loosely referred to as domain, to determine whether messages can be satisfied locally, or whether they must be routed or redirected. In RADIUS, realm names are not necessarily piggybacked on the DNS namespace but may be independent of it.



#### Real-time Accounting

Real-time accounting involves the processing of information on resource usage within a defined time window. Time constraints are typically imposed in order to limit financial risk.

#### Relay Agent or Relay

Relays forward requests and responses based on routing-related AVPs and realm routing table entries. Since relays do not make policy decisions, they do not examine or alter non-routing AVPs. As a result, relays never originate messages, do not need to understand the semantics of messages or non-routing AVPs, and are capable of handling any Diameter application or message type. Since relays make decisions based on information in routing AVPs and realm forwarding tables they do not keep state on NAS resource usage or sessions in progress.

#### Redirect Agent

Rather than forwarding requests and responses between clients and servers, redirect agents refer clients to servers and allow them to communicate directly. Since redirect agents do not sit in the forwarding path, they do not alter any AVPs transiting between client and server. Redirect agents do not originate messages and are capable of handling any message type, although they may be configured only to redirect messages of certain types, while acting as relay or proxy agents for other types. As with proxy agents, redirect agents do not keep state with respect to sessions or NAS resources.

#### Roaming Relationships

Roaming relationships include relationships between companies and ISPs, relationships among peer ISPs within a roaming consortium, and relationships between an ISP and a roaming consortium.

#### Security Association

A security association is an association between two endpoints in a Diameter session which allows the endpoints to communicate with integrity and confidentiality, even in the presence of relays and/or proxies.

#### Session

A session is a related progression of events devoted to a particular activity. Each application SHOULD provide guidelines as to when a session begins and ends. All Diameter packets with the same Session-Identifier are considered to be part of the same session.

#### Session state

A stateful agent is one that maintains session state information, by keeping track of all authorized active sessions. Each authorized session is bound to a particular service, and its state is considered active either until it is notified otherwise, or by expiration.

#### Sub-session

A sub-session represents a distinct service (e.g., QoS or data characteristics) provided to a given session. These services may happen concurrently (e.g., simultaneous voice and data transfer during the same session) or serially. These changes in sessions are tracked with the Accounting-Sub-Session-Id.

#### Transaction state

The Diameter protocol requires that agents maintain transaction state, which is used for failover purposes. Transaction state implies that upon forwarding a request, the Hop-by-Hop identifier is saved; the field is replaced with a locally unique identifier, which is restored to its original value when the corresponding answer is received. The request's state is released upon receipt of the answer. A stateless agent is one that only maintains transaction state.

#### Translation Agent

A translation agent is a stateful Diameter node that performs protocol translation between Diameter and another AAA protocol, such as RADIUS.

#### Transport Connection

A transport connection is a TCP or SCTP connection existing directly between two Diameter peers, otherwise known as a Peer-to-Peer Connection.

#### Upstream

Upstream is used to identify the direction of a particular Diameter message from the access device towards the home server.

#### User

The entity requesting or using some resource, in support of which a Diameter client has generated a request.

## 2. Protocol Overview

The base Diameter protocol may be used by itself for accounting applications, but for use in authentication and authorization it is always extended for a particular application. Two Diameter applications are defined by companion documents: NASREQ [NASREQ],

Mobile IPv4 [DIAMMIP]. These applications are introduced in this document but specified elsewhere. Additional Diameter applications MAY be defined in the future (see Section 11.3).

Diameter Clients MUST support the base protocol, which includes accounting. In addition, they MUST fully support each Diameter application that is needed to implement the client's service, e.g., NASREQ and/or Mobile IPv4. A Diameter Client that does not support both NASREQ and Mobile IPv4, MUST be referred to as "Diameter X Client" where X is the application which it supports, and not a "Diameter Client".

Diameter Servers MUST support the base protocol, which includes accounting. In addition, they MUST fully support each Diameter application that is needed to implement the intended service, e.g., NASREQ and/or Mobile IPv4. A Diameter Server that does not support both NASREQ and Mobile IPv4, MUST be referred to as "Diameter X Server" where X is the application which it supports, and not a "Diameter Server".

Diameter Relays and redirect agents are, by definition, protocol transparent, and MUST transparently support the Diameter base protocol, which includes accounting, and all Diameter applications.

Diameter proxies MUST support the base protocol, which includes accounting. In addition, they MUST fully support each Diameter application that is needed to implement proxied services, e.g., NASREQ and/or Mobile IPv4. A Diameter proxy which does not support also both NASREQ and Mobile IPv4, MUST be referred to as "Diameter X Proxy" where X is the application which it supports, and not a "Diameter Proxy".

The base Diameter protocol concerns itself with capabilities negotiation, how messages are sent and how peers may eventually be abandoned. The base protocol also defines certain rules that apply to all exchanges of messages between Diameter nodes.

Communication between Diameter peers begins with one peer sending a message to another Diameter peer. The set of AVPs included in the message is determined by a particular Diameter application. One AVP that is included to reference a user's session is the Session-Id.

The initial request for authentication and/or authorization of a user would include the Session-Id. The Session-Id is then used in all subsequent messages to identify the user's session (see Section 8 for more information). The communicating party may accept the request, or reject it by returning an answer message with the Result-Code AVP

set to indicate an error occurred. The specific behavior of the Diameter server or client receiving a request depends on the Diameter application employed.

Session state (associated with a Session-Id) MUST be freed upon receipt of the Session-Termination-Request, Session-Termination-Answer, expiration of authorized service time in the Session-Timeout AVP, and according to rules established in a particular Diameter application.

## 2.1. Transport

Transport profile is defined in [AAATrans].

The base Diameter protocol is run on port 3868 of both TCP [TCP] and SCTP [SCTP] transport protocols.

Diameter clients MUST support either TCP or SCTP, while agents and servers MUST support both. Future versions of this specification MAY mandate that clients support SCTP.

A Diameter node MAY initiate connections from a source port other than the one that it declares it accepts incoming connections on, and MUST be prepared to receive connections on port 3868. A given Diameter instance of the peer state machine MUST NOT use more than one transport connection to communicate with a given peer, unless multiple instances exist on the peer in which case a separate connection per process is allowed.

When no transport connection exists with a peer, an attempt to connect SHOULD be periodically made. This behavior is handled via the Tc timer, whose recommended value is 30 seconds. There are certain exceptions to this rule, such as when a peer has terminated the transport connection stating that it does not wish to communicate.

When connecting to a peer and either zero or more transports are specified, SCTP SHOULD be tried first, followed by TCP. See Section 5.2 for more information on peer discovery.

Diameter implementations SHOULD be able to interpret ICMP protocol port unreachable messages as explicit indications that the server is not reachable, subject to security policy on trusting such messages. Diameter implementations SHOULD also be able to interpret a reset from the transport and timed-out connection attempts.

If Diameter receives data up from TCP that cannot be parsed or identified as a Diameter error made by the peer, the stream is compromised and cannot be recovered. The transport connection MUST be closed using a RESET call (send a TCP RST bit) or an SCTP ABORT message (graceful closure is compromised).

#### 2.1.1. SCTP Guidelines

The following are guidelines for Diameter implementations that support SCTP:

1. For interoperability: All Diameter nodes MUST be prepared to receive Diameter messages on any SCTP stream in the association.
2. To prevent blocking: All Diameter nodes SHOULD utilize all SCTP streams available to the association to prevent head-of-the-line blocking.

#### 2.2. Securing Diameter Messages

Diameter clients, such as Network Access Servers (NASes) and Mobility Agents MUST support IP Security [SECARCH], and MAY support TLS [TLS]. Diameter servers MUST support TLS and IPsec. The Diameter protocol MUST NOT be used without any security mechanism (TLS or IPsec).

It is suggested that IPsec can be used primarily at the edges and in intra-domain traffic, such as using pre-shared keys between a NAS a local AAA proxy. This also eases the requirements on the NAS to support certificates. It is also suggested that inter-domain traffic would primarily use TLS. See Sections 13.1 and 13.2 for more details on IPsec and TLS usage.

#### 2.3. Diameter Application Compliance

Application Identifiers are advertised during the capabilities exchange phase (see Section 5.3). For a given application, advertising support of an application implies that the sender supports all command codes, and the AVPs specified in the associated ABNFs, described in the specification.

An implementation MAY add arbitrary non-mandatory AVPs to any command defined in an application, including vendor-specific AVPs. Please refer to Section 11.1.1 for details.

## 2.4. Application Identifiers

Each Diameter application MUST have an IANA assigned Application Identifier (see Section 11.3). The base protocol does not require an Application Identifier since its support is mandatory. During the capabilities exchange, Diameter nodes inform their peers of locally supported applications. Furthermore, all Diameter messages contain an Application Identifier, which is used in the message forwarding process.

The following Application Identifier values are defined:

Diameter Common Messages	0
NASREQ	1 [NASREQ]
Mobile-IP	2 [DIAMMIP]
Diameter Base Accounting	3
Relay	0xffffffff

Relay and redirect agents MUST advertise the Relay Application Identifier, while all other Diameter nodes MUST advertise locally supported applications. The receiver of a Capabilities Exchange message advertising Relay service MUST assume that the sender supports all current and future applications.

Diameter relay and proxy agents are responsible for finding an upstream server that supports the application of a particular message. If none can be found, an error message is returned with the Result-Code AVP set to `DIAMETER_UNABLE_TO_DELIVER`.

## 2.5. Connections vs. Sessions

This section attempts to provide the reader with an understanding of the difference between connection and session, which are terms used extensively throughout this document.

A connection is a transport level connection between two peers, used to send and receive Diameter messages. A session is a logical concept at the application layer, and is shared between an access device and a server, and is identified via the Session-Id AVP

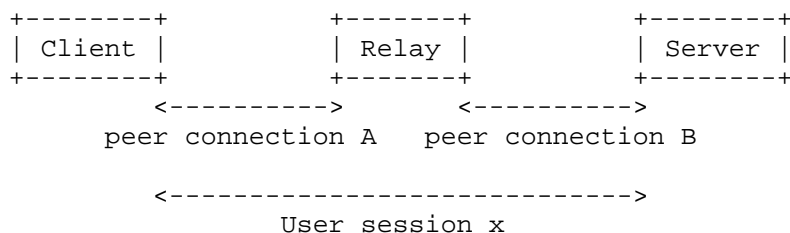


Figure 1: Diameter connections and sessions

In the example provided in Figure 1, peer connection A is established between the Client and its local Relay. Peer connection B is established between the Relay and the Server. User session X spans from the Client via the Relay to the Server. Each "user" of a service causes an auth request to be sent, with a unique session identifier. Once accepted by the server, both the client and the server are aware of the session. It is important to note that there is no relationship between a connection and a session, and that Diameter messages for multiple sessions are all multiplexed through a single connection.

## 2.6. Peer Table

The Diameter Peer Table is used in message forwarding, and referenced by the Realm Routing Table. A Peer Table entry contains the following fields:

### Host identity

Following the conventions described for the DiameterIdentity derived AVP data format in Section 4.4. This field contains the contents of the Origin-Host (Section 6.3) AVP found in the CER or CEA message.

### StatusT

This is the state of the peer entry, and MUST match one of the values listed in Section 5.6.

### Static or Dynamic

Specifies whether a peer entry was statically configured, or dynamically discovered.

### Expiration time

Specifies the time at which dynamically discovered peer table entries are to be either refreshed, or expired.

#### TLS Enabled

Specifies whether TLS is to be used when communicating with the peer.

Additional security information, when needed (e.g., keys, certificates)

### 2.7. Realm-Based Routing Table

All Realm-Based routing lookups are performed against what is commonly known as the Realm Routing Table (see Section 12). A Realm Routing Table Entry contains the following fields:

#### Realm Name

This is the field that is typically used as a primary key in the routing table lookups. Note that some implementations perform their lookups based on longest-match-from-the-right on the realm rather than requiring an exact match.

#### Application Identifier

An application is identified by a vendor id and an application id. For all IETF standards track Diameter applications, the vendor id is zero. A route entry can have a different destination based on the application identification AVP of the message. This field MUST be used as a secondary key field in routing table lookups.

#### Local Action

The Local Action field is used to identify how a message should be treated. The following actions are supported:

1. LOCAL - Diameter messages that resolve to a route entry with the Local Action set to Local can be satisfied locally, and do not need to be routed to another server.
2. RELAY - All Diameter messages that fall within this category MUST be routed to a next hop server, without modifying any non-routing AVPs. See Section 6.1.8 for relaying guidelines
3. PROXY - All Diameter messages that fall within this category MUST be routed to a next hop server. The local server MAY apply its local policies to the message by including new AVPs to the message prior to routing. See Section 6.1.8 for proxying guidelines.
4. REDIRECT - Diameter messages that fall within this category MUST have the identity of the home Diameter server(s) appended, and returned to the sender of the message. See Section 6.1.7 for redirect guidelines.



#### Server Identifier

One or more servers the message is to be routed to. These servers MUST also be present in the Peer table. When the Local Action is set to RELAY or PROXY, this field contains the identity of the server(s) the message must be routed to. When the Local Action field is set to REDIRECT, this field contains the identity of one or more servers the message should be redirected to.

#### Static or Dynamic

Specifies whether a route entry was statically configured, or dynamically discovered.

#### Expiration time

Specifies the time which a dynamically discovered route table entry expires.

It is important to note that Diameter agents MUST support at least one of the LOCAL, RELAY, PROXY or REDIRECT modes of operation. Agents do not need to support all modes of operation in order to conform with the protocol specification, but MUST follow the protocol compliance guidelines in Section 2. Relay agents MUST NOT reorder AVPs, and proxies MUST NOT reorder AVPs.

The routing table MAY include a default entry that MUST be used for any requests not matching any of the other entries. The routing table MAY consist of only such an entry.

When a request is routed, the target server MUST have advertised the Application Identifier (see Section 2.4) for the given message, or have advertised itself as a relay or proxy agent. Otherwise, an error is returned with the Result-Code AVP set to DIAMETER\_UNABLE\_TO\_DELIVER.

## 2.8. Role of Diameter Agents

In addition to client and servers, the Diameter protocol introduces relay, proxy, redirect, and translation agents, each of which is defined in Section 1.3. These Diameter agents are useful for several reasons:

- They can distribute administration of systems to a configurable grouping, including the maintenance of security associations.
- They can be used for concentration of requests from an number of co-located or distributed NAS equipment sets to a set of like user groups.
- They can do value-added processing to the requests or responses.

- They can be used for load balancing.
- A complex network will have multiple authentication sources, they can sort requests and forward towards the correct target.

The Diameter protocol requires that agents maintain transaction state, which is used for failover purposes. Transaction state implies that upon forwarding a request, its Hop-by-Hop identifier is saved; the field is replaced with a locally unique identifier, which is restored to its original value when the corresponding answer is received. The request's state is released upon receipt of the answer. A stateless agent is one that only maintains transaction state.

The Proxy-Info AVP allows stateless agents to add local state to a Diameter request, with the guarantee that the same state will be present in the answer. However, the protocol's failover procedures require that agents maintain a copy of pending requests.

A stateful agent is one that maintains session state information; by keeping track of all authorized active sessions. Each authorized session is bound to a particular service, and its state is considered active either until it is notified otherwise, or by expiration. Each authorized session has an expiration, which is communicated by Diameter servers via the Session-Timeout AVP.

Maintaining session state MAY be useful in certain applications, such as:

- Protocol translation (e.g., RADIUS <-> Diameter)
- Limiting resources authorized to a particular user
- Per user or transaction auditing

A Diameter agent MAY act in a stateful manner for some requests and be stateless for others. A Diameter implementation MAY act as one type of agent for some requests, and as another type of agent for others.

#### 2.8.1. Relay Agents

Relay Agents are Diameter agents that accept requests and route messages to other Diameter nodes based on information found in the messages (e.g., Destination-Realm). This routing decision is performed using a list of supported realms, and known peers. This is known as the Realm Routing Table, as is defined further in Section 2.7.

Relays MAY be used to aggregate requests from multiple Network Access Servers (NASes) within a common geographical area (POP). The use of Relays is advantageous since it eliminates the need for NASes to be configured with the necessary security information they would otherwise require to communicate with Diameter servers in other realms. Likewise, this reduces the configuration load on Diameter servers that would otherwise be necessary when NASes are added, changed or deleted.

Relays modify Diameter messages by inserting and removing routing information, but do not modify any other portion of a message. Relays SHOULD NOT maintain session state but MUST maintain transaction state.

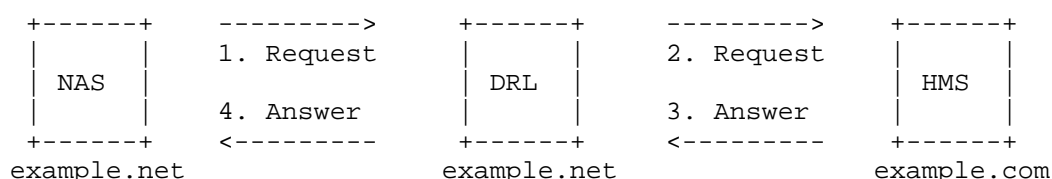


Figure 2: Relaying of Diameter messages

The example provided in Figure 2 depicts a request issued from NAS, which is an access device, for the user bob@example.com. Prior to issuing the request, NAS performs a Diameter route lookup, using "example.com" as the key, and determines that the message is to be relayed to DRL, which is a Diameter Relay. DRL performs the same route lookup as NAS, and relays the message to HMS, which is example.com's Home Diameter Server. HMS identifies that the request can be locally supported (via the realm), processes the authentication and/or authorization request, and replies with an answer, which is routed back to NAS using saved transaction state.

Since Relays do not perform any application level processing, they provide relaying services for all Diameter applications, and therefore MUST advertise the Relay Application Identifier.

### 2.8.2. Proxy Agents

Similarly to relays, proxy agents route Diameter messages using the Diameter Routing Table. However, they differ since they modify messages to implement policy enforcement. This requires that proxies maintain the state of their downstream peers (e.g., access devices) to enforce resource usage, provide admission control, and provisioning.

It is important to note that although proxies MAY provide a value-add function for NASes, they do not allow access devices to use end-to-end security, since modifying messages breaks authentication.

Proxies MAY be used in call control centers or access ISPs that provide outsourced connections, they can monitor the number and types of ports in use, and make allocation and admission decisions according to their configuration.

Proxies that wish to limit resources MUST maintain session state. All proxies MUST maintain transaction state.

Since enforcing policies requires an understanding of the service being provided, Proxies MUST only advertise the Diameter applications they support.

### 2.8.3. Redirect Agents

Redirect agents are useful in scenarios where the Diameter routing configuration needs to be centralized. An example is a redirect agent that provides services to all members of a consortium, but does not wish to be burdened with relaying all messages between realms. This scenario is advantageous since it does not require that the consortium provide routing updates to its members when changes are made to a member's infrastructure.

Since redirect agents do not relay messages, and only return an answer with the information necessary for Diameter agents to communicate directly, they do not modify messages. Since redirect agents do not receive answer messages, they cannot maintain session state. Further, since redirect agents never relay requests, they are not required to maintain transaction state.

The example provided in Figure 3 depicts a request issued from the access device, NAS, for the user bob@example.com. The message is forwarded by the NAS to its relay, DRL, which does not have a routing entry in its Diameter Routing Table for example.com. DRL has a default route configured to DRD, which is a redirect agent that returns a redirect notification to DRL, as well as HMS' contact information. Upon receipt of the redirect notification, DRL establishes a transport connection with HMS, if one doesn't already exist, and forwards the request to it.

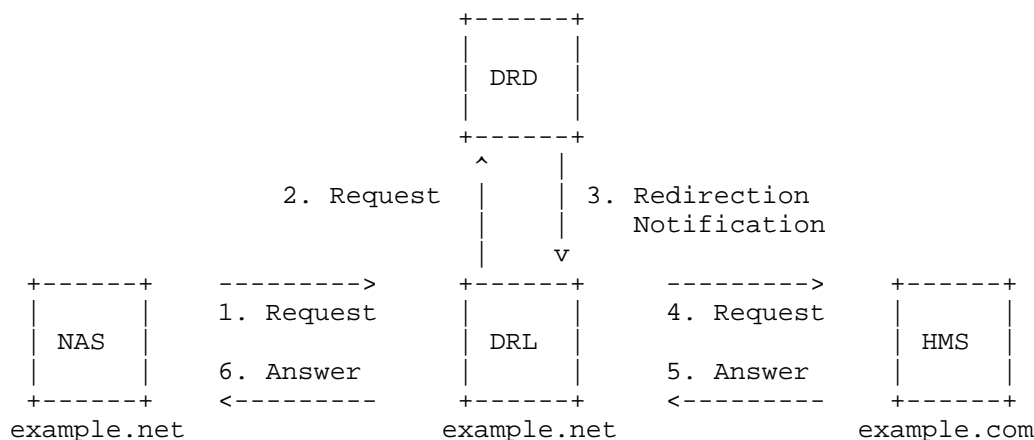


Figure 3: Redirecting a Diameter Message

Since redirect agents do not perform any application level processing, they provide relaying services for all Diameter applications, and therefore MUST advertise the Relay Application Identifier.

#### 2.8.4. Translation Agents

A translation agent is a device that provides translation between two protocols (e.g., RADIUS<->Diameter, TACACS+<->Diameter). Translation agents are likely to be used as aggregation servers to communicate with a Diameter infrastructure, while allowing for the embedded systems to be migrated at a slower pace.

Given that the Diameter protocol introduces the concept of long-lived authorized sessions, translation agents MUST be session stateful and MUST maintain transaction state.

Translation of messages can only occur if the agent recognizes the application of a particular request, and therefore translation agents MUST only advertise their locally supported applications.

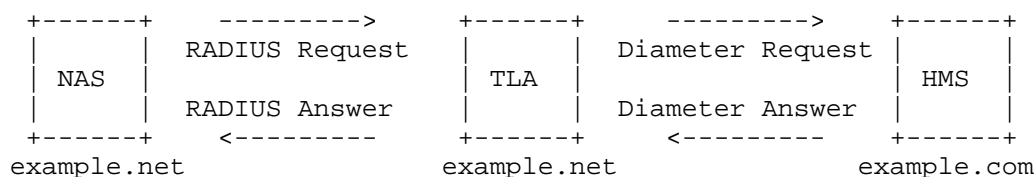


Figure 4: Translation of RADIUS to Diameter

## 2.9. End-to-End Security Framework

End-to-end security services include confidentiality and message origin authentication. These services are provided by supporting AVP integrity and confidentiality between two peers, communicating through agents.

End-to-end security is provided via the End-to-End security extension, described in [AAACMS]. The circumstances requiring the use of end-to-end security are determined by policy on each of the peers. Security policies, which are not the subject of standardization, may be applied by next hop Diameter peer or by destination realm. For example, where TLS or IPsec transmission-level security is sufficient, there may be no need for end-to-end security.

End-to-end security policies include:

- Never use end-to-end security.
- Use end-to-end security on messages containing sensitive AVPs. Which AVPs are sensitive is determined by service provider policy. AVPs containing keys and passwords should be considered sensitive. Accounting AVPs may be considered sensitive. Any AVP for which the P bit may be set or which may be encrypted may be considered sensitive.
- Always use end-to-end security.

It is strongly RECOMMENDED that all Diameter implementations support end-to-end security.

## 2.10. Diameter Path Authorization

As noted in Section 2.2, Diameter requires transmission level security to be used on each connection (TLS or IPsec). Therefore, each connection is authenticated, replay and integrity protected and confidential on a per-packet basis.

In addition to authenticating each connection, each connection as well as the entire session MUST also be authorized. Before initiating a connection, a Diameter Peer MUST check that its peers are authorized to act in their roles. For example, a Diameter peer may be authentic, but that does not mean that it is authorized to act as a Diameter Server advertising a set of Diameter applications.

Prior to bringing up a connection, authorization checks are performed at each connection along the path. Diameter capabilities negotiation (CER/CEA) also MUST be carried out, in order to determine what Diameter applications are supported by each peer. Diameter sessions MUST be routed only through authorized nodes that have advertised support for the Diameter application required by the session.

As noted in Section 6.1.8, a relay or proxy agent MUST append a Route-Record AVP to all requests forwarded. The AVP contains the identity of the peer the request was received from.

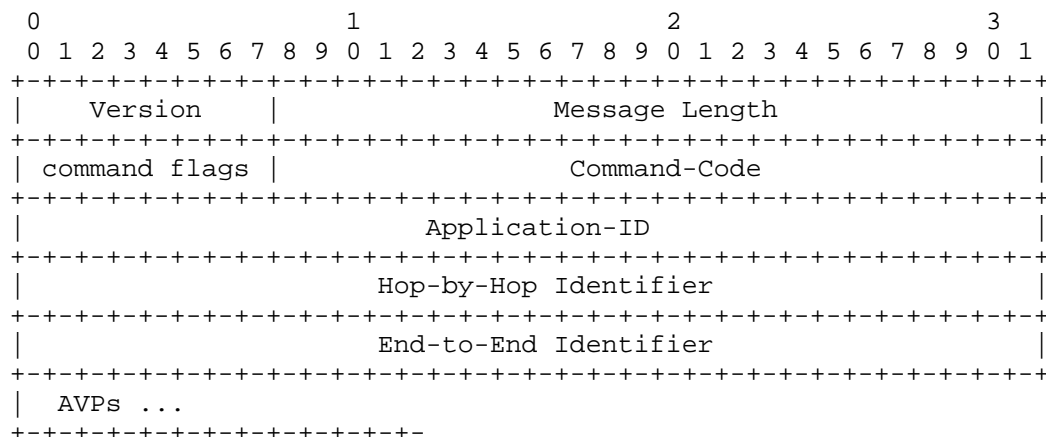
The home Diameter server, prior to authorizing a session, MUST check the Route-Record AVPs to make sure that the route traversed by the request is acceptable. For example, administrators within the home realm may not wish to honor requests that have been routed through an untrusted realm. By authorizing a request, the home Diameter server is implicitly indicating its willingness to engage in the business transaction as specified by the contractual relationship between the server and the previous hop. A `DIAMETER_AUTHORIZATION_REJECTED` error message (see Section 7.1.5) is sent if the route traversed by the request is unacceptable.

A home realm may also wish to check that each accounting request message corresponds to a Diameter response authorizing the session. Accounting requests without corresponding authorization responses SHOULD be subjected to further scrutiny, as should accounting requests indicating a difference between the requested and provided service.

Similarly, the local Diameter agent, on receiving a Diameter response authorizing a session, MUST check the Route-Record AVPs to make sure that the route traversed by the response is acceptable. At each step, forwarding of an authorization response is considered evidence of a willingness to take on financial risk relative to the session. A local realm may wish to limit this exposure, for example, by establishing credit limits for intermediate realms and refusing to accept responses which would violate those limits. By issuing an accounting request corresponding to the authorization response, the local realm implicitly indicates its agreement to provide the service indicated in the authorization response. If the service cannot be provided by the local realm, then a `DIAMETER_UNABLE_TO_COMPLY` error message MUST be sent within the accounting request; a Diameter client receiving an authorization response for a service that it cannot perform MUST NOT substitute an alternate service, and then send accounting requests for the alternate service instead.

### 3. Diameter Header

A summary of the Diameter header format is shown below. The fields are transmitted in network byte order.



#### Version

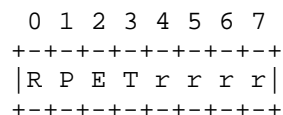
This Version field MUST be set to 1 to indicate Diameter Version 1.

#### Message Length

The Message Length field is three octets and indicates the length of the Diameter message including the header fields.

#### Command Flags

The Command Flags field is eight bits. The following bits are assigned:



R(equest) - If set, the message is a request. If cleared, the message is an answer.

P(roxiable) - If set, the message MAY be proxied, relayed or redirected. If cleared, the message MUST be locally processed.

E(rror) - If set, the message contains a protocol error, and the message will not conform to the ABNF described for this command. Messages with the 'E'



bit set are commonly referred to as error messages. This bit MUST NOT be set in request messages. See Section 7.2.

T(Potentially re-transmitted message)

- This flag is set after a link failover procedure, to aid the removal of duplicate requests. It is set when resending requests not yet acknowledged, as an indication of a possible duplicate due to a link failure. This bit MUST be cleared when sending a request for the first time, otherwise the sender MUST set this flag. Diameter agents only need to be concerned about the number of requests they send based on a single received request; retransmissions by other entities need not be tracked. Diameter agents that receive a request with the T flag set, MUST keep the T flag set in the forwarded request. This flag MUST NOT be set if an error answer message (e.g., a protocol error) has been received for the earlier message. It can be set only in cases where no answer has been received from the server for a request and the request is sent again. This flag MUST NOT be set in answer messages.

r(eserved) - these flag bits are reserved for future use, and MUST be set to zero, and ignored by the receiver.

#### Command-Code

The Command-Code field is three octets, and is used in order to communicate the command associated with the message. The 24-bit address space is managed by IANA (see Section 11.2.1).

Command-Code values 16,777,214 and 16,777,215 (hexadecimal values FFFFFFFE -FFFFFFF) are reserved for experimental use (See Section 11.3).

#### Application-ID

Application-ID is four octets and is used to identify to which application the message is applicable for. The application can be an authentication application, an accounting application or a vendor specific application. See Section 11.3 for the possible values that the application-id may use.

The application-id in the header MUST be the same as what is contained in any relevant AVPs contained in the message.

#### Hop-by-Hop Identifier

The Hop-by-Hop Identifier is an unsigned 32-bit integer field (in network byte order) and aids in matching requests and replies. The sender **MUST** ensure that the Hop-by-Hop identifier in a request is unique on a given connection at any given time, and **MAY** attempt to ensure that the number is unique across reboots. The sender of an Answer message **MUST** ensure that the Hop-by-Hop Identifier field contains the same value that was found in the corresponding request. The Hop-by-Hop identifier is normally a monotonically increasing number, whose start value was randomly generated. An answer message that is received with an unknown Hop-by-Hop Identifier **MUST** be discarded.

#### End-to-End Identifier

The End-to-End Identifier is an unsigned 32-bit integer field (in network byte order) and is used to detect duplicate messages. Upon reboot implementations **MAY** set the high order 12 bits to contain the low order 12 bits of current time, and the low order 20 bits to a random value. Senders of request messages **MUST** insert a unique identifier on each message. The identifier **MUST** remain locally unique for a period of at least 4 minutes, even across reboots. The originator of an Answer message **MUST** ensure that the End-to-End Identifier field contains the same value that was found in the corresponding request. The End-to-End Identifier **MUST NOT** be modified by Diameter agents of any kind. The combination of the Origin-Host (see Section 6.3) and this field is used to detect duplicates. Duplicate requests **SHOULD** cause the same answer to be transmitted (modulo the hop-by-hop Identifier field and any routing AVPs that may be present), and **MUST NOT** affect any state that was set when the original request was processed. Duplicate answer messages that are to be locally consumed (see Section 6.2) **SHOULD** be silently discarded.

#### AVPs

AVPs are a method of encapsulating information relevant to the Diameter message. See Section 4 for more information on AVPs.

### 3.1. Command Codes

Each command Request/Answer pair is assigned a command code, and the sub-type (i.e., request or answer) is identified via the 'R' bit in the Command Flags field of the Diameter header.

Every Diameter message MUST contain a command code in its header's Command-Code field, which is used to determine the action that is to be taken for a particular message. The following Command Codes are defined in the Diameter base protocol:

Command-Name	Abbrev.	Code	Reference
-----	-----	-----	-----
Abort-Session-Request	ASR	274	8.5.1
Abort-Session-Answer	ASA	274	8.5.2
Accounting-Request	ACR	271	9.7.1
Accounting-Answer	ACA	271	9.7.2
Capabilities-Exchange-Request	CER	257	5.3.1
Capabilities-Exchange-Answer	CEA	257	5.3.2
Device-Watchdog-Request	DWR	280	5.5.1
Device-Watchdog-Answer	DWA	280	5.5.2
Disconnect-Peer-Request	DPR	282	5.4.1
Disconnect-Peer-Answer	DPA	282	5.4.2
Re-Auth-Request	RAR	258	8.3.1
Re-Auth-Answer	RAA	258	8.3.2
Session-Termination-Request	STR	275	8.4.1
Session-Termination-Answer	STA	275	8.4.2

### 3.2. Command Code ABNF specification

Every Command Code defined MUST include a corresponding ABNF specification, which is used to define the AVPs that MUST or MAY be present. The following format is used in the definition:

```
command-def      = command-name "::~=" diameter-message
command-name     = diameter-name
diameter-name    = ALPHA *(ALPHA / DIGIT / "-")
diameter-message = header [ *fixed] [ *required] [ *optional]
                  [ *fixed]
header           = "<" Diameter-Header:" command-id
                  [r-bit] [p-bit] [e-bit] [application-id]>"
application-id   = 1*DIGIT
command-id       = 1*DIGIT
                  ; The Command Code assigned to the command
r-bit           = ", REQ"
                  ; If present, the 'R' bit in the Command
                  ; Flags is set, indicating that the message
                  ; is a request, as opposed to an answer.
p-bit           = ", PXY"
                  ; If present, the 'P' bit in the Command
                  ; Flags is set, indicating that the message
                  ; is proxiabile.
e-bit           = ", ERR"
                  ; If present, the 'E' bit in the Command
                  ; Flags is set, indicating that the answer
                  ; message contains a Result-Code AVP in
                  ; the "protocol error" class.
fixed           = [qual] "<" avp-spec ">"
                  ; Defines the fixed position of an AVP
required        = [qual] "{" avp-spec "}"
                  ; The AVP MUST be present and can appear
                  ; anywhere in the message.
```

optional = [qual] "[" avp-name "]"  
; The avp-name in the 'optional' rule cannot  
; evaluate to any AVP Name which is included  
; in a fixed or required rule. The AVP can  
; appear anywhere in the message.

qual = [min] "\*" [max]  
; See ABNF conventions, RFC 2234 Section 6.6.  
; The absence of any qualifiers depends on whether  
; it precedes a fixed, required, or optional  
; rule. If a fixed or required rule has no  
; qualifier, then exactly one such AVP MUST  
; be present. If an optional rule has no  
; qualifier, then 0 or 1 such AVP may be  
; present.  
;  
; NOTE: "[" and "]" have a different meaning  
; than in ABNF (see the optional rule, above).  
; These braces cannot be used to express  
; optional fixed rules (such as an optional  
; ICV at the end). To do this, the convention  
; is '0\*1fixed'.

min = 1\*DIGIT  
; The minimum number of times the element may  
; be present. The default value is zero.

max = 1\*DIGIT  
; The maximum number of times the element may  
; be present. The default value is infinity. A  
; value of zero implies the AVP MUST NOT be  
; present.

avp-spec = diameter-name  
; The avp-spec has to be an AVP Name, defined  
; in the base or extended Diameter  
; specifications.

avp-name = avp-spec / "AVP"  
; The string "AVP" stands for \*any\* arbitrary  
; AVP Name, which does not conflict with the  
; required or fixed position AVPs defined in  
; the command code definition.

The following is a definition of a fictitious command code:

```
Example-Request ::= < "Diameter-Header: 9999999, REQ, PXY" >
    { User-Name }
    * { Origin-Host }
    * [ AVP
```

### 3.3. Diameter Command Naming Conventions

Diameter command names typically includes one or more English words followed by the verb Request or Answer. Each English word is delimited by a hyphen. A three-letter acronym for both the request and answer is also normally provided.

An example is a message set used to terminate a session. The command name is Session-Terminate-Request and Session-Terminate-Answer, while the acronyms are STR and STA, respectively.

Both the request and the answer for a given command share the same command code. The request is identified by the R(equest) bit in the Diameter header set to one (1), to ask that a particular action be performed, such as authorizing a user or terminating a session. Once the receiver has completed the request it issues the corresponding answer, which includes a result code that communicates one of the following:

- The request was successful
- The request failed
- An additional request must be sent to provide information the peer requires prior to returning a successful or failed answer.
- The receiver could not process the request, but provides information about a Diameter peer that is able to satisfy the request, known as redirect.

Additional information, encoded within AVPs, MAY also be included in answer messages.

### 4. Diameter AVPs

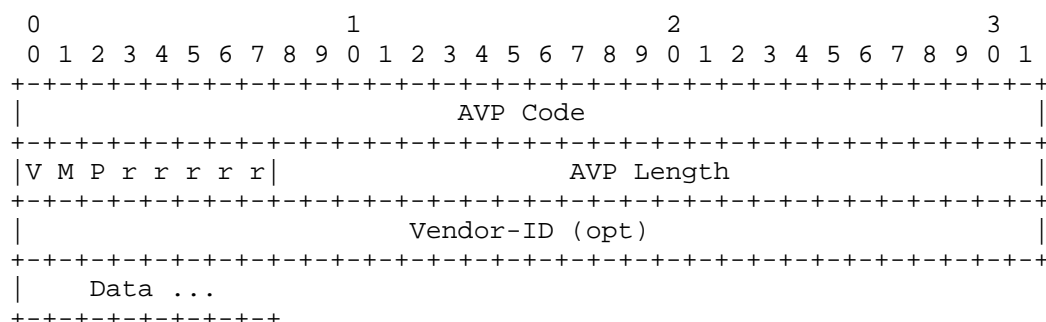
Diameter AVPs carry specific authentication, accounting, authorization, routing and security information as well as configuration details for the request and reply.

Some AVPs MAY be listed more than once. The effect of such an AVP is specific, and is specified in each case by the AVP description.

Each AVP of type OctetString MUST be padded to align on a 32-bit boundary, while other AVP types align naturally. A number of zero-valued bytes are added to the end of the AVP Data field till a word boundary is reached. The length of the padding is not reflected in the AVP Length field.

#### 4.1. AVP Header

The fields in the AVP header MUST be sent in network byte order. The format of the header is:



##### AVP Code

The AVP Code, combined with the Vendor-Id field, identifies the attribute uniquely. AVP numbers 1 through 255 are reserved for backward compatibility with RADIUS, without setting the Vendor-Id field. AVP numbers 256 and above are used for Diameter, which are allocated by IANA (see Section 11.1).

##### AVP Flags

The AVP Flags field informs the receiver how each attribute must be handled. The 'r' (reserved) bits are unused and SHOULD be set to 0. Note that subsequent Diameter applications MAY define additional bits within the AVP Header, and an unrecognized bit SHOULD be considered an error. The 'P' bit indicates the need for encryption for end-to-end security.

The 'M' Bit, known as the Mandatory bit, indicates whether support of the AVP is required. If an AVP with the 'M' bit set is received by a Diameter client, server, proxy, or translation agent and either the AVP or its value is unrecognized, the message MUST be rejected. Diameter Relay and redirect agents MUST NOT reject messages with unrecognized AVPs.

The 'M' bit MUST be set according to the rules defined for the AVP containing it. In order to preserve interoperability, a Diameter implementation MUST be able to exclude from a Diameter message any Mandatory AVP which is neither defined in the base Diameter protocol nor in any of the Diameter Application specifications governing the message in which it appears. It MAY do this in one of the following ways:

- 1) If a message is rejected because it contains a Mandatory AVP which is neither defined in the base Diameter standard nor in any of the Diameter Application specifications governing the message in which it appears, the implementation may resend the message without the AVP, possibly inserting additional standard AVPs instead.
- 2) A configuration option may be provided on a system wide, per peer, or per realm basis that would allow/prevent particular Mandatory AVPs to be sent. Thus an administrator could change the configuration to avoid interoperability problems.

Diameter implementations are required to support all Mandatory AVPs which are allowed by the message's formal syntax and defined either in the base Diameter standard or in one of the Diameter Application specifications governing the message.

AVPs with the 'M' bit cleared are informational only and a receiver that receives a message with such an AVP that is not supported, or whose value is not supported, MAY simply ignore the AVP.

The 'V' bit, known as the Vendor-Specific bit, indicates whether the optional Vendor-ID field is present in the AVP header. When set the AVP Code belongs to the specific vendor code address space.

Unless otherwise noted, AVPs will have the following default AVP Flags field settings:

The 'M' bit MUST be set. The 'V' bit MUST NOT be set.

#### AVP Length

The AVP Length field is three octets, and indicates the number of octets in this AVP including the AVP Code, AVP Length, AVP Flags, Vendor-ID field (if present) and the AVP data. If a message is received with an invalid attribute length, the message SHOULD be rejected.



#### 4.1.1. Optional Header Elements

The AVP Header contains one optional field. This field is only present if the respective bit-flag is enabled.

##### Vendor-ID

The Vendor-ID field is present if the 'V' bit is set in the AVP Flags field. The optional four-octet Vendor-ID field contains the IANA assigned "SMI Network Management Private Enterprise Codes" [ASSIGNNO] value, encoded in network byte order. Any vendor wishing to implement a vendor-specific Diameter AVP MUST use their own Vendor-ID along with their privately managed AVP address space, guaranteeing that they will not collide with any other vendor's vendor-specific AVP(s), nor with future IETF applications.

A vendor ID value of zero (0) corresponds to the IETF adopted AVP values, as managed by the IANA. Since the absence of the vendor ID field implies that the AVP in question is not vendor specific, implementations MUST NOT use the zero (0) vendor ID.

#### 4.2. Basic AVP Data Formats

The Data field is zero or more octets and contains information specific to the Attribute. The format and length of the Data field is determined by the AVP Code and AVP Length fields. The format of the Data field MUST be one of the following base data types or a data type derived from the base data types. In the event that a new Basic AVP Data Format is needed, a new version of this RFC must be created.

##### OctetString

The data contains arbitrary data of variable length. Unless otherwise noted, the AVP Length field MUST be set to at least 8 (12 if the 'V' bit is enabled). AVP Values of this type that are not a multiple of four-octets in length is followed by the necessary padding so that the next AVP (if any) will start on a 32-bit boundary.

##### Integer32

32 bit signed value, in network byte order. The AVP Length field MUST be set to 12 (16 if the 'V' bit is enabled).

##### Integer64

64 bit signed value, in network byte order. The AVP Length field MUST be set to 16 (20 if the 'V' bit is enabled).

**Unsigned32**

32 bit unsigned value, in network byte order. The AVP Length field MUST be set to 12 (16 if the 'V' bit is enabled).

**Unsigned64**

64 bit unsigned value, in network byte order. The AVP Length field MUST be set to 16 (20 if the 'V' bit is enabled).

**Float32**

This represents floating point values of single precision as described by [FLOATPOINT]. The 32-bit value is transmitted in network byte order. The AVP Length field MUST be set to 12 (16 if the 'V' bit is enabled).

**Float64**

This represents floating point values of double precision as described by [FLOATPOINT]. The 64-bit value is transmitted in network byte order. The AVP Length field MUST be set to 16 (20 if the 'V' bit is enabled).

**Grouped**

The Data field is specified as a sequence of AVPs. Each of these AVPs follows - in the order in which they are specified - including their headers and padding. The AVP Length field is set to 8 (12 if the 'V' bit is enabled) plus the total length of all included AVPs, including their headers and padding. Thus the AVP length field of an AVP of type Grouped is always a multiple of 4.

#### 4.3. Derived AVP Data Formats

In addition to using the Basic AVP Data Formats, applications may define data formats derived from the Basic AVP Data Formats. An application that defines new AVP Derived Data Formats MUST include them in a section entitled "AVP Derived Data Formats", using the same format as the definitions below. Each new definition must be either defined or listed with a reference to the RFC that defines the format.

The below AVP Derived Data Formats are commonly used by applications.

**Address**

The Address format is derived from the OctetString AVP Base Format. It is a discriminated union, representing, for example a 32-bit (IPv4) [IPV4] or 128-bit (IPv6) [IPV6] address, most significant octet first. The first two octets of the Address

AVP represents the AddressType, which contains an Address Family defined in [IANAADFAM]. The AddressType is used to discriminate the content and format of the remaining octets.

#### Time

The Time format is derived from the OctetString AVP Base Format. The string MUST contain four octets, in the same format as the first four bytes are in the NTP timestamp format. The NTP Timestamp format is defined in chapter 3 of [SNTP].

This represents the number of seconds since 0h on 1 January 1900 with respect to the Coordinated Universal Time (UTC).

On 6h 28m 16s UTC, 7 February 2036 the time value will overflow. SNTP [SNTP] describes a procedure to extend the time to 2104. This procedure MUST be supported by all DIAMETER nodes.

#### UTF8String

The UTF8String format is derived from the OctetString AVP Base Format. This is a human readable string represented using the ISO/IEC IS 10646-1 character set, encoded as an OctetString using the UTF-8 [UFT8] transformation format described in RFC 2279.

Since additional code points are added by amendments to the 10646 standard from time to time, implementations MUST be prepared to encounter any code point from 0x00000001 to 0x7fffffff. Byte sequences that do not correspond to the valid encoding of a code point into UTF-8 charset or are outside this range are prohibited.

The use of control codes SHOULD be avoided. When it is necessary to represent a new line, the control code sequence CR LF SHOULD be used.

The use of leading or trailing white space SHOULD be avoided.

For code points not directly supported by user interface hardware or software, an alternative means of entry and display, such as hexadecimal, MAY be provided.

For information encoded in 7-bit US-ASCII, the UTF-8 charset is identical to the US-ASCII charset.

UTF-8 may require multiple bytes to represent a single character / code point; thus the length of an UTF8String in octets may be different from the number of characters encoded.

Note that the AVP Length field of an UTF8String is measured in octets, not characters.

### DiameterIdentity

The DiameterIdentity format is derived from the OctetString AVP Base Format.

DiameterIdentity = FQDN

DiameterIdentity value is used to uniquely identify a Diameter node for purposes of duplicate connection and routing loop detection.

The contents of the string MUST be the FQDN of the Diameter node. If multiple Diameter nodes run on the same host, each Diameter node MUST be assigned a unique DiameterIdentity. If a Diameter node can be identified by several FQDNs, a single FQDN should be picked at startup, and used as the only DiameterIdentity for that node, whatever the connection it is sent on.

### DiameterURI

The DiameterURI MUST follow the Uniform Resource Identifiers (URI) syntax [URI] rules specified below:

"aaa://" FQDN [ port ] [ transport ] [ protocol ]

; No transport security

"aaas://" FQDN [ port ] [ transport ] [ protocol ]

; Transport security used

FQDN = Fully Qualified Host Name

port = ":" 1\*DIGIT

; One of the ports used to listen for  
; incoming connections.  
; If absent,  
; the default Diameter port (3868) is  
; assumed.

transport = ";transport=" transport-protocol

; One of the transports used to listen  
; for incoming connections. If absent,  
; the default SCTP [SCTP] protocol is  
; assumed. UDP MUST NOT be used when  
; the aaa-protocol field is set to  
; diameter.

```
transport-protocol = ( "tcp" / "sctp" / "udp" )

protocol           = ";protocol=" aaa-protocol

                   ; If absent, the default AAA protocol
                   ; is diameter.

aaa-protocol       = ( "diameter" / "radius" / "tacacs+" )
```

The following are examples of valid Diameter host identities:

```
aaa://host.example.com;transport=tcp
aaa://host.example.com:6666;transport=tcp
aaa://host.example.com;protocol=diameter
aaa://host.example.com:6666;protocol=diameter
aaa://host.example.com:6666;transport=tcp;protocol=diameter
aaa://host.example.com:1813;transport=udp;protocol=radius
```

#### Enumerated

Enumerated is derived from the Integer32 AVP Base Format. The definition contains a list of valid values and their interpretation and is described in the Diameter application introducing the AVP.

#### IPFilterRule

The IPFilterRule format is derived from the OctetString AVP Base Format. It uses the ASCII charset. Packets may be filtered based on the following information that is associated with it:

Direction	(in or out)
Source and destination IP address	(possibly masked)
Protocol	
Source and destination port	(lists or ranges)
TCP flags	
IP fragment flag	
IP options	
ICMP types	

Rules for the appropriate direction are evaluated in order, with the first matched rule terminating the evaluation. Each packet is evaluated once. If no rule matches, the packet is dropped if the last rule evaluated was a permit, and passed if the last rule was a deny.

IPFilterRule filters MUST follow the format:

action dir proto from src to dst [options]

action        permit - Allow packets that match the rule.  
              deny    - Drop packets that match the rule.

dir            "in" is from the terminal, "out" is to the  
              terminal.

proto         An IP protocol specified by number. The "ip"  
              keyword means any protocol will match.

src and dst   <address/mask> [ports]

The <address/mask> may be specified as:

ipno           An IPv4 or IPv6 number in dotted-  
              quad or canonical IPv6 form. Only  
              this exact IP number will match the  
              rule.

ipno/bits     An IP number as above with a mask  
              width of the form 1.2.3.4/24. In  
              this case, all IP numbers from  
              1.2.3.0 to 1.2.3.255 will match.  
              The bit width MUST be valid for the  
              IP version and the IP number MUST  
              NOT have bits set beyond the mask.  
              For a match to occur, the same IP  
              version must be present in the  
              packet that was used in describing  
              the IP address. To test for a  
              particular IP version, the bits part  
              can be set to zero. The keyword  
              "any" is 0.0.0.0/0 or the IPv6  
              equivalent. The keyword "assigned"  
              is the address or set of addresses  
              assigned to the terminal. For IPv4,  
              a typical first rule is often "deny  
              in ip! assigned"

The sense of the match can be inverted by  
preceding an address with the not modifier (!),  
causing all other addresses to be matched  
instead. This does not affect the selection of  
port numbers.

With the TCP, UDP and SCTP protocols, optional ports may be specified as:

`{port/port-port}[ ,ports[,...]]`

The '-' notation specifies a range of ports (including boundaries).

Fragmented packets that have a non-zero offset (i.e., not the first fragment) will never match a rule that has one or more port specifications. See the frag option for details on matching fragmented packets.

options:

frag Match if the packet is a fragment and this is not the first fragment of the datagram. frag may not be used in conjunction with either tcpflags or TCP/UDP port specifications.

ipoptions spec

Match if the IP header contains the comma separated list of options specified in spec. The supported IP options are:

ssrr (strict source route), lsrr (loose source route), rr (record packet route) and ts (timestamp). The absence of a particular option may be denoted with a '!'.

tcptoptions spec

Match if the TCP header contains the comma separated list of options specified in spec. The supported TCP options are:

mss (maximum segment size), window (tcp window advertisement), sack (selective ack), ts (rfc1323 timestamp) and cc (rfc1644 t/tcp connection count). The absence of a particular option may be denoted with a '!'.

established

TCP packets only. Match packets that have the RST or ACK bits set.

setup TCP packets only. Match packets that have the SYN bit set but no ACK bit.

`tcpflags spec`

TCP packets only. Match if the TCP header contains the comma separated list of flags specified in spec. The supported TCP flags are:

`fin`, `syn`, `rst`, `psh`, `ack` and `urg`. The absence of a particular flag may be denoted with a '!'. A rule that contains a `tcpflags` specification can never match a fragmented packet that has a non-zero offset. See the `frag` option for details on matching fragmented packets.

`icmptypes types`

ICMP packets only. Match if the ICMP type is in the list types. The list may be specified as any combination of ranges or individual types separated by commas. Both the numeric values and the symbolic values listed below can be used. The supported ICMP types are:

`echo reply (0)`, `destination unreachable (3)`, `source quench (4)`, `redirect (5)`, `echo request (8)`, `router advertisement (9)`, `router solicitation (10)`, `time-to-live exceeded (11)`, `IP header bad (12)`, `timestamp request (13)`, `timestamp reply (14)`, `information request (15)`, `information reply (16)`, `address mask request (17)` and `address mask reply (18)`.

There is one kind of packet that the access device MUST always discard, that is an IP fragment with a fragment offset of one. This is a valid packet, but it only has one use, to try to circumvent firewalls.

An access device that is unable to interpret or apply a deny rule MUST terminate the session. An access device that is unable to interpret or apply a permit rule MAY apply a more restrictive rule. An access device MAY apply deny rules of its own before the supplied rules, for example to protect the access device owner's infrastructure.

The rule syntax is a modified subset of `ipfw(8)` from FreeBSD, and the `ipfw.c` code may provide a useful base for implementations.



#### QoSFilterRule

The QoSFilterRule format is derived from the OctetString AVP Base Format. It uses the ASCII charset. Packets may be marked or metered based on the following information that is associated with it:

Direction	(in or out)
Source and destination IP address	(possibly masked)
Protocol	
Source and destination port	(lists or ranges)
DSCP values	(no mask or range)

Rules for the appropriate direction are evaluated in order, with the first matched rule terminating the evaluation. Each packet is evaluated once. If no rule matches, the packet is treated as best effort. An access device that is unable to interpret or apply a QoS rule SHOULD NOT terminate the session.

QoSFilterRule filters MUST follow the format:

action dir proto from src to dst [options]

tag	- Mark packet with a specific DSCP [DIFFSERV]. The DSCP option MUST be included.
meter	- Meter traffic. The metering options MUST be included.

dir           The format is as described under IPFilterRule.

proto           The format is as described under IPFilterRule.

src and dst    The format is as described under IPFilterRule.

#### 4.4. Grouped AVP Values

The Diameter protocol allows AVP values of type 'Grouped.' This implies that the Data field is actually a sequence of AVPs. It is possible to include an AVP with a Grouped type within a Grouped type, that is, to nest them. AVPs within an AVP of type Grouped have the same padding requirements as non-Grouped AVPs, as defined in Section 4.

The AVP Code numbering space of all AVPs included in a Grouped AVP is the same as for non-grouped AVPs. Further, if any of the AVPs encapsulated within a Grouped AVP has the 'M' (mandatory) bit set, the Grouped AVP itself MUST also include the 'M' bit set.

Every Grouped AVP defined MUST include a corresponding grammar, using ABNF [ABNF] (with modifications), as defined below.

```
grouped-avp-def  = name "::=" avp

name-fmt         = ALPHA *(ALPHA / DIGIT / "-")

name             = name-fmt
                  ; The name has to be the name of an AVP,
                  ; defined in the base or extended Diameter
                  ; specifications.

avp              = header [ *fixed] [ *required] [ *optional]
                  [ *fixed]

header           = "<" "AVP-Header:" avpcode [vendor] ">"

avpcode          = 1*DIGIT
                  ; The AVP Code assigned to the Grouped AVP

vendor           = 1*DIGIT
                  ; The Vendor-ID assigned to the Grouped AVP.
                  ; If absent, the default value of zero is
                  ; used.
```

#### 4.4.1. Example AVP with a Grouped Data type

The Example-AVP (AVP Code 999999) is of type Grouped and is used to clarify how Grouped AVP values work. The Grouped Data field has the following ABNF grammar:

```
Example-AVP ::= < AVP Header: 999999 >
              { Origin-Host }
              1*{ Session-Id }
              *[ AVP ]
```

An Example-AVP with Grouped Data follows.

The Origin-Host AVP is required (Section 6.3). In this case:

```
Origin-Host = "example.com".
```

One or more Session-Ids must follow. Here there are two:

```
Session-Id =  
    "grump.example.com:33041;23432;893;0AF3B81"
```

```
Session-Id =  
    "grump.example.com:33054;23561;2358;0AF3B82"
```

optional AVPs included are

```
Recovery-Policy = <binary>  
    2163bcd0ad82371f6bc09484133c3f09ad74a0dd5346d54195a7cf0b35  
    2cabcb881839a4fdcfbc1769e2677a4c1fb499284c5f70b48f58503a45c5  
    c2d6943f82d5930f2b7c1da640f476f0e9c9572a50db8ea6e51e1c2c7bd  
    f8bb43dc995144b8dbe297ac739493946803e1cee3e15d9b765008alb2a  
    cf4ac777c80041d72c01e691cf751dbf86e85f509f3988e5875dc905119  
    26841f00f0e29a6d1ddc1a842289d440268681e052b30fb638045f7779c  
    1d873c784f054f688f5001559ecff64865ef975f3e60d2fd7966b8c7f92
```

```
Futuristic-Acct-Record = <binary>  
    fe19da5802acd98b07a5b86cb4d5d03f0314ab9ef1ad0b67111ff3b90a0  
    57fe29620bf3585fd2dd9fcc38ce62f6cc208c6163c008f4258d1bc88b8  
    17694a74ccad3ec69269461b14b2e7a4c111fb239e33714da207983f58c  
    41d018d56fe938f3cbf089aac12a912a2f0d1923a9390e5f789cb2e5067  
    d3427475e49968f841
```

The data for the optional AVPs is represented in hex since the format of these AVPs is neither known at the time of definition of the Example-AVP group, nor (likely) at the time when the example instance of this AVP is interpreted - except by Diameter implementations which support the same set of AVPs. The encoding example illustrates how padding is used and how length fields are calculated. Also note that AVPs may be present in the Grouped AVP value which the receiver cannot interpret (here, the Recover-Policy and Futuristic-Acct-Record AVPs).

This AVP would be encoded as follows:

	0	1	2	3	4	5	6	7
0	Example AVP Header (AVP Code = 999999), Length = 468							
8	Origin-Host AVP Header (AVP Code = 264), Length = 19							
16	'e'	'x'	'a'	'm'	'p'	'l'	'e'	'.'
24	'c'	'o'	'm'	Padding	Session-Id AVP Header			
32	(AVP Code = 263), Length = 50				'g'	'r'	'u'	'm'
. . .								
64	'A'	'F'	'3'	'B'	'8'	'1'	Padding	Padding
72	Session-Id AVP Header (AVP Code = 263), Length = 51							
80	'g'	'r'	'u'	'm'	'p'	'.'	'e'	'x'
. . .								
104	'0'	'A'	'F'	'3'	'B'	'8'	'2'	Padding
112	Recovery-Policy Header (AVP Code = 8341), Length = 223							
120	0x21	0x63	0xbc	0x1d	0x0a	0xd8	0x23	0x71
. . .								
320	0x2f	0xd7	0x96	0x6b	0x8c	0x7f	0x92	Padding
328	Futuristic-Acct-Record Header (AVP Code = 15930), Length = 137							
336	0xfe	0x19	0xda	0x58	0x02	0xac	0xd9	0x8b
. . .								
464	0x41	Padding	Padding	Padding				

#### 4.5. Diameter Base Protocol AVPs

The following table describes the Diameter AVPs defined in the base protocol, their AVP Code values, types, possible flag values and whether the AVP MAY be encrypted. For the originator of a Diameter message, "Encr" (Encryption) means that if a message containing that AVP is to be sent via a Diameter agent (proxy, redirect or relay) then the message MUST NOT be sent unless there is end-to-end security between the originator and the recipient and integrity / confidentiality protection is offered for this AVP OR the originator has locally trusted configuration that indicates that end-to-end security is not needed. Similarly, for the originator of a Diameter message, a "P" in the "MAY" column means that if a message containing that AVP is to be sent via a Diameter agent (proxy, redirect or relay) then the message MUST NOT be sent unless there is end-to-end security between the originator and the recipient or the originator has locally trusted configuration that indicates that end-to-end security is not needed.

Due to space constraints, the short form DiamIdent is used to represent DiameterIdentity.

Attribute Name	AVP Code	Section Defined	Data Type	AVP Flag rules				Encr
				MUST	MAY	SHLD NOT	MUST NOT	
Acct-Interim-Interval	85	9.8.2	Unsigned32	M	P		V	Y
Accounting-Realtime-Required	483	9.8.7	Enumerated	M	P		V	Y
Acct-Multi-Session-Id	50	9.8.5	UTF8String	M	P		V	Y
Accounting-Record-Number	485	9.8.3	Unsigned32	M	P		V	Y
Accounting-Record-Type	480	9.8.1	Enumerated	M	P		V	Y
Accounting-Session-Id	44	9.8.4	OctetString	M	P		V	Y
Accounting-Sub-Session-Id	287	9.8.6	Unsigned64	M	P		V	Y
Acct-Application-Id	259	6.9	Unsigned32	M	P		V	N
Auth-Application-Id	258	6.8	Unsigned32	M	P		V	N
Auth-Request-Type	274	8.7	Enumerated	M	P		V	N
Authorization-Lifetime	291	8.9	Unsigned32	M	P		V	N
Auth-Grace-Period	276	8.10	Unsigned32	M	P		V	N
Auth-Session-State	277	8.11	Enumerated	M	P		V	N
Re-Auth-Request-Type	285	8.12	Enumerated	M	P		V	N
Class	25	8.20	OctetString	M	P		V	Y
Destination-Host	293	6.5	DiamIdent	M	P		V	N
Destination-Realm	283	6.6	DiamIdent	M	P		V	N
Disconnect-Cause	273	5.4.3	Enumerated	M	P		V	N
E2E-Sequence AVP	300	6.15	Grouped	M	P		V	Y
Error-Message	281	7.3	UTF8String		P		V,M	N
Error-Reporting-Host	294	7.4	DiamIdent		P		V,M	N
Event-Timestamp	55	8.21	Time	M	P		V	N
Experimental-Result	297	7.6	Grouped	M	P		V	N

				AVP Flag rules					
Attribute Name	AVP Code	Section Defined	Data Type	MUST	MAY	SHLD NOT	MUST NOT	MAY Encr	
Experimental-Result-Code	298	7.7	Unsigned32	M	P		V	N	
Failed-AVP	279	7.5	Grouped	M	P		V	N	
Firmware-Revision	267	5.3.4	Unsigned32				P,V,M	N	
Host-IP-Address	257	5.3.5	Address	M	P		V	N	
Inband-Security-Id	299	6.10	Unsigned32	M	P		V	N	
Multi-Round-Time-Out	272	8.19	Unsigned32	M	P		V	Y	
Origin-Host	264	6.3	DiamIdent	M	P		V	N	
Origin-Realm	296	6.4	DiamIdent	M	P		V	N	
Origin-State-Id	278	8.16	Unsigned32	M	P		V	N	
Product-Name	269	5.3.7	UTF8String				P,V,M	N	
Proxy-Host	280	6.7.3	DiamIdent	M			P,V	N	
Proxy-Info	284	6.7.2	Grouped	M			P,V	N	
Proxy-State	33	6.7.4	OctetString	M			P,V	N	
Redirect-Host	292	6.12	DiamURI	M	P		V	N	
Redirect-Host-Usage	261	6.13	Enumerated	M	P		V	N	
Redirect-Max-Cache-Time	262	6.14	Unsigned32	M	P		V	N	
Result-Code	268	7.1	Unsigned32	M	P		V	N	
Route-Record	282	6.7.1	DiamIdent	M			P,V	N	
Session-Id	263	8.8	UTF8String	M	P		V	Y	
Session-Timeout	27	8.13	Unsigned32	M	P		V	N	
Session-Binding	270	8.17	Unsigned32	M	P		V	Y	
Session-Server-Failover	271	8.18	Enumerated	M	P		V	Y	
Supported-Vendor-Id	265	5.3.6	Unsigned32	M	P		V	N	
Termination-Cause	295	8.15	Enumerated	M	P		V	N	
User-Name	1	8.14	UTF8String	M	P		V	Y	
Vendor-Id	266	5.3.3	Unsigned32	M	P		V	N	
Vendor-Specific-Application-Id	260	6.11	Grouped	M	P		V	N	

## 5. Diameter Peers

This section describes how Diameter nodes establish connections and communicate with peers.

### 5.1. Peer Connections

Although a Diameter node may have many possible peers that it is able to communicate with, it may not be economical to have an established connection to all of them. At a minimum, a Diameter node SHOULD have an established connection with two peers per realm, known as the primary and secondary peers. Of course, a node MAY have additional connections, if it is deemed necessary. Typically, all messages for a realm are sent to the primary peer, but in the event that failover procedures are invoked, any pending requests are sent to the secondary peer. However, implementations are free to load balance requests between a set of peers.

Note that a given peer MAY act as a primary for a given realm, while acting as a secondary for another realm.

When a peer is deemed suspect, which could occur for various reasons, including not receiving a DWA within an allotted timeframe, no new requests should be forwarded to the peer, but failover procedures are invoked. When an active peer is moved to this mode, additional connections SHOULD be established to ensure that the necessary number of active connections exists.

There are two ways that a peer is removed from the suspect peer list:

1. The peer is no longer reachable, causing the transport connection to be shutdown. The peer is moved to the closed state.
2. Three watchdog messages are exchanged with accepted round trip times, and the connection to the peer is considered stabilized.

In the event the peer being removed is either the primary or secondary, an alternate peer SHOULD replace the deleted peer, and assume the role of either primary or secondary.

### 5.2. Diameter Peer Discovery

Allowing for dynamic Diameter agent discovery will make it possible for simpler and more robust deployment of Diameter services. In order to promote interoperable implementations of Diameter peer discovery, the following mechanisms are described. These are based



on existing IETF standards. The first option (manual configuration) MUST be supported by all DIAMETER nodes, while the latter two options (SRVLOC and DNS) MAY be supported.

There are two cases where Diameter peer discovery may be performed. The first is when a Diameter client needs to discover a first-hop Diameter agent. The second case is when a Diameter agent needs to discover another agent - for further handling of a Diameter operation. In both cases, the following 'search order' is recommended:

1. The Diameter implementation consults its list of static (manually) configured Diameter agent locations. These will be used if they exist and respond.
2. The Diameter implementation uses SLPv2 [SLP] to discover Diameter services. The Diameter service template [TEMPLATE] is included in Appendix A.

It is recommended that SLPv2 security be deployed (this requires distributing keys to SLPv2 agents). This is discussed further in Appendix A. SLPv2 security SHOULD be used (requiring distribution of keys to SLPv2 agents) in order to ensure that discovered peers are authorized for their roles. SLPv2 is discussed further in Appendix A.

3. The Diameter implementation performs a NAPTR query for a server in a particular realm. The Diameter implementation has to know in advance which realm to look for a Diameter agent in. This could be deduced, for example, from the 'realm' in a NAI that a Diameter implementation needed to perform a Diameter operation on.

- 3.1 The services relevant for the task of transport protocol selection are those with NAPTR service fields with values "AAA+D2x", where x is a letter that corresponds to a transport protocol supported by the domain. This specification defines D2T for TCP and D2S for SCTP. We also establish an IANA registry for NAPTR service name to transport protocol mappings.

These NAPTR records provide a mapping from a domain, to the SRV record for contacting a server with the specific transport protocol in the NAPTR services field. The resource record will contain an empty regular expression and a replacement value, which is the SRV record for that particular transport protocol. If the server supports multiple transport protocols, there will be multiple NAPTR records, each with a different service value. As per RFC 2915 [NAPTR], the client

discards any records whose services fields are not applicable. For the purposes of this specification, several rules are defined.

- 3.2 A client MUST discard any service fields that identify a resolution service whose value is not "D2X", for values of X that indicate transport protocols supported by the client. The NAPTR processing as described in RFC 2915 will result in discovery of the most preferred transport protocol of the server that is supported by the client, as well as an SRV record for the server.

The domain suffixes in the NAPTR replacement field SHOULD match the domain of the original query.

4. If no NAPTR records are found, the requester queries for those address records for the destination address, '\_diameter.\_sctp'.realm or '\_diameter.\_tcp'.realm. Address records include A RR's, AAAA RR's or other similar records, chosen according to the requestor's network protocol capabilities. If the DNS server returns no address records, the requestor gives up.

If the server is using a site certificate, the domain name in the query and the domain name in the replacement field MUST both be valid based on the site certificate handed out by the server in the TLS or IKE exchange. Similarly, the domain name in the SRV query and the domain name in the target in the SRV record MUST both be valid based on the same site certificate. Otherwise, an attacker could modify the DNS records to contain replacement values in a different domain, and the client could not validate that this was the desired behavior, or the result of an attack

Also, the Diameter Peer MUST check to make sure that the discovered peers are authorized to act in its role. Authentication via IKE or TLS, or validation of DNS RRs via DNSSEC is not sufficient to conclude this. For example, a web server may have obtained a valid TLS certificate, and secured RRs may be included in the DNS, but this does not imply that it is authorized to act as a Diameter Server.

Authorization can be achieved for example, by configuration of a Diameter Server CA. Alternatively this can be achieved by definition of OIDs within TLS or IKE certificates so as to signify Diameter Server authorization.

A dynamically discovered peer causes an entry in the Peer Table (see Section 2.6) to be created. Note that entries created via DNS MUST expire (or be refreshed) within the DNS TTL. If a peer is discovered

outside of the local realm, a routing table entry (see Section 2.7) for the peer's realm is created. The routing table entry's expiration MUST match the peer's expiration value.

### 5.3. Capabilities Exchange

When two Diameter peers establish a transport connection, they MUST exchange the Capabilities Exchange messages, as specified in the peer state machine (see Section 5.6). This message allows the discovery of a peer's identity and its capabilities (protocol version number, supported Diameter applications, security mechanisms, etc.)

The receiver only issues commands to its peers that have advertised support for the Diameter application that defines the command. A Diameter node MUST cache the supported applications in order to ensure that unrecognized commands and/or AVPs are not unnecessarily sent to a peer.

A receiver of a Capabilities-Exchange-Req (CER) message that does not have any applications in common with the sender MUST return a Capabilities-Exchange-Answer (CEA) with the Result-Code AVP set to `DIAMETER_NO_COMMON_APPLICATION`, and SHOULD disconnect the transport layer connection. Note that receiving a CER or CEA from a peer advertising itself as a Relay (see Section 2.4) MUST be interpreted as having common applications with the peer.

Similarly, a receiver of a Capabilities-Exchange-Req (CER) message that does not have any security mechanisms in common with the sender MUST return a Capabilities-Exchange-Answer (CEA) with the Result-Code AVP set to `DIAMETER_NO_COMMON_SECURITY`, and SHOULD disconnect the transport layer connection.

CERs received from unknown peers MAY be silently discarded, or a CEA MAY be issued with the Result-Code AVP set to `DIAMETER_UNKNOWN_PEER`. In both cases, the transport connection is closed. If the local policy permits receiving CERs from unknown hosts, a successful CEA MAY be returned. If a CER from an unknown peer is answered with a successful CEA, the lifetime of the peer entry is equal to the lifetime of the transport connection. In case of a transport failure, all the pending transactions destined to the unknown peer can be discarded.

The CER and CEA messages MUST NOT be proxied, redirected or relayed.

Since the CER/CEA messages cannot be proxied, it is still possible that an upstream agent receives a message for which it has no available peers to handle the application that corresponds to the Command-Code. In such instances, the 'E' bit is set in the answer

message (see Section 7.) with the Result-Code AVP set to `DIAMETER_UNABLE_TO_DELIVER` to inform the downstream to take action (e.g., re-routing request to an alternate peer).

With the exception of the Capabilities-Exchange-Request message, a message of type Request that includes the Auth-Application-Id or Acct-Application-Id AVPs, or a message with an application-specific command code, MAY only be forwarded to a host that has explicitly advertised support for the application (or has advertised the Relay Application Identifier).

#### 5.3.1. Capabilities-Exchange-Request

The Capabilities-Exchange-Request (CER), indicated by the Command-Code set to 257 and the Command Flags' 'R' bit set, is sent to exchange local capabilities. Upon detection of a transport failure, this message MUST NOT be sent to an alternate peer.

When Diameter is run over SCTP [SCTP], which allows for connections to span multiple interfaces and multiple IP addresses, the Capabilities-Exchange-Request message MUST contain one Host-IP-Address AVP for each potential IP address that MAY be locally used when transmitting Diameter messages.

##### Message Format

```
<CER> ::= < Diameter Header: 257, REQ >
    { Origin-Host }
    { Origin-Realm }
    1* { Host-IP-Address }
    { Vendor-Id }
    { Product-Name }
    [ Origin-State-Id ]
    * [ Supported-Vendor-Id ]
    * [ Auth-Application-Id ]
    * [ Inband-Security-Id ]
    * [ Acct-Application-Id ]
    * [ Vendor-Specific-Application-Id ]
    [ Firmware-Revision ]
    * [ AVP ]
```

#### 5.3.2. Capabilities-Exchange-Answer

The Capabilities-Exchange-Answer (CEA), indicated by the Command-Code set to 257 and the Command Flags' 'R' bit cleared, is sent in response to a CER message.

When Diameter is run over SCTP [SCTP], which allows connections to span multiple interfaces, hence, multiple IP addresses, the Capabilities-Exchange-Answer message MUST contain one Host-IP-Address AVP for each potential IP address that MAY be locally used when transmitting Diameter messages.

#### Message Format

```
<CEA> ::= < Diameter Header: 257 >
        { Result-Code }
        { Origin-Host }
        { Origin-Realm }
    1* { Host-IP-Address }
        { Vendor-Id }
        { Product-Name }
        [ Origin-State-Id ]
        [ Error-Message ]
    * [ Failed-AVP ]
    * [ Supported-Vendor-Id ]
    * [ Auth-Application-Id ]
    * [ Inband-Security-Id ]
    * [ Acct-Application-Id ]
    * [ Vendor-Specific-Application-Id ]
        [ Firmware-Revision ]
    * [ AVP ]
```

#### 5.3.3. Vendor-Id AVP

The Vendor-Id AVP (AVP Code 266) is of type Unsigned32 and contains the IANA "SMI Network Management Private Enterprise Codes" [ASSIGNNO] value assigned to the vendor of the Diameter application. In combination with the Supported-Vendor-Id AVP (Section 5.3.6), this MAY be used in order to know which vendor specific attributes may be sent to the peer. It is also envisioned that the combination of the Vendor-Id, Product-Name (Section 5.3.7) and the Firmware-Revision (Section 5.3.4) AVPs MAY provide very useful debugging information.

A Vendor-Id value of zero in the CER or CEA messages is reserved and indicates that this field is ignored.

#### 5.3.4. Firmware-Revision AVP

The Firmware-Revision AVP (AVP Code 267) is of type Unsigned32 and is used to inform a Diameter peer of the firmware revision of the issuing device.

For devices that do not have a firmware revision (general purpose computers running Diameter software modules, for instance), the revision of the Diameter software module may be reported instead.

#### 5.3.5. Host-IP-Address AVP

The Host-IP-Address AVP (AVP Code 257) is of type Address and is used to inform a Diameter peer of the sender's IP address. All source addresses that a Diameter node expects to use with SCTP [SCTP] MUST be advertised in the CER and CEA messages by including a Host-IP-Address AVP for each address. This AVP MUST ONLY be used in the CER and CEA messages.

#### 5.3.6. Supported-Vendor-Id AVP

The Supported-Vendor-Id AVP (AVP Code 265) is of type Unsigned32 and contains the IANA "SMI Network Management Private Enterprise Codes" [ASSIGNNO] value assigned to a vendor other than the device vendor. This is used in the CER and CEA messages in order to inform the peer that the sender supports (a subset of) the vendor-specific AVPs defined by the vendor identified in this AVP.

#### 5.3.7. Product-Name AVP

The Product-Name AVP (AVP Code 269) is of type UTF8String, and contains the vendor assigned name for the product. The Product-Name AVP SHOULD remain constant across firmware revisions for the same product.

### 5.4. Disconnecting Peer connections

When a Diameter node disconnects one of its transport connections, its peer cannot know the reason for the disconnect, and will most likely assume that a connectivity problem occurred, or that the peer has rebooted. In these cases, the peer may periodically attempt to reconnect, as stated in Section 2.1. In the event that the disconnect was a result of either a shortage of internal resources, or simply that the node in question has no intentions of forwarding any Diameter messages to the peer in the foreseeable future, a periodic connection request would not be welcomed. The Disconnection-Reason AVP contains the reason the Diameter node issued the Disconnect-Peer-Request message.

The Disconnect-Peer-Request message is used by a Diameter node to inform its peer of its intent to disconnect the transport layer, and that the peer shouldn't reconnect unless it has a valid reason to do so (e.g., message to be forwarded). Upon receipt of the message, the

Disconnect-Peer-Answer is returned, which SHOULD contain an error if messages have recently been forwarded, and are likely in flight, which would otherwise cause a race condition.

The receiver of the Disconnect-Peer-Answer initiates the transport disconnect.

#### 5.4.1. Disconnect-Peer-Request

The Disconnect-Peer-Request (DPR), indicated by the Command-Code set to 282 and the Command Flags' 'R' bit set, is sent to a peer to inform its intentions to shutdown the transport connection. Upon detection of a transport failure, this message MUST NOT be sent to an alternate peer.

Message Format

```
<DPR> ::= < Diameter Header: 282, REQ >
        { Origin-Host }
        { Origin-Realm }
        { Disconnect-Cause }
```

#### 5.4.2. Disconnect-Peer-Answer

The Disconnect-Peer-Answer (DPA), indicated by the Command-Code set to 282 and the Command Flags' 'R' bit cleared, is sent as a response to the Disconnect-Peer-Request message. Upon receipt of this message, the transport connection is shutdown.

Message Format

```
<DPA> ::= < Diameter Header: 282 >
        { Result-Code }
        { Origin-Host }
        { Origin-Realm }
        [ Error-Message ]
        * [ Failed-AVP ]
```

#### 5.4.3. Disconnect-Cause AVP

The Disconnect-Cause AVP (AVP Code 273) is of type Enumerated. A Diameter node MUST include this AVP in the Disconnect-Peer-Request message to inform the peer of the reason for its intention to shutdown the transport connection. The following values are supported:

```
REBOOTING                                0
    A scheduled reboot is imminent.

BUSY                                     1
    The peer's internal resources are constrained, and it has
    determined that the transport connection needs to be closed.

DO_NOT_WANT_TO_TALK_TO_YOU              2
    The peer has determined that it does not see a need for the
    transport connection to exist, since it does not expect any
    messages to be exchanged in the near future.
```

### 5.5. Transport Failure Detection

Given the nature of the Diameter protocol, it is recommended that transport failures be detected as soon as possible. Detecting such failures will minimize the occurrence of messages sent to unavailable agents, resulting in unnecessary delays, and will provide better failover performance. The Device-Watchdog-Request and Device-Watchdog-Answer messages, defined in this section, are used to proactively detect transport failures.

### 5.5.1. Device-Watchdog-Request

The Device-Watchdog-Request (DWR), indicated by the Command-Code set to 280 and the Command Flags' 'R' bit set, is sent to a peer when no traffic has been exchanged between two peers (see Section 5.5.3). Upon detection of a transport failure, this message MUST NOT be sent to an alternate peer.

## Message Format

```
<DWR> ::= < Diameter Header: 280, REQ >
{
  Origin-Host }
{
  Origin-Realm }
[
  Origin-State-Id ]
```

### 5.5.2. Device-Watchdog-Answer

The Device-Watchdog-Answer (DWA), indicated by the Command-Code set to 280 and the Command Flags' 'R' bit cleared, is sent as a response to the Device-Watchdog-Request message.



## Message Format

```
<DWA> ::= < Diameter Header: 280 >
        { Result-Code }
        { Origin-Host }
        { Origin-Realm }
        [ Error-Message ]
        * [ Failed-AVP ]
        [ Original-State-Id ]
```

## 5.5.3. Transport Failure Algorithm

The transport failure algorithm is defined in [AAATrans]. All Diameter implementations MUST support the algorithm defined in the specification in order to be compliant to the Diameter base protocol.

## 5.5.4. Failover and Failback Procedures

In the event that a transport failure is detected with a peer, it is necessary for all pending request messages to be forwarded to an alternate agent, if possible. This is commonly referred to as failover.

In order for a Diameter node to perform failover procedures, it is necessary for the node to maintain a pending message queue for a given peer. When an answer message is received, the corresponding request is removed from the queue. The Hop-by-Hop Identifier field is used to match the answer with the queued request.

When a transport failure is detected, if possible all messages in the queue are sent to an alternate agent with the T flag set. On booting a Diameter client or agent, the T flag is also set on any records still remaining to be transmitted in non-volatile storage. An example of a case where it is not possible to forward the message to an alternate server is when the message has a fixed destination, and the unavailable peer is the message's final destination (see Destination-Host AVP). Such an error requires that the agent return an answer message with the 'E' bit set and the Result-Code AVP set to DIAMETER\_UNABLE\_TO\_DELIVER.

It is important to note that multiple identical requests or answers MAY be received as a result of a failover. The End-to-End Identifier field in the Diameter header along with the Origin-Host AVP MUST be used to identify duplicate messages.

As described in Section 2.1, a connection request should be periodically attempted with the failed peer in order to re-establish the transport connection. Once a connection has been successfully established, messages can once again be forwarded to the peer. This is commonly referred to as failback.

#### 5.6. Peer State Machine

This section contains a finite state machine that MUST be observed by all Diameter implementations. Each Diameter node MUST follow the state machine described below when communicating with each peer. Multiple actions are separated by commas, and may continue on succeeding lines, as space requires. Similarly, state and next state may also span multiple lines, as space requires.

This state machine is closely coupled with the state machine described in [AAATrans], which is used to open, close, failover, probe, and reopen transport connections. Note in particular that [AAATrans] requires the use of watchdog messages to probe connections. For Diameter, DWR and DWA messages are to be used.

I- is used to represent the initiator (connecting) connection, while the R- is used to represent the responder (listening) connection. The lack of a prefix indicates that the event or action is the same regardless of the connection on which the event occurred.

The stable states that a state machine may be in are Closed, I-Open and R-Open; all other states are intermediate. Note that I-Open and R-Open are equivalent except for whether the initiator or responder transport connection is used for communication.

A CER message is always sent on the initiating connection immediately after the connection request is successfully completed. In the case of an election, one of the two connections will shut down. The responder connection will survive if the Origin-Host of the local Diameter entity is higher than that of the peer; the initiator connection will survive if the peer's Origin-Host is higher. All subsequent messages are sent on the surviving connection. Note that the results of an election on one peer are guaranteed to be the inverse of the results on the other.

For TLS usage, a TLS handshake will begin when both ends are in the open state. If the TLS handshake is successful, all further messages will be sent via TLS. If the handshake fails, both ends move to the closed state.

The state machine constrains only the behavior of a Diameter implementation as seen by Diameter peers through events on the wire.

Any implementation that produces equivalent results is considered compliant.

state	event	action	next state
-----			
Closed	Start	I-Snd-Conn-Req	Wait-Conn-Ack
	R-Conn-CER	R-Accept, Process-CER, R-Snd-CEA	R-Open
Wait-Conn-Ack	I-Rcv-Conn-Ack	I-Snd-CER	Wait-I-CEA
	I-Rcv-Conn-Nack	Cleanup	Closed
	R-Conn-CER	R-Accept, Process-CER	Wait-Conn-Ack/ Elect
	Timeout	Error	Closed
Wait-I-CEA	I-Rcv-CEA	Process-CEA	I-Open
	R-Conn-CER	R-Accept, Process-CER, Elect	Wait>Returns
	I-Peer-Disc	I-Disc	Closed
	I-Rcv-Non-CEA	Error	Closed
	Timeout	Error	Closed
Wait-Conn-Ack/ Elect	I-Rcv-Conn-Ack	I-Snd-CER,Elect	Wait>Returns
	I-Rcv-Conn-Nack	R-Snd-CEA	R-Open
	R-Peer-Disc	R-Disc	Wait-Conn-Ack
	R-Conn-CER	R-Reject	Wait-Conn-Ack/ Elect
	Timeout	Error	Closed
Wait>Returns	Win-Election	I-Disc,R-Snd-CEA	R-Open
	I-Peer-Disc	I-Disc, R-Snd-CEA	R-Open
	I-Rcv-CEA	R-Disc	I-Open
	R-Peer-Disc	R-Disc	Wait-I-CEA
	R-Conn-CER	R-Reject	Wait>Returns
	Timeout	Error	Closed
R-Open	Send-Message	R-Snd-Message	R-Open
	R-Rcv-Message	Process	R-Open
	R-Rcv-DWR	Process-DWR, R-Snd-DWA	R-Open
	R-Rcv-DWA	Process-DWA	R-Open
	R-Conn-CER	R-Reject	R-Open
	Stop	R-Snd-DPR	Closing
	R-Rcv-DPR	R-Snd-DPA, R-Disc	Closed

	R-Peer-Disc	R-Disc	Closed
	R-Rcv-CER	R-Snd-CEA	R-Open
	R-Rcv-CEA	Process-CEA	R-Open
I-Open	Send-Message	I-Snd-Message	I-Open
	I-Rcv-Message	Process	I-Open
	I-Rcv-DWR	Process-DWR,	I-Open
		I-Snd-DWA	
	I-Rcv-DWA	Process-DWA	I-Open
	R-Conn-CER	R-Reject	I-Open
	Stop	I-Snd-DPR	Closing
	I-Rcv-DPR	I-Snd-DPA,	Closed
		I-Disc	
	I-Peer-Disc	I-Disc	Closed
	I-Rcv-CER	I-Snd-CEA	I-Open
	I-Rcv-CEA	Process-CEA	I-Open
Closing	I-Rcv-DPA	I-Disc	Closed
	R-Rcv-DPA	R-Disc	Closed
	Timeout	Error	Closed
	I-Peer-Disc	I-Disc	Closed
	R-Peer-Disc	R-Disc	Closed

#### 5.6.1. Incoming connections

When a connection request is received from a Diameter peer, it is not, in the general case, possible to know the identity of that peer until a CER is received from it. This is because host and port determine the identity of a Diameter peer; and the source port of an incoming connection is arbitrary. Upon receipt of CER, the identity of the connecting peer can be uniquely determined from Origin-Host.

For this reason, a Diameter peer must employ logic separate from the state machine to receive connection requests, accept them, and await CER. Once CER arrives on a new connection, the Origin-Host that identifies the peer is used to locate the state machine associated with that peer, and the new connection and CER are passed to the state machine as an R-Conn-CER event.

The logic that handles incoming connections SHOULD close and discard the connection if any message other than CER arrives, or if an implementation-defined timeout occurs prior to receipt of CER.

Because handling of incoming connections up to and including receipt of CER requires logic, separate from that of any individual state machine associated with a particular peer, it is described separately in this section rather than in the state machine above.

### 5.6.2. Events

Transitions and actions in the automaton are caused by events. In this section, we will ignore the -I and -R prefix, since the actual event would be identical, but would occur on one of two possible connections.

Start	The Diameter application has signaled that a connection should be initiated with the peer.
R-Conn-CER	An acknowledgement is received stating that the transport connection has been established, and the associated CER has arrived.
Rcv-Conn-Ack	A positive acknowledgement is received confirming that the transport connection is established.
Rcv-Conn-Nack	A negative acknowledgement was received stating that the transport connection was not established.
Timeout	An application-defined timer has expired while waiting for some event.
Rcv-CER	A CER message from the peer was received.
Rcv-CEA	A CEA message from the peer was received.
Rcv-Non-CEA	A message other than CEA from the peer was received.
Peer-Disc	A disconnection indication from the peer was received.
Rcv-DPR	A DPR message from the peer was received.
Rcv-DPA	A DPA message from the peer was received.
Win-Election	An election was held, and the local node was the winner.
Send-Message	A message is to be sent.
Rcv-Message	A message other than CER, CEA, DPR, DPA, DWR or DWA was received.
Stop	The Diameter application has signaled that a connection should be terminated (e.g., on system shutdown).

### 5.6.3. Actions

Actions in the automaton are caused by events and typically indicate the transmission of packets and/or an action to be taken on the connection. In this section we will ignore the I- and R-prefix, since the actual action would be identical, but would occur on one of two possible connections.

Snd-Conn-Req	A transport connection is initiated with the peer.
Accept	The incoming connection associated with the R-Conn-CER is accepted as the responder connection.
Reject	The incoming connection associated with the R-Conn-CER is disconnected.
Process-CER	The CER associated with the R-Conn-CER is processed.
Snd-CER	A CER message is sent to the peer.
Snd-CEA	A CEA message is sent to the peer.
Cleanup	If necessary, the connection is shutdown, and any local resources are freed.
Error	The transport layer connection is disconnected, either politely or abortively, in response to an error condition. Local resources are freed.
Process-CEA	A received CEA is processed.
Snd-DPR	A DPR message is sent to the peer.
Snd-DPA	A DPA message is sent to the peer.
Disc	The transport layer connection is disconnected, and local resources are freed.
Elect	An election occurs (see Section 5.6.4 for more information).
Snd-Message	A message is sent.
Snd-DWR	A DWR message is sent.
Snd-DWA	A DWA message is sent.
Process-DWR	The DWR message is serviced.

Process-DWA      The DWA message is serviced.

Process          A message is serviced.

#### 5.6.4. The Election Process

The election is performed on the responder. The responder compares the Origin-Host received in the CER sent by its peer with its own Origin-Host. If the local Diameter entity's Origin-Host is higher than the peer's, a Win-Election event is issued locally.

The comparison proceeds by considering the shorter OctetString to be padded with zeros so that its length is the same as the length of the longer, then performing an octet-by-octet unsigned comparison with the first octet being most significant. Any remaining octets are assumed to have value 0x80.

### 6. Diameter message processing

This section describes how Diameter requests and answers are created and processed.

#### 6.1. Diameter Request Routing Overview

A request is sent towards its final destination using a combination of the Destination-Realm and Destination-Host AVPs, in one of these three combinations:

- a request that is not able to be proxied (such as CER) MUST NOT contain either Destination-Realm or Destination-Host AVPs.
- a request that needs to be sent to a home server serving a specific realm, but not to a specific server (such as the first request of a series of round-trips), MUST contain a Destination-Realm AVP, but MUST NOT contain a Destination-Host AVP.
- otherwise, a request that needs to be sent to a specific home server among those serving a given realm, MUST contain both the Destination-Realm and Destination-Host AVPs.

The Destination-Host AVP is used as described above when the destination of the request is fixed, which includes:

- Authentication requests that span multiple round trips
- A Diameter message that uses a security mechanism that makes use of a pre-established session key shared between the source and the final destination of the message.

- Server initiated messages that MUST be received by a specific Diameter client (e.g., access device), such as the Abort-Session-Request message, which is used to request that a particular user's session be terminated.

Note that an agent can forward a request to a host described in the Destination-Host AVP only if the host in question is included in its peer table (see Section 2.7). Otherwise, the request is routed based on the Destination-Realm only (see Sections 6.1.6).

The Destination-Realm AVP MUST be present if the message is proxiable. Request messages that may be forwarded by Diameter agents (proxies, redirects or relays) MUST also contain an Acct-Application-Id AVP, an Auth-Application-Id AVP or a Vendor-Specific-Application-Id AVP. A message that MUST NOT be forwarded by Diameter agents (proxies, redirects or relays) MUST not include the Destination-Realm in its ABNF. The value of the Destination-Realm AVP MAY be extracted from the User-Name AVP, or other application-specific methods.

When a message is received, the message is processed in the following order:

1. If the message is destined for the local host, the procedures listed in Section 6.1.4 are followed.
2. If the message is intended for a Diameter peer with whom the local host is able to directly communicate, the procedures listed in Section 6.1.5 are followed. This is known as Request Forwarding.
3. The procedures listed in Section 6.1.6 are followed, which is known as Request Routing.
4. If none of the above is successful, an answer is returned with the Result-Code set to DIAMETER\_UNABLE\_TO\_DELIVER, with the E-bit set.

For routing of Diameter messages to work within an administrative domain, all Diameter nodes within the realm MUST be peers.

Note the processing rules contained in this section are intended to be used as general guidelines to Diameter developers. Certain implementations MAY use different methods than the ones described here, and still comply with the protocol specification. See Section 7 for more detail on error handling.



#### 6.1.1.1. Originating a Request

When creating a request, in addition to any other procedures described in the application definition for that specific request, the following procedures MUST be followed:

- the Command-Code is set to the appropriate value
- the 'R' bit is set
- the End-to-End Identifier is set to a locally unique value
- the Origin-Host and Origin-Realm AVPs MUST be set to the appropriate values, used to identify the source of the message
- the Destination-Host and Destination-Realm AVPs MUST be set to the appropriate values as described in Section 6.1.
- an Acct-Application-Id AVP, an Auth-Application-Id or a Vendor-Specific-Application-Id AVP must be included if the request is proxiable.

#### 6.1.1.2. Sending a Request

When sending a request, originated either locally, or as the result of a forwarding or routing operation, the following procedures MUST be followed:

- the Hop-by-Hop Identifier should be set to a locally unique value
- The message should be saved in the list of pending requests.

Other actions to perform on the message based on the particular role the agent is playing are described in the following sections.

#### 6.1.1.3. Receiving Requests

A relay or proxy agent MUST check for forwarding loops when receiving requests. A loop is detected if the server finds its own identity in a Route-Record AVP. When such an event occurs, the agent MUST answer with the Result-Code AVP set to `DIAMETER_LOOP_DETECTED`.

#### 6.1.1.4. Processing Local Requests

A request is known to be for local consumption when one of the following conditions occur:

- The Destination-Host AVP contains the local host's identity,

- The Destination-Host AVP is not present, the Destination-Realm AVP contains a realm the server is configured to process locally, and the Diameter application is locally supported, or
- Both the Destination-Host and the Destination-Realm are not present.

When a request is locally processed, the rules in Section 6.2 should be used to generate the corresponding answer.

#### 6.1.5. Request Forwarding

Request forwarding is done using the Diameter Peer Table. The Diameter peer table contains all of the peers that the local node is able to directly communicate with.

When a request is received, and the host encoded in the Destination-Host AVP is one that is present in the peer table, the message SHOULD be forwarded to the peer.

#### 6.1.6. Request Routing

Diameter request message routing is done via realms and applications. A Diameter message that may be forwarded by Diameter agents (proxies, redirects or relays) MUST include the target realm in the Destination-Realm AVP and one of the application identification AVPs Auth-Application-Id, Acct-Application-Id or Vendor-Specific-Application-Id. The realm MAY be retrieved from the User-Name AVP, which is in the form of a Network Access Identifier (NAI). The realm portion of the NAI is inserted in the Destination-Realm AVP.

Diameter agents MAY have a list of locally supported realms and applications, and MAY have a list of externally supported realms and applications. When a request is received that includes a realm and/or application that is not locally supported, the message is routed to the peer configured in the Realm Routing Table (see Section 2.7).

#### 6.1.7. Redirecting requests

When a redirect agent receives a request whose routing entry is set to REDIRECT, it MUST reply with an answer message with the 'E' bit set, while maintaining the Hop-by-Hop Identifier in the header, and include the Result-Code AVP to DIAMETER\_REDIRECT\_INDICATION. Each of the servers associated with the routing entry are added in separate Redirect-Host AVP.

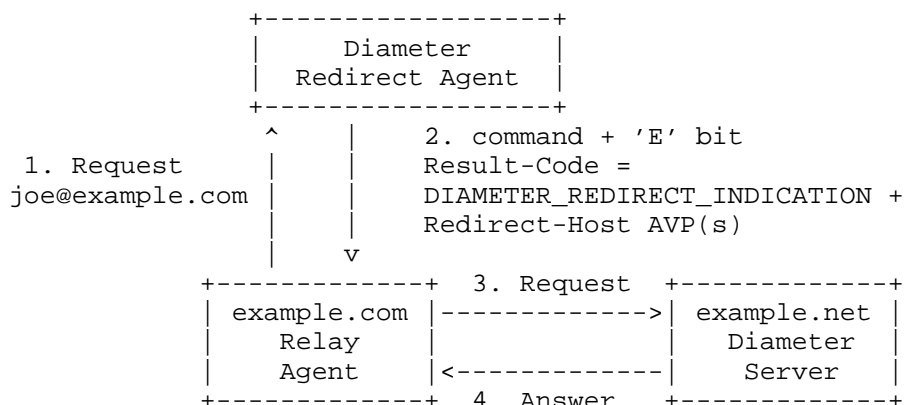


Figure 5: Diameter Redirect Agent

The receiver of the answer message with the 'E' bit set, and the Result-Code AVP set to DIAMETER\_REDIRECT\_INDICATION uses the hop-by-hop field in the Diameter header to identify the request in the pending message queue (see Section 5.3) that is to be redirected. If no transport connection exists with the new agent, one is created, and the request is sent directly to it.

Multiple Redirect-Host AVPs are allowed. The receiver of the answer message with the 'E' bit set selects exactly one of these hosts as the destination of the redirected message.

#### 6.1.8. Relaying and Proxying Requests

A relay or proxy agent MUST append a Route-Record AVP to all requests forwarded. The AVP contains the identity of the peer the request was received from.

The Hop-by-Hop identifier in the request is saved, and replaced with a locally unique value. The source of the request is also saved, which includes the IP address, port and protocol.

A relay or proxy agent MAY include the Proxy-Info AVP in requests if it requires access to any local state information when the corresponding response is received. Proxy-Info AVP has certain security implications and SHOULD contain an embedded HMAC with a node-local key. Alternatively, it MAY simply use local storage to store state information.

The message is then forwarded to the next hop, as identified in the Realm Routing Table.

Figure 6 provides an example of message routing using the procedures listed in these sections.

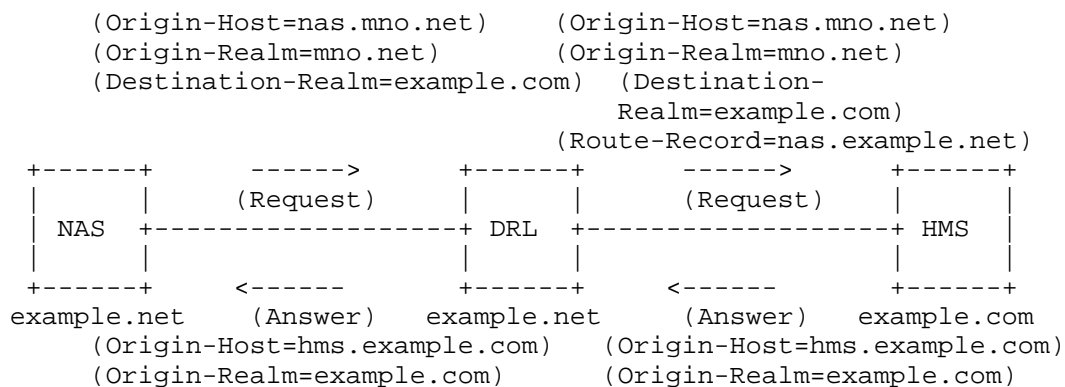


Figure 6: Routing of Diameter messages

## 6.2. Diameter Answer Processing

When a request is locally processed, the following procedures MUST be applied to create the associated answer, in addition to any additional procedures that MAY be discussed in the Diameter application defining the command:

- The same Hop-by-Hop identifier in the request is used in the answer.
- The local host's identity is encoded in the Origin-Host AVP.
- The Destination-Host and Destination-Realm AVPs MUST NOT be present in the answer message.
- The Result-Code AVP is added with its value indicating success or failure.
- If the Session-Id is present in the request, it MUST be included in the answer.
- Any Proxy-Info AVPs in the request MUST be added to the answer message, in the same order they were present in the request.
- The 'P' bit is set to the same value as the one in the request.
- The same End-to-End identifier in the request is used in the answer.

Note that the error messages (see Section 7.3) are also subjected to the above processing rules.

#### 6.2.1. Processing received Answers

A Diameter client or proxy MUST match the Hop-by-Hop Identifier in an answer received against the list of pending requests. The corresponding message should be removed from the list of pending requests. It SHOULD ignore answers received that do not match a known Hop-by-Hop Identifier.

#### 6.2.2. Relaying and Proxying Answers

If the answer is for a request which was proxied or relayed, the agent MUST restore the original value of the Diameter header's Hop-by-Hop Identifier field.

If the last Proxy-Info AVP in the message is targeted to the local Diameter server, the AVP MUST be removed before the answer is forwarded.

If a relay or proxy agent receives an answer with a Result-Code AVP indicating a failure, it MUST NOT modify the contents of the AVP. Any additional local errors detected SHOULD be logged, but not reflected in the Result-Code AVP. If the agent receives an answer message with a Result-Code AVP indicating success, and it wishes to modify the AVP to indicate an error, it MUST modify the Result-Code AVP to contain the appropriate error in the message destined towards the access device as well as include the Error-Reporting-Host AVP and it MUST issue an STR on behalf of the access device.

The agent MUST then send the answer to the host that it received the original request from.

#### 6.3. Origin-Host AVP

The Origin-Host AVP (AVP Code 264) is of type DiameterIdentity, and MUST be present in all Diameter messages. This AVP identifies the endpoint that originated the Diameter message. Relay agents MUST NOT modify this AVP.

The value of the Origin-Host AVP is guaranteed to be unique within a single host.

Note that the Origin-Host AVP may resolve to more than one address as the Diameter peer may support more than one address.

This AVP SHOULD be placed as close to the Diameter header as possible. 6.10

#### 6.4. Origin-Realm AVP

The Origin-Realm AVP (AVP Code 296) is of type DiameterIdentity. This AVP contains the Realm of the originator of any Diameter message and MUST be present in all messages.

This AVP SHOULD be placed as close to the Diameter header as possible.

#### 6.5. Destination-Host AVP

The Destination-Host AVP (AVP Code 293) is of type DiameterIdentity. This AVP MUST be present in all unsolicited agent initiated messages, MAY be present in request messages, and MUST NOT be present in Answer messages.

The absence of the Destination-Host AVP will cause a message to be sent to any Diameter server supporting the application within the realm specified in Destination-Realm AVP.

This AVP SHOULD be placed as close to the Diameter header as possible.

#### 6.6. Destination-Realm AVP

The Destination-Realm AVP (AVP Code 283) is of type DiameterIdentity, and contains the realm the message is to be routed to. The Destination-Realm AVP MUST NOT be present in Answer messages. Diameter Clients insert the realm portion of the User-Name AVP. Diameter servers initiating a request message use the value of the Origin-Realm AVP from a previous message received from the intended target host (unless it is known a priori). When present, the Destination-Realm AVP is used to perform message routing decisions.

Request messages whose ABNF does not list the Destination-Realm AVP as a mandatory AVP are inherently non-routable messages.

This AVP SHOULD be placed as close to the Diameter header as possible.

#### 6.7. Routing AVPs

The AVPs defined in this section are Diameter AVPs used for routing purposes. These AVPs change as Diameter messages are processed by agents, and therefore MUST NOT be protected by end-to-end security.

#### 6.7.1. Route-Record AVP

The Route-Record AVP (AVP Code 282) is of type DiameterIdentity. The identity added in this AVP MUST be the same as the one received in the Origin-Host of the Capabilities Exchange message.

#### 6.7.2. Proxy-Info AVP

The Proxy-Info AVP (AVP Code 284) is of type Grouped. The Grouped Data field has the following ABNF grammar:

```
Proxy-Info ::= < AVP Header: 284 >
               { Proxy-Host }
               { Proxy-State }
               * [ AVP ]
```

#### 6.7.3. Proxy-Host AVP

The Proxy-Host AVP (AVP Code 280) is of type DiameterIdentity. This AVP contains the identity of the host that added the Proxy-Info AVP.

#### 6.7.4. Proxy-State AVP

The Proxy-State AVP (AVP Code 33) is of type OctetString, and contains state local information, and MUST be treated as opaque data.

#### 6.8. Auth-Application-Id AVP

The Auth-Application-Id AVP (AVP Code 258) is of type Unsigned32 and is used in order to advertise support of the Authentication and Authorization portion of an application (see Section 2.4). The Auth-Application-Id MUST also be present in all Authentication and/or Authorization messages that are defined in a separate Diameter specification and have an Application ID assigned.

#### 6.9. Acct-Application-Id AVP

The Acct-Application-Id AVP (AVP Code 259) is of type Unsigned32 and is used in order to advertise support of the Accounting portion of an application (see Section 2.4). The Acct-Application-Id MUST also be present in all Accounting messages. Exactly one of the Auth-Application-Id and Acct-Application-Id AVPs MAY be present.

#### 6.10. Inband-Security-Id AVP

The Inband-Security-Id AVP (AVP Code 299) is of type Unsigned32 and is used in order to advertise support of the Security portion of the application.

Currently, the following values are supported, but there is ample room to add new security Ids.

NO_INBAND_SECURITY	0
This peer does not support TLS. This is the default value, if the AVP is omitted.	
TLS	1
This node supports TLS security, as defined by [TLS].	

#### 6.11. Vendor-Specific-Application-Id AVP

The Vendor-Specific-Application-Id AVP (AVP Code 260) is of type Grouped and is used to advertise support of a vendor-specific Diameter Application. Exactly one of the Auth-Application-Id and Acct-Application-Id AVPs MAY be present.

This AVP MUST also be present as the first AVP in all experimental commands defined in the vendor-specific application.

This AVP SHOULD be placed as close to the Diameter header as possible.

AVP Format

```
<Vendor-Specific-Application-Id> ::= < AVP Header: 260 >
                                   1* [ Vendor-Id ]
                                   0*1{ Auth-Application-Id }
                                   0*1{ Acct-Application-Id }
```

#### 6.12. Redirect-Host AVP

One or more of instances of this AVP MUST be present if the answer message's 'E' bit is set and the Result-Code AVP is set to DIAMETER\_REDIRECT\_INDICATION.

Upon receiving the above, the receiving Diameter node SHOULD forward the request directly to one of the hosts identified in these AVPs. The server contained in the selected Redirect-Host AVP SHOULD be used for all messages pertaining to this session.

#### 6.13. Redirect-Host-Usage AVP

The Redirect-Host-Usage AVP (AVP Code 261) is of type Enumerated. This AVP MAY be present in answer messages whose 'E' bit is set and the Result-Code AVP is set to DIAMETER\_REDIRECT\_INDICATION.



When present, this AVP dictates how the routing entry resulting from the Redirect-Host is to be used. The following values are supported:

- DONT\_CACHE 0  
The host specified in the Redirect-Host AVP should not be cached. This is the default value.
- ALL\_SESSION 1  
All messages within the same session, as defined by the same value of the Session-ID AVP MAY be sent to the host specified in the Redirect-Host AVP.
- ALL\_REALM 2  
All messages destined for the realm requested MAY be sent to the host specified in the Redirect-Host AVP.
- REALM\_AND\_APPLICATION 3  
All messages for the application requested to the realm specified MAY be sent to the host specified in the Redirect-Host AVP.
- ALL\_APPLICATION 4  
All messages for the application requested MAY be sent to the host specified in the Redirect-Host AVP.
- ALL\_HOST 5  
All messages that would be sent to the host that generated the Redirect-Host MAY be sent to the host specified in the Redirect-Host AVP.
- ALL\_USER 6  
All messages for the user requested MAY be sent to the host specified in the Redirect-Host AVP.

#### 6.14. Redirect-Max-Cache-Time AVP

The Redirect-Max-Cache-Time AVP (AVP Code 262) is of type Unsigned32. This AVP MUST be present in answer messages whose 'E' bit is set, the Result-Code AVP is set to DIAMETER\_REDIRECT\_INDICATION and the Redirect-Host-Usage AVP set to a non-zero value.

This AVP contains the maximum number of seconds the peer and route table entries, created as a result of the Redirect-Host, will be cached. Note that once a host created due to a redirect indication is no longer reachable, any associated peer and routing table entries MUST be deleted.

## 6.15. E2E-Sequence AVP

The E2E-Sequence AVP (AVP Code 300) provides anti-replay protection for end to end messages and is of type grouped. It contains a random value (an OctetString with a nonce) and counter (an Integer). For each end-to-end peer with which a node communicates (or remembers communicating) a different nonce value **MUST** be used and the counter is initiated at zero and increases by one each time this AVP is emitted to that peer. This AVP **MUST** be included in all messages which use end-to-end protection (e.g., CMS signing or encryption).

## 7. Error Handling

There are two different types of errors in Diameter; protocol and application errors. A protocol error is one that occurs at the base protocol level, and MAY require per hop attention (e.g., message routing error). Application errors, on the other hand, generally occur due to a problem with a function specified in a Diameter application (e.g., user authentication, Missing AVP).

Result-Code AVP values that are used to report protocol errors MUST only be present in answer messages whose 'E' bit is set. When a request message is received that causes a protocol error, an answer message is returned with the 'E' bit set, and the Result-Code AVP is set to the appropriate protocol error value. As the answer is sent back towards the originator of the request, each proxy or relay agent MAY take action on the message.

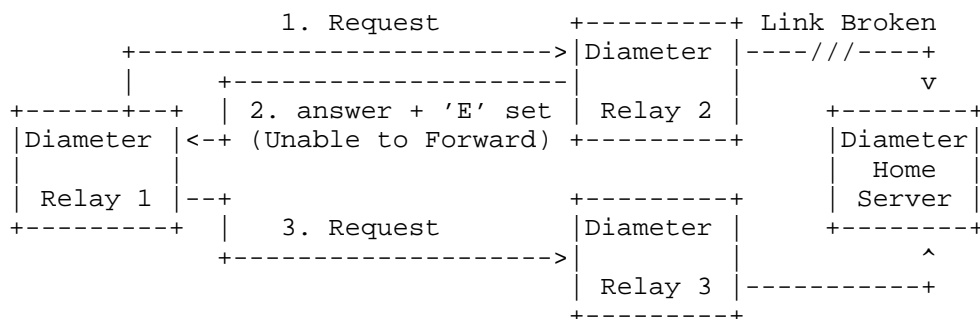


Figure 7: Example of Protocol Error causing answer message

Figure 7 provides an example of a message forwarded upstream by a Diameter relay. When the message is received by Relay 2, and it detects that it cannot forward the request to the home server, an answer message is returned with the 'E' bit set and the Result-Code AVP set to DIAMETER\_UNABLE\_TO\_DELIVER. Given that this error falls

within the protocol error category, Relay 1 would take special action, and given the error, attempt to route the message through its alternate Relay 3.

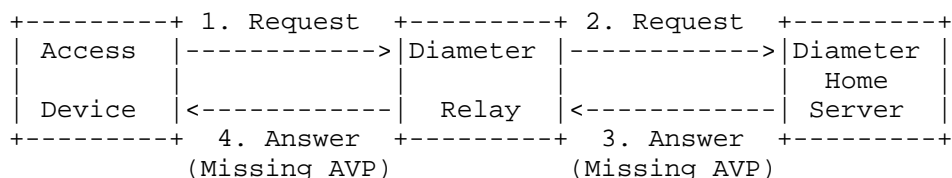


Figure 8: Example of Application Error Answer message

Figure 8 provides an example of a Diameter message that caused an application error. When application errors occur, the Diameter entity reporting the error clears the 'R' bit in the Command Flags, and adds the Result-Code AVP with the proper value. Application errors do not require any proxy or relay agent involvement, and therefore the message would be forwarded back to the originator of the request.

There are certain Result-Code AVP application errors that require additional AVPs to be present in the answer. In these cases, the Diameter node that sets the Result-Code AVP to indicate the error MUST add the AVPs. Examples are:

- An unrecognized AVP is received with the 'M' bit (Mandatory bit) set, causes an answer to be sent with the Result-Code AVP set to `DIAMETER_AVP_UNSUPPORTED`, and the Failed-AVP AVP containing the offending AVP.
- An AVP that is received with an unrecognized value causes an answer to be returned with the Result-Code AVP set to `DIAMETER_INVALID_AVP_VALUE`, with the Failed-AVP AVP containing the AVP causing the error.
- A command is received with an AVP that is omitted, yet is mandatory according to the command's ABNF. The receiver issues an answer with the Result-Code set to `DIAMETER_MISSING_AVP`, and creates an AVP with the AVP Code and other fields set as expected in the missing AVP. The created AVP is then added to the Failed-AVP AVP.

The Result-Code AVP describes the error that the Diameter node encountered in its processing. In case there are multiple errors, the Diameter node MUST report only the first error it encountered

(detected possibly in some implementation dependent order). The specific errors that can be described by this AVP are described in the following section.

### 7.1. Result-Code AVP

The Result-Code AVP (AVP Code 268) is of type Unsigned32 and indicates whether a particular request was completed successfully or whether an error occurred. All Diameter answer messages defined in IETF applications MUST include one Result-Code AVP. A non-successful Result-Code AVP (one containing a non 2xxx value other than `DIAMETER_REDIRECT_INDICATION`) MUST include the Error-Reporting-Host AVP if the host setting the Result-Code AVP is different from the identity encoded in the Origin-Host AVP.

The Result-Code data field contains an IANA-managed 32-bit address space representing errors (see Section 11.4). Diameter provides the following classes of errors, all identified by the thousands digit in the decimal notation:

- 1xxx (Informational)
- 2xxx (Success)
- 3xxx (Protocol Errors)
- 4xxx (Transient Failures)
- 5xxx (Permanent Failure)

A non-recognized class (one whose first digit is not defined in this section) MUST be handled as a permanent failure.

#### 7.1.1. Informational

Errors that fall within this category are used to inform the requester that a request could not be satisfied, and additional action is required on its part before access is granted.

`DIAMETER_MULTI_ROUND_AUTH` 1001

This informational error is returned by a Diameter server to inform the access device that the authentication mechanism being used requires multiple round trips, and a subsequent request needs to be issued in order for access to be granted.

#### 7.1.2. Success

Errors that fall within the Success category are used to inform a peer that a request has been successfully completed.

`DIAMETER_SUCCESS` 2001

The Request was successfully completed.

**DIAMETER\_LIMITED\_SUCCESS** 2002

When returned, the request was successfully completed, but additional processing is required by the application in order to provide service to the user.

**7.1.3. Protocol Errors**

Errors that fall within the Protocol Error category SHOULD be treated on a per-hop basis, and Diameter proxies MAY attempt to correct the error, if it is possible. Note that these and only these errors MUST only be used in answer messages whose 'E' bit is set.

**DIAMETER\_COMMAND\_UNSUPPORTED** 3001

The Request contained a Command-Code that the receiver did not recognize or support. This MUST be used when a Diameter node receives an experimental command that it does not understand.

**DIAMETER\_UNABLE\_TO\_DELIVER** 3002

This error is given when Diameter can not deliver the message to the destination, either because no host within the realm supporting the required application was available to process the request, or because Destination-Host AVP was given without the associated Destination-Realm AVP.

**DIAMETER\_REALM\_NOT\_SERVED** 3003

The intended realm of the request is not recognized.

**DIAMETER\_TOO\_BUSY** 3004

When returned, a Diameter node SHOULD attempt to send the message to an alternate peer. This error MUST only be used when a specific server is requested, and it cannot provide the requested service.

**DIAMETER\_LOOP\_DETECTED** 3005

An agent detected a loop while trying to get the message to the intended recipient. The message MAY be sent to an alternate peer, if one is available, but the peer reporting the error has identified a configuration problem.

**DIAMETER\_REDIRECT\_INDICATION** 3006

A redirect agent has determined that the request could not be satisfied locally and the initiator of the request should direct the request directly to the server, whose contact information has been added to the response. When set, the Redirect-Host AVP MUST be present.

**DIAMETER\_APPLICATION\_UNSUPPORTED** 3007

A request was sent for an application that is not supported.

**DIAMETER\_INVALID\_HDR\_BITS** 3008

A request was received whose bits in the Diameter header were either set to an invalid combination, or to a value that is inconsistent with the command code's definition.

**DIAMETER\_INVALID\_AVP\_BITS** 3009

A request was received that included an AVP whose flag bits are set to an unrecognized value, or that is inconsistent with the AVP's definition.

**DIAMETER\_UNKNOWN\_PEER** 3010

A CER was received from an unknown peer.

**7.1.4. Transient Failures**

Errors that fall within the transient failures category are used to inform a peer that the request could not be satisfied at the time it was received, but MAY be able to satisfy the request in the future.

**DIAMETER\_AUTHENTICATION\_REJECTED** 4001

The authentication process for the user failed, most likely due to an invalid password used by the user. Further attempts MUST only be tried after prompting the user for a new password.

**DIAMETER\_OUT\_OF\_SPACE** 4002

A Diameter node received the accounting request but was unable to commit it to stable storage due to a temporary lack of space.

**ELECTION\_LOST** 4003

The peer has determined that it has lost the election process and has therefore disconnected the transport connection.

**7.1.5. Permanent Failures**

Errors that fall within the permanent failures category are used to inform the peer that the request failed, and should not be attempted again.

**DIAMETER\_AVP\_UNSUPPORTED** 5001

The peer received a message that contained an AVP that is not recognized or supported and was marked with the Mandatory bit. A Diameter message with this error MUST contain one or more Failed-AVP AVP containing the AVPs that caused the failure.

**DIAMETER\_UNKNOWN\_SESSION\_ID** 5002

The request contained an unknown Session-Id.

**DIAMETER\_AUTHORIZATION\_REJECTED 5003**

A request was received for which the user could not be authorized. This error could occur if the service requested is not permitted to the user.

**DIAMETER\_INVALID\_AVP\_VALUE 5004**

The request contained an AVP with an invalid value in its data portion. A Diameter message indicating this error MUST include the offending AVPs within a Failed-AVP AVP.

**DIAMETER\_MISSING\_AVP 5005**

The request did not contain an AVP that is required by the Command Code definition. If this value is sent in the Result-Code AVP, a Failed-AVP AVP SHOULD be included in the message. The Failed-AVP AVP MUST contain an example of the missing AVP complete with the Vendor-Id if applicable. The value field of the missing AVP should be of correct minimum length and contain zeroes.

**DIAMETER\_RESOURCES\_EXCEEDED 5006**

A request was received that cannot be authorized because the user has already expended allowed resources. An example of this error condition is a user that is restricted to one dial-up PPP port, attempts to establish a second PPP connection.

**DIAMETER\_CONTRADICTING\_AVPS 5007**

The Home Diameter server has detected AVPs in the request that contradicted each other, and is not willing to provide service to the user. One or more Failed-AVP AVPs MUST be present, containing the AVPs that contradicted each other.

**DIAMETER\_AVP\_NOT\_ALLOWED 5008**

A message was received with an AVP that MUST NOT be present. The Failed-AVP AVP MUST be included and contain a copy of the offending AVP.

**DIAMETER\_AVP\_OCCURS\_TOO\_MANY\_TIMES 5009**

A message was received that included an AVP that appeared more often than permitted in the message definition. The Failed-AVP AVP MUST be included and contain a copy of the first instance of the offending AVP that exceeded the maximum number of occurrences

**DIAMETER\_NO\_COMMON\_APPLICATION 5010**

This error is returned when a CER message is received, and there are no common applications supported between the peers.

**DIAMETER\_UNSUPPORTED\_VERSION 5011**

This error is returned when a request was received, whose version number is unsupported.

**DIAMETER\_UNABLE\_TO\_COMPLY 5012**

This error is returned when a request is rejected for unspecified reasons.

**DIAMETER\_INVALID\_BIT\_IN\_HEADER 5013**

This error is returned when an unrecognized bit in the Diameter header is set to one (1).

**DIAMETER\_INVALID\_AVP\_LENGTH 5014**

The request contained an AVP with an invalid length. A Diameter message indicating this error MUST include the offending AVPs within a Failed-AVP AVP.

**DIAMETER\_INVALID\_MESSAGE\_LENGTH 5015**

This error is returned when a request is received with an invalid message length.

**DIAMETER\_INVALID\_AVP\_BIT\_COMBO 5016**

The request contained an AVP with which is not allowed to have the given value in the AVP Flags field. A Diameter message indicating this error MUST include the offending AVPs within a Failed-AVP AVP.

**DIAMETER\_NO\_COMMON\_SECURITY 5017**

This error is returned when a CER message is received, and there are no common security mechanisms supported between the peers. A Capabilities-Exchange-Answer (CEA) MUST be returned with the Result-Code AVP set to DIAMETER\_NO\_COMMON\_SECURITY.

## 7.2. Error Bit

The 'E' (Error Bit) in the Diameter header is set when the request caused a protocol-related error (see Section 7.1.3). A message with the 'E' bit MUST NOT be sent as a response to an answer message. Note that a message with the 'E' bit set is still subjected to the processing rules defined in Section 6.2. When set, the answer message will not conform to the ABNF specification for the command, and will instead conform to the following ABNF:



## Message Format

```
<answer-message> ::= < Diameter Header: code, ERR [PXY] >
    0*1< Session-Id >
        { Origin-Host }
        { Origin-Realm }
        { Result-Code }
        [ Origin-State-Id ]
        [ Error-Reporting-Host ]
        [ Proxy-Info ]
    * [ AVP ]
```

Note that the code used in the header is the same than the one found in the request message, but with the 'R' bit cleared and the 'E' bit set. The 'P' bit in the header is set to the same value as the one found in the request message.

## 7.3. Error-Message AVP

The Error-Message AVP (AVP Code 281) is of type UTF8String. It MAY accompany a Result-Code AVP as a human readable error message. The Error-Message AVP is not intended to be useful in real-time, and SHOULD NOT be expected to be parsed by network entities.

## 7.4. Error-Reporting-Host AVP

The Error-Reporting-Host AVP (AVP Code 294) is of type DiameterIdentity. This AVP contains the identity of the Diameter host that sent the Result-Code AVP to a value other than 2001 (Success), only if the host setting the Result-Code is different from the one encoded in the Origin-Host AVP. This AVP is intended to be used for troubleshooting purposes, and MUST be set when the Result-Code AVP indicates a failure.

## 7.5. Failed-AVP AVP

The Failed-AVP AVP (AVP Code 279) is of type Grouped and provides debugging information in cases where a request is rejected or not fully processed due to erroneous information in a specific AVP. The value of the Result-Code AVP will provide information on the reason for the Failed-AVP AVP.

The possible reasons for this AVP are the presence of an improperly constructed AVP, an unsupported or unrecognized AVP, an invalid AVP value, the omission of a required AVP, the presence of an explicitly excluded AVP (see tables in Section 10), or the presence of two or more occurrences of an AVP which is restricted to 0, 1, or 0-1 occurrences.

A Diameter message MAY contain one Failed-AVP AVP, containing the entire AVP that could not be processed successfully. If the failure reason is omission of a required AVP, an AVP with the missing AVP code, the missing vendor id, and a zero filled payload of the minimum required length for the omitted AVP will be added.

#### AVP Format

```
<Failed-AVP> ::= < AVP Header: 279 >
                1* {AVP}
```

### 7.6. Experimental-Result AVP

The Experimental-Result AVP (AVP Code 297) is of type Grouped, and indicates whether a particular vendor-specific request was completed successfully or whether an error occurred. Its Data field has the following ABNF grammar:

#### AVP Format

```
Experimental-Result ::= < AVP Header: 297 >
                      { Vendor-Id }
                      { Experimental-Result-Code }
```

The Vendor-Id AVP (see Section 5.3.3) in this grouped AVP identifies the vendor responsible for the assignment of the result code which follows. All Diameter answer messages defined in vendor-specific applications MUST include either one Result-Code AVP or one Experimental-Result AVP.

### 7.7. Experimental-Result-Code AVP

The Experimental-Result-Code AVP (AVP Code 298) is of type Unsigned32 and contains a vendor-assigned value representing the result of processing the request.

It is recommended that vendor-specific result codes follow the same conventions given for the Result-Code AVP regarding the different types of result codes and the handling of errors (for non 2xxx values).

## 8. Diameter User Sessions

Diameter can provide two different types of services to applications. The first involves authentication and authorization, and can optionally make use of accounting. The second only makes use of accounting.

When a service makes use of the authentication and/or authorization portion of an application, and a user requests access to the network, the Diameter client issues an auth request to its local server. The auth request is defined in a service specific Diameter application (e.g., NASREQ). The request contains a Session-Id AVP, which is used in subsequent messages (e.g., subsequent authorization, accounting, etc) relating to the user's session. The Session-Id AVP is a means for the client and servers to correlate a Diameter message with a user session.

When a Diameter server authorizes a user to use network resources for a finite amount of time, and it is willing to extend the authorization via a future request, it MUST add the Authorization-Lifetime AVP to the answer message. The Authorization-Lifetime AVP defines the maximum number of seconds a user MAY make use of the resources before another authorization request is expected by the server. The Auth-Grace-Period AVP contains the number of seconds following the expiration of the Authorization-Lifetime, after which the server will release all state information related to the user's session. Note that if payment for services is expected by the serving realm from the user's home realm, the Authorization-Lifetime AVP, combined with the Auth-Grace-Period AVP, implies the maximum length of the session the home realm is willing to be fiscally responsible for. Services provided past the expiration of the Authorization-Lifetime and Auth-Grace-Period AVPs are the responsibility of the access device. Of course, the actual cost of services rendered is clearly outside the scope of the protocol.

An access device that does not expect to send a re-authorization or a session termination request to the server MAY include the Auth-Session-State AVP with the value set to NO\_STATE\_MAINTAINED as a hint to the server. If the server accepts the hint, it agrees that since no session termination message will be received once service to the user is terminated, it cannot maintain state for the session. If the answer message from the server contains a different value in the Auth-Session-State AVP (or the default value if the AVP is absent), the access device MUST follow the server's directives. Note that the value NO\_STATE\_MAINTAINED MUST NOT be set in subsequent re-authorization requests and answers.

The base protocol does not include any authorization request messages, since these are largely application-specific and are defined in a Diameter application document. However, the base protocol does define a set of messages that is used to terminate user sessions. These are used to allow servers that maintain state information to free resources.

When a service only makes use of the Accounting portion of the Diameter protocol, even in combination with an application, the Session-Id is still used to identify user sessions. However, the session termination messages are not used, since a session is signaled as being terminated by issuing an accounting stop message.

#### 8.1. Authorization Session State Machine

This section contains a set of finite state machines, representing the life cycle of Diameter sessions, and which MUST be observed by all Diameter implementations that make use of the authentication and/or authorization portion of a Diameter application. The term Service-Specific below refers to a message defined in a Diameter application (e.g., Mobile IPv4, NASREQ).

There are four different authorization session state machines supported in the Diameter base protocol. The first two describe a session in which the server is maintaining session state, indicated by the value of the Auth-Session-State AVP (or its absence). One describes the session from a client perspective, the other from a server perspective. The second two state machines are used when the server does not maintain session state. Here again, one describes the session from a client perspective, the other from a server perspective.

When a session is moved to the Idle state, any resources that were allocated for the particular session must be released. Any event not listed in the state machines MUST be considered as an error condition, and an answer, if applicable, MUST be returned to the originator of the message.

In the state table, the event 'Failure to send X' means that the Diameter agent is unable to send command X to the desired destination. This could be due to the peer being down, or due to the peer sending back a transient failure or temporary protocol error notification DIAMETER\_TOO\_BUSY or DIAMETER\_LOOP\_DETECTED in the Result-Code AVP of the corresponding Answer command. The event 'X successfully sent' is the complement of 'Failure to send X'.

The following state machine is observed by a client when state is maintained on the server:

CLIENT, STATEFUL			
State	Event	Action	New State
Idle	Client or Device Requests access	Send service specific auth req	Pending
Idle	ASR Received for unknown session	Send ASA with Result-Code = UNKNOWN_SESSION_ID	Idle
Pending	Successful Service-specific authorization answer received with default Auth-Session-State value	Grant Access	Open
Pending	Successful Service-specific authorization answer received but service not provided	Sent STR	Discon
Pending	Error processing successful Service-specific authorization answer	Sent STR	Discon
Pending	Failed Service-specific authorization answer received	Cleanup	Idle
Open	User or client device requests access to service	Send service specific auth req	Open
Open	Successful Service-specific authorization answer received	Provide Service	Open
Open	Failed Service-specific authorization answer received.	Discon. user/device	Idle
Open	Session-Timeout Expires on Access Device	Send STR	Discon

Open	ASR Received, client will comply with request to end the session	Send ASA with Result-Code = SUCCESS, Send STR.	Discon
Open	ASR Received, client will not comply with request to end the session	Send ASA with Result-Code != SUCCESS	Open
Open	Authorization-Lifetime + Auth-Grace-Period expires on access device	Send STR	Discon
Discon	ASR Received	Send ASA	Discon
Discon	STA Received	Discon. user/device	Idle

The following state machine is observed by a server when it is maintaining state for the session:

SERVER, STATEFUL			
State	Event	Action	New State
Idle	Service-specific authorization request received, and user is authorized	Send successful serv. specific answer	Open
Idle	Service-specific authorization request received, and user is not authorized	Send failed serv. specific answer	Idle
Open	Service-specific authorization request received, and user is authorized	Send successful serv. specific answer	Open
Open	Service-specific authorization request received, and user is not authorized	Send failed serv. specific answer, Cleanup	Idle
Open	Home server wants to terminate the service	Send ASR	Discon

Open	Authorization-Lifetime (and Auth-Grace-Period) expires on home server.	Cleanup	Idle
Open	Session-Timeout expires on home server	Cleanup	Idle
Discon	Failure to send ASR	Wait, resend ASR	Discon
Discon	ASR successfully sent and ASA Received with Result-Code	Cleanup	Idle
Not Discon	ASA Received	None	No Change.
Any	STR Received	Send STA, Cleanup.	Idle

The following state machine is observed by a client when state is not maintained on the server:

CLIENT, STATELESS			
State	Event	Action	New State
Idle	Client or Device Requests access	Send service specific auth req	Pending
Pending	Successful Service-specific authorization answer received with Auth-Session-State set to NO_STATE_MAINTAINED	Grant Access	Open
Pending	Failed Service-specific authorization answer received	Cleanup	Idle
Open	Session-Timeout Expires on Access Device	Discon. user/device	Idle
Open	Service to user is terminated	Discon. user/device	Idle

The following state machine is observed by a server when it is not maintaining state for the session:

SERVER, STATELESS			
State	Event	Action	New State
-----			
Idle	Service-specific authorization request received, and successfully processed	Send serv. specific answer	Idle

## 8.2. Accounting Session State Machine

The following state machines **MUST** be supported for applications that have an accounting portion or that require only accounting services. The first state machine is to be observed by clients.

See Section 9.7 for Accounting Command Codes and Section 9.8 for Accounting AVPs.

The server side in the accounting state machine depends in some cases on the particular application. The Diameter base protocol defines a default state machine that **MUST** be followed by all applications that have not specified other state machines. This is the second state machine in this section described below.

The default server side state machine requires the reception of accounting records in any order and at any time, and does not place any standards requirement on the processing of these records. Implementations of Diameter **MAY** perform checking, ordering, correlation, fraud detection, and other tasks based on these records. Both base Diameter AVPs as well as application specific AVPs **MAY** be inspected as a part of these tasks. The tasks can happen either immediately after record reception or in a post-processing phase. However, as these tasks are typically application or even policy dependent, they are not standardized by the Diameter specifications. Applications **MAY** define requirements on when to accept accounting records based on the used value of Accounting-Realtime-Required AVP, credit limits checks, and so on.

However, the Diameter base protocol defines one optional server side state machine that **MAY** be followed by applications that require keeping track of the session state at the accounting server. Note that such tracking is incompatible with the ability to sustain long duration connectivity problems. Therefore, the use of this state machine is recommended only in applications where the value of the Accounting-Realtime-Required AVP is `DELIVER_AND_GRANT`, and hence accounting connectivity problems are required to cause the serviced user to be disconnected. Otherwise, records produced by the client



may be lost by the server which no longer accepts them after the connectivity is re-established. This state machine is the third state machine in this section. The state machine is supervised by a supervision session timer *Ts*, which the value should be reasonably higher than the *Acct\_Interim\_Interval* value. *Ts* MAY be set to two times the value of the *Acct\_Interim\_Interval* so as to avoid the accounting session in the Diameter server to change to Idle state in case of short transient network failure.

Any event not listed in the state machines MUST be considered as an error condition, and a corresponding answer, if applicable, MUST be returned to the originator of the message.

In the state table, the event 'Failure to send' means that the Diameter client is unable to communicate with the desired destination. This could be due to the peer being down, or due to the peer sending back a transient failure or temporary protocol error notification *DIAMETER\_OUT\_OF\_SPACE*, *DIAMETER\_TOO\_BUSY*, or *DIAMETER\_LOOP\_DETECTED* in the Result-Code AVP of the Accounting Answer command.

The event 'Failed answer' means that the Diameter client received a non-transient failure notification in the Accounting Answer command.

Note that the action 'Disconnect user/dev' MUST have an effect also to the authorization session state table, e.g., cause the STR message to be sent, if the given application has both authentication/authorization and accounting portions.

The states *PendingS*, *PendingI*, *PendingL*, *PendingE* and *PendingB* stand for pending states to wait for an answer to an accounting request related to a Start, Interim, Stop, Event or buffered record, respectively.

CLIENT, ACCOUNTING				
State	Event	Action	New State	
Idle	Client or device requests access	Send accounting start req.	<i>PendingS</i>	
Idle	Client or device requests a one-time service	Send accounting event req	<i>PendingE</i>	
Idle	Records in storage	Send record	<i>PendingB</i>	

PendingS	Successful accounting start answer received		Open
PendingS	Failure to send and buffer space available and realtime not equal to DELIVER_AND_GRANT	Store Start Record	Open
PendingS	Failure to send and no buffer space available and realtime equal to GRANT_AND_LOSE		Open
PendingS	Failure to send and no buffer space available and realtime not equal to GRANT_AND_LOSE	Disconnect user/dev	Idle
PendingS	Failed accounting start answer received and realtime equal to GRANT_AND_LOSE		Open
PendingS	Failed accounting start answer received and realtime not equal to GRANT_AND_LOSE	Disconnect user/dev	Idle
PendingS	User service terminated	Store stop record	PendingS
Open	Interim interval elapses	Send accounting interim record	PendingI
Open	User service terminated	Send accounting stop req.	PendingL
PendingI	Successful accounting interim answer received		Open
PendingI	Failure to send and (buffer space available or old record can be overwritten) and realtime not equal to DELIVER_AND_GRANT	Store interim record	Open
PendingI	Failure to send and no buffer space available and realtime equal to GRANT_AND_LOSE		Open

PendingI	Failure to send and no buffer space available and realtime not equal to GRANT_AND_LOSE	Disconnect user/dev	Idle
PendingI	Failed accounting interim answer received and realtime equal to GRANT_AND_LOSE		Open
PendingI	Failed accounting interim answer received and realtime not equal to GRANT_AND_LOSE	Disconnect user/dev	Idle
PendingI	User service terminated	Store stop record	PendingI
PendingE	Successful accounting event answer received		Idle
PendingE	Failure to send and buffer space available	Store event record	Idle
PendingE	Failure to send and no buffer space available		Idle
PendingE	Failed accounting event answer received		Idle
PendingB	Successful accounting answer received	Delete record	Idle
PendingB	Failure to send		Idle
PendingB	Failed accounting answer received	Delete record	Idle
PendingL	Successful accounting stop answer received		Idle
PendingL	Failure to send and buffer space available	Store stop record	Idle
PendingL	Failure to send and no buffer space available		Idle
PendingL	Failed accounting stop answer received		Idle

SERVER, STATELESS ACCOUNTING			
State	Event	Action	New State
-----			
Idle	Accounting start request received, and successfully processed.	Send accounting start answer	Idle
Idle	Accounting event request received, and successfully processed.	Send accounting event answer	Idle
Idle	Interim record received, and successfully processed.	Send accounting interim answer	Idle
Idle	Accounting stop request received, and successfully processed	Send accounting stop answer	Idle
Idle	Accounting request received, no space left to store records	Send accounting answer, Result-Code = OUT_OF_SPACE	Idle

SERVER, STATEFUL ACCOUNTING			
State	Event	Action	New State
-----			
Idle	Accounting start request received, and successfully processed.	Send accounting start answer, Start Ts	Open
Idle	Accounting event request received, and successfully processed.	Send accounting event answer	Idle

Idle	Accounting request received, no space left to store records	Send accounting answer, Result-Code = OUT_OF_ SPACE	Idle
Open	Interim record received, and successfully processed.	Send accounting interim answer, Restart Ts	Open
Open	Accounting stop request received, and successfully processed	Send accounting stop answer, Stop Ts	Idle
Open	Accounting request received, no space left to store records	Send accounting answer, Result-Code = OUT_OF_ SPACE, Stop Ts	Idle
Open	Session supervision timer Ts expired	Stop Ts	Idle

### 8.3. Server-Initiated Re-Auth

A Diameter server may initiate a re-authentication and/or re-authorization service for a particular session by issuing a Re-Auth-Request (RAR).

For example, for pre-paid services, the Diameter server that originally authorized a session may need some confirmation that the user is still using the services.

An access device that receives a RAR message with Session-Id equal to a currently active session MUST initiate a re-auth towards the user, if the service supports this particular feature. Each Diameter application MUST state whether service-initiated re-auth is supported, since some applications do not allow access devices to prompt the user for re-auth.

### 8.3.1. Re-Auth-Request

The Re-Auth-Request (RAR), indicated by the Command-Code set to 258 and the message flags' 'R' bit set, may be sent by any server to the access device that is providing session service, to request that the user be re-authenticated and/or re-authorized.

#### Message Format

```
<RAR> ::= < Diameter Header: 258, REQ, PXY >
        < Session-Id >
        { Origin-Host }
        { Origin-Realm }
        { Destination-Realm }
        { Destination-Host }
        { Auth-Application-Id }
        { Re-Auth-Request-Type }
        [ User-Name ]
        [ Origin-State-Id ]
        * [ Proxy-Info ]
        * [ Route-Record ]
        * [ AVP ]
```

### 8.3.2. Re-Auth-Answer

The Re-Auth-Answer (RAA), indicated by the Command-Code set to 258 and the message flags' 'R' bit clear, is sent in response to the RAR. The Result-Code AVP MUST be present, and indicates the disposition of the request.

A successful RAA message MUST be followed by an application-specific authentication and/or authorization message.

## Message Format

```
<RAA> ::= < Diameter Header: 258, PXY >
        < Session-Id >
        { Result-Code }
        { Origin-Host }
        { Origin-Realm }
        [ User-Name ]
        [ Origin-State-Id ]
        [ Error-Message ]
        [ Error-Reporting-Host ]
        * [ Failed-AVP ]
        * [ Redirect-Host ]
        [ Redirect-Host-Usage ]
        [ Redirect-Host-Cache-Time ]
        * [ Proxy-Info ]
        * [ AVP ]
```

## 8.4. Session Termination

It is necessary for a Diameter server that authorized a session, for which it is maintaining state, to be notified when that session is no longer active, both for tracking purposes as well as to allow stateful agents to release any resources that they may have provided for the user's session. For sessions whose state is not being maintained, this section is not used.

When a user session that required Diameter authorization terminates, the access device that provided the service MUST issue a Session-Termination-Request (STR) message to the Diameter server that authorized the service, to notify it that the session is no longer active. An STR MUST be issued when a user session terminates for any reason, including user logoff, expiration of Session-Timeout, administrative action, termination upon receipt of an Abort-Session-Request (see below), orderly shutdown of the access device, etc.

The access device also MUST issue an STR for a session that was authorized but never actually started. This could occur, for example, due to a sudden resource shortage in the access device, or because the access device is unwilling to provide the type of service requested in the authorization, or because the access device does not support a mandatory AVP returned in the authorization, etc.

It is also possible that a session that was authorized is never actually started due to action of a proxy. For example, a proxy may modify an authorization answer, converting the result from success to failure, prior to forwarding the message to the access device. If the answer did not contain an Auth-Session-State AVP with the value

NO\_STATE\_MAINTAINED, a proxy that causes an authorized session not to be started MUST issue an STR to the Diameter server that authorized the session, since the access device has no way of knowing that the session had been authorized.

A Diameter server that receives an STR message MUST clean up resources (e.g., session state) associated with the Session-Id specified in the STR, and return a Session-Termination-Answer.

A Diameter server also MUST clean up resources when the Session-Timeout expires, or when the Authorization-Lifetime and the Auth-Grace-Period AVPs expires without receipt of a re-authorization request, regardless of whether an STR for that session is received. The access device is not expected to provide service beyond the expiration of these timers; thus, expiration of either of these timers implies that the access device may have unexpectedly shut down.

#### 8.4.1. Session-Termination-Request

The Session-Termination-Request (STR), indicated by the Command-Code set to 275 and the Command Flags' 'R' bit set, is sent by the access device to inform the Diameter Server that an authenticated and/or authorized session is being terminated.

##### Message Format

```
<STR> ::= < Diameter Header: 275, REQ, PXY >
          < Session-Id >
          { Origin-Host }
          { Origin-Realm }
          { Destination-Realm }
          { Auth-Application-Id }
          { Termination-Cause }
          [ User-Name ]
          [ Destination-Host ]
          * [ Class ]
          [ Origin-State-Id ]
          * [ Proxy-Info ]
          * [ Route-Record ]
          * [ AVP ]
```



#### 8.4.2. Session-Termination-Answer

The Session-Termination-Answer (STA), indicated by the Command-Code set to 275 and the message flags' 'R' bit clear, is sent by the Diameter Server to acknowledge the notification that the session has been terminated. The Result-Code AVP MUST be present, and MAY contain an indication that an error occurred while servicing the STR.

Upon sending or receipt of the STA, the Diameter Server MUST release all resources for the session indicated by the Session-Id AVP. Any intermediate server in the Proxy-Chain MAY also release any resources, if necessary.

##### Message Format

```
<STA> ::= < Diameter Header: 275, PXY >
        < Session-Id >
        { Result-Code }
        { Origin-Host }
        { Origin-Realm }
        [ User-Name ]
    * [ Class ]
      [ Error-Message ]
      [ Error-Reporting-Host ]
    * [ Failed-AVP ]
      [ Origin-State-Id ]
    * [ Redirect-Host ]
      [ Redirect-Host-Usage ]
        ^
      [ Redirect-Max-Cache-Time ]
    * [ Proxy-Info ]
    * [ AVP ]
```

#### 8.5. Aborting a Session

A Diameter server may request that the access device stop providing service for a particular session by issuing an Abort-Session-Request (ASR).

For example, the Diameter server that originally authorized the session may be required to cause that session to be stopped for credit or other reasons that were not anticipated when the session was first authorized. On the other hand, an operator may maintain a management server for the purpose of issuing ASRs to administratively remove users from the network.

An access device that receives an ASR with Session-ID equal to a currently active session MAY stop the session. Whether the access

device stops the session or not is implementation- and/or configuration-dependent. For example, an access device may honor ASRs from certain agents only. In any case, the access device MUST respond with an Abort-Session-Answer, including a Result-Code AVP to indicate what action it took.

Note that if the access device does stop the session upon receipt of an ASR, it issues an STR to the authorizing server (which may or may not be the agent issuing the ASR) just as it would if the session were terminated for any other reason.

#### 8.5.1. Abort-Session-Request

The Abort-Session-Request (ASR), indicated by the Command-Code set to 274 and the message flags' 'R' bit set, may be sent by any server to the access device that is providing session service, to request that the session identified by the Session-Id be stopped.

##### Message Format

```
<ASR> ::= < Diameter Header: 274, REQ, PXY >
          < Session-Id >
          { Origin-Host }
          { Origin-Realm }
          { Destination-Realm }
          { Destination-Host }
          { Auth-Application-Id }
          [ User-Name ]
          [ Origin-State-Id ]
          * [ Proxy-Info ]
          * [ Route-Record ]
          * [ AVP ]
```

#### 8.5.2. Abort-Session-Answer

The Abort-Session-Answer (ASA), indicated by the Command-Code set to 274 and the message flags' 'R' bit clear, is sent in response to the ASR. The Result-Code AVP MUST be present, and indicates the disposition of the request.

If the session identified by Session-Id in the ASR was successfully terminated, Result-Code is set to DIAMETER\_SUCCESS. If the session is not currently active, Result-Code is set to DIAMETER\_UNKNOWN\_SESSION\_ID. If the access device does not stop the session for any other reason, Result-Code is set to DIAMETER\_UNABLE\_TO\_COMPLY.

## Message Format

```
<ASA> ::= < Diameter Header: 274, PXY >
        < Session-Id >
        { Result-Code }
        { Origin-Host }
        { Origin-Realm }
        [ User-Name ]
        [ Origin-State-Id ]
        [ Error-Message ]
        [ Error-Reporting-Host ]
        * [ Failed-AVP ]
        * [ Redirect-Host ]
          [ Redirect-Host-Usage ]
          [ Redirect-Max-Cache-Time ]
        * [ Proxy-Info ]
        * [ AVP ]
```

## 8.6. Inferring Session Termination from Origin-State-Id

Origin-State-Id is used to allow rapid detection of terminated sessions for which no STR would have been issued, due to unanticipated shutdown of an access device.

By including Origin-State-Id in CER/CEA messages, an access device allows a next-hop server to determine immediately upon connection whether the device has lost its sessions since the last connection.

By including Origin-State-Id in request messages, an access device also allows a server with which it communicates via proxy to make such a determination. However, a server that is not directly connected with the access device will not discover that the access device has been restarted unless and until it receives a new request from the access device. Thus, use of this mechanism across proxies is opportunistic rather than reliable, but useful nonetheless.

When a Diameter server receives an Origin-State-Id that is greater than the Origin-State-Id previously received from the same issuer, it may assume that the issuer has lost state since the previous message and that all sessions that were active under the lower Origin-State-Id have been terminated. The Diameter server MAY clean up all session state associated with such lost sessions, and MAY also issues STRs for all such lost sessions that were authorized on upstream servers, to allow session state to be cleaned up globally.

### 8.7. Auth-Request-Type AVP

The Auth-Request-Type AVP (AVP Code 274) is of type Enumerated and is included in application-specific auth requests to inform the peers whether a user is to be authenticated only, authorized only or both. Note any value other than both MAY cause RADIUS interoperability issues. The following values are defined:

AUTHENTICATE\_ONLY                    1

The request being sent is for authentication only, and MUST contain the relevant application specific authentication AVPs that are needed by the Diameter server to authenticate the user.

AUTHORIZE\_ONLY                        2

The request being sent is for authorization only, and MUST contain the application specific authorization AVPs that are necessary to identify the service being requested/offered.

AUTHORIZE\_AUTHENTICATE               3

The request contains a request for both authentication and authorization. The request MUST include both the relevant application specific authentication information, and authorization information necessary to identify the service being requested/offered.

### 8.8. Session-Id AVP

The Session-Id AVP (AVP Code 263) is of type UTF8String and is used to identify a specific session (see Section 8). All messages pertaining to a specific session MUST include only one Session-Id AVP and the same value MUST be used throughout the life of a session. When present, the Session-Id SHOULD appear immediately following the Diameter Header (see Section 3).

The Session-Id MUST be globally and eternally unique, as it is meant to uniquely identify a user session without reference to any other information, and may be needed to correlate historical authentication information with accounting information. The Session-Id includes a mandatory portion and an implementation-defined portion; a recommended format for the implementation-defined portion is outlined below.

The Session-Id MUST begin with the sender's identity encoded in the DiameterIdentity type (see Section 4.4). The remainder of the Session-Id is delimited by a ";" character, and MAY be any sequence that the client can guarantee to be eternally unique; however, the following format is recommended, (square brackets [] indicate an optional element):

<DiameterIdentity>;<high 32 bits>;<low 32 bits>[;<optional value>]

<high 32 bits> and <low 32 bits> are decimal representations of the high and low 32 bits of a monotonically increasing 64-bit value. The 64-bit value is rendered in two part to simplify formatting by 32-bit processors. At startup, the high 32 bits of the 64-bit value MAY be initialized to the time, and the low 32 bits MAY be initialized to zero. This will for practical purposes eliminate the possibility of overlapping Session-Ids after a reboot, assuming the reboot process takes longer than a second. Alternatively, an implementation MAY keep track of the increasing value in non-volatile memory.

<optional value> is implementation specific but may include a modem's device Id, a layer 2 address, timestamp, etc.

Example, in which there is no optional value:

accesspoint7.acme.com;1876543210;523

Example, in which there is an optional value:

accesspoint7.acme.com;1876543210;523;mobile@200.1.1.88

The Session-Id is created by the Diameter application initiating the session, which in most cases is done by the client. Note that a Session-Id MAY be used for both the authorization and accounting commands of a given application.

#### 8.9. Authorization-Lifetime AVP

The Authorization-Lifetime AVP (AVP Code 291) is of type Unsigned32 and contains the maximum number of seconds of service to be provided to the user before the user is to be re-authenticated and/or re-authorized. Great care should be taken when the Authorization-Lifetime value is determined, since a low, non-zero, value could create significant Diameter traffic, which could congest both the network and the agents.

A value of zero (0) means that immediate re-auth is necessary by the access device. This is typically used in cases where multiple authentication methods are used, and a successful auth response with this AVP set to zero is used to signal that the next authentication method is to be immediately initiated. The absence of this AVP, or a value of all ones (meaning all bits in the 32 bit field are set to one) means no re-auth is expected.

If both this AVP and the Session-Timeout AVP are present in a message, the value of the latter MUST NOT be smaller than the Authorization-Lifetime AVP.

An Authorization-Lifetime AVP MAY be present in re-authorization messages, and contains the number of seconds the user is authorized to receive service from the time the re-auth answer message is received by the access device.

This AVP MAY be provided by the client as a hint of the maximum lifetime that it is willing to accept. However, the server MAY return a value that is equal to, or smaller, than the one provided by the client.

#### 8.10. Auth-Grace-Period AVP

The Auth-Grace-Period AVP (AVP Code 276) is of type Unsigned32 and contains the number of seconds the Diameter server will wait following the expiration of the Authorization-Lifetime AVP before cleaning up resources for the session.

#### 8.11. Auth-Session-State AVP

The Auth-Session-State AVP (AVP Code 277) is of type Enumerated and specifies whether state is maintained for a particular session. The client MAY include this AVP in requests as a hint to the server, but the value in the server's answer message is binding. The following values are supported:

STATE\_MAINTAINED                      0  
This value is used to specify that session state is being maintained, and the access device MUST issue a session termination message when service to the user is terminated. This is the default value.

NO\_STATE\_MAINTAINED                    1  
This value is used to specify that no session termination messages will be sent by the access device upon expiration of the Authorization-Lifetime.

#### 8.12. Re-Auth-Request-Type AVP

The Re-Auth-Request-Type AVP (AVP Code 285) is of type Enumerated and is included in application-specific auth answers to inform the client of the action expected upon expiration of the Authorization-Lifetime. If the answer message contains an Authorization-Lifetime AVP with a positive value, the Re-Auth-Request-Type AVP MUST be present in an answer message. The following values are defined:

**AUTHORIZE\_ONLY** 0

An authorization only re-auth is expected upon expiration of the Authorization-Lifetime. This is the default value if the AVP is not present in answer messages that include the Authorization-Lifetime.

**AUTHORIZE\_AUTHENTICATE** 1

An authentication and authorization re-auth is expected upon expiration of the Authorization-Lifetime.

**8.13. Session-Timeout AVP**

The Session-Timeout AVP (AVP Code 27) [RADIUS] is of type Unsigned32 and contains the maximum number of seconds of service to be provided to the user before termination of the session. When both the Session-Timeout and the Authorization-Lifetime AVPs are present in an answer message, the former MUST be equal to or greater than the value of the latter.

A session that terminates on an access device due to the expiration of the Session-Timeout MUST cause an STR to be issued, unless both the access device and the home server had previously agreed that no session termination messages would be sent (see Section 8.9).

A Session-Timeout AVP MAY be present in a re-authorization answer message, and contains the remaining number of seconds from the beginning of the re-auth.

A value of zero, or the absence of this AVP, means that this session has an unlimited number of seconds before termination.

This AVP MAY be provided by the client as a hint of the maximum timeout that it is willing to accept. However, the server MAY return a value that is equal to, or smaller, than the one provided by the client.

**8.14. User-Name AVP**

The User-Name AVP (AVP Code 1) [RADIUS] is of type UTF8String, which contains the User-Name, in a format consistent with the NAI specification [NAI].

**8.15. Termination-Cause AVP**

The Termination-Cause AVP (AVP Code 295) is of type Enumerated, and is used to indicate the reason why a session was terminated on the access device. The following values are defined:

- DIAMETER\_LOGOUT 1  
The user initiated a disconnect
- DIAMETER\_SERVICE\_NOT\_PROVIDED 2  
This value is used when the user disconnected prior to the receipt of the authorization answer message.
- DIAMETER\_BAD\_ANSWER 3  
This value indicates that the authorization answer received by the access device was not processed successfully.
- DIAMETER\_ADMINISTRATIVE 4  
The user was not granted access, or was disconnected, due to administrative reasons, such as the receipt of a Abort-Session-Request message.
- DIAMETER\_LINK\_BROKEN 5  
The communication to the user was abruptly disconnected.
- DIAMETER\_AUTH\_EXPIRED 6  
The user's access was terminated since its authorized session time has expired.
- DIAMETER\_USER\_MOVED 7  
The user is receiving services from another access device.
- DIAMETER\_SESSION\_TIMEOUT 8  
The user's session has timed out, and service has been terminated.

#### 8.16. Origin-State-Id AVP

The Origin-State-Id AVP (AVP Code 278), of type Unsigned32, is a monotonically increasing value that is advanced whenever a Diameter entity restarts with loss of previous state, for example upon reboot. Origin-State-Id MAY be included in any Diameter message, including CER.

A Diameter entity issuing this AVP MUST create a higher value for this AVP each time its state is reset. A Diameter entity MAY set Origin-State-Id to the time of startup, or it MAY use an incrementing counter retained in non-volatile memory across restarts.

The Origin-State-Id, if present, MUST reflect the state of the entity indicated by Origin-Host. If a proxy modifies Origin-Host, it MUST either remove Origin-State-Id or modify it appropriately as well.



Typically, Origin-State-Id is used by an access device that always starts up with no active sessions; that is, any session active prior to restart will have been lost. By including Origin-State-Id in a message, it allows other Diameter entities to infer that sessions associated with a lower Origin-State-Id are no longer active. If an access device does not intend for such inferences to be made, it **MUST** either not include Origin-State-Id in any message, or set its value to 0.

#### 8.17. Session-Binding AVP

The Session-Binding AVP (AVP Code 270) is of type Unsigned32, and **MAY** be present in application-specific authorization answer messages. If present, this AVP **MAY** inform the Diameter client that all future application-specific re-auth messages for this session **MUST** be sent to the same authorization server. This AVP **MAY** also specify that a Session-Termination-Request message for this session **MUST** be sent to the same authorizing server.

This field is a bit mask, and the following bits have been defined:

##### RE\_AUTH 1

When set, future re-auth messages for this session **MUST NOT** include the Destination-Host AVP. When cleared, the default value, the Destination-Host AVP **MUST** be present in all re-auth messages for this session.

##### STR 2

When set, the STR message for this session **MUST NOT** include the Destination-Host AVP. When cleared, the default value, the Destination-Host AVP **MUST** be present in the STR message for this session.

##### ACCOUNTING 4

When set, all accounting messages for this session **MUST NOT** include the Destination-Host AVP. When cleared, the default value, the Destination-Host AVP, if known, **MUST** be present in all accounting messages for this session.

#### 8.18. Session-Server-Failover AVP

The Session-Server-Failover AVP (AVP Code 271) is of type Enumerated, and **MAY** be present in application-specific authorization answer messages that either do not include the Session-Binding AVP or include the Session-Binding AVP with any of the bits set to a zero value. If present, this AVP **MAY** inform the Diameter client that if a

re-auth or STR message fails due to a delivery problem, the Diameter client SHOULD issue a subsequent message without the Destination-Host AVP. When absent, the default value is REFUSE\_SERVICE.

The following values are supported:

REFUSE\_SERVICE                    0  
If either the re-auth or the STR message delivery fails, terminate service with the user, and do not attempt any subsequent attempts.

TRY\_AGAIN                         1  
If either the re-auth or the STR message delivery fails, resend the failed message without the Destination-Host AVP present.

ALLOW\_SERVICE                    2  
If re-auth message delivery fails, assume that re-authorization succeeded. If STR message delivery fails, terminate the session.

TRY\_AGAIN\_ALLOW\_SERVICE         3  
If either the re-auth or the STR message delivery fails, resend the failed message without the Destination-Host AVP present. If the second delivery fails for re-auth, assume re-authorization succeeded. If the second delivery fails for STR, terminate the session.

#### 8.19. Multi-Round-Time-Out AVP

The Multi-Round-Time-Out AVP (AVP Code 272) is of type Unsigned32, and SHOULD be present in application-specific authorization answer messages whose Result-Code AVP is set to DIAMETER\_MULTI\_ROUND\_AUTH. This AVP contains the maximum number of seconds that the access device MUST provide the user in responding to an authentication request.

#### 8.20. Class AVP

The Class AVP (AVP Code 25) is of type OctetString and is used to by Diameter servers to return state information to the access device. When one or more Class AVPs are present in application-specific authorization answer messages, they MUST be present in subsequent re-authorization, session termination and accounting messages. Class AVPs found in a re-authorization answer message override the ones found in any previous authorization answer message. Diameter server implementations SHOULD NOT return Class AVPs that require more than 4096 bytes of storage on the Diameter client. A Diameter client that receives Class AVPs whose size exceeds local available storage MUST terminate the session.

#### 8.21. Event-Timestamp AVP

The Event-Timestamp (AVP Code 55) is of type Time, and MAY be included in an Accounting-Request and Accounting-Answer messages to record the time that the reported event occurred, in seconds since January 1, 1900 00:00 UTC.

### 9. Accounting

This accounting protocol is based on a server directed model with capabilities for real-time delivery of accounting information. Several fault resilience methods [ACCMGMT] have been built in to the protocol in order minimize loss of accounting data in various fault situations and under different assumptions about the capabilities of the used devices.

#### 9.1. Server Directed Model

The server directed model means that the device generating the accounting data gets information from either the authorization server (if contacted) or the accounting server regarding the way accounting data shall be forwarded. This information includes accounting record timeliness requirements.

As discussed in [ACCMGMT], real-time transfer of accounting records is a requirement, such as the need to perform credit limit checks and fraud detection. Note that batch accounting is not a requirement, and is therefore not supported by Diameter. Should batched accounting be required in the future, a new Diameter application will need to be created, or it could be handled using another protocol. Note, however, that even if at the Diameter layer accounting requests are processed one by one, transport protocols used under Diameter typically batch several requests in the same packet under heavy traffic conditions. This may be sufficient for many applications.

The authorization server (chain) directs the selection of proper transfer strategy, based on its knowledge of the user and relationships of roaming partnerships. The server (or agents) uses the Acct-Interim-Interval and Accounting-Realtime-Required AVPs to control the operation of the Diameter peer operating as a client. The Acct-Interim-Interval AVP, when present, instructs the Diameter node acting as a client to produce accounting records continuously even during a session. Accounting-Realtime-Required AVP is used to control the behavior of the client when the transfer of accounting records from the Diameter client is delayed or unsuccessful.

The Diameter accounting server MAY override the interim interval or the realtime requirements by including the Acct-Interim-Interval or Accounting-Realtime-Required AVP in the Accounting-Answer message. When one of these AVPs is present, the latest value received SHOULD be used in further accounting activities for the same session.

## 9.2. Protocol Messages

A Diameter node that receives a successful authentication and/or authorization messages from the Home AAA server MUST collect accounting information for the session. The Accounting-Request message is used to transmit the accounting information to the Home AAA server, which MUST reply with the Accounting-Answer message to confirm reception. The Accounting-Answer message includes the Result-Code AVP, which MAY indicate that an error was present in the accounting message. A rejected Accounting-Request message MAY cause the user's session to be terminated, depending on the value of the Accounting-Realtime-Required AVP received earlier for the session in question.

Each Diameter Accounting protocol message MAY be compressed, in order to reduce network bandwidth usage. If IPsec and IKE are used to secure the Diameter session, then IP compression [IPComp] MAY be used and IKE [IKE] MAY be used to negotiate the compression parameters. If TLS is used to secure the Diameter session, then TLS compression [TLS] MAY be used.

## 9.3. Application document requirements

Each Diameter application (e.g., NASREQ, MobileIP), MUST define their Service-Specific AVPs that MUST be present in the Accounting-Request message in a section entitled "Accounting AVPs". The application MUST assume that the AVPs described in this document will be present in all Accounting messages, so only their respective service-specific AVPs need to be defined in this section.

## 9.4. Fault Resilience

Diameter Base protocol mechanisms are used to overcome small message loss and network faults of temporary nature.

Diameter peers acting as clients MUST implement the use of failover to guard against server failures and certain network failures. Diameter peers acting as agents or related off-line processing systems MUST detect duplicate accounting records caused by the sending of same record to several servers and duplication of messages

in transit. This detection **MUST** be based on the inspection of the Session-Id and Accounting-Record-Number AVP pairs. Appendix C discusses duplicate detection needs and implementation issues.

Diameter clients **MAY** have non-volatile memory for the safe storage of accounting records over reboots or extended network failures, network partitions, and server failures. If such memory is available, the client **SHOULD** store new accounting records there as soon as the records are created and until a positive acknowledgement of their reception from the Diameter Server has been received. Upon a reboot, the client **MUST** start sending the records in the non-volatile memory to the accounting server with appropriate modifications in termination cause, session length, and other relevant information in the records.

A further application of this protocol may include AVPs to control how many accounting records may at most be stored in the Diameter client without committing them to the non-volatile memory or transferring them to the Diameter server.

The client **SHOULD NOT** remove the accounting data from any of its memory areas before the correct Accounting-Answer has been received. The client **MAY** remove oldest, undelivered or yet unacknowledged accounting data if it runs out of resources such as memory. It is an implementation dependent matter for the client to accept new sessions under this condition.

#### 9.5. Accounting Records

In all accounting records, the Session-Id AVP **MUST** be present; the User-Name AVP **MUST** be present if it is available to the Diameter client. If strong authentication across agents is required, end-to-end security may be used for authentication purposes.

Different types of accounting records are sent depending on the actual type of accounted service and the authorization server's directions for interim accounting. If the accounted service is a one-time event, meaning that the start and stop of the event are simultaneous, then the Accounting-Record-Type AVP **MUST** be present and set to the value `EVENT_RECORD`.

If the accounted service is of a measurable length, then the AVP **MUST** use the values `START_RECORD`, `STOP_RECORD`, and possibly, `INTERIM_RECORD`. If the authorization server has not directed interim accounting to be enabled for the session, two accounting records **MUST** be generated for each service of type session. When the initial

Accounting-Request for a given session is sent, the Accounting-Record-Type AVP MUST be set to the value START\_RECORD. When the last Accounting-Request is sent, the value MUST be STOP\_RECORD.

If the authorization server has directed interim accounting to be enabled, the Diameter client MUST produce additional records between the START\_RECORD and STOP\_RECORD, marked INTERIM\_RECORD. The production of these records is directed by Acct-Interim-Interval as well as any re-authentication or re-authorization of the session. The Diameter client MUST overwrite any previous interim accounting records that are locally stored for delivery, if a new record is being generated for the same session. This ensures that only one pending interim record can exist on an access device for any given session.

A particular value of Accounting-Sub-Session-Id MUST appear only in one sequence of accounting records from a DIAMETER client, except for the purposes of retransmission. The one sequence that is sent MUST be either one record with Accounting-Record-Type AVP set to the value EVENT\_RECORD, or several records starting with one having the value START\_RECORD, followed by zero or more INTERIM\_RECORD and a single STOP\_RECORD. A particular Diameter application specification MUST define the type of sequences that MUST be used.

#### 9.6. Correlation of Accounting Records

The Diameter protocol's Session-Id AVP, which is globally unique (see Section 8.8), is used during the authorization phase to identify a particular session. Services that do not require any authorization still use the Session-Id AVP to identify sessions. Accounting messages MAY use a different Session-Id from that sent in authorization messages. Specific applications MAY require different a Session-ID for accounting messages.

However, there are certain applications that require multiple accounting sub-sessions. Such applications would send messages with a constant Session-Id AVP, but a different Accounting-Sub-Session-Id AVP. In these cases, correlation is performed using the Session-Id. It is important to note that receiving a STOP\_RECORD with no Accounting-Sub-Session-Id AVP when sub-sessions were originally used in the START\_RECORD messages implies that all sub-sessions are terminated.

Furthermore, there are certain applications where a user receives service from different access devices (e.g., Mobile IPv4), each with their own unique Session-Id. In such cases, the Acct-Multi-Session-Id AVP is used for correlation. During authorization, a server that

determines that a request is for an existing session SHOULD include the Acct-Multi-Session-Id AVP, which the access device MUST include in all subsequent accounting messages.

The Acct-Multi-Session-Id AVP MAY include the value of the original Session-Id. It's contents are implementation specific, but MUST be globally unique across other Acct-Multi-Session-Id, and MUST NOT change during the life of a session.

A Diameter application document MUST define the exact concept of a session that is being accounted, and MAY define the concept of a multi-session. For instance, the NASREQ DIAMETER application treats a single PPP connection to a Network Access Server as one session, and a set of Multilink PPP sessions as one multi-session.

### 9.7. Accounting Command-Codes

This section defines Command-Code values that MUST be supported by all Diameter implementations that provide Accounting services.

#### 9.7.1. Accounting-Request

The Accounting-Request (ACR) command, indicated by the Command-Code field set to 271 and the Command Flags' 'R' bit set, is sent by a Diameter node, acting as a client, in order to exchange accounting information with a peer.

One of Acct-Application-Id and Vendor-Specific-Application-Id AVPs MUST be present. If the Vendor-Specific-Application-Id grouped AVP is present, it must have an Acct-Application-Id inside.

The AVP listed below SHOULD include service specific accounting AVPs, as described in Section 9.3.

## Message Format

```
<ACR> ::= < Diameter Header: 271, REQ, PXY >
      < Session-Id >
      { Origin-Host }
      { Origin-Realm }
      { Destination-Realm }
      { Accounting-Record-Type }
      { Accounting-Record-Number }
      [ Acct-Application-Id ]
      [ Vendor-Specific-Application-Id ]
      [ User-Name ]
      [ Accounting-Sub-Session-Id ]
      [ Acct-Session-Id ]
      [ Acct-Multi-Session-Id ]
      [ Acct-Interim-Interval ]
      [ Accounting-Realtime-Required ]
      [ Origin-State-Id ]
      [ Event-Timestamp ]
      * [ Proxy-Info ]
      * [ Route-Record ]
      * [ AVP ]
```

## 9.7.2. Accounting-Answer

The Accounting-Answer (ACA) command, indicated by the Command-Code field set to 271 and the Command Flags' 'R' bit cleared, is used to acknowledge an Accounting-Request command. The Accounting-Answer command contains the same Session-Id and includes the usage AVPs only if CMS is in use when sending this command. Note that the inclusion of the usage AVPs when CMS is not being used leads to unnecessarily large answer messages, and can not be used as a server's proof of the receipt of these AVPs in an end-to-end fashion. If the Accounting-Request was protected by end-to-end security, then the corresponding ACA message MUST be protected by end-to-end security.

Only the target Diameter Server, known as the home Diameter Server, SHOULD respond with the Accounting-Answer command.

One of Acct-Application-Id and Vendor-Specific-Application-Id AVPs MUST be present. If the Vendor-Specific-Application-Id grouped AVP is present, it must have an Acct-Application-Id inside.

The AVP listed below SHOULD include service specific accounting AVPs, as described in Section 9.3.



## Message Format

```

<ACA> ::= < Diameter Header: 271, PXY >
        < Session-Id >
        { Result-Code }
        { Origin-Host }
        { Origin-Realm }
        { Accounting-Record-Type }
        { Accounting-Record-Number }
        [ Acct-Application-Id ]
        [ Vendor-Specific-Application-Id ]
        [ User-Name ]
        [ Accounting-Sub-Session-Id ]
        [ Acct-Session-Id ]
        [ Acct-Multi-Session-Id ]
        [ Error-Reporting-Host ]
        [ Acct-Interim-Interval ]
        [ Accounting-Realtime-Required ]
        [ Origin-State-Id ]
        [ Event-Timestamp ]
        * [ Proxy-Info ]
        * [ AVP ]

```

## 9.8. Accounting AVPs

This section contains AVPs that describe accounting usage information related to a specific session.

## 9.8.1. Accounting-Record-Type AVP

The Accounting-Record-Type AVP (AVP Code 480) is of type Enumerated and contains the type of accounting record being sent. The following values are currently defined for the Accounting-Record-Type AVP:

## EVENT\_RECORD 1

An Accounting Event Record is used to indicate that a one-time event has occurred (meaning that the start and end of the event are simultaneous). This record contains all information relevant to the service, and is the only record of the service.

## START\_RECORD 2

An Accounting Start, Interim, and Stop Records are used to indicate that a service of a measurable length has been given. An Accounting Start Record is used to initiate an accounting session, and contains accounting information that is relevant to the initiation of the session.

**INTERIM\_RECORD**

3

An Interim Accounting Record contains cumulative accounting information for an existing accounting session. Interim Accounting Records SHOULD be sent every time a re-authentication or re-authorization occurs. Further, additional interim record triggers MAY be defined by application-specific Diameter applications. The selection of whether to use INTERIM\_RECORD records is done by the Acct-Interim-Interval AVP.

**STOP\_RECORD**

4

An Accounting Stop Record is sent to terminate an accounting session and contains cumulative accounting information relevant to the existing session.

**9.8.2. Acct-Interim-Interval**

The Acct-Interim-Interval AVP (AVP Code 85) is of type Unsigned32 and is sent from the Diameter home authorization server to the Diameter client. The client uses information in this AVP to decide how and when to produce accounting records. With different values in this AVP, service sessions can result in one, two, or two+N accounting records, based on the needs of the home-organization. The following accounting record production behavior is directed by the inclusion of this AVP:

1. The omission of the Acct-Interim-Interval AVP or its inclusion with Value field set to 0 means that EVENT\_RECORD, START\_RECORD, and STOP\_RECORD are produced, as appropriate for the service.
2. The inclusion of the AVP with Value field set to a non-zero value means that INTERIM\_RECORD records MUST be produced between the START\_RECORD and STOP\_RECORD records. The Value field of this AVP is the nominal interval between these records in seconds. The Diameter node that originates the accounting information, known as the client, MUST produce the first INTERIM\_RECORD record roughly at the time when this nominal interval has elapsed from the START\_RECORD, the next one again as the interval has elapsed once more, and so on until the session ends and a STOP\_RECORD record is produced.

The client MUST ensure that the interim record production times are randomized so that large accounting message storms are not created either among records or around a common service start time.

#### 9.8.3. Accounting-Record-Number AVP

The Accounting-Record-Number AVP (AVP Code 485) is of type Unsigned32 and identifies this record within one session. As Session-Id AVPs are globally unique, the combination of Session-Id and Accounting-Record-Number AVPs is also globally unique, and can be used in matching accounting records with confirmations. An easy way to produce unique numbers is to set the value to 0 for records of type EVENT\_RECORD and START\_RECORD, and set the value to 1 for the first INTERIM\_RECORD, 2 for the second, and so on until the value for STOP\_RECORD is one more than for the last INTERIM\_RECORD.

#### 9.8.4. Acct-Session-Id AVP

The Acct-Session-Id AVP (AVP Code 44) is of type OctetString is only used when RADIUS/Diameter translation occurs. This AVP contains the contents of the RADIUS Acct-Session-Id attribute.

#### 9.8.5. Acct-Multi-Session-Id AVP

The Acct-Multi-Session-Id AVP (AVP Code 50) is of type UTF8String, following the format specified in Section 8.8. The Acct-Multi-Session-Id AVP is used to link together multiple related accounting sessions, where each session would have a unique Session-Id, but the same Acct-Multi-Session-Id AVP. This AVP MAY be returned by the Diameter server in an authorization answer, and MUST be used in all accounting messages for the given session.

#### 9.8.6. Accounting-Sub-Session-Id AVP

The Accounting-Sub-Session-Id AVP (AVP Code 287) is of type Unsigned64 and contains the accounting sub-session identifier. The combination of the Session-Id and this AVP MUST be unique per sub-session, and the value of this AVP MUST be monotonically increased by one for all new sub-sessions. The absence of this AVP implies no sub-sessions are in use, with the exception of an Accounting-Request whose Accounting-Record-Type is set to STOP\_RECORD. A STOP\_RECORD message with no Accounting-Sub-Session-Id AVP present will signal the termination of all sub-sessions for a given Session-Id.

#### 9.8.7. Accounting-Realtime-Required AVP

The Accounting-Realtime-Required AVP (AVP Code 483) is of type Enumerated and is sent from the Diameter home authorization server to the Diameter client or in the Accounting-Answer from the accounting server. The client uses information in this AVP to decide what to do if the sending of accounting records to the accounting server has been temporarily prevented due to, for instance, a network problem.

**DELIVER\_AND\_GRANT**

1

The AVP with Value field set to DELIVER\_AND\_GRANT means that the service MUST only be granted as long as there is a connection to an accounting server. Note that the set of alternative accounting servers are treated as one server in this sense. Having to move the accounting record stream to a backup server is not a reason to discontinue the service to the user.

**GRANT\_AND\_STORE**

2

The AVP with Value field set to GRANT\_AND\_STORE means that service SHOULD be granted if there is a connection, or as long as records can still be stored as described in Section 9.4.

This is the default behavior if the AVP isn't included in the reply from the authorization server.

**GRANT\_AND\_LOSE**

3

The AVP with Value field set to GRANT\_AND\_LOSE means that service SHOULD be granted even if the records can not be delivered or stored.

**10. AVP Occurrence Table**

The following tables presents the AVPs defined in this document, and specifies in which Diameter messages they MAY, or MAY NOT be present. Note that AVPs that can only be present within a Grouped AVP are not represented in this table.

The table uses the following symbols:

- 0 The AVP MUST NOT be present in the message.
- 0+ Zero or more instances of the AVP MAY be present in the message.
- 0-1 Zero or one instance of the AVP MAY be present in the message. It is considered an error if there are more than one instance of the AVP.
- 1 One instance of the AVP MUST be present in the message.
- 1+ At least one instance of the AVP MUST be present in the message.

**10.1. Base Protocol Command AVP Table**

The table in this section is limited to the non-accounting Command Codes defined in this specification.

Attribute Name	Command-Code											
	CER	CEA	DPR	DPA	DWR	DWA	RAR	RAA	ASR	ASA	STR	STA
Acct-Interim-Interval	0	0	0	0	0	0	0-1	0	0	0	0	0
Accounting-Realtime-Required	0	0	0	0	0	0	0-1	0	0	0	0	0
Acct-Application-Id	0+	0+	0	0	0	0	0	0	0	0	0	0
Auth-Application-Id	0+	0+	0	0	0	0	1	0	1	0	1	0
Auth-Grace-Period	0	0	0	0	0	0	0	0	0	0	0	0
Auth-Request-Type	0	0	0	0	0	0	0	0	0	0	0	0
Auth-Session-State	0	0	0	0	0	0	0	0	0	0	0	0
Authorization-Lifetime	0	0	0	0	0	0	0	0	0	0	0	0
Class	0	0	0	0	0	0	0	0	0	0	0+	0+
Destination-Host	0	0	0	0	0	0	1	0	1	0	0-1	0
Destination-Realm	0	0	0	0	0	0	1	0	1	0	1	0
Disconnect-Cause	0	0	1	0	0	0	0	0	0	0	0	0
Error-Message	0	0-1	0	0-1	0	0-1	0	0-1	0	0-1	0	0-1
Error-Reporting-Host	0	0	0	0	0	0	0	0-1	0	0-1	0	0-1
Failed-AVP	0	0+	0	0+	0	0+	0	0+	0	0+	0	0+
Firmware-Revision	0-1	0-1	0	0	0	0	0	0	0	0	0	0
Host-IP-Address	1+	1+	0	0	0	0	0	0	0	0	0	0
Inband-Security-Id	0+	0+	0	0	0	0	0	0	0	0	0	0
Multi-Round-Time-Out	0	0	0	0	0	0	0	0	0	0	0	0
Origin-Host	1	1	1	1	1	1	1	1	1	1	1	1
Origin-Realm	1	1	1	1	1	1	1	1	1	1	1	1
Origin-State-Id	0-1	0-1	0	0	0-1	0-1	0-1	0-1	0-1	0-1	0-1	0-1
Product-Name	1	1	0	0	0	0	0	0	0	0	0	0
Proxy-Info	0	0	0	0	0	0	0+	0+	0+	0+	0+	0+
Redirect-Host	0	0	0	0	0	0	0	0+	0	0+	0	0+
Redirect-Host-Usage	0	0	0	0	0	0	0	0-1	0	0-1	0	0-1
Redirect-Max-Cache-Time	0	0	0	0	0	0	0	0-1	0	0-1	0	0-1
Result-Code	0	1	0	1	0	1	0	1	0	0	0	1
Re-Auth-Request-Type	0	0	0	0	0	0	1	0	0	0	0	0
Route-Record	0	0	0	0	0	0	0+	0	0+	0	0+	0
Session-Binding	0	0	0	0	0	0	0	0	0	0	0	0
Session-Id	0	0	0	0	0	0	1	1	1	1	1	1
Session-Server-Failover	0	0	0	0	0	0	0	0	0	0	0	0
Session-Timeout	0	0	0	0	0	0	0	0	0	0	0	0
Supported-Vendor-Id	0+	0+	0	0	0	0	0	0	0	0	0	0
Termination-Cause	0	0	0	0	0	0	0	0	0	0	1	0
User-Name	0	0	0	0	0	0	0-1	0-1	0-1	0-1	0-1	0-1
Vendor-Id	1	1	0	0	0	0	0	0	0	0	0	0

Vendor-Specific-Application-Id	0+	0+	0	0	0	0	0	0	0	0	0	0	0	0
-----	+	+	+	+	+	+	+	+	+	+	+	+	+	+

## 10.2. Accounting AVP Table

The table in this section is used to represent which AVPs defined in this document are to be present in the Accounting messages. These AVP occurrence requirements are guidelines, which may be expanded, and/or overridden by application-specific requirements in the Diameter applications documents.

Attribute Name	Command Code	
	ACR	ACA
Acct-Interim-Interval	0-1	0-1
Acct-Multi-Session-Id	0-1	0-1
Accounting-Record-Number	1	1
Accounting-Record-Type	1	1
Acct-Session-Id	0-1	0-1
Accounting-Sub-Session-Id	0-1	0-1
Accounting-Realtime-Required	0-1	0-1
Acct-Application-Id	0-1	0-1
Auth-Application-Id	0	0
Class	0+	0+
Destination-Host	0-1	0
Destination-Realm	1	0
Error-Reporting-Host	0	0+
Event-Timestamp	0-1	0-1
Origin-Host	1	1
Origin-Realm	1	1
Proxy-Info	0+	0+
Route-Record	0+	0+
Result-Code	0	1
Session-Id	1	1
Termination-Cause	0-1	0-1
User-Name	0-1	0-1
Vendor-Specific-Application-Id	0-1	0-1

## 11. IANA Considerations

This section provides guidance to the Internet Assigned Numbers Authority (IANA) regarding registration of values related to the Diameter protocol, in accordance with BCP 26 [IANA]. The following policies are used here with the meanings defined in BCP 26: "Private Use", "First Come First Served", "Expert Review", "Specification Required", "IETF Consensus", "Standards Action".

This section explains the criteria to be used by the IANA for assignment of numbers within namespaces defined within this document.

Diameter is not intended as a general purpose protocol, and allocations SHOULD NOT be made for purposes unrelated to authentication, authorization or accounting.

For registration requests where a Designated Expert should be consulted, the responsible IESG area director should appoint the Designated Expert. For Designated Expert with Specification Required, the request is posted to the AAA WG mailing list (or, if it has been disbanded, a successor designated by the Area Director) for comment and review, and MUST include a pointer to a public specification. Before a period of 30 days has passed, the Designated Expert will either approve or deny the registration request and publish a notice of the decision to the AAA WG mailing list or its successor. A denial notice must be justified by an explanation and, in the cases where it is possible, concrete suggestions on how the request can be modified so as to become acceptable.

### 11.1. AVP Header

As defined in Section 4, the AVP header contains three fields that requires IANA namespace management; the AVP Code, Vendor-ID and Flags field.

#### 11.1.1. AVP Codes

The AVP Code namespace is used to identify attributes. There are multiple namespaces. Vendors can have their own AVP Codes namespace which will be identified by their Vendor-ID (also known as Enterprise-Number) and they control the assignments of their vendor-specific AVP codes within their own namespace. The absence of a Vendor-ID or a Vendor-ID value of zero (0) identifies the IETF IANA controlled AVP Codes namespace. The AVP Codes and sometimes also possible values in an AVP are controlled and maintained by IANA.

AVP Code 0 is not used. AVP Codes 1-255 are managed separately as RADIUS Attribute Types [RADTYPE]. This document defines the AVP Codes 257-274, 276-285, 287, 291-300, 480, 483 and 485-486. See Section 4.5 for the assignment of the namespace in this specification.

AVPs may be allocated following Designated Expert with Specification Required [IANA]. Release of blocks of AVPs (more than 3 at a time for a given purpose) should require IETF Consensus.

Note that Diameter defines a mechanism for Vendor-Specific AVPs, where the Vendor-Id field in the AVP header is set to a non-zero value. Vendor-Specific AVPs codes are for Private Use and should be encouraged instead of allocation of global attribute types, for functions specific only to one vendor's implementation of Diameter, where no interoperability is deemed useful. Where a Vendor-Specific AVP is implemented by more than one vendor, allocation of global AVPs should be encouraged instead.

#### 11.1.2. AVP Flags

There are 8 bits in the AVP Flags field of the AVP header, defined in Section 4. This document assigns bit 0 ('V'endor Specific), bit 1 ('M'andatory) and bit 2 ('P'rotected). The remaining bits should only be assigned via a Standards Action [IANA].

#### 11.2. Diameter Header

As defined in Section 3, the Diameter header contains two fields that require IANA namespace management; Command Code and Command Flags.

##### 11.2.1. Command Codes

The Command Code namespace is used to identify Diameter commands. The values 0-255 are reserved for RADIUS backward compatibility, and are defined as "RADIUS Packet Type Codes" in [RADTYPE]. Values 256-16,777,213 are for permanent, standard commands, allocated by IETF Consensus [IANA]. This document defines the Command Codes 257, 258, 271, 274-275, 280 and 282. See Section 3.1 for the assignment of the namespace in this specification.

The values 16,777,214 and 16,777,215 (hexadecimal values 0xfffffe - 0xffffff) are reserved for experimental commands. As these codes are only for experimental and testing purposes, no guarantee is made for interoperability between Diameter peers using experimental commands, as outlined in [IANA-EXP].



### 11.2.2. Command Flags

There are eight bits in the Command Flags field of the Diameter header. This document assigns bit 0 ('R'equest), bit 1 ('P'roxy), bit 2 ('E'rror) and bit 3 ('T'). Bits 4 through 7 MUST only be assigned via a Standards Action [IANA].

### 11.3. Application Identifiers

As defined in Section 2.4, the Application Identifier is used to identify a specific Diameter Application. There are standards-track application ids and vendor specific application ids.

IANA [IANA] has assigned the range 0x00000001 to 0x00ffffff for standards-track applications; and 0x01000000 - 0xffffffff for vendor specific applications, on a first-come, first-served basis. The following values are allocated.

Diameter Common Messages	0
NASREQ	1 [NASREQ]
Mobile-IP	2 [DIAMMIP]
Diameter Base Accounting	3
Relay	0xffffffff

Assignment of standards-track application IDs are by Designated Expert with Specification Required [IANA].

Both Application-Id and Acct-Application-Id AVPs use the same Application Identifier space.

Vendor-Specific Application Identifiers, are for Private Use. Vendor-Specific Application Identifiers are assigned on a First Come, First Served basis by IANA.

### 11.4. AVP Values

Certain AVPs in Diameter define a list of values with various meanings. For attributes other than those specified in this section, adding additional values to the list can be done on a First Come, First Served basis by IANA.

#### 11.4.1. Result-Code AVP Values

As defined in Section 7.1, the Result-Code AVP (AVP Code 268) defines the values 1001, 2001-2002, 3001-3010, 4001-4002 and 5001-5017.

All remaining values are available for assignment via IETF Consensus [IANA].

#### 11.4.2. Accounting-Record-Type AVP Values

As defined in Section 9.8.1, the Accounting-Record-Type AVP (AVP Code 480) defines the values 1-4. All remaining values are available for assignment via IETF Consensus [IANA].

#### 11.4.3. Termination-Cause AVP Values

As defined in Section 8.15, the Termination-Cause AVP (AVP Code 295) defines the values 1-8. All remaining values are available for assignment via IETF Consensus [IANA].

#### 11.4.4. Redirect-Host-Usage AVP Values

As defined in Section 6.13, the Redirect-Host-Usage AVP (AVP Code 261) defines the values 0-5. All remaining values are available for assignment via IETF Consensus [IANA].

#### 11.4.5. Session-Server-Failover AVP Values

As defined in Section 8.18, the Session-Server-Failover AVP (AVP Code 271) defines the values 0-3. All remaining values are available for assignment via IETF Consensus [IANA].

#### 11.4.6. Session-Binding AVP Values

As defined in Section 8.17, the Session-Binding AVP (AVP Code 270) defines the bits 1-4. All remaining bits are available for assignment via IETF Consensus [IANA].

#### 11.4.7. Disconnect-Cause AVP Values

As defined in Section 5.4.3, the Disconnect-Cause AVP (AVP Code 273) defines the values 0-2. All remaining values are available for assignment via IETF Consensus [IANA].

#### 11.4.8. Auth-Request-Type AVP Values

As defined in Section 8.7, the Auth-Request-Type AVP (AVP Code 274) defines the values 1-3. All remaining values are available for assignment via IETF Consensus [IANA].

#### 11.4.9. Auth-Session-State AVP Values

As defined in Section 8.11, the Auth-Session-State AVP (AVP Code 277) defines the values 0-1. All remaining values are available for assignment via IETF Consensus [IANA].

#### 11.4.10. Re-Auth-Request-Type AVP Values

As defined in Section 8.12, the Re-Auth-Request-Type AVP (AVP Code 285) defines the values 0-1. All remaining values are available for assignment via IETF Consensus [IANA].

#### 11.4.11. Accounting-Realtime-Required AVP Values

As defined in Section 9.8.7, the Accounting-Realtime-Required AVP (AVP Code 483) defines the values 1-3. All remaining values are available for assignment via IETF Consensus [IANA].

#### 11.4.12. Inband-Security-Id AVP (code 299)

As defined in Section 6.10, the Inband-Security-Id AVP (AVP Code 299) defines the values 0-1. All remaining values are available for assignment via IETF Consensus [IANA].

#### 11.5. Diameter TCP/SCTP Port Numbers

The IANA has assigned TCP and SCTP port number 3868 to Diameter.

#### 11.6. NAPTR Service Fields

The registration in the RFC MUST include the following information:

Service Field: The service field being registered. An example for a new fictitious transport protocol called NCTP might be "AAA+D2N".

Protocol: The specific transport protocol associated with that service field. This MUST include the name and acronym for the protocol, along with reference to a document that describes the transport protocol. For example - "New Connectionless Transport Protocol (NCTP), RFC 5766".

Name and Contact Information: The name, address, email address and telephone number for the person performing the registration.

The following values have been placed into the registry:

Services Field	Protocol
AAA+D2T	TCP
AAA+D2S	SCTP

#### 12. Diameter protocol related configurable parameters

This section contains the configurable parameters that are found throughout this document:

#### Diameter Peer

A Diameter entity MAY communicate with peers that are statically configured. A statically configured Diameter peer would require that either the IP address or the fully qualified domain name (FQDN) be supplied, which would then be used to resolve through DNS.

#### Realm Routing Table

A Diameter proxy server routes messages based on the realm portion of a Network Access Identifier (NAI). The server MUST have a table of Realm Names, and the address of the peer to which the message must be forwarded to. The routing table MAY also include a "default route", which is typically used for all messages that cannot be locally processed.

#### Tc timer

The Tc timer controls the frequency that transport connection attempts are done to a peer with whom no active transport connection exists. The recommended value is 30 seconds.

### 13. Security Considerations

The Diameter base protocol assumes that messages are secured by using either IPsec or TLS. This security mechanism is acceptable in environments where there is no untrusted third party agent. In other situations, end-to-end security is needed.

Diameter clients, such as Network Access Servers (NASes) and Mobility Agents MUST support IP Security [SECARCH] and MAY support TLS [TLS]. Diameter servers MUST support TLS and IPsec. Diameter implementations MUST use transmission-level security of some kind (IPsec or TLS) on each connection.

If a Diameter connection is not protected by IPsec, then the CER/CEA exchange MUST include an Inband-Security-ID AVP with a value of TLS. For TLS usage, a TLS handshake will begin when both ends are in the open state, after completion of the CER/CEA exchange. If the TLS handshake is successful, all further messages will be sent via TLS. If the handshake fails, both ends move to the closed state.

It is suggested that IPsec be used primarily at the edges for intra-domain exchanges. For NAS devices without certificate support, pre-shared keys can be used between the NAS and a local AAA proxy.

For protection of inter-domain exchanges, TLS is recommended. See Sections 13.1 and 13.2 for more details on IPsec and TLS usage.

### 13.1. IPsec Usage

All Diameter implementations MUST support IPsec ESP [IPsec] in transport mode with non-null encryption and authentication algorithms to provide per-packet authentication, integrity protection and confidentiality, and MUST support the replay protection mechanisms of IPsec.

Diameter implementations MUST support IKE for peer authentication, negotiation of security associations, and key management, using the IPsec DOI [IPSECDOI]. Diameter implementations MUST support peer authentication using a pre-shared key, and MAY support certificate-based peer authentication using digital signatures. Peer authentication using the public key encryption methods outlined in IKE's Sections 5.2 and 5.3 [IKE] SHOULD NOT be used.

Conformant implementations MUST support both IKE Main Mode and Aggressive Mode. When pre-shared keys are used for authentication, IKE Aggressive Mode SHOULD be used, and IKE Main Mode SHOULD NOT be used. When digital signatures are used for authentication, either IKE Main Mode or IKE Aggressive Mode MAY be used.

When digital signatures are used to achieve authentication, an IKE negotiator SHOULD use IKE Certificate Request Payload(s) to specify the certificate authority (or authorities) that are trusted in accordance with its local policy. IKE negotiators SHOULD use pertinent certificate revocation checks before accepting a PKI certificate for use in IKE's authentication procedures.

The Phase 2 Quick Mode exchanges used to negotiate protection for Diameter connections MUST explicitly carry the Identity Payload fields (IDci and IDcr). The DOI provides for several types of identification data. However, when used in conformant implementations, each ID Payload MUST carry a single IP address and a single non-zero port number, and MUST NOT use the IP Subnet or IP Address Range formats. This allows the Phase 2 security association to correspond to specific TCP and SCTP connections.

Since IPsec acceleration hardware may only be able to handle a limited number of active IKE Phase 2 SAs, Phase 2 delete messages may be sent for idle SAs, as a means of keeping the number of active Phase 2 SAs to a minimum. The receipt of an IKE Phase 2 delete message SHOULD NOT be interpreted as a reason for tearing down a Diameter connection. Rather, it is preferable to leave the connection up, and if additional traffic is sent on it, to bring up another IKE Phase 2 SA to protect it. This avoids the potential for continually bringing connections up and down.

### 13.2. TLS Usage

A Diameter node that initiates a connection to another Diameter node acts as a TLS client according to [TLS], and a Diameter node that accepts a connection acts as a TLS server. Diameter nodes implementing TLS for security MUST mutually authenticate as part of TLS session establishment. In order to ensure mutual authentication, the Diameter node acting as TLS server must request a certificate from the Diameter node acting as TLS client, and the Diameter node acting as TLS client MUST be prepared to supply a certificate on request.

Diameter nodes MUST be able to negotiate the following TLS cipher suites:

```
TLS_RSA_WITH_RC4_128_MD5
TLS_RSA_WITH_RC4_128_SHA
TLS_RSA_WITH_3DES_EDE_CBC_SHA
```

Diameter nodes SHOULD be able to negotiate the following TLS cipher suite:

```
TLS_RSA_WITH_AES_128_CBC_SHA
```

Diameter nodes MAY negotiate other TLS cipher suites.

### 13.3. Peer-to-Peer Considerations

As with any peer-to-peer protocol, proper configuration of the trust model within a Diameter peer is essential to security. When certificates are used, it is necessary to configure the root certificate authorities trusted by the Diameter peer. These root CAs are likely to be unique to Diameter usage and distinct from the root CAs that might be trusted for other purposes such as Web browsing. In general, it is expected that those root CAs will be configured so as to reflect the business relationships between the organization hosting the Diameter peer and other organizations. As a result, a Diameter peer will typically not be configured to allow connectivity with any arbitrary peer. When certificate authentication Diameter peers may not be known beforehand, and therefore peer discovery may be required.

Note that IPsec is considerably less flexible than TLS when it comes to configuring root CAs. Since use of Port identifiers is prohibited within IKE Phase 1, within IPsec it is not possible to uniquely configure trusted root CAs for each application individually; the same policy must be used for all applications. This implies, for example, that a root CA trusted for use with Diameter must also be

trusted to protect SNMP. These restrictions can be awkward at best. Since TLS supports application-level granularity in certificate policy, TLS SHOULD be used to protect Diameter connections between administrative domains. IPsec is most appropriate for intra-domain usage when pre-shared keys are used as a security mechanism.

When pre-shared key authentication is used with IPsec to protect Diameter, unique pre-shared keys are configured with Diameter peers, who are identified by their IP address (Main Mode), or possibly their FQDN (Aggressive Mode). As a result, it is necessary for the set of Diameter peers to be known beforehand. Therefore, peer discovery is typically not necessary.

The following is intended to provide some guidance on the issue.

It is recommended that a Diameter peer implement the same security mechanism (IPsec or TLS) across all its peer-to-peer connections. Inconsistent use of security mechanisms can result in redundant security mechanisms being used (e.g., TLS over IPsec) or worse, potential security vulnerabilities. When IPsec is used with Diameter, a typical security policy for outbound traffic is "Initiate IPsec, from me to any, destination port Diameter"; for inbound traffic, the policy would be "Require IPsec, from any to me, destination port Diameter".

This policy causes IPsec to be used whenever a Diameter peer initiates a connection to another Diameter peer, and to be required whenever an inbound Diameter connection occurs. This policy is attractive, since it does not require policy to be set for each peer or dynamically modified each time a new Diameter connection is created; an IPsec SA is automatically created based on a simple static policy. Since IPsec extensions are typically not available to the sockets API on most platforms, and IPsec policy functionality is implementation dependent, use of a simple static policy is the often the simplest route to IPsec-enabling a Diameter implementation.

One implication of the recommended policy is that if a node is using both TLS and IPsec, there is not a convenient way in which to use either TLS or IPsec, but not both, without reserving an additional port for TLS usage. Since Diameter uses the same port for TLS and non-TLS usage, where the recommended IPsec policy is put in place, a TLS-protected connection will match the IPsec policy, and both IPsec and TLS will be used to protect the Diameter connection. To avoid this, it would be necessary to plumb peer-specific policies either statically or dynamically.

If IPsec is used to secure Diameter peer-to-peer connections, IPsec policy SHOULD be set so as to require IPsec protection for inbound connections, and to initiate IPsec protection for outbound connections. This can be accomplished via use of inbound and outbound filter policy.

## 14. References

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## Appendix A. Diameter Service Template

The following service template describes the attributes used by Diameter servers to advertise themselves. This simplifies the process of selecting an appropriate server to communicate with. A Diameter client can request specific Diameter servers based on characteristics of the Diameter service desired (for example, an AAA server to use for accounting.)

Name of submitter: "Erik Guttman" <Erik.Guttman@sun.com> Language of service template: en

## Security Considerations:

Diameter clients and servers use various cryptographic mechanisms to protect communication integrity, confidentiality as well as perform end-point authentication. It would thus be difficult if not impossible for an attacker to advertise itself using SLPv2 and pose as a legitimate Diameter peer without proper preconfigured secrets or cryptographic keys. Still, as Diameter services are vital for network operation it is important to use SLPv2 authentication to prevent an attacker from modifying or eliminating service advertisements for legitimate Diameter servers.

## Template text:

-----template begins here-----

template-type=service:diameter

template-version=0.0

template-description=

The Diameter protocol is defined by RFC 3588.

template-url-syntax=

url-path= ; The Diameter URL format is described in Section 2.9.

; Example: 'aaa://aaa.example.com:1812;transport=tcp

supported-auth-applications= string L M

# This attribute lists the Diameter applications supported by the  
# AAA implementation. The applications currently defined are:

# Application Name Defined by

# -----

# NASREQ Diameter Network Access Server Application

# MobileIP Diameter Mobile IP Application

#

# Notes:

# . Diameter implementations support one or more applications.

# . Additional applications may be defined in the future.

# An updated service template will be created at that time.

```

#
NASREQ,MobileIP

supported-acct-applications= string L M
# This attribute lists the Diameter applications supported by the
# AAA implementation.  The applications currently defined are:
# Application Name      Defined by
# -----
# NASREQ                Diameter Network Access Server Application
# MobileIP              Diameter Mobile IP Application
#
# Notes:
#   . Diameter implementations support one or more applications.
#   . Additional applications may be defined in the future.
#   . An updated service template will be created at that time.
#
NASREQ,MobileIP

supported-transports= string L M
SCTP
# This attribute lists the supported transports that the Diameter
# implementation accepts.  Note that a compliant Diameter
# implementation MUST support SCTP, though it MAY support other
# transports, too.
SCTP,TCP

```

-----template ends here-----

## Appendix B. NAPTR Example

As an example, consider a client that wishes to resolve `aaa:ex.com`. The client performs a NAPTR query for that domain, and the following NAPTR records are returned:

```

;;          order pref flags service          regexp replacement
IN NAPTR 50  50  "s"  "AAA+D2S"              ""
_diameter._sctp.example.com IN NAPTR 100  50  "s"  "AAA+D2T"
"" _aaa._tcp.example.com

```

This indicates that the server supports SCTP, and TCP, in that order. If the client supports over SCTP, SCTP will be used, targeted to a host determined by an SRV lookup of `_diameter._sctp.ex.com`. That lookup would return:

```

;;          Priority Weight Port   Target
IN SRV  0      1      5060  server1.example.com IN SRV  0
2      5060  server2.example.com

```

## Appendix C. Duplicate Detection

As described in Section 9.4, accounting record duplicate detection is based on session identifiers. Duplicates can appear for various reasons:

- Failover to an alternate server. Where close to real-time performance is required, failover thresholds need to be kept low and this may lead to an increased likelihood of duplicates. Failover can occur at the client or within Diameter agents.
- Failure of a client or agent after sending of a record from non-volatile memory, but prior to receipt of an application layer ACK and deletion of the record. record to be sent. This will result in retransmission of the record soon after the client or agent has rebooted.
- Duplicates received from RADIUS gateways. Since the retransmission behavior of RADIUS is not defined within [RFC2865], the likelihood of duplication will vary according to the implementation.
- Implementation problems and misconfiguration.

The T flag is used as an indication of an application layer retransmission event, e.g., due to failover to an alternate server. It is defined only for request messages sent by Diameter clients or agents. For instance, after a reboot, a client may not know whether it has already tried to send the accounting records in its non-volatile memory before the reboot occurred. Diameter servers MAY use the T flag as an aid when processing requests and detecting duplicate messages. However, servers that do this MUST ensure that duplicates are found even when the first transmitted request arrives at the server after the retransmitted request. It can be used only in cases where no answer has been received from the Server for a request and the request is sent again, (e.g., due to a failover to an alternate peer, due to a recovered primary peer or due to a client re-sending a stored record from non-volatile memory such as after reboot of a client or agent).

In some cases the Diameter accounting server can delay the duplicate detection and accounting record processing until a post-processing phase takes place. At that time records are likely to be sorted according to the included User-Name and duplicate elimination is easy in this case. In other situations it may be necessary to perform real-time duplicate detection, such as when credit limits are imposed or real-time fraud detection is desired.

In general, only generation of duplicates due to failover or re-sending of records in non-volatile storage can be reliably detected by Diameter clients or agents. In such cases the Diameter client or agents can mark the message as possible duplicate by setting the T flag. Since the Diameter server is responsible for duplicate detection, it can choose to make use of the T flag or not, in order to optimize duplicate detection. Since the T flag does not affect interoperability, and may not be needed by some servers, generation of the T flag is REQUIRED for Diameter clients and agents, but MAY be implemented by Diameter servers.

As an example, it can be usually be assumed that duplicates appear within a time window of longest recorded network partition or device fault, perhaps a day. So only records within this time window need to be looked at in the backward direction. Secondly, hashing techniques or other schemes, such as the use of the T flag in the received messages, may be used to eliminate the need to do a full search even in this set except for rare cases.

The following is an example of how the T flag may be used by the server to detect duplicate requests.

A Diameter server MAY check the T flag of the received message to determine if the record is a possible duplicate. If the T flag is set in the request message, the server searches for a duplicate within a configurable duplication time window backward and forward. This limits database searching to those records where the T flag is set. In a well run network, network partitions and device faults will presumably be rare events, so this approach represents a substantial optimization of the duplicate detection process. During failover, it is possible for the original record to be received after the T flag marked record, due to differences in network delays experienced along the path by the original and duplicate transmissions. The likelihood of this occurring increases as the failover interval is decreased. In order to be able to detect out of order duplicates, the Diameter server should use backward and forward time windows when performing duplicate checking for the T flag marked request. For example, in order to allow time for the original record to exit the network and be recorded by the accounting server, the Diameter server can delay processing records with the T flag set until a time period `TIME_WAIT + RECORD_PROCESSING_TIME` has elapsed after the closing of the original transport connection. After this time period has expired, then it may check the T flag marked records against the database with relative assurance that the original records, if sent, have been received and recorded.



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