

Network Working Group
Internet-Draft
Intended status: Informational
Expires: 10 August 2026

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6 February 2026

A Simple BGP-based Mobile Routing System for the Aeronautical
Telecommunications Network
draft-ietf-rtgwg-atn-bgp-30

Abstract

The International Civil Aviation Organization (ICAO) is investigating mobile routing solutions for a worldwide Aeronautical Telecommunications Network with Internet Protocol Services (ATN/IPS). The ATN/IPS will eventually replace existing communication services with an IP-based service supporting pervasive Air Traffic Management (ATM) for Air Traffic Controllers (ATC), Airline Operations Controllers (AOC), and all commercial aircraft worldwide. This informational document describes a simple and extensible mobile routing service based on the industry-standard Border Gateway Protocol (BGP) to address the ATN/IPS requirements.

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1. Introduction

The worldwide Air Traffic Management (ATM) system today uses a service known as Aeronautical Telecommunications Network based on Open Systems Interconnection (ATN/OSI). The service is used to augment controller to pilot voice communications with rudimentary short text command and control messages. The service has seen successful deployment in some worldwide ATM domains.

The International Civil Aviation Organization (ICAO) is now engaged in the development of a next-generation replacement for ATN/OSI known as Aeronautical Telecommunications Network with Internet Protocol Services (ATN/IPS) [ATN][ATN-IPS]. ATN/IPS will eventually provide an Internetworking service supporting pervasive ATM for Air Traffic Controllers (ATC), Airline Operations Controllers (AOC), and all commercial aircraft worldwide. The ATN/IPS will require a new mobile routing service as a core element of the architecture. This document presents an approach based on the Border Gateway Protocol (BGP) [RFC4271].

Aircraft communicate via wireless aviation data links that typically support much lower data rates than terrestrial wireless and wired-line communications. For example, some Very High Frequency (VHF)-based data links only support data rates on the order of 32Kbps and an emerging L-Band data link that is expected to play a key role in future aeronautical communications only supports rates on the order of 1Mbps. Although satellite data links can provide much higher data rates during optimal conditions, like any other aviation data link they are subject to errors, delay, disruption, signal intermittence, degradation due to atmospheric conditions, etc.

The well-connected ground domain ATN/IPS network should therefore treat each safety-of-flight critical packet produced by (or destined to) an aircraft as a precious commodity and strive for an optimized service that provides the highest possible degree of reliability. Furthermore, continuous performance-intensive control messaging services such as BGP peering sessions must be carried only over the well-connected secured ground domain ATN/IPS network and never over low-end aviation data links.

The ATN/IPS is an IP-based overlay network configured over one or more Internetworking underlays ("INETs") maintained by aeronautical network service providers such as ARINC, SITA and Inmarsat. The Overlay Multilink Network Interface (OMNI) [I-D.templin-6man-omni3] uses an adaptation layer encapsulation to create a Non-Broadcast, Multiple Access (NBMA) virtual link spanning the entire ATN/IPS. Each aircraft connects to the OMNI link via an OMNI interface configured over the aircraft's underlying physical and/or virtual access network interfaces.

Each underlying INET comprises one or more "partitions" where all nodes within a partition can exchange packets with all other nodes, i.e., the partition is connected internally. There is no requirement that each INET partition uses the same IP protocol version nor has consistent IP addressing plans in relation to other partitions. Instead, the OMNI link sees each partition as a "segment" of a lower layer topology concatenated by a service known as the OMNI Adaptation Layer (OAL) [I-D.templin-6man-omni3] based on IPv6 encapsulation [RFC2473].

The IPv6 addressing architecture provides different classes of addresses, including Global Unicast Addresses (GUAs) and Link-Local Addresses (LLAs) [RFC4291]. A new address class known as Multilink-Local Addresses (MLAs) [I-D.templin-6man-mla] is also necessary to support the OMNI virtual link abstraction over an end system's multiple underlying interfaces. All nodes on the OMNI link configure a unique MLA used for multilink local routing and addressing within the OMNI limited domain.

The ATN/IPS receives IPv6 GUA Mobility Service Prefixes (MSPs) from an Internet assigned numbers authority, and each aircraft will receive a MSP-based Mobile Network Prefix (MNP) delegation that travels with the aircraft. ATCs and AOCs will similarly configure addresses from an MNP or Foreign Network Prefix (FNP) exposed by different networks. (Note that while IPv6 GUAs are assumed for ATN/IPS, IPv4 with public/private addresses could also be used.)

Each aircraft Client configures the Subnet Router Anycast (SRA) address for each of its delegated MNPs and associates them with First-Hop Segment (FHS) Proxy/Servers on their underlying interfaces while coordinating with a Mobility Anchor Point (MAP) Proxy/Server to register their MNPs and MLAs in the routing system. Due to MNP/MLA delegation policies and random node mobility properties, MNPs and MLAs are generally not aggregable in the BGP routing service and are often represented as many more-specific prefixes instead of a smaller number of aggregated prefixes.

BGP routing service infrastructure nodes additionally configure unique MLAs that are persistently present and unchanging in the routing system. The BGP routing services therefore establish forwarding table entries based on these MNPs and MLAs which support adaptation layer routing and addressing.

When a source Client's OMNI interface forwards an original IP packet, the OAL performs IPv6 encapsulation using source and target MLAs as Source and Destination Addresses and with a Segment Routing Header (SRH) which includes the MLAs of intermediate systems. The OAL source next subjects each resulting "OAL packet" to IPv6

fragmentation and header compression, then encapsulates each fragment in INET headers specific to the underlying Internetwork. These resulting "carrier packets" then traverse a series of Internetworks connected by OAL intermediate systems which re-encapsulate them in new INET headers for traversal of the next Internetwork in succession. The carrier packets finally arrive at the OAL destination which performs adaptation layer decapsulation and reassembly to restore the original IP packet. A high-level ATN/IPS network diagram is shown in Figure 1:

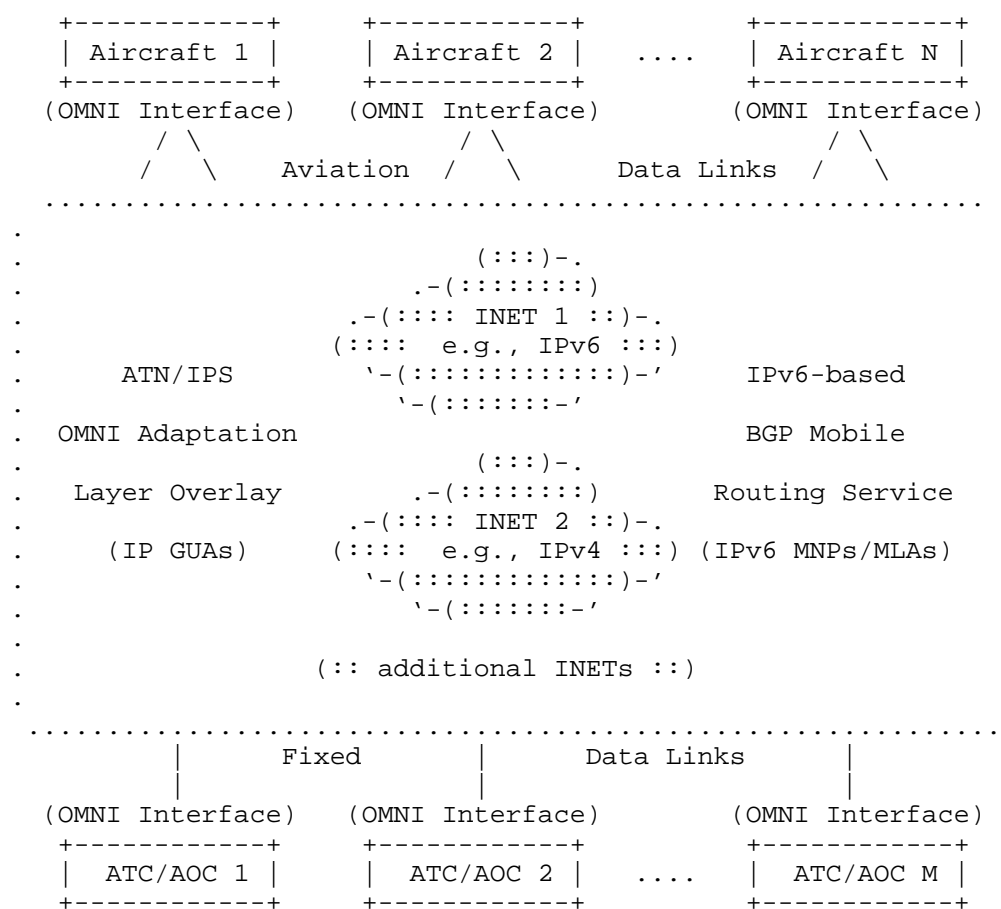


Figure 1: ATN/IPS Network Diagram

Connexion By Boeing [CBB] was an early aviation mobile routing service based on dynamic updates in the global public Internet BGP routing system. Practical experience with the approach has shown

that frequent injections and withdrawals of prefixes in the Internet routing system can result in excessive BGP update messaging, slow routing table convergence times, and extended outages when no route is available. This is due to both conservative default BGP protocol timing parameters (see Section 6) and the complex peering interconnections of BGP routers within the global Internet infrastructure. The situation is further exacerbated by frequent aircraft mobility events that each result in BGP updates which must be propagated to all BGP routers in the Internet with full routing tables.

We therefore consider a routing system using an overlay network that maintains a private BGP routing protocol instance between ATN/IPS Autonomous System (AS) Border Routers (ASBRs). The private BGP instance does not interact with the native BGP routing systems in underlying INETs, and BGP updates are unidirectional from "stub" ASBRs (s-ASBRs) to a small set of "core" ASBRs (c-ASBRs) in a hub-and-spokes topology. No extensions to the BGP protocol are necessary, and BGP routing is based on adaptation layer MNPs and MLAs instead of network or data link layer public/private IP prefixes. This allows ASBRs to perform adaptation layer forwarding based on intermediate layer IPv6 header information instead of network layer forwarding based on upper layer IP header information or link layer forwarding based on lower layer IP header information.

The s-ASBRs for each stub AS connect to a small number of c-ASBRs via secured direct point-to-point links and/or secured tunnels (e.g., IPsec [RFC4301], etc.) over the underlying INET. Neighboring c-ASBRs should also use IP layer or lower-layer security services over their connecting links to ensure INET layer security.

The s-ASBRs engage in external BGP (eBGP) peerings with their respective c-ASBRs, and only maintain routing table entries for the MNPs/MLAs currently active within the stub AS. The s-ASBRs send BGP updates for MNP/MLA injections or withdrawals to c-ASBRs but do not receive any BGP updates from c-ASBRs. Instead, the s-ASBRs maintain default routes with their c-ASBRs as the next hop, and therefore hold only partial topology information.

The c-ASBRs connect to other c-ASBRs within the same partition using internal BGP (iBGP) peerings over which they collaboratively maintain a full routing table for all active MNPs/MLAs currently in service within the partition. Therefore, only the c-ASBRs maintain a full BGP routing table and never send any BGP updates to s-ASBRs. This simple routing model therefore greatly reduces the number of BGP updates that need to be synchronized among peers, and the number is reduced further still when intradomain routing changes within stub ASes are processed within the AS instead of being propagated to the core. BGP Route Reflectors (RRs) [RFC4456] can also be used to support increased scaling properties.

With these intra- and inter-INET BGP peerings in place, a forwarding plane spanning tree is established that properly covers the entire operating domain. All nodes in the network can be visited using strict spanning tree hops, but in many instances this may result in longer paths than are necessary. AERO [I-D.templin-6man-aero3] provides an example service for discovering and utilizing (route-optimized) shortcuts that do not always follow strict spanning tree paths.

The remainder of this document discusses the proposed BGP-based ATN/IPS mobile routing service.

2. Terminology

The terms Autonomous System (AS) and Autonomous System Border Router (ASBR) are the same as defined in [RFC4271]. The term Internet Protocol (IP) refers generically to either protocol version unless specifically qualified as IPv4 [RFC0791] or IPv6 [RFC8200].

The following terms are defined for the purposes of this document:

Air Traffic Management (ATM)

The worldwide service for coordinating safe aviation operations.

Air Traffic Controller (ATC)

A government agent responsible for coordinating with aircraft within a defined operational region via voice and/or data Command and Control messaging.

Airline Operations Controller (AOC)

An airline agent responsible for tracking and coordinating with aircraft within their fleet.

Aeronautical Telecommunications Network with Internet Protocol Services (ATN/IPS)

A future aviation network for ATCs and AOCs to coordinate with all

aircraft operating worldwide. The ATN/IPS will provide an IP-based overlay network service that connects access networks via encapsulation and forwarding over one or more Internetworking underlays.

Internetworking underlay ("INET")

A wide-area network that supports overlay network encapsulation/forwarding and connects Radio Access Networks to the rest of the ATN/IPS. Example INET service providers for civil aviation include ARINC, SITA and Inmarsat.

(Radio) Access Network ("ANET")

An aviation radio data link service provider's network, including radio transmitters and receivers as well as supporting ground-domain infrastructure needed to convey a customer's data packets to outside INETs. The term ANET is intended in the same spirit as for radio-based Internet service provider networks (e.g., cellular operators), but can also refer to ground-domain networks that connect AOCs and ATCs.

partition (or "segment")

A fully-connected internal subnetwork of an INET in which all nodes can communicate with all other nodes within the same partition using the same IP protocol version and addressing plan. Each INET consists of one or more partitions.

Overlay Multilink Network Interface (OMNI)

A virtual layer 2 bridging service that presents an ATN/IPS overlay unified link view even though the underlay may consist of multiple INET partitions. The OMNI virtual link is manifested through nested encapsulation in which original IP packets from ATN/IPS Clients are first encapsulated in adaptation layer IPv6 headers which are then forwarded to the next hop using INET encapsulation if necessary. Forwarding over the OMNI virtual link is therefore based on the adaptation layer addresses instead of the original IP addresses. In this way, packets sent from a source can be conveyed over the OMNI virtual link even though there may be many underlying INET partitions in the path to the destination.

OMNI Adaptation Layer (OAL)

A middle layer below the IP layer but above the INET layer that forwards original IP packets by applying IPv6 encapsulation, fragmentation and header compression followed by INET encapsulation to form "carrier packets". End systems that configure OMNI interfaces act as the OAL source and destination, while intermediate systems with OMNI interfaces act as OAL forwarding nodes. Regardless of the number of intermediate

systems in the path, the network layer IP TTL/Hop Limit is not decremented during (OAL layer) forwarding. Further details on OMNI and the OAL are found in [I-D.templin-6man-omni3].

OAL Autonomous System (OAL AS)

A "hub-of-hubs" autonomous system maintained through peerings between the core autonomous systems of different OMNI virtual link partitions.

Core Autonomous System Border Router (c-ASBR)

A BGP router located in the hub of the INET partition hub-and-spokes overlay network topology.

Core Autonomous System (Core AS)

The "hub" autonomous system maintained by all c-ASBRs within the same partition.

Stub Autonomous System Border Router (s-ASBR)

A BGP router configured as a spoke in the INET partition hub-and-spokes overlay network topology.

Stub Autonomous System (Stub AS)

A logical grouping that includes all Clients currently associated with a given s-ASBR.

Client

An ATC, AOC or aircraft that connects to the ATN/IPS as a leaf node. The Client could be a singleton host, or a router that connects a mobile or fixed network.

Proxy/Server

An ANET/INET border node that acts as a transparent intermediary between Clients and s-ASBRs. From the Client's perspective, the Proxy/Server presents the appearance that the Client is communicating directly with the s-ASBR. From the s-ASBR's perspective, the Proxy/Server presents the appearance that the s-ASBR is communicating directly with the Client.

Mobile Network Prefix (MNP)

An IP prefix that is delegated to an ATN/IPS end system, including mobile aircraft and in some cases also ATCs, AOCs, etc. in fixed configurations. Each end system configures and responds to a Subnet Router Anycast (SRA) GUA taken from the MNP.

Mobility Service Prefix (MSP)

An aggregated IP prefix assigned to the ATN/IPS by an Internet assigned numbers authority, and from which all MNPs are delegated (e.g., up to 2^{32} IPv6 /56 MNPs could be delegated from a /24 MSP).

Foreign Network Prefix (FNP)

An IP prefix assigned to another link or network outside of the OMNI link domain. Relay s-ASBRs advertise the FNPs into the BGP routing service and act as routers to forward packets between MNP correspondents on the OMNI link and FNP correspondents on other links.

Multilink Local Address (MLA)

A unique IPv6 address assigned by each node on the OMNI link according to [I-D.templin-6man-mla]. MLAs are assured globally unique but routable only within the OMNI link limited domain. MLAs have a lesser IPv6 address selection preference than MNP/FNP addresses.

3. ATN/IPS Routing System

The ATN/IPS routing system comprises a private BGP instance coordinated in an overlay network via tunnels between neighboring ASBRs over one or more underlying INETs. The ATN/IPS routing system interacts with underlying INET BGP routing systems only through the static advertisement of a small and unchanging set of MSPs instead of the full dynamically changing set of MNPs.

Within each INET partition, each s-ASBR connects a stub AS to the INET partition core using a distinct stub AS Number (ASN). Each s-ASBR further uses eBGP to peer with one or more c-ASBRs. All c-ASBRs are members of the INET partition core AS, and use a shared core ASN. Unique ASNs are assigned according to the standard 32-bit ASN format [RFC4271][RFC6793]. ASNs are allocated and managed by an ATN/IPS assigned numbers authority established by ICAO, which must ensure that ASNs are responsibly distributed without duplication and/or overlap.

The c-ASBRs use iBGP to maintain a synchronized consistent view of all active MNPs/MLAs currently in service within the INET partition. Figure 2 below represents the reference INET partition deployment. (Note that the figure shows details for only two s-ASBRs (s-ASBR1 and s-ASBR2) due to space constraints, but the other s-ASBRs should be understood to have similar Stub AS, MNP and eBGP peering arrangements.) The solution described in this document is flexible enough to extend to these topologies.

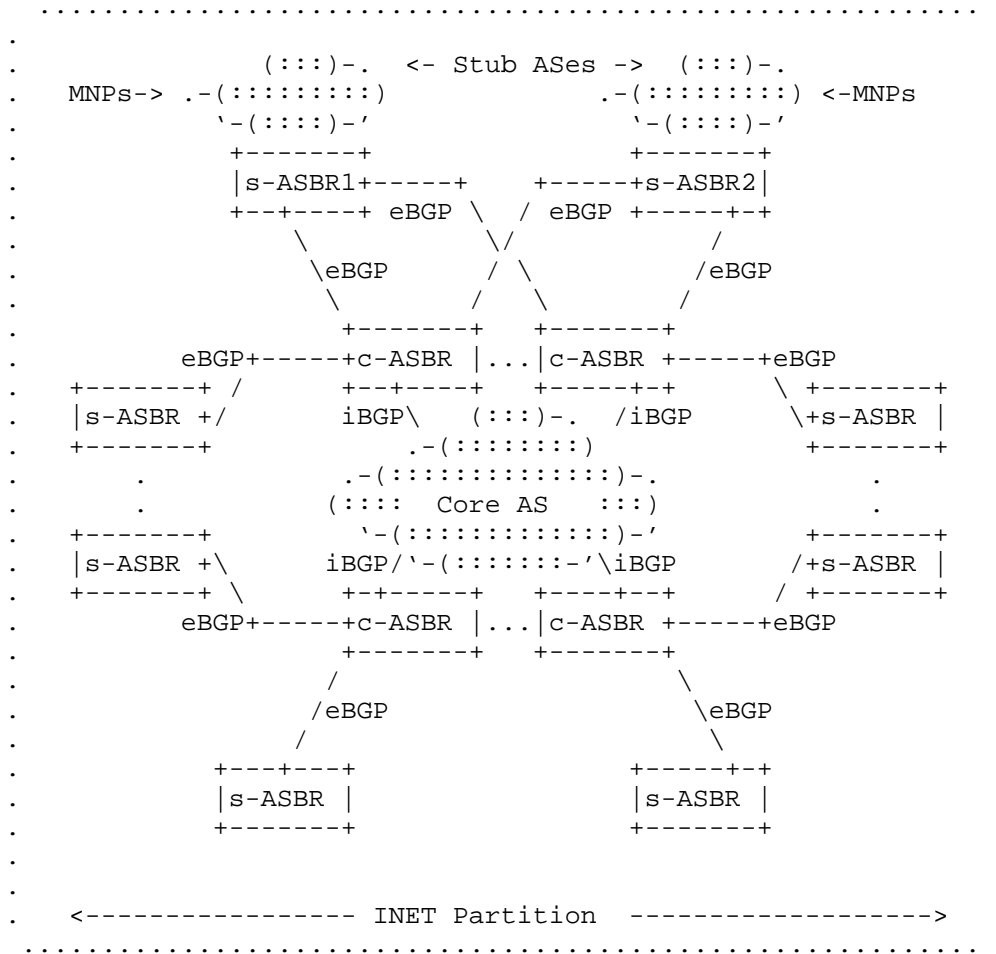


Figure 2: INET Partition Reference Deployment

In the reference deployment, each s-ASBR maintains routes for active MNPs/MLAs that currently belong to its stub AS. In response to "Inter-domain" mobility events, each s-ASBR dynamically announces new MNPs/MLAs and withdraws departed MNPs/MLAs in its eBGP updates to c-ASBRs. Since ATN/IPS end systems are expected to remain within the same stub AS for extended timeframes, however, intra-domain mobility events (such as an aircraft handing off between cell towers) are handled within the stub AS instead of being propagated as inter-domain eBGP updates.

Each c-ASBR configures a black-hole route for each of its MSPs. By black-holing the MSPs, the c-ASBR maintains forwarding table entries only for the MNPs that are currently active. If an arriving packet matches a black-hole route without matching an MNP, the c-ASBR should drop the packet and may also generate an ICMPv6 Destination Unreachable message [RFC4443], i.e., without forwarding the packet outside of the ATN/IPS overlay based on a less-specific route.

The c-ASBRs do not send BGP updates for MNPs/MLAs to s-ASBRs, but instead originate a default route. In this way, s-ASBRs have only partial topology knowledge (i.e., they know only about the active MNPs/MLAs currently within their stub ASes) and they forward all other packets to c-ASBRs which have full topology knowledge.

Each s-ASBR and c-ASBR configures an MLA that is permanently announced into the routing system. The core ASes of each INET partition are joined together through external BGP peerings. The c-ASBRs of each partition establish external peerings with the c-ASBRs of other partitions to form a "core-of-cores" OMNI link AS. The OMNI link AS contains the global knowledge of all MNPs/MLAs deployed worldwide, and supports ATN/IPS overlay communications between nodes located in different INET partitions by virtue of OAL encapsulation. Figure 3 shows a reference OAL topology.

distributed across multiple c-ASBR sets instead of concentrated in a single c-ASBR set. For example, a first c-ASBR set could aggregate an MSP segment A::/32, a second set could aggregate B::/32, a third could aggregate C::/32, etc. The union of all MSP segments would then constitute the collective MSP(s) for the entire ATN/IPS, with potential for supporting many millions of mobile networks or more.

In this design, each set of c-ASBRs services a specific set of MSPs, and each s-ASBR configures MSP-specific routes that list the correct set of c-ASBRs as next hops. This also allows for natural incremental deployment, and can support initial medium-scale deployments followed by dynamic deployment of additional ATN/IPS infrastructure elements without disturbing the already-deployed base. For example, additional c-ASBRs can be added later if service demand ever outgrows the initial deployment. For larger-scale applications (such as unmanned air vehicles and terrestrial vehicles) even larger scales can be accommodated by adding more c-ASBRs.

Consider now that the c-ASBRs provide adaptation layer gateways between independent Internetworks to form a true network-of-networks supporting the ATN/IPS overlay. A similar arrangement was first envisioned by the "catenet" model [POUZIN] [IEN48].

4. ATN/IPS (Radio) Access Network (ANET) Model

(Radio) Access Networks (ANETs) connect end system Clients such as aircraft, ATCs, AOCs etc. to the ATN/IPS routing system. Clients may connect to multiple ANETs at once, for example, when they have satellite, cellular, WiFi and/or other data links activated simultaneously. Each Client configures an OMNI interface [I-D.templin-6man-omni3] over its underlying ANET interfaces as a connection to an NBMA virtual link (manifested by the OAL) that spans the entire ATN/IPS. Clients may further move between ANETs in a manner that is perceived as a network layer mobility event. Clients should therefore employ a multilink/mobility routing service such as those discussed in Section 8.

Clients and s-ASBRs assign MLAs to their OMNI interfaces for use as adaptation layer Source and Destination Addresses during encapsulation. Clients and s-ASBRs use the MLAs to establish Neighbor Cache Entries (NCEs) via the OMNI interface as well as for (multihop) forwarding within the Access Networks. The s-ASBRs in turn rewrite Client MLAs to their own MLAs to support forwarding over the ATN/IPS interdomain region. Clients register their active data link connections with their serving s-ASBRs as discussed in Section 3. Clients may connect to s-ASBRs either directly, or via a Proxy/Server at the ANET/INET boundary.

```

+-----+ Client +-----+
|         |       |         |
+-----+ OMNI Interface +-----+
|         |       |         |
+-----+       +-----+

Data Link "A"      Data Link "B"
+-----+          +-----+
|         |        |         |
+-----+          +-----+

( ::: )-.          ( ::: )-.
.- ( :::::::::: ) <- (Radio) Access Networks -> .- ( :::::::::: )
'-( ::: )-'        '-( ::: )-'

+-----+          +-----+
|   P/S   |        |   P/S   |
+-----+          +-----+

...                ...

^ ^               ^ ^
||                ||
||                ||
++===== > | s-ASBR | <===== ++
            +-----+
              eBGP
              ( ::: )-.
              .- ( :::::::::: )
              .- ( :: ATN/IPS ::: )-.
              ( :::: BGP Routing :::: )
              '- ( :: System :::: )-'
              '- ( :::::::::: )-'

<----- OMNI virtual link bridged by the OAL ----->

```

Figure 4: ATN/IPS ANET Architecture

When a Client connects to an ANET it specifies a nearby s-ASBR that it has selected to connect to the ATN/IPS. The login process is transparently brokered by a Proxy/Server at the border of the ANET which then conveys the connection request to the s-ASBR via adaptation layer encapsulation and forwarding across the OMNI virtual link. Each ANET border Proxy/Server is also equally capable of serving in the s-ASBR role so that a first on-link Proxy/Server can be selected as the s-ASBR while all others perform the Proxy/Server role in a hub-and-spokes arrangement. An on-link Proxy/Server is selected to serve the s-ASBR role when it receives a control message from a Client specifically requesting that service.

The Client can coordinate with a network-based s-ASBR over additional ANETs after it has already coordinated with a first-hop Proxy/Server over a first ANET. If the Client connects to multiple ANETs, the s-ASBR will register the individual ANET Proxy/Servers as conduits through which the Client can be reached. The Client then sees the s-ASBR as the "hub" in a "hub-and-spokes" arrangement with the first-hop Proxy/Servers as spokes. Selection of a network-based s-ASBR is through the discovery methods specified in relevant mobility and virtual link coordination specifications (e.g., see AERO [I-D.templin-6man-aero3] and OMNI [I-D.templin-6man-omni3]).

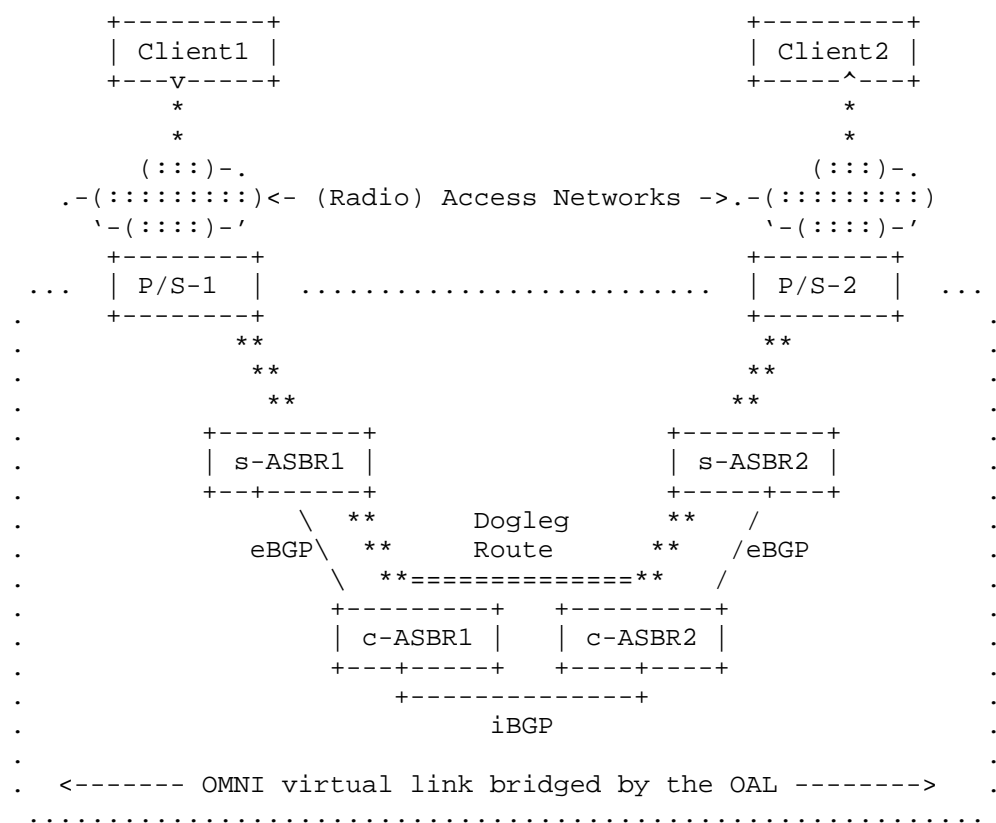
The s-ASBR represents all of its active Clients as MNP/MLA routes in the ATN/IPS BGP routing system. The s-ASBR's stub AS is therefore used only to advertise the set of MNPs/MLAs of all its active Clients to its BGP peer c-ASBRs and not to peer with other s-ASBRs (i.e., the stub AS is a logical construct and not a physical one). The s-ASBR injects the MNPs/MLAs of its active Clients and withdraws those of its departed Clients via BGP updates to c-ASBRs, which further propagate them to other c-ASBRs within the OAL AS. Since Clients are expected to remain associated with their current s-ASBR for extended periods, the level of MNP/MLA injections and withdrawals in the BGP routing system will be on the order of the numbers of network joins, leaves and s-ASBR handovers for aircraft operations (see: Section 6). It is important to observe that fine-grained events such as Client mobility and Quality of Service (QoS) signaling are coordinated only by the Client's current s-ASBRs, and do not involve other ASBRs in the routing system. In this way, intradomain routing changes within the stub AS are not propagated into the rest of the ATN/IPS BGP routing system.

5. ATN/IPS Route Optimization

ATN/IPS end systems will frequently need to communicate with correspondents associated with other s-ASBRs. In the BGP peering topology discussed in Section 3, this can initially only be accommodated by including multiple extraneous hops and/or spanning tree segments in the forwarding path. In many cases, it would be desirable to establish a "short cut" around this "dogleg" route so that packets can traverse a minimum number of forwarding hops across the OMNI virtual link. ATN/IPS end systems could therefore employ a route optimization service according to the mobility service employed (see: Section 8).

Each s-ASBR provides designated routing services for only a subset of all active Clients, and instead acts as a simple Proxy/Server for other Clients. As a designated router, the s-ASBR advertises the MNPs/MLAs of each of its active Clients into the ATN/IPS routing system and provides basic (unoptimized) forwarding services when

These route optimized paths are established through secured control plane messaging (i.e., over secured tunnels and/or using higher-layer control message authentications) but do not provide lower-layer security for the data plane. Data communications over these route optimized paths should therefore employ higher-layer security.



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+-----+
| Client1 |
+-----V-----+
*
*
( ::: )-.
.- ( :::::::::: ) <- (Radio) Access Networks -> .- ( :::::::::: )
\' ( :::::::::: )-\'
+-----+
| P/S-1 | .....
+-----V-----+
*
*=====*
+-----+
| s-ASBR1 |
+-----+
      \
eBGP  \
      \
+-----+ +-----+
| c-ASBR1 | | c-ASBR2 |
+-----+ +-----+
      +-----+
      |
iBGP
+-----+
| Client2 |
+-----^-----+
*
*
( ::: )-.
.- ( :::::::::: )
\' ( :::::::::: )-\'
+-----+
| P/S-2 | .....
+-----^-----+
*
*=====*
+-----+
| s-ASBR2 |
+-----+
      /
/eBGP  /
      /
+-----+ +-----+
| c-ASBR1 | | c-ASBR2 |
+-----+ +-----+
      +-----+
      |
iBGP
+-----+
<----- OMNI virtual link bridged by the OAL ----->

```

Figure 6: Optimized Route

6. BGP Protocol Considerations

The number of eBGP peering sessions that each c-ASBR must service is proportional to the number of s-ASBRs in its local partition. Network emulations with lightweight virtual machines have shown that a single c-ASBR can service at least 100 eBGP peerings from s-ASBRs that each advertise 10K MNP/MLA routes, i.e., $O(10^6)$ total. It is expected that robust c-ASBRs can service many more peerings than this - possibly by multiple orders of magnitude. But even assuming a conservative limit, the number of s-ASBRs could be increased by also increasing the number of c-ASBRs. Since c-ASBRs also peer with each other using iBGP, however, larger-scale c-ASBR deployments may need to employ an adjunct facility such as BGP Route Reflectors (RRs)[RFC4456].

The number of aircraft in operation at a given time worldwide is likely to be significantly less than $O(10^6)$, but we will assume this number for a worst-case analysis. Assuming a worst-case average 1 hour flight profile from gate-to-gate with 10 service region transitions per flight, the entire system will need to service at most $O(10^7)$ BGP updates per hour (2778 updates per second). This number is within the realm of the peak BGP update messaging seen in the global public Internet today [BGP2]. Assuming a BGP update message size of 100 octets (800bits), the total amount of BGP control message traffic to a single c-ASBR will be less than 2.5Mbps which is a nominal rate for modern data links.

Industry standard BGP routers provide configurable parameters with conservative default values. For example, the default hold time is 90 seconds, the default keepalive time is 1/3 of the hold time, and the default MinRouteAdvertisementInterval is 30 seconds for eBGP peers and 5 seconds for iBGP peers (see Section 10 of [RFC4271]). For the simple mobile routing system described herein, these parameters can be set to more aggressive values to support faster neighbor/link failure detection and faster routing protocol convergence times. For example, a hold time of 3 seconds and a MinRouteAdvertisementInterval of 0 seconds for both iBGP and eBGP.

Instead of adjusting BGP default time values, BGP routers can use the Bidirectional Forwarding Detection (BFD) protocol [RFC5880] to quickly detect link failures that do not result in interface state changes, BGP peer failures, and administrative state changes. BFD is important in environments where rapid response to failures is required for routing reconvergence and, hence, communications continuity.

Each c-ASBR will be using eBGP both in the ATN/IPS and the INET with the ATN/IPS unicast IP routes resolving over INET routes. Consequently, c-ASBRs and potentially s-ASBRs will need to support separate local ASes for the two BGP routing domains and routing policy or assure routes are not propagated between the two BGP routing domains. From a conceptual, operational and correctness standpoint, the implementation should provide isolation between the two BGP routing domains (e.g., separate BGP instances).

This gives rise to a BGP routing system that must accommodate large numbers of long and non-aggregable MNP prefixes and MLA addresses. The system is kept stable and scalable through the s-ASBR / c-ASBR hub-and-spokes topology which ensures that mobility-related churn is not exposed to the core.

7. Scalable Mapping

Scaling properties for the worldwide civil aviation airplane population are likely to remain within reasonable bounds for the pure BGP routing system discussed in [I-D.templin-6man-aero3] for the foreseeable future. However, the advent of unmanned air systems and all other manners of mobile nodes will soon present multiple orders of magnitude more mobility targets which may exceed the carrying capacity of a BGP-only service.

In order to support unbounded scaling, the BGP routing system can be limited to carry only the MLAs of all ASBRs on the OMNI link (and possibly also the MNPs/FNPs/MLAs of a limited number of mobile nodes) without carrying the entire population of mobile node MNP/FNP/MLA information. Each s-ASBR then registers the MNP/FNP routes and MLA addresses of its dependent mobile nodes with a scalable mapping system that can be used to resolve a target address based on longest prefix match into an s-ASBR MLA address. The global Domain Name System (DNS) ip6.arpa reverse zone can be used for this purpose.

Address resolution then becomes a two-phase operation where the address resolution request is first forwarded (e.g., via a default or more-specific route) to a c-ASBR which determines the MLA of the current s-ASBR for the address resolution target. The mapping system agent then changes the OAL destination to the discovered MLA and forwards the address resolution request to the s-ASBR which returns a fully qualified address resolution response.

By maintaining mobile node to s-ASBR mappings in a scalable ancillary lookup directory database, the BGP routing system only needs to scale to the total population of ASBRs that make up the OMNI link (plus possibly also a limited number of mobile nodes). This is likely to remain within acceptable scaling limits even for extremely large mobile node populations for the foreseeable future.

Note that the c-ASBR performs these mapping system lookups only for subject prefixes associated with the OMNI link, e.g., those covered by MSPs or any well-known FNPs. For all other subject prefixes the c-ASBR returns an address resolution response with a /64 prefix that covers the target address while including its own MLA as the resolved address.

Note also that s-ASBRs need only add or replace DNS resource records with the IP prefix mappings as they receive registration requests from new Clients. If the Client moves to a new s-ASBR, the new s-ASBR simply replaces the old resource records with fresh information.

8. Stub AS Mobile Routing Services

Stub ASes maintain intradomain routing information for mobile node clients, and are responsible for all localized mobility signaling without disturbing the BGP routing system. Clients can enlist the services of a candidate mobility service such as Mobile IPv6 (MIPv6) [RFC6275], LISP [I-D.ietf-lisp-rfc6830bis] or AERO [I-D.templin-6man-aero3] according to the service offered by the stub AS. Further details of mobile routing services are out of scope for this document.

9. Implementation Status

The BGP routing topology described in this document has been modeled in realistic network emulations showing that at least 1 million MNPs can be propagated to each c-ASBR even on lightweight virtual machines. No BGP routing protocol extensions need to be adopted.

10. IANA Considerations

This document does not introduce any IANA considerations.

11. Security Considerations

ATN/IPS ASBRs on the open Internet are susceptible to the same attack profiles as for any Internet nodes. For this reason, ASBRs should employ physical security, data link layer security [MACSEC] and/or IP securing mechanisms such as IPsec [RFC4301], etc.

ATN/IPS ASBRs present targets for Distributed Denial of Service (DDoS) attacks. This concern is no different than for any node on the open Internet, where attackers could send spoofed packets to the node at high data rates. This can be mitigated by connecting ATN/IPS ASBRs over dedicated links with no connections to the Internet and/or when ASBR connections to the Internet are only permitted through well-managed firewalls.

ATN/IPS s-ASBRs should institute rate limits to protect low data rate aviation data links from receiving DDoS packet floods.

BGP protocol message exchanges and control message exchanges used for route optimization must be secured to ensure the integrity of the system-wide routing information base. Security is based on IP layer security associations between peers which ensure confidentiality, integrity and authentication over secured tunnels (see above). Higher layer security protection such as TCP-AO [RFC5926] is therefore optional, since it would be redundant with the security provided at lower layers.

Data communications over route optimized paths should employ end-to-end higher-layer security since only the control plane and unoptimized paths are protected by lower-layer security. End-to-end higher-layer security mechanisms include QUIC-TLS [RFC9001], TLS [RFC8446], DTLS [RFC6347], SSH [RFC4251], etc. applied in a manner outside the scope of this document.

This document does not include any new specific requirements for mitigation of DDoS.

11.1. Public Key Infrastructure (PKI) Considerations

In development of the overall ATN/IPS operational concept, ICAO addressed the security concerns in multiple ways to ensure coordination and consistency across the various groups. This also avoided potential duplicative work. Technical provisions related specifically to the operation of ATN/IPS are specified in supporting ATN/IPS standards. However, other considerations such as the establishment of a PKI, were determined to have an impact beyond ATN/IPS. ICAO created a Trust Framework Study Group (TFSG) to define various governance, policy, procedures and overall technical performance requirements for system connectivity and interoperability.

As part of their charter, the TSFG is specifically developing a concept of operations for a common aviation digital trust framework and principles to facilitate an interoperable secure, cyber resilient and seamless exchange of information in a digitally connected environment. They are also developing governance principles, policy, procedures and requirements for establishing digital identity for a global trust framework that will consider any exchange of information among users of the aviation ecosystem, and to promote these concepts with all relevant stakeholders.

ATN/IPS will take advantage of the developments of TFSG within the overall ATN/IPS operational concept. As such, this will include the usage of the PKI specification resulting from the TFSG.

12. Acknowledgements

This work is aligned with the FAA as per the SE2025 contract number DTFAWA-15-D-00030.

This work is aligned with the NASA Safe Autonomous Systems Operation (SASO) program under NASA contract number NNA16BD84C.

This work is aligned with the Boeing Commercial Airplanes (BCA) Internet of Things (IoT) and autonomy programs.

This work is aligned with the Boeing Information Technology (BIT) MobileNet program.

The following individuals contributed insights that have improved the document: Ahmad Amin, Mach Chen, Russ Housley, Erik Kline, Hubert Kuenig, Tony Li, Gyan Mishra, Alexandre Petrescu, Dave Thaler, Pascal Thubert, Michael Tuxen, Eric Vyncke, Tony Whyman. Vaughn Maiolla is further remembered for his support and guidance.

Honoring life, liberty and the pursuit of happiness.

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Appendix A. BGP Convergence Considerations

Experimental evidence has shown that BGP convergence time required after an MNP is asserted at a new location or withdrawn from an old location can be several hundred milliseconds even under optimal AS peering arrangements. This means that packets in flight destined to an MNP route that has recently been changed can be (mis)delivered to an old s-ASBR after a Client has moved to a new s-ASBR.

To address this issue, the old s-ASBR can maintain temporary state for a "departed" Client that includes an OAL address for the new s-ASBR. The OAL address never changes since ASBRs are fixed infrastructure elements that never move. Hence, packets arriving at the old s-ASBR can be forwarded to the new s-ASBR while the BGP routing system is still undergoing reconvergence. Therefore, as long as the Client associates with the new s-ASBR before it departs from the old s-ASBR (while informing the old s-ASBR of its new location) packets in flight during the BGP reconvergence window are accommodated without loss.

Appendix B. AERO/OMNI Spanning Tree Properties

The AERO/OMNI services establish an adaptation layer for the OSI reference model appearing as "the layer below the network layer but above the data link layer". The adaptation layer presents a virtual bridging service from the network layer's perspective and a BGP-based IPv6 routing service from the data link layer's perspective.

AERO/OMNI overlay networks include s-ASBRs (aka Proxy/Servers) and c-ASBRs (aka Gateways) as the vertices in a graph with a typically sparse collection of edges between pairwise vertices. The graph

minimally comprises a true spanning tree connecting all vertices but may also include additional edges in the spirit of IEEE 802.1aq Shortest Path Bridging. BGP routing then ensures loop freedom even if this "augmented" spanning tree includes cycles.

Each edge in the spanning tree corresponds to one or more connecting links which may be physical (e.g., fiber/free-space optics, etc.) or virtual (an IP tunnel over an underlying Internetwork). The OMNI interface provides a nexus for link selection, where control messages between vertices must be forwarded over edge links that are secured at the network layer or below while data messages may be forwarded over unsecured links. AERO/OMNI refer to these distinct forwarding contexts as the "secured spanning tree" and "unsecured spanning tree", respectively.

AERO route optimization provides an essential service to establish shortcut paths in the data plane that do not necessarily follow strict spanning tree paths. The shortcuts are formed through secured spanning tree control message exchanges which establish shortest path forwarding state in Clients, Proxy/Servers and intermediate Gateways. Route optimization therefore reduces spanning tree traffic concentration and instead distributes traffic over a diversity of underlying Internetwork paths.

Appendix C. Change Log

<< RFC Editor - remove prior to publication >>

Differences from -27 to -28:

- * Introduced MLAs and explained their role in the architecture.
- * Included reference to MLA spec.

Differences from -26 to -27:

- * Update references.

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