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Network File System (NFS) Version 4 Minor Version 1 Protocol  
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## Abstract

This document describes the Network File System (NFS) version 4 minor version 1, including features retained from the base protocol (NFS version 4 minor version 0, which is specified in RFC 7530) and protocol extensions made and part of Minor Version 1. The later minor version has no dependencies on NFS version 4 minor version 0, and was, until recently, documented as a completely separate protocol.

This document is part of a set of documents which collectively obsolete RFCs 8881 and 8434. In addition to many corrections and clarifications, it will rely on NFSv4-wide documents to substantially revise the treatment of protocol extension, internationalization, and security, superseding the descriptions of those aspects of the protocol appearing in RFCs 5661 and 8881.

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## 1. Introduction to this Update

This document is intended to be the basis for a revised and updated specification of NFSv4.1. Unlike [RFC8881], which provided a limited-function update to [RFC5661], this document has a broader mandate and will do the following:

- \* Revise any text known to be wrong or otherwise inappropriate. Such text will not be retained as it has been merely because it is outside a limited pre-specified change scope.

This includes changes in some errata reports with the status REJECTED, where there is a Working Group Consensus that change is necessary.

- \* Correct protocol defects, that, by their nature, can and should be addressed via a limited use of the extension mechanism described in Section 9 of [RFC8178].

For more discussion of the correction of existing protocol defects, see Section 1.4. This discussion, which covers protocol defects addressed in NFSv4-wide documents in addition to this document, will focus on how compatibility issues are addressed when defects are corrected. It will also pay attention to the question of when XDR extensions in a bis document are necessary as part of providing features already included (albeit incorrectly) in the current minor version as opposed to deferring such extensions to later minor versions.

- \* Incorporate the necessary changes to the description of NFSv4.1 that were published in documents marked as updating [RFC8881], in this case [RFC8434].



This incorporation involved a major advance in our understanding of the structure and benefits of pNFS and the possibilities of its further extension. See Section 1.5.

### 1.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as specified in BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

The above differs from the corresponding statement in earlier specification versions ([RFC5661] and [RFC8881]) which only referred to [RFC2119]. For further discussion of this change, see Appendix B.3

In some cases, such keywords will appear in all capitals within quotations, direct or indirect, from earlier documents. In such cases, these terms are to be interpreted just as above, except where it is explicitly noted that such an interpretation is not to be inferred. In some cases, it might be that this document's approach to the matter would not use those key words for reasons explained in the text. Such a shift might cause compatibility issues, if the previous keyword were actually relied upon but it also possible that it was not relied upon while the implications of that use were ignored for various reasons.

The reader should be aware that, as discussed in Appendix D.1, there are uses of the keywords listed above in RFCs 5661 and 8881 which might not have been appropriate, even though the interpretation specified above was intended when the text was written and submitted for publication. In some cases, the text in this document has been updated to correct the issue but it should be understood that not all such questionable uses have been addressed and that this state of affairs might continue to exist until a later draft of this document is submitted for publication.

It should not be assumed that all statements that do not include these keywords are inherently non-normative and can be safely ignored as constituting implementation advice. In some cases, where guidelines are given for the specification of future documents, the intent is normative, as that word is normally used, outside the IETF. An important example is the specification of layout types, as described in Section 19. This is important because many aspects of the pNFS feature depend on the specific layout type(s) implemented, essentially making pNFS an extensible feature suite. For detailed discussion, see Section 18.3.

## 1.2. The Changed Role of this Specification

Previous specifications for this minor version ([RFC5661], [RFC8881]) have purported to describe the protocol in its entirety, without reference to features common to all minor versions of NFS Version 4. In contrast, this update relies on a set of base documents describing common aspects of the NFSv4 protocol that applies to all minor versions.

- \* Rules for extensions and creation of new minor versions appear only in [RFC8178], unlike previously in which they appeared in the NFSv4.1 specification. This eliminates the unfortunate situation in which each minor version was allowed to create its own extension rules.
- \* Handling of internationalization-related matters (for all minor versions) is now discussed in its own document, which is expected to be an RFC derived from [I-D.ietf-nfsv4-internationalization].

That document, based in large part on the handling of internationalization for NFSv4 minor version zero outlined in [RFC7530], has been extended to cover all minor versions and enhanced to fully support case-insensitive handling of internationalized file names.

This corrects the unfortunate situation in which internationalization for minor version one and subsequent minor versions (in [RFC5661] and [RFC8881]) had never been implemented and could never have been implemented by NFS Version four clients and servers.

- \* Handling of core security-related matters for all NFSv4 minor versions will be consolidated in a set of documents that are expected to be RFCs derived from [I-D.dnoveck-nfsv4-security] and [I-D.ietf-nfsv4-acls-update].

This shift is made necessary by the following issues, many of which are of long standing and make the continuation of previous approaches to these issues insupportable:

- The lack of a substantial threat analysis in the Security Considerations section of any existing minor version specification.

- The unfortunate designation of AUTH\_SYS as "OPTIONAL means of authentication" which had the effect of obscuring the severe security problems with its common use. The use of "OPTIONAL" suggested that use of AUTH\_SYS had no harmful consequences while the phrase "means of authentication" ignored the fact that no actual authentication took place.
- The assumption that data confidentiality could be satisfactorily addressed as an occasionally-used optional facility.
- The neglect of the need for dependable semantic description of the protocol's authorization semantics.

Earlier versions of NFS had avoided the need for such descriptions by relying on POSIX semantics. The addition of non-POSIX semantic elements, including named attributes and ACLs, interfered with that approach although the necessity was not recognized as NFSv4 was being formulated.

The availability of new transport-level security features such as those provided by RPC-with-TLS [RFC9289] provides a basis to correct many of the above issues.

In providing that sort of correction, we need to be careful not to declare existing implementations non-compliant post facto, while still providing adequate warning of the security consequences of continuing to use the NFS Version 4 protocol insecurely, as described in previous specifications.

### 1.3. Possibility of Compatibility Issues

Because this document or other documents that are part of this specification update might make changes to matters previously discussed in [RFC8881], the possibility of compatibility issues cannot be excluded.

However, the likelihood of such issues arising is expected to be low because:

- \* In many cases, descriptions within [RFC8881] differed from the existing implementations and change was necessary to avoid directing implementers to create implementations that could not interoperate with existing implementations.
- \* There are cases in which attributes were specified to be OPTIONAL even though there were no good reasons to implement a server that did not otherwise provide support for them.

- \* In some cases, it was necessary to change the definition of OPTIONAL features that had never been implemented, often because problems in the existing specification made the features impossible to implement as defined. In such cases, the absence of implementations makes it impossible for compatibility problems to arise.
- \* In some cases, problems in the protocol as previously are being corrected as described in Section 9 of [RFC8178], by creating OPTIONAL protocol extensions. In such cases, implementations unaware of the extensions will function as they did before, while others have the option of using the extensions when interacting with an implementation supporting the extension.

Specific updates are dealt with in the following subsections:

- \* Updates in RFCTBD10 are addressed in Section 1.3.1.
- \* Updates in new NFSv4-wide documents are discussed in Section 1.3.2.
- \* Updates in RFCTBD30 are addressed in Section 1.3.3.

#### 1.3.1. Compatibility issues for RFCTBD10

This document introduces a number of potential incompatibilities that either cannot be realized or could only be realized in situations highly unlikely to exist. For discussion of the reasons such incompatibilities are highly unlikely to be problematic, see the rest of this section which also explains the needs for the changes.

One important example concerns the conversion of the attributes mode, owner, and owner\_group from OPTIONAL to REQUIRED. Since the vast majority of servers support these attributes and the vast majority of clients use them, any incompatibilities would be very unlikely to exist.

If there did exist servers not supporting any of these attributes and clients capable of interacting with them, this unexpected situation would not result in any incompatibility. The only change would be that the client and the server would become non-compliant when they previously had been considered compliant. This change of status has no consequences for client-server compatibility.

Overall, the need for this change illustrates previous specifications' lack of serious attention to security issues, which it is part of the job of the NFsv4.1 respecification effort to correct.

Another important class of potential compatibility issues in which we had to change the specification derives from situations in which the previous text designate normal, often unavoidable, behavior as non-compliant:

- \* Previously, clients were required to wait forever for an RPC response while many (e.g. Linux) clients stopped waiting in the event of a control-C. This common behavior, formerly non-compliant, is now allowed but there is no incompatibility resulting from this change even though much existing behavior has gone from non-compliant to compliant.
- \* Previously the use of RoCE for NFS/RDMA had been disallowed (formally) but is now allowed. Even so, there are clients and servers that do this and work together doing so and making them compliant raises no compatibility issues.
- \* Previously retries were not allowed ("MUST NOT" was used) unless there was a connection drop. This was changed to a recommendation since there us no way this could be a "fundamental requirement of the specification" given that the reply cache was designed to avoid any negative consequences of retries. In any case there was no possibility of an incompatibility since the effect of the changes is only to make previously non-compliant clients compliant.

There are also other changes made to make specification text match the actual implementations. Because care was used in making sure that the change reflected implementation changes documented in errata reports, we do not expect compatibility issues to arise. Some important examples concern changes to pNFS-related operations that have been documented in errata report and discussed within the group so that existing implementations all follow the changed approach. These include the following:

- \* In LAYOUTCOMMIT4args, loca\_offset and loca\_length are present but not used. Previously, they had been used to effect partial layout commits, which are no longer provided for outside of explicitly layout-type-specific features.

- \* In struct `notify_deviceid_change4`, `ndc_immediate` is present but not used. Previously, this field had been used to indicate whether or not a device change notification was being done immediately rather than being delayed.

A number of new features that are highly unlikely to have been implemented have been respecified so that they are effectively implementable. Although the Working Group's has limits on its ability to investigate the matter, in each case, it appears very unlikely that implementations exist. See below for details. In the absence of such implementations, we can expect new implementations to follow this specification, avoiding compatibility issues.

- \* The persistent reply cache feature was added to NFSv4.1 in [RFC5661] but there is good reason to believe it has never been implemented based on that document or on [RFC8881]. Although an implementer would have no obligation to inform the Working Group of the existence of any implementation, the nature of the current description makes it very unlikely that any such implementations exist.

For reasons discussed in Section 8.1, the feature was essentially implementable as described, while it is unlikely that any attempting an approximate implementation would do so without consulting the Working Group.

As a result, there is no real possibility of compatibility issues arising.

- \* The directory delegation feature was added to NFSv4.1 in [RFC5661] but there has been no discussion of implementations even though the performance effects of the feature would be of interest to the Working Group, since this feature was included in NFSv4.1 because of performance concerns.

As things turned out, there were discussions of NFSv4 performance on metadata-intensive workloads at Working Group meetings with the effect of directory delegations not being commented on leading to a strong belief that no such implementations existed.

When the reasons for the presumed non-implementation were analyzed, changes were made to the specification of this feature. While the probability of implementations existing is non-zero, their possible existence will not give rise to compatibility issues due to the way in which changes were made.

All of the XDR changes took the form of extensions to correct defects as discussed in Section 9 of [RFC8178]. In addition the added requirement for delegation recall on certain changes to the directory will not lead to incompatibility since the server is free to recall delegations as it chooses. In the case of possible existing servers that fail to recall the delegation in such cases, they will become non-compliant. However, this fact is not troublesome as we will be going from a compliant server that does not work correctly to a non-compliant one, with the non-compliance being an important cue to fix the server.

### 1.3.2. Compatibility issues for NFSv4-wide Documents

For these three NFSv4-wide documents, a major implementation area was respecified because of major problem in the existing NFSv4.1 descriptions of these areas. Because the replacement text was intended to reflect the actual implementation, the likelihood of incompatibilities is quite low:

- \* RFCTBD20 respecified internationalization in a fashion quite different from the way it is specified in [RFC5661] and [RFC8881].

This major shift did not give rise to compatibility issues because the approach to internationalization presented in [RFC5661] and [RFC8881] had never been implemented. We can be sure of this because any attempt to implement this stringprep-based approach would have generated comments/questions due to the complexity involved in, for example, rejecting unassigned code points.

Existing implementations were compatible with the approach discussed in [RFC7530] and the treatment in RFCTBD20 maintained that compatibility.

- \* RFCTBD21 respecified security in a fashion markedly different in tone from the way it had been discussed in previous specifications. This included a different approach to the use of AUTH\_SYS, more concern about issues related to the security of data in flight, and the inclusion of a discussion of authorization semantics.

Despite these major differences, care has been taken to avoid the creation of incompatibilities. The new treatment avoided incompatibilities by avoiding any disallowing of previously allowed behavior, even in cases in which there might be grounds for doing so.

For example, although the use of AUTH\_SYS in the clear has enough associated security issues that saying it MUST NOT or SHOULD NOT be used, this cannot be done for existing minor versions. To deal with this situation, its use is not disallowed or recommended against. However, unlike in previous specifications, implementers are given appropriate notice of the security issues resulting from use of AUTH\_SYS in the clear.

The addition of discussion of RPC-with-TLS is another important difference which cannot cause incompatibility issues. Despite the optionality of the use of this facility, it is important to discuss in connection with NFSv4 security and its handling of the security of data in flight.

- \* The new treatment of ACLs in RFCTBD22, while extensive, does not give rise to major incompatibilities because of the way we were forced to adapt to the previous approach while using a clearer explanatory framework that made the POSIX-oriented subset the basis of the description and explicitly making the Windows-ACL-based extensions OPTIONAL.

For the most part, this approach does not generate incompatibilities in that it moves formerly allowed behaviors (implicitly optional) to be formally OPTIONAL, without changing the set of allowed behaviors. However, with the addition of the OPTIONAL attribute Aclchoice, the client is able to determine what extensions to the base ACL mode have been implemented.

The one possible formal incompatibility arises from the three ACE mask bits corresponding to the POSIX R, W, and X bits. support for these are now REQUIRED whereas previously it had been suggested that they were effectively optional along with all the other ACE mask bits. The chance of this being a problem is low since, we have not heard any discussion of server without support for all of these bits and the difficulty of adopting to POSIX semantics, of such implementations existed would certainly have led to Working Group discussion of the issue.

### 1.3.3. Compatibility issues for RFCTBD30

These changes can be divided into three groups:

- \* In many cases, the XDR changes consist of changes in the typedef naming structure, with no change to the structure of messages described by the XDR. As a result, compatibility issues are not expected to arise.



- \* Many changes cannot create incompatibilities because the change consists only of changed names of typedefs. Such changes were necessary to clarify how naming interpretation rules applied to various strings. Because all such strings have types that evaluate to opaque<>;, there can be no incompatibilities generated.
- \* In some cases, fields defined in various messages have become unused. Space is still reserved for these fields but the on-use is indicated in comments, so there is no XDR change to generate an incompatibility.
- \* All substantive XDR changes take the form of XDR extensions as provided for by Section 9 of [RFC8178]. As a result, implementations unaware of the change will function as they did previously, while those that are aware of extension have the option of adapting to the support or non-support by the peer implementation.

#### 1.4. Addressing Protocol Defects Via XDR Changes

This section provides an overview of situations in which it was necessary to change the protocol XDR to correct specification issues that needed to be addressed as part of the respecification effort. Although all these issues can, from today's perspective, be viewed as arising from mistakes, it is not always clear whether that description is appropriate given the knowledge available when the decisions were made. In any case, the fact that we would have done something different if the decision were made today does not imply that the bis document should adopt that approach. The following considerations will often determine if we choose to defer such work to a later minor version. This is sometimes necessary despite the fact that the extension technique within [RFC8178] would allow the same sort of work to be done in a bis document rather than an in extension.

- \* The need to avoid adding new features that are substantial, since they would need a period of review and refinement more easily done in an extension document.

In general, we would try to avoid adding new attributes or operations in the bis document to provide additional functionality not considered during the original protocol development.

- \* In features structured to be extensible, the XDR is often enhanced easily to expand the functionality available.

- \* If a new attribute is necessary to make a feature usable, that will make necessary to specify as part of the bis document, rather than live with a severely compromised feature in a minor version being respecified.

The correction of protocol defects often gives rise to compatibility issues and their possible presence will be discussed below. In addition, the question of when it is appropriate to address such issues using the protocol extension mechanism described in [RFC8178] needs to be considered. Section 9 of that document alludes to this possibility but we have to decide when defects are best addressed in that way.

These defects can be divided into two groups based on their origin.

- \* Defects that originated in minor version zero

Many of these defects are addressed in the new NFSv4-wide documents ([I-D.dnoveck-nfsv4-security], [I-D.ietf-nfsv4-acls-update], and [I-D.ietf-nfsv4-internationalization]). For these defects, the greater change scope require more attention to compatibility issues. In addition, that greater scope limits the degree to which protocol extension can be used in providing a correction since that extension would need to be propagated to two non-extensible minor versions.

- \* Defects that originated in minor version one.

A major source of defects was the result of the addition of a set of OPTIONAL features in v4.1, that have never been implemented, making it important to eliminate issues in the specification that have led to this situation.

The presumptive non-implementation of these features will limit interoperability concerns. However, since we cannot be sure about the possible existence of implementations under development, we will try to provide for the possibility of interoperating with earlier implementations, even if that interoperation is hypothetical. Only in the case of features whose current specification makes implementation impossible can we ignore the possibility of interoperating with such implementations.

The following defects were addressed as part of this update effort:

- 1: Internationalization is being thoroughly respecified in the NFSv4-wide document [I-D.ietf-nfsv4-internationalization].

There were no NFSv4.1 compatibility issues to deal with since the handling of internationalization approach mandated by [RFC8881] had never been implemented.

Potential NFSv4.0 compatibility issues were very limited since the approach followed in [RFC7530] was continued and that approach had been used by all implementations.

In one particular case, there is a potential compatibility issue arising from the transition of one potential troubling server behavior from being discouraged using SHOULD NOT to being prohibited using MUST NOT. However, in view of the unlikelihood of ongoing use of the discouraged behavior, this has not been considered problematic.

The respecification of the fs\_charset\_cap attribute raises the possibility of within-NFS4.1 compatibility issues. However the very limited use of this attribute by clients combine with the lack of clarity in the previous definition makes it unlikely that the use of protocol extension to support previous uses would be justified.

- 2: Because the Persistent reply cache feature could not be implemented as described in [RFC8881], the entire area was respecified in Section 8.

The reasons for this respecification are discussed in Section 8.1.

Because the existing feature specification was unimplementable there were no compatibility issues to deal with.

No protocol extensions were needed since the small set of bits defined for the earlier feature could be coopted.

- 3: Security for all minor versions is being thoroughly respecified in the NFSv4-wide document [I-D.dnoveck-nfsv4-security]. In this discussion, issues related to authorization semantics and to ACLs are being dealt with separately.

This respecification was made necessary by the lack of threat analyses for all minor versions, the absence of any discussion of the security problems associated with the use of AUTH\_SYS, and the half-hearted approach to the security of over-the-wire transmission in which transmission in the clear was the default and the provision of secure transmission was an option requiring per-fs configuration.

- 4: As part of the new handling of security, a more serious treatment of authorization semantics was necessary. As part of effecting this, the attributes mode, owner, and owner\_group became REQUIRED, as it is impossible to effect security without them.

This change was not connected to the shift in terminology in which the attributes incorrectly described as "RECOMMENDED" became "OPTIONAL".

No significant compatibility issue are expected, since the existence of servers not supporting these attributes or of clients interacting with such servers, while possible theoretically, has to be considered extremely unlikely and none are known to the working group.

- 5: A number of gaps in the description of authorization semantics needed to be addressed. These include the lack of a clear description of authorization for operations on named attribute directories and potential use of the "sticky" bit in controlling authorization of file deletion.

These matters are being addressed within [I-D.dnoveck-nfsv4-security] where they are being tracked as Consensus Items #66 and #6 respectively.

Working group consideration of the security document will involve resolving those two Consensus Items, as well as others.

- 6: The handling of ACLs for all minor versions is being thoroughly respecified in the NFSv4-wide document [I-D.ietf-nfsv4-acls-update].

Such a respecification was made necessary by the profound underspecification of the ACL feature that arose from a misguided attempt to support two very different approaches to the provision of ACLs. The problems posed by the different semantics of these two were never clearly addressed since it was erroneously assumed that semantic description could be avoided. As a result, potential interoperability was compromised since there was no way for the client to determine what ACL-based facilities were supported by a particular server, given that the specification treated these differences as if they were quality-of-implementation issues.

The development of [I-D.ietf-nfsv4-acls-update] has included a respecification of the area in which support for a subset of draft POSIX ACLs, termed UNIX ACLs, was the core and the various

additions to that core were considered additional OPTIONAL features. These included the features that motivated the extensions in the NFSv4 ACL model.

The treatment within [I-D.ietf-nfsv4-acls-update] now makes it clear that the current ACL model is only one of a set of ACL formats, implying that future minor versions could define new ones.

The development of [I-D.ietf-nfsv4-acls-update] has involved use of protocol extension within NFSv4.1 in addition to necessary structural changes that did not involve XDR changes.

The development of [I-D.ietf-nfsv4-acls-update] to support the rfc8881bis effort will most likely be limited to providing interoperability for those using the facilities within the UNIX ACL core or within the draft POSIX acl model. Interoperability for features beyond that set is likely to be delayed to later ACL bis, while the deletion of unneeded proposed features will have to wait for a later minor version, e.g., NFSv4.3.

- 7: It has been necessary to define a new read-only per-fs OPTIONAL attribute that will allow clients to determine which of the OPTIONAL extensions to the core UNIX ACL model are supported by the server.

While this was essential to make the NFSv4 extensions usable, it also has a critical role in making POSIX ACL support available within NFSv4, albeit with some client mapping/filtering,

- 8: The defects described in Appendix D.2.1 needed to be addressed together, in connection with making it clear that the term "Exactly-once Semantics" ignored the fact that there were valid reasons to give up on requests which could leave them unexecuted.

This change did not give rise to compatibility issues since the specification was changed to match existing implementations, and these are expected to remain as they are.

- 9: The inadvertent prohibition of the use of RoCE in implementing NFSv4.1 using RPC-over-RDMA was removed.

This is another case in which compatibility issue are not expected because the spec has been changed to match existing implementations.

- 10: The corrections discussed in Appendix D.2.3 had to be made since most of the worries expressed within it were the result of misunderstandings.

Although no compatibility issues are expected we will need to review the changes and reach consensus on them.

- 11: There were a number of issues in the earlier specification of the directory delegation feature that need to be addressed to enable implementations of this needed feature to be produced. Given the lack of implementation during the long period since they were introduced in [RFC5661] many years ago.

While a significant part of the problems could be ascribed to clarity issues, there were also a set of defects, some of which required protocol extensions, as provided for in Section 9 of [RFC8178].

The defects which contributed substantially to this long-lasting lack of implementation includes the failure to fully address authorization issues for the use of cached directory data, implementability issues regarding the maintenance of cached attribute data, and the assumption that clients could maintain the cached directory contents only in the same format as used by the server. For more discussion of these defects, see Appendix D.2.2.

These issues were addressed by a major rewrite of Section 16.2 in which protocol extension was necessary, including the addition of new values to the enum `notify_type4`. In addition, there are complementary changes made to Section 25.39 and Section 27.4) and to operations that might result in notifications being sent.

Although no client and server implementations of this feature are known to exist, the possibility of them existing cannot be excluded. As a result, the revised specification takes care to deal appropriately with such hypothetical implementations, and to not prohibit their use unless that is necessary to avoid unacceptable system behavior.

- 12: Change in the recommendations regarding handling of numeric strings to represent users and groups.

Formerly considered troublesome even in the AUTH\_SYS case despite the fact that there is no explanation given as to how to effect mapping between numeric ids and strings. Instead, it is assumed that client and server will somehow agree to do this without the specification making it possible or giving a convincing reason that such mapping is needed.

Of the above, only the items 7, 8, 9, 10, and 14 required protocol extension to resolve. All will to be incorporated in RFCTBD30.

#### 1.5. Incorporation of RFC8434

The incorporation of [RFC8434], originally written to clarify the requirements for new layout types, resulted in a deep and significant improvement in our understanding of pNFS-related features of the protocol and how various layout types deal with similar issues.

As a consequence, the original description of pNFS, formerly in Sections 12 and 13 of [RFC8881] has been restructured as three top-level sections in the revised NFsv4.1 specification:

- \* Section 18 is based on Section 12 of [RFC8881] but the material has undergone considerable revision to fit in the new three-section organization.

Most previous references to various layout types have been eliminated except for a few clearly marked as examples. As a result, this section now makes clear that pNFS is a framework making deliberate allowance for per-layout-type differences in a number of areas.

- \* Section 19 is based on [RFC8434] but the requirements have been expanded to make clearer the division between the aspects common to all layout types (the subject of Section 18) and the choices that individual layout types might make to address the gaps deliberately left by the treatment in Section 18.
- \* Section 20 is based on Section 13 of [RFC8881] but the material has been adapted to meet the requirements of Section 19.

As part of these changes it was necessary to fill in gaps left previously regarding the validity of layout revocation and how conflicts arising as a result of server-multipathing could be addressed by layout/recall and with layouts re-established once the two targets were again in sync. For details, see Section 20.1.

## 2. Introduction to this Minor Version Specification

### 2.1. The NFS Version 4 Minor Version 1 Protocol

The NFS version 4 minor version 1 (NFSv4.1) protocol is the second minor version of the NFS version 4 (NFSv4) protocol. The first minor version, NFSv4.0, is now described in [RFC7530], as modified by [RFC7931] and [RFC8587]. Minor version 1 follows the guidelines for minor versioning presented in [RFC8178].

As a minor version, NFSv4.1 is consistent with the overall goals for NFSv4, but extends the protocol so as to better meet those goals, based on experiences with NFSv4.0. In addition, NFSv4.1 has adopted some additional goals, which motivate some of the major extensions in NFSv4.1, such as the use of the sessions model.

This minor version adds a considerable number of new operations including some that are not OPTIONAL and makes a number of NFSv4.0 operations MANDATORY to NOT implement. As a result, the vast majority of NFSv4.0 requests are not valid in NFSv4.1 and vice versa. While clients and servers that support both minor versions are common, such implementations treat the two versions as distinct protocols sharing a substantial common heritage.

### 2.2. Scope of This Document

This document describes the NFSv4.1 protocol. With respect to NFSv4.0, this document does not:

- \* Describe the NFSv4.0 protocol, except where needed to contrast it with NFSv4.1.
- \* Modify the specification of the NFSv4.0 protocol.
- \* Clarify the NFSv4.0 protocol.

### 2.3. NFSv4 Goals

The NFSv4 protocol is a further revision of the NFS protocol defined already by NFSv3 [RFC1813]. It retains the essential characteristics of previous versions: easy recovery; independence of transport protocols, operating systems, and file systems; simplicity; and good performance. NFSv4 had the following goals:

- \* Improved access and good performance on the Internet.



The protocol is designed to transit firewalls easily, perform well where latency is high and bandwidth is low, and scale to very large numbers of clients per server.

- \* Strong security with facilities for negotiation of security handling built into the protocol.

The protocol has been built on the work of the ONCRPC working group by supporting the RPCSEC\_GSS protocol. Additionally, the NFSv4.1 protocol provides a mechanism to allow clients and servers the ability to negotiate security and has provisions requiring or recommending client and server support for a minimal set of security schemes.

The protocol now takes advantage of the ability of RPC to make confidentiality available by using TLS-based encryption on connections to be used for NFSv4.1, which may limit the need for negotiation regarding facilities such as privacy.

- \* Good cross-platform interoperability.

The protocol embraces a file system model that provides a useful, common set of features that does not unduly favor one file system or operating system over another.

- \* Designed for protocol extensions via minor versioning.

The protocol is designed to accept standard extensions within a framework that enables and encourages backward compatibility.

When extensions are OPTIONAL, they can be added to an existing extensible minor version.

## 2.4. NFSv4.1 Goals

NFSv4.1 has the following goals, within the framework established by the overall NFSv4 goals.

- \* To correct significant structural weaknesses and oversights discovered in the base protocol.
- \* To add clarity and specificity to areas left unaddressed or not addressed in sufficient detail in the base protocol. However, as stated in Section 2.2, it is not a goal to clarify the NFSv4.0 protocol in the NFSv4.1 specification.
- \* To add specific features based on experience with the existing protocol and recent industry developments.

- \* To provide protocol support to take advantage of clustered server deployments including the ability to provide scalable parallel access to sets of files distributed among multiple servers, using a range of data access protocols.

This parallel access may involve striping of single files among a set of servers within the cluster.

Alternatively, parallel access may be provided by distributing unstriped files within the cluster allowing each client to contact the server holding each particular file directly.

## 2.5. General Definitions

The following definitions provide an appropriate context for the reader.

Byte: In this document, a byte is an octet, i.e., a datum exactly eight bits in length.

Client: The client is the entity that accesses the NFS server's resources. The client may be an application that contains the logic to access the NFS server directly. The client may also be a traditional operating system client that provides remote file system services for a set of applications.

A client is uniquely identified by a client owner.

With reference to byte-range locking, the client is also the entity that maintains a set of locks on behalf of one or more applications. This client is responsible for providing crash or failure recovery for those locks it manages, in order to deal with the possibility of server reboot.

Note that multiple clients may share the same transport and connection and multiple clients may exist on the same network node.

Client ID: The client ID is a 64-bit quantity used as a unique, short-hand reference to a client-supplied client owner, consisting of a verifier and a client owner id. The server is responsible for supplying the client ID.

Client Owner: The client owner includes a unique string, the client owner id, opaque to the server, that identifies a client. It also includes a verifier to enable successive instances of the same client to be distinguished. Multiple network connections and source network addresses originating from those connections may

share a client owner. The server is expected to treat requests from connections with the same client owner as coming from the same client. See Section 5.5 for more detail.

**File System:** A file system is the collection of objects accessed using a particular server (as identified by the major identifier of a server owner, which is defined later in this section) that share the same value of the fsid attribute (see Section 11.12.1.9).

**Lease:** A lease is an interval of time defined by the server for which the client is granted locks. At the end of a lease period, unrecallable locks may be revoked if the lease has not been extended. Such a lock must be revoked if a conflicting lock has been granted after the lease interval. Revocation of unrecallable locks within the lease interval is expected to be an unusual event and clients normally expect such revocations to be rare.

A server grants a client a single lease for all of its associated locking state. Recallable locks such as layouts and delegations can be revoked within the lease period and are generally not affected by the state of the lease.

**Lock:** The term "lock" is used to refer to byte-range (in UNIX environments, also known as record) locks, share reservations, delegations, or layouts unless specifically stated otherwise.

**Secret State Verifier (SSV):** The SSV is a unique secret key shared between a client and server. The SSV serves as the secret key for an internal (that is, internal to NFSv4.1) Generic Security Services (GSS) mechanism (the SSV GSS mechanism; see Section 7.9). The SSV GSS mechanism uses the SSV to compute message integrity code (MIC) and Wrap tokens. See Section 7.8.3 for more details on how NFSv4.1 uses the SSV and the SSV GSS mechanism.

**Server:** The Server is the entity responsible for coordinating client access to a set of file systems and is identified by a server owner. A server can span multiple network addresses.

**Server Owner:** The server owner identifies the server to the client. The server owner consists of a major identifier and a minor identifier. When the client has two connections each to a peer with the same major identifier, the client assumes that both peers are the same server (the server namespace is the same via each connection) and that lock state is sharable across both connections. When each peer has both the same major and minor identifiers, the client assumes that each connection might be associated with the same session.

**Stable Storage:** Stable storage is storage from which data stored by an NFSv4.1 server can be recovered without data loss from multiple power failures (including cascading power failures, that is, several power failures in quick succession), operating system failures, and/or hardware failure of components other than the storage medium itself (such as disk, nonvolatile RAM, flash memory, etc.).

Some examples of stable storage that are allowable for an NFS server include:

1. Media commit of data; that is, the modified data has been successfully written to the disk media, for example, the disk platter.
2. An immediate reply disk drive with battery-backed, on-drive intermediate storage or uninterruptible power system (UPS).
3. Server commit of data with battery-backed intermediate storage and recovery software.
4. Cache commit with uninterruptible power system (UPS) and recovery software.

**Stateid:** A stateid is a 128-bit quantity returned by a server that uniquely defines the open and locking states provided by the server for a specific open-owner or lock-owner/open-owner pair for a specific file and type of lock.

**Verifier:** A verifier is a 64-bit quantity that is changed to indicate that a corresponding change on one of the peers has occurred requiring the other peer to adjust to possibility of change. There are a number of 64-bit quantities identified as verifiers:

- \* In many cases, a verifier is a 64-bit quantity generated by the server that the client can use to determine if the client has restarted and potentially lost writes that had not been reliably committed to stable storage.
- \* Another type of verifier is used to indicate whether the server mapping between directory cookies and directory entries has changed, requiring the client to restart directory interrogations that normally may be continued across multiple REaddir requests.
- \* As described above, client owners include a verifier, allowing the server to determine when a client reboot has occurred

The Secret State Verifier is not a verifier in the sense given in this definition.

## 2.6. Overview of NFSv4.1 Features

The major features of the NFSv4.1 protocol will be reviewed in brief. This will be done to provide an appropriate context for both the reader who is familiar with the previous versions of the NFS protocol and the reader who is new to the NFS protocols. For the reader new to the NFS protocols, there is still a set of fundamental knowledge that is expected. The reader should be familiar with the External Data Representation (XDR) and Remote Procedure Call (RPC) protocols as described in [RFC4506] and [RFC5531]. A basic knowledge of file systems and distributed file systems is expected as well.

In general, this specification of NFSv4.1 will not distinguish those features added in minor version 1 from those present in the base protocol but will treat NFSv4.1 as a unified whole. See Section 4 for a summary of the differences between NFSv4.0 and NFSv4.1.

## 2.7. RPC and Security

As with previous versions of NFS, the External Data Representation (XDR) and Remote Procedure Call (RPC) mechanisms used for the NFSv4.1 protocol are those defined in [RFC4506] and [RFC5531], as extended by [RFC9289]) to provide TLS-based encryption and client-host authentication. NFSv4.1 security. A description of the basics of NFSv4.1 security will appear in an NFSv4-wide security document, to be derived from [I-D.dnoveck-nfsv4-security].

NFSv4.1 introduces parallel access (see Section 2.8.2), through the use of pNFS. The security framework described above is significantly modified by the introduction of pNFS (see Section 18.11), because of the addition of additional actors and because data access sometimes does not rely on RPC for principal identification/authentication. The appropriate handling depends on the data access protocol used (see Section 18.4.5) which depends in turn on the layout type (see Section 18.2.) The sections 18.11.1 through 18.11.3 discuss the security implications of using different sorts of data access protocols.

## 2.8. Protocol Structure

### 2.8.1. Core Protocol

Unlike NFSv3, which relied on a series of ancillary protocols (e.g., NLM, NSM (Network Status Monitor), MOUNT), within all minor versions of NFSv4 a single RPC protocol is available to make requests to the server. Facilities that had used separate protocols, such as locking, are now integrated within a single unified protocol, although, to implement pNFS, different data access protocols may also be used.

### 2.8.2. Parallel Access

Minor version 1 supports high-performance data access to a clustered server implementation by enabling a separation of metadata access and data access, with the latter able to be done to multiple servers in parallel.

Such parallel data access is controlled by recallable objects known as "layouts", which are integrated into the protocol locking model. Clients direct requests for data access to a set of data servers specified by the layout via a data storage protocol which may be NFSv4.1 or may be another protocol.

Because the protocols used for parallel data access are not necessarily RPC-based, the RPC-based security model (Section 2.7) is impacted (see Section 18.11). The degree of impact varies with the protocol (see Section 18.4.5) used for data access, and can be as low as zero for some RPC-based data access protocols (see Section 20.25).

## 2.9. File System Model

The general file system model used for the NFSv4.1 protocol is the same as for previous minor versions of NFSv4. The server file system is hierarchical with the regular files contained within being treated as opaque byte streams. File names MAY be restricted to UTF-8-encoded strings of Unicode characters or treated as opaque. In addition, for some file systems, name handling MAY reflect the UTF-8 canonical equivalence relation, and in some cases, case-based equivalence relations as well.

The NFSv4.1 protocol does not rely on a separate protocol to provide for the initial mapping between path name and filehandle. All file systems exported by a server are presented as a tree so that all file systems are reachable from a special per-server global root filehandle. This allows LOOKUP operations to be used to perform functions previously provided by the MOUNT protocol. The server is responsible for providing any necessary pseudo file systems to bridge any gaps that arise due to unexported portions of the server-local name space that are between exported file systems.

#### 2.9.1. Filehandles

As in previous versions of the NFS protocol, opaque filehandles are used to identify individual files and directories. Lookup-type and create operations translate file and directory names to filehandles, which are then used to identify objects in subsequent operations.

The NFSv4.1 protocol provides support for persistent filehandles, guaranteed to be valid for the lifetime of the file system object designated and never to be reused after that. In addition, it provides support to allow servers to provide filehandles with more limited validity guarantees, referred to as volatile filehandles.

#### 2.9.2. Numbered File Attributes

The NFSv4.1 protocol has a rich and extensible file object attribute structure in which each attribute is assigned an attribute number. The set of such attributes can be usefully divided in a number of ways, in order to provide helpful context for server implementers choosing to implement or not implement particular attributes which are not REQUIRED and for client implementers deciding how to deal with non-support of particular attributes which are not REQUIRED.

- \* Attributes differ as to their scope, with only a subset applicable to a particular file object, while others apply to an entire file system.

As a practical matter, attributes applicable to a single file object, often require support within the file system proper. While this functionality is most often provided by a file system created initially for local access and only later adapted to remote use through an NFSv4.1 server, there are also file systems purpose-built for remote access.

Attributes applicable to an entire file system do not typically require support within the file system proper. One possible exception is when such attributes can be set to indicate the client's desire for some particular feature that inherently require file system support.

Note that, in all these cases, applications are most likely to be adapted to features that can be accessed using existing file access facilities. As a result, implementers are unlikely to devote efforts to implementation of OPTIONAL features and attributes which require interactions with applications while being more open to attributes usable by the client and server to communicate to optimize data flows without requiring application involvement.

- \* Attributes differ as to their mutability characteristics, including whether the attribute is question can be modified explicitly by the client and whether the attribute modification happens as a result of performing other operations, such as modifying a file or directory
- \* In this update of the NFSv4.1 specification, the details for handling authorization-related attributes are the responsibility of the NFSv4-wide security documents expected to be derived from [I-D.dnoveck-nfsv4-security] and [I-D.ietf-nfsv4-acls-update].

Although the detailed categorization of such attributes will be the responsibility of the security documents, this document will, in Section 11.3, provide a brief summary and make clear that some of these are more necessary than others, and that they all cannot be reasonably treated as having the being of the same class regarding the need for server support.

Attributes are divided into a number of classes based on the protocol's requirements/recommendations for server implementation and the client's expected response to a server's non-support of those attributes. This categorization differs from that appearing previously for a number of reasons with the specific differences explained in Section 11.2.

- \* A significant number of attributes are described as REQUIRED so that servers MUST provide support for them.



These include the same set of attributes described in this way in RFCs [RFC7530] and [RFC8881]. In addition, there are a set of authorization-related attributes that need to be included in this group for reasons explained in [I-D.dnoveck-nfsv4-security]. The inclusion of all these attributes is discussed in more detail in Section 11.4.

- \* Many attributes are truly OPTIONAL, even though such attributes have been erroneously categorized as "RECOMMENDED" in the past. These attributes are discussed in more detail in Section 11.5.

For a more detailed explanation of these shifts in terminology, see Section 11.2.

- \* It appears necessary to designate certain authorization-related attributes as Experimental.

These attributes are discussed in more detail in Sections 11.3 and 11.6.

Descriptions of each specific attribute appears in the following places:

- \* Those which are authorization-related are described in [I-D.dnoveck-nfsv4-security] and [I-D.ietf-nfsv4-acls-update].
- \* Others are described sections 11.14 through 11.17.5

These descriptions can be found using the three sources described below:

- \* The list of Authorization-related attributes appears in Section 11.18
- \* The list of REQUIRED Attributes appears in Table 4 within Section 11.10.
- \* The list of attributes which are not REQUIRED) appears in Table 5 within Section 11.11.

"Named attributes", despite this designation, which will be retained, differ substantially from file attributes per se and are explained in Section 2.9.3. The entire named attribute feature is OPTIONAL and such objects, when they exist, are not part of the categorization above.

### 2.9.3. Named Attributes

A named attribute is an opaque byte stream that is associated with a directory or file and referred to by a string name. The named attribute feature was originally intended to be used to provide extended attributes providing more flexibility than is possible for attributes defined within a standards-track protocol.

Servers providing support for the named attribute feature, which is OPTIONAL, allow a number of opaque byte streams to be associated with a directory or a file. This feature allows applications to define multiple data streams associated with a file object, each of which can be opened, read, and written just as files are.

This approach to the issue provided a level of functionality beyond that typically available for extended attributes which typically restrict the length of such attributes in order to make them available in the same way that other attributes are. Nevertheless, despite this mismatch, NFSv4.1 makes these data streams available for that purpose and they have been used to support extended attribute functionality. The possible development of extended attribute protocol support more tailored to that specific application is a matter to be dealt with in future minor versions.

For further information about the use of named attributes, see Section 11.7

### 2.9.4. Multi-Server Namespace

NFSv4.1 contains a number of features to allow implementation of namespaces that cross server boundaries and that allow and facilitate a nondisruptive transfer of support for individual file systems between servers. They are all based upon attributes that allow one file system to specify alternate, additional, and new location information that specifies how the client may access that file system.

These attributes can be used to provide for individual active file systems:

- \* Alternate network addresses to access the current file system instance.
- \* The locations of alternate file system instances or replicas to be used in the event that the current file system instance becomes unavailable.

These file system location attributes may be used together with the concept of absent file systems, in which a position in the server namespace is associated with locations on other servers without there being any corresponding file system instance on the current server. For example,

- \* These attributes may be used with absent file systems to implement referrals whereby one server may direct the client to a file system provided by another server. This allows extensive multi-server namespaces to be constructed.
- \* These attributes may be provided when a previously present file system becomes absent. This allows nondisruptive migration of file systems to alternate servers.

### 3. Locking Facilities

As mentioned previously, NFSv4.1 is a single protocol that includes locking facilities. These locking facilities include support for many types of locks including a number of sorts of recallable locks. Recallable locks such as delegations allow the client to be assured that certain events will not occur so long as that lock is held. When circumstances change, the lock is recalled via a callback request. The assurances provided by delegations allow more extensive caching to be done safely when circumstances allow it.

The types of locks are:

- \* Share reservations as established by OPEN operations.
- \* Byte-range locks.
- \* File delegations, which are recallable locks that assure the holder that inconsistent opens and file changes cannot occur so long as the delegation is held.
- \* Directory delegations, which are recallable locks that assure the holder that inconsistent directory modifications cannot occur so long as the delegation is held.
- \* Layouts, which are recallable objects that assure the holder that direct access to the file data may be performed directly by the client and that no change to the data's location that is inconsistent with that access may be made so long as the layout is held.

All non-recallable locks for a given client are tied together under a single client-wide lease. All requests made on sessions associated with the client renew that lease. When the client's lease is not promptly renewed, the client's locks are subject to revocation. In the event of server restart, clients have the opportunity to safely reclaim their locks within a special grace period.

Recallable locks are subject to revocation irrespective of lease state. Servers often need to revoke such locks when recalling them does not result in their prompt return.

#### 4. Differences from NFSv4.0

The following summarizes the major differences between minor version One and the base protocol:

- \* Implementation of the sessions model (Section 7).
- \* Parallel access to data (Section 18).
- \* Addition of the RECLAIM\_COMPLETE operation to better structure the lock reclamation process (Section 25.51).
- \* Enhanced delegation support as follows.
  - Delegations on directories and other file types in addition to regular files (Section 25.39, Section 25.49).
  - Operations to optimize acquisition of recalled or denied delegations (Section 25.49, Section 27.5, Section 27.7).
  - Notifications of changes to files and directories (Section 25.39, Section 27.4).
  - A method to allow a server to indicate that it is recalling one or more delegations for resource management reasons, and thus a method to allow the client to pick which delegations to return (Section 27.6).
- \* Attributes can be set atomically during exclusive file create via the OPEN operation (See the new EXCLUSIVE4\_1 creation method in Section 25.16).

- \* Open files can be preserved if removed and the hard link count ("hard link" is defined in an Open Group [hardlink] standard) goes to zero, thus obviating the need for clients to rename deleted files to partially hidden names, a technique colloquially called "silly rename" (See the new OPEN4\_RESULT\_PRESERVE\_UNLINKED reply flag in Section 25.16).
- \* Improved compatibility with Microsoft Windows for Access Control Lists (sACL and dacl attributes, aCL inheritance).
- \* Data retention (Section 11.17).
- \* Identification of the implementation of the NFS client and server (Section 25.35).
- \* Support for notification of the availability of byte-range locks (See the new OPEN4\_RESULT\_MAY\_NOTIFY\_LOCK reply flag in Section 25.16 and see Section 27.11).
- \* In NFSv4.1, LIPKEY and SPKM-3 are not required security mechanisms [RFC2847].

## 5. Core Infrastructure

### 5.1. Introduction

NFSv4.1 relies on core infrastructure common to nearly every operation. This core infrastructure is described in the remainder of this section.

### 5.2. RPC and XDR

The NFSv4.1 protocol is a Remote Procedure Call (RPC) application that uses RPC version 2 and the corresponding external Data Representation (XDR) as defined in [RFC5531] and [RFC4506]. The transport-level encryption and client-host authentication facilities described in [RFC9289] can also be used.

### 5.3. RPC-Based Security

In addition to the above, as discussed in Section 5.3.1, some security flavors provide additional security services.

NFSv4.1 clients and servers MUST implement RPCSEC\_GSS. (This requirement to implement is not a requirement to use.) Other flavors, such as AUTH\_NONE and AUTH\_SYS, can be implemented as well, although the security implications of doing so need to be carefully considered, particularly when the client host is not itself

authenticated. In particular, it is RECOMMENDED by rpc-tls [RFC9289] that AUTH\_SYS not be used when client host authentication is not in effect.

#### 5.3.1. RPCSEC\_GSS and Security Services

RPCSEC\_GSS [RFC2203] uses the functionality of GSS-API [RFC2743]. This allows for the use of various security mechanisms by the RPC layer without the additional implementation overhead of adding RPC security flavors.

##### 5.3.1.1. GSS Server Principal

Regardless of what security mechanism under RPCSEC\_GSS is being used, the NFS server MUST identify itself in GSS-API via a GSS\_C\_NT\_HOSTBASED\_SERVICE name type. GSS\_C\_NT\_HOSTBASED\_SERVICE names are of the form:

service@hostname

For NFS, the "service" element is

nfs

Implementations of security mechanisms will convert nfs@hostname to various different forms. For Kerberos V5, the following form is RECOMMENDED:

nfs/hostname

#### 5.4. COMPOUND and CB\_COMPOUND

A significant departure from the versions of the NFS protocol before NFSv4 is the introduction of the COMPOUND procedure. For the NFSv4 protocol, in all minor versions, there are exactly two RPC procedures, NULL and COMPOUND. The COMPOUND procedure is defined as a series of individual operations and these operations perform the sorts of functions performed by traditional NFS procedures.

The operations combined within a COMPOUND request are evaluated in order by the server, without any atomicity guarantees. A limited set of facilities exist to pass results from one operation to another. Once an operation returns a failing result, the evaluation ends and the results of all evaluated operations are returned to the client.

With the use of the COMPOUND procedure, the client is able to build simple or complex requests. These COMPOUND requests allow for a reduction in the number of RPCs needed for logical file system

operations. For example, multi-component look up requests can be constructed by combining multiple LOOKUP operations. Those can be further combined with operations such as GETATTR, READDIR, or OPEN plus READ to do more complicated sets of operation without incurring additional latency.

NFSv4.1 also contains a considerable set of callback operations in which the server makes an RPC directed at the client. Callback RPCs have a similar structure to that of the normal server requests. In all minor versions of the NFSv4 protocol, there are two callback RPC procedures: CB\_NULL and CB\_COMPOUND. The CB\_COMPOUND procedure is defined in an analogous fashion to that of COMPOUND with its own set of callback operations.

The addition of new server and callback operations within the COMPOUND and CB\_COMPOUND request framework provides a means of extending the protocol in subsequent minor versions.

Except for a small number of operations needed for session creation, server requests and callback requests are performed within the context of a session. Sessions provide a client context for every request and support robust replay protection, which provides the original response to the caller in the case of non-idempotent requests.

#### 5.5. Client Identifiers and Client Owners

For each operation that obtains or depends on locking state, the specific client needs to be identifiable by the server.

Each distinct client instance is represented by a client ID. A client ID is a 64-bit identifier representing a specific client at a given time. The client ID is changed whenever the client re-initializes, and may change when the server re-initializes. Client IDs are used to support lock identification and crash recovery.

During steady state operation, the client ID associated with each operation is derived from the session (See Section 7) on which the operation is sent. A session is associated with a client ID when the session is created.

Unlike NFSv4.0, the only NFSv4.1 operations possible before a client ID is established are those needed to establish the client ID.

A sequence of an EXCHANGE\_ID operation followed by a CREATE\_SESSION operation using that client ID (eir\_clientid as returned from EXCHANGE\_ID) is required to establish and confirm the client ID on the server. Establishment of identification by a new incarnation of

the client also has the effect of immediately releasing any locking state that a previous incarnation of that same client might have had on the server. Such released state would include all byte-range lock, share reservation, layout state. Also, where the server supports neither the CLAIM\_DELEGATE\_PREV nor the CLAIM\_DELEG\_PREV\_FH claim types, all delegation state associated with the same client is released as well. For discussion of delegation state recovery, see Section 15.2.1. For discussion of layout state recovery, see Section 18.9.1.

Releasing such state requires that the server be able to determine that one client instance is the successor of another. Where this cannot be done, for any of a number of reasons, the locking state will remain for a time subject to lease expiration (See Section 13.3) and the new client will need to wait for such state to be removed, if it makes conflicting lock requests.

Client identification is encapsulated by the following client owner data type:

```
struct client_owner4 {  
    verifier4      co_verifier;  
    opaque         co_ownerid<NFS4_OPAQUE_LIMIT>;  
};
```

The first field, `co_verifier`, is a client incarnation verifier, allowing the server to distinguish successive incarnations (e.g., reboots) of the same client. The server will start the process of canceling the client's leased state if `co_verifier` is different than what the server has previously recorded for the identified client (as specified in the `co_ownerid` field).

The second field, `co_ownerid`, contains the client owner id. This is a variable-length string that uniquely defines the client so that subsequent instances of the same client bear the same `co_ownerid` with a different verifier.

There are several considerations for how the client generates the `co_ownerid` string:

- \* The string should be unique so that multiple clients do not present the same string. The consequences of two clients presenting the same string range from one client getting an error to one client having its leased state abruptly and unexpectedly cancelled.



- \* The string should be selected so that subsequent incarnations (e.g., restarts) of the same client cause the client to present the same string. The implementer is cautioned from an approach that requires the string to be recorded in a local file because this precludes the use of the implementation in an environment where there is no local disk and all file access is from an NFSv4.1 server.
- \* The string should be the same for each server network address that the client accesses. This way, if a server has multiple interfaces, the client can trunk traffic over multiple network paths as described in Section 7.5. (Note: the precise opposite was advised in the NFSv4.0 specification [RFC3530].)
- \* The algorithm for generating the string should not assume that the client's network address will not change, unless the client implementation knows it is using statically assigned network addresses. This includes changes between client incarnations and even changes while the client is still running in its current incarnation. Thus, with dynamic address assignment, if the client includes just the client's network address in the co\_ownerid string, there is a real risk that after the client gives up the network address, another client, using a similar algorithm for generating the co\_ownerid string, would generate a conflicting co\_ownerid string.

Given the above considerations, an example of a well-generated co\_ownerid string is one that includes:

- \* If applicable, the client's statically assigned network address.
- \* Additional information that tends to be unique, such as one or more of:
  - The client machine's serial number (for privacy reasons, it is best to perform some one-way function on the serial number).
  - A Media Access Control (MAC) address (again, a one-way function should be performed).
  - The timestamp of when the NFSv4.1 software was first installed on the client (though this is subject to the previously mentioned caution about using information that is stored in a file, because the file might only be accessible over NFSv4.1).

- A true random number. However, since this number ought to be the same between client incarnations, this shares the same problem as that of using the timestamp of the software installation.
- \* For a user-level NFSv4.1 client, it should contain additional information to distinguish the client from other user-level clients running on the same host, such as a process identifier or other unique sequence.

The client ID is assigned by the server (the `eir_clientid` result from `EXCHANGE_ID`) and should be chosen so that it will not conflict with a client ID previously assigned by the server. This applies across server restarts.

In the event of a server restart, a client may find out that its current client ID is no longer valid when it receives an `NFS4ERR_STALE_CLIENTID` error. The precise circumstances depend on the characteristics of the sessions involved, specifically whether the session is persistent (See Section 8), but in each case the client will receive this error when it attempts to establish a new session with the existing client ID and receives the error `NFS4ERR_STALE_CLIENTID`, indicating that a new client ID needs to be obtained via `EXCHANGE_ID` and the new session established with that client ID.

When a session is not persistent, the client will find out that it needs to create a new session as a result of getting an `NFS4ERR_BADSESSION`, since the session in question was lost as part of a server restart. When the existing client ID is presented to a server as part of creating a session and that client ID is not recognized, as would happen after a server restart, the server will reject the request with the error `NFS4ERR_STALE_CLIENTID`.

In the case of the session being persistent, the client will re-establish communication using the existing session after the restart. This session will be associated with the existing client ID but may only be used to retransmit operations that the client previously transmitted and did not see replies to. Replies to operations that the server previously performed will come from the reply cache; otherwise, `NFS4ERR_DEADSESSION` will be returned. Hence, such a session is referred to as "dead". In this situation, in order to perform new operations, the client needs to establish a new session. If an attempt is made to establish this new session with the existing client ID, the server will reject the request with `NFS4ERR_STALE_CLIENTID`.

When NFS4ERR\_STALE\_CLIENTID is received in either of these situations, the client needs to obtain a new client ID by use of the EXCHANGE\_ID operation, then use that client ID as the basis of a new session, and then proceed to any other necessary recovery for the server restart case (See Section 13.4.2).

See the descriptions of EXCHANGE\_ID (Section 25.35) and CREATE\_SESSION (Section 25.36) for a complete specification of these operations.

#### 5.5.1. Upgrade from NFSv4.0 to NFSv4.1

To facilitate upgrade from NFSv4.0 to NFSv4.1, a server may compare a value of data type client\_owner4 in an EXCHANGE\_ID with a value of data type nfs\_client\_id4 that was established using the SETCLIENTID operation of NFSv4.0. A server that does so will allow an upgraded client to avoid waiting until the lease (i.e., the lease established by the NFSv4.0 instance client) expires. This requires that the value of data type client\_owner4 be constructed the same way as the value of data type nfs\_client\_id4. If the latter's contents included the server's network address (per the recommendations of the NFSv4.0 specification [RFC3530]), and the NFSv4.1 client does not wish to use a client ID that prevents trunking, it should send two EXCHANGE\_ID operations. The first EXCHANGE\_ID will have a client\_owner4 equal to the nfs\_client\_id4. This will clear the state created by the NFSv4.0 client. The second EXCHANGE\_ID will not have the server's network address. The state created for the second EXCHANGE\_ID will not have to wait for lease expiration, because there will be no state to expire.

#### 5.5.2. Server Release of Client ID

NFSv4.1 introduces a new operation called DESTROY\_CLIENTID (Section 25.50), which the client uses to destroy a client ID it no longer needs. This permits graceful, bilateral release of a client ID. The operation cannot be used if there are sessions associated with the client ID, or state with an unexpired lease.

If the server determines that the client holds no associated state for its client ID (associated state includes unrevoked sessions, opens, locks, delegations, layouts, and wants), the server MAY choose to unilaterally release the client ID in order to conserve resources. If the client contacts the server after this release, the server MUST ensure that the client receives the appropriate error so that it will use the EXCHANGE\_ID/CREATE\_SESSION sequence to establish a new client ID. The server ought to be very hesitant to release a client ID since the resulting work on the client to recover from such an event will be the same burden as if the server had failed and restarted.

Typically, a server would not release a client ID unless there had been no activity from that client for many minutes. As long as there are sessions, opens, locks, delegations, layouts, or wants, the server **MUST NOT** release the client ID. See Section 7.13.1.4 for discussion on releasing inactive sessions.

### 5.5.3. Resolving Client Owner Conflicts

When the server gets an `EXCHANGE_ID` for a client owner that currently has no state, or that has state but the lease has expired, the server **MUST** allow the `EXCHANGE_ID` and confirm the new client ID if followed by the appropriate `CREATE_SESSION`.

When the server gets an `EXCHANGE_ID` for a new incarnation of a client owner that currently has an old incarnation with state and an unexpired lease, the server is allowed to dispose of the state of the previous incarnation of the client owner if one of the following is true:

- \* The client ID was created without explicit state protection (i.e., `SP4_NONE` was used) and without client host authentication while the current `EXCHANGE_ID` shares those characteristics and the principals used for client ID creation and the current `EXCHANGE_ID` match as well.
- \* The client ID was created without explicit state protection (i.e., `SP4_NONE` was used) and with client host authentication while the current `EXCHANGE_ID` shares those characteristics with the `EXCHANGE_ID` used to create the client ID while the authenticated client hosts match as well.
- \* The principal that created the client ID for the client owner is the same as the principal that is sending the `EXCHANGE_ID` operation. Note that if the client ID was created with `SP4_MACH_CRED` state protection (Section 25.35), the principal **MUST** be based on `RPCSEC_GSS` authentication, the `RPCSEC_GSS` service used **MUST** be integrity or privacy, and the same GSS mechanism and principal **MUST** be used as that used when the client ID was created.
- \* The client ID was established with `SP4_SSV` protection (Section 25.35, Section 7.8.3) and the client sends the `EXCHANGE_ID` with the security flavor set to `RPCSEC_GSS` using the GSS SSV mechanism (Section 7.9).
- \* The client ID was established with `SP4_SSV` protection, and under the conditions described herein, the `EXCHANGE_ID` was sent with `SP4_MACH_CRED` state protection. Because the SSV might not persist

across client and server restart, and because the first time a client sends EXCHANGE\_ID to a server it does not have an SSV, the client MAY send the subsequent EXCHANGE\_ID without an SSV RPCSEC\_GSS handle. Instead, as with SP4\_MACH\_CRED protection, the principal MUST be based on RPCSEC\_GSS authentication, the RPCSEC\_GSS service used MUST be integrity or privacy, and the same GSS mechanism and principal MUST be used as that used when the client ID was created.

If none of the above situations apply, the server MUST return NFS4ERR\_CLID\_INUSE.

If the server accepts the principal and co\_ownerid as matching that which created the client ID, and the co\_verifier in the EXCHANGE\_ID differs from the co\_verifier used when the client ID was created, then after the server receives a CREATE\_SESSION that confirms the client ID, the server deletes state. If the co\_verifier values are the same (e.g., the client either is updating properties of the client ID (Section 25.35) or is attempting trunking (Section 7.5), the server MUST NOT delete state.

## 5.6. Server Owners

The server owner is similar to a client owner (Section 5.5), but unlike the client owner, there is no shorthand server ID. The server owner is defined in the following data type:

```
struct server_owner4 {  
    uint64_t      so_minor_id;  
    opaque        so_major_id<NFS4_OPAQUE_LIMIT>;  
};
```

The server owner is returned from EXCHANGE\_ID. When the so\_major\_id fields are the same in two EXCHANGE\_ID results, the connections that each EXCHANGE\_ID were sent over can be assumed to address the same server (as defined in Section 2.5). If the so\_minor\_id fields are also the same, then not only do both connections connect to the same server, but the session can be shared across both connections. The reader is cautioned that multiple servers may deliberately or accidentally claim to have the same so\_major\_id or so\_major\_id/so\_minor\_id; the reader should examine Sections 7.5 and 25.35 in order to avoid acting on falsely matching server owner values.

The considerations for generating an `so_major_id` are similar to that for generating a `co_ownerid` string (see Section 5.5). The consequences of two servers generating conflicting `so_major_id` values are less dire than they are for `co_ownerid` conflicts because the client can use `RPCSEC_GSS` to compare the authenticity of each server (See Section 7.5).

## 5.7. Transport Layers

### 5.7.1. REQUIRED and RECOMMENDED Properties of Transports

NFSv4.1 works over Remote Direct Memory Access (RDMA) and non-RDMA-based transports with the following attributes:

- \* The transport supports reliable delivery of data, which NFSv4.1 requires. However the possibility of connections breaking is addressed in NFSv4.1 by a session-based replay cache to prevent the spurious re-execution of non-idempotent requests or modifying idempotent requests.
- \* The transport delivers data in the order it was sent. Ordered delivery simplifies detection of transmit errors, and simplifies the sending of arbitrary sized requests and responses via the record marking protocol [RFC5531].

Because efficient handling is required when sending large amounts of data, congestion control facilities are a significant concern.

- \* When NFSv4.1 is used over an IP-based network protocol, it is REQUIRED that the transport provide congestion control.
- \* When NFSv4.1 is used over a non-IP network protocol, it is RECOMMENDED that the transport provide congestion control.

To enhance the possibilities for interoperability, it is strongly recommended that NFSv4.1 client and server implementations support operation over the TCP transport protocol.

It is permissible for a connectionless transport to be used under NFSv4.1; however, reliable and in-order delivery of data combined with congestion control by the connectionless transport is REQUIRED. As a consequence, UDP by itself MUST NOT be used as an NFSv4.1 transport, although transports to be used for NFSv4.1 may be layered on UDP. NFSv4.1 assumes that a client transport address and server transport address used to send data over a transport together constitute a connection, even if the underlying transport eschews the concept of a connection.

### 5.7.2. Client and Server Transport Behavior

[Author Aside]: Section substantially revised to address unjustified use of RFC2119-defined keywords regarding retries and replace that with appropriate implementation advice.

When a connection-oriented transport (e.g., TCP) is used, the client and server are normally expected to maintain the use of connections already established for a considerable length of time. This is for a number reasons:

- \* This will prevent the weakening of the transport's congestion control mechanisms by the confusion resulting from dispersing a single burst of network traffic into multiple connections.
- \* This will improve performance for the WAN environment by eliminating the need for connection setup handshakes. This is particularly so given that each new connection requires the re-establishment of a connection id, as well as setting up new connections.
- \* The NFSv4.1 callback model requires the client and server to maintain a client-created backchannel (See Section 7.3.1) for the server to use so that any dropping of the connection will interfere with the use of the backchannel.

Although it is not fatal for a requester to retry without a disconnect between the request and retry, there are good reasons to avoid this practice. The retry does consume resources, especially with RDMA, where each request, retry or not, consumes a credit. Retries for no reason, especially retries sent shortly after the previous attempt, are a poor use of network bandwidth and defeat the purpose of a transport's inherent congestion control system.

There is no situation in which a replier is allowed to silently drop a request, whether the request is a retry or not. (The silent drop behavior of RPCSEC\_GSS [RFC2203] is not relevant here since this behavior happens at the RPCSEC\_GSS layer, which is at a lower layer in the request processing.) While the replier MAY disconnect the connection, if it does not do so, it is obligated to execute the request or return an appropriate error based on the contents of the reply cache (see Section 7.6.1).

When sending a reply, the replier MUST send the reply to the same full network address (e.g., if using an IP-based transport, the source port of the requester is part of the full network address) from which the requester sent the request. If using a connection-oriented transport, replies MUST be sent on the connection from which the request was received.

If a connection is dropped after the replier receives the request but before the replier sends the reply, the replier might have a pending reply. If a connection is established with the same source and destination full network address as the dropped connection, then the replier MUST NOT send the reply until the requester retries the request. The reason for this prohibition is that the requester MAY retry a request over a different connection (provided that connection is associated with the original request's session).

When using RDMA transports, there are other reasons for avoiding retries over the same connection:

- \* RDMA transports use "credits" to enforce flow control, where a credit is a right to a peer to transmit a message. If one peer were to retransmit a request (or reply), it would consume an additional credit. If the replier retransmitted a reply, it would certainly result in an RDMA connection loss, since the requester would typically only post a single receive buffer for each request. If the requester retransmitted a request, the additional credit consumed on the server might lead to RDMA connection failure unless the client accounted for it and decreased its available credit, leading to wasted resources.
- \* RDMA credits present a new issue to the reply cache in NFSv4.1. The reply cache may be used when a connection within a session is lost, such as after the client reconnects. Credit information is a dynamic property of the RDMA connection, and stale values must not be replayed from the cache. This implies that the reply cache contents must not be blindly used when replies are sent from it, and credit information appropriate to the channel must be refreshed by the RPC layer.

In addition, as described in Section 7.6.2, while a session is active, an NFSv4.1 requester that ceases to wait for an outstanding reply MUST take appropriate care to avoid that situation vitiating guarantees needed to maintain the exactly-once semantics needed for the successful operation of the session-based reply cache.



### 5.7.3. Ports

Historically, NFSv3 servers have listened over TCP port 2049. The registered port 2049 [RFC3232] for the NFS protocol should be the default configuration for NFSv4.1, although the port 20049 is used for NFSv4.1 layered on RPC-over-RDMA.

The use of a reserved port has been common for NFS implementations and it is expected that this will apply to NFSv4.1 as well. While the use of RPC binding protocols as described in [RFC1833] is a possibility, there is no requirement that servers provide support for such use.

In light of this, a client should avoid this sort of use unless it has good reason to expect such support to be present on the server, while accessing NFS services at the appropriate well-known port depending on the transport to be used.

## 6. Security-related Infrastructure

NFSv4.1 relies on the Infrastructure described by the NFSv4-wide security-related documents, currently [I-D.dnoveck-nfsv4-security] and [I-D.ietf-nfsv4-acls-update]. This infrastructure includes:

- \* The RPC-based facilities to provide authentication, privacy and integrity, including facilities provided by the various authentication flavors and those provided at the transport layer.
- \* The security negotiation facilities described in section 16 of the security document, as they have been enhanced to support selection of transport-layer facilities, as well as authentication flavors and associated services.
- \* The authorization facilities described in Sections 5.3, 5.4, and 7 of the security document.
- \* The audit and alarm facilities described in Section 13 of the security document.

There are, however, a number of places where the NFSv4-wide treatment needs to be supplemented to deal with NFSv4.1-specific features, requirements, and recommendations as discussed below:

- \* The parallel NFS feature is described in Section 18, with security for it dealt with in Section 18.11. The case of the file layout type is described in Section 20 with security for it dealt with in Section 20.25.

The security for parallel NFS is dealt with in this specification even though it relies on the security infrastructure described in the NFSv4-wide security document. As a result, it will be discussed in a restructured Section 28.

- \* The handling of SECINFO and SECINFO\_NO\_NAME is complicated because both share the extensions to the negotiation process described in Section 12 of the NFSv4-wide security document, currently [I-D.dnoveck-nfsv4-security], while the former operation is present in all minor versions while the latter is specific to NFSv4.1.
- \* The requirements and recommendations regarding associated security services are discussed in Section 6.1. The discussion had been modified to include the possibility that encryption at the RPC transport layer might obviate the need for these services, although the existing requirements and recommendations still stand.

[Author Aside]: Further work in this area is likely and there should be working group discussion of possible changes. Of particular concern is the use of "SHOULD" in connection with support for privacy, as it is not clear what might be valid reasons not to support this. The provision of confidentiality using transport-based encryption further complicates the matter, although it needs to be clear that the need that confidentiality be available in some form is strongly recommended.

#### 6.1. NFSv4.1-specific Recommendations and Requirements Regarding Security Services

[Author Aside]: Significant revisions have been made to address the hole created by the fact that the discussion of client support of data privacy uses the word "SHOULD".

Via the GSS-API, RPCSEC\_GSS can be used to identify and authenticate users on clients to servers, and servers to users. Authentication of the client itself is not provided but can be provided by RPC independently of the use of RPCSEC\_GSS.

GSS-API can also perform integrity checking on the entire RPC message, including the RPC header, and on the arguments or results. Finally, privacy/confidentiality, usually via encryption, is a service available with RPCSEC\_GSS. Privacy is provided for the arguments and results. Note that if privacy is selected, integrity, authentication, and identification are enabled. If privacy is not selected, but integrity is selected, authentication and identification are enabled. If integrity and privacy are not selected, but authentication is enabled, identification is enabled. RPCSEC\_GSS does not provide identification as a separate service.

Although GSS-API has an authentication service distinct from its privacy and integrity services, GSS-API's authentication service is not used for RPCSEC\_GSS's authentication service. Instead, each RPC request and response header is integrity protected with the GSS-API integrity service, and this allows RPCSEC\_GSS to offer per-RPC authentication and identity. See [RFC2203] for more information.

NFSv4.1 client and servers MUST support RPCSEC\_GSS's integrity and authentication service. NFSv4.1 servers MUST support RPCSEC\_GSS's privacy service.

NFSv4.1 clients SHOULD support RPCSEC\_GSS's privacy service. Given that it has never been made clear, as suggested by the definition of "SHOULD" in [RFC2119], it has to be assumed that this statement, appearing in previous specifications has been treated as providing permission for clients not to support RPCSEC\_GSS privacy. In light of this situation, it needs to be understood that, with regard to the use of "SHOULD" above, valid reasons to bypass this recommendation are limited to the reliance of implementors on those previous specifications and the difficulty of changing them now.

The following consequences need to be kept in mind by those not providing such support:

- \* Any data accessed on connections for which rpc-tls support is not provided will be available, in the clear, to those with the ability to monitor network traffic on a network segment used to effect the access.
- \* The existence of clients without privacy support would make it difficult or impossible to enforce privacy constraints that would otherwise be straightforward. The effect is to undercut the process of security negotiation, which is the only possible way to provide confidentiality when rpc-tls encryption is not in effect.

The reader is directed to Section 18.3.1 of [I-D.dnoveck-nfsv4-security] for a more complete discussion of security issues regarding data in flight.

## 6.2. NFSv4.1-specific Details of Security Negotiation

Unlike NFSv4.0, which only has the SECINFO operation, NFSv4.1 has the SECINFO\_NO\_NAME operation as well. As a result, many of the details of performing security negotiation will be different from those in other minor versions and need to be discussed in this document, in the sections below.

### 6.2.1. Put Filehandle Operations

The term "put filehandle operation" refers to PUTROOTFH, PUTPUBFH, PUTFH, and RESTOREFH. Each of the subsections herein describes how the server handles a subseries of operations that starts with a put filehandle operation.

#### 6.2.1.1. Put Filehandle Operation + SAVEFH

The client is saving a filehandle for a future RESTOREFH, LINK, or RENAME. SAVEFH MUST NOT return NFS4ERR\_WRONGSEC. To determine whether or not the put filehandle operation returns NFS4ERR\_WRONGSEC, the server implementation pretends SAVEFH is not in the series of operations and examines which of the situations described in the other subsections of Section 6.2.1 apply.

#### 6.2.1.2. Two or More Put Filehandle Operations

For a series of N put filehandle operations, the server MUST NOT return NFS4ERR\_WRONGSEC to the first N-1 put filehandle operations. The Nth put filehandle operation is handled as if it is the first in a subseries of operations. For example, if the server received a COMPOUND request with this series of operations -- PUTFH, PUTROOTFH, LOOKUP -- then the PUTFH operation is ignored for NFS4ERR\_WRONGSEC purposes, and the PUTROOTFH, LOOKUP subseries is processed as according to Section 6.2.1.3.

#### 6.2.1.3. Put Filehandle Operation + LOOKUP (or OPEN of an Existing Name)

This situation also applies to a put filehandle operation followed by a LOOKUP or an OPEN operation that specifies an existing component name.

In this situation, the client is potentially crossing a security policy boundary, and the set of security tuples the parent directory supports may differ from those of the child. The server implementation may decide whether to impose any restrictions on security policy administration. There are at least three approaches (`sec_policy_child` is the tuple set of the child export, `sec_policy_parent` is that of the parent).

- (a) `sec_policy_child <= sec_policy_parent` (`<=` for subset). This means that the set of security tuples specified on the security policy of a child directory is always a subset of its parent directory.
- (b) `sec_policy_child ^ sec_policy_parent != {}` (`^` for intersection, `{}` for the empty set). This means that the set of security tuples specified on the security policy of a child directory always has a non-empty intersection with that of the parent.
- (c) `sec_policy_child ^ sec_policy_parent == {}`. This means that the set of security tuples specified on the security policy of a child directory may not intersect with that of the parent. In other words, there are no restrictions on how the system administrator may set up these tuples.

In order for a server to support approaches (b) (for the case when a client chooses a flavor that is not a member of `sec_policy_parent`) and (c), the put filehandle operation cannot return `NFS4ERR_WRONGSEC` when there is a security tuple mismatch. Instead, it should be returned from the LOOKUP (or OPEN by existing component name) that follows.

Since the above guideline does not contradict approach (a), it should be followed in general. Even if approach (a) is implemented, it is possible for the security tuple used to be acceptable for the target of LOOKUP but not for the filehandles used in the put filehandle operation. The put filehandle operation could be a `PUTROOTFH` or `PUTPUBFH`, where the client cannot know the security tuples for the root or public filehandle. Or the security policy for the filehandle used by the put filehandle operation could have changed since the time the filehandle was obtained.

Therefore, an NFSv4.1 server MUST NOT return `NFS4ERR_WRONGSEC` in response to the put filehandle operation if the operation is immediately followed by a LOOKUP or an OPEN by component name.

#### 6.2.1.4. Put Filehandle Operation + LOOKUPP

Since SECINFO only works its way down, there is no way LOOKUPP can return NFS4ERR\_WRONGSEC without SECINFO\_NO\_NAME. SECINFO\_NO\_NAME solves this issue via style SECINFO\_STYLE4\_PARENT, which works in the opposite direction as SECINFO. As with Section 6.2.1.3, a put filehandle operation that is followed by a LOOKUPP MUST NOT return NFS4ERR\_WRONGSEC. If the server does not support SECINFO\_NO\_NAME, the client's only recourse is to send the put filehandle operation, LOOKUPP, GETFH sequence of operations with every security tuple it supports.

Regardless of whether SECINFO\_NO\_NAME is supported, an NFSv4.1 server MUST NOT return NFS4ERR\_WRONGSEC in response to a put filehandle operation if the operation is immediately followed by a LOOKUPP.

#### 6.2.1.5. Put Filehandle Operation + SECINFO/SECINFO\_NO\_NAME

A security-sensitive client is allowed to choose a strong security tuple when querying a server to determine a file object's permitted security tuples. The security tuple chosen by the client does not have to be included in the tuple list of the security policy of either the parent directory indicated in the put filehandle operation or the child file object indicated in SECINFO (or any parent directory indicated in SECINFO\_NO\_NAME). Of course, the server has to be configured for whatever security tuple the client selects; otherwise, the request will fail at the RPC layer with an appropriate authentication error.

In theory, there is no connection between the security flavor used by SECINFO or SECINFO\_NO\_NAME and those supported by the security policy. But in practice, the client may start looking for strong flavors from those supported by the security policy, followed by those in the REQUIRED set.

The NFSv4.1 server MUST NOT return NFS4ERR\_WRONGSEC to a put filehandle operation that is immediately followed by SECINFO or SECINFO\_NO\_NAME. The NFSv4.1 server MUST NOT return NFS4ERR\_WRONGSEC from SECINFO or SECINFO\_NO\_NAME.

#### 6.2.1.6. Put Filehandle Operation + Nothing

The NFSv4.1 server MUST NOT return NFS4ERR\_WRONGSEC.

#### 6.2.1.7. Put Filehandle Operation + Anything Else

"Anything Else" includes OPEN by filehandle.

The security policy enforcement applies to the filehandle specified in the put filehandle operation. Therefore, the put filehandle operation MUST return NFS4ERR\_WRONGSEC when there is a security tuple mismatch. This avoids the complexity of adding NFS4ERR\_WRONGSEC as an allowable error to every other operation.

A COMPOUND containing the series put filehandle operation + SECINFO\_NO\_NAME (style SECINFO\_STYLE4\_CURRENT\_FH) is an efficient way for the client to recover from NFS4ERR\_WRONGSEC.

The NFSv4.1 server MUST NOT return NFS4ERR\_WRONGSEC to any operation other than a put filehandle operation, LOOKUP, LOOKUPP, and OPEN (by component name).

#### 6.2.1.8. Operations after SECINFO and SECINFO\_NO\_NAME

Suppose a client sends a COMPOUND procedure containing the series SEQUENCE, PUTFH, SECINFO\_NO\_NAME, READ, and suppose the security tuple used does not match that required for the target file. By rule (See Section 6.2.1.5), neither PUTFH nor SECINFO\_NO\_NAME can return NFS4ERR\_WRONGSEC. By rule (See Section 6.2.1.7), READ cannot return NFS4ERR\_WRONGSEC. The issue is resolved by the fact that SECINFO and SECINFO\_NO\_NAME consume the current filehandle (note that this is a change from NFSv4.0). This leaves no current filehandle for READ to use so that READ returns NFS4ERR\_NOFILEHANDLE.

#### 6.2.2. LINK and RENAME

The LINK and RENAME operations use both the current and saved filehandles. If the security policy of the saved filehandle rejects the security flavor used in the COMPOUND request's credentials, the server MAY return NFS4ERR\_WRONGSEC from LINK or RENAME. When the server does so, if there is no intersection between the security policies of saved and current filehandles, this implies that it will be impossible for the client to perform the intended LINK or RENAME operation.

For example, suppose the client sends this COMPOUND request: SEQUENCE, PUTFH bFH, SAVEFH, PUTFH aFH, RENAME "c" "d", where filehandles bFH and aFH refer to different directories. Suppose no common security tuple exists between the security policies of aFH and bFH. If the client sends the request using credentials acceptable to bFH's security policy but not aFH's policy, then the PUTFH aFH operation will fail with NFS4ERR\_WRONGSEC. After a SECINFO\_NO\_NAME request, the client sends SEQUENCE, PUTFH bFH, SAVEFH, PUTFH aFH, RENAME "c" "d", using credentials acceptable to aFH's security policy but not bFH's policy. The server returns NFS4ERR\_WRONGSEC on the RENAME operation.

To prevent a client from an endless sequence of a request containing LINK or RENAME, followed by a request containing SECINFO\_NO\_NAME or SECINFO, the server MUST detect when the security policies of the current and saved filehandles have no mutually acceptable security tuple, and MUST NOT return NFS4ERR\_WRONGSEC from LINK or RENAME in that situation. Instead, the server MUST do one of two things:

- \* The server can return NFS4ERR\_XDEV.
- \* The server can allow the security policy of the current filehandle to override that of the saved filehandle, and so return NFS4\_OK.

## 7. Session

NFSv4.1 clients and servers MUST support and MUST use the session feature as described in this section.

### 7.1. Motivation and Overview

Previous versions and minor versions of NFS have suffered from the following:

- \* Lack of support for Exactly Once Semantics (EOS). This includes lack of support for EOS through server failure and recovery.
- \* Limited callback support, including no support for sending callbacks through firewalls, and races between replies to normal requests and callbacks.
- \* Limited trunking over multiple network paths.
- \* Requiring machine credentials for fully secure operation.

Through the introduction of a session, NFSv4.1 addresses the above shortfalls with practical solutions:

- \* EOS is enabled by a reply cache with a bounded size, making it feasible to keep the cache in persistent storage and enable EOS through server failure and recovery. One reason than previous revisions of NFS did not support EOS was because some EOS approaches often limited parallelism. As will be explained in Section 7.6, NFSv4.1 supports EOS without unduly limiting parallelism.
- \* The NFSv4.1 client (defined in Section 2.5) creates transport connections and provides them to the server to use for sending callback requests, thus solving the firewall issue (Section 25.34). Races between responses from client requests and



callbacks caused by the requests are detected via the session's sequencing properties that are a consequence of EOS (Section 7.6.3).

- \* The NFSv4.1 client can associate an arbitrary number of connections with the session, and thus provide trunking (Section 7.5).
- \* The NFSv4.1 client and server produce a session key independent of client and server machine credentials which can be used to compute a digest for protecting critical session management operations (Section 7.8.3).
- \* The NFSv4.1 client can also create secure RPCSEC\_GSS contexts for use by the session's backchannel that do not require the server to authenticate to a client machine principal (Section 7.8.2).

A session is a dynamically created, long-lived server object created by a client and used over time from one or more transport connections. Its function is to maintain the server's state relative to the connection(s) belonging to a client instance. This state is entirely independent of the connection itself, and indeed the state exists whether or not the connection exists. A client may have one or more sessions associated with it so that client-associated state may be accessed using any of the sessions associated with that client's client ID, when connections are associated with those sessions. When no connections are associated with any of a client ID's sessions for an extended time, such objects as locks, opens, delegations, layouts, etc. are subject to expiration. The session serves as an object representing a means of access by a client to the associated client state on the server, independent of the physical means of access to that state.

A single client may create multiple sessions. A single session MUST NOT serve multiple clients.

## 7.2. NFSv4 Integration

Sessions are part of NFSv4.1 and not NFSv4.0. Normally, a major infrastructure change such as sessions would require a new major version number to an Open Network Computing (ONC) RPC program like NFS. However, because NFSv4 encapsulates its functionality in a single procedure, COMPOUND, and because COMPOUND can support an arbitrary number of operations, sessions have been added to NFSv4.1 with little difficulty. COMPOUND includes a minor version number field, and for NFSv4.1 this minor version is set to 1. When the NFSv4 server processes a COMPOUND with the minor version set to 1, it expects a different set of operations than it does for NFSv4.0.

NFSv4.1 defines the SEQUENCE operation, which is required for every COMPOUND that operates over an established session, with the exception of some session administration operations, such as DESTROY\_SESSION (Section 25.37).

#### 7.2.1. SEQUENCE and CB\_SEQUENCE

In NFSv4.1, when the SEQUENCE operation is present, it MUST be the first operation in the COMPOUND procedure. The primary purpose of SEQUENCE is to carry the session identifier. The session identifier associates all other operations in the COMPOUND procedure with a particular session. SEQUENCE also contains required information for maintaining EOS (See Section 7.6). Session-enabled NFSv4.1 COMPOUND requests thus have the form:

```
+-----+-----+-----+-----+-----+-----+
| tag | minorversion | numops | SEQUENCE op | op + args | ...
|      | (= 1)      | (limited) | + args   |           |
+-----+-----+-----+-----+-----+-----+
```

and the replies have the form:

```
+-----+-----+-----+-----+-----+-----+--//
|last status | tag | numres |status + SEQUENCE op + results | //
+-----+-----+-----+-----+-----+-----+--//
//-----+-----
// status + op + results | ...
//-----+-----
```

A CB\_COMPOUND procedure request and reply has a similar form to COMPOUND, but instead of a SEQUENCE operation, there is a CB\_SEQUENCE operation. CB\_COMPOUND also has an additional field called "callback\_ident", which is superfluous in NFSv4.1 and MUST be ignored by the client. CB\_SEQUENCE has the same information as SEQUENCE, and also includes other information needed to resolve callback races (Section 7.6.3).

#### 7.2.2. Client ID and Session Association

Each client ID (Section 5.5) can have zero or more active sessions. A client ID and associated session are required to perform file access in NFSv4.1. Each time a session is used (whether by a client sending a request to the server or the client replying to a callback request from the server), the state leased to its associated client ID is automatically renewed.

State (which can consist of share reservations, locks, delegations, and layouts (Section 3)) is tied to the client ID. Client state is not tied to any individual session. Successive state changing operations from a given state owner MAY go over different sessions, provided the session is associated with the same client ID. A callback MAY arrive over a different session than that of the request that originally acquired the state pertaining to the callback. For example, if session A is used to acquire a delegation, a request to recall the delegation MAY arrive over session B if both sessions are associated with the same client ID. Sections 7.8.1 and 7.8.2 discuss the security considerations around callbacks.

### 7.3. Channels

A channel is not a connection. A channel represents a single direction in which ONC RPC requests are sent as part of a session..

Each session has one or two channels: the fore channel and the backchannel. Because there are at most two channels per session, and because each channel has a distinct purpose, channels are not assigned identifiers.

The fore channel is used for ordinary requests from the client to the server, and carries COMPOUND requests and responses. A session always has a fore channel.

The backchannel is used for callback requests from server to client, and carries CB\_COMPOUND requests and responses. Whether or not there is a backchannel is decided by the client; however, many features of NFSv4.1 require a backchannel. NFSv4.1 servers MUST support backchannels.

Each session has resources for each channel, including separate reply caches (see Section 7.6.1). Note that even the backchannel requires a reply cache (or, at least, a slot table in order to detect retries). This is because it is necessary to avoid re-execution of modifying requests, even if they are idempotent.

#### 7.3.1. Association of Connections, Channels, and Sessions

Each channel is associated with zero or more transport connections (whether of the same transport protocol or different transport protocols). A connection can be associated with one channel or both channels of a session; the client and server negotiate whether a connection will carry traffic for one channel or both channels via the CREATE\_SESSION (Section 25.36) and the BIND\_CONN\_TO\_SESSION (Section 25.34) operations. When a session is created via CREATE\_SESSION, the connection that transported the CREATE\_SESSION

request is automatically associated with the fore channel, and optionally the backchannel. If the client specifies no state protection (Section 25.35) when the session is created, then when SEQUENCE is transmitted on a different connection, the connection is automatically associated with the fore channel of the session specified in the SEQUENCE operation.

A connection's association with a session is not exclusive. A connection associated with the channel(s) of one session may be simultaneously associated with the channel(s) of other sessions including sessions associated with other client IDs.

It is permissible for connections of multiple transport types to be associated with the same channel. For example, both TCP and RDMA connections can be associated with the fore channel. In the event an RDMA and non-RDMA connection are associated with the same channel, it is desirable for the maximum number slots to be at least one more than the total number of RDMA credits (Section 7.6.1). This way, if all RDMA credits are used, the non-RDMA connection can have at least one outstanding request. If a server supports multiple transport types, it **MUST** allow a client to associate connections from each transport to a channel.

It is permissible for a connection of one type of transport to be associated with the fore channel, and a connection of a different type to be associated with the backchannel.

#### 7.4. Server Scope

Servers each specify a server scope value in the form of an opaque string `eir_server_scope` returned as part of the results of an `EXCHANGE_ID` operation. The purpose of the server scope is to allow a group of servers to indicate to clients that a set of servers sharing the same server scope value has arranged to use distinct values of opaque identifiers so that the two servers never assign the same value to two distinct objects. Thus, the identifiers generated by two servers within that set can be assumed compatible so that, in certain important cases, identifiers generated by one server in that set may be presented to another server of the same scope.

The use of such compatible values does not imply that a value generated by one server will always be accepted by another. In most cases, it will not. However, a server will not inadvertently accept a value generated by another server. When it does accept it, it will be because it is recognized as valid and carrying the same meaning as on another server of the same scope.

When servers are of the same server scope, this compatibility of values applies to the following identifiers:

- \* Filehandle values. A filehandle value accepted by two servers of the same server scope denotes the same object. A WRITE operation sent to one server is reflected immediately in a READ sent to the other.
- \* Server owner values. When the server scope values are the same, server owner value may be validly compared. In cases where the server scope values are different, server owner values are treated as different even if they contain identical strings of bytes.

The coordination among servers required to provide such compatibility can be quite minimal, and limited to a simple partition of the ID space. The recognition of common values requires additional implementation, but this can be tailored to the specific situations in which that recognition is desired.

Clients will have occasion to compare the server scope values of multiple servers under a number of circumstances, each of which will be discussed under the appropriate functional section:

- \* When server owner values received in response to EXCHANGE\_ID operations sent to multiple network addresses are compared for the purpose of determining the validity of various forms of trunking, as described in Section 17.5.2.
- \* When network or server reconfiguration causes the same network address to possibly be directed to different servers, with the necessity for the client to determine when lock reclaim should be attempted, as described in Section 13.4.2.1.

When two replies from EXCHANGE\_ID, each from two different server network addresses, have the same server scope, there are a number of ways a client can validate that the common server scope is due to two servers cooperating in a group.

- \* If both EXCHANGE\_ID requests were sent with RPCSEC\_GSS ([RFC2203], [RFC5403], [RFC7861]) authentication and the server principal is the same for both targets, the equality of server scope is validated. It is RECOMMENDED that two servers intending to share the same server scope and server\_owner major\_id also share the same principal name. In some cases, this simplifies the client's task of validating server scope.

- \* The client may accept the appearance of the second server in the `fs_locations` or `fs_locations_info` attribute for a relevant file system. For example, if there is a migration event for a particular file system or there are locks to be reclaimed on a particular file system, the attributes for that particular file system may be used. The client sends the GETATTR request to the first server for the `fs_locations` or `fs_locations_info` attribute with RPCSEC\_GSS authentication. It may need to do this in advance of the need to verify the common server scope. If the client successfully authenticates the reply to GETATTR, and the GETATTR request and reply containing the `fs_locations` or `fs_locations_info` attribute refers to the second server, then the equality of server scope is supported. A client may choose to limit the use of this form of support to information relevant to the specific file system involved (e.g. a file system being migrated).

### 7.5. Trunking

Trunking is the use of multiple connections between a client and server in order to increase the speed of data transfer. NFSv4.1 supports two types of trunking: session trunking and client ID trunking.

In the context of a single server network address, it can be assumed that all connections are accessing the same server, and NFSv4.1 servers **MUST** support both forms of trunking. When multiple connections use a set of network addresses to access the same server, the server **MUST** support both forms of trunking. NFSv4.1 servers in a clustered configuration **MAY** allow network addresses for different servers to use client ID trunking.

Clients may use either form of trunking as long as they do not, when trunking between different server network addresses, violate the servers' mandates as to the kinds of trunking to be allowed (See below). With regard to callback channels, the client **MUST** allow the server to choose among all callback channels valid for a given client ID and **MUST** support trunking when the connections supporting the backchannel allow session or client ID trunking to be used for callbacks.

Session trunking is essentially the association of multiple connections, each with potentially different target and/or source network addresses, to the same session. When the target network addresses (server addresses) of the two connections are the same, the server **MUST** support such session trunking. When the target network addresses are different, the server **MAY** indicate such support using the data returned by the EXCHANGE\_ID operation (See below).

Client ID trunking is the association of multiple sessions to the same client ID. Servers MUST support client ID trunking for two target network addresses whenever they allow session trunking for those same two network addresses. In addition, a server MAY, by presenting the same major server owner ID (Section 5.6) and server scope (Section 7.4), allow an additional case of client ID trunking. When two servers return the same major server owner and server scope, it means that the two servers are cooperating on locking state management, which is a prerequisite for client ID trunking.

Distinguishing when the client is allowed to use session and client ID trunking requires understanding how the results of the EXCHANGE\_ID (Section 25.35) operation identify a server. Suppose a client sends EXCHANGE\_IDs over two different connections, each with a possibly different target network address, but each EXCHANGE\_ID operation has the same value in the eia\_clientowner field. If the same NFSv4.1 server is listening over each connection, then each EXCHANGE\_ID result MUST return the same values of eir\_clientid, eir\_server\_owner.so\_major\_id, and eir\_server\_scope. The client can then treat each connection as referring to the same server (subject to verification; see Section 7.5.1 below), and it can use each connection to trunk requests and replies. The client's choice is whether session trunking or client ID trunking applies.

Session Trunking. If the eia\_clientowner argument is the same in two different EXCHANGE\_ID requests, and the eir\_clientid, eir\_server\_owner.so\_major\_id, eir\_server\_owner.so\_minor\_id, and eir\_server\_scope results match in both EXCHANGE\_ID results, then the client is permitted to perform session trunking. If the client has no session mapping to the tuple of eir\_clientid, eir\_server\_owner.so\_major\_id, eir\_server\_scope, and eir\_server\_owner.so\_minor\_id, then it creates the session via a CREATE\_SESSION operation over one of the connections, which associates the connection to the session. If there is a session for the tuple, the client can send BIND\_CONN\_TO\_SESSION to associate the connection to the session.

Of course, if the client does not desire to use session trunking, it is not required to do so. It can invoke CREATE\_SESSION on the connection. This will result in client ID trunking as described below. It can also decide to drop the connection if it does not choose to use trunking.

Client ID Trunking. If the eia\_clientowner argument is the same in two different EXCHANGE\_ID requests, and the eir\_clientid, eir\_server\_owner.so\_major\_id, and eir\_server\_scope results match in both EXCHANGE\_ID results, then the client is permitted to perform client ID trunking (regardless of whether the

eir\_server\_owner.so\_minor\_id results match). The client can associate each connection with different sessions, where each session is associated with the same server.

The client completes the act of client ID trunking by invoking CREATE\_SESSION on each connection, using the same client ID that was returned in eir\_clientid. These invocations create two sessions and also associate each connection with its respective session. The client is free to decline to use client ID trunking by simply dropping the connection at this point.

When doing client ID trunking, locking state is shared across sessions associated with that same client ID. This requires the server to coordinate state across sessions and the client to be able to associate the same locking state with multiple sessions.

It is always possible that, as a result of various sorts of reconfiguration events, eir\_server\_scope and eir\_server\_owner values may be different on subsequent EXCHANGE\_ID requests made to the same network address.

In most cases, such reconfiguration events will be disruptive and indicate that an IP address formerly connected to one server is now connected to an entirely different one.

Some guidelines on client handling of such situations follow:

- \* When eir\_server\_scope changes, the client has no assurance that any IDs that it obtained previously (e.g., filehandles) can be validly used on the new server, and, even if the new server accepts them, there is no assurance that this is not due to accident. Thus, it is best to treat all such state as lost or stale, although a client may assume that the probability of inadvertent acceptance is low and treat this situation as within the next case.
- \* When eir\_server\_scope remains the same and eir\_server\_owner.so\_major\_id changes, the client can use the filehandles it has, consider its locking state lost, and attempt to reclaim or otherwise re-obtain its locks. It might find that its filehandle is now stale. However, if NFS4ERR\_STALE is not returned, it can proceed to reclaim or otherwise re-obtain its open locking state.



- \* When `eir_server_scope` and `eir_server_owner.so_major_id` remain the same, the client has to use the now-current values of `eir_server_owner.so_minor_id` in deciding on appropriate forms of trunking. This may result in connections being dropped or new sessions being created.

#### 7.5.1. Verifying Claims of Matching Server Identity

When the server responds using two different connections that claim matching or partially matching `eir_server_owner`, `eir_server_scope`, and `eir_clientid` values, the client does not have to trust the servers' claims. The client may verify these claims before trunking traffic in the following ways:

- \* For session trunking to work as desired connections to different network paths need to be associated with the same NFSv4.1 server and usable on the same session. Clients need to verify that that is the case and servers MUST allow clients to perform reliable verification. When a client ID is created, the client needs to, unless client host authentication is in effect, specify that `BIND_CONN_TO_SESSION` is to be verified according to the `SP4_SSV` or `SP4_MACH_CRED` (Section 25.35) state protection options. For `SP4_SSV`, reliable verification depends on a shared secret (the SSV) that is established via the `SET_SSV` (see Section 25.47) operation.

When a new connection is associated with the session (via the `BIND_CONN_TO_SESSION` operation, see Section 25.34), if the client specified `SP4_SSV` state protection for the `BIND_CONN_TO_SESSION` operation, the client MUST send the `BIND_CONN_TO_SESSION` with `RPCSEC_GSS` protection, using integrity or privacy, and an `RPCSEC_GSS` handle created with the `GSS SSV` mechanism (See Section 7.9).

If the client mistakenly tries to associate a connection to a session of a wrong server, the server will either reject the attempt because it is not aware of the session identifier of the `BIND_CONN_TO_SESSION` arguments, or it will reject the attempt because the `RPCSEC_GSS` authentication fails. Even if the server mistakenly or maliciously accepts the connection association attempt, the `RPCSEC_GSS` verifier it computes in the response will not be verified by the client, so the client will know it cannot use the connection for trunking the specified session.

If the client specified SP4\_MACH\_CRED state protection, the BIND\_CONN\_TO\_SESSION operation will use RPCSEC\_GSS integrity or privacy, using the same credential that was used when the client ID was created. Mutual authentication via RPCSEC\_GSS assures the client that the connection is associated with the correct session of the correct server.

- \* For client ID trunking, the client has at least two options for verifying that the same client ID obtained from two different EXCHANGE\_ID operations came from the same server. The first option is to use RPCSEC\_GSS authentication when sending each EXCHANGE\_ID operation. Each time an EXCHANGE\_ID is sent with RPCSEC\_GSS authentication, the client notes the principal name of the GSS target. If the EXCHANGE\_ID results indicate that client ID trunking is possible, and the GSS targets' principal names are the same, the servers are the same and client ID trunking is allowed.

The second option for verification is to use SP4\_SSV protection. When the client sends EXCHANGE\_ID, it specifies SP4\_SSV protection. The first EXCHANGE\_ID the client sends always has to be confirmed by a CREATE\_SESSION call. The client then sends SET\_SSV. Later, the client sends EXCHANGE\_ID to a second destination network address different from the one the first EXCHANGE\_ID was sent to. The client checks that each EXCHANGE\_ID reply has the same `eir_clientid`, `eir_server_owner.so_major_id`, and `eir_server_scope`. If so, the client verifies the claim by sending a CREATE\_SESSION operation to the second destination address, protected with RPCSEC\_GSS integrity using an RPCSEC\_GSS handle returned by the second EXCHANGE\_ID. If the server accepts the CREATE\_SESSION request, and if the client verifies the RPCSEC\_GSS verifier and integrity codes, then the client has proof the second server knows the SSV, and thus the two servers are cooperating for the purposes of specifying server scope and client ID trunking.

## 7.6. Exactly Once Semantics

[Author Aside]: This section, including some subsections, has been substantially modified from the corresponding section appearing in previous specifications [RFC5661] [RFC8881] and earlier drafts of this document. Change has been driven primarily by the incorrect use of RFC2119-defined keywords, most importantly in the case in which RPC requests need to be aborted, leading to some related changes to clarify the appropriate level of checking for the possibility of false retry. As part of this revised description, it is explained that, given the possibility of requests being aborted, the term "Exactly-once semantics" describes an aspiration and that what is really provided would better be called "at-most-once semantics."

Also, the description of retry has been revised to properly use RFC2119 keywords. For more detailed information regarding changes which have been made, see Appendix D.2.1.

Via the session, NFSv4.1 offers what is termed "exactly once semantics" (EOS) for requests sent over a channel. EOS is supported on both the fore channel and backchannel.

Although this term is well-established and will not be changed, it should be noted that what is actually provided is at-most-once semantics to accommodate the possibility that the client will need to abort RPC requests, remaining unsure about whether the requested actions have been performed one time or not at all.

Each COMPOUND or CB\_COMPOUND request that is sent with a leading SEQUENCE or CB\_SEQUENCE operation needs to be executed by the receiver at most once. This requirement holds regardless of whether the request is sent with reply caching specified (See Section 7.6.1.3). The requirement also holds in the case in which NFSv4.1 is a pNFS data access protocol and the requester is sending the request over a session created between a pNFS data client and pNFS data server. To help understand the need for this requirement, we divide the requests sent to be executed into three categories:

- \* Non-idempotent requests.
- \* Idempotent modifying requests.
- \* Idempotent non-modifying requests.

An example of a non-idempotent request is RENAME. Obviously, if a replier executes the same RENAME request twice, and the first execution succeeds, the re-execution will fail. If the replier returns the result from the re-execution, this result is incorrect. For this reason, EOS is required for non-idempotent requests and it is supplemented by returning the original response to the requester.

An example of an idempotent modifying request is a COMPOUND request containing a WRITE operation. Repeated execution of the same WRITE has the same effect as execution of that WRITE a single time. Nevertheless, enforcing EOS for WRITES and other idempotent modifying requests is necessary to avoid data corruption, which could result from executing the same write request multiple times including some executions that occur after the completion of the first is noted by the requester.

Suppose a client sends WRITE A to a noncompliant server that does not enforce EOS, and receives no response, perhaps due to a network partition. The client reconnects to the server and re-sends WRITE A. Now, the server has outstanding two instances of A. The server can be in a situation in which it executes and replies to the retry of A, while the first A is still waiting in the server's internal I/O system for some resource. Upon receiving the reply to the second attempt of WRITE A, the client believes its WRITE is done so it is free to send WRITE B, which overlaps the byte-range of A. When the original A is dispatched from the server's I/O system and executed (thus, the second time A will have been written), then what has been written by B can be overwritten and thus corrupted.

An example of an idempotent non-modifying request is a COMPOUND containing SEQUENCE, PUTFH, READLINK, and nothing else. The re-execution of such a request will not cause data corruption or produce an incorrect result. Nonetheless, the session feature's design provides that the replier **MUST** enforce EOS for all requests, whether or not they are idempotent or modifying.

Note that fully complete EOS is not possible unless the server persists the reply cache in stable storage, and unless the server is somehow implemented to never require a restart (indeed, if such a server exists, the distinction between a reply cache kept in stable storage versus one that is not is one without meaning). See Section 8 for a discussion of persistence in the reply cache. Regardless, even if the server does not persist the reply cache, EOS improves robustness and correctness relative to previous versions of NFS because the earlier duplicate request/reply caches were based on the ONC RPC transaction identifier (XID). Section 7.6.1 explains the shortcomings of the XID as a basis for a reply cache and describes how NFSv4.1 sessions improve upon the XID.

#### 7.6.1. Slot Identifiers and Reply Cache

The RPC layer provides a transaction ID (XID), which, while required to be unique, is not convenient for tracking requests for two reasons. First, the XID is only meaningful to the requester; it cannot be interpreted by the replier except to test for equality with previously sent requests. When consulting an RPC-based duplicate request cache, the opaqueness of the XID requires a computationally expensive look up (often via a hash that includes XID and source address). NFSv4.1 requests include a non-opaque slot ID, which can be used as an index into a slot table, which is far more efficient. Second, because RPC requests can be executed by the replier in any order, there is no bound on the number of requests that may be outstanding at any time. To achieve perfect EOS, using ONC RPC would require storing all replies in the reply cache. XIDs are 32 bits;

storing over four billion ( $2^{32}$ ) replies in the reply cache is not practical. In practice, previous versions of NFS have chosen to store a fixed number of replies in the cache, and to use a least recently used (LRU) approach to replacing cache entries with new entries when the cache is full. In NFSv4.1, the number of outstanding requests is bounded by the size of the slot table, and a sequence ID per slot is used to tell the replier when it is safe to delete a cached reply.

In the NFSv4.1 reply cache, when the requester sends a new request, it selects a slot ID in the range  $0..N$ , where  $N$  is the replier's current maximum slot ID granted to the requester on the session over which the request is to be sent. The value of  $N$  starts out as equal to `ca_maxrequests - 1` (Section 25.36), but can be adjusted by the response to `SEQUENCE` or `CB_SEQUENCE` as described later in this section. The slot ID must be unused by any of the requests that the requester already has active on the session. "Unused" here means the requester has no outstanding request for that slot ID.

A slot contains a sequence ID and the cached reply corresponding to the request sent with that sequence ID. The sequence ID is a 32-bit unsigned value and is therefore in the range  $0..0xFFFFFFFF$  ( $2^{32} - 1$ ). The first time a slot is used, the requester MUST specify a sequence ID of one (Section 25.36). Each time a slot is reused, the request MUST specify a sequence ID that is one greater than that of the previous request on the slot. If the previous sequence ID was `0xFFFFFFFF`, then the next request for the slot MUST have the sequence ID set to zero (i.e.,  $(2^{32} - 1) + 1 \bmod 2^{32}$ ).

The sequence ID accompanies the slot ID in each request. It is for the critical check at the replier: it used to efficiently determine whether a request using a certain slot ID is a retransmit or a new, never-before-seen request. It is not feasible for the requester to assert that it is retransmitting to implement this, because for any given request the requester cannot know whether the replier has seen it unless the replier actually replies. Of course, if the requester has seen the reply, the requester would not retransmit.

The replier compares each received request's sequence ID with the last one previously received for that slot ID, to see if the new request is:

- \* A new request, in which the sequence ID is one greater than that previously seen in the slot (accounting for sequence wraparound). The replier proceeds to execute the new request, and the replier MUST increase the slot's sequence ID by one.

- \* A retransmitted request, in which the sequence ID is equal to that currently recorded in the slot. If the original request has executed to completion, the replier returns the cached reply. See Section 7.6.2 for direction on how the replier deals with retries of requests that are still in progress.
- \* A misordered retry, in which the sequence ID is less than (accounting for sequence wraparound) that previously seen in the slot. The replier MUST return NFS4ERR\_SEQ\_MISORDERED (as the result from SEQUENCE or CB\_SEQUENCE).
- \* A misordered new request, in which the sequence ID is two or more than (accounting for sequence wraparound) that previously seen in the slot. Note that because the sequence ID MUST wrap around to zero once it reaches 0xFFFFFFFF, a misordered new request and a misordered retry cannot be distinguished. Thus, the replier MUST return NFS4ERR\_SEQ\_MISORDERED (as the result from SEQUENCE or CB\_SEQUENCE).

Unlike the XID, the slot ID is always within a specific range; this has two implications. The first implication is that for a given session, the replier need only cache the results of a limited number of COMPOUND requests. The second implication derives from the first, which is that unlike XID-indexed reply caches (also known as duplicate request caches - DRCs), the slot ID-based reply cache cannot be overflowed. Through use of the sequence ID to identify retransmitted requests, the replier does not need to actually cache the request itself, reducing the storage requirements of the reply cache further. These facilities make it practical to maintain all the required entries for an effective reply cache.

As a result, the slot ID, sequence ID, and session ID take over the traditional role of the XID and source network address in the replier's reply cache implementation. This approach is considerably more portable and completely robust -- it is not subject to the reassignment of ports as clients reconnect over IP networks. In addition, the RPC XID is not used in the reply cache, enhancing robustness of the cache in the face of any rapid reuse of XIDs by the requester. While the replier does not care about the XID for the purposes of reply cache management (but the replier MUST return the same XID that was in the request), nonetheless there are considerations for the XID in NFSv4.1 that are the same as all other previous versions of NFS. The RPC XID remains in each message and needs to be formulated in NFSv4.1 requests as in any other ONC RPC request. The reasons include:

- \* The RPC layer retains its existing semantics and implementation.

- \* The requester and replier must be able to interoperate at the RPC layer, prior to the NFSv4.1 decoding of the SEQUENCE or CB\_SEQUENCE operation.
- \* If an operation is being used that does not start with SEQUENCE or CB\_SEQUENCE (e.g., BIND\_CONN\_TO\_SESSION), then the RPC XID is needed for correct operation to match the reply to the request.
- \* The SEQUENCE or CB\_SEQUENCE operation may generate an error. If so, the embedded slot ID, sequence ID, and session ID (if present) in the request will not be in the reply, and the requester has only the XID to match the reply to the request.

Given that well-formulated XIDs continue to be required, this raises the question: why do SEQUENCE and CB\_SEQUENCE replies have a session ID, slot ID, and sequence ID? Having the session ID in the reply means that the requester does not have to use the XID to look up the session ID, which would be necessary if the connection were associated with multiple sessions. Having the slot ID and sequence ID in the reply means that the requester does not have to use the XID to look up the slot ID and sequence ID. Furthermore, since the XID is only 32 bits, it is too small to guarantee the re-association of a reply with its request (See [rpc\_xid\_issues]); having session ID, slot ID, and sequence ID in the reply allows the client to validate that the reply in fact belongs to the matched request.

The SEQUENCE (and CB\_SEQUENCE) operation also carries a "highest\_slotid" value, which carries additional requester slot usage information. The requester MUST always indicate the slot ID representing the outstanding request with the highest-numbered slot value. The requester should in all cases provide the most conservative value possible, although it can be increased somewhat above the actual instantaneous usage to maintain some minimum or optimal level. This provides a way for the requester to yield unused request slots back to the replier, which in turn can use the information to reallocate resources.

The replier responds with both a new target highest\_slotid and an enforced highest\_slotid, described as follows:

- \* The target highest\_slotid is an indication to the requester of the highest\_slotid the replier wishes the requester to be using. This permits the replier to withdraw (or add) resources from a requester that has been found to not be using them, in order to more fairly share resources among a varying level of demand from other requesters. The requester must always comply with the replier's value updates, since they indicate newly established hard limits on the requester's access to session resources.

However, because of request pipelining, the requester might have active requests in flight reflecting prior values; therefore, the replier cannot immediately require the requester to comply.

- \* The enforced highest\_slotid indicates the highest slot ID the requester is permitted to use on a subsequent SEQUENCE or CB\_SEQUENCE operation. The replier's enforced highest\_slotid should be no less than the highest\_slotid the requester indicated in the SEQUENCE or CB\_SEQUENCE arguments.

A requester can be intransigent with respect to lowering its highest\_slotid argument to a Sequence operation, i.e. the requester continues to ignore the target highest\_slotid in the response to a Sequence operation, and continues to set its highest\_slotid argument to be higher than the target highest\_slotid. This can be considered particularly egregious behavior when the replier knows there are no outstanding requests with slot IDs higher than its target highest\_slotid. When faced with such intransigence, the replier is free to take more forceful action, and MAY reply with a new enforced highest\_slotid that is less than its previous enforced highest\_slotid. Thereafter, if the requester continues to send requests with a highest\_slotid that is greater than the replier's new enforced highest\_slotid, the server MAY return NFS4ERR\_BAD\_HIGH\_SLOT, unless the slot ID in the request is greater than the new enforced highest\_slotid and the request is a retry.

The replier should retain the slots it wants to retire until the requester sends a request with a highest\_slotid less than or equal to the replier's new enforced highest\_slotid.

The requester can also be intransigent with respect to sending non-retry requests that have a slot ID that exceeds the replier's highest\_slotid. Once the replier has forcibly lowered the enforced highest\_slotid, the requester is only allowed to send retries on slots that exceed the replier's highest\_slotid. If a request is received with a slot ID that is higher than the new enforced highest\_slotid, and the sequence ID is one higher than what is in the slot's reply cache, then the server can both retire the slot and return NFS4ERR\_BADSLOT (however, the server MUST NOT do one and not the other). The reason it is safe to retire the slot is because by using the next sequence ID, the requester is indicating it has received the previous reply for the slot.

- \* The requester is better off using the lowest available slot when sending a new request. This way, the replier may be able to retire slot entries faster. However, where the replier is actively adjusting its granted highest\_slotid, it will not be able



to use only the receipt of the slot ID and highest\_slotid in the request. Neither the slot ID nor the highest\_slotid used in a request may reflect the replier's current idea of the requester's session limit because the request may have been sent from the requester before the update was received. Therefore, in the downward adjustment case, the replier may have to retain a number of reply cache entries at least as large as the old value of maximum requests outstanding, until it can infer that the requester has seen a reply containing the new granted highest\_slotid. The replier can infer that the requester has seen such a reply when it receives a new request with the same slot ID as the request replied to and the next higher sequence ID.

#### 7.6.1.1. Caching of SEQUENCE and CB\_SEQUENCE Replies

When a SEQUENCE or CB\_SEQUENCE operation is successfully executed, its reply MUST always be cached. Specifically, session ID, sequence ID, and slot ID MUST be cached in the reply cache. The reply from SEQUENCE also includes the highest slot ID, target highest slot ID, and status flags. Instead of caching these values, the server MAY re-compute the values from the current state of the fore channel, session, and/or client ID as appropriate. Similarly, the reply from CB\_SEQUENCE includes a highest slot ID and target highest slot ID. The client MAY re-compute the values from the current state of the session as appropriate.

Regardless of whether or not a replier is re-computing highest slot ID, target slot ID, and status on replies to retries, the requester cannot assume that the values are being re-computed whenever it receives a reply after a retry is sent, since it has no way of knowing whether the reply it has received was sent by the replier in response to the retry or is a delayed response to the original request. Therefore, it may be the case that highest slot ID, target slot ID, or status bits may reflect the state of affairs when the request was first executed. Although acting based on such delayed information is valid, it may cause the receiver of the reply to do unneeded work. Requesters MAY choose to send additional requests to get the current state of affairs or use the state of affairs reported by subsequent requests, in preference to acting immediately on data that might be out of date.

#### 7.6.1.2. Errors from SEQUENCE and CB\_SEQUENCE

Any time SEQUENCE or CB\_SEQUENCE returns an error, the sequence ID of the slot MUST NOT change. The replier MUST NOT modify the reply cache entry for the slot whenever an error is returned from SEQUENCE or CB\_SEQUENCE.

### 7.6.1.3. Optional Reply Caching

On a per-request basis, the requester can choose to direct the replier to cache the reply to all operations after the first operation (SEQUENCE or CB\_SEQUENCE) via the `sa_cachethis` or `csa_cachethis` fields of the arguments to SEQUENCE or CB\_SEQUENCE. The reason it needs to direct the replier to cache the entire reply is that the request contains a non-idempotent operations [Chet]. Caching the reply may offer little benefit. If the reply is too large (see Section 7.6.4), it may not be cacheable anyway. Even if the reply to an idempotent request is small enough to cache, unnecessarily caching the reply slows down the server and increases RPC latency.

Whether or not the requester requests the reply to be cached has no effect on the slot processing. If the result of SEQUENCE or CB\_SEQUENCE is NFS4\_OK, then the slot's sequence ID MUST be incremented by one. If a requester does not direct the replier to cache the reply, the replier MUST do one of following:

- \* The replier can cache the entire original reply. Even though `sa_cachethis` or `csa_cachethis` is FALSE, the replier is always free to cache. It may choose this approach in order to simplify implementation.
- \* The replier enters into its reply cache a reply consisting of the original results to the SEQUENCE or CB\_SEQUENCE operation, and with the next operation in COMPOUND or CB\_COMPOUND having the error NFS4ERR\_RETRY\_UNCACHED\_REP. Thus, if the requester later retries the request, it will get NFS4ERR\_RETRY\_UNCACHED\_REP. If a replier receives a retried Sequence operation where the reply to the COMPOUND or CB\_COMPOUND was not cached, then the replier,
  - MAY return NFS4ERR\_RETRY\_UNCACHED\_REP in reply to a Sequence operation if the Sequence operation is not the first operation (granted, a requester that does so is in violation of the NFSv4.1 protocol).
  - MUST NOT return NFS4ERR\_RETRY\_UNCACHED\_REP in reply to a Sequence operation if the Sequence operation is the first operation.
- \* If the second operation is an illegal operation, or an operation that was legal in a previous minor version of NFSv4 and MUST NOT be supported in the current minor version (e.g., SETCLIENTID), the replier MUST NOT ever return NFS4ERR\_RETRY\_UNCACHED\_REP. Instead the replier MUST return NFS4ERR\_OP\_ILLEGAL or NFS4ERR\_BADXDR or NFS4ERR\_NOTSUPP as appropriate.

- \* If the second operation can result in another error status, the replier MAY return a status other than NFS4ERR\_RETRY\_UNCACHED\_REP, provided the operation is not executed in such a way that the state of the replier is changed. Examples of such error statuses include NFS4ERR\_SEQUENCE\_POS, NFS4ERR\_REQ\_TOO\_BIG, and NFS4ERR\_NOTSUPP returned for an operation that is legal but not REQUIRED in the current minor version and but is not supported by the replier or its file system.

The discussion above assumes that the retried request matches the original one. Section 7.6.1.3.1 discusses what the replier might do, and MUST do when it is aware that original and retried requests do not match. Since the replier might only cache a small amount of the information that would be required to determine whether this is a case of a false retry, the replier may send to the client any of the following responses:

- \* The cached reply to the original request. This is sent if users of the original request and retry match, and there is no evidence that there is in fact a mismatch between the original request and retry.

This can occur if the server caches the entire request and compares it to the retry but also in situations in which only a limited comparison or no comparison is possible. For details see Section 7.6.1.3.1

- \* A reply that consists only of the Sequence operation with the error NFS4ERR\_SEQ\_FALSE\_RETRY.

This is sent if the users of the original request and putative retry do not match, or, if they do, the server has sufficient data to indicate that the supposed retry does not match the original request.

- \* A reply consisting of the response to Sequence with the status NFS4\_OK, together with the second operation as it appeared in the retried request with an error of NFS4ERR\_RETRY\_UNCACHED\_REP or other error as described above.
- \* A reply that consists of the response to Sequence with the status NFS4\_OK, together with the second operation as it appeared in the original request with an error of NFS4ERR\_RETRY\_UNCACHED\_REP or other error as described above.

#### 7.6.1.3.1. False Retry

[Author Aside]: Section substantially revised to explain why false retries can occur, even though EOS is designed to avoid them. This is used as a basis for explaining the potential need for false retry detection while avoiding a level of checking that would be a performance issue.

The mechanisms described in Section 7.6 are designed to ensure that if a Sequence operation is sent and matches a request in the reply cache with the same slot ID and sequence ID then, it is a retry of that original request. However, it is possible, although quite unlikely, that servers will encounter requests where this is not the case, in which case the request is considered a "false retry".

- \* False retries can occur if the client does not implement request sequencing as described in Section 7.6.
- \* They can also occur as a result of situations in which large number of requests are aborted and considered complete, even though no response has been received by the requester. However, for this situation to result in a false retry there would have to be a sequence of over four billion such requests being processed using the same slot ID with that sequence followed by a long-delayed transmission of an abandoned request.

If a requester sent a Sequence operation with a slot ID and sequence ID that are in the reply cache but the replier detects that the retried request is not the same as the original request, including a retry that has different operations or different arguments in the operations from the original and a retry that uses a different principal in the RPC request's credential field that translates to a different user, then this is a false retry.

Given the low expected frequency of such false retries, the replier is not obligated to check for their existence although it is prudent to do so with requesters whose implementation of EOD is any way suspect or where the requests are transmitted over a network capable of delivering a request a very long time after it was sent. When the replier does detect a false retry, it is permitted (but not always obligated) to return `NFS4ERR_SEQ_FALSE_RETRY` in response to the Sequence operation when it detects a false retry.

Translations of particularly privileged user values to other users due to the lack of appropriately secure credentials, as configured on the replier, should be applied before determining whether the users are the same or different. If the replier determines the users are different between the original request and a retry, then the replier **MUST** return `NFS4ERR_SEQ_FALSE_RETRY`.

Regardless of whether such user mismatches do occur, the occurrence of false retries is an indication that the EOS logic is faulty, has not been implemented correctly, or that there is an extraordinary frequency of aborted requests. In light of this fact, there are practical limits to the information that might be saved in order to determine whether a particular request is a false retry. In the case of large requests recording the entire request might not be practical while a recording a compact form in the form of a checksum might unacceptably limit performance.

In the case of requests for which the reply is cached, comparing the operations in the cached response to those in the putative retry can serve to detect interactions with clients not properly implementing EOS or aborting requests inappropriately. In other cases, recording the operation count and the identity of the first non-SEQUENCE operation can make a simple check for false retry feasible.

If an operation of the retry is an illegal operation, or an operation that was legal in a previous minor version of NFSv4 and **MUST NOT** be supported in the current minor version (e.g., `SETCLIENTID`), the replier **MAY** return `NFS4ERR_SEQ_FALSE_RETRY` (and **MUST** do so if the users of the original request and retry differ). Otherwise, the replier **MAY** return `NFS4ERR_OP_ILLEGAL` or `NFS4ERR_BADXDR` or `NFS4ERR_NOTSUPP` as appropriate. Note that the handling is in contrast for how the replier deals with retries requests with no cached reply. The difference is due to `NFS4ERR_SEQ_FALSE_RETRY` being a valid error for only Sequence operations, whereas `NFS4ERR_RETRY_UNCACHED_REP` is a valid error for all operations except illegal operations and operations that **MUST NOT** be supported in the current minor version of NFSv4.

#### 7.6.2. Retry and Replay of Reply

Because NFSv4.1 is used on transports providing reliable delivery, retrying requests within an existing connection is unlikely to be helpful. Requesters will not normally retry a request, unless the connection it used to send the request disconnects. The requester can then reconnect and re-send the request, or it can re-send the request over a different connection that is associated with the same session, to deal with the possibility that the original connection is no longer functioning appropriately.

If the requester is a server wanting to re-send a callback operation over the backchannel of a session, the requester of course cannot reconnect because only the client can associate connections with the backchannel. The server can re-send the request over another connection that is bound to the same session's backchannel. If there is no such connection, the server is forced to indicate that the session has no backchannel by setting the SEQ4\_STATUS\_CB\_PATH\_DOWN\_SESSION flag bit in the response to the next SEQUENCE operation from the client. The client then has no option but to associate a new connection with the session (or destroy the session).

Note that it is not, in general, fatal for a requester to retry without a disconnect between the request and retry. However, in order to prevent false retries (see Section 7.6.1.3.1), the requester MUST NOT retry a request once the slot used to send that request has been used to send a new request.

Nevertheless, the retry does consume resources, especially with RDMA, where each request, retry or not, consumes a credit. Retries for no reason, especially retries sent shortly after the previous attempt, are a poor use of network bandwidth and defeat the purpose of a transport's inherent congestion control system.

A requester will normally wait for a reply to a request before using the slot for another request and MUST do so unless events such as termination of the issuing process makes it impossible to do so. If no such situation were to arise, then the protocol design would ensure no false retry situation could occur (See Section 7.6.1.3.1 for details). When it does not wait for a reply, the requester cannot be sure that using the next sequence ID for the slot chosen, as it normally does, will always be accepted. For example, suppose a requester sends a request with sequence ID 1, and does not wait for the response. The next time it uses the slot, it sends the new request with sequence ID 2. If the replier has not seen the request with sequence ID 1, then the replier is not expecting sequence ID 2, and rejects the requester's new request with NFS4ERR\_SEQ\_MISORDERED (as the result from SEQUENCE or CB\_SEQUENCE).

In light of the above, clients that do not wait for a reply before reusing the slot need to be aware of the possibility of receiving NFS4ERRR\_SEQ\_MISORDERED as a result and infer the probable existence of a request not received by the server. The client will then adjust the current sequence id sent, using successful execution as an indication that seqids on that slot are again correctly aligned.

RDMA fabrics do not guarantee that the memory handles (Steering Tags) within each RPC/RDMA "chunk" [RFC8166] are valid on a scope outside that of a single connection. Therefore, handles used by the direct operations become invalid after connection loss. The server must ensure that any RDMA operations that must be replayed from the reply cache use the newly provided handle(s) from the most recent request.

A retry might be sent while the original request is still in progress on the replier. In this case, the replier should deal with the issue by returning NFS4ERR\_DELAY as the reply to SEQUENCE or CB\_SEQUENCE operation, but implementations MAY return NFS4ERR\_MISORDERED. Since errors from SEQUENCE and CB\_SEQUENCE are never recorded in the reply cache, this approach allows the results of the execution of the original request to be properly recorded in the reply cache (assuming that the requester specified the reply to be cached).

### 7.6.3. Resolving Server Callback Races

It is possible for server callbacks to arrive at the client before the reply from related forward channel operations. For example, a client may have been granted a delegation to a file it has opened, but the reply to the OPEN (informing the client of the granting of the delegation) may be delayed in the network. If a conflicting operation arrives at the server, it will recall the delegation using the backchannel, which may be on a different transport connection, perhaps even a different network, or even a different session associated with the same client ID.

The presence of a session between the client and server alleviates this issue. When a session is in place, each client request is uniquely identified by its { session ID, slot ID, sequence ID } triple. By the rules under which slot entries (reply cache entries) are retired, the server has knowledge whether the client has "seen" each of the server's replies. The server can therefore provide sufficient information to the client to allow it to disambiguate between an erroneous or conflicting callback race condition.

For each client operation that might result in some sort of server callback, the server needs to keep track of the { session ID, slot ID, sequence ID } triple of the client request until the slot ID retirement rules allow the server to determine that the client has, in fact, seen the server's reply. Until the time the { session ID, slot ID, sequence ID } request triple can be retired, any recalls of the associated object MUST carry an array of these referring identifiers (in the CB\_SEQUENCE operation's arguments), for the benefit of the client. After this time, it is not necessary for the server to provide this information in related callbacks, since it is certain that a race condition can no longer occur.

The CB\_SEQUENCE operation that begins each server callback carries a list of "referring" { session ID, slot ID, sequence ID } triples. If the client finds the request corresponding to the referring session ID, slot ID, and sequence ID to be currently outstanding (i.e., the server's reply has not been seen by the client), it can determine that the callback has raced the reply, and act accordingly. If the client does not find the request corresponding to the referring triple to be outstanding (including the case of a session ID referring to a destroyed session), then there is no race with respect to this triple. The server should limit the referring triples to requests that refer to those that apply to the objects referred to in the CB\_COMPOUND procedure.

The client must not simply wait forever for the expected server reply to arrive before responding to the CB\_COMPOUND that won the race, because it is possible that it will be delayed indefinitely. The client should assume the likely case that the reply will arrive within the average round-trip time for COMPOUND requests to the server, and wait that period of time. If that period of time expires, it can respond to the CB\_COMPOUND with NFS4ERR\_DELAY. There are other scenarios under which callbacks may race replies. Among them are pNFS layout recalls as described in Section 18.7.5.2.

#### 7.6.4. COMPOUND and CB\_COMPOUND Construction Issues

Very large requests and replies may pose both buffer management issues (especially with RDMA) and reply cache issues. When the session is created (Section 25.36), for each channel (fore and back), the client and server negotiate the maximum-sized request they will send or process (ca\_maxrequestsize), the maximum-sized reply they will return or process (ca\_maxresponsesize), and the maximum-sized reply they will store in the reply cache (ca\_maxresponsesize\_cached). For discussion related to the selections of appropriate values for these quantities, see Section 7.6.5

If a request exceeds ca\_maxrequestsize, the reply will have the status NFS4ERR\_REQ\_TOO\_BIG. A replier MAY return NFS4ERR\_REQ\_TOO\_BIG as the status for the first operation (SEQUENCE or CB\_SEQUENCE) in the request (which means that no operations in the request executed and that the state of the slot in the reply cache is unchanged), or it MAY opt to return it on a subsequent operation in the same COMPOUND or CB\_COMPOUND request (which means that at least one operation did execute and that the state of the slot in the reply cache does change). The replier MUST set NFS4ERR\_REQ\_TOO\_BIG on the operation that exceeds ca\_maxrequestsize.



If a reply exceeds `ca_maxresponsesize`, the reply will have the status `NFS4ERR_REP_TOO_BIG`. A replier MAY return `NFS4ERR_REP_TOO_BIG` as the status for the first operation (`SEQUENCE` or `CB_SEQUENCE`) in the request, or it MAY opt to return it on a subsequent operation (in the same `COMPOUND` or `CB_COMPOUND` reply). A replier MAY return `NFS4ERR_REP_TOO_BIG` in the reply to `SEQUENCE` or `CB_SEQUENCE`, even if the response would still exceed `ca_maxresponsesize`.

If `sa_cachethis` or `csa_cachethis` is `TRUE`, then the replier MUST cache a reply except if an error is returned by the `SEQUENCE` or `CB_SEQUENCE` operation (see Section 7.6.1.2). If the reply exceeds `ca_maxresponsesize_cached` (and `sa_cachethis` or `csa_cachethis` is `TRUE`), then the server MUST return `NFS4ERR_REP_TOO_BIG_TO_CACHE`. Even if `NFS4ERR_REP_TOO_BIG_TO_CACHE` (or any other error for that matter) is returned on an operation other than the first operation (`SEQUENCE` or `CB_SEQUENCE`), then the reply MUST be cached if `sa_cachethis` or `csa_cachethis` is `TRUE`. For example, if a `COMPOUND` has eleven operations, including `SEQUENCE`, the fifth operation is a `RENAME`, and the tenth operation is a `READ` for one million bytes, the server may return `NFS4ERR_REP_TOO_BIG_TO_CACHE` on the tenth operation. Since the server executed several operations, especially the non-idempotent `RENAME`, the client's request to cache the reply needs to be honored in order for the correct operation of exactly once semantics. If the client retries the request, the server will have cached a reply that contains results for ten of the eleven requested operations, with the tenth operation having a status of `NFS4ERR_REP_TOO_BIG_TO_CACHE`.

A client needs to take care that, when sending operations that change the current filehandle (except for `PUTFH`, `PUTPUBFH`, `PUTROOTFH`, and `RESTOREFH`), it does not exceed the maximum reply buffer before the `GETFH` operation. Otherwise, the client will have to retry the operation that changed the current filehandle, in order to obtain the desired filehandle. For the `OPEN` operation (See Section 25.16), retry is not always available as an option. The following guidelines for the handling of filehandle-changing operations are advised:

- \* Within the same `COMPOUND` procedure, a client should send `GETFH` immediately after a current filehandle-changing operation. A client MUST send `GETFH` after a current filehandle-changing operation that is also non-idempotent (e.g., the `OPEN` operation), unless the operation is `RESTOREFH`. `RESTOREFH` is an exception, because even though it is non-idempotent, the filehandle `RESTOREFH` produced originated from an operation that is either idempotent (e.g., `PUTFH`, `LOOKUP`), or non-idempotent (e.g., `OPEN`, `CREATE`). If the origin is non-idempotent, then because the client MUST send `GETFH` after the origin operation, the client can recover if `RESTOREFH` returns an error.

- \* A server MAY return NFS4ERR\_REP\_TOO\_BIG or NFS4ERR\_REP\_TOO\_BIG\_TO\_CACHE (if sa\_cachethis is TRUE) on a filehandle-changing operation if the reply would be too large on the next operation.
- \* A server MUST return NFS4ERR\_REP\_TOO\_BIG or NFS4ERR\_REP\_TOO\_BIG\_TO\_CACHE (if sa\_cachethis is TRUE) on a filehandle-changing, non-idempotent operation if the reply would be too large on the next operation, especially if the operation is OPEN.
- \* A server MAY return NFS4ERR\_UNSAFE\_COMPOUND to a non-idempotent current filehandle-changing operation, if it looks at the next operation (in the same COMPOUND procedure) and finds it is not GETFH. The server MAY do this if it is unable to determine in advance whether the total response size would exceed ca\_maxresponsesize\_cached or ca\_maxresponsesize.

#### 7.6.5. Setting Size limits for Sessions

The following issues need to be taken account of in setting the size limits introduced in Section 7.6.4.

- \* The primary factor governing the choice of ca\_maxrequestsize is the need to make WRITE requests of sufficient size, so that IO overhead does not become excessive. In addition. SETATTR requests may be long and require a high ca\_maxrequestsize when they are used to set attributes such as acl, sacl and dacl, which can be of substantial size.

In both cases, allowances also needs to be made for smaller operations that will need to be included in the COMPOUND.

- \* The primary factor governing the choice of ca\_maxresponsesize is the need to make READ requests of sufficient size, so that IO overhead does not become excessive. In addition. GETATTR request may be long and require a high ca\_maxresponsesize when they are used to get the value attributes such as acl, sacl and dacl, which can be of substantial size.

In both cases, allowances also need to be made for smaller operation that will need to be included in the COMPOUND.

- \* In determining the appropriate value for `ca_maxresponsesize_cached`, the response sizes of idempotent requests such as `READ` and `GETATTR` can be ignored unless the client needs to issue `COMPOUNDS` that combine non-idempotent operation together with operations that return long responses. When a client chooses to do this, a large `ca_maxresponsesize_cached` is likely to be needed.

When new operations are proposed the effect on the requirements on `ca_maxresponsesize_cached` needs to be considered since it might be impractical to require large reply caches.

## 7.7. RDMA Considerations

A complete discussion of the operation of RPC-based protocols over RDMA transports is in [RFC8166]. A discussion of the operation of NFSv4, including NFSv4.1, over RDMA is in [RFC8267]. Where RDMA is considered, this specification assumes the use of such a layering; it addresses only the upper-layer issues relevant to making best use of RPC/RDMA.

### 7.7.1. RDMA Connection Resources

RDMA requires its consumers to register memory and post buffers of a specific size and number for receive operations.

Registration of memory can be a relatively high-overhead operation, since it requires pinning of buffers, assignment of attributes (e.g., readable/writable), and initialization of hardware translation. Preregistration is desirable to reduce overhead. These registrations are specific to hardware interfaces and even to RDMA connection endpoints; therefore, negotiation of their limits is desirable to manage resources effectively.

Following basic registration, these buffers must be posted by the RPC layer to handle receives. These buffers remain in use by the RPC/NFSv4.1 implementation; the size and number of them must be known to the remote peer in order to avoid RDMA errors that would cause a fatal error on the RDMA connection.

NFSv4.1 manages slots as resources on a per-session basis (See Section 7), while RDMA connections manage credits on a per-connection basis. This means that in order for a peer to send data over RDMA to a remote buffer, it has to have both an NFSv4.1 slot and an RDMA credit. If multiple RDMA connections are associated with a session, then if the total number of credits across all RDMA connections associated with the session is  $X$ , and the number of slots in the session is  $Y$ , then the maximum number of outstanding requests is the lesser of  $X$  and  $Y$ .

#### 7.7.2. Flow Control

Previous versions of NFS do not provide flow control; instead, they rely on the windowing provided by transports like TCP to throttle requests. This does not work with RDMA, which provides no operation flow control and will terminate a connection in error when limits are exceeded. Limits such as maximum number of requests outstanding are therefore negotiated when a session is created (See the `ca_maxrequests` field in Section 25.36). These limits then provide the maxima within which each connection associated with the session's channel(s) must remain. RDMA connections are managed within these limits as described in Section 3.3 of [RFC8166]; if there are multiple RDMA connections, then the maximum number of requests for a channel will be divided among the RDMA connections. Put a different way, the onus is on the replier to ensure that the total number of RDMA credits across all connections associated with the replier's channel does exceed the channel's maximum number of outstanding requests.

The limits may also be modified dynamically at the replier's choosing by manipulating certain parameters present in each NFSv4.1 reply. In addition, the `CB_RECALL_SLOT` callback operation (see Section 27.8) can be sent by a server to a client to return RDMA credits to the server, thereby lowering the maximum number of requests a client can have outstanding to the server.

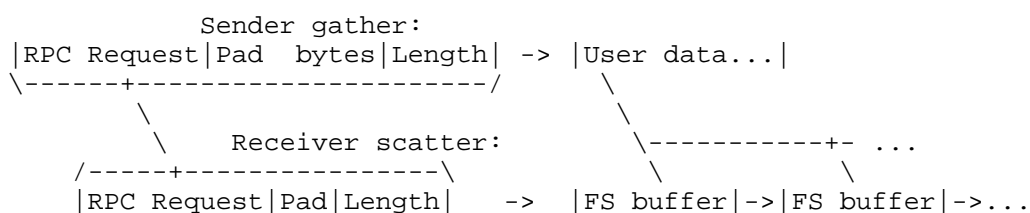
#### 7.7.3. Padding

Header padding is requested by each peer at session initiation (See the `ca_headerpadsize` argument to `CREATE_SESSION` in Section 25.36), and subsequently used by the RPC RDMA layer, as described in [RFC8166]. Zero padding is permitted.

Padding leverages the useful property that RDMA preserve alignment of data, even when they are placed into anonymous (untagged) buffers. If requested, client inline writes will insert appropriate pad bytes within the request header to align the data payload on the specified boundary. The client is encouraged to add sufficient padding (up to

the negotiated size) so that the "data" field of the WRITE operation is aligned. Most servers can make good use of such padding, which allows them to chain receive buffers in such a way that any data carried by client requests will be placed into appropriate buffers at the server, ready for file system processing. The receiver's RPC layer encounters no overhead from skipping over pad bytes, and the RDMA layer's high performance makes the insertion and transmission of padding on the sender a significant optimization. In this way, the need for servers to perform RDMA Read to satisfy all but the largest client writes is obviated. An added benefit is the reduction of message round trips on the network -- a potentially good trade, where latency is present.

The value to choose for padding is subject to a number of criteria. A primary source of variable-length data in the RPC header is the authentication information, the form of which is client-determined, possibly in response to server specification. The contents of COMPOUNDS, sizes of strings such as those passed to RENAME, etc. all go into the determination of a maximal NFSv4.1 request size and therefore minimal buffer size. The client must select its offered value carefully, so as to avoid overburdening the server, and vice versa. The benefit of an appropriate padding value is higher performance.



In the above case, the server may recycle unused buffers to the next posted receive if unused by the actual received request, or may pass the now-complete buffers by reference for normal write processing. For a server that can make use of it, this removes any need for data copies of incoming data, without resorting to complicated end-to-end buffer advertisement and management. This includes most kernel-based and integrated server designs, among many others. The client may perform similar optimizations, if desired.

#### 7.7.4. Dual RDMA and Non-RDMA Transports

Some RDMA transports (e.g. [RFC5040]) permit a "streaming" (non-RDMA) phase, where ordinary traffic might flow before "stepping up" to RDMA mode, commencing RDMA traffic. Some RDMA transports start connections always in RDMA mode. NFSv4.1 allows, but does not assume, a streaming phase before RDMA mode. When a connection is associated with a session, the client and server negotiate whether the connection is used in RDMA or non-RDMA mode (See Sections 25.36 and 25.34).

### 7.8. Session Security

#### 7.8.1. Session Callback Security

Via session/connection association, NFSv4.1 improves security over that provided by NFSv4.0 for the backchannel. The connection is client-initiated (see Section 25.34) and subject to the same firewall and routing checks as the fore channel. At the client's option (See Section 25.35), connection association is fully authenticated before being activated (See Section 25.34). Traffic from the server over the backchannel is authenticated exactly as the client specifies (See Section 7.8.2).

#### 7.8.2. Backchannel RPC Security

When the NFSv4.1 client establishes the backchannel, it informs the server of the security flavors and principals to use when sending requests. If the security flavor is RPCSEC\_GSS, the client expresses the principal in the form of an established RPCSEC\_GSS context. The server is free to use any of the flavor/principal combinations the client offers, but it MUST NOT use combinations not offered. This way, the client need not provide a target GSS principal for the backchannel as it did with NFSv4.0, nor does the server have to implement an RPCSEC\_GSS initiator as it did with NFSv4.0 [RFC3530].

The CREATE\_SESSION (Section 25.36) and BACKCHANNEL\_CTL (Section 25.33) operations allow the client to specify flavor/principal combinations.

Also note that the SP4\_SSV state protection mode (See Sections 25.35 and 7.8.3) has the side benefit of providing SSV-derived RPCSEC\_GSS contexts (Section 7.9).

### 7.8.3. Protection from Unauthorized State Changes

As described to this point in the specification, the state model of NFSv4.1 is vulnerable to an attacker that sends a SEQUENCE operation with a forged session ID and with a slot ID that it expects the legitimate client to use next. When the legitimate client uses the slot ID with the same sequence number, the server returns the attacker's result from the reply cache, which disrupts the legitimate client and thus denies service to it. Similarly, an attacker could send a CREATE\_SESSION with a forged client ID to create a new session associated with the client ID. The attacker could send requests using the new session that change locking state, such as LOCKU operations to release locks the legitimate client has acquired. Setting a security policy on the file that requires RPCSEC\_GSS credentials when manipulating the file's state is one potential work around, but has the disadvantage of preventing a legitimate client from releasing state when RPCSEC\_GSS is required to do so, but a GSS context cannot be obtained (possibly because the user has logged off the client).

NFSv4.1 provides three options to a client for state protection, which are specified when a client creates a client ID via EXCHANGE\_ID (Section 25.35).

The first (SP4\_NONE) is to simply waive state protection, except for that provided by client host authentication.

The other two options (SP4\_MACH\_CRED and SP4\_SSV) share several traits:

- \* An RPCSEC\_GSS-based credential is used to authenticate client ID and session maintenance operations, including creating and destroying a session, associating a connection with the session, and destroying the client ID.
- \* Because RPCSEC\_GSS is used to authenticate client ID and session maintenance, the attacker cannot associate a rogue connection with a legitimate session, or associate a rogue session with a legitimate client ID in order to maliciously alter the client ID's lock state via CLOSE, LOCKU, DELEGRETURN, LAYOUTRETURN, etc.
- \* In cases where the server's security policies on a portion of its namespace require RPCSEC\_GSS authentication, a client may have to use an RPCSEC\_GSS credential to remove per-file state (e.g., LOCKU, CLOSE, etc.). The server may require that the principal that removes the state match certain criteria (e.g., the principal might have to be the same as the one that acquired the state). However, the client might not have an RPCSEC\_GSS context for such

a principal, and might not be able to create such a context (perhaps because the user has logged off). When the client establishes SP4\_MACH\_CRED or SP4\_SSV protection, it can specify a list of operations that the server MUST allow using the machine credential (if SP4\_MACH\_CRED is used) or the SSV credential (if SP4\_SSV is used).

The SP4\_MACH\_CRED state protection option uses a machine credential where the principal that creates the client ID MUST also be the principal that performs client ID and session maintenance operations. The security of the machine credential state protection approach depends entirely on safeguarding the per-machine credential. Assuming a proper safeguard using the per-machine credential for operations like CREATE\_SESSION, BIND\_CONN\_TO\_SESSION, DESTROY\_SESSION, and DESTROY\_CLIENTID will prevent an attacker from associating a rogue connection with a session, or associating a rogue session with a client ID.

There are at least three scenarios for the SP4\_MACH\_CRED option:

1. The system administrator configures a unique, permanent per-machine credential for one of the mandated GSS mechanisms (e.g., if Kerberos V5 is used, a "keytab" containing a principal derived from a client host name could be used).
2. The client is used by a single user, and so the client ID and its sessions are used by just that user. If the user's credential expires, then session and client ID maintenance cannot occur, but since the client has a single user, only that user is inconvenienced.
3. The physical client has multiple users, but the client implementation has a unique client ID for each user. This is effectively the same as the second scenario, but a disadvantage is that each user needs to be allocated at least one session each, so the approach suffers from lack of economy.



The SP4\_SSV protection option uses the SSV (Section 2.5), via RPCSEC\_GSS and the SSV GSS mechanism (Section 7.9), to protect state from attack. The SP4\_SSV protection option is intended for the situation comprised of a client that has multiple active users and a system administrator who wants to avoid the burden of installing a permanent machine credential on each client. The SSV is established and updated on the server via SET\_SSV (See Section 25.47). To prevent eavesdropping, a client should send SET\_SSV via RPCSEC\_GSS with the privacy service or use tls encryption on the connection making the request. Several aspects of the SSV make it intractable for an attacker to guess the SSV, and thus associate rogue connections with a session, and rogue sessions with a client ID:

- \* The arguments to and results of SET\_SSV include digests of the old and new SSV, respectively.
- \* Because the initial value of the SSV is zero, therefore known, the client that opts for SP4\_SSV protection and opts to apply SP4\_SSV protection to BIND\_CONN\_TO\_SESSION and CREATE\_SESSION MUST send at least one SET\_SSV operation before the first BIND\_CONN\_TO\_SESSION operation or before the second CREATE\_SESSION operation on a client ID. If it does not, the SSV mechanism will not generate tokens (Section 7.9). A client should send SET\_SSV as soon as a session is created.
- \* A SET\_SSV request does not replace the SSV with the argument to SET\_SSV. Instead, the current SSV on the server is logically exclusive Ored (XORed) with the argument to SET\_SSV. For the protection to be in effect, each time a new principal uses a client ID for the first time, the client MUST send a SET\_SSV with that principal's RPCSEC\_GSS credentials, with RPCSEC\_GSS service set to RPC\_GSS\_SVC\_PRIVACY.

Here are the types of attacks that can be attempted by an attacker named Eve on a victim named Bob, and how SP4\_SSV protection foils each attack:

- \* Suppose Eve is the first user to log into a legitimate client. Eve's use of an NFSv4.1 file system will cause the legitimate client to create a client ID with SP4\_SSV protection, specifying that the BIND\_CONN\_TO\_SESSION operation MUST use the SSV credential. Eve's use of the file system also causes an SSV to be created. The SET\_SSV operation that creates the SSV will be protected by the RPCSEC\_GSS context created by the legitimate client, which uses Eve's GSS principal and credentials. Eve can eavesdrop on the network while her RPCSEC\_GSS context is created and the SET\_SSV using her context is sent. Even if the legitimate client sends the SET\_SSV with RPC\_GSS\_SVC\_PRIVACY, because Eve

knows her own credentials, she can decrypt the SSV. Eve can compute an RPCSEC\_GSS credential that BIND\_CONN\_TO\_SESSION will accept, and so associate a new connection with the legitimate session. Eve can change the slot ID and sequence state of a legitimate session, and/or the SSV state, in such a way that when Bob accesses the server via the same legitimate client, the legitimate client will be unable to use the session.

The client's only recourse is to create a new client ID for Bob to use, and establish a new SSV for the client ID. The client will be unable to delete the old client ID, and will let the lease on the old client ID expire.

Once the legitimate client establishes an SSV over the new session using Bob's RPCSEC\_GSS context, Eve can use the new session via the legitimate client, but she cannot disrupt Bob. Moreover, because the client should have modified the SSV due to Eve using the new session, Bob cannot get revenge on Eve by associating a rogue connection with the session.

The question is how did the legitimate client detect that Eve has hijacked the old session? When the client detects that a new principal, Bob, wants to use the session, it should have sent a SET\_SSV, which leads to the following sub-scenarios:

- Let us suppose that from the rogue connection, Eve sent a SET\_SSV with the same slot ID and sequence ID that the legitimate client later uses. The server will assume the SET\_SSV sent with Bob's credentials is a retry, and return to the legitimate client the reply it sent Eve. However, unless Eve can correctly guess the SSV the legitimate client will use, the digest verification checks in the SET\_SSV response will fail. That is an indication to the client that the session has apparently been hijacked.
- Alternatively, Eve sent a SET\_SSV with a different slot ID than the legitimate client uses for its SET\_SSV. Then the digest verification of the SET\_SSV sent with Bob's credentials fails on the server, and the error returned to the client makes it apparent that the session has been hijacked.
- Alternatively, Eve sent an operation other than SET\_SSV, but with the same slot ID and sequence that the legitimate client uses for its SET\_SSV. The server returns to the legitimate client the response it sent Eve. The client sees that the response is not at all what it expects. The client assumes either session hijacking or a server bug, and either way destroys the old session.

- \* Eve associates a rogue connection with the session as above, and then destroys the session. Again, Bob goes to use the server from the legitimate client, which sends a SET\_SSV using Bob's credentials. The client receives an error that indicates that the session does not exist. When the client tries to create a new session, this will fail because the SSV it has does not match that which the server has, and now the client knows the session was hijacked. The legitimate client establishes a new client ID.
- \* If Eve creates a connection before the legitimate client establishes an SSV, because the initial value of the SSV is zero and therefore known, Eve can send a SET\_SSV that will pass the digest verification check. However, because the new connection has not been associated with the session, the SET\_SSV is rejected for that reason.

In summary, an attacker's disruption of state when SP4\_SSV protection is in use is limited to the formative period of a client ID, its first session, and the establishment of the SSV. Once a non-malicious user uses the client ID, the client quickly detects any hijack and rectifies the situation. Once a non-malicious user successfully modifies the SSV, the attacker cannot use NFSv4.1 operations to disrupt the non-malicious user.

Note that neither the SP4\_MACH\_CRED nor SP4\_SSV protection approaches prevent hijacking of a transport connection that has previously been associated with a session. If the goal of a counter-threat strategy is to prevent connection hijacking, the use of IPsec or TLS is RECOMMENDED.

If a connection hijack occurs, the hijacker could in theory change locking state and negatively impact the service to legitimate clients. However, if the server is configured to require the use of RPCSEC\_GSS with integrity or privacy on the affected file objects, and if EXCHGID4\_FLAG\_BIND\_PRINC\_STATEID capability (Section 25.35) is in force, this will thwart unauthorized attempts to change locking state.

#### 7.9. The Secret State Verifier (SSV) GSS Mechanism

The SSV provides the secret key for a GSS mechanism internal to NFSv4.1 that NFSv4.1 uses for state protection. Contexts for this mechanism are not established via the RPCSEC\_GSS protocol. Instead, the contexts are automatically created when EXCHANGE\_ID specifies SP4\_SSV protection. The only tokens defined are the PerMsgToken (emitted by GSS\_GetMIC) and the SealedMessage token (emitted by GSS\_Wrap).

The mechanism OID for the SSV mechanism is `iso.org.dod.internet.private.enterprise.Michael Eisler.nfs.ssv_mech (1.3.6.1.4.1.28882.1.1)`. While the SSV mechanism does not define any initial context tokens, the OID can be used to let servers indicate that the SSV mechanism is acceptable whenever the client sends a `SECINFO` or `SECINFO_NO_NAME` operation (see Section 6.2).

The SSV mechanism defines four subkeys derived from the SSV value. Each time `SET_SSV` is invoked, the subkeys are recalculated by the client and server. The calculation of each of the four subkeys depends on each of the four respective `ssv_subkey4` enumerated values. The calculation uses the HMAC [RFC2104] algorithm, using the current SSV as the key, the one-way hash algorithm as negotiated by `EXCHANGE_ID`, and the input text as represented by the XDR encoded enumeration value for that subkey of data type `ssv_subkey4`. If the length of the output of the HMAC algorithm exceeds the length of key of the encryption algorithm (which is also negotiated by `EXCHANGE_ID`), then the subkey MUST be truncated from the HMAC output, i.e., if the subkey is of N bytes long, then the first N bytes of the HMAC output MUST be used for the subkey. The specification of `EXCHANGE_ID` states that the length of the output of the HMAC algorithm MUST NOT be less than the length of subkey needed for the encryption algorithm (See Section 25.35).

```
/* Input for computing subkeys */
enum ssv_subkey4 {
    SSV4_SUBKEY_MIC_I2T      = 1,
    SSV4_SUBKEY_MIC_T2I      = 2,
    SSV4_SUBKEY_SEAL_I2T     = 3,
    SSV4_SUBKEY_SEAL_T2I     = 4
};
```

The subkey derived from `SSV4_SUBKEY_MIC_I2T` is used for calculating message integrity codes (MICs) that originate from the NFSv4.1 client, whether as part of a request over the fore channel or a response over the backchannel. The subkey derived from `SSV4_SUBKEY_MIC_T2I` is used for MICs originating from the NFSv4.1 server. The subkey derived from `SSV4_SUBKEY_SEAL_I2T` is used for encryption text originating from the NFSv4.1 client, and the subkey derived from `SSV4_SUBKEY_SEAL_T2I` is used for encryption text originating from the NFSv4.1 server.

The `PerMsgToken` description is based on an XDR definition:

```
/* Input for computing smt_hmac */
struct ssv_mic_plain_tkn4 {
    uint32_t      smpt_ssv_seq;
    opaque        smpt_orig_plain<>;
};

/* SSV GSS PerMsgToken token */
struct ssv_mic_tkn4 {
    uint32_t      smt_ssv_seq;
    opaque        smt_hmac<>;
};
```

The field `smt_hmac` is an HMAC calculated by using the subkey derived from `SSV4_SUBKEY_MIC_I2T` or `SSV4_SUBKEY_MIC_T2I` as the key, the one-way hash algorithm as negotiated by `EXCHANGE_ID`, and the input text as represented by data of type `ssv_mic_plain_tkn4`. The field `smpt_ssv_seq` is the same as `smt_ssv_seq`. The field `smpt_orig_plain` is the "message" input passed to `GSS_GetMIC()` (See Section 2.3.1 of [RFC2743]). The caller of `GSS_GetMIC()` provides a pointer to a buffer containing the plain text. The SSV mechanism's entry point for `GSS_GetMIC()` encodes this into an opaque array, and the encoding will include an initial four-byte length, plus any necessary padding. Prepended to this will be the XDR encoded value of `smpt_ssv_seq`, thus making up an XDR encoding of a value of data type `ssv_mic_plain_tkn4`, which in turn is the input into the HMAC.

The token emitted by `GSS_GetMIC()` is XDR encoded and of XDR data type `ssv_mic_tkn4`. The field `smt_ssv_seq` comes from the SSV sequence number, which is equal to one after `SET_SSV` (Section 25.47) is called the first time on a client ID. Thereafter, the SSV sequence number is incremented on each `SET_SSV`. Thus, `smt_ssv_seq` represents the version of the SSV at the time `GSS_GetMIC()` was called. As noted in Section 25.35, the client and server can maintain multiple concurrent versions of the SSV. This allows the SSV to be changed without serializing all RPC calls that use the SSV mechanism with `SET_SSV` operations. Once the HMAC is calculated, it is XDR encoded into `smt_hmac`, which will include an initial four-byte length, and any necessary padding. Prepended to this will be the XDR encoded value of `smt_ssv_seq`.

The `SealedMessage` description is based on an XDR definition:

```
/* Input for computing ssct_encr_data and ssct_hmac */
struct ssv_seal_plain_tkn4 {
    opaque          sspt_confounder<>;
    uint32_t        sspt_ssv_seq;
    opaque          sspt_orig_plain<>;
    opaque          sspt_pad<>;
};

/* SSV GSS SealedMessage token */
struct ssv_seal_cipher_tkn4 {
    uint32_t        ssct_ssv_seq;
    opaque          ssct_iv<>;
    opaque          ssct_encr_data<>;
    opaque          ssct_hmac<>;
};
```

The token emitted by GSS\_Wrap() is XDR encoded and of XDR data type ssv\_seal\_cipher\_tkn4.

The ssct\_ssv\_seq field has the same meaning as smt\_ssv\_seq.

The ssct\_encr\_data field is the result of encrypting a value of the XDR encoded data type ssv\_seal\_plain\_tkn4. The encryption key is the subkey derived from SSV4\_SUBKEY\_SEAL\_I2T or SSV4\_SUBKEY\_SEAL\_T2I, and the encryption algorithm is that negotiated by EXCHANGE\_ID.

The ssct\_iv field is the initialization vector (IV) for the encryption algorithm (if applicable) and is sent in clear text. The content and size of the IV MUST comply with the specification of the encryption algorithm. For example, the id-aes256-CBC algorithm MUST use a 16-byte initialization vector (IV), which MUST be unpredictable for each instance of a value of data type ssv\_seal\_plain\_tkn4 that is encrypted with a particular SSV key.

The ssct\_hmac field is the result of computing an HMAC using the value of the XDR encoded data type ssv\_seal\_plain\_tkn4 as the input text. The key is the subkey derived from SSV4\_SUBKEY\_MIC\_I2T or SSV4\_SUBKEY\_MIC\_T2I, and the one-way hash algorithm is that negotiated by EXCHANGE\_ID.

The sspt\_confounder field is a random value.

The sspt\_ssv\_seq field is the same as ssvt\_ssv\_seq.

The field sspt\_orig\_plain field is the original plaintext and is the "input\_message" input passed to GSS\_Wrap() (See Section 2.3.3 of [RFC2743]). As with the handling of the plaintext by the SSV mechanism's GSS\_GetMIC() entry point, the entry point for GSS\_Wrap()

expects a pointer to the plaintext, and will XDR encode an opaque array into `sspt_orig_plain` representing the plain text, along with the other fields of an instance of data type `ssv_seal_plain_tkn4`.

The `sspt_pad` field is present to support encryption algorithms that require inputs to be in fixed-sized blocks. The content of `sspt_pad` is zero filled except for the length. Beware that the XDR encoding of `ssv_seal_plain_tkn4` contains three variable-length arrays, and so each array consumes four bytes for an array length, and each array that follows the length is always padded to a multiple of four bytes per the XDR standard.

For example, suppose the encryption algorithm uses 16-byte blocks, and the `sspt_confounder` is three bytes long, and the `sspt_orig_plain` field is 15 bytes long. The XDR encoding of `sspt_confounder` uses eight bytes (4 + 3 + 1-byte pad), the XDR encoding of `sspt_ssv_seq` uses four bytes, the XDR encoding of `sspt_orig_plain` uses 20 bytes (4 + 15 + 1-byte pad), and the smallest XDR encoding of the `sspt_pad` field is four bytes. This totals 36 bytes. The next multiple of 16 is 48; thus, the length field of `sspt_pad` needs to be set to 12 bytes, or a total encoding of 16 bytes. The total number of XDR encoded bytes is thus  $8 + 4 + 20 + 16 = 48$ .

`GSS_Wrap()` emits a token that is an XDR encoding of a value of data type `ssv_seal_cipher_tkn4`. Note that regardless of whether or not the caller of `GSS_Wrap()` requests confidentiality, the token always has confidentiality. This is because the SSV mechanism is for `RPCSEC_GSS`, and `RPCSEC_GSS` never produces `GSS_wrap()` tokens without confidentiality.

There is one SSV per client ID. There is a single GSS context for a client ID / SSV pair. All SSV mechanism `RPCSEC_GSS` handles of a client ID / SSV pair share the same GSS context. SSV GSS contexts do not expire except when the SSV is destroyed (causes would include the client ID being destroyed or a server restart). Since one purpose of context expiration is to replace keys that have been in use for "too long", hence vulnerable to compromise by brute force or accident, the client can replace the SSV key by sending periodic `SET_SSV` operations, which is done by cycling through different users' `RPCSEC_GSS` credentials. This way, the SSV is replaced without destroying the SSV's GSS contexts.

SSV `RPCSEC_GSS` handles can be expired or deleted by the server at any time, and the `EXCHANGE_ID` operation can be used to create more SSV `RPCSEC_GSS` handles. Expiration of SSV `RPCSEC_GSS` handles does not imply that the SSV or its GSS context has expired.

The client MUST establish an SSV via SET\_SSV before the SSV GSS context can be used to emit tokens from GSS\_Wrap() and GSS\_GetMIC(). If SET\_SSV has not been successfully called, attempts to emit tokens MUST fail.

The SSV mechanism does not support replay detection and sequencing in its tokens because RPCSEC\_GSS does not use those features (see "Context Creation Requests", Section 5.2.2 of [RFC2203]). However, Section 7.10 discusses special considerations for the SSV mechanism when used with RPCSEC\_GSS.

#### 7.10. Security Considerations for RPCSEC\_GSS When Using the SSV Mechanism

When a client ID is created with SP4\_SSV state protection (See Section 25.35), the client is permitted to associate multiple RPCSEC\_GSS handles with the single SSV GSS context (See Section 7.9). Because of the way RPCSEC\_GSS (both version 1 and version 2, see [RFC2203] and [RFC5403]) calculate the verifier of the reply, special care must be taken by the implementation of the NFSv4.1 client to prevent attacks by a man-in-the-middle. The verifier of an RPCSEC\_GSS reply is the output of GSS\_GetMIC() applied to the input value of the seq\_num field of the RPCSEC\_GSS credential (data type rpc\_gss\_cred\_ver\_1\_t) (See Section 5.3.3.2 of [RFC2203]). If multiple RPCSEC\_GSS handles share the same GSS context, then if one handle is used to send a request with the same seq\_num value as another handle, an attacker could block the reply, and replace it with the verifier used for the other handle.

There are multiple ways to prevent the attack on the SSV RPCSEC\_GSS verifier in the reply. The simplest is believed to be as follows.

- \* Each time one or more new SSV RPCSEC\_GSS handles are created via EXCHANGE\_ID, the client sends a SET\_SSV operation to modify the SSV. By changing the SSV, the new handles will not result in the re-use of an SSV RPCSEC\_GSS verifier in a reply.
- \* When a requester decides to use N SSV RPCSEC\_GSS handles, it assigns a unique and non-overlapping range of seq\_nums to each SSV RPCSEC\_GSS handle. The size of each range needs to be equal to MAXSEQ / N (See Section 5 of [RFC2203] for the definition of MAXSEQ). When an SSV RPCSEC\_GSS handle reaches its maximum, it should force the replier to destroy the handle by sending a NULL RPC request with seq\_num set to MAXSEQ + 1 (See Section 5.3.3.3 of [RFC2203]).



- \* When the requester wants to increase or decrease N, it forces the replier to destroy all N handles by sending a NULL RPC request on each handle with seq\_num set to MAXSEQ + 1. If the requester is the client, it needs to send a SET\_SSV operation before using new handles. If the requester is the server, then the client needs to send a SET\_SSV operation when it detects that the server has forced it to destroy a backchannel's SSV RPCSEC\_GSS handle. By sending a SET\_SSV operation, the SSV will change, and so the attacker will be unavailable to successfully replay a previous verifier in a reply to the requester.

Note that if the replier carefully creates the SSV RPCSEC\_GSS handles, the related risk of a man-in-the-middle splicing a forged SSV RPCSEC\_GSS credential with a verifier for another handle does not exist. This is because the verifier in an RPCSEC\_GSS request is computed from input that includes both the RPCSEC\_GSS handle and seq\_num (See Section 5.3.1 of [RFC2203]). Provided the replier takes care to avoid re-using the value of an RPCSEC\_GSS handle that it creates, such as by including a generation number in the handle, the man-in-the-middle will not be able to successfully replay a previous verifier in the request to a replier.

## 7.11. Session Mechanics - Steady State

### 7.11.1. Obligations of the Server

The server has the primary obligation to monitor the state of backchannel resources that the client has created for the server (RPCSEC\_GSS contexts and backchannel connections). If these resources vanish, the server takes action as specified in Section 7.13.2.

### 7.11.2. Obligations of the Client

The client needs to honor the following obligations in order to utilize the session:

- \* Keep a necessary session from going idle on the server. A client that requires a session but nonetheless is not sending operations risks having the session be destroyed by the server. This is because sessions consume resources, and resource limitations may force the server to cull an inactive session. A server MAY consider a session to be inactive if the client has not used the session before the session inactivity timer (Section 7.12) has expired.

- \* Destroy the session when not needed. If a client has multiple sessions, one of which has no requests waiting for replies, and has been idle for some period of time, it should destroy the session.
- \* Maintain GSS contexts and RPCSEC\_GSS handles for the backchannel. If the client requires the server to use the RPCSEC\_GSS security flavor for callbacks, then it needs to be sure the RPCSEC\_GSS handles and/or their GSS contexts that are handed to the server via BACKCHANNEL\_CTL or CREATE\_SESSION are unexpired.
- \* Preserve a connection for a backchannel. The server requires a backchannel in order to gracefully recall recallable state or notify the client of certain events. Note that if the connection is not being used for the fore channel, there is no way for the client to tell if the connection is still alive (e.g., the server restarted without sending a disconnect). The onus is on the server, not the client, to determine if the backchannel's connection is alive, and to indicate in the response to a SEQUENCE operation when the last connection associated with a session's backchannel has disconnected.

#### 7.11.3. Steps the Client Takes to Establish a Session

If the client does not have a client ID, the client sends EXCHANGE\_ID to establish a client ID. If it opts for SP4\_MACH\_CRED or SP4\_SSV protection, in the spo\_must\_enforce list of operations, it need to, at a minimum, specify CREATE\_SESSION, DESTROY\_SESSION, BIND\_CONN\_TO\_SESSION, BACKCHANNEL\_CTL, and DESTROY\_CLIENTID. If it opts for SP4\_SSV protection, the client needs to ask for SSV-based RPCSEC\_GSS handles.

The client uses the client ID to send a CREATE\_SESSION on a connection to the server. The results of CREATE\_SESSION indicate whether or not the server undertakes to persist the session reply cache in which a server restarts, and the client notes this for future reference.

If the client specified SP4\_SSV state protection when the client ID was created, then it needs to then it send SET\_SSV in the first COMPOUND after the session is created. Each time a new principal goes to use the client ID, it needs to send a SET\_SSV again.

If the client wants to use delegations, layouts, directory notifications, or any other state that requires a backchannel, then it needs to add a connection to the backchannel if CREATE\_SESSION did not already do so. The client creates a connection, and calls BIND\_CONN\_TO\_SESSION to associate the connection with the session and

the session's backchannel. If CREATE\_SESSION did not already do so, the client MUST tell the server what security is required in order for the client to accept callbacks. The client does this via BACKCHANNEL\_CTL. If the client selected SP4\_MACH\_CRED or SP4\_SSV protection when it called EXCHANGE\_ID, then the client needs to specify that the backchannel use RPCSEC\_GSS contexts for security.

If the client wants to use additional connections for the backchannel, then it needs to call BIND\_CONN\_TO\_SESSION on each connection it wants to use with the session. If the client wants to use additional connections for the fore channel, then it needs to call BIND\_CONN\_TO\_SESSION if it specified SP4\_SSV or SP4\_MACH\_CRED state protection when the client ID was created.

At this point, the session has reached steady state.

#### 7.12. Session Inactivity Timer

The server MAY maintain a session inactivity timer for each session. If the session inactivity timer expires, then the server MAY destroy the session. To avoid losing a session due to inactivity, the client MUST renew the session inactivity timer. The length of session inactivity timer MUST NOT be less than the lease\_time attribute (Section 11.12.1.11). As with lease renewal (Section 13.3), when the server receives a SEQUENCE operation, it resets the session inactivity timer, and MUST NOT allow the timer to expire while the rest of the operations in the COMPOUND procedure's request are still executing. Once the last operation has finished, the server MUST set the session inactivity timer to expire no sooner than the sum of the current time and the value of the lease\_time attribute.

#### 7.13. Session Mechanics - Recovery

##### 7.13.1. Events Requiring Client Action

The following events require client action to recover.

##### 7.13.1.1. RPCSEC\_GSS Context Loss by Callback Path

If all RPCSEC\_GSS handles granted by the client to the server for callback use have expired, the client MUST establish a new handle via BACKCHANNEL\_CTL. The sr\_status\_flags field of the SEQUENCE results indicates when callback handles are nearly expired, or fully expired (See Section 25.46.3).

#### 7.13.1.2. Connection Loss

If the client loses the last connection of the session and wants to retain the session, then it needs to create a new connection, and if, when the client ID was created, `BIND_CONN_TO_SESSION` was specified in the `spo_must_enforce` list, the client **MUST** use `BIND_CONN_TO_SESSION` to associate the connection with the session.

If there was a request outstanding at the time of connection loss, then if the client wants to continue to use the session, it **MUST** retry the request, as described in Section 7.6.2. Note that it is not necessary to retry requests over a connection with the same source network address or the same destination network address as the lost connection. As long as the session ID, slot ID, and sequence ID in the retry match that of the original request, the server will recognize the request as a retry if it executed the request prior to disconnect.

If the connection that was lost was the last one associated with the backchannel, and the client wants to retain the backchannel and/or prevent revocation of recallable state, the client needs to reconnect, and if it does, it **MUST** associate the connection to the session and backchannel via `BIND_CONN_TO_SESSION`. The server indicates when it has no callback connection via the `sr_status_flags` result from `SEQUENCE`.

#### 7.13.1.3. Backchannel GSS Context Loss

Via the `sr_status_flags` result of the `SEQUENCE` operation or other means, the client will learn if some or all of the `RPCSEC_GSS` contexts it assigned to the backchannel have been lost. If the client wants to retain the backchannel and/or not put recallable state subject to revocation, the client needs to use `BACKCHANNEL_CTL` to assign new contexts.

#### 7.13.1.4. Loss of Session

The replier might lose a record of the session. Causes include:

- \* Replier failure and restart.
- \* A catastrophe that causes the reply cache to be corrupted or lost on the media on which it was stored. This applies even if the replier indicated in the `CREATE_SESSION` results that it would persist the cache.
- \* The server purges the session of a client that has been inactive for a very extended period of time.

- \* As a result of configuration changes among a set of clustered servers, a network address previously connected to one server becomes connected to a different server that has no knowledge of the session in question. Such a configuration change will generally only happen when the original server ceases to function for a time.

Loss of reply cache often leads to loss of session. The replier indicates loss of session to the requester by returning NFS4ERR\_BADSESSION on the next operation that uses the session ID that refers to the lost session.

Although loss of session is often associated with loss of the associated clientid and corresponding locking state, this is not always the case. A session can be lost without loss of the corresponding clientid-based locking state in the event of clientid trunking, or when locking state is stored persistently but the reply cache is not. See Section 8 for details.

In the event of server restart, in the absence of clientid trunking, the following situations can arise:

- \* If neither the reply cache nor locking state is being stored persistently both the session and clientid are lost and new ones need to be established to continue operation.
- \* If the reply cache is persistent, it is possible that existing locking state is available so the existing session id and clientid can be tried going forward to determine if operation can be continued with existing locking state or a new clientid needs to be established and locks reclaimed.
- \* If the reply cache is not persistent, and the locking state is available in persistent storage the session is lost and a new session can be created for the existing clientid.

After an event like a server restart, the client may have lost its connections. The client assumes for the moment that the session has not been lost. It reconnects, and if it specified connection association enforcement when the session was created, it invokes BIND\_CONN\_TO\_SESSION using the session ID. Otherwise, it invokes SEQUENCE. If BIND\_CONN\_TO\_SESSION or SEQUENCE returns NFS4ERR\_BADSESSION, the client knows the session is not available to it when communicating with that network address. If the connection survives session loss, then the next SEQUENCE operation the client sends over the connection will get back NFS4ERR\_BADSESSION. The client again knows the session was lost.

Here is one suggested algorithm for the client when it gets NFS4ERR\_BADSESSION. It is not obligatory in that, if a client does not want to take advantage of such features as trunking, it may omit parts of it. However, it is a useful example that draws attention to various possible recovery issues:

1. If the client has other connections to other server network addresses associated with the same session, attempt a COMPOUND with a single operation, SEQUENCE, on each of the other connections.
2. If the attempts succeed, the session is still alive, and this is a strong indicator that the server's network address has moved. The client might send an EXCHANGE\_ID on the connection that returned NFS4ERR\_BADSESSION to see if there are opportunities for client ID trunking (i.e., the same client ID and so\_major\_id value are returned). The client might use DNS to see if the moved network address was replaced with another, so that the performance and availability benefits of session trunking can continue.
3. If the SEQUENCE requests fail with NFS4ERR\_BADSESSION, then the session no longer exists on any of the server network addresses for which the client has connections associated with that session ID. It is possible that the session is still alive and available on other network addresses. The client sends an EXCHANGE\_ID on all the connections to see if the server owner is still listening on those network addresses. If the same server owner is returned but a new client ID is returned, this is a strong indicator of a server restart. If both the same server owner and same client ID are returned, then this is a strong indication that the server did delete the session, and the client will need to send a CREATE\_SESSION if it has no other sessions for that client ID. If a different server owner is returned, the client can use DNS to find other network addresses. If it does not, or if DNS does not find any other addresses for the server, then the client will be unable to provide NFSv4.1 service, and fatal errors should be returned to processes that were using the server. If the client is using a "mount" paradigm, unmounting the server is advised.
4. If the client knows of no other connections associated with the session ID and server network addresses that are, or have been, associated with the session ID, then the client can use DNS to find other network addresses. If it does not, or if DNS does not find any other addresses for the server, then the client will be unable to provide NFSv4.1 service, and fatal errors should be returned to processes that were using the server. If the client is using a "mount" paradigm, unmounting the server is advised.

If there is a reconfiguration event that results in the same network address being assigned to servers where the `eir_server_scope` value is different, it cannot be guaranteed that a session ID generated by the first will be recognized as invalid by the first. Therefore, in managing server reconfigurations among servers with different server scope values, it is necessary to make sure that all clients have disconnected from the first server before effecting the reconfiguration. Nonetheless, clients should not assume that servers will always adhere to this requirement; clients **MUST** be prepared to deal with unexpected effects of server reconfigurations. Even where a session ID is inappropriately recognized as valid, it is likely either that the connection will not be recognized as valid or that a sequence value for a slot will not be correct. Therefore, when a client receives results indicating such unexpected errors, the use of `EXCHANGE_ID` to determine the current server configuration is **RECOMMENDED**.

A variation on the above is that after a server's network address moves, there is no NFSv4.1 server listening, e.g., no listener on port 2049. In this example, one of the following occur: the NFSv4 server returns `NFS4ERR_MINOR_VERSION_MISMATCH`, the NFS server returns a `PROG_MISMATCH` error, the RPC listener on 2049 returns `PROG_UNVAIL`, or attempts to reconnect to the network address timeout. These are best treated as equivalent to `SEQUENCE` returning `NFS4ERR_BADSESSION` for these purposes.

When the client detects session loss, it needs to call `CREATE_SESSION` to recover. Any non-idempotent operations that were in progress might have been performed on the server at the time of session loss. The client has no general way to recover from this.

Note that loss of session does not imply loss of byte-range lock, open, delegation, or layout state because locks, opens, delegations, and layouts are tied to the client ID and depend on the client ID, not the session. Nor does loss of byte-range lock, open, delegation, or layout state imply loss of session state, because the session depends on the client ID; loss of client ID however does imply loss of session, byte-range lock, open, delegation, and layout state. See Section 13.4.2. A session can survive a server restart, but lock recovery may still be needed.

It is possible that `CREATE_SESSION` will fail with `NFS4ERR_STALE_CLIENTID` (e.g., the server restarts and does not preserve client ID state). If so, the client needs to call `EXCHANGE_ID`, followed by `CREATE_SESSION`.

### 7.13.2. Events Requiring Server Action

The following events require server action to recover.

#### 7.13.2.1. Client Crash and Restart

As described in Section 25.35, a restarted client sends `EXCHANGE_ID` in such a way that it causes the server to delete any sessions it had.

#### 7.13.2.2. Client Crash with No Restart

If a client crashes and never comes back, it will never send `EXCHANGE_ID` with its old client owner. Thus, the server has session state that will never be used again. After an extended period of time, and if the server has resource constraints, it MAY destroy the old session as well as locking state.

#### 7.13.2.3. Extended Network Partition

To the server, the extended network partition may be no different from a client crash with no restart (see Section 7.13.2.2). Unless the server can discern that there is a network partition, it is free to treat the situation as if the client has crashed permanently.

#### 7.13.2.4. Backchannel Connection Loss

If there were callback requests outstanding at the time of a connection loss, then the server MUST retry the requests, as described in Section 7.6.2. Note that it is not necessary to retry requests over a connection with the same source network address or the same destination network address as the lost connection. As long as the session ID, slot ID, and sequence ID in the retry match that of the original request, the callback target will recognize the request as a retry even if it did see the request prior to disconnect.

If the connection lost is the last one associated with the backchannel, then the server MUST indicate that in the `sr_status_flags` field of every `SEQUENCE` reply until the backchannel is re-established. There are two situations, each of which uses different status flags: no connectivity for the session's backchannel and no connectivity for any session backchannel of the client. See Section 25.46 for a description of the appropriate flags in `sr_status_flags`.



#### 7.13.2.5. GSS Context Loss

The server needs to monitor when the number of RPCSEC\_GSS handles assigned to the backchannel reaches one, and when that one handle is near expiry (i.e., between one and two periods of lease time), and indicate so in the `sr_status_flags` field of all SEQUENCE replies. The server MUST indicate when all of the backchannel's assigned RPCSEC\_GSS handles have expired via the `sr_status_flags` field of all SEQUENCE replies.

#### 7.14. Parallel NFS and Sessions

A client and server can potentially be a non-pNFS implementation, a metadata server implementation, a data server implementation, or two or three types of implementations. The `EXCHGID4_FLAG_USE_NON_PNFS`, `EXCHGID4_FLAG_USE_PNFS_MDS`, and `EXCHGID4_FLAG_USE_PNFS_DS` flags (not mutually exclusive) are passed in the `EXCHANGE_ID` arguments and results to allow the client to indicate how it wants to use sessions created under the client ID, and to allow the server to indicate how it will allow the sessions to be used. See Section 20.5 for sessions considerations regarding the pNFS files layout type.

### 8. Persistence

[Author Aside]: This is a new top-level section which is based on the Persistence section previously within the discussion of Exactly-once Semantics. Essentially, it deletes the feature described in [RFC8881] which could never be implemented in that form and addresses the need with a new feature having the same goals.

While file data and metadata are typically stored persistently and are not affected by server restart, with the exception of certain optimizations for writing data, there are two sorts of data not normally stored persistently, that often are affected by server restart. Since [RFC8881] did not address either of these in a way that could be implemented, the entire area has been respecified for reasons discussed in Section 8.1.

For each of these types of data, the protocol provides an OPTIONAL feature whereby the server can provide persistent storage to eliminate functional problems when the data is lost or to simplify the process of reconstructing the data based on the client's knowledge.

- \* Reply caches may be stored persistently, as described in Section 8.2, allowing the same at-most-once semantics (often called "EOS") provided by the session-based reply cache to be maintained across server restarts.

As discussed below, the server may provide a persistent reply cache allowing EOS across server restarts or fully persistent sessions that allow the use of existing sessions to be continued across server restart.

- \* Per-client locking state may be stored persistently, as described in Section 8.3, allowing clients to continue after server restart without the delay caused by interposing a grace period during which all new lock requests are to be rejected.

If per-client locking state is not stored persistently, a grace period is provided to allow clients time to reclaim their locks. When this period is needed, requests to obtain new locks (e.g. when opening a file) are delayed until all clients have had a chance to reclaim their locks.

Although the incremental cost of supporting lock persistence is generally low enough that servers providing persistent sessions would provide persistent locking state as well, these two features are independent and the client cannot always assume lock persistence is available when an associated session is persistent and successfully recovered. For a discussion of how the client would be able to determine what state has been stored persistently and continue operation without unnecessary disruption, see Section 8.4

#### 8.1. Need for Feature Respecification

The original material has been modified substantially and extended in order address the three items listed below. As a result, the focus of the section has shifted to include all elements relevant to persistence across server failure, rather than dealing only with reply cache issues.

- \* Eliminate elements of the description that made the feature essentially unimplementable. These include overbroad requirements for atomicity and the assumption that all requests needed to be continued across server restart.
- \* Appropriately discuss lock persistence and its relation to reply cache persistence and session persistence.
- \* Provide new material describing the process by which the client finds out about the presence of persistence- related features in the event of server restart.

## 8.2. Persistence of Reply Cache

Since the reply cache is bounded, it is possible for the reply cache to be maintained in persistent storage so that it can be made available across server restarts. When the server undertakes to provide this support when the session is created (see Section 25.36 for details), it is uncertain whether what will be provided is either:

- \* Persistence of the reply cache only.
- \* Persistent of the session including its membership within the clientid of which it is a part.

The replier needs to persist the following information if it agreed to provide persistence for the session (when the session was created;

- \* The session ID.
- \* The slot table

This need to include the sequence ID and cached reply for each slot.

- \* Information about the connection(s) used by the server with is sufficient to determine whether a client attempting to connect after a server Restart.

This sort of information can be used to provide either of the two distinct sorts of session-based persistence. The server provides no specific commitment to provide either of these, although, as described in Section 8.4, the client will be able to determine which form, if any, has actually been provided, and respond appropriately

In describing persistence-related semantics it will be helpful to define the following two terms:

- \* An operation is said "reply-caching relevant" if it is either non-idempotent, modifying, or is the final operation (including the case of request termination because of an error) of a request that is specifically requested to be cached (i.e., has a SEQUENCE operation with sa\_cachethis set to true).
- \* A request is said "reply-caching relevant" if it contains one or more operations which are non-idempotent or modifying or it is specifically requested to be cached (i.e., has a SEQUENCE operation with sa\_cachethis set to true).

Whichever form of session-based persistence is provided by the server, any requests the client retries after the server restarts will return the results that are cached in the reply cache. However, these two forms differ with regard to the handling of new requests and the possible use of clientid-based persistence facilities:

- \* If only reply cache persistence is provided, any new requests will fail with NFS4ERR\_DEADSESSION being returned as the result of the initial SEQUENCE operation.

Because there is no need to use the sequence id to order future request the server does not need to update persistent storage, if two successive requests using the same slot are both not reply-caching relevant, although it does if one or both of the request is reply-cache relevant.

- \* If session persistence is provided, the existing session can be used after connection re-establishment to support the execution of new requests so that the client will be able to continue just as it would have if no session restart had occurred.

A persistent reply cache places certain demands on the server. Although it is not necessary to execute successive operations within a COMPOUND atomically, the transfer of the results of a set of operations and their installation in the persistent cache must immediately follow the execution of any reply-cache relevant operation so that it is impossible for operations to be executed or have other visible effects while not appearing in persistent reply cache.

If a client were to retry a sequence of operations that was issued to the server, the only acceptable outcomes are:

- \* an indication that the request is still being processed.
- \* a cached reply reflecting the completion of the request,
- \* a cached reply reflecting the interruption of the request due to server failure.
- \* an indication that the client ID or session has been lost (indicating a catastrophic loss of the reply cache or a session that has been deleted because the client failed to use the session for an extended period of time).

The possibility exists of situations in which a server could fail and restart in the middle of a COMPOUND procedure that contains one or more non-idempotent or idempotent-but-modifying operations. If the

server allows COMPOUND procedures to be continued after server failure, it creates significantly greater challenges for the execution of such requests and the atomic placement of results in the reply cache.

When a server providing a persistent reply cache does not continue a COMPOUND procedure that was interrupted by a server failure, the error NFS4ERR\_DEADSESSION is returned on the last operation which was executed.

### 8.3. Persistence of Locking State

Servers may make locking state available across a server restart in a number of ways including the following:

- \* Data related to the existence of locks and their corresponding characteristics can be stored in persistent RAM and then used after restart if the address of that storage can be reliably obtained after restart.
- \* The storage of locking-related state can be integrated with the file system by treating locking state in the same fashion used for other metadata.
- \* Locking state information may be periodically logged to block-based low-latency persistent storage with logging of individual updates.

Although the details will vary with the means of providing persistence that is adopted, it is important that locking state made available across the server restart be consistent with locking state reflected in the results of requests made by clients.

The simplest part of this is to ensure that all locking state changes are effectively made available persistently before returning to the requester. In addition, when lock state additions or deletions are reflected in the processing of other operations, the state changes must be available persistently before allowing or denying some operation done by another client. For example, when opens denying write prevent file removal, granting such opens or doing corresponding closes need to be reflected persistently before denying or allowing corresponding file removal. Similar consideration apply to doing IO when mandatory byte-range locks are supported

The following facts need to be kept in mind:

- \* There is no commitment by the server to provide this persistence and it may be dropped if for a particular client if unusual situations make it advisable.

This decision is made separately for each client so that it is possible there will be server restarts where some, but not all, clients have persistent locking state available.

- \* While the fact that a reclaim on a reclaimable lock is part of the locking state which is to be persistent, the client's state of awareness of that need not be.

There is thus no need for the reclaiming client to inform the server that it has completed specific individual reclaims after receiving the response.

#### 8.4. Client Handling of Server Failure When Persistence Can be Used

When server failure occurs, the connection to the client will be disconnected and the client can then find out, as described below, whether server failure has occurred and what steps are necessary to continue use of the client with minimal disruption to those using the client.

This process includes the potential use of a persistent reply cache, as described in Section 8.5. The same process is followed depending on whether the server provided only a persistent reply cache or full session persistence.

If the server did not promise any session persistence, the client instead immediately does an `EXCHANGE_ID` followed by a `CREATE_SESSION`. On the other hand, if there was a possible use of a persistent reply cache, the use of `EXCHANGE_ID/CREATE_SESSION` is conditional and only happens if a new request has been completed with the error `NFS4ERR_STALECLIENTID`.

In either case, the next step depends on whether the `clientid` is the same as the one before the disconnection. If it is, then recovery is complete and new requests can be issued. This could happen if there were no server restart but also could if a combination of session-based and `clientid`-based persistence allowed the server failure to be dealt with essentially transparently.

In the case in which the `clientid` is different, the client need to reclaim its locks, as described in Section 8.6.

Even in the case in which lock persistence is available for a client, it is still possible that attempts to obtain new locks will fail with NFS4ERR\_GRACE if other clients do not have their locks made available persistently.

#### 8.5. Client Use of Session-based Persistence

After the connection to the server is re-established, the server will try to re-establish the connection, as the connection breakage occurred at a lower layer, without server restart. Although it is theoretically possible for an intermediary to hide such a disconnection, it would cause problems if it were to do so and the client had no knowledge of the server failure. The discussion here assumes that no such disconnection-hiding implementation is in effect.

After re-establishing the connection to the server, the client would initially attempt to continue use of the session, since it has no knowledge of whether the disconnection was the result of a server restart. If persistence was not requested when creating the session or the server indicated it was not present, then the client can legitimately conclude that EOS semantics was not available across server restart and needs to operate in that environment.

The continued use of the existing session could include both retries of requests issued before the disconnection and issuing new requests. As a result, the discussion below will deal with both type of requests. Given that context, one needs to note the following:

- \* Whether a given request is a retry or a new one may be judged differently by the client and the server.

While it is virtually certain that a new request issued by the client will be perceived as such by the server, the reverse is not the case. Retries issued by the client might be perceived as new requests, if the original request was lost before it was executed or its existence was noted in persistent storage.

- \* Although it might be desirable for a client to obtain information about existing requests before issuing a new one, the discussion will not assume that clients take steps to prevent new requests from being issued.

Since retries, as perceived by the client, may be considered as new requests by the server, the prevention of new requests by the client does not ensure that the server will not see and respond to such requests.

After re-establishing the connection, the client will be able to issue requests, including retries of requests already issued before the disconnection occurred. These retries need to be issued since there is no way the results of these requests could be communicated back to the client in the absence of a retry since the connection on which it was received no longer exists.

When responses to these requests are received, what is to be done depends primarily about the error, if any, associated with the response:

- \* In all the cases except the two special error codes noted in the bulleted items below, including receiving no error, the client can conclude that the request was executed to completion as reflected in the response. By design, the client is not aware of whether the execution occurred before or after the server restart, or whether a server restart, in fact, occurred. However, if persistence was requested when the session was created and the server indicated it was present, the client can assume that the request was executed exactly once with the result reflected in the response.

When this is the result that is returned for new requests, it can be because the server has provided full session persistence or because no server restart has occurred. In the former case, it must be true that the server has provided persistent storage of locking state for the `d` associated `clientid` since, if it had not, the error `NFS4ERR_STALECLIENTID` would have been returned.

- \* In the case that `NFS4ERR_DEADSESSION` is returned on the `SEQUENCE` operation, the most likely cause is that the request was, from the server's point of view, a new request and that session persistence was not provided by the server. In this case, the current request should be deferred until the results of all retried requests known to the client have been resolved. Others that are considered new by the server also need to be deferred until a reply cache information is obtained.

In the case that `NFS4ERR_DEADSESSION` is returned on another operation, the request is one that was discontinued as a result of server restart. It is most likely that the request was one that contained more than one non-idempotent or modifying operations, with the server failing after one had been completed but before later operations were started. In this case the client has been informed of a partially complete request and needs to issue a new request to include the operations that were not performed as part of the initial request.



- \* In the case that the error NFS4ERR\_STALECLIENTID is returned, the server has recognized a new request but was unable to continue its execution because the locking information it would use has been destroyed as part of the server restart. This can occur if no persistence was provided for the session, if the persistence was limited to the reply cache or if there was session persistence and client locking state was not maintained persistently.

In this case lock recovery will be required but it will need to be delayed until all requests that were issued before the disconnection have been marked completed using the persisted reply cache.

Once the existing pending requests are disposed of, the client can proceed to doing new requests, although it might have to do lock recovery first. This can occur after a persistent reply cache is used to provide EOS or after it is found that there is no session persistence provided by the server.

#### 8.6. Client Use of Clientid-based Persistence

At this point, lock recovery needs to begin if a new request is processed and completes returning the error NFS4ERR\_STALECLIENTID. If no new requests have been issued at this point, the client can issue a request consisting only of a SEQUENCE operation to provide a test. If NFS4ERR\_STALECLIENTID is not returned then the client will assume either that there has been no server restart or that server restart has been accompanied with the recovery of locking state for the current clientid. Otherwise, lock recovery can be done as part of a server-provided grace period. The following three steps need to be taken:

When lock recovery is necessary, the client needs to inform the new server of the existence of its locks before using stateids it obtained before the server restart. This process is referred to as reclaiming the client's locks, which is accomplished using the method listed below, depending on the type of lock to be reclaimed.

- \* Opens can generally be reclaimed by doing an OPEN with the claim type CLAIM\_PREVIOUS.

This includes the case of opens associated with delegation. For details, see Section 15.2.1,

There is no specific way to reclaim delegations that have no associated open. In such cases, the client can open the file asking for an associated delegation, and return it immediately

- \* To reclaim byte-range locks, a LOCK operation with the reclaim parameter set to true is used.

The associated open will need to be reclaimed first.

- \* There is no provision regarding reclaiming of layouts and thus no way to obtain them during a grace period.

As a result, in case in which locking state is not made available by the server across a server failure, use of the data server is not immediately available and the client is best off doing IO through the MDS until obtaining needed layouts once the rest of lock reclamation is complete.

Once all reclaimable locks have been reclaimed, the client needs to do a global RECLAIM\_COMPLETE to indicate that process is complete. This is necessary to allow new locks to be obtained. However, even after this done, such requests might still be rejected with NFS4ERR\_GRACE if other clients have not completed their lock reclamations.

## 9. Protocol Constants and Data Types

The syntax and semantics to describe the data types of the NFSv4.1 protocol are defined in the XDR ([RFC4506]) and RPC ([RFC5531]) documents. The next sections build upon the XDR data types to define constants, types, and structures specific to this protocol. The full list of XDR data types is in RFCTBD30.

### 9.1. Basic Constants

```
const NFS4_FH_SIZE           = 128;
const NFS4_VERIFIER_SIZE     = 8;
const NFS4_OPAQUE_LIMIT      = 1024;
const NFS4_SESSIONID_SIZE    = 16;

const NFS4_INT64_MAX          = 0x7fffffffffffffff;
const NFS4_UINT64_MAX         = 0xfffffffffffffff;
const NFS4_INT32_MAX          = 0x7fffffff;
const NFS4_UINT32_MAX         = 0xffffffff;

const NFS4_MAXFILELEN        = 0xfffffffffffffff;
const NFS4_MAXFILEOFF        = 0xfffffffffffffff;
```

Except where noted, all these constants are defined in bytes.

- \* NFS4\_FH\_SIZE is the maximum size of a filehandle.

- \* NFS4\_VERIFIER\_SIZE is the fixed size of a verifier.
- \* NFS4\_OPAQUE\_LIMIT is the maximum size of certain opaque information.
- \* NFS4\_SESSIONID\_SIZE is the fixed size of a session identifier.
- \* NFS4\_INT64\_MAX is the maximum value of a signed 64-bit integer.
- \* NFS4\_UINT64\_MAX is the maximum value of an unsigned 64-bit integer.
- \* NFS4\_INT32\_MAX is the maximum value of a signed 32-bit integer.
- \* NFS4\_UINT32\_MAX is the maximum value of an unsigned 32-bit integer.
- \* NFS4\_MAXFILELEN is the maximum length of a regular file.
- \* NFS4\_MAXFILEOFF is the maximum offset into a regular file.

## 9.2. Basic Data Types

These are the base NFSv4.1 data types.

Data Type	Definition
int32_t	typedef int int32_t;
uint32_t	typedef unsigned int uint32_t;
int64_t	typedef hyper int64_t;
uint64_t	typedef unsigned hyper uint64_t;
attrlist4	typedef opaque attrlist4<>; Used for file/directory attributes.
bitmap4	typedef uint32_t bitmap4<>; Used in attribute array encoding.
changeid4	typedef uint64_t changeid4; Used in the definition of change_info4.

clientid4	typedef uint64_t clientid4;  Shorthand reference to client identification.
count4	typedef uint32_t count4;  Various count parameters (READ, WRITE, COMMIT).
length4	typedef uint64_t length4;  The length of a byte-range within a file.
mode4	typedef uint32_t mode4;  Mode attribute data type.
nfs_cookie4	typedef uint64_t nfs_cookie4;  Opaque cookie value for READDIR.
nfs_fh4	typedef opaque nfs_fh4<NFS4_FHSIZE>;  Filehandle definition.
nfs_ftype4	enum nfs_ftype4;  Various defined file types.
nfsstat4	enum nfsstat4;  Return value for operations.
offset4	typedef uint64_t offset4;  Various offset designations (READ, WRITE, LOCK, COMMIT).
qop4	typedef uint32_t qop4;  Quality of protection designation in SECINFO.
sec_oid4	typedef opaque sec_oid4<>;  Security Object Identifier. The sec_oid4 data type is not really opaque. Instead, it contains an ASN.1 OBJECT IDENTIFIER as used by GSS-API in the mech_type argument to GSS_Init_sec_context. See [RFC2743] for details.

sequenceid4	typedef uint32_t sequenceid4;  Sequence number used for various session operations (EXCHANGE_ID, CREATE_SESSION, SEQUENCE, CB_SEQUENCE).
seqid4	typedef uint32_t seqid4;  Sequence identifier used for locking.
sessionid4	typedef opaque sessionid4[NFS4_SESSIONID_SIZE];  Session identifier.
slotid4	typedef uint32_t slotid4;  Sequencing artifact for various session operations (SEQUENCE, CB_SEQUENCE).
utf8string	typedef opaque utf8string<>;  UTF-8 encoding for strings.
utf8str_cis	typedef utf8string utf8str_cis;  Case-insensitive UTF-8 string.
utf8str_cs	typedef utf8string utf8str_cs;  Case-sensitive UTF-8 string.
utf8str_mixed	typedef utf8string utf8str_mixed;  UTF-8 strings with a domain or host prefix and an server or file name suffix. Domains can be internationalized as described in [I-D.ietf-nfsv4-internationalization].
utf8pref	typedef opaque utf8pref<>;  String for which UTF-8 encoding is preferred, although other encodings can be used,
component4	typedef utf8pref component4;  Represents pathname components, which may be either encoded using UTF-8 or nor, with use of

	UTF-8 needed to support normalization and case-insensitivity.
linktext4	typedef opaque linktext4<>  Symbolic link contents ("symbolic link" is defined in an Open Group [symlink] standard).
pathname4	typedef component4 pathname4<>;  Represents pathname for fs_locations.
verifier4	typedef opaque verifier4[NFS4_VERIFIER_SIZE];  Verifier used for various operations (COMMIT, CREATE, EXCHANGE_ID, OPEN, READDIR, WRITE) NFS4_VERIFIER_SIZE is defined as 8.

Table 1

## End of Base Data Types

## 9.3. Structured Data Types

## 9.3.1. nfstime4

```
struct nfstime4 {
    int64_t      seconds;
    uint32_t     nseconds;
};
```

The nfstime4 data type gives the number of seconds and nanoseconds since midnight or zero hour January 1, 1970 Coordinated Universal Time (UTC). Values greater than zero for the seconds field denote dates after the zero hour January 1, 1970. Values less than zero for the seconds field denote dates before the zero hour January 1, 1970. In both cases, the nseconds field is to be added to the seconds field for the final time representation. For example, if the time to be represented is one-half second before zero hour January 1, 1970, the seconds field would have a value of negative one (-1) and the nseconds field would have a value of one-half second (500000000). Values greater than 999,999,999 for nseconds are invalid.

This data type is used to pass time and date information. A server converts to and from its local representation of time when processing time values, preserving as much accuracy as possible. If the precision of timestamps stored for a file system object is less than

defined, loss of precision can occur. An adjunct time maintenance protocol is RECOMMENDED to reduce skew between client and server times.

#### 9.3.2. time\_how4

```
enum time_how4 {
    SET_TO_SERVER_TIME4 = 0,
    SET_TO_CLIENT_TIME4 = 1
};
```

#### 9.3.3. settime4

```
union settime4 switch (time_how4 set_it) {
    case SET_TO_CLIENT_TIME4:
        nfstime4      time;
    default:
        void;
};
```

The time\_how4 and settime4 data types are used for setting timestamps in file object attributes. If set\_it is SET\_TO\_SERVER\_TIME4, then the server uses its local representation of time for the time value.

#### 9.3.4. specdata4

```
struct specdata4 {
    uint32_t specdata1; /* major device number */
    uint32_t specdata2; /* minor device number */
};
```

This data type represents the device numbers for the device file types NF4CHR and NF4BLK.

#### 9.3.5. fsid4

```
struct fsid4 {
    uint64_t      major;
    uint64_t      minor;
};
```

#### 9.3.6. change\_policy4

```
struct change_policy4 {
    uint64_t      cp_major;
    uint64_t      cp_minor;
};
```

The `change_policy4` data type is used for the `change_policy` OPTIONAL attribute. It provides change sequencing indication analogous to the `change` attribute. To enable the server to present a value valid across server re-initialization without requiring persistent storage, two 64-bit quantities are used, allowing one to be a server instance ID and the second to be incremented non-persistently, within a given server instance.

#### 9.3.7. `fattr4`

```
struct fattr4 {
    bitmap4      attrmask;
    attrlist4    attr_vals;
};
```

The `fattr4` data type is used to represent sets of protocol-defined attributes.

The `bitmap` is a counted array of 32-bit integers used to contain bit values. The position of the integer in the array that contains bit `n` can be computed from the expression  $(n / 32)$ , and its bit within that integer is  $(n \bmod 32)$ .

```

              0              1
+-----+-----+-----+
| count  | 31 .. 0 | 63 .. 32 |
+-----+-----+-----+
```

#### 9.3.8. `change_info4`

```
struct change_info4 {
    bool      atomic;
    changeid4 before;
    changeid4 after;
};
```

This data type is used with the `CREATE`, `LINK`, `OPEN`, `REMOVE`, and `RENAME` operations to let the client know the value of the `change` attribute for the directory in which the target file system object resides.

#### 9.3.9. `netaddr4`

```
struct netaddr4 {
    /* see struct rpcb in RFC 1833 */
    string na_r_netid<>; /* network id */
    string na_r_addr<>;  /* universal address */
};
```



The `netaddr4` data type is used to identify network transport endpoints. The `na_r_netid` and `na_r_addr` fields respectively contain a `netid` and `uaddr`. The `netid` and `uaddr` concepts are defined in [RFC5665]. The `netid` and `uaddr` formats for TCP over IPv4 and TCP over IPv6 are defined in [RFC5665], specifically Tables 2 and 3 and in Sections 5.2.3.3 and 5.2.3.4.

#### 9.3.10. `state_owner4`

```
struct state_owner4 {
    clientid4      clientid;
    opaque         owner<NFS4_OPAQUE_LIMIT>;
};

typedef state_owner4 open_owner4;
typedef state_owner4 lock_owner4;
```

The `state_owner4` data type is the base type for the `open_owner4` (Section 9.3.10.1) and `lock_owner4` (Section 9.3.10.2).

##### 9.3.10.1. `open_owner4`

This data type is used to identify the owner of OPEN state.

##### 9.3.10.2. `lock_owner4`

This structure is used to identify the owner of byte-range locking state.

#### 9.3.11. `open_to_lock_owner4`

```
struct open_to_lock_owner4 {
    seqid4         open_seqid;
    stateid4       open_stateid;
    seqid4         lock_seqid;
    lock_owner4    lock_owner;
};
```

This data type is used for the first LOCK operation done for an `open_owner4`. It provides both the `open_stateid` and `lock_owner`, such that the transition is made from a valid `open_stateid` sequence to that of the new `lock_stateid` sequence. Using this mechanism avoids the confirmation of the `lock_owner/lock_seqid` pair since it is tied to established state in the form of the `open_stateid/open_seqid`.

#### 9.3.12. `stateid4`

```
struct stateid4 {
    uint32_t      seqid;
    opaque        other[12];
};
```

This data type is used for the various state sharing mechanisms between the client and server. The client never modifies a value of data type stateid. The starting value of the "seqid" field is undefined. The server is required to increment the "seqid" field by one at each transition of the stateid. This is important since the client will inspect the seqid in OPEN stateids to determine the order of OPEN processing done by the server.

#### 9.3.13. layouttype4

```
enum layouttype4 {
    LAYOUT4_NFSV4_1_FILES    = 0x1,
    LAYOUT4_OSD2_OBJECTS     = 0x2,
    LAYOUT4_BLOCK_VOLUME     = 0x3
};
```

This data type indicates what type of layout is being used. The file server advertises the layout types it supports through the fs\_layout\_type file system attribute (Section 11.16.1). A client asks for layouts of a particular type in LAYOUTGET, and processes those layouts using layout-type-specific logic.

The layouttype4 data type is 32 bits in length. The range represented by the layout type is split into three parts. Type 0x0 is reserved. Types within the range 0x00000001-0x7FFFFFFF are globally unique and are assigned according to the description in Section 29.5; they are maintained by IANA. Types within the range 0x80000000-0xFFFFFFFF are site specific and for private use only.

The LAYOUT4\_NFSV4\_1\_FILES enumeration specifies that the NFSv4.1 file layout type, as defined in Section 20, is to be used. The LAYOUT4\_OSD2\_OBJECTS enumeration specifies that the object layout, as defined in [RFC5664], is to be used. Similarly, the LAYOUT4\_BLOCK\_VOLUME enumeration specifies that the block/volume layout, as defined in [RFC5663], is to be used.

#### 9.3.14. deviceid4

```
const NFS4_DEVICEID4_SIZE = 16;

typedef opaque deviceid4[NFS4_DEVICEID4_SIZE];
```

Layout information includes device IDs that specify a storage device through a compact handle. Addressing and type information is obtained with the GETDEVICEINFO operation. Device IDs are not guaranteed to be valid across metadata server restarts. A device ID is unique per client ID and layout type. See Section 18.4.9 for more details.

#### 9.3.15. device\_addr4

```
struct device_addr4 {  
    layouttype4      da_layout_type;  
    opaque           da_addr_body<>;  
};
```

The device address is used to set up a communication channel with the storage device. Different layout types will require different data types to define how they communicate with storage devices. The nominally opaque `da_addr_body` field is interpreted based on the specified `da_layout_type` field. As required by Section 19.1.2, the layout type specification provides the actual data within this field.

This document defines the device address for the NFSv4.1 file layout (See Section 20.7.2), which identifies a storage device by network IP address and port number. Device types for other layout types are defined by their respective layout specifications.

#### 9.3.16. layout\_content4

```
struct layout_content4 {  
    layouttype4 loc_type;  
    opaque      loc_body<>;  
};
```

The `loc_body` field is interpreted based on the layout type (`loc_type`). This document defines the `loc_body` for the NFSv4.1 file layout type; see Section 20.7 for its definition.

#### 9.3.17. layout4

```
struct layout4 {  
    offset4      lo_offset;  
    length4      lo_length;  
    layoutiomode4 lo_iomode;  
    layout_content4 lo_content;  
};
```

The layout4 data type defines a layout for a file. The layout type specific data is opaque within lo\_content. Since layouts are subdividable, the offset and length together with the file's filehandle, the client ID, iomode, and layout type identify the layout.

#### 9.3.18. layoutupdate4

```
struct layoutupdate4 {
    layouttype4      lou_type;
    opaque           lou_body<>;
};
```

The layoutupdate4 data type is used by the client to return updated layout information to the metadata server via the LAYOUTCOMMIT (Section 25.42) operation. This data type provides a channel to pass layout-type-specific information (in field lou\_body) back to the metadata server. The actual contents of the nominally opaque field lou\_body is the responsibility of the layout type specification for lou\_type. See Section 19.1.8 for details about how this is used by LAYOUTCOMMIT for individual layout type. For example, for the block/volume layout type, this could include the list of reserved blocks that were written.

#### 9.3.19. layouthint4

```
struct layouthint4 {
    layouttype4      loh_type;
    opaque           loh_body<>;
};
```

The layouthint4 data type is used by the client to pass in a hint about the type of layout it would like created for a particular file. It is the data type specified by the layout\_hint attribute described in Section 11.16.4. The metadata server may ignore the hint or may selectively ignore fields within the hint. This hint should be provided at create time as part of the initial attributes within OPEN. The actual contents loh\_body field is specific to the type of layout (loh\_type) and is described by the layout type specification.

#### 9.3.20. layoutiomode4

```
enum layoutiomode4 {
    LAYOUTIOMODE4_READ      = 1,
    LAYOUTIOMODE4_RW        = 2,
    LAYOUTIOMODE4_ANY        = 3
};
```

The iomode specifies whether the client intends to just read or both read and write the data represented by the layout. While the LAYOUTIOMODE4\_ANY iomode MUST NOT be used in the arguments to the LAYOUTGET operation, it MAY be used in the arguments to the LAYOUTRETURN and CB\_LAYOUTRECALL operations. The LAYOUTIOMODE4\_ANY iomode specifies that layouts pertaining to both LAYOUTIOMODE4\_READ and LAYOUTIOMODE4\_RW iomodes are being returned or recalled, respectively. The metadata server's use of the iomode may depend on the layout type being used. The storage devices MAY validate I/O accesses against the iomode and reject invalid accesses.

#### 9.3.21. nfs\_impl\_id4

```
struct nfs_impl_id4 {  
    utf8str_cis    nii_domain;  
    utf8str_cs     nii_name;  
    nfstime4       nii_date;  
};
```

This data type is used to identify client and server implementation details. The nii\_domain field is the DNS domain name with which the implementer is associated. The nii\_name field is the product name of the implementation and is completely free form. It is RECOMMENDED that the nii\_name be used to distinguish machine architecture, machine platforms, revisions, versions, and patch levels. The nii\_date field is the timestamp of when the software instance was published or built.

#### 9.3.22. threshold\_item4

```
struct threshold_item4 {  
    layouttype4    thi_layout_type;  
    bitmap4        thi_hintset;  
    opaque         thi_hintlist<>;  
};
```

This data type contains a list of hints specific to a layout type for helping the client determine when it should send I/O directly through the metadata server versus the storage devices. The data type consists of the layout type (thi\_layout\_type), a bitmap (thi\_hintset) describing the set of hints supported by the server (they may differ based on the layout type), and a list of hints (thi\_hintlist) whose content is determined by the hintset bitmap. See the mdsthreshold attribute for more details.

The thi\_hintset field is a bitmap of the following values:

name	#	Data Type	Description
threshold4_read_size	0	length4	If a file's length is less than the value of threshold4_read_size, then it is RECOMMENDED that the client read from the file via the MDS and not a storage device.
threshold4_write_size	1	length4	If a file's length is less than the value of threshold4_write_size, then it is RECOMMENDED that the client write to the file via the MDS and not a storage device.
threshold4_read_iosize	2	length4	For read I/O sizes below this threshold, it is RECOMMENDED that the client read data using the MDS.
threshold4_write_iosize	3	length4	For write I/O sizes below this threshold, it is RECOMMENDED that the client write data using the MDS.

Table 2

## 9.3.23. mdsthreshold4

```
struct mdsthreshold4 {
    threshold_item4 mth_hints<>;
};
```

This data type holds an array of elements of data type threshold\_item4, each of which is valid for a particular layout type. An array is necessary because a server can support multiple layout types for a single file.

## 10. Filehandles

The filehandle in the NFS protocol is a per-server unique identifier for a file system object. The contents of the filehandle are opaque to the client. Therefore, the server is responsible for translating the filehandle to an internal representation of the file system object.

### 10.1. Obtaining the First Filehandle

The operations of the NFS protocol are defined in terms of one or more filehandles. Therefore, the client needs a filehandle to initiate communication with the server. With the NFSv3 protocol ([RFC1813]), there exists an ancillary protocol to obtain this first filehandle. The MOUNT protocol, RPC program number 100005, provides the mechanism of translating a string-based file system pathname to a filehandle, which can then be used by the NFS protocols.

The MOUNT protocol has deficiencies in the area of security and use via firewalls. This is one reason that the use of the public filehandle was introduced in [RFC2054] and [RFC2055]. With the use of the public filehandle in combination with the LOOKUP operation in the NFSv3 protocol, it has been demonstrated that the MOUNT protocol is unnecessary for viable interaction between NFS client and server.

Therefore, the NFSv4.1 protocol will not use an ancillary protocol for translation from string-based pathnames to a filehandle. Two special filehandles will be used as starting points for the NFS client.

#### 10.1.1. Root Filehandle

The first of the special filehandles is the ROOT filehandle. The ROOT filehandle is the "conceptual" root of the file system namespace at the NFS server. The client uses or starts with the ROOT filehandle by employing the PUTROOTFH operation. The PUTROOTFH operation instructs the server to set the "current" filehandle to the ROOT of the server's file tree. Once this PUTROOTFH operation is used, the client can then traverse the entirety of the server's file tree with the LOOKUP operation. A complete discussion of the server namespace is in Section 12.

### 10.1.2. Public Filehandle

The second special filehandle is the PUBLIC filehandle. Unlike the ROOT filehandle, the PUBLIC filehandle may be bound or represent an arbitrary file system object at the server. The server is responsible for this binding. It may be that the PUBLIC filehandle and the ROOT filehandle refer to the same file system object. However, it is up to the administrative software at the server and the policies of the server administrator to define the binding of the PUBLIC filehandle and server file system object. The client may not make any assumptions about this binding. The client uses the PUBLIC filehandle via the PUTPUBFH operation.

### 10.2. Filehandle Types

In the NFSv3 protocol, there was one type of filehandle with a single set of semantics. This type of filehandle is termed "persistent" in NFSv4.1. The semantics of a persistent filehandle remain the same as before. A new type of filehandle introduced in NFSv4.1 is the "volatile" filehandle, which attempts to accommodate certain server environments.

The volatile filehandle type was introduced to address server functionality or implementation issues that make correct implementation of a persistent filehandle infeasible. Some server environments do not provide a file-system-level invariant that can be used to construct a persistent filehandle. The underlying server file system may not provide the invariant or the server's file system programming interfaces may not provide access to the needed invariant. Volatile filehandles may ease the implementation of server functionality such as hierarchical storage management or file system reorganization or migration. However, the volatile filehandle increases the implementation burden for the client.

Since the client will need to handle persistent and volatile filehandles differently, a file attribute is defined that may be used by the client to determine the filehandle types being returned by the server.

#### 10.2.1. General Properties of a Filehandle

The filehandle contains all the information the server needs to distinguish an individual file. To the client, the filehandle is opaque. The client stores filehandles for use in a later request and can compare two filehandles from the same server for equality by doing a byte-by-byte comparison. However, the client MUST NOT otherwise interpret the contents of filehandles. If two filehandles from the same server are equal, they MUST refer to the same file.



Servers should try to maintain a one-to-one correspondence between filehandles and files, but this is not required. Clients MUST use filehandle comparisons only to improve performance, not for correct behavior. All clients need to be prepared for situations in which it cannot be determined whether two filehandles denote the same object and in such cases, avoid making invalid assumptions that might cause incorrect behavior. Further discussion of filehandle and attribute comparison in the context of data caching is presented in Section 15.3.6.

As an example, in the case that two different pathnames when traversed at the server terminate at the same file system object, the server will return the same filehandle for each path. This can occur if a hard link (See [hardlink]) is used to create two file names that refer to the same underlying file object and associated data. For example, if paths /a/b/c and /a/d/c refer to the same file, the server will return the same filehandle for both pathnames' traversals.

#### 10.2.2. Persistent Filehandle

A persistent filehandle is defined as having a fixed value for the lifetime of the file system object to which it refers. Once the server creates the filehandle for a file system object, the server MUST accept the same filehandle for the object for the lifetime of the object. If the server restarts, the NFS server MUST honor the same filehandle value as it did in the server's previous instantiation. Similarly, if the file system is migrated, the new NFS server MUST honor the same filehandle as the old NFS server.

The persistent filehandle will become stale or invalid when the file system object is removed. When the server is presented with a persistent filehandle that refers to a deleted object, it MUST return an error of NFS4ERR\_STALE. A filehandle may become stale when the file system containing the object is no longer available. The file system may become unavailable if it exists on removable media and the media is no longer available at the server or the file system in whole has been destroyed or the file system has simply been removed from the server's namespace (i.e., unmounted in a UNIX environment).

#### 10.2.3. Volatile Filehandle

A volatile filehandle does not share the same longevity characteristics of a persistent filehandle. The server may determine that a volatile filehandle is no longer valid at many different points in time. If the server can definitively determine that a volatile filehandle refers to an object that has been removed, the server should return NFS4ERR\_STALE to the client (as is the case for

persistent filehandles). In all other cases where the server determines that a volatile filehandle can no longer be used, it should return an error of NFS4ERR\_FHEXPIRED.

The REQUIRED attribute "fh\_expire\_type" is used by the client to determine what type of filehandle the server is providing for a particular file system. This attribute is a bitmask with the following values:

**FH4\_PERSISTENT** The value of FH4\_PERSISTENT is used to indicate a persistent filehandle, which is valid until the object is removed from the file system. The server will not return NFS4ERR\_FHEXPIRED for this filehandle. FH4\_PERSISTENT is defined as a value in which none of the bits specified below are set.

**FH4\_VOLATILE\_ANY** The filehandle may expire at any time, except as specifically excluded (i.e., FH4\_NO\_EXPIRE\_WITH\_OPEN).

**FH4\_NOEXPIRE\_WITH\_OPEN** May only be set when FH4\_VOLATILE\_ANY is set. If this bit is set, then the meaning of FH4\_VOLATILE\_ANY is qualified to exclude any expiration of the filehandle when it is open.

**FH4\_VOL\_MIGRATION** The filehandle will expire as a result of a file system transition (migration or replication), in those cases in which the continuity of filehandle use is not specified by handle class information within the fs\_locations\_info attribute. When this bit is set, clients without access to fs\_locations\_info information should assume that filehandles will expire on file system transitions.

**FH4\_VOL\_RENAME** The filehandle will expire during rename. This includes a rename by the requesting client or a rename by any other client. If FH4\_VOL\_ANY is set, FH4\_VOL\_RENAME is redundant.

Servers that provide volatile filehandles that can expire while open require special care as regards handling of RENAMEs and REMOVEs. This situation can arise if FH4\_VOL\_MIGRATION or FH4\_VOL\_RENAME is set, if FH4\_VOLATILE\_ANY is set and FH4\_NOEXPIRE\_WITH\_OPEN is not set, or if a non-read-only file system has a transition target in a different handle class. In these cases, the server should deny a RENAME or REMOVE that would affect an OPEN file of any of the components leading to the OPEN file. In addition, the server should deny all RENAME or REMOVE requests during the grace period, in order to make sure that reclaims of files where filehandles may have expired do not do a reclaim for the wrong file.

Volatile filehandles are especially suitable for implementation of the pseudo file systems used to bridge exports. See Section 12.5 for a discussion of this.

### 10.3. One Method of Constructing a Volatile Filehandle

A volatile filehandle, while opaque to the client, could contain:

```
[volatile bit = 1 | server boot time | slot | generation number]
```

- \* slot is an index in the server volatile filehandle table
- \* generation number is the generation number for the table entry/  
slot

When the client presents a volatile filehandle, the server makes the following checks, which assume that the check for the volatile bit has passed. If the server boot time is less than the current server boot time, return NFS4ERR\_FHEXPIRED. If slot is out of range, return NFS4ERR\_BADHANDLE. If the generation number does not match, return NFS4ERR\_FHEXPIRED.

When the server restarts, the table is gone (it is volatile).

If the volatile bit is 0, then it is a persistent filehandle with a different structure following it.

### 10.4. Client Recovery from Filehandle Expiration

If possible, the client should recover from the receipt of an NFS4ERR\_FHEXPIRED error. The client must take on additional responsibility so that it may prepare itself to recover from the expiration of a volatile filehandle. If the server returns persistent filehandles, the client does not need these additional steps.

For volatile filehandles, most commonly the client will need to store the component names leading up to and including the file system object in question. With these names, the client should be able to recover by finding a filehandle in the namespace that is still available or by starting at the root of the server's file system namespace.

If the expired filehandle refers to an object that has been removed from the file system, obviously the client will not be able to recover from the expired filehandle.

It is also possible that the expired filehandle refers to a file that has been renamed. If the file was renamed by another client, again it is possible that the original client will not be able to recover. However, in the case that the client itself is renaming the file and the file is open, it is possible that the client may be able to recover. The client can determine the new pathname based on the processing of the rename request. The client can then regenerate the new filehandle based on the new pathname. The client could also use the COMPOUND procedure to construct a series of operations like:

```
RENAME A B
LOOKUP B
GETFH
```

Note that the COMPOUND procedure does not provide atomicity. This example only reduces the overhead of recovering from an expired filehandle.

## 11. File Attributes

To meet the requirements of extensibility and increased interoperability with non-UNIX platforms, attributes are being handled in a more flexible manner than NFSv3. The NFSv3 `fattr3` structure consists of a fixed list of attributes some of which that might not all be supported by some potential servers and includes some attributes that not all clients have an interest in. The `fattr3` structure and similar fixed structures cannot be extended as new needs arise and provide no way to indicate non-support of particular attributes. Within the NFSv4.1 protocol, the client is able to query what attributes the server supports and construct requests that deal only with those supported attributes (or a subset thereof). This raises the issues, discussed in Sections 11.1 through 11.3 and 11.5 through 11.6, of determining how the non-support of particular attributes is to be dealt with.

### 11.1. Categorization of File Attributes

In order to clarify the requirements for server support of particular attributes, and to provide guidance for clients dealing with non-support of particular attributes, all NFSv4.1 attributes are divided into the groups listed below:

All of these attributes are accommodated in the NFSv4.1 protocol by a specific, well-defined encoding and are identified by a number. They are interrogated by setting a bit in the bit vector sent in a `GETATTR`, request. The server response includes a bit vector to indicate which attributes were returned in the response.

The following attribute categories are defined:

- \* REQUIRED attributes, as discussed in Section 11.4.
- \* OPTIONAL attributes, as discussed in Section 11.5.
- \* Experimental attributes, as discussed in Section 11.6.

New attributes of any of these categories may be added to the NFSv4 protocol as part of a new minor version by publishing a Standards Track RFC that allocates a new attribute number value and defines the encoding for the attribute. In addition, new minor versions can move attributes between categories or make formerly OPTIONAL and Experimental attributes MANDATORY to NOT implement. Similarly, OPTIONAL attributes may be added to an existing extensible version by publishing a Standards Track RFC that allocates a new attribute number value and defines the encoding for the attribute. See [RFC8178] for further details

## 11.2. Changes in the Categorization of File Attributes

The categorization of file attributes appearing in this specification differs from that previously published for a number of reasons:

- \* The description of the attributes for which support is not REQUIRED no longer uses the RFC2119 keyword "RECOMMENDED" as this is not in accord with the definition of that term in [RFC2119].

We now describe such attributes as OPTIONAL, leaving it to server to decide which are worthy of support and to clients to decide whether they wish to use servers on which they are not supported.

- \* The categorization of requirements/recommendation as to support for authorization-related attributes is now the responsibility of the NFSv4-wide security documents, to be derived from [I-D.dnoveck-nfsv4-security] and [I-D.ietf-nfsv4-acls-update].

Currently, given the likely lack of agreement on the semantics of ACLs, it is likely that acl would best be described as an Experimental attribute. See Section 11.3 for further discussion.

As one illustration of the new approach to these matters, and its differences from older approaches, let us consider the following statement from Section 5.1 of [RFC8881]. Referring to the REQUIRED attributes, it states:

The client is expected to be able to function with an attribute set limited to these attributes. With just the REQUIRED attributes some client functionality may be impaired or limited in some ways. In the case of servers not supporting the owner, mode, or acl-related attributes, there would be no ability to provide substantial security-related functionality.

This expectation was not a reasonable one when first formulated and as the NFSv4 protocols have been developed, there have never been any cases of it being realized. There is no reason to implement a server without the minimal authorization-related attributes derived from NFSv3 and no point in working to develop clients capable of interoperating with it. There is no motivation for the working group to devote any time to defining how such a combination is to operate or for implementers to experiment to try to implement remote file access without any meaningful authorization process.

Further, the above also seems to conflict with the following, appearing in Section 5.2 of [RFC8881]:

It is expected that servers will support all attributes they comfortably can and only fail to support attributes that are difficult to support in their operating environments.

Together, these imply that there are operating environments in which it difficult to support all of mode, owner, group, and acl attributes. It is hard to believe that any such environments exist or that there would be any point in implementing an NFSv4.1 server using then, if they did exist.

### 11.3. Categorization of Authorization-related Attributes

This section provides an overview of the issues involved in appropriately categorizing the authorization-related attributes, although the final categorization of these will appear in NFSv4-wide security documents, expected to be based on [I-D.dnoveck-nfsv4-security] and [I-D.ietf-nfsv4-acls-update].

Authorization-related attributes that are part of NFSv4.1 can be divided into those connected to the POSIX-based authorization model used in NFSv3 and those related to the use of ACLs to provide a more flexible authorization model. Within the context of NFSv4.1, the following should be noted:

- \* The attributes mode, owner, and owner\_group need to be considered REQUIRED, as they are in [I-D.dnoveck-nfsv4-security].

This is despite the fact that previous specifications have considered these attributes as OPTIONAL, although the word "RECOMMENDED" was sometimes used. In any case, the new categorization in [I-D.ietf-nfsv4-acls-update] has to be considered dispositive both with regard to NFSv4.1 and other minor versions.

- \* The attributes acl, sacl, and dacl, although designated as OPTIONAL, have never been documented in a manner allowing effective client-server interoperability, suggesting that they would more appropriately be designated as "Experimental".

While it is possible that tightening of the specifications being done in [I-D.dnoveck-nfsv4-security] and [I-D.ietf-nfsv4-acls-update] as part of the rfc8881bis effort might allow this to change, that is not yet assured

In any case, efforts to provide a path to interoperability will continue and might affect this categorization in later minor versions, even if NFSv4.1 is not affected. See [I-D.ietf-nfsv4-acls-update] for details,

- \* The attribute aclsupport is appropriately designated as OPTIONAL, as it is in [I-D.dnoveck-nfsv4-security].

#### 11.4. REQUIRED Attributes

These MUST be supported by every NFSv4.1 client and server in order to ensure a minimum level of interoperability. The server MUST store and return these attributes when requested. A client may ask for the value of any of these attributes to be returned by setting a bit in the GETATTR request, and the server MUST return their value.

The client is expected to be able to function with an attribute set limited to these attributes. With just the REQUIRED attributes some client functionality may be unavailable or functionally limited .

#### 11.5. OPTIONAL Attributes

These attributes are understood well enough to warrant support in the NFSv4.1 protocol. However, they might not be supported on all servers or used by all clients. A client may ask for any of these attributes to be returned by setting a bit in the GETATTR request but need to handle the case where the server does not return them. A client MAY ask for the set of attributes the server supports within a given file system and has no reason to request attributes the server does not support. A server is REQUIRED to be deal with requests for unsupported attributes by not returning values for them rather than

by considering the request an error.

Previous versions of the NFSv4.1 specification [RFC5661] [RFC8881] have described these attributes as "RECOMMENDED" even though that description is not accord with [RFC2119]. The NFSv4.0 specification [RFC7530] still uses "RECOMMENDED" although explicitly disclaiming the assumption that the RFC2119 definition applies in this case. The description of these attribute as OPTIONAL connects them appropriately to provisions for protocol extension and minor versioning in which attributes are to be treated as OPTIONAL.

#### 11.6. Experimental Attributes

While the vast majority of attributes are, as described in Section 11.5, "understood well enough to warrant support in the NFSv4.1 protocol", it appears to be the case that, for several attributes, that understanding was never properly recorded in existing NFSv4.1 specification documents. While it might be possibly to rectify that issue before eventual publication of this document, the likely existence of multiple incompatible implementations of such attributes make that unlikely

Although the existence of such attributes has never been acknowledged before as part of the categorization of NFSv4 attributes. Nevertheless, such attributes have existed in all NFSv4 minor versions and the necessary clarification, if it occurs, is not likely to be complete for some time.

While the intention has always been that attribute not be included in Proposed Standards unless they are described adequately to allow interoperable implementation to be developed. Despite that intention, such attributes have been included in multiple minor versions. Given the need to correct that situation, we need to be clear about the issues that have led to these unfortunate situations, so that we can, over time, address them.

#### 11.7. Named Attributes

These attributes are not supported by direct encoding in the NFSv4 protocol but are accessed using string names rather than numbers and each corresponds to an uninterpreted stream of bytes that is stored in its own file system object. The namespace for these attributes may be accessed by using the OPENATTR operation as described below. The OPENATTR operation returns a filehandle for a "named attribute directory", and further perusal and modification of the namespace may be done using operations that work on more typical directories, subject to restrictions discussed below. In particular, REaddir may be used to get a list of such named attributes, and LOOKUP and OPEN



may select a particular attribute. Creation of a new named attribute can be accomplished using an OPEN specifying file creation.

OPENATTR takes a filehandle for the object and returns the filehandle for the attribute directory. The filehandle for the named attributes designates a directory object accessible by LOOKUP or READDIR and contains files whose names identify the named attributes and whose data bytes are the value of those attributes. For example:

LOOKUP	"foo"	; look up file
GETATTR	attrbits	
OPENATTR		; access foo's named attributes
LOOKUP	"xllicon"	; look up specific attribute
READ	0,4096	; read stream of bytes

Table 3

Named attributes are intended for data needed by applications rather than by NFS client implementations. NFS implementers who wish to define new attributes need to specify them as OPTIONAL attributes using the protocol extension facilities specified in [RFC8178].

Once an OPEN is done, named attributes may be examined and changed using READ and WRITE operations referencing the filehandles and stateids returned by OPEN.

Named attributes may have their own (non-named) attributes. Each of these objects MUST have all of the REQUIRED attributes and may have additional attributes which are not REQUIRED. However, the sets of supported attributes for named attributes need not be, and typically will not be, as large as that for other objects in that file system. Nevertheless, the value of the supported\_attrs attribute should reflect the supported attributes for the file system and will not reflect the restricted attribute sets for these special objects.

Named attributes and the named attribute directories can be the target of delegations (in the case of the named attribute directory, these will be directory delegations). However, since granting of delegations is at the server's discretion, a server need not support delegations on named attributes or on named attribute directories.

Support for named attributes is OPTIONAL and clients need to be prepared to deal with servers that do not support them. However, clients are entitled to assume that if OPENATTR is supported, there will be support for arbitrarily named attributes, rather than support for a few specific names known to the server. If a server does support named attributes, a client that is also able to handle them should be able to copy a file's data and metadata with complete transparency from one location to another since names allowed for regular directory entries are expected to be valid for named attribute names as well.

In NFSv4.1, the structure of named attribute directories is restricted in a number of ways, in order to prevent the development of non-interoperable implementations in which some servers support a fully general hierarchical directory structure for named attributes while others support a limited, non-hierarchical structure for named attributes. In such a mixed environment, clients or applications might come to depend on non-portable extensions. The restrictions are:

- \* CREATE is not allowed in a named attribute directory. Thus, such objects as symbolic links and special files are not allowed to be named attributes. Further, directories may not be created in a named attribute directory, so a hierarchical structure of named attributes for a single object is not allowed.
- \* If OPENATTR is done on a named attribute directory or on a named attribute, the server MUST return NFS4ERR\_WRONG\_TYPE.
- \* Doing a RENAME of a named attribute to a different named attribute directory or to an ordinary (i.e., non-named-attribute) directory is not allowed.
- \* Creating hard links between named attribute directories or between named attribute directories and ordinary directories is not allowed.

Names of attributes will not be controlled by this document or other IETF Standards Track documents, beyond what is necessary to regulate the names of files within directories to handle internationalization and case-insensitivity. See Section 29.2 for further discussion.

## 11.8. Classification of Attributes

Each of the protocol-defined attributes can be classified in one of three categories: per server (i.e., the value of the attribute will be the same for all file objects that share the same server owner; see Section 5.6 for a definition of server owner), per file system (i.e., the value of the attribute will be the same for some or all file objects that share the same fsid attribute (Section 11.12.1.9) and server owner), or per file system object. Note that it is possible that some per file system attributes may vary within the file system, depending on the value of the "homogeneous" (Section 11.12.2.16) attribute. Note that the attributes `time_access_set` and `time_modify_set` are not listed in this section because they are write-only attributes corresponding to `time_access` and `time_modify`, and are used in a special instance of `SETATTR`.

- \* The per-server attribute is:

`lease_time`

- \* The per-file system attributes are:

`supported_attrs`, `suppattr_exclcreat`, `fh_expire_type`,  
`link_support`, `symlink_support`, `unique_handles`, `aclsupport`,  
`cansettime`, `case_insensitive`, `case_preserving`,  
`chown_restricted`, `files_avail`, `files_free`, `files_total`,  
`fs_locations`, `homogeneous`, `maxfilesize`, `maxname`, `maxread`,  
`maxwrite`, `no_trunc`, `space_avail`, `space_free`, `space_total`,  
`time_delta`, `change_policy`, `fs_status`, `fs_layout_type`,  
`fs_locations_info`, `fs_charset_cap`

- \* The per-file system object attributes are:

`type`, `change`, `size`, `named_attr`, `fsid`, `rdattr_error`, `filehandle`,  
`acl`, `archive`, `fileid`, `hidden`, `maxlink`, `mimetype`, `mode`,  
`numlinks`, `owner`, `owner_group`, `rawdev`, `space_used`, `system`,  
`time_access`, `time_backup`, `time_create`, `time_metadata`,  
`time_modify`, `mounted_on_fileid`, `dir_notif_delay`,  
`dirent_notif_delay`, `dacl`, `sacl`, `layout_type`, `layout_hint`,  
`layout_blksize`, `layout_alignment`, `mdsthreshold`, `retention_get`,  
`retention_set`, `retentevt_get`, `retentevt_set`, `retention_hold`,  
`mode_set_masked`

For `quota_avail_hard`, `quota_avail_soft`, and `quota_used`, see their definitions below for the appropriate classification.

### 11.9. Set-Only and Get-Only Attributes

Some of the protocol-defined attributes are set-only; i.e., they can be set via SETATTR but not retrieved via GETATTR. Similarly, some protocol-defined attributes are get-only; i.e., they can be retrieved via GETATTR but not set via SETATTR. If a client attempts to set a get-only attribute or get a set-only attributes, the server MUST return NFS4ERR\_INVALID.

### 11.10. REQUIRED Attributes - List and Definition References

The list of REQUIRED attributes appears in Table 4. The meaning of the columns of the table are:

Name: The name of the attribute.

Id: The number assigned to the attribute. In the event of conflicts between the number assigned here and in RFCTBD30, the latter is likely authoritative, but should be resolved with Errata to this document and/or RFCTBD30. See [errata] for the Errata process.

Data Type: The XDR data type of the attribute.

Acc: Access allowed to the attribute. R means read-only (GETATTR may retrieve, SETATTR may not set). W means write-only (SETATTR may set, GETATTR may not retrieve). R W means read/write (GETATTR may retrieve, SETATTR may set).

Defined in: The section of this specification that describes the attribute.

Name	Id	Data Type	Acc	Defined in:
supported_attrs	0	bitmap4	R	Section 11.12.1.1
type	1	nfs_ftype4	R	Section 11.12.1.2
fh_expire_type	2	uint32_t	R	Section 11.12.1.3
change	3	uint64_t	R	Section 11.12.1.4
size	4	uint64_t	R W	Section 11.12.1.5

link_support	5	bool	R	Section 11.12.1.6
symlink_support	6	bool	R	Section 11.12.1.7
named_attr	7	bool	R	Section 11.12.1.8
fsid	8	fsid4	R	Section 11.12.1.9
unique_handles	9	bool	R	Section 11.12.1.10
lease_time	10	nfs_lease4	R	Section 11.12.1.11
rdattr_error	11	enum	R	Section 11.12.1.12
filehandle	19	nfs_fh4	R	Section 11.12.1.13
mode	33	mode4	R W	Section 11.18
owner	36	utf8str_mixed	R W	Section 11.18
owner_group	37	utf8str_mixed	R W	Section 11.18
suppattr_exclcreat	75	bitmap4	R	Section 11.12.1.14

Table 4

#### 11.11. OPTIONAL Attributes - List and Definition References

The OPTIONAL attributes are defined in Table 5. The meanings of the column headers are the same as Table 4; see Section 11.10 for the meanings.

Name	Id	Data Type	Acc	Defined in:
acl	12	nfsace4<>	R W	Section 11.18
aclsupport	13	uint32_t	R	Section 11.18
archive	14	bool	R W	Section 11.12.2.1
cansettime	15	bool	R	Section 11.12.2.2
case_insensitive	16	bool	R	Section 11.12.2.3
case_preserving	17	bool	R	Section 11.12.2.4
change_policy	60	chg_policy4	R	Section 11.12.2.5
chown_restricted	18	bool	R	Section 11.12.2.6
dacl	58	nfsacl41	R W	Section 11.18
dir_notif_delay	56	nfstime4	R	Section 11.15.1
dirent_notif_delay	57	nfstime4	R	Section 11.15.2
fileid	20	uint64_t	R	Section 11.12.2.7
files_avail	21	uint64_t	R	Section 11.12.2.8
files_free	22	uint64_t	R	Section 11.12.2.9
files_total	23	uint64_t	R	Section 11.12.2.10

fs_charset_cap	76	uint32_t	R	Section 11.12.2.11
fs_layout_type	62	layouttype4<>	R	Section 11.16.1
fs_locations	24	fs_locations	R	Section 11.12.2.12
fs_locations_info	67	fs_locations_info4	R	Section 11.12.2.13
fs_status	61	fs4_status	R	Section 11.12.2.14
hidden	25	bool	R W	Section 11.12.2.15
homogeneous	26	bool	R	Section 11.12.2.16
layout_alignment	66	uint32_t	R	Section 11.16.2
layout_blksize	65	uint32_t	R	Section 11.16.3
layout_hint	63	layouthint4	W	Section 11.16.4
layout_type	64	layouttype4<>	R	Section 11.16.5
maxfilesize	27	uint64_t	R	Section 11.12.2.17
maxlink	28	uint32_t	R	Section 11.12.2.18
maxname	29	uint32_t	R	Section 11.12.2.19
maxread	30	uint64_t	R	Section 11.12.2.20
maxwrite	31	uint64_t	R	Section 11.12.2.21

mdsthreshold	68	mdsthreshold4	R	Section 11.16.6
mimetype	32	utf8str_cs	R W	Section 11.12.2.22
mode	33	mode4	R W	Section 11.18
mode_set_masked	74	mode_masked4	W	Section 11.18
mounted_on_fileid	55	uint64_t	R	Section 11.12.2.23
no_trunc	34	bool	R	Section 11.12.2.24
numlinks	35	uint32_t	R	Section 11.12.2.25
quota_avail_hard	38	uint64_t	R	Section 11.12.2.26
quota_avail_soft	39	uint64_t	R	Section 11.12.2.27
quota_used	40	uint64_t	R	Section 11.12.2.28
rawdev	41	specdata4	R	Section 11.12.2.29
retentevt_get	71	retention_get4	R	Section 11.17.3
retentevt_set	72	retention_set4	W	Section 11.17.4
retention_get	69	retention_get4	R	Section 11.17.1
retention_hold	73	uint64_t	R W	Section 11.17.5
retention_set	70	retention_set4	W	Section 11.17.2



sacl	59	nfsacl41	R W	Section 11.18
space_avail	42	uint64_t	R	Section 11.12.2.30
space_free	43	uint64_t	R	Section 11.12.2.31
space_total	44	uint64_t	R	Section 11.12.2.32
space_used	45	uint64_t	R	Section 11.12.2.33
system	46	bool	R W	Section 11.12.2.34
time_access	47	nfstime4	R	Section 11.12.2.35
time_access_set	48	settime4	W	Section 11.12.2.36
time_backup	49	nfstime4	R W	Section 11.12.2.37
time_create	50	nfstime4	R W	Section 11.12.2.38
time_delta	51	nfstime4	R	Section 11.12.2.39
time_metadata	52	nfstime4	R	Section 11.12.2.40
time_modify	53	nfstime4	R	Section 11.12.2.41
time_modify_set	54	settime4	W	Section 11.12.2.42

Table 5

## 11.12. Attribute Definitions

### 11.12.1. Definitions of REQUIRED Attributes

#### 11.12.1.1. Attribute 0: supported\_attrs

The bit vector that would retrieve all protocol-defined attributes that are supported for this object. The scope of this attribute applies to all objects with a matching fsid.

#### 11.12.1.2. Attribute 1: type

Designates the type of an object in terms of one of a number of special constants:

- \* NF4REG designates a regular file.
- \* NF4DIR designates a directory.
- \* NF4BLK designates a block device special file.
- \* NF4CHR designates a character device special file.
- \* NF4LNK designates a symbolic link.
- \* NF4SOCK designates a named socket special file.
- \* NF4FIFO designates a fifo special file.
- \* NF4ATTRDIR designates a named attribute directory.
- \* NF4NAMEDATTR designates a named attribute.

Within the explanatory text and operation descriptions, the following phrases will be used with the meanings given below:

- \* The phrase "is a directory" means that the object's type attribute is NF4DIR or NF4ATTRDIR.
- \* The phrase "is a special file" means that the object's type attribute is NF4BLK, NF4CHR, NF4SOCK, or NF4FIFO.
- \* The phrases "is an ordinary file" and "is a regular file" mean that the object's type attribute is NF4REG or NF4NAMEDATTR.

#### 11.12.1.3. Attribute 2: fh\_expire\_type

Server uses this to specify filehandle expiration behavior to the client. See Section 10 for additional description.

#### 11.12.1.4. Attribute 3: change

A value created by the server that the client can use to determine if file data, directory contents, or attributes of the object have been modified. The server may return the object's time\_metadata attribute for this attribute's value, but only if the file system object cannot be updated more frequently than the resolution of time\_metadata.

#### 11.12.1.5. Attribute 4: size

The size of the object in bytes.

#### 11.12.1.6. Attribute 5: link\_support

TRUE, if the object's file system supports hard links.

#### 11.12.1.7. Attribute 6: symlink\_support

TRUE, if the object's file system supports symbolic links.

#### 11.12.1.8. Attribute 7: named\_attr

TRUE, if this object has named attributes. In other words, object has a non-empty named attribute directory.

#### 11.12.1.9. Attribute 8: fsid

Unique file system identifier for the file system holding this object. The fsid attribute has major and minor components, each of which are of data type uint64\_t.

#### 11.12.1.10. Attribute 9: unique\_handles

TRUE, if two distinct filehandles are guaranteed to refer to two different file system objects.

#### 11.12.1.11. Attribute 10: lease\_time

Duration of the lease at server in seconds.

#### 11.12.1.12. Attribute 11: rdattrib\_error

Error returned from an attempt to retrieve attributes during a READDIR operation.

#### 11.12.1.13. Attribute 19: filehandle

The filehandle of this object (primarily for use by READDIR requests).

#### 11.12.1.14. Attribute 75: suppattrib\_exclcreat

The bit vector that would set all protocol-defined attributes that are supported by the EXCLUSIVE4\_1 method of file creation via the OPEN operation. The scope of this attribute applies to all objects with a matching fsid.

### 11.12.2. Definitions of Uncategorized OPTIONAL Attributes

The definitions of most of the OPTIONAL attributes follow. Collections that share a common category are defined in other sections.

#### 11.12.2.1. Attribute 14: archive

TRUE, if this file has been archived since the time of last modification (deprecated in favor of time\_backup).

#### 11.12.2.2. Attribute 15: cansettime

TRUE, if the server is able to change the times for a file system object as specified in a SETATTR operation.

#### 11.12.2.3. Attribute 16: case\_insensitive

TRUE, if file name comparisons on this file system are case insensitive.

#### 11.12.2.4. Attribute 17: case\_preserving

TRUE, if file name case on this file system is preserved.

#### 11.12.2.5. Attribute 60: `change_policy`

A value created by the server that the client can use to determine if some server policy related to the current file system has been subject to change. If the value remains the same, then the client can be sure that the values of the attributes related to `fs_location` and the `fss_type` field of the `fs_status` attribute have not changed. On the other hand, a change in this value does necessarily imply a change in policy. It is up to the client to interrogate the server to determine if some policy relevant to it has changed. See Section 9.3.6 for details.

This attribute **MUST** change when the value returned by the `fs_locations` or `fs_locations_info` attribute changes, when a file system goes from read-only to writable or vice versa, or when the allowable set of security flavors for the file system or any part thereof is changed.

#### 11.12.2.6. Attribute 18: `chown_restricted`

If **TRUE**, the server will reject any request to change either the owner or the group associated with a file if the caller is not a privileged user (for example, "root" in UNIX operating environments or, in Windows 2000, the "Take Ownership" privilege).

#### 11.12.2.7. Attribute 20: `fileid`

A number uniquely identifying the file within the file system.

#### 11.12.2.8. Attribute 21: `files_avail`

File slots available to this user on the file system containing this object -- this should be the smallest relevant limit.

#### 11.12.2.9. Attribute 22: `files_free`

Free file slots on the file system containing this object -- this should be the smallest relevant limit.

#### 11.12.2.10. Attribute 23: `files_total`

Total file slots on the file system containing this object.

#### 11.12.2.11. Attribute 76: `fs_charset_cap`

Character set capabilities for this file system. See Section 21.1.

#### 11.12.2.12. Attribute 24: fs\_locations

Locations where this file system may be found. If the server returns NFS4ERR\_MOVED as an error, this attribute MUST be supported. See Section 17.16 for more details.

#### 11.12.2.13. Attribute 67: fs\_locations\_info

Full function file system location. See Section 17.17.2 for more details.

#### 11.12.2.14. Attribute 61: fs\_status

Generic file system type information. See Section 17.18 for more details.

#### 11.12.2.15. Attribute 25: hidden

TRUE, if the file is considered hidden with respect to the Windows API.

#### 11.12.2.16. Attribute 26: homogeneous

TRUE, if this object's file system is homogeneous; i.e., all objects in the file system (all objects on the server with the same fsid) have common values for all per-file-system attributes.

#### 11.12.2.17. Attribute 27: maxfilesize

Maximum supported file size for the file system of this object.

#### 11.12.2.18. Attribute 28: maxlink

Maximum number of links for this object.

#### 11.12.2.19. Attribute 29: maxname

Maximum file name size supported for this object.

#### 11.12.2.20. Attribute 30: maxread

Maximum amount of data the READ operation will return for this object.

#### 11.12.2.21. Attribute 31: maxwrite

Maximum amount of data the WRITE operation will accept for this object. This attribute should be supported if the file is writable. Lack of this attribute can lead to the client either wasting bandwidth or not receiving the best performance.

#### 11.12.2.22. Attribute 32: mimetype

MIME body type/subtype of this object.

#### 11.12.2.23. Attribute 55: mounted\_on\_fileid

Like fileid, but if the target filehandle is the root of a file system, this attribute represents the fileid of the underlying directory.

UNIX-based operating environments connect a file system into the namespace by connecting (mounting) the file system onto the existing file object (the mount point, usually a directory) of an existing file system. When the mount point's parent directory is read via an API like readdir(), the return results are directory entries, each with a component name and a fileid. The fileid of the mount point's directory entry will be different from the fileid that the stat() system call returns. The stat() system call is returning the fileid of the root of the mounted file system, whereas readdir() is returning the fileid that stat() would have returned before any file systems were mounted on the mount point.

Unlike NFSv3, NFSv4.1 allows a client's LOOKUP request to cross other file systems. The client detects the file system crossing whenever the filehandle argument of LOOKUP has an fsid attribute different from that of the filehandle returned by LOOKUP. A UNIX-based client will consider this a "mount point crossing". UNIX has a legacy scheme for allowing a process to determine its current working directory. This relies on readdir() of a mount point's parent and stat() of the mount point returning fileids as previously described. The mounted\_on\_fileid attribute corresponds to the fileid that readdir() would have returned as described previously.

While the NFSv4.1 client could simply fabricate a fileid corresponding to what mounted\_on\_fileid provides (and if the server does not support mounted\_on\_fileid, the client has no choice), there is a risk that the client will generate a fileid that conflicts with one that is already assigned to another object in the file system. Instead, if the server can provide the mounted\_on\_fileid, the potential for client operational problems in this area is eliminated.

If the server detects that there is no mounted point at the target file object, then the value for `mounted_on_fileid` that it returns is the same as that of the `fileid` attribute.

The `mounted_on_fileid` attribute is OPTIONAL, and the server should provide it if possible. For a UNIX-based server, this is straightforward. Usually, `mounted_on_fileid` will be requested as part of a `READDIR` operation, in which case it is trivial (at least for UNIX-based servers) to return `mounted_on_fileid` since it is equal to the `fileid` of a directory entry returned by `readdir()`. If `mounted_on_fileid` is requested in a `GETATTR` operation, the server should obey an invariant that has it returning a value that is equal to the file object's entry in the object's parent directory, i.e., what `readdir()` would have returned. Some operating environments allow a series of two or more file systems to be mounted onto a single mount point. In this case, for the server to obey the aforementioned invariant, it will need to find the base mount point, and not the intermediate mount points.

#### 11.12.2.24. Attribute 34: `no_trunc`

If this attribute is `TRUE`, then if the client uses a file name longer than `name_max`, an error will be returned instead of the name being truncated.

#### 11.12.2.25. Attribute 35: `numlinks`

Number of hard links to this object.

#### 11.12.2.26. Attribute 38: `quota_avail_hard`

The value in bytes that represents the amount of additional disk space beyond the current allocation that can be allocated to this file or directory before further allocations will be refused. It is understood that this space may be consumed by allocations to other files or directories.

#### 11.12.2.27. Attribute 39: `quota_avail_soft`

The value in bytes that represents the amount of additional disk space that can be allocated to this file or directory before the user may reasonably be warned. It is understood that this space may be consumed by allocations to other files or directories though there is a rule as to which other files or directories.



#### 11.12.2.28. Attribute 40: quota\_used

The value in bytes that represents the amount of disk space used by this file or directory and possibly a number of other similar files or directories, where the set of "similar" meets at least the criterion that allocating space to any file or directory in the set will reduce the "quota\_avail\_hard" of every other file or directory in the set.

Note that there may be a number of distinct but overlapping sets of files or directories for which a quota\_used value is maintained, e.g., "all files with a given owner", "all files with a given group owner", etc. The server is at liberty to choose any of those sets when providing the content of the quota\_used attribute, but should do so in a repeatable way. The rule may be configured per file system or may be "choose the set with the smallest quota".

#### 11.12.2.29. Attribute 41: rawdev

Raw device number of file of type NF4BLK or NF4CHR. The device number is split into major and minor numbers. If the file's type attribute is not NF4BLK or NF4CHR, the value returned is bot considered useful.

#### 11.12.2.30. Attribute 42: space\_avail

Disk space in bytes available to this user on the file system containing this object -- this should be the smallest relevant limit.

#### 11.12.2.31. Attribute 43: space\_free

Free disk space in bytes on the file system containing this object -- this should be the smallest relevant limit.

#### 11.12.2.32. Attribute 44: space\_total

Total disk space in bytes on the file system containing this object.

#### 11.12.2.33. Attribute 45: space\_used

Number of file system bytes allocated to this object.

#### 11.12.2.34. Attribute 46: system

This attribute is TRUE if this file is a "system" file with respect to the Windows operating environment.

#### 11.12.2.35. Attribute 47: time\_access

The time\_access attribute represents the time of last access to the object by a READ operation sent to the server. The notion of what is an "access" depends on the server's operating environment and/or the server's file system semantics. For example, for servers obeying Portable Operating System Interface (POSIX) semantics, time\_access would be updated only by the READ and READDIR operations and not any of the operations that modify the content of the object [read\_atime], [readdir\_atime], [write\_atime]. Of course, setting the corresponding time\_access\_set attribute is another way to modify the time\_access attribute.

Whenever the file object resides on a writable file system, the server should make its best efforts to record time\_access into stable storage. However, to mitigate the performance effects of doing so, and most especially whenever the server is satisfying the read of the object's content from its cache, the server MAY cache access time updates and lazily write them to stable storage. It is also acceptable to give administrators of the server the option to disable time\_access updates.

#### 11.12.2.36. Attribute 48: time\_access\_set

Sets the time of last access to the object. SETATTR use only.

#### 11.12.2.37. Attribute 49: time\_backup

The time of last backup of the object.

#### 11.12.2.38. Attribute 50: time\_create

The time of creation of the object. This attribute does not have any relation to the traditional UNIX file attribute "ctime" or "change time".

#### 11.12.2.39. Attribute 51: time\_delta

Smallest useful server time granularity.

#### 11.12.2.40. Attribute 52: time\_metadata

The time of last metadata modification of the object.

#### 11.12.2.41. Attribute 53: time\_modify

The time of last modification to the object.

#### 11.12.2.42. Attribute 54: time\_modify\_set

Sets the time of last modification to the object. SETATTR use only.

#### 11.13. Interpreting owner and owner\_group

The attributes "owner" and "owner\_group" (and also users and groups within the "acl" attribute) are transferred in the form of a UTF-8 string. This string can be used to identify users and groups in several ways:

- \* A string of the form "name@domain" can be used to give a user or group name together with a domain in which those are defined.

This form provides greater degree of extensibility than was possible in NFSv3 which limited these identifiers to 32-bit unsigned integers whose values are all centrally administered as members within a common domain.

- \* Numeric ids converted to string form.

Using this format maintains the strengths and weaknesses of the NFSv3 approach.

The following issues are relevant in selected the form to use.

- \* The use of the form "name@domain" provides greater flexibility, both with regard to the number of users that can be accommodated and to the management of multiple sets of users in separate domains.

Taking advantage of this flexibility often requires extensive work because of limitations of the API's used to reference users and groups.

- \* The use of the form "name@domain" allows clients and servers to work together even if they have different internal formats for user and groups.

In many cases, there is no need for such mapping.

Providing this mapping requires extra implementation and raises potential security issues.

For detailed discussions regarding which of the forms clients and server are to use for these values, see Section 5.1 of [I-D.dnoveck-nfsv4-security].

#### 11.14. Character Case Attributes

With respect to the `case_insensitive` and `case_preserving` attributes, each UCS-4 character (which UTF-8 encodes) can be mapped to an equivalent character of different case or compared in a case-insensitive manner. The details vary based on the Unicode version implemented by the server for the current file system. Details of the process and how the client can best deal with uncertainty about the process will be discussed in the NFSv4-wide internationalization document (See [I-D.ietf-nfsv4-internationalization] for the latest version)

#### 11.15. Directory Notification Attributes

As described in Section 25.39, the client can request a minimum delay for notifications of changes to attributes, but the server is free to ignore what the client requests. The client can determine in advance what notification delays the server will accept by sending a `GETATTR` operation for either or both of two directory notification attributes. When the client calls the `GET_DIR_DELEGATION` operation and asks for attribute change notifications, it should request notification delays that are no less than the values in the server-provided attributes.

##### 11.15.1. Attribute 56: `dir_notif_delay`

The `dir_notif_delay` attribute is the minimum number of seconds the server will delay before notifying the client of a change to the directory's attributes.

##### 11.15.2. Attribute 57: `dirent_notif_delay`

The `dirent_notif_delay` attribute is the minimum number of seconds the server will delay before notifying the client of a change to a file object that has an entry in the directory.

#### 11.16. pNFS Attribute Definitions

##### 11.16.1. Attribute 62: `fs_layout_type`

The `fs_layout_type` attribute (see Section 9.3.13) applies to a file system and indicates what layout types are supported by the file system. In this context, "supported" means that the server on which the interrogation is done is able to act as the metadata server for the specified layout type within that file system by providing layouts of that layout type in response to `GETLAYOUT`.

When the client encounters a new fsid, the client should obtain the value for the `fs_layout_type` attribute associated with the new file system. This attribute is used by the client to determine if the layout types supported by the server match any of the client's supported layout types.

#### 11.16.2. Attribute 66: `layout_alignment`

When a client holds layouts on files of a file system, the `layout_alignment` attribute indicates the preferred alignment for I/O to files on that file system. Where possible, the client should send READ and WRITE operations with offsets that are whole multiples of the `layout_alignment` attribute.

#### 11.16.3. Attribute 65: `layout_blksize`

When a client holds layouts on files of a file system, the `layout_blksize` attribute indicates the preferred block size for I/O to files on that file system. Where possible, the client should send READ operations with a count argument that is a whole multiple of `layout_blksize`, and WRITE operations with a data argument of size that is a whole multiple of `layout_blksize`.

#### 11.16.4. Attribute 63: `layout_hint`

The `layout_hint` attribute (see Section 9.3.19) may be set on newly created files to influence the metadata server's choice for the file's layout. If possible, this attribute is one of those set in the initial attributes within the OPEN operation. The metadata server may choose to ignore this attribute. The `layout_hint` attribute is a subset of the layout structure returned by LAYOUTGET. For example, instead of specifying particular devices, this would be used to suggest the stripe width of a file. The server implementation determines which fields within the layout will be used.

#### 11.16.5. Attribute 64: `layout_type`

This attribute lists the layout type(s) available for a file. The value returned by the server is for informational purposes only. The client will use the LAYOUTGET operation to obtain the information needed in order to perform I/O, for example, the specific device information for the file and its layout.

#### 11.16.6. Attribute 68: mdsthreshold

This attribute is a server-provided hint used to communicate to the client when it is more efficient to send READ and WRITE operations to the metadata server or the data server. The two types of thresholds described are file size thresholds and I/O size thresholds. If a file's size is smaller than the file size threshold, data accesses are to be sent to the metadata server. If an I/O request has a length that is below the I/O size threshold, the I/O should be sent to the metadata server. Each threshold type is specified separately for read and write.

The server MAY provide both types of thresholds for a file. If both file size and I/O size are provided, the client needs to reach or exceed both thresholds before sending its read or write requests to the data server. Alternatively, if only one of the specified thresholds is reached or exceeded, the I/O requests are sent to the metadata server.

For each threshold type, a value of zero indicates no READ or WRITE should be sent to the metadata server, while a value of all ones indicates that all READs or WRITEs should be sent to the metadata server.

The attribute is available on a per-filehandle basis. If the current filehandle refers to a non-pNFS file or directory, the metadata server should return an attribute that is representative of the filehandle's file system. It is suggested that this attribute is queried as part of the OPEN operation. Due to dynamic system changes, the client should not assume that the attribute will remain constant for any specific time period; thus, it should be periodically refreshed.

#### 11.17. Retention Attributes

Retention is a concept whereby a file object can be placed in an immutable, undeletable, unrenamable state for a fixed or infinite duration of time. Once in this "retained" state, the file cannot be moved out of the state until the duration of retention has been reached.

When retention is enabled, retention MUST extend to the data of the file, and the name of file. The server MAY extend retention to any other property of the file, including any subset of REQUIRED, OPTIONAL, and named attributes, with the exceptions noted in this section.

Servers MAY support or not support retention on any file object type.

The five retention attributes are explained in the next subsections.

#### 11.17.1. Attribute 69: retention\_get

If retention is enabled for the associated file, this attribute's value represents the retention begin time of the file object. This attribute's value is only readable with the GETATTR operation and MUST NOT be modified by the SETATTR operation (Section 11.9). The value of the attribute consists of:

```
const ET4_DURATION_INFINITE    = 0xffffffffffffffff;
struct retention_get4 {
    uint64_t      rg_duration;
    nfstime4      rg_begin_time<1>;
};
```

The field `rg_duration` is the duration in seconds indicating how long the file will be retained once retention is enabled. The field `rg_begin_time` is an array of up to one absolute time value. If the array is zero length, no beginning retention time has been established, and retention is not enabled. If `rg_duration` is equal to `RET4_DURATION_INFINITE`, the file, once retention is enabled, will be retained for an infinite duration.

If (as soon as) `rg_duration` is zero, then `rg_begin_time` will be of zero length, and again, retention is not (no longer) enabled.

#### 11.17.2. Attribute 70: retention\_set

This attribute is used to set the retention duration and optionally enable retention for the associated file object. This attribute is only modifiable via the SETATTR operation and MUST NOT be retrieved by the GETATTR operation (Section 11.9). This attribute corresponds to `retention_get`. The value of the attribute consists of:

```
struct retention_set4 {
    bool          rs_enable;
    uint64_t      rs_duration<1>;
};
```

If the client sets `rs_enable` to `TRUE`, then it is enabling retention on the file object with the begin time of retention starting from the server's current time and date. The duration of the retention can also be provided if the `rs_duration` array is of length one. The duration is the time in seconds from the begin time of retention, and if set to `RET4_DURATION_INFINITE`, the file is to be retained forever. If retention is enabled, with no duration specified in either this SETATTR or a previous SETATTR, the duration defaults to zero seconds.

The server MAY restrict the enabling of retention or the duration of retention on the basis of the ACE4\_WRITE\_RETENTION ACL permission. The enabling of retention MUST NOT prevent the enabling of event-based retention or the modification of the retention\_hold attribute.

The following rules apply to both the retention\_set and retentevt\_set attributes.

- \* As long as retention is not enabled, the client is permitted to decrease the duration.
- \* The duration can always be set to an equal or higher value, even if retention is enabled. Note that once retention is enabled, the actual duration (as returned by the retention\_get or retentevt\_get attributes; see Section 11.17.1 or Section 11.17.3) is constantly counting down to zero (one unit per second), unless the duration was set to RET4\_DURATION\_INFINITE. Thus, it will not be possible for the client to precisely extend the duration on a file that has retention enabled.
- \* While retention is enabled, attempts to disable retention or decrease the retention's duration MUST fail with the error NFS4ERR\_INVALID.
- \* If the principal attempting to change retention\_set or retentevt\_set does not have ACE4\_WRITE\_RETENTION permissions, the attempt MUST fail with NFS4ERR\_ACCESS.

#### 11.17.3. Attribute 71: retentevt\_get

Gets the event-based retention duration, and if enabled, the event-based retention begin time of the file object. This attribute is like retention\_get, but refers to event-based retention. The event that triggers event-based retention is not defined by the NFSv4.1 specification.

#### 11.17.4. Attribute 72: retentevt\_set

Sets the event-based retention duration, and optionally enables event-based retention on the file object. This attribute corresponds to retentevt\_get and is like retention\_set, but refers to event-based retention. When event-based retention is set, the file MUST be retained even if non-event-based retention has been set, and the duration of non-event-based retention has been reached. Conversely, when non-event-based retention has been set, the file MUST be retained even if event-based retention has been set, and the duration of event-based retention has been reached. The server MAY restrict the enabling of event-based retention or the duration of event-based



retention on the basis of the ACE4\_WRITE\_RETENTION ACL permission. The enabling of event-based retention MUST NOT prevent the enabling of non-event-based retention or the modification of the retention\_hold attribute.

#### 11.17.5. Attribute 73: retention\_hold

Gets or sets administrative retention holds, one hold per bit position.

This attribute allows one to 64 administrative holds, one hold per bit on the attribute. If retention\_hold is not zero, then the file MUST NOT be deleted, renamed, or modified, even if the duration on enabled event or non-event-based retention has been reached. The server MAY restrict the modification of retention\_hold on the basis of the ACE4\_WRITE\_RETENTION\_HOLD ACL permission. The enabling of administration retention holds does not prevent the enabling of event-based or non-event-based retention.

If the principal attempting to change retention\_hold does not have ACE4\_WRITE\_RETENTION\_HOLD permissions, the attempt MUST fail with NFS4ERR\_ACCESS.

#### 11.18. Access Control Attributes

The use of the access control attributes are fully described in various sections of the NFSv4-wide security documents [I-D.dnoveck-nfsv4-security] [I-D.ietf-nfsv4-acls-update].

- \* The mode, mode\_set\_masked, owner, and owner\_group attributes are described in Sections 5.3.1 through 5.3.4 of [I-D.dnoveck-nfsv4-security].
- \* The acl, aclsupport, sacl, and dacl attributes are described in Sections 3.4, 3.5, 3.6, and 3.8 of [I-D.ietf-nfsv4-acls-update].

#### 12. Single-Server Namespace

This section describes the NFSv4 single-server namespace. Single-server namespaces may be presented directly to clients, or they may be used as a basis to form larger multi-server namespaces (e.g., site-wide or organization-wide) to be presented to clients, as described in Section 17.

### 12.1. Server Exports

On a UNIX server, the namespace describes all the files reachable by pathnames under the root directory or "/". On a Windows server, the namespace constitutes all the files on disks named by mapped disk letters. NFS server administrators rarely make the entire server's file system namespace available to NFS clients. More often, portions of the namespace are made available via an "export" feature. In previous versions of the NFS protocol, the root filehandle for each export is obtained through the MOUNT protocol; the client sent a string that identified the export name within the namespace and the server returned the root filehandle for that export. The MOUNT protocol also provided an EXPORTS procedure that enumerated the server's exports.

### 12.2. Browsing Exports

The NFSv4.1 protocol provides a root filehandle that clients can use to obtain filehandles for the exports of a particular server, via a series of LOOKUP operations within a COMPOUND, to traverse a path. A common user experience is to use a graphical user interface (perhaps a file "Open" dialog window) to find a file via progressive browsing through a directory tree. The client must be able to move from one export to another export via single-component, progressive LOOKUP operations.

This style of browsing is not well supported by the NFSv3 protocol. In NFSv3, the client expects all LOOKUP operations to remain within a single server file system. For example, the device attribute will not change. This prevents a client from taking namespace paths that span exports.

In the case of NFSv3, an automounter on the client can obtain a snapshot of the server's namespace using the EXPORTS procedure of the MOUNT protocol. If it understands the server's pathname syntax, it can create an image of the server's namespace on the client. The parts of the namespace that are not exported by the server are filled in with directories that might be constructed similarly to an NFSv4.1 "pseudo file system" (See Section 12.3) that allows the user to browse from one mounted file system to another. There is a drawback to this representation of the server's namespace on the client: it is static. If the server administrator adds a new export, the client will be unaware of it.

### 12.3. Server Pseudo File System

NFSv4.1 servers avoid this namespace inconsistency by presenting all the exports for a given server within the framework of a single namespace for that server. An NFSv4.1 client uses LOOKUP and READDIR operations to browse seamlessly from one export to another.

Where there are portions of the server namespace that are not exported, clients require some way of traversing those portions to reach actual exported file systems. A technique that servers may use to provide for this is to bridge the unexported portion of the namespace via a "pseudo file system" that provides a view of exported directories only. A pseudo file system has a unique fsid and behaves like a normal, read-only file system.

Based on the construction of the server's namespace, it is possible that multiple pseudo file systems may exist. For example,

/a	pseudo file system
/a/b	real file system
/a/b/c	pseudo file system
/a/b/c/d	real file system

Each of the pseudo file systems is considered a separate entity and therefore MUST have its own fsid, unique among all the fsids for that server.

### 12.4. Multiple Roots

Certain operating environments are sometimes described as having "multiple roots". In such environments, individual file systems are commonly represented by disk or volume names. NFSv4 servers for these platforms can construct a pseudo file system above these root names so that disk letters or volume names are simply directory names in the pseudo root.

### 12.5. Filehandle Volatility

The nature of the server's pseudo file system is that it is a logical representation of file system(s) available from the server. Therefore, the pseudo file system is most likely constructed dynamically when the server is first instantiated. It is expected that the pseudo file system may not have an on-disk counterpart from which persistent filehandles could be constructed. Even though it is preferable that the server provide persistent filehandles for the pseudo file system, the NFS client should expect that pseudo file system filehandles are volatile. This can be confirmed by checking the associated "fh\_expire\_type" attribute for those filehandles in

question. If the filehandles are volatile, the NFS client must be prepared to recover a filehandle value (e.g., with a series of LOOKUP operations) when receiving an error of NFS4ERR\_FHEXPIRED.

Because it is quite likely that servers will implement pseudo file systems using volatile filehandles, clients need to be prepared for them, rather than assuming that all filehandles will be persistent.

#### 12.6. Exported Root

If the server's root file system is exported, one might conclude that a pseudo file system is unneeded. This is not necessarily so. Assume the following file systems on a server:

```
/          fs1  (exported)
/a         fs2  (not exported)
/a/b       fs3  (exported)
```

Because fs2 is not exported, fs3 cannot be reached with simple LOOKUPS. The server must bridge the gap with a pseudo file system.

#### 12.7. Mount Point Crossing

The server file system environment may be constructed in such a way that one file system contains a directory that is 'covered' or mounted upon by a second file system. For example:

```
/a/b          (file system 1)
/a/b/c/d      (file system 2)
```

The pseudo file system for this server may be constructed to look like:

```
/          (place holder/not exported)
/a/b       (file system 1)
/a/b/c/d   (file system 2)
```

It is the server's responsibility to present the pseudo file system that is complete to the client. If the client sends a LOOKUP request for the path /a/b/c/d, the server's response is the filehandle of the root of the file system /a/b/c/d. In previous versions of the NFS protocol, the server would respond with the filehandle of directory /a/b/c/d within the file system /a/b.

The NFS client will be able to determine if it crosses a server mount point by a change in the value of the "fsid" attribute.

## 12.8. Security Policy and Namespace Presentation

Because NFSv4 clients possess the ability to change the security mechanisms used, after determining what is allowed, by using SECINFO and SECINFO\_NO\_NAME, the server should not present a different view of the namespace based on the security mechanism being used by a client. Instead, it should present a consistent view and return NFS4ERR\_WRONGSEC if an attempt is made to access data with an inappropriate security mechanism.

If security considerations make it necessary to hide the existence of a particular file system, as opposed to all of the data within it, the server can apply the security policy of a shared resource in the server's namespace to components of the resource's ancestors. For example:

/	(place holder/not exported)
/a/b	(file system 1)
/a/b/MySecretProject	(file system 2)

The /a/b/MySecretProject directory is a real file system and is the shared resource. Suppose the security policy for /a/b/MySecretProject is Kerberos with integrity and it is desired to limit knowledge of the existence of this file system. In this case, the server should apply the same security policy to /a/b. This allows for knowledge of the existence of a file system to be secured when desirable.

For the case of the use of multiple, disjoint security mechanisms in the server's resources, applying that sort of policy would result in the higher-level file system not being accessible using any security flavor. Therefore, that sort of configuration is not compatible with hiding the existence (as opposed to the contents) from clients using multiple disjoint sets of security flavors.

In other circumstances, a desirable policy is for the security of a particular object in the server's namespace to include the union of all security mechanisms of all direct descendants. A common and convenient practice, unless strong security requirements dictate otherwise, is to make the entire the pseudo file system accessible by all of the valid security mechanisms.

Where there is concern about the security of data on the network, clients should use strong security mechanisms to access the pseudo file system in order to prevent man-in-the-middle attacks.

### 13. State Management

Integrating locking into the NFS protocol necessarily causes it to be stateful. With the inclusion of such features as share reservations, file and directory delegations, recallable layouts, and support for mandatory byte-range locking, the protocol becomes substantially more dependent on proper management of state than the combination of NFS and NLM (Network Lock Manager) [xnfs]. used for locking within previous NFS versions. The new features include expanded locking facilities, which provide some measure of inter-client exclusion, but the state also offers features not readily providable using a stateless model. There are three components to making this state manageable:

- \* clear division between client and server
- \* ability to reliably detect inconsistency in state between client and server
- \* simple and robust recovery mechanisms

In this model, the server owns the state information. The client requests changes in locks and the server responds with the changes made. Non-client-initiated changes in locking state are infrequent. The client receives prompt notification of such changes and can adjust its view of the locking state to reflect the server's changes.

Individual pieces of state created by the server and passed to the client at its request are represented by 128-bit stateids. These stateids may represent a particular open file, a set of byte-range locks held by a particular owner, or a recallable delegation of privileges to access a file in particular ways or at a particular location.

In all cases, there is a transition from the most general information that represents a client as a whole to the eventual lightweight stateid used for most client and server locking interactions. The details of this transition will vary with the type of object but it always starts with a client ID.

#### 13.1. Client and Session ID

A client must establish a client ID (See Section 5.5) and then one or more sessionids (See Section 7) before performing any operations to open, byte-range lock, delegate, or obtain a layout for a file object. Each session ID is associated with a specific client ID, and thus serves as a shorthand reference to an NFSv4.1 client.

For some types of locking interactions, the client will represent some number of internal locking entities called "owners", which normally correspond to processes internal to the client. For other types of locking-related objects, such as delegations and layouts, no such intermediate entities are provided for, and the locking-related objects are considered to be transferred directly between the server and a unitary client.

### 13.2. Stateid Definition

When the server grants a lock of any type (including opens, byte-range locks, delegations, and layouts), it responds with a unique stateid that represents a set of locks (often a single lock) for the same file, of the same type, and sharing the same ownership characteristics. Thus, opens of the same file by different open-owners each have an identifying stateid. Similarly, each set of byte-range locks on a file owned by a specific lock-owner has its own identifying stateid. Delegations and layouts also have associated stateids by which they may be referenced. The stateid is used as a shorthand reference to a lock or set of locks, and given a stateid, the server can determine the associated state-owner or state-owners (in the case of an open-owner/lock-owner pair) and the associated filehandle. When stateids are used, the current filehandle must be the one associated with that stateid.

All stateids associated with a given client ID are associated with a common lease that represents the claim of those stateids and the objects they represent to be maintained by the server. See Section 13.3 for a discussion of the lease.

The server may assign stateids independently for different clients. A stateid with the same bit pattern for one client may designate an entirely different set of locks for a different client. The stateid is always interpreted with respect to the client ID associated with the current session. Stateids apply to all sessions associated with the given client ID, and the client may use a stateid obtained from one session on another session associated with the same client ID.

#### 13.2.1. Stateid Types

With the exception of special stateids (See Section 13.2.3), each stateid represents locking objects of one of a set of types defined by the NFSv4.1 protocol. Note that in all these cases, where we speak of guarantee, it is understood there are situations such as a client restart, or lock revocation, that allow the guarantee to be voided.

- \* Stateids may represent opens of files.

Each stateid in this case represents the OPEN state for a given client ID/open-owner/filehandle triple. Such stateids are subject to change (with consequent incrementing of the stateid's seqid) in response to OPENS that result in upgrade and OPEN\_DOWNGRADE operations.

- \* Stateids may represent sets of byte-range locks.

All locks held on a particular file by a particular owner and gotten under the aegis of a particular open file are associated with a single stateid with the seqid being incremented whenever LOCK and LOCKU operations affect that set of locks.

- \* Stateids may represent file delegations, which are recallable guarantees by the server to the client that other clients will not reference or modify a particular file, until the delegation is returned. In NFSv4.1, file delegations may be obtained on both regular and non-regular files.

A stateid represents a single delegation held by a client for a particular filehandle.

- \* Stateids may represent directory delegations, which are recallable guarantees by the server to the client that other clients will not modify the directory without appropriate notice to the holder, until the delegation is returned.

A stateid represents a single delegation held by a client for a particular directory filehandle.

- \* Stateids may represent layouts, which are recallable guarantees by the server to the client that particular files may be accessed via an alternate data access protocol at specific locations. Such access is limited to particular sets of byte-ranges and may proceed until those byte-ranges are reduced or the layout is returned.

A stateid represents the set of all layouts held by a particular client for a particular filehandle with a given layout type. The seqid is updated as the layouts of that set of byte-ranges change, via layout stateid changing operations such as LAYOUTGET and LAYOUTRETURN.



### 13.2.2. Stateid Structure

Stateids are divided into two fields, a 96-bit "other" field identifying the specific set of locks and a 32-bit "seqid" sequence value. Except in the case of special stateids (See Section 13.2.3), a particular value of the "other" field denotes a set of locks of the same type (for example, byte-range locks, opens, delegations, or layouts), for a specific file or directory, and sharing the same ownership characteristics. The seqid designates a specific instance of such a set of locks, and is incremented to indicate changes in such a set of locks, either by the addition or deletion of locks from the set, a change in the byte-range they apply to, or an upgrade or downgrade in the type of one or more locks.

When such a set of locks is first created, the server returns a stateid with seqid value of one. On subsequent operations that modify the set of locks, the server is required to increment the "seqid" field by one whenever it returns a stateid for the same state-owner/file/type combination and there is some change in the set of locks actually designated. In this case, the server will return a stateid with an "other" field the same as previously used for that state-owner/file/type combination, with an incremented "seqid" field. This pattern continues until the seqid is incremented past NFS4\_UINT32\_MAX, and one (not zero) is the next seqid value.

The purpose of the incrementing of the seqid is to allow the server to communicate to the client the order in which operations that modified locking state associated with a stateid have been processed and to make it possible for the client to send requests that are conditional on the set of locks not having changed since the stateid in question was returned.

Except for layout stateids (Section 18.7.3), when a client sends a stateid to the server, it has two choices with regard to the seqid sent. It may set the seqid to zero to indicate to the server that it wishes the most up-to-date seqid for that stateid's "other" field to be used. This would be the common choice in the case of a stateid sent with a READ or WRITE operation. It also may set a non-zero value, in which case the server checks if that seqid is the correct one. In that case, the server is required to return NFS4ERR\_OLD\_STATEID if the seqid is lower than the most current value and NFS4ERR\_BAD\_STATEID if the seqid is greater than the most current value. This would be the common choice in the case of stateids sent with a CLOSE or OPEN\_DOWNGRADE. Because OPENS may be sent in parallel for the same owner, a client might close a file without knowing that an OPEN upgrade had been done by the server, changing the lock in question. If CLOSE were sent with a zero seqid, the OPEN upgrade would be cancelled before the client even received an indication that an upgrade had happened.

When a stateid is sent by the server to the client as part of a callback operation, it is not subject to checking for a current seqid and returning NFS4ERR\_OLD\_STATEID. This is because the client is not in a position to know the most up-to-date seqid and thus cannot verify it. Unless specially noted, the seqid value for a stateid sent by the server to the client as part of a callback is required to be zero with NFS4ERR\_BAD\_STATEID returned if it is not.

In making comparisons between seqids, both by the client in determining the order of operations and by the server in determining whether the NFS4ERR\_OLD\_STATEID is to be returned, the possibility of the seqid being swapped around past the NFS4\_UINT32\_MAX value needs to be taken into account. When two seqid values are being compared, the total count of slots for all sessions associated with the current client is used to do this. When one seqid value is less than this total slot count and another seqid value is greater than NFS4\_UINT32\_MAX minus the total slot count, the former is to be treated as lower than the latter, despite the fact that it is numerically greater.

### 13.2.3. Special Stateids

Stateid values whose "other" field is either all zeros or all ones are reserved. They may not be assigned by the server but have special meanings defined by the protocol. The particular meaning depends on whether the "other" field is all zeros or all ones and the specific value of the "seqid" field.

The following combinations of "other" and "seqid" are defined in NFSv4.1:

- \* When "other" and "seqid" are both zero, the stateid is treated as a special anonymous stateid, which can be used in READ, WRITE, and SETATTR requests to indicate the absence of any OPEN state associated with the request. When an anonymous stateid value is used and an existing open denies the form of access requested, then access will be denied to the request. This stateid MUST NOT be used on operations to data servers (Section 20.10).
- \* When "other" and "seqid" are both all ones, the stateid is a special READ bypass stateid. When this value is used in WRITE or SETATTR, it is treated like the anonymous value. When used in READ, the server MAY grant access, even if access would normally be denied to READ operations. This stateid MUST NOT be used on operations to data servers.
- \* When "other" is zero and "seqid" is one, the stateid represents the current stateid, which is whatever value is the last stateid returned by an operation within the COMPOUND. In the case of an OPEN, the stateid returned for the open file and not the delegation is used. The stateid passed to the operation in place of the special value has its "seqid" value set to zero, except when the current stateid is used by the operation CLOSE or OPEN\_DOWNGRADE. If there is no operation in the COMPOUND that has returned a stateid value, the server MUST return the error NFS4ERR\_BAD\_STATEID. As illustrated in Figure 6, if the value of a current stateid is a special stateid and the stateid of an operation's arguments has "other" set to zero and "seqid" set to one, then the server MUST return the error NFS4ERR\_BAD\_STATEID.
- \* When "other" is zero and "seqid" is NFS4\_UINT32\_MAX, the stateid represents a reserved stateid value defined to be invalid. When this stateid is used, the server MUST return the error NFS4ERR\_BAD\_STATEID.

If a stateid value is used that has all zeros or all ones in the "other" field but does not match one of the cases above, the server MUST return the error NFS4ERR\_BAD\_STATEID.

Special stateids, unlike other stateids, are not associated with individual client IDs or filehandles and can be used with all valid client IDs and filehandles. In the case of a special stateid designating the current stateid, the current stateid value substituted for the special stateid is associated with a particular client ID and filehandle, and so, if it is used where the current filehandle does not match that associated with the current stateid, the operation to which the stateid is passed will return NFS4ERR\_BAD\_STATEID.

#### 13.2.4. Stateid Lifetime and Validation

Stateids must remain valid until either a client restart or a server restart or until the client returns all of the locks associated with the stateid by means of an operation such as CLOSE or DELEGRETURN. If the locks are lost due to revocation, as long as the client ID is valid, the stateid remains a valid designation of that revoked state until the client frees it by using FREE\_STATEID. Stateids associated with byte-range locks are an exception. They remain valid even if a LOCKU frees all remaining locks, so long as the open file with which they are associated remains open, unless the client frees the stateids via the FREE\_STATEID operation.

It should be noted that there are situations in which the client's locks become invalid, without the client requesting they be returned. These include lease expiration and a number of forms of lock revocation within the lease period. It is important to note that in these situations, the stateid remains valid and the client can use it to determine the disposition of the associated lost locks.

An "other" value must never be reused for a different purpose (i.e., different filehandle, owner, or type of locks) within the context of a single client ID. A server may retain the "other" value for the same purpose beyond the point where it may otherwise be freed, but if it does so, it must maintain "seqid" continuity with previous values.

One mechanism that may be used to satisfy the requirement that the server recognize invalid and out-of-date stateids is for the server to divide the "other" field of the stateid into two fields.

- \* an index into a table of locking-state structures.
- \* a generation number that is incremented on each allocation of a table entry for a particular use.

And then store in each table entry,

- \* the client ID with which the stateid is associated.
- \* the current generation number for the (at most one) valid stateid sharing this index value.
- \* the filehandle of the file on which the locks are taken.
- \* an indication of the type of stateid (open, byte-range lock, file delegation, directory delegation, layout).

- \* the last "seqid" value returned corresponding to the current "other" value.
- \* an indication of the current status of the locks associated with this stateid, in particular, whether these have been revoked and if so, for what reason.

With this information, an incoming stateid can be validated and the appropriate error returned when necessary. Special and non-special stateids are handled separately. (See Section 13.2.3 for a discussion of special stateids.)

Note that stateids are implicitly qualified by the current client ID, as derived from the client ID associated with the current session. Note, however, that the semantics of the session will prevent stateids associated with a previous client or server instance from being analyzed by this procedure.

If server restart has resulted in an invalid client ID or a session ID that is invalid, SEQUENCE will return an error and the operation that takes a stateid as an argument will never be processed.

If there has been a server restart where there is a persistent session and all leased state has been lost, then the session in question will, although valid, be marked as dead, and any operation not satisfied by means of the reply cache will receive the error NFS4ERR\_DEADSESSION, and thus not be processed as indicated below.

When a stateid is being tested and the "other" field is all zeros or all ones, a check that the "other" and "seqid" fields match a defined combination for a special stateid is done and the results determined as follows:

- \* If the "other" and "seqid" fields do not match a defined combination associated with a special stateid, the error NFS4ERR\_BAD\_STATEID is returned.
- \* If the special stateid is one designating the current stateid and there is a current stateid, then the current stateid is substituted for the special stateid and the checks appropriate to non-special stateids are performed.
- \* If the combination is valid in general but is not appropriate to the context in which the stateid is used (e.g., an all-zero stateid is used when an OPEN stateid is required in a LOCK operation), the error NFS4ERR\_BAD\_STATEID is also returned.

- \* Otherwise, the check is completed and the special stateid is accepted as valid.

When a stateid is being tested, and the "other" field is neither all zeros nor all ones, the following procedure could be used to validate an incoming stateid and return an appropriate error, when necessary, assuming that the "other" field would be divided into a table index and an entry generation.

- \* If the table index field is outside the range of the associated table, return NFS4ERR\_BAD\_STATEID.
- \* If the selected table entry is of a different generation than that specified in the incoming stateid, return NFS4ERR\_BAD\_STATEID.
- \* If the selected table entry does not match the current filehandle, return NFS4ERR\_BAD\_STATEID.
- \* If the client ID in the table entry does not match the client ID associated with the current session, return NFS4ERR\_BAD\_STATEID.
- \* If the stateid represents revoked state, then return NFS4ERR\_EXPIRED, NFS4ERR\_ADMIN\_REVOKED, or NFS4ERR\_DELEG\_REVOKED, as appropriate.
- \* If the stateid type is not valid for the context in which the stateid appears, return NFS4ERR\_BAD\_STATEID. Note that a stateid may be valid in general, as would be reported by the TEST\_STATEID operation, but be invalid for a particular operation, as, for example, when a stateid that doesn't represent byte-range locks is passed to the non-from\_open case of LOCK or to LOCKU, or when a stateid that does not represent an open is passed to CLOSE or OPEN\_DOWNGRADE. In such cases, the server MUST return NFS4ERR\_BAD\_STATEID.
- \* If the "seqid" field is not zero and it is greater than the current sequence value corresponding to the current "other" field, return NFS4ERR\_BAD\_STATEID.
- \* If the "seqid" field is not zero and it is less than the current sequence value corresponding to the current "other" field, return NFS4ERR\_OLD\_STATEID.

- \* Otherwise, the stateid is valid and the table entry should contain any additional information about the type of stateid and information associated with that particular type of stateid, such as the associated set of locks, e.g., open-owner and lock-owner information, as well as information on the specific locks, e.g., open modes and byte-ranges.

#### 13.2.5. Stateid Use for I/O Operations

Clients performing I/O operations need to select an appropriate stateid based on the locks (including opens and delegations) held by the client and the various types of state-owners sending the I/O requests. SETATTR operations that change the file size are treated like I/O operations in this regard.

The following rules, applied in order of decreasing priority, govern the selection of the appropriate stateid. In following these rules, the client will only consider locks of which it has actually received notification by an appropriate operation response or callback. Note that the rules are slightly different in the case of I/O to data servers when file layouts are being used (See Section 20.14.1).

- \* If the client holds a delegation for the file in question, the delegation stateid should be used.
- \* Otherwise, if the entity corresponding to the lock-owner (e.g., a process) sending the I/O has a byte-range lock stateid for the associated open file, then the byte-range lock stateid for that lock-owner and open file should be used.
- \* If there is no byte-range lock stateid, then the OPEN stateid for the open file in question is to be used.
- \* Finally, if none of the above apply, then a special stateid can sometimes be used.

Ignoring these rules may result in situations in which the server does not have information necessary to properly process the request. For example, when mandatory byte-range locks are in effect, if the stateid does not indicate the proper lock-owner, via a lock stateid, a request might be avoidably rejected.

The server however should not try to enforce these ordering rules and should use whatever information is available to properly process I/O requests. In particular, when a client has a delegation for a given file, it should take note of this fact in processing a request, even if it is sent with a special stateid.

### 13.2.6. Stateid Use for SETATTR Operations

Because each operation is associated with a session ID and from that the clientid can be determined, operations do not need to include a stateid for the server to be able to determine whether they should cause a delegation to be recalled or are to be treated as done within the scope of the delegation.

In the case of SETATTR operations, a stateid is present. In cases other than those that set the file size, the client may send either a special stateid or, when a delegation is held for the file in question, a delegation stateid. While the server should validate the stateid and may use the stateid to optimize the determination as to whether a delegation is held, it should note the presence of a delegation even when a special stateid is sent, and MUST accept a valid delegation stateid when sent.

### 13.3. Lease Renewal

Each client/server pair, as represented by a client ID, has a single lease. The purpose of the lease is to allow the client to indicate to the server, in a low-overhead way, that it is active, and thus that the server is to retain the client's locks. This arrangement allows the server to remove stale locking-related objects that are held by a client that has crashed or is otherwise unreachable, once the relevant lease expires. This in turn allows other clients to obtain conflicting locks without being delayed indefinitely by inactive or unreachable clients. It is not a mechanism for cache consistency and lease renewals may not be denied if the lease interval has not expired.

Since each session is associated with a specific client (identified by the client's client ID), any operation sent on that session is an indication that the associated client is reachable. When a request is sent for a given session, successful execution of a SEQUENCE operation (or successful retrieval of the result of SEQUENCE from the reply cache) on an unexpired lease will result in the lease being implicitly renewed, for the standard renewal period (equal to the lease\_time attribute).

If the client ID's lease has not expired when the server receives a SEQUENCE operation, then the server MUST renew the lease. If the client ID's lease has expired when the server receives a SEQUENCE operation, the server MAY renew the lease; this depends on whether any state was revoked as a result of the client's failure to renew the lease before expiration.



Absent other activity that would renew the lease, a COMPOUND consisting of a single SEQUENCE operation will suffice. The client should also take communication-related delays into account and take steps to ensure that the renewal messages actually reach the server in good time. For example:

- \* When trunking is in effect, the client should consider sending multiple requests on different connections, in order to ensure that renewal occurs, even in the event of blockage in the path used for one of those connections.
- \* Transport retransmission delays might become so large as to approach or exceed the length of the lease period. This may be particularly likely when the server is unresponsive due to a restart; see Section 13.4.2.1. If the client implementation is not careful, transport retransmission delays can result in the client failing to detect a server restart before the grace period ends. The scenario is that the client is using a transport with exponential backoff, such that the maximum retransmission timeout exceeds both the grace period and the lease\_time attribute. A network partition causes the client's connection's retransmission interval to back off, and even after the partition heals, the next transport-level retransmission is sent after the server has restarted and its grace period ends.

The client MUST either recover from the ensuing NFS4ERR\_NO\_GRACE errors or it MUST ensure that, despite transport-level retransmission intervals that exceed the lease\_time, a SEQUENCE operation is sent that renews the lease before expiration. The client can achieve this by associating a new connection with the session, and sending a SEQUENCE operation on it. However, if the attempt to establish a new connection is delayed for some reason (e.g., exponential backoff of the connection establishment packets), the client will have to abort the connection establishment attempt before the lease expires, and attempt to reconnect.

If the server renews the lease upon receiving a SEQUENCE operation, the server MUST NOT allow the lease to expire while the rest of the operations in the COMPOUND procedure's request are still executing. Once the last operation has finished, and the response to COMPOUND has been sent, the server MUST set the lease to expire no sooner than the sum of current time and the value of the lease\_time attribute.

A client ID's lease can expire when it has been at least the lease interval (lease\_time) since the last lease-renewing SEQUENCE operation was sent on any of the client ID's sessions and there are no active COMPOUND operations on any such sessions.

Because the SEQUENCE operation is the basic mechanism to renew a lease, and because it must be done at least once for each lease period, it is the natural mechanism whereby the server will inform the client of changes in the lease status that the client needs to be informed of. The client should inspect the status flags (sr\_status\_flags) returned by sequence and take the appropriate action (See Section 25.46.3 for details).

- \* The status bits SEQ4\_STATUS\_CB\_PATH\_DOWN and SEQ4\_STATUS\_CB\_PATH\_DOWN\_SESSION indicate problems with the backchannel that the client may need to address in order to receive callback requests.
- \* The status bits SEQ4\_STATUS\_CB\_GSS\_CONTEXTS\_EXPIRING and SEQ4\_STATUS\_CB\_GSS\_CONTEXTS\_EXPIRED indicate problems with GSS contexts or RPCSEC\_GSS handles for the backchannel that the client might have to address in order to allow callback requests to be sent.
- \* The status bits SEQ4\_STATUS\_EXPIRED\_ALL\_STATE\_REVOKED, SEQ4\_STATUS\_EXPIRED\_SOME\_STATE\_REVOKED, SEQ4\_STATUS\_ADMIN\_STATE\_REVOKED, and SEQ4\_STATUS\_RECALLABLE\_STATE\_REVOKED notify the client of lock revocation events. When these bits are set, the client should use TEST\_STATEID to find what stateids have been revoked and use FREE\_STATEID to acknowledge loss of the associated state.
- \* The status bit SEQ4\_STATUS\_LEASE\_MOVE indicates that responsibility for lease renewal has been transferred to one or more new servers.
- \* The status bit SEQ4\_STATUS\_RESTART\_RECLAIM\_NEEDED indicates that due to server restart the client must reclaim locking state.
- \* The status bit SEQ4\_STATUS\_BACKCHANNEL\_FAULT indicates that the server has encountered an unrecoverable fault with the backchannel (e.g., it has lost track of a sequence ID for a slot in the backchannel).

#### 13.4. Crash Recovery

A critical requirement in crash recovery is that both the client and the server know when the other has failed. Additionally, it is required that a client sees a consistent view of data across server restarts. All READ and WRITE operations that may have been queued within the client or network buffers must wait until the client has successfully recovered the locks protecting the READ and WRITE operations. Any that reach the server before the server can safely

determine that the client has recovered enough locking state to be sure that such operations can be safely processed must be rejected. This will happen because either:

- \* The state presented is no longer valid since it is associated with a now invalid client ID. In this case, the client will receive either an NFS4ERR\_BADSESSION or NFS4ERR\_DEADSESSION error, and any attempt to attach a new session to that invalid client ID will result in an NFS4ERR\_STALE\_CLIENTID error.
- \* Subsequent recovery of locks may make execution of the operation inappropriate (NFS4ERR\_GRACE).

#### 13.4.1. Client Failure and Recovery

In the event that a client fails, the server may release the client's locks when the associated lease has expired. Conflicting locks from another client may only be granted after this lease expiration. As discussed in Section 13.3, when a client has not failed and re-establishes its lease before expiration occurs, requests for conflicting locks will not be granted.

To minimize client delay upon restart, lock requests are associated with an instance of the client by a client-supplied verifier. This verifier is part of the client\_owner4 sent in the initial EXCHANGE\_ID call made by the client. The server returns a client ID as a result of the EXCHANGE\_ID operation. The client then confirms the use of the client ID by establishing a session associated with that client ID (see Section 25.36.3 for a description of how this is done). All locks, including opens, byte-range locks, delegations, and layouts obtained by sessions using that client ID, are associated with that client ID.

Since the verifier will be changed by the client upon each initialization, the server can compare a new verifier to the verifier associated with currently held locks and determine that they do not match. This signifies the client's new instantiation and subsequent loss (upon confirmation of the new client ID) of locking state. As a result, the server is free to release all locks held that are associated with the old client ID that was derived from the old verifier. At this point, conflicting locks from other clients, kept waiting while the lease had not yet expired, can be granted. In addition, all stateids associated with the old client ID can also be freed, as they are no longer reference-able.

Note that the verifier must have the same uniqueness properties as the verifier for the COMMIT operation.

#### 13.4.2. Server Failure and Recovery

If the server loses locking state (usually as a result of a restart), it must allow clients time to discover this fact and re-establish the lost locking state. The client must be able to re-establish the locking state without having the server deny valid requests because the server has granted conflicting access to another client. Likewise, if there is a possibility that clients have not yet re-established their locking state for a file and that such locking state might make it invalid to perform READ or WRITE operations. For example, if mandatory locks are a possibility, the server must disallow READ and WRITE operations for that file.

A client can determine that loss of locking state has occurred via several methods.

1. When a SEQUENCE (most common) or other operation returns NFS4ERR\_BADSESSION, this may mean that the session has been destroyed but the client ID is still valid. The client sends a CREATE\_SESSION request with the client ID to re-establish the session. If CREATE\_SESSION fails with NFS4ERR\_STALE\_CLIENTID, the client must establish a new client ID (see Section 13.1) and re-establish its lock state with the new client ID, after the CREATE\_SESSION operation succeeds (See Section 13.4.2.1).
2. When a SEQUENCE (most common) or other operation on a persistent session returns NFS4ERR\_DEADSESSION, this indicates that a session is no longer usable for new, i.e., not satisfied from the reply cache, operations. Once all pending operations are determined to be either performed before the retry or not performed, the client sends a CREATE\_SESSION request with the client ID to re-establish the session. If CREATE\_SESSION fails with NFS4ERR\_STALE\_CLIENTID, the client must establish a new client ID (see Section 13.1) and re-establish its lock state after the CREATE\_SESSION, with the new client ID, succeeds (Section 13.4.2.1).
3. When an operation, neither SEQUENCE nor preceded by SEQUENCE (for example, CREATE\_SESSION, DESTROY\_SESSION), returns NFS4ERR\_STALE\_CLIENTID, the client MUST establish a new client ID (Section 13.1) and re-establish its lock state (Section 13.4.2.1).

#### 13.4.2.1. State Reclaim

When state information and the associated locks are lost as a result of a server restart, the protocol must provide a way to cause that state to be re-established. The approach used is to define, for most types of locking state (layouts are an exception), a request whose function is to allow the client to re-establish on the server a lock first obtained from a previous instance. Generally, these requests are variants of the requests normally used to create locks of that type and are referred to as "reclaim-type" requests, and the process of re-establishing such locks is referred to as "reclaiming" them.

Because each client must have an opportunity to reclaim all of the locks that it has without the possibility that some other client will be granted a conflicting lock, a "grace period" is devoted to the reclaim process. During this period, requests creating client IDs and sessions are handled normally, but locking requests are subject to special restrictions. Only reclaim-type locking requests are allowed, unless the server can reliably determine (through state persistently maintained across restart instances) that granting any such lock cannot possibly conflict with a subsequent reclaim. When a request is made to obtain a new lock (i.e., not a reclaim-type request) during the grace period and such a determination cannot be made, the server must return the error NFS4ERR\_GRACE.

Once a session is established using the new client ID, the client will use reclaim-type locking requests (e.g., LOCK operations with reclaim set to TRUE and OPEN operations with a claim type of CLAIM\_PREVIOUS; see Section 14.11) to re-establish its locking state. Once this is done, or if there is no such locking state to reclaim, the client sends a global RECLAIM\_COMPLETE operation, i.e., one with the rca\_one\_fs argument set to FALSE, to indicate that it has reclaimed all of the locking state that it will reclaim. Once a client sends such a RECLAIM\_COMPLETE operation, it may attempt non-reclaim locking operations, although it might get an NFS4ERR\_GRACE status result from each such operation until the period of special handling is over. See Section 17.11.9 for a discussion of the analogous handling lock reclamation in the case of file systems transitioning from server to server.

During the grace period, the server must reject any non-reclaim locking requests (i.e., other LOCK and OPEN operations) with an error of NFS4ERR\_GRACE, unless it can guarantee that these may be done safely, as described below. In addition, READ and WRITE requests that are not associated with a reclaimed OPEN need to be rejected as well.

The grace period may last until all clients that are known to have possibly had locks have done a global RECLAIM\_COMPLETE operation, indicating that they have finished reclaiming the locks they held before the server restart. This means that a client that has done a RECLAIM\_COMPLETE must be prepared to receive an NFS4ERR\_GRACE when attempting to acquire new locks. In order for the server to know that all clients with possible prior lock state have done a RECLAIM\_COMPLETE, the server must maintain in stable storage a list of clients that may have such locks. The server may also terminate the grace period before all clients have done a global RECLAIM\_COMPLETE. The server should not terminate the grace period without all expected RECLAIM\_COMPLETES before a time equal to the lease period in order to give clients an opportunity to find out about the server restart, as a result of sending requests on associated sessions with a frequency governed by the lease time. Note that when a client does not send such requests (or they are sent by the client but not received by the server), it is possible for the grace period to expire before the client finds out that the server restart has occurred.

Some additional time in order to allow a client to establish a new client ID and session and to effect lock reclaims may be added to the lease time. Note that analogous rules apply to file system-specific grace periods discussed in Section 17.11.9.

If the server can reliably determine that granting a non-reclaim request will not conflict with reclamation of locks by other clients, the NFS4ERR\_GRACE error does not have to be returned even within the grace period, although NFS4ERR\_GRACE must always be returned to clients attempting a non-reclaim lock request before doing their own global RECLAIM\_COMPLETE. For the server to be able to service READ and WRITE operations during the grace period, it must again be able to guarantee that no possible conflict could arise between a potential reclaim locking request and the READ or WRITE operation. If the server is unable to offer that guarantee, the NFS4ERR\_GRACE error must be returned to the client.

For a server to provide simple, valid handling during the grace period, the easiest method is to simply reject all non-reclaim locking requests and READ and WRITE operations not subsumed within reclaimed OPENS by returning the NFS4ERR\_GRACE error. However, a server may keep information about granted locks in stable storage. With this information, the server could determine if a locking operation, or a READ or WRITE outside a reclaimed OPEN can be safely processed.

For example, if the server maintained on stable storage summary information on whether mandatory locks exist, either mandatory byte-range locks, or share reservations specifying deny modes, many

requests could be allowed during the grace period. If it is known that no such share reservations exist, OPEN request that do not specify deny modes can be safely granted. If, in addition, it is known that no mandatory byte-range locks exist, either through information stored on stable storage or simply because the server does not support such locks, READ and WRITE operations may be safely processed during the grace period. Another important case is where it is known that no mandatory byte-range locks exist, either because the server does not provide support for them or because their absence is known from persistently recorded data. In this case, READ and WRITE operations specifying stateids derived from reclaim-type operations may be validly processed during the grace period because of the fact that the valid reclaim ensures that no lock subsequently granted can prevent the I/O.

To reiterate, for a server that allows non-reclaim lock and I/O requests to be processed during the grace period, it MUST determine that no lock subsequently reclaimed will be rejected and that no lock subsequently reclaimed would have prevented any I/O operation processed during the grace period.

Clients should be prepared for the return of NFS4ERR\_GRACE errors for non-reclaim lock and I/O requests. In this case, the client should employ a retry mechanism for the request. A delay (on the order of several seconds) between retries should be used to avoid overwhelming the server. Further discussion of the general issue is included in [Floyd]. The client must account for the server that can perform I/O and non-reclaim locking requests within the grace period as well as those that cannot do so.

A reclaim-type locking request outside the server's grace period can only succeed if the server can guarantee that no conflicting lock or I/O request has been granted since restart.

A server may, upon restart, establish a new value for the lease period. Therefore, clients should, once a new client ID is established, refetch the lease\_time attribute and use it as the basis for lease renewal for the lease associated with that server. However, the server must establish, for this restart event, a grace period at least as long as the lease period for the previous server instantiation. This allows the client state obtained during the previous server instance to be reliably re-established.

The possibility exists that, because of server configuration events, the client will be communicating with a server different than the one on which the locks were obtained, as shown by the combination of `eir_server_scope` and `eir_server_owner`. This leads to the issue of if and when the client should attempt to reclaim locks previously obtained on what is being reported as a different server. The rules to resolve this question are as follows:

- \* If the server scope is different, the client should not attempt to reclaim locks. In this situation, no lock reclaim is possible. Any attempt to re-obtain the locks with non-reclaim operations is problematic since there is no guarantee that the existing filehandles will be recognized by the new server, or that if recognized, they denote the same objects. It is best to treat the locks as having been revoked by the reconfiguration event.
- \* If the server scope is the same, the client should attempt to reclaim locks, even if the `eir_server_owner` value is different. In this situation, it is the responsibility of the server to return `NFS4ERR_NO_GRACE` if it cannot provide correct support for lock reclaim operations, including the prevention of edge conditions.

The `eir_server_owner` field is not used in making this determination. Its function is to specify trunking possibilities for the client (See Section 7.5) and not to control lock reclaim.

#### 13.4.2.1.1. Security Issues for State Reclaim

During the grace period, a client can reclaim state that it believes or asserts it had before the server restarted. Unless the server has maintained a complete record of all the state the client had, the server has little choice but to trust the client's requests. (Of course, if the server maintained a complete record, then there would be no need to force the client to reclaim state after server restart.) While the server has to trust the client to tell the truth, the negative consequences for security are limited to enabling denial-of-service attacks in situations in which `AUTH_SYS`, particularly `AUTH_SYS` in the clear, is supported. The fundamental rule for the server when processing reclaim requests is that it **MUST NOT** grant the reclaim if an equivalent non-reclaim request would not be granted during steady state due to access control or access conflict issues. For example, an `OPEN` request during a reclaim will be refused with `NFS4ERR_ACCESS` if the principal making the request does not have sufficient access to open the file according to the `acl`, `dacl`, or mode attributes of the file.



Nonetheless, it is possible that a client operating in error or maliciously could, during reclaim, prevent another client from reclaiming access to state. For example, an attacker could send an OPEN reclaim operation with a deny mode that prevents another client from reclaiming the OPEN state it had before the server restarted. The attacker could perform the same denial of service during steady state prior to server restart, as long as the attacker had permissions. Given that the attack vectors are equivalent, the grace period does not offer any additional opportunity for denial of service, and any concerns about this attack vector, whether during grace or steady state, are addressed in the same way, by using RPCSEC\_GSS for authentication and limiting access to the file only to principals that the owner of the file trusts.

Note that if prior to restart the server had client IDs with the EXCHGID4\_FLAG\_BIND\_PRINC\_STATEID (Section 25.35) capability set, then the server needs to record in stable storage the client owner id and the principal that established the client ID via EXCHANGE\_ID. If the server does not do so, then there is a risk a client will be unable to reclaim state if it does not have a credential for a principal that was originally authorized to establish the state.

#### 13.4.3. Network Partitions and Recovery

If the duration of a network partition is greater than the lease period provided by the server, the server will not have received a lease renewal from the client. If this occurs, the server may free all locks held for the client or it may allow the lock state to remain for a considerable period, subject to the constraint that if a request for a conflicting lock is made, locks associated with an expired lease do not prevent such a conflicting lock from being granted but MUST be revoked as necessary so as to avoid interfering with such conflicting requests.

If the server chooses to delay freeing of lock state until there is a conflict, it may either free all of the client's locks once there is a conflict or it may only revoke the minimum set of locks necessary to allow conflicting requests. When it adopts the finer-grained approach, it must revoke all locks associated with a given stateid, even if the conflict is with only a subset of locks.

When the server chooses to free all of a client's lock state, either immediately upon lease expiration or as a result of the first attempt to obtain a conflicting a lock, the server may report the loss of lock state in a number of ways.

The server may choose to invalidate the session and the associated client ID. In this case, once the client can communicate with the server, it will receive an NFS4ERR\_BADSESSION error. Upon attempting to create a new session, it would get an NFS4ERR\_STALE\_CLIENTID. Upon creating the new client ID and new session, the client will attempt to reclaim locks. Normally, the server will not allow the client to reclaim locks, because the server will not be in its recovery grace period.

Another possibility is for the server to maintain the session and client ID but for all stateids held by the client to become invalid or stale. Once the client can reach the server after such a network partition, the status returned by the SEQUENCE operation will indicate a loss of locking state; i.e., the flag SEQ4\_STATUS\_EXPIRED\_ALL\_STATE\_REVOKED will be set in sr\_status\_flags. In addition, all I/O submitted by the client with the now invalid stateids will fail with the server returning the error NFS4ERR\_EXPIRED. Once the client learns of the loss of locking state, it will suitably notify the applications that held the invalidated locks. The client should then take action to free invalidated stateids, either by establishing a new client ID using a new verifier or by doing a FREE\_STATEID operation to release each of the invalidated stateids.

When the server adopts a finer-grained approach to revocation of locks when a client's lease has expired, only a subset of stateids will normally become invalid during a network partition. When the client can communicate with the server after such a network partition heals, the status returned by the SEQUENCE operation will indicate a partial loss of locking state (SEQ4\_STATUS\_EXPIRED\_SOME\_STATE\_REVOKED). In addition, operations, including I/O submitted by the client, with the now invalid stateids will fail with the server returning the error NFS4ERR\_EXPIRED. Once the client learns of the loss of locking state, it will use the TEST\_STATEID operation on all of its stateids to determine which locks have been lost and then suitably notify the applications that held the invalidated locks. The client can then release the invalidated locking state and acknowledge the revocation of the associated locks by doing a FREE\_STATEID operation on each of the invalidated stateids.

When a network partition is combined with a server restart, there are edge conditions that place requirements on the server in order to avoid silent data corruption following the server restart. Two of these edge conditions are known, and are discussed below.

The first edge condition arises as a result of the scenarios such as the following:

1. Client A acquires a lock.
2. Client A and server experience mutual network partition, such that client A is unable to renew its lease.
3. Client A's lease expires, and the server releases the lock.
4. Client B acquires a lock that would have conflicted with that of client A.
5. Client B releases its lock.
6. Server restarts.
7. Network partition between client A and server heals.
8. Client A connects to a new server instance and finds out about server restart.
9. Client A reclaims its lock within the server's grace period.

Thus, at the final step, the server has erroneously granted client A's lock reclaim. If client B modified the object the lock was protecting, client A will experience object corruption.

The second known edge condition arises in situations such as the following:

1. Client A acquires one or more locks.
2. Server restarts.
3. Client A and server experience mutual network partition, such that client A is unable to reclaim all of its locks within the grace period.
4. Server's reclaim grace period ends. Client A has either no locks or an incomplete set of locks known to the server.
5. Client B acquires a lock that would have conflicted with a lock of client A that was not reclaimed.
6. Client B releases the lock.
7. Server restarts a second time.
8. Network partition between client A and server heals.

9. Client A connects to new server instance and finds out about server restart.
10. Client A reclaims its lock within the server's grace period.

As with the first edge condition, the final step of the scenario of the second edge condition has the server erroneously granting client A's lock reclaim.

Solving the first and second edge conditions requires either that the server always assumes after it restarts that some edge condition occurs, and thus returns NFS4ERR\_NO\_GRACE for all reclaim attempts, or that the server record some information in stable storage. The amount of information the server records in stable storage is in inverse proportion to how harsh the server intends to be whenever edge conditions arise. The server that is completely tolerant of all edge conditions will record in stable storage every lock that is acquired, removing the lock record from stable storage only when the lock is released. For the two edge conditions discussed above, the harshest a server can be, and still support a grace period for reclaims, requires that the server record in stable storage some minimal information. For example, a server implementation could, for each client, save in stable storage a record containing:

- \* the co\_ownerid field from the client\_owner4 presented in the EXCHANGE\_ID operation.
- \* a boolean that indicates if the client's lease expired or if there was administrative intervention (see Section 13.5) to revoke a byte-range lock, share reservation, or delegation and there has been no acknowledgment, via FREE\_STATEID, of such revocation.
- \* a boolean that indicates whether the client may have locks that it believes to be reclaimable in situations in which the grace period was terminated, making the server's view of lock reclaimability suspect. The server will set this for any client record in stable storage where the client has not done a suitable RECLAIM\_COMPLETE (global or file system-specific depending on the target of the lock request) before it grants any new (i.e., not reclaimed) lock to any client.

Assuming the above record keeping, for the first edge condition, after the server restarts, the record that client A's lease expired means that another client could have acquired a conflicting byte-range lock, share reservation, or delegation. Hence, the server must reject a reclaim from client A with the error NFS4ERR\_NO\_GRACE.

For the second edge condition, after the server restarts for a second time, the indication that the client had not completed its reclaims at the time at which the grace period ended means that the server must reject a reclaim from client A with the error NFS4ERR\_NO\_GRACE.

When either edge condition occurs, the client's attempt to reclaim locks will result in the error NFS4ERR\_NO\_GRACE. When this is received, or after the client restarts with no lock state, the client will send a global RECLAIM\_COMPLETE. When the RECLAIM\_COMPLETE is received, the server and client are again in agreement regarding reclaimable locks and both booleans in persistent storage can be reset, to be set again only when there is a subsequent event that causes lock reclaim operations to be questionable.

Regardless of the level and approach to record keeping, the server MUST implement one of the following strategies (which apply to reclaims of share reservations, byte-range locks, and delegations):

1. Reject all reclaims with NFS4ERR\_NO\_GRACE. This is extremely unforgiving, but necessary if the server does not record lock state in stable storage.
2. Record sufficient state in stable storage such that all known edge conditions involving server restart, including the two noted in this section, are detected. It is acceptable to erroneously recognize an edge condition and not allow a reclaim, when, with sufficient knowledge, it would be allowed. The error the server would return in this case is NFS4ERR\_NO\_GRACE. Note that it is not known if there are other edge conditions.

In the event that, after a server restart, the server determines there is unrecoverable damage or corruption to the information in stable storage, then for all clients and/or locks that may be affected, the server MUST return NFS4ERR\_NO\_GRACE.

A mandate for the client's handling of the NFS4ERR\_NO\_GRACE error is outside the scope of this specification, since the strategies for such handling are very dependent on the client's operating environment. However, one potential approach is described below.

When the client receives NFS4ERR\_NO\_GRACE, it could examine the change attribute of the objects for which the client is trying to reclaim state, and use that to determine whether to re-establish the state via normal OPEN or LOCK operations. This is acceptable provided that the client's operating environment allows it. In other words, the client implementer is advised to document for his users the behavior. The client could also inform the application that its byte-range lock or share reservations (whether or not they were

delegated) have been lost, such as via a UNIX signal, a Graphical User Interface (GUI) pop-up window, etc. See Section 15.5 for a discussion of what the client should do for dealing with unreclaimed delegations on client state.

For further discussion of revocation of locks, see Section 13.5.

### 13.5. Server Revocation of Locks

At any point, the server can revoke locks held by a client, and the client must be prepared for this event. When the client detects that its locks have been or may have been revoked, the client is responsible for validating the state information between itself and the server. Validating locking state for the client means that it must verify or reclaim state for each lock currently held.

The first occasion of lock revocation is upon server restart. Note that this includes situations in which sessions are persistent and locking state is lost. In this class of instances, the client will receive an error (NFS4ERR\_STALE\_CLIENTID) on an operation that takes client ID, usually as part of recovery in response to a problem with the current session), and the client will proceed with normal crash recovery as described in the Section 13.4.2.1.

The second occasion of lock revocation is the inability to renew the lease before expiration, as discussed in Section 13.4.3. While this is considered a rare or unusual event, the client must be prepared to recover. The server is responsible for determining the precise consequences of the lease expiration, informing the client of the scope of the lock revocation decided upon. The client then uses the status information provided by the server in the SEQUENCE results (field `sr_status_flags`, see Section 25.46.3) to synchronize its locking state with that of the server, in order to recover.

The third occasion of lock revocation can occur as a result of revocation of locks within the lease period, either because of administrative intervention or because a recallable lock (a delegation or layout) was not returned within the lease period after having been recalled. While these are considered rare events, they are possible, and the client must be prepared to deal with them. When either of these events occurs, the client finds out about the situation through the status returned by the SEQUENCE operation. Any use of stateids associated with locks revoked during the lease period will receive the error NFS4ERR\_ADMIN\_REVOKED or NFS4ERR\_DELEG\_REVOKED, as appropriate.

In all situations in which a subset of locking state may have been revoked, which include all cases in which locking state is revoked within the lease period, it is up to the client to determine which locks have been revoked and which have not. It does this by using the TEST\_STATEID operation on the appropriate set of stateids. Once the set of revoked locks has been determined, the applications can be notified, and the invalidated stateids can be freed and lock revocation acknowledged by using FREE\_STATEID.

### 13.6. Short and Long Leases

When determining the time period for the server lease, the usual lease trade-offs apply. A short lease is good for fast server recovery at a cost of increased operations to effect lease renewal (when there are no other operations during the period to effect lease renewal as a side effect). A long lease is certainly kinder and gentler to servers trying to handle very large numbers of clients. The number of extra requests to effect lock renewal drops in inverse proportion to the lease time. The disadvantages of a long lease include the possibility of slower recovery after certain failures. After server failure, a longer grace period may be required when some clients do not promptly reclaim their locks and do a global RECLAIM\_COMPLETE. In the event of client failure, the longer period for a lease to expire will force conflicting requests to wait longer.

A long lease is practical if the server can store lease state in stable storage. Upon recovery, the server can reconstruct the lease state from its stable storage and continue operation with its clients.

### 13.7. Clocks, Propagation Delay, and Calculating Lease Expiration

To avoid the need for synchronized clocks, lease times are granted by the server as a time delta. However, there is a requirement that the client and server clocks do not drift excessively over the duration of the lease. There is also the issue of propagation delay across the network, which could easily be several hundred milliseconds, as well as the possibility that requests will be lost and need to be retransmitted.

To take propagation delay into account, the client should subtract it from lease times (e.g., if the client estimates the one-way propagation delay as 200 milliseconds, then it can assume that the lease is already 200 milliseconds old when it gets it). In addition, it will take another 200 milliseconds to get a response back to the server. So the client must send a lease renewal or write data back to the server at least 400 milliseconds before the lease would expire. If the propagation delay varies over the life of the lease

(e.g., the client is on a mobile host), the client will need to continuously subtract the increase in propagation delay from the lease times.

The server's lease period configuration should take into account the network distance of the clients that will be accessing the server's resources. It is expected that the lease period will take into account the network propagation delays and other network delay factors for the client population. Since the protocol does not allow for an automatic method to determine an appropriate lease period, the server's administrator may have to tune the lease period.

### 13.8. Obsolete Locking Infrastructure from NFSv4.0

There are a number of operations and fields within existing operations that no longer have a function in NFSv4.1. In one way or another, these changes are all due to the implementation of sessions that provide client context and exactly once semantics as a base feature of the protocol, separate from locking itself.

The following NFSv4.0 operations MUST NOT be implemented in NFSv4.1. The server MUST return NFS4ERR\_NOTSUPP if these operations are found in an NFSv4.1 COMPOUND.

- \* SETCLIENTID since its function has been replaced by EXCHANGE\_ID.
- \* SETCLIENTID\_CONFIRM since client ID confirmation now happens by means of CREATE\_SESSION.
- \* OPEN\_CONFIRM because state-owner-based seqids have been replaced by the sequence ID in the SEQUENCE operation.
- \* RELEASE\_LOCKOWNER because lock-owners with no associated locks do not have any sequence-related state and so can be deleted by the server at will.
- \* RENEW because every SEQUENCE operation for a session causes lease renewal, making a separate operation superfluous.

Also, there are a number of fields, present in existing operations, related to locking that have no use in minor version 1. They were used in minor version 0 to perform functions now provided in a different fashion.

- \* Sequence ids used to sequence requests for a given state-owner and to provide retry protection, now provided via sessions.



- \* Client IDs used to identify the client associated with a given request. Client identification is now available using the client ID associated with the current session, without needing an explicit client ID field.

Such vestigial fields in existing operations have no function in NFSv4.1 and are ignored by the server. Note that client IDs in operations new to NFSv4.1 (such as CREATE\_SESSION and DESTROY\_CLIENTID) are not ignored.

## 14. File Locking and Share Reservations

To support Win32 share reservations, it is necessary to provide operations that atomically open or create files. Having a separate share/unshare operation would not allow correct implementation of the Win32 OpenFile API. In order to correctly implement share semantics, the previous NFS protocol mechanisms used when a file is opened or created (LOOKUP, CREATE, ACCESS) need to be replaced. The NFSv4.1 protocol defines an OPEN operation that is capable of atomically looking up, creating, and locking a file on the server.

### 14.1. Opens and Byte-Range Locks

It is assumed that manipulating a byte-range lock is rare when compared to READ and WRITE operations. It is also assumed that server restarts and network partitions are relatively rare. Therefore, it is important that the READ and WRITE operations have a lightweight mechanism to indicate if they possess a held lock. A LOCK operation contains the heavyweight information required to establish a byte-range lock and uniquely define the owner of the lock.

#### 14.1.1. State-Owner Definition

When opening a file or requesting a byte-range lock, the client must specify an identifier that represents the owner of the requested lock. This identifier is in the form of a state-owner, represented in the protocol by a `state_owner4`, a variable-length opaque array that, when concatenated with the current client ID, uniquely defines the owner of a lock managed by the client. This may be a thread ID, process ID, or other unique value.

Owners of opens and owners of byte-range locks are separate entities and remain separate even if the same opaque arrays are used to designate owners of each. The protocol distinguishes between open-owners (represented by `open_owner4` structures) and lock-owners (represented by `lock_owner4` structures).

Each open is associated with a specific open-owner while each byte-range lock is associated with a lock-owner and an open-owner, the latter being the open-owner associated with the open file under which the LOCK operation was done. Delegations and layouts, on the other hand, are not associated with a specific owner but are associated with the client as a whole (identified by a client ID).

#### 14.1.2. Use of the Stateid and Locking

All READ, WRITE, and SETATTR operations contain a stateid. For the purposes of this section, SETATTR operations that change the size attribute of a file are treated as if they are writing the area between the old and new sizes (i.e., the byte-range truncated or added to the file by means of the SETATTR), even where SETATTR is not explicitly mentioned in the text. The stateid passed to one of these operations must be one that represents an open, a set of byte-range locks, or a delegation, or it may be a special stateid representing anonymous access or the special bypass stateid.

If the state-owner performs a READ or WRITE operation in a situation in which it has established a byte-range lock or share reservation on the server (any OPEN constitutes a share reservation), the stateid (previously returned by the server) must be used to indicate what locks, including both byte-range locks and share reservations, are held by the state-owner. If no state is established by the client, either a byte-range lock or a share reservation, a special stateid for anonymous state (zero as the value for "other" and "seqid") is used. (See Section 13.2.3 for a description of 'special' stateids in general.) Regardless of whether a stateid for anonymous state or a stateid returned by the server is used, if there is a conflicting share reservation or mandatory byte-range lock held on the file, the server MUST refuse to service the READ or WRITE operation.

Share reservations are established by OPEN operations and by their nature are mandatory in that when the OPEN denies READ or WRITE operations, that denial results in such operations being rejected with error NFS4ERR\_LOCKED. Byte-range locks may be implemented by the server as either mandatory or advisory, or the choice of mandatory or advisory behavior may be determined by the server on the basis of the file being accessed (for example, some UNIX-based servers support a "mandatory lock bit" on the mode attribute such that if set, byte-range locks are required on the file before I/O is possible). When byte-range locks are advisory, they only prevent the granting of conflicting lock requests and have no effect on READs or WRITEs. Mandatory byte-range locks, however, prevent conflicting I/O operations. When they are attempted, they are rejected with NFS4ERR\_LOCKED. When the client gets NFS4ERR\_LOCKED on a file for which it knows it has the proper share reservation, it will need to

send a LOCK operation on the byte-range of the file that includes the byte-range the I/O was to be performed on, with an appropriate locktype field of the LOCK operation's arguments (i.e., READ\*\_LT for a READ operation, WRITE\*\_LT for a WRITE operation).

Note that for UNIX environments that support mandatory byte-range locking, the distinction between advisory and mandatory locking is subtle. In fact, advisory and mandatory byte-range locks are exactly the same as far as the APIs and requirements on implementation. If the mandatory lock attribute is set on the file, the server checks to see if the lock-owner has an appropriate shared (READ\_LT) or exclusive (WRITE\_LT) byte-range lock on the byte-range it wishes to READ from or WRITE to. If there is no appropriate lock, the server checks if there is a conflicting lock (which can be done by attempting to acquire the conflicting lock on behalf of the lock-owner, and if successful, release the lock after the READ or WRITE operation is done), and if there is, the server returns NFS4ERR\_LOCKED.

For Windows environments, byte-range locks are always mandatory, so the server always checks for byte-range locks during I/O requests.

Thus, the LOCK operation does not need to distinguish between advisory and mandatory byte-range locks. It is the server's processing of the READ and WRITE operations that introduces the distinction.

Every stateid that is validly passed to READ, WRITE, or SETATTR, with the exception of special stateid values, defines an access mode for the file (i.e., OPEN4\_SHARE\_ACCESS\_READ, OPEN4\_SHARE\_ACCESS\_WRITE, or OPEN4\_SHARE\_ACCESS\_BOTH).

- \* For stateids associated with opens, this is the mode defined by the original OPEN that caused the allocation of the OPEN stateid and as modified by subsequent OPENS and OPEN\_DOWNGRADES for the same open-owner/file pair.
- \* For stateids returned by byte-range LOCK operations, the appropriate mode is the access mode for the OPEN stateid associated with the lock set represented by the stateid.
- \* For delegation stateids, the access mode is based on the type of delegation.

When a READ, WRITE, or SETATTR (that specifies the size attribute) operation is done, the operation is subject to checking against the access mode to verify that the operation is appropriate given the stateid with which the operation is associated.

In the case of WRITE-type operations (i.e., WRITES and SETATTRs that set size), the server MUST verify that the access mode allows writing and MUST return an NFS4ERR\_OPENMODE error if it does not. In the case of READ, the server may perform the corresponding check on the access mode, or it may choose to allow READ on OPENS for OPEN4\_SHARE\_ACCESS\_WRITE, to accommodate clients whose WRITE implementation may unavoidably do reads (e.g., due to buffer cache constraints). However, even if READs are allowed in these circumstances, the server MUST still check for locks that conflict with the READ (e.g., another OPEN specified OPEN4\_SHARE\_DENY\_READ or OPEN4\_SHARE\_DENY\_BOTH). Note that a server that does enforce the access mode check on READs need not explicitly check for conflicting share reservations since the existence of OPEN for OPEN4\_SHARE\_ACCESS\_READ guarantees that no conflicting share reservation can exist.

The READ bypass special stateid (all bits of "other" and "seqid" set to one) indicates a desire to bypass locking checks. The server MAY allow READ operations to bypass locking checks at the server, when this special stateid is used. However, WRITE operations with this special stateid value MUST NOT bypass locking checks and are treated exactly the same as if a special stateid for anonymous state were used.

A lock may not be granted while a READ or WRITE operation using one of the special stateids is being performed and the scope of the lock to be granted would conflict with the READ or WRITE operation. This can occur when:

- \* A mandatory byte-range lock is requested with a byte-range that conflicts with the byte-range of the READ or WRITE operation. For the purposes of this paragraph, a conflict occurs when a shared lock is requested and a WRITE operation is being performed, or an exclusive lock is requested and either a READ or a WRITE operation is being performed.
- \* A share reservation is requested that denies reading and/or writing and the corresponding operation is being performed.
- \* A delegation is to be granted and the delegation type would prevent the I/O operation, i.e., READ and WRITE conflict with an OPEN\_DELEGATE\_WRITE delegation and WRITE conflicts with an OPEN\_DELEGATE\_READ delegation.

When a client holds a delegation, it needs to ensure that the stateid sent conveys the association of operation with the delegation, to avoid the delegation from being avoidably recalled. When the delegation stateid, a stateid open associated with that delegation,

or a stateid representing byte-range locks derived from such an open is used, the server knows that the READ, WRITE, or SETATTR does not conflict with the delegation but is sent under the aegis of the delegation. Even though it is possible for the server to determine from the client ID (via the session ID) that the client does in fact have a delegation, the server is not obliged to check this, so using a special stateid can result in avoidable recall of the delegation.

#### 14.2. Lock Ranges

The protocol allows a lock-owner to request a lock with a byte-range and then either upgrade, downgrade, or unlock a sub-range of the initial lock, or a byte-range that overlaps -- fully or partially -- either with that initial lock or a combination of a set of existing locks for the same lock-owner. It is expected that this will be an uncommon type of request. In any case, servers or server file systems may not be able to support sub-range lock semantics. In the event that a server receives a locking request that represents a sub-range of current locking state for the lock-owner, the server is allowed to return the error NFS4ERR\_LOCK\_RANGE to signify that it does not support sub-range lock operations. Therefore, the client should be prepared to receive this error and, if appropriate, report the error to the requesting application.

The client is discouraged from combining multiple independent locking ranges that happen to be adjacent into a single request since the server may not support sub-range requests for reasons related to the recovery of byte-range locking state in the event of server failure. As discussed in Section 13.4.2, the server may employ certain optimizations during recovery that work effectively only when the client's behavior during lock recovery is similar to the client's locking behavior prior to server failure.

#### 14.3. Upgrading and Downgrading Locks

If a client has a WRITE\_LT lock on a byte-range, it can request an atomic downgrade of the lock to a READ\_LT lock via the LOCK operation, by setting the type to READ\_LT. If the server supports atomic downgrade, the request will succeed. If not, it will return NFS4ERR\_LOCK\_NOTSUPP. The client should be prepared to receive this error and, if appropriate, report the error to the requesting application.

If a client has a READ\_LT lock on a byte-range, it can request an atomic upgrade of the lock to a WRITE\_LT lock via the LOCK operation by setting the type to WRITE\_LT or WRITEW\_LT. If the server does not support atomic upgrade, it will return NFS4ERR\_LOCK\_NOTSUPP. If the upgrade can be achieved without an existing conflict, the request

will succeed. Otherwise, the server will return either NFS4ERR\_DENIED or NFS4ERR\_DEADLOCK. The error NFS4ERR\_DEADLOCK is returned if the client sent the LOCK operation with the type set to WRITE\_LT and the server has detected a deadlock. The client should be prepared to receive such errors and, if appropriate, report the error to the requesting application.

#### 14.4. Stateid Seqid Values and Byte-Range Locks

When a LOCK or LOCKU operation is performed, the stateid returned has the same "other" value as the argument's stateid, and a "seqid" value that is incremented (relative to the argument's stateid) to reflect the occurrence of the LOCK or LOCKU operation. The server MUST increment the value of the "seqid" field whenever there is any change to the locking status of any byte offset as described by any of the locks covered by the stateid. A change in locking status includes a change from locked to unlocked or the reverse or a change from being locked for READ\_LT to being locked for WRITE\_LT or the reverse.

When there is no such change, as, for example, when a range already locked for WRITE\_LT is locked again for WRITE\_LT, the server MAY increment the "seqid" value.

#### 14.5. Issues with Multiple Open-Owners

When the same file is opened by multiple open-owners, a client will have multiple OPEN stateids for that file, each associated with a different open-owner. In that case, there can be multiple LOCK and LOCKU requests for the same lock-owner sent using the different OPEN stateids, and so a situation may arise in which there are multiple stateids, each representing byte-range locks on the same file and held by the same lock-owner but each associated with a different open-owner.

In such a situation, the locking status of each byte (i.e., whether it is locked, the READ\_LT or WRITE\_LT type of the lock, and the lock-owner holding the lock) MUST reflect the last LOCK or LOCKU operation done for the lock-owner in question, independent of the stateid through which the request was sent.

When a byte is locked by the lock-owner in question, the open-owner to which that byte-range lock is assigned should be that of the open-owner associated with the stateid through which the last LOCK of that byte was done. When there is a change in the open-owner associated with locks for the stateid through which a LOCK or LOCKU was done, the "seqid" field of the stateid MUST be incremented, even if the locking, in terms of lock-owners has not changed. When there is a change to the set of locked bytes associated with a different stateid for the same lock-owner, i.e., associated with a different open-owner, the "seqid" value for that stateid MUST NOT be incremented.

#### 14.6. Blocking Locks

Some clients require the support of blocking locks. While NFSv4.1 provides a callback when a previously unavailable lock becomes available, this is an OPTIONAL feature and clients cannot depend on its presence. Clients need to be prepared to continually poll for the lock. This presents a fairness problem. Two of the lock types, READW\_LT and WRITEW\_LT, are used to indicate to the server that the client is requesting a blocking lock. When the callback is not used, the server should maintain an ordered list of pending blocking locks. When the conflicting lock is released, the server may wait for the period of time equal to lease\_time for the first waiting client to re-request the lock. After the lease period expires, the next waiting client request is allowed the lock. Clients are required to poll at an interval sufficiently small that it is likely to acquire the lock in a timely manner. The server is not required to maintain a list of pending blocked locks as it is used to increase fairness and not correct operation. Because of the unordered nature of crash recovery, storing of lock state to stable storage would be required to guarantee ordered granting of blocking locks.

Servers may also note the lock types and delay returning denial of the request to allow extra time for a conflicting lock to be released, allowing a successful return. In this way, clients can avoid the burden of needless frequent polling for blocking locks. The server should take care in the length of delay in the event the client retransmits the request.

If a server receives a blocking LOCK operation, denies it, and then later receives a nonblocking request for the same lock, which is also denied, then it should remove the lock in question from its list of pending blocking locks. Clients should use such a nonblocking request to indicate to the server that this is the last time they intend to poll for the lock, as may happen when the process requesting the lock is interrupted. This is a courtesy to the server, to prevent it from unnecessarily waiting a lease period before granting other LOCK operations. However, clients are not

required to perform this courtesy, and servers must not depend on them doing so. Also, clients must be prepared for the possibility that this final locking request will be accepted.

When a server indicates, via the flag `OPEN4_RESULT_MAY_NOTIFY_LOCK`, that `CB_NOTIFY_LOCK` callbacks might be done for the current open file, the client should take notice of this, but, since this is a hint, cannot rely on a `CB_NOTIFY_LOCK` always being done. A client may reasonably reduce the frequency with which it polls for a denied lock, since the greater latency that might occur is likely to be eliminated given a prompt callback, but it still needs to poll. When it receives a `CB_NOTIFY_LOCK`, it should promptly try to obtain the lock, but it should be aware that other clients may be polling and that the server is under no obligation to reserve the lock for that particular client.

#### 14.7. Share Reservations

A share reservation is a mechanism to control access to a file. It is a separate and independent mechanism from byte-range locking. When a client opens a file, it sends an `OPEN` operation to the server specifying the type of access required (`READ`, `WRITE`, or `BOTH`) and the type of access to deny others (`OPEN4_SHARE_DENY_NONE`, `OPEN4_SHARE_DENY_READ`, `OPEN4_SHARE_DENY_WRITE`, or `OPEN4_SHARE_DENY_BOTH`). If the `OPEN` fails, the client will fail the application's open request.

Pseudo-code definition of the semantics:

```
if (request.access == 0) {
    return (NFS4ERR_INVALID)
} else {
    if ((request.access & file_state.deny) ||
        (request.deny & file_state.access)) {
        return (NFS4ERR_SHARE_DENIED)
    }
    return (NFS4ERR_OK);
}
```

When doing this checking of share reservations on `OPEN`, the current `file_state` used in the algorithm includes bits that reflect all current opens, including those for the open-owner making the new `OPEN` request.

The constants used for the `OPEN` and `OPEN_DOWNGRADE` operations for the access and deny fields are as follows:



```
const OPEN4_SHARE_ACCESS_READ    = 0x00000001;
const OPEN4_SHARE_ACCESS_WRITE   = 0x00000002;
const OPEN4_SHARE_ACCESS_BOTH    = 0x00000003;

const OPEN4_SHARE_DENY_NONE      = 0x00000000;
const OPEN4_SHARE_DENY_READ      = 0x00000001;
const OPEN4_SHARE_DENY_WRITE     = 0x00000002;
const OPEN4_SHARE_DENY_BOTH      = 0x00000003;
```

#### 14.8. OPEN/CLOSE Operations

To provide correct share semantics, a client MUST use the OPEN operation to obtain the initial filehandle and indicate the desired access and what access, if any, to deny. Even if the client intends to use a special stateid for anonymous state or READ bypass, it must still obtain the filehandle for the regular file with the OPEN operation so the appropriate share semantics can be applied. Clients that do not have a deny mode built into their programming interfaces for opening a file should request a deny mode of OPEN4\_SHARE\_DENY\_NONE.

The OPEN operation with the CREATE flag also subsumes the CREATE operation for regular files as used in previous versions of the NFS protocol. This allows a create with a share to be done atomically.

The CLOSE operation removes all share reservations held by the open-owner on that file. If byte-range locks are held, the client should release all locks before sending a CLOSE operation. The server MAY free all outstanding locks on CLOSE, but some servers might not support the CLOSE of a file that still has byte-range locks held. The server MUST return failure, NFS4ERR\_LOCKS\_HELD, if any locks would exist after the CLOSE.

The LOOKUP operation will return a filehandle without establishing any lock state on the server. Without a valid stateid, the server will assume that the client has the least access. For example, if one client opened a file with OPEN4\_SHARE\_DENY\_BOTH and another client accesses the file via a filehandle obtained through LOOKUP, the second client could only read the file using the special read bypass stateid. The second client could not WRITE the file at all because it would not have a valid stateid from OPEN and the special anonymous stateid would not be allowed access.

#### 14.9. Open Upgrade and Downgrade

When an OPEN is done for a file and the open-owner for which the OPEN is being done already has the file open, the result is to upgrade the open file status maintained on the server to include the access and deny bits specified by the new OPEN as well as those for the existing OPEN. The result is that there is one open file, as far as the protocol is concerned, and it includes the union of the access and deny bits for all of the OPEN requests completed. The OPEN is represented by a single stateid whose "other" value matches that of the original open, and whose "seqid" value is incremented to reflect the occurrence of the upgrade. The increment is required in cases in which the "upgrade" results in no change to the open mode (e.g., an OPEN is done for read when the existing open file is opened for OPEN4\_SHARE\_ACCESS\_BOTH). Only a single CLOSE will be done to reset the effects of both OPENs. The client may use the stateid returned by the OPEN effecting the upgrade or with a stateid sharing the same "other" field and a seqid of zero, although care needs to be taken as far as upgrades that happen while the CLOSE is pending. Note that the client, when sending the OPEN, may not know that the same file is in fact being opened. For reasons clarified later, the above only applies if both OPENs result in the OPENED object being designated by the same filehandle.

There are a number of situations in which the open-owner and file are the same but the upgrade, as described above, is inappropriate:

- \* If the principal doing the later OPEN is different from the one doing the preceding OPEN, the upgrade SHOULD NOT be done. The only valid reason known to allow this recommendation to be bypassed is that previous specifications including [RFC8881] did not deal with the issue at all, resulting in implementors doing the upgrade described above despite the fact the authorization of IO operation done by these two OPENs becomes confused since authorization checking for IO done within OPEN needs to allow the operations authorized by the OPEN.
- \* When the server chooses to export multiple filehandles corresponding to the same file object and returns different filehandles on two different OPENs of the same file object, the server MUST NOT "OR" together the access and deny bits and coalesce the two open files. Instead, the server needs to maintain separate OPENs with separate stateids and will require separate CLOSEs to free them.

When multiple open files on the client are merged into a single OPEN file object on the server, the close of one of the open files (on the client) may necessitate change of the access and deny

status of the open file on the server. This is because the union of the access and deny bits for the remaining opens may be smaller (i.e., a proper subset) than previously. The OPEN\_DOWNGRADE operation is used to make the necessary change and the client should use it to update the server so that share reservation requests by other clients are handled properly. The stateid returned has the same "other" field as that passed to the server. The "seqid" value in the returned stateid MUST be incremented, even in situations in which there is no change to the access and deny bits for the file.

#### 14.10. Parallel OPENS

Unlike the case of NFSv4.0, in which OPEN operations for the same open-owner are inherently serialized because of the owner-based seqid, multiple OPENS for the same open-owner may be done in parallel. When clients do this, they may encounter situations in which, because of the existence of hard links, two OPEN operations may turn out to open the same file, with a later OPEN performed being an upgrade of the first, with this fact only visible to the client once the operations complete.

In this situation, clients may determine the order in which the OPENS were performed by examining the stateids returned by the OPENS. Stateids that share a common value of the "other" field can be recognized as having opened the same file, with the order of the operations determinable from the order of the "seqid" fields, mod any possible wraparound of the 32-bit field.

When the possibility exists that the client will send multiple OPENS for the same open-owner in parallel, it may be the case that an open upgrade may happen without the client knowing beforehand that this could happen. Because of this possibility, CLOSEs and OPEN\_DOWNGRADES should generally be sent with a non-zero seqid in the stateid, to avoid the possibility that the status change associated with an open upgrade is not inadvertently lost.

#### 14.11. Reclaim of Open and Byte-Range Locks

Special forms of the LOCK and OPEN operations are provided when it is necessary to re-establish byte-range locks or opens after a server failure.

- \* To reclaim existing opens, an OPEN operation is performed using a CLAIM\_PREVIOUS. Because the client, in this type of situation, will have already opened the file and have the filehandle of the target file, this operation requires that the current filehandle be the target file, rather than a directory, and no file name is specified.
- \* To reclaim byte-range locks, a LOCK operation with the reclaim parameter set to true is used.

Reclaims of opens associated with delegations are discussed in Section 15.2.1.

## 15. Client-Side Caching for Files

Client-side caching of data, and of file attributes, is essential to providing good performance with the NFS protocol. Given the difficulties of providing distributed cache coherence, NFSv4 has focused on a number of techniques to reduce the need for such facilities:

- \* Using the state provided by OPEN to avoid coherence- related checks, in the common cases in which files are not shared among client unless all access is read-only.
- \* Taking advantage of the structure of existing file access API's to limit the need for inter-file coherence to be limited to that provided by the name and directory caching described in Section 16.
- \* Using share reservation to prevent the sort of sharing which would require inter-client cache coherence.

Unfortunately, the lack of support for share reservations has led to caching being used without the coherence support that would make it effective. As a result, there remains a need that will need to be addressed in future minor versions to provide ways of advising clients not to use data and attribute caching in sharing situations in which coherence would be needed.

In order to deal effectively with the lack of cache coherence, several NFS client implementation techniques have been used to deal the problems that a lack of coherence poses for users. These techniques have not been clearly defined by earlier protocol specifications, and it is often unclear what is valid or invalid client behavior.

Cases in which access to the same file is shared among multiple files including at least one writer pose particular difficulties given the absence of cache coherence. In such cases it may be necessary to avoid caching to deal with the issue. NFSv4.1 has no standard means to notify clients of the need to avoid such caching although the possibility of extensions to provide for this case is discussed in Appendix D.2.19.

The NFSv4.1 protocol uses many techniques similar to those that have been used in previous protocol versions. The NFSv4.1 protocol does not provide distributed cache coherence. However, it defines a more limited set of caching guarantees to allow locks and share reservations to be used without destructive interference from client-side caching. This leaves difficulties that have yet to be addressed satisfactorily for some environments in which share reservations are not available.

Although numerous forms of caching are dealt within the included subsections it is important to note how some of these commonly addressed separately are related in important ways. In particular, file data caching (discussed in Section 15.3 and attribute caching (discussed in Section 15.6) are more closely related than might be thought of at first, since:

- \* There are an important set of attributes that are modified by IO operations with access or modify file data.

The possibility of changes to file data and attributes tied to file data is constrained similarly by the potential existence of share reservations

In at least one important case, discussion of the rules associated with data caching is discussed only with regard to their effect on file attributes. This happens with the treatment of write-behind caching which has turned out to be quite consequential for many workloads that are important for NFSv4 to address without the difficulties result from the use of incoherent caches.

- \* They are both affected importantly by the existence of file delegations, as discussed in Section 15.4.

While delegations are helpful in dealing with the common cases in which sharing is infrequent, problems remain in dealing with files shared among multiple clients

### 15.1. Performance Challenges for Client-Side Caching

Caching techniques used in previous versions of the NFS protocol have been successful in providing good performance. However, several scalability challenges can arise when those techniques are used with very large numbers of clients. This is particularly true when clients are geographically distributed, which classically increases the latency for cache revalidation requests.

The previous versions of the NFS protocol repeat their file data cache validation requests at the time the file is opened. This behavior can have serious performance drawbacks. A common case is one in which a file is only accessed by a single client. Therefore, sharing is infrequent.

In this case, repeated references to the server to find that no conflicts exist are expensive. A better option with regards to performance is to allow a client that repeatedly opens a file to do so without reference to the server. This is done until potentially conflicting operations from another client actually occur.

A similar situation arises in connection with byte-range locking. Sending LOCK and LOCKU operations as well as the READ and WRITE operations necessary to make data caching consistent with the locking semantics (See Section 15.3.4) can severely limit performance. When locking is used to provide protection against infrequent conflicts, a large penalty is incurred. This penalty may discourage the use of byte-range locking by applications.

The NFSv4.1 protocol provides more aggressive caching strategies with the following design goals:

- \* Compatibility with a large range of server semantics.
- \* Providing the same caching benefits as previous versions of the NFS protocol when unable to support the more aggressive model.
- \* Requirements for aggressive caching are organized so that a large portion of the benefit can be obtained even when not all of the requirements can be met.

The appropriate requirements for the server are discussed in later sections in which specific forms of caching are covered (See Section 15.4).

## 15.2. Delegation and Callbacks

Recallable delegation of server responsibilities for a file to a client improves performance by avoiding repeated requests to the server in the absence of inter-client conflict. With the use of a "callback" RPC from server to client, a server recalls delegated responsibilities when another client engages in sharing of a delegated file.

A delegation is passed from the server to the client, specifying the object of the delegation and the type of delegation. There are different types of delegations, but each type contains a stateid to be used to represent the delegation when performing operations that depend on the delegation. This stateid is similar to those associated with locks and share reservations but differs in that the stateid for a delegation is associated with a client ID and may be used on behalf of all the open-owners for the given client. A delegation is made to the client as a whole and not to any specific process or thread of control within it.

The backchannel is established by CREATE\_SESSION and BIND\_CONN\_TO\_SESSION, and the client is required to maintain it. Because the backchannel may be down, even temporarily, correct protocol operation does not depend on them. Preliminary testing of backchannel functionality by means of a CB\_COMPOUND procedure with a single operation, CB\_SEQUENCE, can be used to check the continuity of the backchannel. A server avoids delegating responsibilities until it has determined that the backchannel exists. Because the granting of a delegation is always conditional upon the absence of conflicting access, clients MUST NOT assume that a delegation will be granted and they MUST always be prepared for OPENS, WANT\_DELEGATIONS, and GET\_DIR\_DELEGATIONS to be processed without any delegations being granted.

Unlike locks, an operation by a second client to a delegated file will cause the server to recall a delegation through a callback. For individual operations, we will describe, under IMPLEMENTATION, when such operations are required to effect a recall. A number of points should be noted, however.

- \* The server is free to recall a delegation whenever it feels it is desirable and may do so even if no operations requiring recall are being done.
- \* Operations done outside the NFSv4.1 protocol, due to, for example, access by other protocols including other minor version of NFSv4, or by local access, also need to result in delegation recall when they make analogous changes to file system data, including the

delegated file's contents, its attributes and the set of names linked to that file. What is crucial is if the change would invalidate the guarantees provided by the delegation. When this is possible, the delegation needs to be recalled and **MUST** be returned or revoked before allowing the operation to proceed.

- \* The semantics of the file system are crucial in defining when delegation recall is required. If a particular change within a specific implementation causes change to a file attribute, then delegation recall is required, whether that operation has been specifically listed as requiring delegation recall. Again, what is critical is whether the guarantees provided by the delegation are being invalidated.

Despite those caveats, the implementation sections for a number of operations describe situations in which delegation recall would be required under some common circumstances:

- \* For GETATTR, see Section 25.7.4.
- \* For LINK, see Section 25.9.4.
- \* For OPEN, see Section 25.16.4.
- \* For READ, see Section 25.22.4.
- \* For REMOVE, see Section 25.25.4.
- \* For RENAME, see Section 25.26.4.
- \* For SETATTR, see Section 25.30.4.
- \* For WRITE, see Section 25.32.4.

On recall, the client holding the delegation needs to flush modified state (such as modified data) to the server and return the delegation. The conflicting request will not be acted on until the recall is complete. The recall is considered complete when the client returns the delegation or the server times its wait for the delegation to be returned and revokes the delegation as a result of the timeout. In the interim, the server will either delay responding to conflicting requests or respond to them with NFS4ERR\_DELAY. Following the resolution of the recall, the server has the information necessary to grant or deny the second client's request.

At the time the client receives a delegation recall, it may have substantial state that needs to be flushed to the server. Therefore, the server should allow sufficient time for the delegation to be



returned since it may involve numerous RPCs to the server. If the server is able to determine that the client is diligently flushing state to the server as a result of the recall, the server may extend the usual time allowed for a recall. However, the time allowed for recall completion should not be unbounded.

An example of this is when responsibility to mediate opens on a given file is delegated to a client (See Section 15.4). The server will not know what opens are in effect on the client. Without this knowledge, the server will be unable to determine if the access and deny states for the file allow any particular open until the delegation for the file has been returned.

A client failure or a network partition can result in failure to respond to a recall callback. In this case, the server will revoke the delegation, which in turn will render useless any modified state still on the client.

#### 15.2.1. Delegation Recovery

There are three situations that delegation recovery needs to deal with:

- \* client restart
- \* server restart
- \* network partition (full or backchannel-only)

In the event the client restarts, establishment of a new clientid associated with the new client instance or failure to renew the lease will result in the revocation of byte-range locks and share reservations. Delegations, however, may be treated somewhat differently. It is also possible for the same sorts of revocation to occur as a result of lease non-renewal.

There will be situations in which delegations will need to be re-established after a client restarts. The reason for this is that the client may have file data stored locally and this data was associated with the previously held delegations. The client will need to re-establish the appropriate file state on the server.

To allow for this type of client recovery, the server MAY provide a special period to allow the clients to recover the delegations obtained before the restart. This special period will often be longer the typical lease expiration period. As a result, requests from other clients that conflict with these delegations would need to wait. Because the normal recall process may require significant time

for the client to flush changed state to the server, other clients need be prepared for delays that occur because of a conflicting delegation. Such a longer interval would increase the window for clients to restart and consult stable storage so that the delegations can be returned after the data is appropriately flushed to the server.

This special period, although analogous to the grace period used after server restart, is distinct from it. For OPEN delegations, such delegations are reclaimed using OPEN with a claim type of CLAIM\_DELEGATE\_PREV or CLAIM\_DELEG\_PREV\_FH (See Sections 15.5 and 25.16 for discussion of OPEN delegation and the details of OPEN, respectively). Although these types of OPENS are considered reclaim-type operations they are not, like other sorts of reclaims limited to the grace period. They are intended for use during the special delegation recovery period, and are not directly affected by possible existence of a server grace period.

A server MAY support claim types of CLAIM\_DELEGATE\_PREV and CLAIM\_DELEG\_PREV\_FH, and if it does, it MUST NOT remove delegations upon a CREATE\_SESSION that confirm a client ID created by EXCHANGE\_ID. Instead, the server MUST, for a period of time no less than that of the value of the lease\_time attribute, maintain the client's delegations to allow time for the client to send CLAIM\_DELEGATE\_PREV and/or CLAIM\_DELEG\_PREV\_FH requests. The server that supports CLAIM\_DELEGATE\_PREV and/or CLAIM\_DELEG\_PREV\_FH MUST support the DELEGPURGE operation.

When the server restarts, delegations are reclaimed (using the OPEN operation with CLAIM\_PREVIOUS) in a similar fashion to byte-range locks and share reservations. However, there is a slight semantic difference. In the normal case, if the server decides that a delegation should not be granted, it performs the requested action (e.g., OPEN) without granting any delegation. For reclaim, the server grants the delegation but a special designation is applied so that the client treats the delegation as having been granted but recalled by the server. Because of this, the client has the duty to write all modified state to the server and then return the delegation. This process of handling delegation reclaim reconciles three principles of the NFSv4.1 protocol:

- \* Upon reclaim, a client reporting resources assigned to it by an earlier server instance must be granted those resources.
- \* The server has unquestionable authority to determine whether delegations are to be granted and, once granted, whether they are to be continued.

- \* The use of callbacks should not be depended upon until the client has proven its ability to receive them.

When a client needs to reclaim a delegation and there is no associated open, the client may use the CLAIM\_PREVIOUS variant of the WANT\_DELEGATION operation. However, since the server is not required to support this operation, an alternative is to reclaim via a dummy OPEN together with the delegation using an OPEN of type CLAIM\_PREVIOUS. The dummy open file can be released using a CLOSE to re-establish the original state to be reclaimed, a delegation without an associated open.

When a client has more than a single open associated with a delegation, state for those additional opens can be established using OPEN operations of type CLAIM\_DELEGATE\_CUR. When these are used to establish opens associated with reclaimed delegations, the server MUST allow them when made within the grace period.

When a network partition occurs, delegations are subject to freeing by the server when the lease renewal period expires. This is similar to the behavior for locks and share reservations. For delegations, however, the server may extend the period in which conflicting requests are held off. Eventually, the occurrence of a conflicting request from another client will cause revocation of the delegation. A loss of the backchannel (e.g., by later network configuration change) will have the same effect. A recall request will fail and revocation of the delegation will result.

A client normally finds out about revocation of a delegation when it uses a stateid associated with a delegation and receives one of the errors NFS4ERR\_EXPIRED, NFS4ERR\_ADMIN\_REVOKED, or NFS4ERR\_DELEG\_REVOKED. It also may find out about delegation revocation after a client restart when it attempts to reclaim a delegation and receives that same error. Note that in the case of a revoked OPEN\_DELEGATE\_WRITE delegation, there are issues because data may have been modified by the client whose delegation is revoked and separately by other clients. See Section 15.5.1 for a discussion of such issues. Note also that when delegations are revoked, information about the revoked delegation will be written by the server to stable storage (as described in Section 13.4.3). This is done to deal with the case in which a server restarts after revoking a delegation but before the client holding the revoked delegation is notified about the revocation.

### 15.3. Data Caching

When applications share access to a set of files, they need to be implemented so as to take account of the possibility of conflicting access by another application. This is true whether the applications in question execute on different clients or reside on the same client.

Share reservations and byte-range locks are the facilities the NFSv4.1 protocol provides to allow applications to coordinate access by using mutual exclusion facilities. The NFSv4.1 protocol's data caching must be implemented such that it does not invalidate the assumptions on which those using these facilities depend.

#### 15.3.1. Data Caching and OPENS

In order to avoid invalidating the sharing assumptions on which applications rely, NFSv4.1 clients should not provide cached data to applications or modify it on behalf of an application when it would not be valid to obtain or modify that same data via a READ or WRITE operation.

Furthermore, in the absence of an OPEN delegation (See Section 15.4), two additional rules apply. Note that these rules are obeyed in practice by many NFSv3 clients.

- \* First, cached data present on a client must be revalidated after doing an OPEN. Revalidating means that the client fetches the change attribute from the server, compares it with the cached change attribute, and if different, treats the cached data (as well as the cached attributes) as invalid. This is to ensure that the data for the OPENed file is still correctly reflected in the client's cache. This validation must be done at least when the client's OPEN operation includes a deny of OPEN4\_SHARE\_DENY\_WRITE or OPEN4\_SHARE\_DENY\_BOTH, thus terminating a period in which other clients may have had the opportunity to open the file with OPEN4\_SHARE\_ACCESS\_WRITE/OPEN4\_SHARE\_ACCESS\_BOTH access. Clients may choose to do the revalidation more often (i.e., at OPENS specifying a deny mode of OPEN4\_SHARE\_DENY\_NONE) to parallel the NFSv3 protocol's practice for the benefit of users assuming this degree of cache revalidation.

Since the change attribute is updated for data and metadata modifications, some client implementers may be tempted to use the time\_modify attribute and not the change attribute to validate cached data, so that metadata changes do not spuriously invalidate clean data. The implementer is cautioned in this approach. The change attribute is guaranteed to change for each update to the

file, whereas `time_modify` is guaranteed to change only at the granularity of the `time_delta` attribute. Use by the client's data cache validation logic of `time_modify` and not `change` runs the risk of the client incorrectly marking stale data as valid. Thus, any cache validation approach by the client MUST include the use of the `change` attribute.

- \* Second, modified data must be flushed to the server before closing a file OPENed for access including `OPEN4_SHARE_ACCESS_WRITE`. This is complementary to the first rule. If the data is not flushed at `CLOSE`, the revalidation done after the client `OPENs` a file is unable to achieve its purpose. The other aspect to flushing the data before close is that the data must be committed to stable storage, at the server, before the `CLOSE` operation is requested by the client. In the case of a server restart and a `CLOSEd` file, it may not be possible to retransmit the data to be written to the file, hence, this requirement.

The two rules above are often described as providing "Close-to-open consistency". See Section 15.3.2 for a thorough discussion of this term and its role during the development of NFS.

#### 15.3.2. Data Caching and Close-to-open Consistency

The two rules in the section above are often described as providing "Close-to-open consistency". While this term has not been mentioned in previous NFSv4 minor version specifications, the reference in [RFC8881] to "rules (that) are obeyed in practice by many NFSv3 clients" is clearly meant to include these rules as, normally described as providing "clos-to-open consistency" as part of NFSv4's way of dealing with caching and it has been widely adopted by NFSv4 clients. It is now necessary to provide a clearer treatment of these issues for the following reasons:

- \* Previous treatments have not provided descriptions of or references to the sharing problems that result if additional care is not taken to avoid forms of caching that are problematic when inter-client sharing of open files occurs.
- \* The narrative by which the requirements are explained obscures the fact that sharing issues are not addressed.

Such descriptions cite what needs to be done after the `CLOSE` and before the next `OPEN`, without considering the possibility that the `OPEN` happens before the `CLOSE`, or addressing the issues that would arise in that case, i.e. when the file is open by multiple clients.

Handling of such sharing is discussed in Section 15.3.3 together with a discussion of issues involved in moving the NFSv4 protocol from its current approach based on NFSv3's approach to caching in which such issues were ignored to something more able to deal with inter-client sharing.

#### 15.3.3. Data Caching and Files Open for Write

When a file is being written, it is possible, because of the lack of data cache coherence facilities, for otherwise unexceptionable data caching arrangements to create problems because each client is unaware of changes made by other clients:

- \* The provision for write-behind caching discussed in Section 15.6.1 remains of dubious utility and a source of potential difficulty, even though it is commonly used and cannot be disallowed at this time.

Fortunately, because there is no longer a performance-based imperative to adopt write-behind caching, clients are free to avoid this, using a mount option, as they have been doing. It is possible for the server, using a potential extension, to provide advice specifying that this needs to be avoided, at least when the file is opened for write..

- \* Many other instances of read caching become troublesome when used in situations in which the file is being written unexpectedly, often by multiple clients.

Although, in theory, this could be prevented by readers opening with write being denied, the lack of appropriate provisions to specify this option in many common environments prevents the issue from being addressed using existing facilities.

It is possible to address this issue using a client mount option suppressing write caching or via an extension providing information about the existence of clients who have the file open for write. Because of the negative effect of such suppression on workloads in which the sharing of access to files being written is rare or non-existent, it is desirable to be able to provide updates to readers about the existence of file open for write.

The discussion of data caching in this and earlier ( 15.3.1, 15.3.2) sections has left us with two different approaches to data caching:

- \* One, derived from NFSv3 practice and implicitly endorsed by [RFC8881], is commonly implemented and works in many environments in which sharing of open files is avoided, as it was in NFSv3. Despite this, the reliance on this approach exacerbates NFS's inability to deal correctly with inter-client sharing
- \* The other, outlined in Section 15.3.2, is more suited to cases in which inter-client sharing is present but imposes performance penalties in many important environments in which inter-client sharing is typically avoided as it was in NFSv3.

As a result, despite our obligation to explain these issues clearly, we face difficulties in providing useful guidance to implementers going forward, with the likely result that addressing inter-client sharing in NFSv4 is likely to be delayed to later minor versions. These difficulties are not immediately addressable because:

- \* The incorrect treatment in previous specifications, is pretty much beside the point, since the NFSv3 practice would be followed whether the specification said so or not, as long it did not clearly forbid it.
- \* The work to fully address sharing issues is of a far larger scope than could be done in a bis document
- \* Even the work to address caching issues where sharing is common is likely to require a NFSv4.2 extension.

While addressing this issue is not possible in this document, there is useful discussion of future possibilities in Appendices D.2.18 through D.2.20.

#### 15.3.4. Data Caching and File Locking

For those applications that choose to use byte-range locking instead of share reservations to exclude inconsistent file access, there is an analogous set of constraints that apply to client-side data caching. These rules are effective only if the byte-range locking is used in a way that matches in an equivalent way the actual READ and WRITE operations executed. This is as opposed to byte-range locking that is based on pure convention. For example, it is possible to manipulate a two-megabyte file by dividing the file into two one-megabyte ranges and protecting access to the two byte-ranges by byte-range locks on bytes zero and one. A WRITE\_LT lock on byte zero of the file would represent the right to perform READ and WRITE operations on the first byte-range. A WRITE\_LT lock on byte one of the file would represent the right to perform READ and WRITE operations on the second byte-range. As long as all applications

manipulating the file obey this convention, they will work on a local file system. However, they may not work with the NFSv4.1 protocol unless clients refrain from data caching.

The rules for data caching in the byte-range locking environment are:

- \* First, when a client obtains a byte-range lock for a particular byte-range, the data cache corresponding to that byte-range (if any cache data exists) must be revalidated. If the change attribute indicates that the file may have been updated since the cached data was obtained, the client must flush or invalidate the cached data for the newly locked byte-range. A client might choose to invalidate all of the non-modified cached data that it has for the file, but the only requirement for correct operation is to invalidate all of the data in the newly locked byte-range.
- \* Second, before releasing a `WRITE_LT` lock for a byte-range, all modified data for that byte-range must be flushed to the server. The modified data must also be written to stable storage.

Note that flushing data to the server and the invalidation of cached data must reflect the actual byte-ranges locked or unlocked. Rounding these up or down to reflect client cache block boundaries will cause problems if not carefully done. For example, writing a modified block when only half of that block is within an area being unlocked may cause invalid modification to the byte-range outside the unlocked area. This, in turn, may be part of a byte-range locked by another client. Clients can avoid this situation by synchronously performing portions of `WRITE` operations that overlap that portion (initial or final) that is not a full block. Similarly, invalidating a locked area that is not an integral number of full buffer blocks would require the client to read one or two partial blocks from the server if the revalidation procedure shows that the data that the client possesses may not be valid.

The data that is written to the server as a prerequisite to the unlocking of a byte-range must be written, at the server, to stable storage. The client may accomplish this either with synchronous writes or by following asynchronous writes with a `COMMIT` operation. This is required because retransmission of the modified data after a server restart might conflict with a lock held by another client.

A client implementation may choose to accommodate applications that use byte-range locking in non-standard ways (e.g., using a byte-range lock as a global semaphore) by flushing to the server more data upon a `LOCKU` than is covered by the locked range. This may include modified data within files other than the one for which the unlocks are being done. In such cases, the client must not interfere with



applications whose READs and WRITEs are being done only within the bounds of byte-range locks that the application holds. For example, an application locks a single byte of a file and proceeds to write that single byte. A client that chose to handle a LOCKU by flushing all modified data to the server could validly write that single byte in response to an unrelated LOCKU operation. However, it would not be valid to write the entire block in which that single written byte was located since it includes an area that is not locked and might be locked by another client. Client implementations can avoid this problem by dividing files with modified data into those for which all modifications are done to areas covered by an appropriate byte-range lock and those for which there are modifications not covered by a byte-range lock. Any writes done for the former class of files must not include areas not locked and thus not modified on the client.

#### 15.3.5. Data Caching and Mandatory File Locking

Client-side data caching needs to respect mandatory byte-range locking when it is in effect. The presence of mandatory byte-range locking for a given file is indicated when the client gets back NFS4ERR\_LOCKED from a READ or WRITE operation on a file for which it has an appropriate share reservation. When mandatory locking is in effect for a file, the client must check for an appropriate byte-range lock for data being read or written. If a byte-range lock exists for the range being read or written, the client may satisfy the request using the client's validated cache. If an appropriate byte-range lock is not held for the range of the read or write, the read or write request must not be satisfied by the client's cache and the request must be sent to the server for processing. When a read or write request partially overlaps a locked byte-range, the request should be subdivided into multiple pieces with each byte-range (locked or not) treated appropriately.

#### 15.3.6. Data Caching and File Identity

When clients cache data, the file data needs to be organized according to the file system object to which the data belongs. For NFSv3 clients, the typical practice has been to assume for the purpose of caching that distinct filehandles represent distinct file system objects. The client then has the choice to organize and maintain the data cache on this basis.

In the NFSv4.1 protocol, there is now the possibility to have significant deviations from a "one filehandle per object" model because a filehandle may be constructed on the basis of the object's pathname. Therefore, clients need a reliable method to determine if two filehandles designate the same file system object. If clients were simply to assume that all distinct filehandles denote distinct

objects and proceed to do data caching on this basis, caching inconsistencies would arise between the distinct client-side objects that mapped to the same server-side object.

By providing a method to differentiate filehandles, the NFSv4.1 protocol alleviates a potential functional regression in comparison with the NFSv3 protocol. Without this method, caching inconsistencies within the same client could occur, and this has not been present in previous versions of the NFS protocol. Note that it is possible to have such inconsistencies with applications executing on multiple clients, but that is not the issue being addressed here.

For the purposes of data caching, the following steps allow an NFSv4.1 client to determine whether two distinct filehandles denote the same server-side object:

- \* If GETATTR directed to two filehandles returns different values of the fsid attribute, then the filehandles represent distinct objects.
- \* If GETATTR for any file with an fsid that matches the fsid of the two filehandles in question returns a unique\_handles attribute with a value of TRUE, then the two objects are distinct.
- \* If GETATTR directed to the two filehandles does not return the fileid attribute for both of the handles, then it cannot be determined whether the two objects are the same. Therefore, operations that depend on that knowledge (e.g., client-side data caching) cannot be done reliably. Note that if GETATTR does not return the fileid attribute for both filehandles, it will return it for neither of the filehandles, since the fsid for both filehandles is the same.
- \* If GETATTR directed to the two filehandles returns different values for the fileid attribute, then they are distinct objects.
- \* Otherwise, they are the same object.

#### 15.4. Open Delegation

When a file is being OPENed, the server may delegate further handling of opens and closes for that file to the opening client. Any such delegation is recallable since the circumstances that allowed for the delegation are subject to change. In particular, if the server receives a conflicting OPEN from another client, the server must recall the delegation before deciding whether the OPEN from the other client may be granted. Making a delegation is up to the server, and clients should not assume that any particular OPEN either will or

will not result in an OPEN delegation. The following is a typical set of conditions that servers might use in deciding whether an OPEN should be delegated:

- \* The client must be able to respond to the server's callback requests. If a backchannel has been established, the server will send a CB\_COMPOUND request, containing a single operation, CB\_SEQUENCE, for a test of backchannel availability.
- \* The client must have responded properly to previous recalls.
- \* There must be no current OPEN conflicting with the requested delegation.
- \* There should be no current delegation that conflicts with the delegation being requested.
- \* The probability of future conflicting open requests should be low based on the recent history of the file.
- \* The existence of any server-specific semantics of OPEN/CLOSE that would make the required handling incompatible with the prescribed handling that the delegated client would apply (See below).

There are two types of OPEN delegations: OPEN\_DELEGATE\_READ and OPEN\_DELEGATE\_WRITE. An OPEN\_DELEGATE\_READ delegation allows a client to handle, on its own, requests to open a file for reading that do not deny OPEN4\_SHARE\_ACCESS\_READ access to others. Multiple OPEN\_DELEGATE\_READ delegations may be outstanding simultaneously and do not conflict. An OPEN\_DELEGATE\_WRITE delegation allows the client to handle, on its own, all opens. Only one OPEN\_DELEGATE\_WRITE delegation may exist for a given file at a given time, and it is inconsistent with any OPEN\_DELEGATE\_READ delegations.

When a client has either type of open delegation, it is assured that neither the contents, the attributes (with the exception of time\_access), nor the names of any links to the file will change without its knowledge, so long as the delegation is held. When a client has an OPEN\_DELEGATE\_WRITE delegation, it may modify the file data locally since no other client will be accessing the file's data. The client holding an OPEN\_DELEGATE\_WRITE delegation may only locally affect file attributes that are intimately connected with the file data: size, change, time\_access, time\_metadata, and time\_modify. All other attributes must be reflected on the server.

When a client has an OPEN delegation, it does not need to send OPENS or CLOSEs to the server. Instead, the client may update the appropriate status internally. For an OPEN\_DELEGATE\_READ delegation,

opens that cannot be handled locally (opens that are for OPEN4\_SHARE\_ACCESS\_WRITE/OPEN4\_SHARE\_ACCESS\_BOTH or that deny OPEN4\_SHARE\_ACCESS\_READ access) must be sent to the server.

When an OPEN delegation is made, the reply to the OPEN contains an OPEN delegation structure that specifies the following:

- \* the type of delegation (OPEN\_DELEGATE\_READ or OPEN\_DELEGATE\_WRITE).
- \* space limitation information to control flushing of data on close (OPEN\_DELEGATE\_WRITE delegation only; see Section 15.4.1)
- \* an nfsace4 specifying read and write permissions
- \* a stateid to represent the delegation

The delegation stateid is separate and distinct from the stateid for the OPEN proper. The standard stateid, unlike the delegation stateid, is associated with a particular lock-owner and will continue to be valid after the delegation is recalled and the file remains open.

When a request internal to the client is made to open a file and an OPEN delegation is in effect, it will be accepted or rejected solely on the basis of the following conditions. Any requirement for other checks to be made by the delegate should result in the OPEN delegation being denied so that the checks can be made by the server itself.

- \* The access and deny bits for the request and the file as described in Section 14.7.
- \* The read and write permissions as determined below.

The nfsace4 passed with delegation can be used to avoid frequent ACCESS calls. The permission check should be as follows:

- \* If the nfsace4 indicates that the open may be done, then it should be granted without reference to the server.
- \* If the nfsace4 indicates that the open may not be done, then an ACCESS request must be sent to the server to obtain the definitive answer.

The server may return an nfsace4 that is more restrictive than the actual ACL of the file. This includes an nfsace4 that specifies denial of all access. Note that some common practices such as

mapping the traditional user "root" to the user "nobody" (See Section 11.13) may make it incorrect to return the actual ACL of the file in the delegation response.

The use of a delegation together with various other forms of caching creates the possibility that no server authentication and authorization will ever be performed for a given user since all of the user's requests might be satisfied locally. Where the client is depending on the server for authentication and authorization, the client should be sure authentication and authorization occurs for each user by use of the ACCESS operation. This should be the case even if an ACCESS operation would not be required otherwise. As mentioned before, the server may enforce frequent authentication by returning an nfsace4 denying all access with every OPEN delegation.

#### 15.4.1. Open Delegation and Data Caching

An OPEN delegation allows much of the message overhead associated with the opening and closing files to be eliminated. An open when an OPEN delegation is in effect does not require that a validation message be sent to the server. The continued endurance of the "OPEN\_DELEGATE\_READ delegation" provides a guarantee that no OPEN for OPEN4\_SHARE\_ACCESS\_WRITE/OPEN4\_SHARE\_ACCESS\_BOTH, and thus no write, has occurred. Similarly, when closing a file opened for OPEN4\_SHARE\_ACCESS\_WRITE/OPEN4\_SHARE\_ACCESS\_BOTH and if an OPEN\_DELEGATE\_WRITE delegation is in effect, the data written does not have to be written to the server until the OPEN delegation is recalled. The continued endurance of the OPEN delegation provides a guarantee that no open, and thus no READ or WRITE, has been done by another client.

For the purposes of OPEN delegation, READS and WRITES done without an OPEN are treated as the functional equivalents of a corresponding type of OPEN. Although a client should not use special stateids when an open exists, delegation handling on the server can use the client ID associated with the current session to determine if the operation has been done by the holder of the delegation (in which case, no recall is necessary) or by another client (in which case, the delegation must be recalled and I/O not proceed until the delegation is returned or revoked).

With delegations, a client is able to avoid writing data to the server when the CLOSE of a file is serviced. The file close system call is the usual point at which the client is notified of a lack of stable storage for the modified file data generated by the application. At the close, file data is written to the server and, through normal accounting, the server is able to determine if the available file system space for the data has been exceeded (i.e., the

server returns NFS4ERR\_NOSPC or NFS4ERR\_DQUOT). This accounting includes quotas. The introduction of delegations requires that an alternative method be in place for the same type of communication to occur between client and server.

In the delegation response, the server provides either the limit of the size of the file or the number of modified blocks and associated block size. The server must ensure that the client will be able to write modified data to the server of a size equal to that provided in the original delegation. The server must make this assurance for all outstanding delegations. Therefore, the server must be careful in its management of available space for new or modified data, taking into account available file system space and any applicable quotas. The server can recall delegations as a result of managing the available file system space. The client should abide by the server's state space limits for delegations. If the client exceeds the stated limits for the delegation, the server's behavior is undefined.

Based on server conditions, quotas, or available file system space, the server may grant OPEN\_DELEGATE\_WRITE delegations with very restrictive space limitations. The limitations may be defined in a way that will always force modified data to be flushed to the server on close.

With respect to authentication, flushing modified data to the server after a CLOSE has occurred may be problematic. For example, the user of the application may have logged off the client, and unexpired authentication credentials may not be present. In this case, the client may need to take special care to ensure that local unexpired credentials will in fact be available. This may be accomplished by tracking the expiration time of credentials and flushing data well in advance of their expiration or by making private copies of credentials to assure their availability when needed.

#### 15.4.2. Open Delegation and File Locks

When a client holds an OPEN\_DELEGATE\_WRITE delegation, lock operations are performed locally. This includes those required for mandatory byte-range locking. This can be done since the delegation implies that there can be no conflicting locks. Similarly, all of the revalidations that would normally be associated with obtaining locks and the flushing of data associated with the releasing of locks need not be done.

When a client holds an OPEN\_DELEGATE\_READ delegation, lock operations are not performed locally. All lock operations, including those requesting non-exclusive locks, are sent to the server for resolution.

#### 15.4.3. Handling of CB\_GETATTR

The server needs to employ special handling for a GETATTR where the target is a file that has an OPEN\_DELEGATE\_WRITE delegation in effect. The reason for this is that the client holding the OPEN\_DELEGATE\_WRITE delegation may have modified the data, and the server needs to reflect this change to the second client that submitted the GETATTR. Therefore, the client holding the OPEN\_DELEGATE\_WRITE delegation needs to be interrogated. The server will use the CB\_GETATTR operation. The only attributes that the server can reliably query via CB\_GETATTR are size and change.

Since CB\_GETATTR is being used to satisfy another client's GETATTR request, the server only needs to know if the client holding the delegation has a modified version of the file. If the client's copy of the delegated file is not modified (data or size), the server can satisfy the second client's GETATTR request from the attributes stored locally at the server. If the file is modified, the server only needs to know about this modified state. If the server determines that the file is currently modified, it will respond to the second client's GETATTR as if the file had been modified locally at the server.

Since the form of the change attribute is determined by the server and is opaque to the client, the client and server need to agree on a method of communicating the modified state of the file. For the size attribute, the client will report its current view of the file size. For the change attribute, the handling is more involved.

For the client, the following steps will be taken when receiving an OPEN\_DELEGATE\_WRITE delegation:

- \* The value of the change attribute will be obtained from the server and cached. Let this value be represented by c.
- \* The client will create a value greater than c that will be used for communicating that modified data is held at the client. Let this value be represented by d.
- \* When the client is queried via CB\_GETATTR for the change attribute, it checks to see if it holds modified data. If the file is modified, the value d is returned for the change attribute value. If this file is not currently modified, the client returns the value c for the change attribute.

For simplicity of implementation, the client MAY for each CB\_GETATTR return the same value d. This is true even if, between successive CB\_GETATTR operations, the client again modifies the file's data or

metadata in its cache. The client can return the same value because the only requirement is that the client be able to indicate to the server that the client holds modified data. Therefore, the value of `d` may always be `c + 1`.

While the change attribute is opaque to the client in the sense that it has no idea what units of time, if any, the server is counting change with, it is not opaque in that the client has to treat it as an unsigned integer, and the server has to be able to see the results of the client's changes to that integer. Therefore, the server MUST encode the change attribute in network order when sending it to the client. The client MUST decode it from network order to its native order when receiving it, and the client MUST encode it in network order when sending it to the server. For this reason, change is defined as an unsigned integer rather than an opaque array of bytes.

For the server, the following steps will be taken when providing an `OPEN_DELEGATE_WRITE` delegation:

- \* Upon providing an `OPEN_DELEGATE_WRITE` delegation, the server will cache a copy of the change attribute in the data structure it uses to record the delegation. Let this value be represented by `sc`.
- \* When a second client sends a `GETATTR` operation on the same file to the server, the server obtains the change attribute from the first client. Let this value be `cc`.
- \* If the value `cc` is equal to `sc`, the file is not modified and the server returns the current values for change, `time_metadata`, and `time_modify` (for example) to the second client.
- \* If the value `cc` is NOT equal to `sc`, the file is currently modified at the first client and most likely will be modified at the server at a future time. The server then uses its current time to construct attribute values for `time_metadata` and `time_modify`. A new value of `sc`, which we will call `nsc`, is computed by the server, such that `nsc >= sc + 1`. The server then returns the constructed `time_metadata`, `time_modify`, and `nsc` values to the requester. The server replaces `sc` in the delegation record with `nsc`. To prevent the possibility of `time_modify`, `time_metadata`, and change from appearing to go backward (which would happen if the client holding the delegation fails to write its modified data to the server before the delegation is revoked or returned), the server needs to update the file's attributes reflect constructed attribute values. For reasons of reasonable performance, committing the constructed attribute values to stable storage is OPTIONAL.



As discussed earlier in this section, the client MAY return the same cc value on subsequent CB\_GETATTR calls, even if the file was modified in the client's cache yet again between successive CB\_GETATTR calls. Therefore, the server must assume that the file has been modified yet again, and MUST take care to ensure that the new nsc it constructs and returns is greater than the previous nsc it returned. An example implementation's delegation record would satisfy this mandate by including a boolean field (let us call it "modified") that is set to FALSE when the delegation is granted, and an sc value set at the time of grant to the change attribute value. The modified field would be set to TRUE the first time cc != sc, and would stay TRUE until the delegation is returned or revoked. The processing for constructing nsc, time\_modify, and time\_metadata would use this pseudo code:

```
if (!modified) {
    do CB_GETATTR for change and size;

    if (cc != sc)
        modified = TRUE;
} else {
    do CB_GETATTR for size;
}

if (modified) {
    sc = sc + 1;
    time_modify = time_metadata = current_time;
    update sc, time_modify, time_metadata into file's metadata;
}
```

This would return to the client (that sent GETATTR) the attributes it requested, but make sure size comes from what CB\_GETATTR returned. The server would not update the file's metadata with the client's modified size.

In the case that the file attribute size is different than the server's current value, the server treats this as a modification regardless of the value of the change attribute retrieved via CB\_GETATTR and responds to the second client as in the last step.

This methodology resolves issues of clock differences between client and server and other scenarios where the use of CB\_GETATTR break down.

It should be noted that the server is under no obligation to use CB\_GETATTR, and therefore the server MAY simply recall the delegation to avoid its use.

#### 15.4.4. Recall of Open Delegation

The following events necessitate recall of an OPEN delegation:

- \* potentially conflicting OPEN request (or a READ or WRITE operation done with a special stateid)
- \* SETATTR sent by another client
- \* REMOVE request for the file
- \* RENAME request for the file as either the source or target of the RENAME

Whether a RENAME of a directory in the path leading to the file results in recall of an OPEN delegation depends on the semantics of the server's file system. If that file system denies such RENAMES when a file is open, the recall must be performed to determine whether the file in question is, in fact, open.

In addition to the situations above, the server may choose to recall OPEN delegations at any time if resource constraints make it advisable to do so. Clients should always be prepared for the possibility of recall.

When a client receives a recall for an OPEN delegation, it needs to update state on the server before returning the delegation. These same updates must be done whenever a client chooses to return a delegation voluntarily. The following items of state need to be dealt with:

- \* If the file associated with the delegation is no longer open and no previous CLOSE operation has been sent to the server, a CLOSE operation must be sent to the server.
- \* If a file has other open references at the client, then OPEN operations must be sent to the server. The appropriate stateids will be provided by the server for subsequent use by the client since the delegation stateid will no longer be valid. These OPEN requests are done with the claim type of CLAIM\_DELEGATE\_CUR. This will allow the presentation of the delegation stateid so that the client can establish the appropriate rights to perform the OPEN. (See Section 25.16, which describes the OPEN operation, for details.)
- \* If there are granted byte-range locks, the corresponding LOCK operations need to be performed. This applies to the OPEN\_DELEGATE\_WRITE delegation case only.

- \* For an OPEN\_DELEGATE\_WRITE delegation, if at the time of recall the file is not open for OPEN4\_SHARE\_ACCESS\_WRITE/OPEN4\_SHARE\_ACCESS\_BOTH, all modified data for the file must be flushed to the server. If the delegation had not existed, the client would have done this data flush before the CLOSE operation.
- \* For an OPEN\_DELEGATE\_WRITE delegation when a file is still open at the time of recall, any modified data for the file needs to be flushed to the server.
- \* With the OPEN\_DELEGATE\_WRITE delegation in place, it is possible that the file was truncated during the duration of the delegation. For example, the truncation could have occurred as a result of an OPEN\_UNCHECKED with a size attribute value of zero. Therefore, if a truncation of the file has occurred and this operation has not been propagated to the server, the truncation must occur before any modified data is written to the server.

In the case of OPEN\_DELEGATE\_WRITE delegation, byte-range locking imposes some additional requirements. To precisely maintain the associated invariant, it is required to flush any modified data in any byte-range for which a WRITE\_LT lock was released while the OPEN\_DELEGATE\_WRITE delegation was in effect. However, because the OPEN\_DELEGATE\_WRITE delegation implies no other locking by other clients, a simpler implementation is to flush all modified data for the file (as described just above) if any WRITE\_LT lock has been released while the OPEN\_DELEGATE\_WRITE delegation was in effect.

An implementation need not wait until delegation recall (or the decision to voluntarily return a delegation) to perform any of the above actions, if implementation considerations (e.g., resource availability constraints) make that desirable. Generally, however, the fact that the actual OPEN state of the file may continue to change makes it not worthwhile to send information about opens and closes to the server, except as part of delegation return. An exception is when the client has no more internal opens of the file. In this case, sending a CLOSE is useful because it reduces resource utilization on the client and server. Regardless of the client's choices on scheduling these actions, all must be performed before the delegation is returned, including (when applicable) the close that corresponds to the OPEN that resulted in the delegation. These actions can be performed either in previous requests or in previous operations in the same COMPOUND request.

#### 15.4.5. Clients That Fail to Honor Delegation Recalls

A client may fail to respond to a recall for various reasons, such as a failure of the backchannel from server to the client. The client may be unaware of a failure in the backchannel. This lack of awareness could result in the client finding out long after the failure that its delegation has been revoked, and another client has modified the data for which the client had a delegation. This is especially a problem for the client that held an OPEN\_DELEGATE\_WRITE delegation.

Status bits returned by SEQUENCE operations help to provide an alternate way of informing the client of issues regarding the status of the backchannel and of recalled delegations. When the backchannel is not available, the server returns the status bit SEQ4\_STATUS\_CB\_PATH\_DOWN on SEQUENCE operations. The client can react by attempting to re-establish the backchannel and by returning recallable objects if a backchannel cannot be successfully re-established.

Whether the backchannel is functioning or not, it may be that the recalled delegation is not returned. Note that the client's lease might still be renewed, even though the recalled delegation is not returned. In this situation, servers need to revoke delegations that are not returned in a period of time equal to the lease period. This period of time should allow the client time to note the backchannel-down status and re-establish the backchannel.

When delegations are revoked, the server will return with the SEQ4\_STATUS\_RECALLABLE\_STATE\_REVOKED status bit set on subsequent SEQUENCE operations. The client should note this and then use TEST\_STATEID to find which delegations have been revoked.

#### 15.4.6. Delegation Revocation

At the point a delegation is revoked, if there are associated opens on the client, these opens may or may not be revoked. If no byte-range lock or open is granted that is inconsistent with the existing open, the stateid for the open may remain valid and be disconnected from the revoked delegation, just as would be the case if the delegation were returned.

For example, if an OPEN for OPEN4\_SHARE\_ACCESS\_BOTH with a deny of OPEN4\_SHARE\_DENY\_NONE is associated with the delegation, granting of another such OPEN to a different client will revoke the delegation but need not revoke the OPEN, since the two OPENs are consistent with each other. On the other hand, if an OPEN denying write access is granted, then the existing OPEN must be revoked.

When opens and/or locks are revoked, the applications holding these opens or locks need to be notified. This notification usually occurs by returning errors for READ/WRITE operations or when a close is attempted for the open file.

If no opens exist for the file at the point the delegation is revoked, then notification of the revocation is unnecessary. However, if there is modified data present at the client for the file, the user of the application should be notified. Unfortunately, it may not be possible to notify the user since active applications may not be present at the client. See Section 15.5.1 for additional details.

#### 15.4.7. Delegations via WANT\_DELEGATION

In addition to providing delegations as part of the reply to OPEN operations, servers MAY provide delegations separate from open, via the OPTIONAL WANT\_DELEGATION operation. This allows delegations to be obtained in advance of an OPEN that might benefit from them, for objects that are not a valid target of OPEN, or to deal with cases in which a delegation has been recalled and the client wants to make an attempt to re-establish it if the absence of use by other clients allows that.

The WANT\_DELEGATION operation may be performed on any type of file object other than a directory.

When a delegation is obtained using WANT\_DELEGATION, any open files for the same filehandle held by that client are to be treated as subordinate to the delegation, just as if they had been created using an OPEN of type CLAIM\_DELEGATE\_CUR. They are otherwise unchanged as to seqid, access and deny modes, and the relationship with byte-range locks. Similarly, because existing byte-range locks are subordinate to an open, those byte-range locks also become indirectly subordinate to that new delegation.

The WANT\_DELEGATION operation provides for delivery of delegations via callbacks, when the delegations are not immediately available. When a requested delegation is available, it is delivered to the client via a CB\_PUSH\_DELEG operation. When this happens, open files for the same filehandle become subordinate to the new delegation at the point at which the delegation is delivered, just as if they had been created using an OPEN of type CLAIM\_DELEGATE\_CUR. Similarly, this occurs for existing byte-range locks subordinate to an open.

### 15.5. Data Caching and Revocation

When locks and delegations are revoked, the assumptions upon which successful caching depends are no longer guaranteed. For any locks or share reservations that have been revoked, the corresponding state-owner needs to be notified. This notification includes applications with a file open that has a corresponding delegation that has been revoked. Cached data associated with the revocation must be removed from the client. In the case of modified data existing in the client's cache, that data must be removed from the client without being written to the server. As mentioned, the assumptions made by the client are no longer valid at the point when a lock or delegation has been revoked. For example, another client may have been granted a conflicting byte-range lock after the revocation of the byte-range lock at the first client. Therefore, the data within the lock range may have been modified by the other client. Obviously, the first client is unable to guarantee to the application what has occurred to the file in the case of revocation.

Notification to a state-owner will in many cases consist of simply returning an error on the next and all subsequent READs/WRITEs to the open file or on the close. Where the methods available to a client make such notification impossible because errors for certain operations may not be returned, more drastic action such as signals or process termination may be appropriate. The justification here is that an invariant on which an application depends may be violated. Depending on how errors are typically treated for the client-operating environment, further levels of notification including logging, console messages, and GUI pop-ups may be appropriate.

#### 15.5.1. Revocation Recovery for Write Open Delegation

Revocation recovery for an OPEN\_DELEGATE\_WRITE delegation poses the special issue of modified data in the client cache while the file is not open. In this situation, any client that does not flush modified data to the server on each close must ensure that the user receives appropriate notification of the failure as a result of the revocation. Since such situations may require human action to correct problems, notification schemes in which the appropriate user or administrator is notified may be necessary. Logging and console messages are typical examples.

If there is modified data on the client, it must not be flushed normally to the server. A client may attempt to provide a copy of the file data as modified during the delegation under a different name in the file system namespace to ease recovery. Note that when the client can determine that the file has not been modified by any other client, or when the client has a complete cached copy of the

file in question, such a saved copy of the client's view of the file may be of particular value for recovery. In another case, recovery using a copy of the file based partially on the client's cached data and partially on the server's copy as modified by other clients will be anything but straightforward, so clients may avoid saving file contents in these situations or specially mark the results to warn users of possible problems.

Saving of such modified data in delegation revocation situations may be limited to files of a certain size or might be used only when sufficient disk space is available within the target file system. Such saving may also be restricted to situations when the client has sufficient buffering resources to keep the cached copy available until it is properly stored to the target file system.

#### 15.6. Attribute Caching

This section focuses on the caching of a file's attributes by clients that do not hold a delegation on the file.

The attributes discussed in this section do not include named attributes. Individual named attributes are treated similarly to files, and caching of the data for these files needs to be handled just as data caching is for ordinary files. Similarly, LOOKUP results from an OPENATTR directory (as well as the directory's contents) can be cached on the same basis as other pathnames.

Within this section and included subsections, attributes intimately connected with the file's data, termed "data-related attributes" often have a special role and there are useful constraints on the circumstances allowing these attributes to be changed, beyond those provided by delegations. These attributes consist of the following:

- \* Size
- \* Modify time
- \* Change

Clients may cache file attributes obtained from the server and use them to avoid subsequent GETATTR requests. However, except when the constraints listed below are relevant, the client needs to understand that because of the lack of a coherence mechanism, such values can become outdated due to actions requested by other clients.

- \* The holder of a write delegation is assured that other clients cannot cause modification of any file attributes.

- \* The holder of a read delegation is assured that other clients cannot cause modification of any file attributes other than access time.
- \* The client associated with an open denying WRITE is assured that other clients cannot cause modification of any of the data-related attributes.
- \* Certain attributes, are by their nature, not subject to change. These include `fsid` and `file_id`.

Clients caching attributes need to be aware of the possibility that cached attributes may become invalid in the common cases in which none of constraints cited above prevents their modification. This makes it necessary that it be able to determine whether the cached attributes are still valid.

A client can validate its cached version of attributes for a file by fetching both the `change` and `time_access` attributes and assuming that if the `change` attribute has the same value as it did when the attributes were cached, then no attributes other than `time_access` have changed. The inclusion of `time_access` is desirable since it is often costless to add this to the set of attributes being fetched. In this way, a current value of `access_time` is made available to deal with the common case in which it is the only attribute modified.

#### 15.6.1. Attribute Caching Coherence Between Requesting Client and Server

The caching as discussed above in Section 15.6 proper is write-through in that modification to file attributes is always done by means of requests to the server and should not be done locally and should not be cached. One potential exception to this are modifications to data-related attributes that occur as a result of write-behind caching. When this option is used, extending a file by writing data to the local data cache is reflected immediately in the size as seen on the client without this change being immediately reflected on the server or visible to other clients. In this case, such changes are not propagated directly to the server, but when the modified data is flushed to the server, corresponding attribute changes are made on the server. When a write delegation is in held by the client, the modified attributes may be returned to the server in response to a `CB_RECALL` call.

Note that if the full set of attributes to be cached is requested by `REaddir` and returned by the server, the results can be cached by the client on the same basis as attributes obtained via `GETATTR`.



The client may maintain a cache of modified attributes for those attributes intimately connected with data of modified regular files (size, time\_modify, and change). Other than those three attributes, the client MUST NOT maintain a cache of modified attributes. Instead, attribute changes are immediately sent to the server.

In some operating environments, the equivalent to time\_access is expected to be implicitly updated by each read of the content of the file object. If an NFS client is caching the content of a file object, whether it is a regular file, directory, or symbolic link, the client need not update the time\_access attribute (via SETATTR or a small READ or READDIR request) on the server with each read that is satisfied from cache. The reason is that this can defeat the performance benefits of caching content, especially since an explicit SETATTR of time\_access may alter the change attribute on the server. If the change attribute changes, clients that are caching the content will think the content has changed, and will re-read unmodified data from the server. Nor is the client encouraged to maintain a modified version of time\_access in its cache, since the client either would eventually have to write the access time to the server with bad performance effects or never update the server's time\_access, thereby resulting in a situation where an application that caches access time between a close and open of the same file observes the access time oscillating between the past and present. The time\_access attribute always means the time of last access to a file by a read that was satisfied by the server. This way clients will tend to see only time\_access changes that go forward in time.

#### 15.6.2. Attribute Caching Coherence Between Requesting Client and Other Clients

The changes made by various clients to file attributes is likely to result in multiple clients having different values for file attributes. Because typical file system application programming interfaces do not provide means to atomically modify or interrogate attributes for multiple files at the same time this incoherence can be dealt with as discussed below. The following guidelines provide an environment where the potential incoherencies alluded to above can be reasonably managed. These guidelines are derived from the practice of previous NFS protocols, modified to deal with the larger set of NFSv4.1 file attributes.

- \* A large set of attributes for a given file (not including per-fsid attributes excepted) are cached as a unit at the client to avoid non-serializability can arise within the context of a single file.

- \* An upper time boundary is maintained on how long a client cache entry can be kept without being refreshed from the server. Actual refresh is unnecessary if it can be determined that no attributes have changed
- \* When operations are performed that change attributes at the server, the updated attribute set is requested as part of the containing RPC. This includes directory operations that result in the update of attributes. This is accomplished by following the modifying operation with a GETATTR operation and then using the results of the GETATTR to update the client's cached attributes. Including the change attribute in this set simplified the client's task of periodically validating these attributes.

#### 15.6.3. Attribute Caching Coherence Between Requesting Client and Remote Applications

The return of attributes modified by one client to applications running on other clients is often problematic because:

- \* The application or the operating environment might have been implemented without consideration of the effects of simultaneous updates by multiple clients.

This includes environments that were defined for local access and have no support for open deny modes.

- \* The applications might have been designed and implemented assuming that certain file attributes could be accessed without worry about remote access delays in accessing attributes.

Given the above situation, the applications will not have the same ability to validate attributes that the client does. This makes it important to prevent the incorporation of outdated value since:

- \* The client is not likely to be able to validate if they become outdated.
- \* Applications, unless special efforts are undertaken, are unlikely to be aware of the possibility of simultaneous access.
- \* The normal way to prevent such simultaneous access, the use of share reservations, might not be available, due to hard-to-address limitations of the operating environment.

The best alternative is likely to be elimination of attribute caching in such cases.

### 15.7. Data and Metadata Caching and Memory Mapped Files

Some operating environments include the capability for an application to map a file's content into the application's address space. Each time the application accesses a memory location that corresponds to a block that has not been loaded into the address space, a page fault occurs and the file is read (or if the block does not exist in the file, the block is allocated and then instantiated in the application's address space).

As long as each memory-mapped access to the file requires a page fault, the relevant attributes of the file that are used to detect access and modification (`time_access`, `time_metadata`, `time_modify`, and `change`) will be updated. However, in many operating environments, when page faults are not required, these attributes will not be updated on reads or updates to the file via memory access (regardless of whether the file is local or is accessed remotely). A client or server MAY fail to update attributes of a file that is being accessed via memory-mapped I/O. This has several implications:

- \* If there is an application on the server that has memory mapped a file that a client is also accessing, the client may not be able to get a consistent value of the change attribute to determine whether or not its cache is stale. A server that knows that the file is memory-mapped could always pessimistically return updated values for change so as to force the application to always get the most up-to-date data and metadata for the file. However, due to the negative performance implications of this, such behavior is OPTIONAL.
- \* If the memory-mapped file is not being modified on the server, and instead is just being read by an application via the memory-mapped interface, the client will not see an updated `time_access` attribute. However, in many operating environments, neither will any process running on the server. Thus, NFS clients are at no disadvantage with respect to local processes.
- \* If there is another client that is memory mapping the file, and if that client is holding an `OPEN_DELEGATE_WRITE` delegation, the same set of issues as discussed in the previous two bullet points apply. However, it should be noted that it is very unlikely that such a delegation will be held since it is normally required that the file be open for read to be mapped into memory. Only if the file were not open and accessed using a special stateid could the delegation be retained while the file in question is mapped into another client's memory. For this reason, such use is highly undesirable.

In this situation, when a server does a CB\_GETATTR to a file that the client has modified in its cache, the reply from CB\_GETATTR would not necessarily be accurate, assuming the delegation is not recalled at this point. As discussed earlier, the client's obligation is to report that the file has been modified since the delegation was granted, not whether it has been modified again between successive CB\_GETATTR calls, and the server MUST assume that any file the client has modified in cache has been modified again between successive CB\_GETATTR calls. Depending on the nature of the client's memory management system, it might not be possible to live up to this weak obligation. A client MAY return stale information in CB\_GETATTR whenever the file is memory-mapped, if another client is accessing the file without opening it.

## 16. Client-side Name and Directory Caching

Although there are important parallels, name and directory caching, discussed in this section as a whole, need to be considered separately from caching for files, discussed in Section 15 as a whole. This is because:

- \* Directories are structured objects with relatively infrequent entry changes effected on the server, the need for client-side data storage is reduced, eliminating one important barrier to effective coherence.
- \* The directory delegation feature, because of its notification features is in a better position to provide inter-client coherence, although achievement of consistency is likely to be delayed by transfer delays

Although directory notification include facilities for attribute updates, the delayed character of the notifications and the lack of directory-oriented constraints on changes implies that these are provided as a helpful convenience feature that leaves the caching architecture substantially unchanged.

### 16.1. Name and Directory Caching without Directory Delegations

The NFSv4.1 directory delegation facility (described in Section 16.2 below) is OPTIONAL for servers to implement. Even where it is implemented, it may not always be functional because of resource availability issues or other constraints. Thus, it is important to understand how name and directory caching are done in the absence of directory delegations. These topics are discussed in the next two subsections.

#### 16.1.1.1. Name Caching

The results of LOOKUP and READDIR operations may be cached to avoid the cost of subsequent LOOKUP operations. Just as in the case of attribute caching, inconsistencies may arise among the various client caches. To mitigate the effects of these inconsistencies and given the context of typical file system APIs, an upper time boundary is maintained for how long a client name cache entry can be kept without verifying that the entry has not been made invalid by a directory change operation performed by another client.

When a client is not making changes to a directory for which there exist name cache entries, the client needs to periodically fetch attributes for that directory to ensure that it is not being modified. After determining that no modification has occurred, the expiration time for the associated name cache entries can be updated to be the current time plus the name cache staleness bound.

When a client is making changes to a given directory, it needs to determine whether there have been changes made to the directory by other clients. It does this by using the change attribute as reported before and after the directory operation in the associated `change_info4` value returned for the operation. The server is able to communicate to the client whether the `change_info4` data is provided atomically with respect to the directory operation. If the change values are provided atomically, the client has a basis for determining, given proper care, whether other clients are modifying the directory in question.

The simplest way to enable the client to make this determination is for the client to serialize all changes made to a specific directory. When this is done, and the server provides before and after values of the change attribute atomically, the client can simply compare the after value of the change attribute from one operation on a directory with the before value on the subsequent operation modifying that directory. When these are equal, the client is assured that no other client is modifying the directory in question.

When such serialization is not used, and there may be multiple simultaneous outstanding operations modifying a single directory sent from a single client, making this sort of determination can be more complicated. If two such operations complete in a different order than they were actually performed, that might give an appearance consistent with modification being made by another client. Where this appears to happen, the client needs to await the completion of all such modifications that were started previously, to see if the outstanding before and after change numbers can be sorted into a chain such that the before value of one change number matches the after value of a previous one, in a chain consistent with this client being the only one modifying the directory.

In either of these cases, the client is able to determine whether the directory is being modified by another client. If the comparison indicates that the directory was updated by another client, the name cache associated with the modified directory is purged from the client. If the comparison indicates no modification, the name cache can be updated on the client to reflect the directory operation and the associated timeout can be extended. The post-operation change value needs to be saved as the basis for future change\_info4 comparisons.

As demonstrated by the scenario above, name caching requires that the client revalidate name cache data by inspecting the change attribute of a directory at the point when the name cache item was cached. This requires that the server update the change attribute for directories when the contents of the corresponding directory is modified. For a client to use the change\_info4 information appropriately and correctly, the server must report the pre- and post-operation change attribute values atomically. When the server is unable to report the before and after values atomically with respect to the directory operation, the server must indicate that fact in the change\_info4 return value. When the information is not atomically reported, the client should not assume that other clients have not changed the directory.

#### 16.1.2. Directory Caching

The results of READDIR operations may be used to avoid subsequent READDIR operations. Just as in the cases of attribute and name caching, inconsistencies may arise among the various client caches. To mitigate the effects of these inconsistencies, and given the context of typical file system APIs, the following rules should be followed:

- \* Cached READDIR information for a directory that is not obtained in a single READDIR operation must always be a consistent snapshot of directory contents. This is determined by using a GETATTR before the first READDIR and after the last READDIR that contributes to the cache.
- \* An upper time boundary is maintained to indicate the length of time a directory cache entry is considered valid before the client must revalidate the cached information.

The revalidation technique parallels that discussed in the case of name caching. When the client is not changing the directory in question, checking the change attribute of the directory with GETATTR is adequate. The lifetime of the cache entry can be extended at these checkpoints. When a client is modifying the directory, the client needs to use the change\_info4 data to determine whether there are other clients modifying the directory. If it is determined that no other client modifications are occurring, the client may update its directory cache to reflect its own changes.

As demonstrated previously, directory caching requires that the client revalidate directory cache data by inspecting the change attribute of a directory at the point when the directory was cached. This requires that the server update the change attribute for directories when the contents of the corresponding directory is modified. For a client to use the change\_info4 information appropriately and correctly, the server must report the pre- and post-operation change attribute values atomically. When the server is unable to report the before and after values atomically with respect to the directory operation, the server must indicate that fact in the change\_info4 return value. When the information is not atomically reported, the client should not assume that other clients have not changed the directory.

## 16.2. Directory Delegations and Notifications

### 16.2.1. Motivation for Directory Delegations

Directory caching for the NFSv4.1 protocol when directory delegations are not available, is similar to file and directory caching in previous versions. Clients typically cache directory information for a duration determined by the client. At the end of that predefined period, the client will query the server to see if the directory has been updated. By caching attributes, clients reduce the number of GETATTR calls made to the server to validate attributes. As a result, frequently accessed files and directories, such as the current working directory, have their attributes cached on the client so that some NFS operations can be performed without making an RPC

call. By caching name and attributes information about most recently looked up entries in a Directory Name Lookup Cache (DNLC), clients are able to avoid sending LOOKUP/GETATTR calls to the server every time such files are accessed.

This caching approach works reasonably well at reducing network traffic in many environments. However, it does not address environments where there are numerous queries for files that do not exist. In these cases of "misses", the client sends requests to the server in order to provide reasonable application semantics and promptly detect the creation of new directory entries. Examples of high miss activity are compilation in software development environments. The current behavior of NFS limits its potential scalability and wide-area sharing effectiveness in these types of environments.

Since, other distributed stateful file system architectures such as AFS and DFS have proven that adding state around directory contents can greatly reduce network traffic in high-miss environments, it is sensible to define and implement such facilities in NFSv4.1.

#### 16.2.2. Directory Caching Features

Delegation of directory contents is an OPTIONAL feature of NFSv4.1. Possession of a delegation can be taken advantage of in a number of ways:

- \* It can be used to provide a recallable assurance that the directory contents have not changed, allowing LOOKUP results (whether successful or not) and REaddir results to be cached, in order to enable these operations to be performed locally.

This mode of operation in which directory contents are fixed is often referred to as the "pure recall" model since any change in the directory contents results in the delegation being recalled. This mode of operation is most effectively used on large directories which are infrequently changed.

- \* The client can request, as part of requesting a delegation, that notifications be provided to update the clients view of the directory contents to match that of the server. See Section 16.2.11 for details. This mode of operation allows directory delegations to be effectively used in handling large directories that experience a significant stream of updates.
- \* Independently of the mode of operation selected, notifications to inform the client of attribute changes can be requested. See Section 16.2.11 for details.



### 16.2.3. Directory Delegation Notifications

As discussed below in Section 16.2.5, the caller had the ability to request various notifications in the form of callbacks (See Section 27.4

These callbacks aid the client by:

- \* Providing ways to update the client's view of the directory contents to match that of the server, instead of recalling the delegation.
- \* Provide Information about changes in directory attributes or the attributes of file objects named by directory entries.
- \* Provide assistance in authorizing LOOKUPS and READDIRS based on names cached by the client.

Each notification identifies the directory and delegation with which it is associated and contains an array of notification elements (each of type notify4).

Each notify4 is defined in a manner similar to that used for fattr4, in that the data is presented in XDR as consisting of a number of opaque elements, that are not actually opaque but are beyond the capacity of XDR to explicitly describe. In the case of the notify4 there are two nominally opaque arrays.

- \* The first is a bitmap4 which has an analogous function to the attrmask in a fattr4.

The bit positions within this mask are selected from the enum notify\_type4.

- \* The second nominally opaque element is the concatenation of selected elements corresponding to set bit in the bit mask described above.

This is analogous to attrlist4 within an fattr4.

The table in Section 27.4.1 indicates the data segment corresponding to each bit set.

#### 16.2.4. Directory Delegation Choices

Unlike other forms of delegation, the directory delegation feature has associated OPTIONAL features that allow clients to maintain correct local copies of directory contents even in the face of ongoing changes to the delegated directories. The client has three basic choices listed below in choosing a system of directory content change notifications. See Section 16.2.11 for details.

A major consideration for clients in making these choices concerns the operations to be performed locally using cached data provided by or assured correct by the directory delegation:

- \* Local LOOKUP is expected to be always used and is the core motivation for the feature.
- \* Local READDIR is an important possibility but is not universally provided because some clients and servers are not prepared to maintain client-side images of READDIR response and deal properly with the ordering of entries and the cookie values associated with them in order to keep the server and the clients fully in sync.
- \* Local GETATTR (and READDIR with attributes) is provided for in the design of directory delegation but is expected to require extra support facilities to be effective, as described in Section 16.2.8.

The need for these support facilities arises principally from the delayed/batched character of directory entry attribute notifications, while prompt notifications are hard to implement given the possibility that certain files will be linked to from multiple directories, making it difficult to promptly generate attribute change notifications. The usability of this feature without prompt notification depends on the need for up-to-date values of various file attributes, which are not currently standardized.

- \* To provide for delegation loss as a result of any directory change.

Delegation recall results from changes to directory contents if directory content notifications are not requested.

Because recalls halt directory changes until the delegation is returned, it is often advantageous to request change notifications (asynchronous) and respond to them by returning the delegation.

- \* To receive notifications about changes to directory contents, considered as an unordered set of names and update the client's view of the directory in accord with the notification.

This can be used to support both positive and negative name caching. However, it does not fully support retention of an accurate client copy in the face of changes because of lack of support for tracking ordering changes and entry cookie changes.

- \* To receive notifications about changes to directory contents, considered as an ordered set with associated identifying cookies.

Because of server-side difficulties in providing some of this information and client-side difficulties providing it, this choice might not be available, but, if so, the previous option will support positive and negative name caching even if repeated READDIRS cannot be eliminated.

In addition there are possible additional forms of notifications to aid the client:

- \* Notifications regarding changes in directory attributes .

Knowledge of such authorization-related attributes can be helpful to the client in validating the authorization to use cached name entries or directories on its own instead of using ACCESS requests.

In some cases, additional facilities are available to simplify the task of meeting authorization requirements. See Section 16.2.6 for details.

Knowledge of modified-time or the change attribute can be helpful in potentially keeping a valid view of the directory when dealing with the consequences of delegation recall. See Section 16.2.11 to see how attribute notifications can be used together with content notifications to associate a set of directory contents and the corresponding directory attributes.

- \* Notification of changes of attributes of objects within the directory.

#### 16.2.5. Directory Delegation Mechanics

The GET\_DIR\_DELEGATION (Section 25.39) operation is used by clients to request directory delegation. The delegation is read-only and the client is not provided any means to make changes to the directory other than by performing NFSv4.1 operations that modify the directory.

As part of obtaining a delegation, the client specifies, using the bit numbers within the notify\_type4 enum that appears below, its choices regarding notification of events related to the reporting of events affecting the delegation. These bits originally specified a set of notification types but have been extended into a wider set of flag values specified when delegations are requested and returned when they are granted:

- \* Some bits request that particular notifications be provided to the client, instead of recalling the delegation.

These bits are used to indicate that the requested notifications will be provided by the server.

- \* Other bits have no associated notification message and provide request flags specifying the requested handling of the delegation and/or response flags indicating useful information about the handling of the delegation and associated notifications.

It is important to note that this enum is subject to extension and has been extended relative to the set of bits defined in [RFC8881]. The distinction between bits that were defined earlier and those added later is important to enable interoperation between clients and servers when one might have been written based on the earlier specification. Although no implementations based on the earlier specification are known, the possibility of their existence cannot be excluded.

```
/*
 * Directory notification types and associated flags
 */
enum notify_type4 {
    /*
     * Present in RFCs 5661, 8881
     */
    NOTIFY4_CHANGE_CHILD_ATTRS = 0,
    NOTIFY4_CHANGE_DIR_ATTRS = 1,
    NOTIFY4_REMOVE_ENTRY = 2,
    NOTIFY4_ADD_ENTRY = 3,
    NOTIFY4_RENAME_ENTRY = 4,
    NOTIFY4_CHANGE_COOKIE_VERIFIER = 5,
    /*
     * Added in NFSv4.1 bis document
     */
    NOTIFY4_GFLAG_EXTEND = 6,
    NOTIFY4_AUFLAG_VALID = 7,
    NOTIFY4_AUFLAG_USER = 8,
    NOTIFY4_AUFLAG_GROUP = 9,
    NOTIFY4_AUFLAG_OTHER = 10,
    NOTIFY4_CHANGE_AUTH = 11,
    NOTIFY4_CFLAG_ORDER = 12,
    NOTIFY4_AUFLAG_GANOW = 13,
    NOTIFY4_AUFLAG_GALATER = 14,
    NOTIFY4_CHANGE_GA = 15,
    NOTIFY4_CHANGE_AMASK = 16,
    NOTIFY4_ACCFLAG_USER = 17,
    NOTIFY4_ACCFLAG_GROUP = 18,
    NOTIFY4_ACCFLAG_OTHER = 19,
    NOTIFY4_PRAN_CONTENT = 20,
    NOTIFY4_PRAN_AUTH = 21
};
```

The following table summarizes the handling of the bit numbers specified in the enum above:

The bits within the enum are organized into the following groups:

- \* Content change notifications which provide updates regarding the set of directory entries (or sometimes the order of those entries).
- \* General flags which affect the process of other in bits in multiple groups.

- \* Attribute notifications which provide updates regarding attributes of directories objects identified by directory entries.
- \* Lookup authorization infrastructure which provides the client assistance in avoiding repeated use of ACCESS in satisfying requests using name caching.
- \* Getattr authorization infrastructure which provides the client assistance in avoiding repeated use of ACCESS in satisfying requests for cached file attributes.

Such help is made necessary by the existence of possible restrictions on attribute fetches in the NFSv4 ACL model.

The group has a reference to an appropriate explanatory section and is followed by the bits included within the group together with, for each bit:

- \* The abbreviated bit name.
- \* The bit number.
- \* The bit's extension status with "old" indicating something known to all client and defined in [RFC8881] and "Ext" indicating an extension added as part of the respecification effort, to correct a defect.
- \* The bit type using one of the codes listed below.
- \* A reference to one or more sections with more detailed description.

The following types of notifications and flags are defined.

PrN Prompt Notification -- One sent immediately as part of the event being reported.

DlN Delayed Notify -- One sent lazily at some later time based on a associated attribute.

SmN Specially managed notification -- One sent promptly or in delayed fashion based on factors explained in the detailed reference.

Bdf Bidirectional flag -- One specified on delegation request and interrogated in the corresponding response.

RgF Request Flag -- One used as part of the delegation request with

no corresponding value in the response.

RsF Response Flag -- One having no use as part of the delegation request while conveying useful information in the response.

Group	Discussion				Ref
Item	#	Ext	Type	Description	
Content	Content Change Notifies				S. 16.2.11
REMOVE_ENTRY	2	Old	PrN	Entry Removal Ntfy	S. 27.4.5
ADD_ENTRY	3	Old	PrN	Entry Addition Ntfy	S. 27.4.4
RENAME_ENTRY	4	Old	PrN	Entry Rename Ntfy	S. 27.4.6
CHANGE_COOKIE_VERIFIER	5	Old	PrN	Change Cookie Verifier	S. 27.4.9
Gen Flags	General Flags				S. 16.2.10
GFLAG_EXTEND	6	Ext	BdF	General Extension Flag	S. 16.2.10
Content Flags	Content Description Flags				S. 16.2.11.3
CFLAG_ORDER	12	Ext	RqF	Controls Content Ordering Needs	S. 16.2.11.3
Attribute	Attribute Change Notifies				S. 16.2.12
CHANGE_CHILD_ATTRS	0	Old	DlN	Changes to Dir Entry Attributes	S. 16.2.12.2
CHANGE_DIR_ATTRS	1	Old	SmN	Changes to	S.

				Dir Attributes	16.2.12.1
CHANGE_AMASK	16	Ext	PrN	Change to Attribute Masks	S. 27.4.10
PRAN_CONTENT	20	Ext	RsF	Prompt Not. for Attr. Changes with Content Changes	S. 16.2.12.1
PRAN_AUTH	21	Ext	RsF	Prompt Not. for Attr. Changes re Authorization	S. 16.2.12.1
Lookup Authorization	Lookup/readdir Authorization Support				S. 16.2.13.1
AUFLAG_VALID	7	Ext	RsF	Validity of Authorization Flags	S. 16.2.13.1
AUFLAG_OWNER	8	Ext	RsF	Authorization for Owner	S. 16.2.13.1
AUFLAG_GROUP	9	Ext	RsF	Authorization for Owning Group	S. 16.2.13.1
AUFLAG_OTHERS	10	Ext	RsF	Authorization for Others	S. 16.2.13.1
ACCFLAG_OWNER	17	Ext	RsF	Need for ACCESS for Owner	S. 16.2.13.1
ACCFLAG_GROUP	18	Ext	RsF	Need for ACCESS for Owning Group	S. 16.2.13.1
ACCFLAG_OTHERS	19	Ext	RsF	Need for ACCESS for Others	S. 16.2.13.1
CHANGE_AUTH	11	Ext	PrN	Changes to	S.



				Lookup Authorization	27.4.11
Getattr Authorization	Getattr Authorization Support				S. 16.2.13.2
AUFLAG_GANOW	13	Ext	RsF	Getattr processing now	S. 16.2.13.2
AUFLAG_GALATER	14	Ext	RsF	Getattr processing deferred	S. 16.2.13.2
CHANGE_GA	15	Ext	PrN	Getattr processing change	S. 27.4.12

Table 6

The holder of a delegation is assured of certain things not being changed while the directory delegation is held, as described below.

- \* That the set of entries within the directory not be changed without sending a requested notification to the client, informing the client of the change.
- \* That the order of directory entries or the cookie values associated with specific directory entry with the client being informed (via a NOTIFY4\_CHANGE\_COOKIE\_VERIFIER notification) of the possibility of change.

Delegations can be recalled by the server at any time and are always recalled before a directory is removed.

#### 16.2.6. Directory Delegation Authorization Extensions

When cached data is used locally in place of LOOKUP, GETATTR, or, READDIR operations, the authorization constraints that would normally be imposed by the server have to be applied by the client instead of by the server. This is because requiring server-based authorization would undercut the performance benefits that caching is intended to provide.

The discussion of these matters is complicated by the fact that, while facilities have been designed to avoid the overhead requirements of authorization and these are described in this

document, the treatment of the matter in earlier specifications did not provide facilities to help the client do this correctly and had little to say on the issue. As a result, clients were faced with the choice of:

- \* ignoring some difficult authorization issues.
- \* implicitly deciding that certain features of the NFSv4 ACL Model are not to be implemented within filesystems for which directory delegation is to be used.
- \* burdening the implementation of directory delegation with authorization checking that would undercut the performance benefits of the feature, when used with filesystems that implement NFSv4 ACLs, especially where the implementation includes certain troublesome features.

Of particular concern are ACLs for which AUDIT or ALARM ACEs are implemented, making it possible that there are directory ACLs where authorization checking cannot be accomplished by the client even if the client is willing and able to process ACLs itself in order to make authorization decisions.

In considering whether the extensions described in Section 16.2.9 are worth implementing within client and server implementation of directory delegation, the following factors need to be considered:

- \* Whether the cached directories are used to prevent remote access when implementing READDIR and GETATTR, or only when avoiding LOOKUPS.
- \* Whether ACLs are implemented making the decision of search permissions for a given user more complicated than it is when this decision is based solely on the mode, owner, and owner-group attributes derived from the POSIX authorization model.
- \* The support for and use of the ACE mask bit controlling authorization to read attribute values.
- \* The preparedness of the client implementer to fetch and interpret directory ACLs and use them to derive search/ read authorization. Beyond the reluctance of client implementers to take on this new task, especially given the uncertain state of the specification of ACL semantics, an important consideration relates to the removal of previous recommendations to avoid this practice. In particular, the direction in [RFC8881] stating that this "SHOULD NOT" be done, will no longer appear in the NFSv4.1bis documents,

As dealt with in this document and related ones, the determination of search/read authorization for directories is not considered a troublesome usage to be discouraged. However, efforts will be made in the definition of the features defined in Section 16.2.9 to accommodate clients that do not have this capability.

When certain combinations of the above issues are present, we will be faced with a situation in which we need to decide how to deal with authorization mistakes that implementers might have been encouraged to make by previous specs, now considered superseded. See Section 16.2.7 for a discussion of these matters. Note that while this could be considered analogous to the situation in which implementers were told that use of AUTH\_SYS in the clear, is an "OPTIONAL means of authentication", with the consequence that such use needed to be allowed in the future while being cautioned against, our situation is different in that:

- \* This disregard of authorization issues was implicit and there was no statement that they could be ignored.
- \* We are dealing with a feature that is unlikely to have troublesome legacy implementations.
- \* While the previous use of "SHOULD NOT in connections with client processing of ACLs might have been relied on, there is unlikely to be specific implementations embodying this approach, so that our accommodation of implementors not doing this is motivated by providing an easier implementation path, rather than accommodating a previous mistaken specification.

#### 16.2.7. Former Authorization Practices and Their Current Validity

There are a number of troubling authorization issues that arise because of features that were included as OPTIONAL extensions to the NFSv4 ACL model.

- \* The potential use of AUDIT and ALARM ACEs, which make the requirements for a client-side authorization check more complicated than a single, cacheable, yes-no decision, making it possible that server-specific implementation is necessary in the event of successful or failed authorization,
- \* The appearance of ACE's that include the ACE mask bit controlling authorization to read attributes of a specified file.

This feature is unlikely to be supported by servers because it is outside the scope of the POSIX authorization model. Nevertheless, there are likely to be servers supporting it, especially if it is needed by Windows clients.

This feature is not troublesome if the client implementations of directory delegation do not support client-local provision of GETATTR or READDIR with associated attributes.

The existence of these features together with the following issues means that some implementors might need the server-side assistance described in Section 16.2.9, even when a client could do authorization checking on its own.

- \* The potential transfer of NFSv3-based code (which could not deal with ACLs to an NFSv4 environment and its subsequent incorporation in a client-side directory delegation implementation.
- \* The lack of attention to authorization issue in previous specifications combined with the sense that a client doing its own authorization checks was to be discouraged.
- \* The preparedness of the client implementer to fetch and interpret directory ACLs and use them to derive search authorization. Beyond the reluctance of a client implementer to take on this new task, especially given the uncertain state of the specification of ACL semantics, the direction in [RFC8881] stating that this "SHOULD NOT" be done, even though, in the NFSv4.1 bis documents, this statement will no longer appear

When, for whatever reason, a client is unwilling to provide this support, implementation of directory delegations is not practical unless the facilities described in Section 16.2.9

As a result, we are faced with the issue of how to accommodate implementations that are now known to have troubling problems that were not recognized when the feature was first described in a Proposed Standard. Normally, one tries to accommodate such situations by recommending against approaches now known to be flawed while considering, as a valid reason to bypass the recommendation, the reliance of the implementer on an approved Proposed Standard at the time. In this case we have a different approach, because of the following distinctive factors:

- \* Unlike the case of implementers being told that use of AUTH\_SYS in the clear, is an "OPTIONAL means of authentication" with the implication that such use does not result in potentially unacceptable security vulnerabilities, there is no direct

suggestion that neglecting these difficulties is acceptable. Instead, while lack of attention to security issues might have led people astray, they were not specifically asked to adopt a flawed approach to security but chose to adopt one on their own. As a result, while we will make certain allowances to accommodate such early implementations, there is no known paradigm that could be cited as valid but discouraged.

- \* The set of implementations involved is likely to be quite small and might be empty or only consist of experimental implementations not widely distributed.

The approach we take here is the same one we take to servers that do not support the extensions described in Section 16.2.13 and to clients that interact with such servers. It has the following elements:

- \* Clients are free in deciding whether to use directory delegations to take account of the problems with earlier approaches in deciding whether to use this feature.
- \* Neglecting the possibility of authorization failure on GETATTR when directory entry attributes are cached is not to be considered valid despite the development of code that ignored this issue based on previous specifications together with the widely held impression that authorization of GETATTR is not required.

If the server implements this feature, clients, if they implement this directory delegation have no basis to justify undercutting it.

- \* When clients use ACCESS to do authorization checks, as they should when they have no ability to process ACLs on their own, allowance needs to be made for them to cache positive and negative results, since without that ability, you might as well fetch the data over the wire anyway, negating the value of directory delegation.

When AUDIT and ALARM ACEs are present, there will be situations in which local processing of the ACL is not sufficient to avoid remote use of ACCESS.

#### 16.2.8. Alternatives to Use of Directory Delegation Authorization Support

The need for implementations of the extensions described in Section 16.2.9 is affected by the following considerations:

- \* The willingness/ability of the client to do authorization checking on its own, including the ability to read and interpret ACLs fetched from the directory or reported via attribute change notifications.
- \* The frequency of encountering ACLs for which local checking is not possible.
- \* The ability of the facilities within Section 16.2.9 to avoid otherwise necessary remote operations.

We will deal separately with the cases of clients prepared or not prepared to use ACLs locally.

For clients that are not able to process ACLs locally, there are a number of features that server filesystems might need support to help clients deal with situations they cannot effectively deal with in supporting directory delegations. Clients need to either suppress directory delegation when these situation exist, or rely on the assistance provided by the authorization support extensions described in the next session.

- \* Support of AUDIT and ALARM ACEs. The ACLSUPPORT attribute for a filesystem can be checked to determine if this feature is present.

When it is present and the client is not able to process ACLs, each LOOKUP would require a remote ACCESS request, effectively negating the performance benefits of directory delegation.

- \* When a server and client work together to support local use of cached attribute data and the filesystem in question supports the ACE mask bit ACE4\_READ\_ATTRIBUTES. Each local execution of GETATTR, using cached data when clients support this, requires authorization that the client, if it cannot process ACLs, is not able to do.

The ACLCHOICE attribute, when it is supported, allows the client to inhibit support for local GETATTR operations using cached data for filesystems that support suppression of the authorization of GETATTR. When ACLCHOICE is not supported or this suppression is possible support for GETATTR cannot be done.

In situations in which these restrictions are not acceptable, the authorization support features can help the client by allowing it to distinguish the directories (expected to be common) where these features do not occur. For example, even when AUDIT and ALARM ACEs are present within a filesystem, they are unlikely to be associated with permission to search directories.

For clients who can process ACLs on their own the situation is different since the client can check for these troublesome situations on its own. Nevertheless, the authorization support facilities, while not necessarily required can optimize the process by avoiding ACL processing in situations in which it can be avoided for a common class of requests.

One important class of authorization requests consists of repeated requests for authorization of the same user, assuming that the authorization returned will typically not change when successive requests to the same directory are made. While the client can deal with this situation on its own by flushing the associated authorization cache when the directory's metadata is changed, the authorization support facilities deal with the issue more completely because:

- \* Changes that do not affect permissions to do LOOKUP (and READDIR where local READDIR is supported) need not flush cached indications of authorization success or failure.
- \* When there are such changes, they can be limited in their effects since certain commonly used classes of users often are not changed by common sort of ACL changes.

While the client could do that analysis on its own, it is preferable to rely on the server since the server needs to support these facilities to allow effective use of directory delegations.

For clients that intend to support local GETATTR, the need for additional support facilities is not affected by the ability of the client to read and interpret ACLs on its own. Given the use of delayed/batched entry attribute change notifications, the important issues concern whether the attributes returned include those that require prompt notification when updated.

#### 16.2.9. Directory Delegation Authorization Support

In providing support for authorization of local operations effecting, using cached data, the equivalents of LOOKUP and READDIR operations, the following issue needs to be dealt with:

The possibility of change in authorization-related attributes would make repeated ACCESS calls necessary, unless facilities are provided to avoid these when possible.

Note that the same issue applies to authorization of GETATTR-equivalent local operations, but that, in that case, there are the following additional issues to deal with:

- \* The only potential reason to not grant such access derives from the possible use of the ACE mask bit `ACE4_READ_ATTRIBUTES`.

Only where that bit is supported in ACLs and used to either deny access or require audit or alarm on this operation is there any possibility of not letting this operation be done unconditionally.

- \* Since permissions would need to be checked for each individual object rather than for the directory as a whole, it is harder to avoid unnecessary `ACCESS` calls in situations where the possibility of denial exists.

The possible existence of multiply-linked file adds further difficulty since it is possible that an ACL could be changed for such an object in case in which the affected directories might not be known.

- \* There are very few ACL implementations supporting use of the ACE mask bit `ACE4_READ_ATTRIBUTES` and no known uses of it.

As a result we have to be prepared to efficiently deal with the simple case where sophisticated support is unnecessary, as well as providing reasonable support to deal with the possibility of it becoming more widely implemented.

To provide better support for authorization of `LOOKUP/READDIR`, we do the following:

- \* When the delegation is created, the server returns information about pre-defined sets of users for which explicit authorization checks can be avoided or for which client-side authorization cannot validly be done. These sets are a natural division of the principal space with POSIX authorization semantics, and remain an appropriate division even though certain other special user groups might be defined and used within ACLs.

The flags `NOTIFY4_AUFLAG_OWNER`, `NOTIFY4_AUFLAG_GROUP`, and `NOTIFY4_AUFLAG_OTHERS` indicate the ability to avoid authorization checks for `LOOKUP` and `READDIR` by the owner of the file, other members of the owning group, and others, respectively. The presence of each of these flags indicate that all principals in the corresponding group can be assumed authorized to perform `LOOKUPS`, and, when `NOTIFY4_CFLAG_ORDER` is set, `READDIRS` as well. When multiple users are included in a given group, if some could fail authorization, the corresponding bit will not be set and authorization needs to be checked using `ACCESS`, by the client doing its own authorization check, or by caching the results of a previous check.



The flags NOTIFY4\_ACCFLAG\_OWNER, NOTIFY4\_ACCFLAG\_GROUP, and NOTIFY4\_ACCFLAG\_OTHERS indicate the need to perform an ACCESS check when the principal performing the operation (LOOKUP and possibly READDIR as in the previous paragraph) is within the corresponding group. The need for such checks, implying a client-side check is not acceptable, arises from the possible inclusion of AUDIT and ALARM ACEs which might require server-side processing of the authorization request in either the successful or unsuccessful authorization cases.

- \* When there is a change in one or more of the directory's authorization-related attributes, the client is notified of the new authorization handling scheme using the NOTIFY4\_CHANGE\_AUTH notification.

The notification provides changes that apply separately to the owning user, other users in the owning group, and other users. For each such group, there are separate bits controlling the need for explicit ACCESS checks for LOOKUP and for READDIR, and directing the client whether to flush cached results for previous ACCESS checks.

To provide adequate support for authorization of local GETATTR and associated attribute caching the facility described in Sections 27.4.12 and 16.2.13.2.

#### 16.2.10. Directory Delegation Feature Version Management

As part work undertaken to respecify NFSv4 minor version one to reflect implementation experience since the publication of [RFC5661], it was necessary to make certain protocol extensions in order to correct problems that had resulted in a lack of implementation of the Directory Delegation feature in the years since its initial introduction.

These extensions took the form of additions to the enum notify\_type4 as described in [RFC8178]. These new values,

- \* Provide new notifications including a set focused on providing authorization support to allow operations to be performed locally without impacting needed authorization semantics.

Provided a new notification to allow server to deal with excessive backchannel traffic for attribute updates without delegation recall.

- \* Created flags to be sent by the client as part of delegation request and by the server as part of delegation creation.

These flags allowed necessary version control, improved authorization handling and a more flexible approach to the provision of position information in content update notifications.

These new notifications and flags are described, together with the older ones, in Sections 16.2.11 through 16.2.13

The flag NOTIFY4\_GFLAG\_EXTEND has a special role in the management of versions, in order to support interoperation of implementations written to conform to [RFC8881] and to the those written to conform to the updated definition:

- \* When NOTIFY4\_GFLAG\_EXTEND is set in a request, the client is indicating that it is aware of the additional flags and notifications.
- \* When NOTIFY4\_GFLAG\_EXTEND is set in a response, the server is indicating that it is aware of the additional flags and notifications. and that the delegation is to be handled in accord with the updated specification of the feature.

#### 16.2.11. Directory Content Notifications

When a directory delegation is held, it is normally expected to remain unmodified. However, to avoid the possibility that directory content changes will force much of the work done in establishing a delegation to be redone, the user can request notifications be sent regarding certain changes, allowing the client to update its view of the directory, while keeping the delegation unrecalled.

Support is present for clients who are interested in the order of directory entries and for those that are not. While the former will have the ability to cache and reuse READDIR responses and synthesize new ones in response to content modification notifications, the latter only need the ability to maintain positive and negative name caches for entries in the directory. The need for order-related information is specified using the flag NOTIFY4\_CFLAG\_ORDER (only specifiable when NOTIFY4\_GFLAG\_EXTEND is specified as well).

##### 16.2.11.1. Core Content Notifications

Notification of directory content changes is requested by requesting any of the notification types NOTIFY4\_ADD\_ENTRY, NOTIFY4\_REMOVE\_ENTRY, NOTIFY4\_RENAME\_ENTRY, although typically a client would request all of these to avoid situations where common directory changes result in directory recall.

These notifications are used as follows. For details regarding the specifics of the relevant notification messages, see the appropriate subsection of Section 27.4.

- \* The creation of a new directory entry, as a result of an OPEN creating a new file, a CREATE operation, or a LINK operation is signaled using a NOTIFY4\_ADD\_ENTRY notification.

It is described in Section 27.4.4.

- \* The deletion of an existing directory entry as a result of a REMOVE operation, or the deletion of a file renamed-over by a cross-directory RENAME operation is signaled using a NOTIFY4\_REMOVE\_ENTRY notification.

It is described in Section 27.4.5.

In the cross-directory rename-over case, the NOTIFY4\_ADD\_ENTRY and NOTIFY4\_REMOVE\_ENTRY notification are to appear in the same notify4, with both notification bits set.

- \* The renaming of a directory entry and the possible deletion of a renamed-over entry as part of a within-directory RENAME operation is signaled by a NOTIFY4\_RENAME-ENTRY notification.

It is described in Section 27.4.6.

#### 16.2.11.2. Order and Format Choices

Based on the client's preference regarding order information and the server's choices, there are three forms of content notification that can be sent:

- \* The order-agnostic format is used when the client is not concerned with directory entry order, as indicated by the NOTIFY4\_CFLAG\_ORDER being zero when the NOTIFY4\_GFLAG\_EXTEND flag is set.

This form is requested by the client if it does not want to maintain an ordered image of the directory locally and is unconcerned with directory cookies.

When this format is requested by a client the notification is not sent to the client that requested the changes, since it is presumed to be aware of the change, and has no need for the additional ordering-related information to be provided.

If the server does not support this option, additional information can be provided using one of the other forms, with unneeded information ignored by the client.

- \* The derived-order format provides a directory entry ordering based on the cookies associated with the directory entries, it being assumed by the client that the order of cookies and the order of directory entries being the same.

A `nad_new_entry_cookie` of length 1 with the new cookie in it provides the necessary position for added files. Any `nad_prev_entry` is of length 1 with an invalid `notify_entry4` indicated by the `ne_file` component being 0 of length 0. The `nad_last_entry` value conveys and should be ignored.

The presence of a `nad_new_entry` not in this form indicates that, if the client is order-aware, the notification is to be treated as in the explicit-order format.

If it is ever the case that the cookie order and the entry order for a directory becomes different, then this format cannot be used and the change should be signaled using `NOTIFY4_CHANGE_COOKIE_VERIFIER` (See Section 16.2.11.3).

- \* The explicit-order format fully uses and sets cookie values `nad_prev_entry` field and `nad_last` entry flags

The above rules change the actual form of the notifications. even though the same XDR definitions are used for the structure associated with the `ADD_ENTRY`, `REMOVE_ENTRY`, and `RENAME_ENTRY` notifications

#### 16.2.11.3. Order-related Content Change Notifications

Notifications of type `NOTIFY4_CHANGE_COOKIE_VERIFIER` are used to signal a number of changes in the order of directory entries and their associated cookies as listed below. The notification are only sent if the client is concerned with directory entry order. If the client is concerned with entry order and these notifications are not requested or cannot be sent for any other reason, then the delegation is recalled.

The changes requiring these notifications include, in addition to cookie verifier changes, any of the following:

- \* Changes in cookies for cached entries.
- \* Changes in directory entry order that do not arise in connection with the content changes discussed in Section 16.2.11.1.

This notification is described in Section 27.4.9.

It is often sent with the old and new verifiers being the same.

#### 16.2.11.4. Associated Directory Attribute Changes"

There a number of attributes that, by their nature, are changed whenever the directory contents is changed. These attributes are `modified_time` and `change`.

When change notification for any of these attributes are requested, they are delivered promptly, without regard to the specification of a delay time in the delegation request.

The notifications appear in the same `notify4` providing information about the corresponding contents change. These include:

- \* Changes signaled by `NOTIFY4_ADD_ENTRY`, `NOTIFY4_REMOVE_ENTRY`, or `NOTIFY4_RENAME_ENTRY` notifications.
- \* Changes due to cross-directory `RENAME` signaled using a combination of `NOTIFY4_ADD_ENTRY` and `NOTIFY4_REMOVE_ENTRY` notifications.
- \* Ordering or cookie change signaled by a `NOTIFY4_ADD_CHANGE_VERIFIER` notification.

#### 16.2.11.5. Possible Recall Due to Directory Changes"

The implementation sections for a number of operations describe situations in which notification or delegation recall would be required under some common circumstances. When these events result in delegation recall, a set of caveats similar to those listed in Section 15.2 apply. Note that in these cases, the operation does not wait for the delegation to be returned or revoked, as it does in other cases of delegation recall.

- \* For `CREATE`, see Section 25.4.4.
- \* For `LINK`, see Section 25.9.4.
- \* For `OPEN`, see Section 25.16.4.
- \* For `REMOVE`, see Section 25.25.4.
- \* For `RENAME`, see Section 25.26.4.
- \* For `SETATTR`, see Section 25.30.4.

#### 16.2.12. Directory Attribute Notifications

Those holding directory delegations may be notified of attribute changes as described below:

- \* Information about changes to directory attributes are provided via the NOTIFY4\_CHANGE\_DIR\_ATTRS notification, as discussed in Section 16.2.12.1.

Most often these changes are provided using a Batch/Deferred approach, although prompt notification can be requested by specifying a delay of zero.

Special prompt handling is available for attributes whose values are inherently tied to directory content changes. See Section 16.2.11.4 for details.

- \* Information about changes to attribute of file objects designated by directory entry are provided by the NOTIFY4\_CHANGE\_CHILD\_ATTRS notification, as discussed in Section 16.2.12.2.
- \* When the server finds it unduly burdensome to provide the notifications requested above, it has the option of recalling the delegation but also can, if requested, effect a change in attributes reported when using a notification.

NOTIFY4\_CHANGE\_AMASK is used to provide updated sets of masks for the attribute updates being provided. This notification is described in Section 27.4.10

These notifications, while asynchronous, are not subject to delay or batching.

This form of notification can be used by the server to reduce or eliminate child attribute notifications, without delegation recall.

##### 16.2.12.1. Directory Attribute Change Notifications

When attributes change for a directory for which a directory delegation is held, the client is notified by a NOTIFY4\_CHANGE\_DIR\_ATTRS notification as described in Section 27.4.7.

When these notifications are not requested, the client is not notified and the delegation is not recalled.

The delivery of such notifications generally uses a batch/delay paradigm, although prompt notification can be requested by specifying a `gdda_dir_attr_delay` value of zero (allowed if the `dir_notif_delay` is zero as well).

These prompt notifications are expected to be available in the following situations:

- \* Notifications for attributes that are expected to change as part of directory content modifications are always presented promptly, if the above condition are met. See Section 16.2.11.4 for details.

This information can be useful by clients that try to cache complete directories to allow provision of a local `READDIR` response for APIs which expect it to be present (e.g., as a directory entry for `"."`).

- \* Notifications for changes of directory attributes that affect the authorization of `LOOKUP` and `READDIR` will be presented promptly when the attributes for which changes are requested and generated are limited to `mode`, `owner`, `owner_group` and `ACL`.

The prompt reporting of changes to these attributes allows clients prepared to do the authorization checking locally to do so reliably, even when the directory attributes are subject to change. See Section 16.2.8 for further discussion.

Whether prompt notification for the abovementioned changes is provided is affected by the attributes for which change is requested and by the delay requested and in effect for attribute notifications. When such notifications are requested and in effect, the server signals that fact by the use of response flags:

- \* `PRAN_CONTENT`, when set, indicate that prompt notification will be used in cases in which attribute changes are an inherent consequence of directory content changes.
- \* `PRAN_AUTH`, when set, indicate that prompt notification will be used in cases in which attribute changes are affecting the attributes associated with authorizing `LOOKUP` and `READDIR` (i.e., `mode`, `owner`, `owner_group`, and `ACL`).

While it might be supposed that prompt notification of changes in authorization-related attributes could be used in eliminating the need for a remote invocation of authorization via use of `ACCESS`, there are cases in which such server-based checks are needed (e.g. due to the presence of `AUDIT` and `ALARM` ACEs).

The use of this approach requires that the client be able to fetch and interpret ACLs locally, a practice formerly discouraged.

#### 16.2.12.2. Entry Attribute Change Notifications

When attributes change for file objects in a directory for which a directory delegation is held, the client is notified by a NOTIFY4\_CHANGE\_DIR\_ATTRS notification as described in Section 27.4.8.

These notification are subject to delay and batching so as to provide reasonably up-to-date attribute caches without excessive network traffic. While prompt delivery can be requested by specifying a `gdda_child_attr_delay` value of zero, server are unlikely to be able to provide this because of the difficulty of find all the delegated directories associated with a file whose attributes are being modified.

#### 16.2.13. Directory Delegation Authorization-related Information

The flag-based indications and notification types discussed below are used to inform the of the proper approach to use in authorizing LOOKUP, READDIR, and GETATTR request to be satisfied using cached data obtained via directory delegation and associated notifications.

Because of the greater complexity of authorization resulting from the addition of various form of ACLs, it is rarely practical for the client to make the authorization itself while the prospect of using an ACCESS request (remote) could have the effect of vitiating the benefits of caching that the directory delegations feature is intended to avoid.

We discuss authorization for these operations below:

- \* Authorization support for LOOKUP is discussed in Section 16.2.13.1.

In the case in which the client is order-aware (NOTIFY4\_CFLAG\_ORDER set or NOTIFY4\_GFLAG\_EXTEND reset), this section also provides authorization support for READDIR satisfied from a local cache.

- \* Authorization support for GETATTR is discussed in Section 16.2.13.2.

Because GETATTR is always authorized unless the server supports NFSv4 ACLs and also supports a separate permission bit for reading attributes these feature might not be supported, requiring the client to use ACCESS when satisfying GETATTR requests locally.



#### 16.2.13.1. Authorization Support for LOOKUP/READDIR

Initial control of LOOKUP authorization is controlled by the setting of the bit NOTIFY4\_AUFLAG\_VALID. If the client is order aware, this bit controls the handling of READDIR authorization as well.

If bit is set, the flags NOTIFY4\_AUFLAG\_OWNER, NOTIFY4\_AUFLAG\_GROUP, and NOTIFY4\_AUFLAG\_OTHERS allow LOOKUP to proceed without further authorization for the owner of the directory, the owning group, or other users, respectively. If the client is order aware, these bits allow the handling of READDIR without additional authorization.

If the bit is not set, authorization of each user/operation pair must be done (using ACCESS) and the results can be cached to avoid repeated ACCESS calls. The applies cases in which the overall bit was set but use of ACCESS was required for a particular user.

This situation continues until a NOTIFY4\_CHANGE\_AUTH notification is received. This notification is described in Section 27.4.11.

This notification is used to modify authorization handling for LOOKUP (and READDIR if the client is order-aware) for the following reasons:

- \* There is a need to treat authorization of LOOKUP and READIR differently for some users.
- \* There is a changes to the directory's authorization-related attributes.

The notification contains three fields modifying authorization handling.

- \* There is a three-bit field each whose its allow use of LOOKUP without additional authorization. The bits concern use by the directory owner, owning group, and others.
- \* There is another three-bit field each whose its allow use READIR without additional authorization. The bits concern use by the directory owner, owning group, and others.
- \* There are three bits, each of which, when set, requires the deletion of cached entries allowing or denying authorization of LOOKUP or READIR for a particular user. The bits control the directory owner, all users in the owning group, and others.

## 16.2.13.2. Authorization Support for GETTATTR

When a directory delegation is granted, the client uses the flags returned to establish an initial authorization state with regard to GETATTR as follows:

- \* If the flag NOTIFY4\_AUFLAG\_GANOW is set, the client is being told that GETATTRs can now be done without explicit ACCESS checks. This status continues until changed by a possible NOTIFY4\_CHANGE\_GA notification.

The server can do this if ACE4\_READ\_ATTRIBUTES is not supported and also if it has scanned the directory to make sure that no current ACEs use that mask and that there are no multiple-linked files that make it possible that such ACEs will be set without the directory delegation holder being notified.

- \* Otherwise, if the flag NOTIFY4\_AUFLAG\_GALATER is set, the client is being told that GETATTRs now require explicit ACCESS checks, but that the situation is expected to change and it will be notified of that using a NOTIFY4\_CHANGE\_GA notification. In this case, GETATTRs based on cached attributes require explicit authorization until changed by a possible NOTIFY4\_CHANGE\_GA notification.

The server can do this to avoid waiting for a scan of the directory looking for troublesome ACLs or multiply-linked linked file that might get troublesome ACLs using one of the other links. The scan can go on with the client being notified of the new status later.

- \* If neither of these bits is set, then the server is indicating the absence of support for avoiding use of ACCESS to check for GETATTR authorization. In this case, GETATTRs based on cached attributes require explicit authorization without the possibility of further change in this situation.

Regardless of the settings of the flags NOTIFY4\_AUFLAG\_GANOW and NOTIFY4\_AUFLAG\_GALATER, accurate reporting of current values of many attributes associated with directory entries cannot be expected because directory entry attribute changes can only be delivered on batched/delayed basis. In order to help ameliorate this situation the NOTIFY4\_CHANGE\_GA notification can provide information allowing full attribute refetch to be avoided in many cases. For this reason clients that request this notification are well advised to accept these notifications and use them to optimize processing as described in Section 27.4.12.

Once this initial state is set, it can be modified as described below, as the server's knowledge of the set of files that require explicit authorization checks changes in response to file system changes.

- \* When a file within the directory is assigned an ACL that can interfere with the client providing cached attributes without ACCESS checks, the client can be notified of that change of status using a NOTIFY4\_CHANGE\_GA notification.
- \* Similarly, when a file within the directory is becomes reachable via an additional link, making it possible that it will subsequently be assigned an ACL without being aware of the directory delegation, there is also a need for the client to be notified. Since such an ACL could interfere with the client providing cached attributes without ACCESS checks, the client is also notified of that change of status using a NOTIFY4\_CHANGE\_GA notification.

#### 16.2.14. Directory Delegation Recall

When necessary the server will recall the directory delegation by sending a callback to the client. It uses the same callback procedure as used for recalling file delegations. The server will recall the delegation in the following situations:

- \* If there is a need to send a content update notification or an authorization update and it is not possible to send that type of notification.

The server will wait for the delegation to be returned or revoked if the notification was one that needed to be sent synchronously.

- \* If a client removes a directory for which a delegation has been granted.

If the server determines the existence of a delegation for a directory is causing too many notifications to be sent out, it may decide to not hand out delegations for that directory and/or recall those already granted. In the case of attribute update notifications, it also has the option of reducing update frequency or limiting set of attributes about which the client is to be notified.

#### 16.2.15. Directory Delegation Recovery

Recovery from client or server restart for state on regular files has two main goals: avoiding the necessity of breaking application guarantees with respect to locked files and delivery of updates cached at the client. Neither of these goals applies to directories protected by OPEN\_DELEGATE\_READ delegations and notifications. As a result, no provision is made for reclaiming directory delegations in the event of client or server restart. The client needs to establish a directory delegation in the same fashion as was done initially.

### 17. Multi-Server Namespace

NFSv4.1 supports attributes that allow a namespace to extend beyond the boundaries of a single server. It is desirable that clients and servers support construction of such multi-server namespaces. Use of such multi-server namespaces is OPTIONAL; however, and for many purposes, single-server namespaces are perfectly acceptable. The use of multi-server namespaces can provide many advantages by separating a file system's logical position in a namespace from the (possibly changing) logistical and administrative considerations that cause a particular file system to be located on a particular server via a single network access path that has to be known in advance or determined using DNS.

#### 17.1. Terminology

In this section as a whole (i.e., within all of Section 17), the phrase "client ID" always refers to the 64-bit shorthand identifier assigned by the server (a `clientid4`) and never to the structure that the client uses to identify itself to the server (called an `nfs_client_id4` or `client_owner` in NFSv4.0 and NFSv4.1, respectively). The opaque identifier within those structures is referred to as a "client id string".

##### 17.1.1. Terminology Related to Trunking

It is particularly important to clarify the distinction between trunking detection and trunking discovery. The definitions we present are applicable to all minor versions of NFSv4, but we will focus on how these terms apply to NFS version 4.1.

- \* Trunking detection refers to ways of deciding whether two specific network addresses are connected to the same NFSv4 server. The means available to make this determination depends on the protocol version, and, in some cases, on the client implementation.

In the case of NFS version 4.1 and later minor versions, the means of trunking detection are as described in this document and are available to every client. Two network addresses connected to the same server can always be used together to access a particular server but cannot necessarily be used together to access a single session. See below for definitions of the terms "server-trunkable" and "session-trunkable".

- \* Trunking discovery is a process by which a client using one network address can obtain other addresses that are connected to the same server. Typically, it builds on a trunking detection facility by providing one or more methods by which candidate addresses are made available to the client, who can then use trunking detection to appropriately filter them.

Despite the support for trunking detection, there was no description of trunking discovery provided in [RFC5661] or [RFC8881], making it necessary to provide those means in this document.

The combination of a server network address and a particular connection type to be used by a connection is referred to as a "server endpoint". Although using different connection types may result in different ports being used, the use of different ports by multiple connections to the same network address in such cases is not the essence of the distinction between the two endpoints used. This is in contrast to the case of port-specific endpoints, in which the explicit specification of port numbers within network addresses is used to allow a single server node to support multiple NFS servers.

Two network addresses connected to the same server are said to be server-trunkable. Two such addresses support the use of client ID trunking, as described in Section 7.5.

Two network addresses connected to the same server such that those addresses can be used to support a single common session are referred to as session-trunkable. Note that two addresses may be server-trunkable without being session-trunkable, and that, when two connections of different connection types are made to the same network address and are based on a single file system location entry, they are always session-trunkable, independent of the connection type, as specified by Section 7.5, since their derivation from the same file system location entry, together with the identity of their network addresses, assures that both connections are to the same server and will return server-owner information, allowing session trunking to be used.

### 17.1.2. Terminology Related to File System Location

Regarding the terminology that relates to the construction of multi-server namespaces out of a set of local per-server namespaces:

- \* Each server has a set of exported file systems that may be accessed by NFSv4 clients. Typically, this is done by assigning each file system a name within the pseudo-fs associated with the server, although the pseudo-fs may be dispensed with if there is only a single exported file system. Each such file system is part of the server's local namespace, and can be considered as a file system instance within a larger multi-server namespace.
- \* The set of all exported file systems for a given server constitutes that server's local namespace.
- \* In some cases, a server will have a namespace more extensive than its local namespace by using features associated with attributes that provide file system location information. These features, which allow construction of a multi-server namespace, are all described in individual sections below and include referrals (Section 17.5.6), migration (Section 17.5.5), and replication (Section 17.5.4).
- \* A file system present in a server's pseudo-fs may have multiple file system instances on different servers associated with it. All such instances are considered replicas of one another. Whether such replicas can be used simultaneously is discussed in Section 17.11.1, while the level of coordination between them (important when switching between them) is discussed in Sections 17.11.2 through 17.11.8 below.
- \* When a file system is present in a server's pseudo-fs, but there is no corresponding local file system, it is said to be "absent". In such cases, all associated instances will be accessed on other servers.

Regarding the terminology that relates to attributes used in trunking discovery and other multi-server namespace features:

- \* File system location attributes include the `fs_locations` and `fs_locations_info` attributes.
- \* File system location entries provide the individual file system locations within the file system location attributes. Each such entry specifies a server, in the form of a hostname or an address, and an fs name, which designates the location of the file system within the server's local namespace. A file system location entry

designates a set of server endpoints to which the client may establish connections. There may be multiple endpoints because a hostname may map to multiple network addresses and because multiple connection types may be used to communicate with a single network address. However, except where explicit port numbers are used to designate a set of servers within a single server node, all such endpoints MUST designate a way of connecting to a single server. The exact form of the location entry varies with the particular file system location attribute used, as described in Section 17.2.

The network addresses used in file system location entries typically appear without port number indications and are used to designate a server at one of the standard ports for NFS access, e.g., 2049 for TCP or 20049 for use with RPC-over-RDMA. Port numbers may be used in file system location entries to designate servers (typically user-level ones) accessed using other port numbers. In the case where network addresses indicate trunking relationships, the use of an explicit port number is inappropriate since trunking is a relationship between network addresses. See Section 17.5.2 for details.

- \* File system location elements are derived from location entries, and each describes a particular network access path consisting of a network address and a location within the server's local namespace. Such location elements need not appear within a file system location attribute, but the existence of each location element derives from a corresponding location entry. When a location entry specifies an IP address, there is only a single corresponding location element. File system location entries that contain a hostname are resolved using DNS, and may result in one or more location elements. All location elements consist of a location address that includes the IP address of an interface to a server and an fs name, which is the location of the file system within the server's local namespace. The fs name can be empty if the server has no pseudo-fs and only a single exported file system at the root filehandle.
- \* Two file system location elements are said to be server-trunkable if they specify the same fs name and the location addresses are such that the location addresses are server-trunkable. When the corresponding network paths are used, the client will always be able to use client ID trunking, but will only be able to use session trunking if the paths are also session-trunkable.

- \* Two file system location elements are said to be session-trunkable if they specify the same fs name and the location addresses are such that the location addresses are session-trunkable. When the corresponding network paths are used, the client will be able to able to use either client ID trunking or session trunking.

Discussion of the term "replica" is complicated by the fact that the term was used in [RFC5661] with a meaning different from that used in this document. In short, in [RFC5661] each replica is identified by a single network access path, while in the current document, a set of network access paths that have server-trunkable network addresses and the same root-relative file system pathname is considered to be a single replica with multiple network access paths.

Each set of server-trunkable location elements defines a set of available network access paths to a particular file system. When there are multiple such file systems, each of which containing the same data, these file systems are considered replicas of one another. Logically, such replication is symmetric, since the fs currently in use and an alternate fs are replicas of each other. Often, in other documents, the term "replica" is not applied to the fs currently in use, despite the fact that the replication relation is inherently symmetric.

## 17.2. File System Location Attributes

NFSv4.1 contains attributes that provide information about how a given file system may be accessed (i.e., at what network address and namespace position). As a result, file systems in the namespace of one server can be associated with one or more instances of that file system on other servers. These attributes contain file system location entries specifying a server address target (either as a DNS name representing one or more IP addresses or as a specific IP address) together with the pathname of that file system within the associated single-server namespace.

The `fs_locations_info` attribute allows specification of one or more file system instance locations where the data corresponding to a given file system may be found. In addition to the specification of file system instance locations, this attribute provides helpful information to do the following:

- \* Guide choices among the various file system instances provided (e.g., priority for use, writability, currency, etc.).
- \* Help the client efficiently effect as seamless a transition as possible among multiple file system instances, when and if that should be necessary.



- \* Guide the selection of the appropriate connection type to be used when establishing a connection.

Within the `fs_locations_info` attribute, each `fs_locations_server4` entry corresponds to a file system location entry: the `fls_server` field designates the server, and the `fl_rootpath` field of the encompassing `fs_locations_item4` gives the location pathname within the server's pseudo-fs.

The `fs_locations` attribute defined in NFSv4.0 is also a part of NFSv4.1. This attribute only allows specification of the file system locations where the data corresponding to a given file system may be found. Servers should make this attribute available whenever `fs_locations_info` is supported, but client use of `fs_locations_info` is preferable because it provides more information.

Within the `fs_locations` attribute, each `fs_location4` contains a file system location entry with the `server` field designating the server and the `rootpath` field giving the location pathname within the server's pseudo-fs.

### 17.3. File System Presence or Absence

A given location in an NFSv4.1 namespace (typically but not necessarily a multi-server namespace) can have a number of file system instance locations associated with it (via the `fs_locations` or `fs_locations_info` attribute). There may also be an actual current file system at that location, accessible via normal namespace operations (e.g., LOOKUP). In this case, the file system is said to be "present" at that position in the namespace, and clients will typically use it, reserving use of additional locations specified via the location-related attributes to situations in which the principal location is no longer available.

When there is no actual file system at the namespace location in question, the file system is said to be "absent". An absent file system contains no files or directories other than the root. Any reference to it, except to access a small set of attributes useful in determining alternate locations, will result in an error, `NFS4ERR_MOVED`. Note that if the server ever returns the error `NFS4ERR_MOVED`, it MUST support the `fs_locations` attribute and should support the `fs_locations_info` and `fs_status` attributes.

While the error name suggests that we have a case of a file system that once was present, and has only become absent later, this is only one possibility. A position in the namespace may be permanently absent with the set of file system(s) designated by the location attributes being the only realization. The name `NFS4ERR_MOVED`

reflects an earlier, more limited conception of its function, but this error will be returned whenever the referenced file system is absent, whether it has moved or not.

Except in the case of GETATTR-type operations (to be discussed later), when the current filehandle at the start of an operation is within an absent file system, that operation is not performed and the error NFS4ERR\_MOVED is returned, to indicate that the file system is absent on the current server.

Because a GETFH cannot succeed if the current filehandle is within an absent file system, filehandles within an absent file system cannot be transferred to the client. When a client does have filehandles within an absent file system, it is the result of obtaining them when the file system was present, and having the file system become absent subsequently.

It should be noted that because the check for the current filehandle being within an absent file system happens at the start of every operation, operations that change the current filehandle so that it is within an absent file system will not result in an error. This allows such combinations as PUTFH-GETATTR and LOOKUP-GETATTR to be used to get attribute information, particularly location attribute information, as discussed below.

The file system attribute `fs_status` can be used to interrogate the present/absent status of a given file system.

#### 17.4. Getting Attributes for an Absent File System

When a file system is absent, most attributes are not available, but it is necessary to allow the client access to the small set of attributes that are available, and most particularly those that give information about the correct current locations for this file system: `fs_locations` and `fs_locations_info`.

##### 17.4.1. GETATTR within an Absent File System

As mentioned above, an exception is made for GETATTR in that attributes may be obtained for a filehandle within an absent file system. This exception only applies if the attribute mask contains at least one attribute bit that indicates the client is interested in a result regarding an absent file system: `fs_locations`, `fs_locations_info`, or `fs_status`. If none of these attributes is requested, GETATTR will result in an NFS4ERR\_MOVED error.

When a GETATTR is done on an absent file system, the set of supported attributes is very limited. Many attributes, including those that are normally REQUIRED, will not be available on an absent file system. In addition to the attributes mentioned above (fs\_locations, fs\_locations\_info, fs\_status), the following attributes should be available on absent file systems. In the case of OPTIONAL attributes, they should be available at least to the same degree that they are available on present file systems.

**change\_policy:** This attribute is useful for absent file systems and can be helpful in summarizing to the client when any of the location-related attributes change.

**fsid:** This attribute should be provided so that the client can determine file system boundaries, including, in particular, the boundary between present and absent file systems. This value must be different from any other fsid on the current server and need have no particular relationship to fsids on any particular destination to which the client might be directed.

**mounted\_on\_fileid:** For objects at the top of an absent file system, this attribute needs to be available. Since the fileid is within the present parent file system, there should be no need to reference the absent file system to provide this information.

Other attributes should not be made available for absent file systems, even when it is possible to provide them. The server should not assume that more information is always better and should avoid gratuitously providing additional information.

When a GETATTR operation includes a bit mask for one of the attributes fs\_locations, fs\_locations\_info, or fs\_status, but where the bit mask includes attributes that are not supported, GETATTR will not return an error, but will return the mask of the actual attributes supported with the results.

Handling of VERIFY/NVERIFY is similar to GETATTR in that if the attribute mask does not include fs\_locations, fs\_locations\_info, or fs\_status, the error NFS4ERR\_MOVED will result. It differs in that any appearance in the attribute mask of an attribute not supported for an absent file system (and note that this will include some normally REQUIRED attributes) will also cause an NFS4ERR\_MOVED result.

#### 17.4.2. READDIR and Absent File Systems

A READDIR performed when the current filehandle is within an absent file system will result in an NFS4ERR\_MOVED error, since, unlike the case of GETATTR, no such exception is made for READDIR.

Attributes for an absent file system may be fetched via a READDIR for a directory in a present file system, when that directory contains the root directories of one or more absent file systems. In this case, the handling is as follows:

- \* If the attribute set requested includes one of the attributes `fs_locations`, `fs_locations_info`, or `fs_status`, then fetching of attributes proceeds normally and no NFS4ERR\_MOVED indication is returned, even when the `rdattr_error` attribute is requested.
- \* If the attribute set requested does not include one of the attributes `fs_locations`, `fs_locations_info`, or `fs_status`, then if the `rdattr_error` attribute is requested, each directory entry for the root of an absent file system will report NFS4ERR\_MOVED as the value of the `rdattr_error` attribute.
- \* If the attribute set requested does not include any of the attributes `fs_locations`, `fs_locations_info`, `fs_status`, or `rdattr_error`, then the occurrence of the root of an absent file system within the directory will result in the READDIR failing with an NFS4ERR\_MOVED error.
- \* The unavailability of an attribute because of a file system's absence, even one that is ordinarily REQUIRED, does not result in any error indication. The set of attributes returned for the root directory of the absent file system in that case is simply restricted to those actually available.

#### 17.5. Uses of File System Location Information

The file system location attributes (i.e., `fs_locations` and `fs_locations_info`), together with the possibility of absent file systems, provide a number of important facilities for reliable, manageable, and scalable data access.

When a file system is present, these attributes can provide the following:

- \* The locations of alternative replicas to be used to access the same data in the event of server failures, communications problems, or other difficulties that make continued access to the current replica impossible or otherwise impractical. Provisioning and use of such alternate replicas is referred to as "replication" and is discussed in Section 17.5.4 below.
- \* The network address(es) to be used to access the current file system instance or replicas of it. Client use of this information is discussed in Section 17.5.2 below.

Under some circumstances, multiple replicas may be used simultaneously to provide higher-performance access to the file system in question, although the lack of state sharing between servers may be an impediment to such use.

When a file system is present but becomes absent, clients can be given the opportunity to have continued access to their data using a different replica. In this case, a continued attempt to use the data in the now-absent file system will result in an NFS4ERR\_MOVED error, and then the successor replica or set of possible replica choices can be fetched and used to continue access. Transfer of access to the new replica location is referred to as "migration" and is discussed in Section 17.5.4 below.

When a file system is currently absent, specification of file system location provides a means by which file systems located on one server can be associated with a namespace defined by another server, thus allowing a general multi-server namespace facility. A designation of such a remote instance, in place of a file system not previously present, is called a "pure referral" and is discussed in Section 17.5.6 below.

Because client support for attributes related to file system location is OPTIONAL, a server may choose to take action to hide migration and referral events from such clients, by acting as a proxy, for example. The server can determine the presence of client support from the arguments of the EXCHANGE\_ID operation (See Section 25.35.3).

#### 17.5.1. Combining Multiple Uses in a Single Attribute

A file system location attribute will sometimes contain information relating to the location of multiple replicas, which may be used in different ways:

- \* File system location entries that relate to the file system instance currently in use provide trunking information, allowing the client to find additional network addresses by which the instance may be accessed.
- \* File system location entries that provide information about replicas to which access is to be transferred.
- \* Other file system location entries that relate to replicas that are available to use in the event that access to the current replica becomes unsatisfactory.

In order to simplify client handling and to allow the best choice of replicas to access, the server should adhere to the following guidelines:

- \* All file system location entries that relate to a single file system instance should be adjacent.
- \* File system location entries that relate to the instance currently in use should appear first.
- \* File system location entries that relate to replica(s) to which migration is occurring should appear before replicas that are available for later use if the current replica should become inaccessible.

#### 17.5.2. File System Location Attributes and Trunking

Trunking is the use of multiple connections between a client and server in order to increase the speed of data transfer. A client may determine the set of network addresses to use to access a given file system in a number of ways:

- \* When the name of the server is known to the client, it may use DNS to obtain a set of network addresses to use in accessing the server.

- \* The client may fetch the file system location attribute for the file system. This will provide either the name of the server (which can be turned into a set of network addresses using DNS) or a set of server-trunkable location entries. Using the latter alternative, the server can provide addresses it regards as desirable to use to access the file system in question. Although these entries can contain port numbers, these port numbers are not used in determining trunking relationships. Once the candidate addresses have been determined and EXCHANGE\_ID done to the proper server, only the value of the so\_major\_id field returned by the servers in question determines whether a trunking relationship actually exists.

When the client fetches a location attribute for a file system, it should be noted that the client may encounter multiple entries for a number of reasons, such that when it determines trunking information, it may need to bypass addresses not trunkable with one already known.

The server can provide location entries that include either names or network addresses. It might use the latter form because of DNS-related security concerns or because the set of addresses to be used might require active management by the server.

Location entries used to discover candidate addresses for use in trunking are subject to change, as discussed in Section 17.5.7 below. The client may respond to such changes by using additional addresses once they are verified or by ceasing to use existing ones. The server can force the client to cease using an address by returning NFS4ERR\_MOVED when that address is used to access a file system. This allows a transfer of client access that is similar to migration, although the same file system instance is accessed throughout.

### 17.5.3. File System Location Attributes and Connection Type Selection

Because of the need to support multiple types of connections, clients face the issue of determining the proper connection type to use when establishing a connection to a given server network address. In some cases, this issue can be addressed through the use of the connection "step-up" facility described in Section 25.36. However, because there are cases in which that facility is not available, the client may have to choose a connection type with no possibility of changing it within the scope of a single connection.

The two file system location attributes differ as to the information made available in this regard. The `fs_locations` attribute provides no information to support connection type selection. As a result, clients supporting multiple connection types would need to attempt to establish connections using multiple connection types until the one preferred by the client is successfully established.

The `fs_locations_info` attribute includes the `FSLI4TF_RDMA` flag, which is convenient for a client wishing to use RDMA. When this flag is set, it indicates that RPC-over-RDMA support is available using the specified location entry. A client can establish a TCP connection and then convert that connection to use RDMA by using the step-up facility.

Irrespective of the particular attribute used, when there is no indication that a step-up operation can be performed, a client supporting RDMA operation can establish a new RDMA connection, and it can be bound to the session already established by the TCP connection, allowing the TCP connection to be dropped and the session converted to further use in RDMA mode, if the server supports that.

#### 17.5.4. File System Replication

The `fs_locations` and `fs_locations_info` attributes provide alternative file system locations, to be used to access data in place of or in addition to the current file system instance. On first access to a file system, the client should obtain the set of alternate locations by interrogating the `fs_locations` or `fs_locations_info` attribute, with the latter being preferred.

In the event that the occurrence of server failures, communications problems, or other difficulties make continued access to the current file system impossible or otherwise impractical, the client can use the alternate locations as a way to get continued access to its data.

The alternate locations may be physical replicas of the (typically read-only) file system data supplemented by possible asynchronous propagation of updates. Alternatively, they may provide for the use of various forms of server clustering in which multiple servers provide alternate ways of accessing the same physical file system. How the difference between replicas affects file system transitions can be represented within the `fs_locations` and `fs_locations_info` attributes, and how the client deals with file system transition issues will be discussed in detail in later sections.

Although the location attributes provide some information about the nature of the inter-replica transition, many aspects of the semantics of possible asynchronous updates are not currently described by the



protocol, which makes it necessary for clients using replication to switch among replicas undergoing change to familiarize themselves with the semantics of the update approach used. Due to this lack of specificity, many applications may find the use of migration more appropriate because a server can propagate all updates made before an established point in time to the new replica as part of the migration event.

#### 17.5.4.1. File System Trunking Presented as Replication

In some situations, a file system location entry may indicate a file system access path to be used as an alternate location, where trunking, rather than replication, is to be used. The situations in which this is appropriate are limited to those in which both of the following are true:

- \* The two file system locations (i.e., the one on which the location attribute is obtained and the one specified in the file system location entry) designate the same locations within their respective single-server namespaces.
- \* The two server network addresses (i.e., the one being used to obtain the location attribute and the one specified in the file system location entry) designate the same server (as indicated by the same value of the `so_major_id` field of the `eir_server_owner` field returned in response to `EXCHANGE_ID`).

When these conditions hold, operations using both access paths are generally trunked, although trunking may be disallowed when the attribute `fs_locations_info` is used:

- \* When the `fs_locations_info` attribute shows the two entries as not having the same simultaneous-use class, trunking is inhibited, and the two access paths cannot be used together.

In this case, the two paths can be used serially with no transition activity required on the part of the client, and any transition between access paths is transparent. In transferring access from one to the other, the client acts as if communication were interrupted, establishing a new connection and possibly a new session to continue access to the same file system.

- \* Note that for two such location entries, any information within the `fs_locations_info` attribute that indicates the need for special transition activity, i.e., the appearance of the two file system location entries with different `handle`, `fileid`, `write-verifier`, `change`, and `readdir` classes, indicates a serious problem. The client, if it allows transition to the file system

instance at all, must not treat any transition as a transparent one. The server should not indicate that these two entries (for the same file system on the same server) belong to different handle, fileid, write-verifier, change, and readdir classes, whether or not the two entries are shown belonging to the same simultaneous-use class.

These situations were recognized by [RFC5661], even though that document made no explicit mention of trunking:

- \* It treated the situation that we describe as trunking as one of simultaneous use of two distinct file system instances, even though, in the explanatory framework now used to describe the situation, the case is one in which a single file system is accessed by two different trunked addresses.
- \* It treated the situation in which two paths are to be used serially as a special sort of "transparent transition". However, in the descriptive framework now used to categorize transition situations, this is considered a case of a "network endpoint transition" (See Section 17.9).

#### 17.5.5. File System Migration

When a file system is present and becomes inaccessible using the current access path, the NFSv4.1 protocol provides a means by which clients can be given the opportunity to have continued access to their data. This may involve using a different access path to the existing replica or providing a path to a different replica. The new access path or the location of the new replica is specified by a file system location attribute. The ensuing migration of access includes the ability to retain locks across the transition. Depending on circumstances, this can involve:

- \* The continued use of the existing clientid when accessing the current replica using a new access path.
- \* Use of lock reclaim, taking advantage of a per-fs grace period.
- \* Use of Transparent State Migration.

Typically, a client will be accessing the file system in question, get an NFS4ERR\_MOVED error, and then use a file system location attribute to determine the new access path for the data. When fs\_locations\_info is used, additional information will be available that will define the nature of the client's handling of the transition to a new server.

In most instances, servers will choose to migrate all clients using a particular file system to a successor replica at the same time to avoid cases in which different clients are updating different replicas. However, migration of an individual client can be helpful in providing load balancing, as long as the replicas in question are such that they represent the same data as described in Section 17.11.8.

- \* In the case in which there is no transition between replicas (i.e., only a change in access path), there are no special difficulties in using of this mechanism to effect load balancing.
- \* In the case in which the two replicas are sufficiently coordinated as to allow a single client coherent, simultaneous access to both, there is, in general, no obstacle to the use of migration of particular clients to effect load balancing. Generally, such simultaneous use involves cooperation between servers to ensure that locks granted on two coordinated replicas cannot conflict and can remain effective when transferred to a common replica.
- \* In the case in which a large set of clients is accessing a file system in a read-only fashion, it can be helpful to migrate all clients with writable access simultaneously, while using load balancing on the set of read-only copies, as long as the rules in Section 17.11.8, which are designed to prevent data reversion, are followed.

In other cases, the client might not have sufficient guarantees of data similarity or coherence to function properly (e.g., the data in the two replicas is similar but not identical), and the possibility that different clients are updating different replicas can exacerbate the difficulties, making the use of load balancing in such situations a perilous enterprise.

The protocol does not specify how the file system will be moved between servers or how updates to multiple replicas will be coordinated. It is anticipated that a number of different server-to-server coordination mechanisms might be used, with the choice left to the server implementer. The NFSv4.1 protocol specifies the method used to communicate the migration event between client and server.

In the case of various forms of server clustering, the new location may be another server providing access to the same physical file system. The client's responsibilities in dealing with this transition will depend on whether a switch between replicas has occurred and the means the server has chosen to provide continuity of locking state. These issues will be discussed in detail below.

Although a single successor location is typical, multiple locations may be provided. When multiple locations are provided, the client will typically use the first one provided. If that is inaccessible for some reason, later ones can be used. In such cases, the client might consider the transition to the new replica to be a migration event, even though some of the servers involved might not be aware of the use of the server that was inaccessible. In such a case, a client might lose access to locking state as a result of the access transfer.

When an alternate location is designated as the target for migration, it must designate the same data (with metadata being the same to the degree indicated by the `fs_locations_info` attribute). Where file systems are writable, a change made on the original file system must be visible on all migration targets. Where a file system is not writable but represents a read-only copy (possibly periodically updated) of a writable file system, similar requirements apply to the propagation of updates. Any change visible in the original file system must already be effected on all migration targets, to avoid any possibility that a client, in effecting a transition to the migration target, will see any reversion in file system state.

#### 17.5.6. Referrals

Referrals allow the server to associate a file system namespace entry located on one server with a file system located on another server. When this includes the use of pure referrals, servers are provided a way of placing a file system in a location within the namespace essentially without respect to its physical location on a particular server. This allows a single server or a set of servers to present a multi-server namespace that encompasses file systems located on a wider range of servers. Some likely uses of this facility include establishment of site-wide or organization-wide namespaces, with the eventual possibility of combining such together into a truly global namespace, such as the one provided by AFS (the Andrew File System) [AFS].

Referrals occur when a client determines, upon first referencing a position in the current namespace, that it is part of a new file system and that the file system is absent. When this occurs, typically upon receiving the error `NFS4ERR_MOVED`, the actual location or locations of the file system can be determined by fetching a `locations` attribute.

The file system location attribute may designate a single file system location or multiple file system locations, to be selected based on the needs of the client. The server, in the `fs_locations_info` attribute, may specify priorities to be associated with various file

system location choices. The server may assign different priorities to different locations as reported to individual clients, in order to adapt to client physical location or to effect load balancing. When both read-only and read-write file systems are present, some of the read-only locations might not be absolutely up-to-date (as they would have to be in the case of replication and migration). Servers may also specify file system locations that include client-substituted variables so that different clients are referred to different file systems (with different data contents) based on client attributes such as CPU architecture.

If the `fs_locations_info` attribute lists multiple possible targets, the relationships among them may be important to the client in selecting which one to use. The same rules specified in Section 17.5.5 below regarding multiple migration targets apply to these multiple replicas as well. For example, the client might prefer a writable target on a server that has additional writable replicas to which it subsequently might switch. Note that, as distinguished from the case of replication, there is no need to deal with the case of propagation of updates made by the current client, since the current client has not accessed the file system in question.

Use of multi-server namespaces is enabled by NFSv4.1 but is not required. The use of multi-server namespaces and their scope will depend on the applications used and system administration preferences.

Multi-server namespaces can be established by a single server providing a large set of pure referrals to all of the included file systems. Alternatively, a single multi-server namespace may be administratively segmented with separate referral file systems (on separate servers) for each separately administered portion of the namespace. The top-level referral file system or any segment may use replicated referral file systems for higher availability.

Generally, multi-server namespaces are for the most part uniform, in that the same data made available to one client at a given location in the namespace is made available to all clients at that namespace location. However, there are facilities provided that allow different clients to be directed to different sets of data, for reasons such as enabling adaptation to such client characteristics as CPU architecture. These facilities are described in Section 17.17.3.

Note that it is possible, when providing a uniform namespace, to provide different location entries to different clients in order to provide each client with a copy of the data physically closest to it or otherwise optimize access (e.g., provide load balancing).

#### 17.5.7. Changes in File System Location Attributes

Although clients will typically fetch a file system location attribute when first accessing a file system and when NFS4ERR\_MOVED is returned, a client can choose to fetch the attribute periodically, in which case, the value fetched may change over time.

For clients not prepared to access multiple replicas simultaneously (see Section 17.11.1), the handling of the various cases of location change are as follows:

- \* Changes in the list of replicas or in the network addresses associated with replicas do not require immediate action. The client will typically update its list of replicas to reflect the new information.
- \* Additions to the list of network addresses for the current file system instance need not be acted on promptly. However, to prepare for a subsequent migration event, the client can choose to take note of the new address and then use it whenever it needs to switch access to a new replica.
- \* Deletions from the list of network addresses for the current file system instance do not require the client to immediately cease use of existing access paths, although new connections are not to be established on addresses that have been deleted. However, clients can choose to act on such deletions by preparing for an eventual shift in access, which becomes unavoidable as soon as the server returns NFS4ERR\_MOVED to indicate that a particular network access path is not usable to access the current file system.

For clients that are prepared to access several replicas simultaneously, the following additional cases need to be addressed. As in the cases discussed above, changes in the set of replicas need not be acted upon promptly, although the client has the option of adjusting its access even in the absence of difficulties that would lead to the selection of a new replica.

- \* When a new replica is added, which may be accessed simultaneously with one currently in use, the client is free to use the new replica immediately.
- \* When a replica currently in use is deleted from the list, the client need not cease using it immediately. However, since the server may subsequently force such use to cease (by returning NFS4ERR\_MOVED), clients might decide to limit the need for later state transfer. For example, new opens might be done on other replicas, rather than on one not present in the list.

#### 17.6. Trunking without File System Location Information

In situations in which a file system is accessed using two server-trunkable addresses (as indicated by the same value of the `so_major_id` field of the `eir_server_owner` field returned in response to `EXCHANGE_ID`), trunked access is allowed even though there might not be any location entries specifically indicating the use of trunking for that file system.

This situation was recognized by [RFC5661], although that document made no explicit mention of trunking and treated the situation as one of simultaneous use of two distinct file system instances. In the explanatory framework now used to describe the situation, the case is one in which a single file system is accessed by two different trunked addresses.

#### 17.7. Users and Groups in a Multi-Server Namespace

As in the case of a single-server environment (see Section 11.13), when an owner or group name of the form "id@domain" is assigned to a file, there is an implicit promise to return that same string when the corresponding attribute is interrogated subsequently. In the case of a multi-server namespace, that same promise applies even if server boundaries have been crossed. Similarly, when the owner attribute of a file is derived from the security principal that created the file, that attribute should have the same value even if the interrogation occurs on a different server from the file creation.

Similarly, the set of security principals recognized by all the participating servers needs to be the same, with each such principal having the same credentials, regardless of the particular server being accessed.

In order to meet these requirements, those setting up multi-server namespaces will need to limit the servers included so that:

- \* In all cases in which more than a single domain is supported, the requirements stated in [RFC8000] are to be respected.
- \* All servers support a common set of domains that includes all of the domains clients use and expect to see returned as the domain portion of an owner or group in the form "id@domain". Note that, although this set most often consists of a single domain, it is possible for multiple domains to be supported.
- \* All servers, for each domain that they support, accept the same set of user and group ids as valid.

- \* All servers recognize the same set of security principals. For each principal, the same credential is required, independent of the server being accessed. In addition, the group membership for each such principal is to be the same, independent of the server accessed.

Note that there is no requirement in general that the users corresponding to particular security principals have the same local representation on each server, even though it is most often the case that this is so.

When AUTH\_SYS is used, the following additional requirements must be met:

- \* Only a single NFSv4 domain can be supported through the use of AUTH\_SYS.
- \* The "local" representation of all owners and groups must be the same on all servers. The word "local" is used here since that is the way that numeric user and group ids are described in Section 11.13. However, when AUTH\_SYS or stringified numeric owners or groups are used, these identifiers are not truly local, since they are known to the clients as well as to the server.

Similarly, when stringified numeric user and group ids are used, the "local" representation of all owners and groups must be the same on all servers, even when AUTH\_SYS is not used.

#### 17.8. Additional Client-Side Considerations

When clients make use of servers that implement referrals, replication, and migration, care should be taken that a user who mounts a given file system that includes a referral or a relocated file system continues to see a coherent picture of that user-side file system despite the fact that it contains a number of server-side file systems that may be on different servers.

One important issue is upward navigation from the root of a server-side file system to its parent (specified as ".." in UNIX), in the case in which it transitions to that file system as a result of referral, migration, or a transition as a result of replication. When the client is at such a point, and it needs to ascend to the parent, it must go back to the parent as seen within the multi-server namespace rather than sending a LOOKUPP operation to the server, which would result in the parent within that server's single-server namespace. In order to do this, the client needs to remember the filehandles that represent such file system roots and use these instead of sending a LOOKUPP operation to the current server. This



will allow the client to present to applications a consistent namespace, where upward navigation and downward navigation are consistent.

Another issue concerns refresh of referral locations. When referrals are used extensively, they may change as server configurations change. It is expected that clients will cache information related to traversing referrals so that future client-side requests are resolved locally without server communication. This is usually rooted in client-side name look up caching. Clients should periodically purge this data for referral points in order to detect changes in location information. When the `change_policy` attribute changes for directories that hold referral entries or for the referral entries themselves, clients should consider any associated cached referral information to be out of date.

#### 17.9. Overview of File Access Transitions

File access transitions are of two types:

- \* Those that involve a transition from accessing the current replica to another one in connection with either replication or migration. How these are dealt with is discussed in Section 17.11.
- \* Those in which access to the current file system instance is retained, while the network path used to access that instance is changed. This case is discussed in Section 17.10.

#### 17.10. Effecting Network Endpoint Transitions

The endpoints used to access a particular file system instance may change in a number of ways, as listed below. In each of these cases, the same `fsid`, client IDs, filehandles, and `stateids` are used to continue access, with a continuity of lock state. In many cases, the same sessions can also be used.

The appropriate action depends on the set of replacement addresses that are available for use (i.e., server endpoints that are server-trunkable with one previously being used).

- \* When use of a particular address is to cease, and there is also another address currently in use that is server-trunkable with it, requests that would have been issued on the address whose use is to be discontinued can be issued on the remaining address(es). When an address is server-trunkable but not session-trunkable with the address whose use is to be discontinued, the request might need to be modified to reflect the fact that a different session will be used.

- \* When use of a particular connection is to cease, as indicated by receiving NFS4ERR\_MOVED when using that connection, but that address is still indicated as accessible according to the appropriate file system location entries, it is likely that requests can be issued on a new connection of a different connection type once that connection is established. Since any two non-port-specific server endpoints that share a network address are inherently session-trunkable, the client can use BIND\_CONN\_TO\_SESSION to access the existing session with the new connection.
- \* When there are no potential replacement addresses in use, but there are valid addresses session-trunkable with the one whose use is to be discontinued, the client can use BIND\_CONN\_TO\_SESSION to access the existing session using the new address. Although the target session will generally be accessible, there may be rare situations in which that session is no longer accessible when an attempt is made to bind the new connection to it. In this case, the client can create a new session to enable continued access to the existing instance using the new connection, providing for the use of existing filehandles, stateids, and client ids while supplying continuity of locking state.
- \* When there is no potential replacement address in use, and there are no valid addresses session-trunkable with the one whose use is to be discontinued, other server-trunkable addresses may be used to provide continued access. Although the use of CREATE\_SESSION is available to provide continued access to the existing instance, servers have the option of providing continued access to the existing session through the new network access path in a fashion similar to that provided by session migration (See Section 17.12). To take advantage of this possibility, clients can perform an initial BIND\_CONN\_TO\_SESSION, as in the previous case, and use CREATE\_SESSION only if that fails.

#### 17.11. Effecting File System Transitions

There are a range of situations in which there is a change to be effected in the set of replicas used to access a particular file system. Some of these may involve an expansion or contraction of the set of replicas used as discussed in Section 17.11.1 below.

For reasons explained in that section, most transitions will involve a transition from a single replica to a corresponding replacement replica. When effecting replica transition, some types of sharing between the replicas may affect handling of the transition as described in Sections 17.11.2 through 17.11.8 below. The attribute `fs_locations_info` provides helpful information to allow the client to determine the degree of inter-replica sharing.

With regard to some types of state, the degree of continuity across the transition depends on the occasion prompting the transition, with transitions initiated by the servers (i.e., migration) offering much more scope for a nondisruptive transition than cases in which the client on its own shifts its access to another replica (i.e., replication). This issue potentially applies to locking state and to session state, which are dealt with below as follows:

- \* An introduction to the possible means of providing continuity in these areas appears in Section 17.11.9 below.
- \* Transparent State Migration is introduced in Section 17.12. The possible transfer of session state is addressed there as well.
- \* The client handling of transitions, including determining how to deal with the various means that the server might take to supply effective continuity of locking state, is discussed in Section 17.13.
- \* The source and destination servers' responsibilities in effecting Transparent State Migration of locking and session state are discussed in Section 17.14.

#### 17.11.1. File System Transitions and Simultaneous Access

The `fs_locations_info` attribute (described in Section 17.17) may indicate that two replicas may be used simultaneously, although some situations in which such simultaneous access is permitted are more appropriately described as instances of trunking (See Section 17.5.4.1). Although situations in which multiple replicas may be accessed simultaneously are somewhat similar to those in which a single replica is accessed by multiple network addresses, there are important differences since locking state is not shared among multiple replicas.

Because of this difference in state handling, many clients will not have the ability to take advantage of the fact that such replicas represent the same data. Such clients will not be prepared to use multiple replicas simultaneously but will access each file system using only a single replica, although the replica selected might make multiple server-trunkable addresses available.

Clients who are prepared to use multiple replicas simultaneously can divide opens among replicas however they choose. Once that choice is made, any subsequent transitions will treat the set of locking state associated with each replica as a single entity.

For example, if one of the replicas become unavailable, access will be transferred to a different replica, which is also capable of simultaneous access with the one still in use.

When there is no such replica, the transition may be to the replica already in use. At this point, the client has a choice between merging the locking state for the two replicas under the aegis of the sole replica in use or treating these separately until another replica capable of simultaneous access presents itself.

#### 17.11.2. Filehandles and File System Transitions

There are a number of ways in which filehandles can be handled across a file system transition. These can be divided into two broad classes depending upon whether the two file systems across which the transition happens share sufficient state to effect some sort of continuity of file system handling.

When there is no such cooperation in filehandle assignment, the two file systems are reported as being in different handle classes. In this case, all filehandles are assumed to expire as part of the file system transition. Note that this behavior does not depend on the `fh_expire_type` attribute and supersedes the specification of the `FH4_VOL_MIGRATION` bit, which only affects behavior when `fs_locations_info` is not available.

When there is cooperation in filehandle assignment, the two file systems are reported as being in the same handle classes. In this case, persistent filehandles remain valid after the file system transition, while volatile filehandles (excluding those that are only volatile due to the `FH4_VOL_MIGRATION` bit) are subject to expiration on the target server.

### 17.11.3. Fileids and File System Transitions

In NFSv4.0, the issue of continuity of fileids in the event of a file system transition was not addressed. The general expectation had been that in situations in which the two file system instances are created by a single vendor using some sort of file system image copy, fileids would be consistent across the transition, while in the analogous multi-vendor transitions they would not. This poses difficulties, especially for the client without special knowledge of the transition mechanisms adopted by the server. Note that although fileid is not a REQUIRED attribute, many servers support fileids and many clients provide APIs that depend on fileids.

It is important to note that while clients themselves may have no trouble with a fileid changing as a result of a file system transition event, applications do typically have access to the fileid (e.g., via stat). The result is that an application may work perfectly well if there is no file system instance transition or if any such transition is among instances created by a single vendor, yet be unable to deal with the situation in which a multi-vendor transition occurs at the wrong time.

Providing the same fileids in a multi-vendor (multiple server vendors) environment has generally been held to be quite difficult. While there is work to be done, it needs to be pointed out that this difficulty is partly self-imposed. Servers have typically identified fileid with inode number, i.e. with a quantity used to find the file in question. This identification poses special difficulties for migration of a file system between vendors where assigning the same index to a given file may not be possible. Note here that a fileid is not required to be useful to find the file in question, only that it is unique within the given file system. Servers prepared to accept a fileid as a single piece of metadata and store it apart from the value used to index the file information can relatively easily maintain a fileid value across a migration event, allowing a truly transparent migration event.

In any case, where servers can provide continuity of fileids, they should, and the client should be able to find out that such continuity is available and take appropriate action. Information about the continuity (or lack thereof) of fileids across a file system transition is represented by specifying whether the file systems in question are of the same fileid class.

Note that when consistent fileids do not exist across a transition (either because there is no continuity of fileids or because fileid is not a supported attribute on one of instances involved), and there are no reliable filehandles across a transition event (either because

there is no filehandle continuity or because the filehandles are volatile), the client is in a position where it cannot verify that files it was accessing before the transition are the same objects. It is forced to assume that no object has been renamed, and, unless there are guarantees that provide this (e.g., the file system is read-only), problems for applications may occur. Therefore, use of such configurations should be limited to situations where the problems that this may cause can be tolerated.

#### 17.11.4. Fsid and File System Transitions

Since fsids are generally only unique on a per-server basis, it is likely that they will change during a file system transition. Clients should not make the fsids received from the server visible to applications since they may not be globally unique, and because they may change during a file system transition event. Applications are best served if they are isolated from such transitions to the extent possible.

Although normally a single source file system will transition to a single target file system, there is a provision for splitting a single source file system into multiple target file systems, by specifying the FSLI4F\_MULTI\_FS flag.

##### 17.11.4.1. File System Splitting

When a file system transition is made and the fs\_locations\_info indicates that the file system in question might be split into multiple file systems (via the FSLI4F\_MULTI\_FS flag), the client needs to do GETATTRs to determine the fsid attribute on all known objects within the file system undergoing transition to determine the new file system boundaries.

Clients might choose to maintain the fsids passed to existing applications by mapping all of the fsids for the descendant file systems to the common fsid used for the original file system.

Splitting a file system can be done on a transition between file systems of the same fileid class, since the fact that fileids are unique within the source file system ensure they will be unique in each of the target file systems.

#### 17.11.5. The Change Attribute and File System Transitions

Since the change attribute is defined as a server-specific one, change attributes fetched from one server are normally presumed to be invalid on another server. Such a presumption is troublesome since it would invalidate all cached change attributes, requiring refetching. Even more disruptive, the absence of any assured continuity for the change attribute means that even if the same value is retrieved on refetch, no conclusions can be drawn as to whether the object in question has changed. The identical change attribute could be merely an artifact of a modified file with a different change attribute construction algorithm, with that new algorithm just happening to result in an identical change value.

When the two file systems have consistent change attribute formats, and this fact is communicated to the client by reporting in the same change class, the client may assume a continuity of change attribute construction and handle this situation just as it would be handled without any file system transition.

#### 17.11.6. Write Verifiers and File System Transitions

In a file system transition, the two file systems might be cooperating in the handling of unstably written data. Clients can determine if this is the case by seeing if the two file systems belong to the same write-verifier class. When this is the case, write verifiers returned from one system may be compared to those returned by the other and superfluous writes can be avoided.

When two file systems belong to different write-verifier classes, any verifier generated by one must not be compared to one provided by the other. Instead, the two verifiers should be treated as not equal even when the values are identical.

#### 17.11.7. READDIR Cookies and Verifiers and File System Transitions

In a file system transition, the two file systems might be consistent in their handling of READDIR cookies and verifiers. Clients can determine if this is the case by seeing if the two file systems belong to the same readdir class. When this is the case, readdir class, READDIR cookies, and verifiers from one system will be recognized by the other, and READDIR operations started on one server can be validly continued on the other simply by presenting the cookie and verifier returned by a READDIR operation done on the first file system to the second.

When two file systems belong to different readdir classes, any READDIR cookie and verifier generated by one is not valid on the second and must not be presented to that server by the client. The client should act as if the verifier were rejected.

#### 17.11.8. File System Data and File System Transitions

When multiple replicas exist and are used simultaneously or in succession by a client, applications using them will normally expect that they contain either the same data or data that is consistent with the normal sorts of changes that are made by other clients updating the data of the file system (with metadata being the same to the degree indicated by the `fs_locations_info` attribute). However, when multiple file systems are presented as replicas of one another, the precise relationship between the data of one and the data of another is not, as a general matter, specified by the NFSv4.1 protocol. It is quite possible to present as replicas file systems where the data of those file systems is sufficiently different that some applications have problems dealing with the transition between replicas. The namespace will typically be constructed so that applications can choose an appropriate level of support, so that in one position in the namespace, a varied set of replicas might be listed, while in another, only those that are up-to-date would be considered replicas. The protocol does define three special cases of the relationship among replicas to be specified by the server and relied upon by clients:

- \* When multiple replicas exist and are used simultaneously by a client (See the `FSLIB4_CLSIMUL` definition within `fs_locations_info`), they must designate the same data. Where file systems are writable, a change made on one instance must be visible on all instances at the same time, regardless of whether the interrogated instance is the one on which the modification was done. This allows a client to use these replicas simultaneously without any special adaptation to the fact that there are multiple replicas, beyond adapting to the fact that locks obtained on one replica are maintained separately (i.e., under a different client ID). In this case, locks (whether share reservations or byte-range locks) and delegations obtained on one replica are immediately reflected on all replicas, in the sense that access from all other servers is prevented regardless of the replica used. However, because the servers are not required to treat two associated client IDs as representing the same client, it is best to access each file using only a single client ID.
- \* When one replica is designated as the successor instance to another existing instance after the return of `NFS4ERR_MOVED` (i.e., the case of migration), the client may depend on the fact that all



changes written to stable storage on the original instance are written to stable storage of the successor (uncommitted writes are dealt with in Section 17.11.6 above).

- \* Where a file system is not writable but represents a read-only copy (possibly periodically updated) of a writable file system, clients have similar requirements with regard to the propagation of updates. They may need a guarantee that any change visible on the original file system instance must be immediately visible on any replica before the client transitions access to that replica, in order to avoid any possibility that a client, in effecting a transition to a replica, will see any reversion in file system state. The specific means of this guarantee varies based on the value of the `fss_type` field that is reported as part of the `fs_status` attribute (See Section 17.18). Since these file systems are presumed to be unsuitable for simultaneous use, there is no specification of how locking is handled; in general, locks obtained on one file system will be separate from those on others. Since these are expected to be read-only file systems, this is not likely to pose an issue for clients or applications.

When none of these special situations applies, there is no basis within the protocol for the client to make assumptions about the contents of a replica file system or its relationship to previous file system instances. Thus, switching between nominally identical read-write file systems would not be possible because either the client does not use the `fs_locations_info` attribute, or the server does not support it.

#### 17.11.9. Lock State and File System Transitions

While accessing a file system, clients obtain locks enforced by the server, which may prevent actions by other clients that are inconsistent with those locks.

When access is transferred between replicas, clients need to be assured that the actions disallowed by holding these locks cannot have occurred during the transition. This can be ensured by the methods below. Unless at least one of these is implemented, clients will not be assured of continuity of lock possession across a migration event:

- \* Providing the client an opportunity to re-obtain his locks via a per-fs grace period on the destination server, denying all clients using the destination file system the opportunity to obtain new locks that conflict with those held by the transferred client as long as that client has not completed its per-fs grace period. Because the lock reclaim mechanism was originally defined to

support server reboot, it implicitly assumes that filehandles will, upon reclaim, be the same as those at open. In the case of migration, this requires that source and destination servers use the same filehandles, as evidenced by using the same server scope (See Section 7.4) or by showing this agreement using `fs_locations_info` (See Section 17.11.2 above).

Note that such a grace period can be implemented without interfering with the ability of non-transferred clients to obtain new locks while it is going on. As long as the destination server is aware of the transferred locks, it can distinguish requests to obtain new locks that contrast with existing locks from those that do not, allowing it to treat such client requests without reference to the ongoing grace period.

- \* Locking state can be transferred as part of the transition by providing Transparent State Migration as described in Section 17.12.

Of these, Transparent State Migration provides the smoother experience for clients in that there is no need to go through a reclaim process before new locks can be obtained; however, it requires a greater degree of inter-server coordination. In general, the servers taking part in migration are free to provide either facility. However, when the filehandles can differ across the migration event, Transparent State Migration is the only available means of providing the needed functionality.

It should be noted that these two methods are not mutually exclusive and that a server might well provide both. In particular, if there is some circumstance preventing a specific lock from being transferred transparently, the destination server can allow it to be reclaimed by implementing a per-fs grace period for the migrated file system.

#### 17.11.9.1. Security Issues Related to Reclaiming Lock State after File System Transitions

Although it is possible for a client reclaiming state to misrepresent its state in the same fashion as described in Section 13.4.2.1.1, most implementations providing for such reclamation in the case of file system transitions will have the ability to detect such misrepresentations. This limits the ability of unauthenticated clients to execute denial-of-service attacks in these circumstances. Nevertheless, the rules stated in Section 13.4.2.1.1 regarding principal verification for reclaim requests apply in this situation as well.

Typically, implementations that support file system transitions will have extensive information about the locks to be transferred. This is because of the following:

- \* Since failure is not involved, there is no need to store locking information in persistent storage.
- \* There is no need, as there is in the failure case, to update multiple repositories containing locking state to keep them in sync. Instead, there is a one-time communication of locking state from the source to the destination server.
- \* Providing this information avoids potential interference with existing clients using the destination file system by denying them the ability to obtain new locks during the grace period.

When such detailed locking information, not necessarily including the associated stateids, is available:

- \* It is possible to detect reclaim requests that attempt to reclaim locks that did not exist before the transfer, rejecting them with NFS4ERR\_RECLAIM\_BAD (Section 22.1.9.4). SN
- \* It is possible when dealing with non-reclaim requests, to determine whether they conflict with existing locks, eliminating the need to return NFS4ERR\_GRACE (Section 22.1.9.2) on non-reclaim requests.

It is possible for implementations of grace periods in connection with file system transitions not to have detailed locking information available at the destination server, in which case, the security situation is exactly as described in Section 13.4.2.1.1.

#### 17.11.9.2. Leases and File System Transitions

In the case of lease renewal, the client may not be submitting requests for a file system that has been transferred to another server. This can occur because of the lease renewal mechanism. The client renews the lease associated with all file systems when submitting a request on an associated session, regardless of the specific file system being referenced.

In order for the client to schedule renewal of its lease where there is locking state that may have been relocated to the new server, the client must find out about lease relocation before that lease expire. To accomplish this, the SEQUENCE operation will return the status bit SEQ4\_STATUS\_LEASE\_MOVED if responsibility for any of the renewed locking state has been transferred to a new server. This will continue until the client receives an NFS4ERR\_MOVED error for each of the file systems for which there has been locking state relocation.

When a client receives an SEQ4\_STATUS\_LEASE\_MOVED indication from a server, for each file system of the server for which the client has locking state, the client should perform an operation. For simplicity, the client may choose to reference all file systems, but what is important is that it must reference all file systems for which there was locking state where that state has moved. Once the client receives an NFS4ERR\_MOVED error for each such file system, the server will clear the SEQ4\_STATUS\_LEASE\_MOVED indication. The client can terminate the process of checking file systems once this indication is cleared (but only if the client has received a reply for all outstanding SEQUENCE requests on all sessions it has with the server), since there are no others for which locking state has moved.

A client may use GETATTR of the fs\_status (or fs\_locations\_info) attribute on all of the file systems to get absence indications in a single (or a few) request(s), since absent file systems will not cause an error in this context. However, it still must do an operation that receives NFS4ERR\_MOVED on each file system, in order to clear the SEQ4\_STATUS\_LEASE\_MOVED indication.

Once the set of file systems with transferred locking state has been determined, the client can follow the normal process to obtain the new server information (through the fs\_locations and fs\_locations\_info attributes) and perform renewal of that lease on the new server, unless information in the fs\_locations\_info attribute shows that no state could have been transferred. If the server has not had state transferred to it transparently, the client will receive NFS4ERR\_STALE\_CLIENTID from the new server, as described above, and the client can then reclaim locks as is done in the event of server failure.

#### 17.11.9.3. Transitions and the Lease\_time Attribute

In order that the client may appropriately manage its lease in the case of a file system transition, the destination server must establish proper values for the lease\_time attribute.

When state is transferred transparently, that state should include the correct value of the `lease_time` attribute. The `lease_time` attribute on the destination server must never be less than that on the source, since this would result in premature expiration of a lease granted by the source server. Upon transitions in which state is transferred transparently, the client is under no obligation to refetch the `lease_time` attribute and may continue to use the value previously fetched (on the source server).

If state has not been transferred transparently, either because the associated servers are shown as having different `eir_server_scope` strings or because the client ID is rejected when presented to the new server, the client should fetch the value of `lease_time` on the new (i.e., destination) server, and use it for subsequent locking requests. However, the server must respect a grace period of at least as long as the `lease_time` on the source server, in order to ensure that clients have ample time to reclaim their lock before potentially conflicting non-reclaimed locks are granted.

#### 17.12. Transferring State upon Migration

When the transition is a result of a server-initiated decision to transition access, and the source and destination servers have implemented appropriate cooperation, it is possible to do the following:

- \* Transfer locking state from the source to the destination server in a fashion similar to that provided by Transparent State Migration in NFSv4.0, as described in [RFC7931]. Server responsibilities are described in Section 17.14.2.
- \* Transfer session state from the source to the destination server. Server responsibilities in effecting such a transfer are described in Section 17.14.3.

The means by which the client determines which of these transfer events has occurred are described in Section 17.13.

##### 17.12.1. Transparent State Migration and pNFS

When pNFS is involved, the protocol is capable of supporting:

- \* Migration of the Metadata Server (MDS), leaving the Data Servers (DSs) in place.
- \* Migration of the file system as a whole, including the MDS and associated DSs.

- \* Replacement of one DS by another.
- \* Migration of a pNFS file system to one in which pNFS is not used.
- \* Migration of a file system not using pNFS to one in which layouts are available.

Note that migration, per se, is only involved in the transfer of the MDS function. Although the servicing of a layout may be transferred from one data server to another, this not done using the file system location attributes. The MDS can effect such transfers by recalling or revoking existing layouts and granting new ones on a different data server.

Migration of the MDS function is directly supported by Transparent State Migration. Layout state will normally be transparently transferred, just as other state is. As a result, Transparent State Migration provides a framework in which, given appropriate inter-MDS data transfer, one MDS can be substituted for another.

Migration of the file system function as a whole can be accomplished by recalling all layouts as part of the initial phase of the migration process. As a result, I/O will be done through the MDS during the migration process, and new layouts can be granted once the client is interacting with the new MDS. An MDS can also effect this sort of transition by revoking all layouts as part of Transparent State Migration, as long as the client is notified about the loss of locking state.

In order to allow migration to a file system on which pNFS is not supported, clients need to be prepared for a situation in which layouts are not available or supported on the destination file system and so direct I/O requests to the destination server, rather than depending on layouts being available.

Replacement of one DS by another is not addressed by migration as such but can be effected by an MDS recalling layouts for the DS to be replaced and issuing new ones to be served by the successor DS.

Migration may transfer a file system from a server that does not support pNFS to one that does. In order to properly adapt to this situation, clients that support pNFS, but function adequately in its absence, should check for pNFS support when a file system is migrated and be prepared to use pNFS when support is available on the destination.

### 17.13. Client Responsibilities When Access Is Transitioned

For a client to respond to an access transition, it must become aware of it. The ways in which this can happen are discussed in Section 17.13.1, which discusses indications that a specific file system access path has transitioned as well as situations in which additional activity is necessary to determine the set of file systems that have been migrated. Section 17.13.2 goes on to complete the discussion of how the set of migrated file systems might be determined. Sections 17.13.3 through 17.13.5 discuss how the client should deal with each transition it becomes aware of, either directly or as a result of migration discovery.

The following terms are used to describe client activities:

- \* "Transition recovery" refers to the process of restoring access to a file system on which NFS4ERR\_MOVED was received.
- \* "Migration recovery" refers to that subset of transition recovery that applies when the file system has migrated to a different replica.
- \* "Migration discovery" refers to the process of determining which file system(s) have been migrated. It is necessary to avoid a situation in which leases could expire when a file system is not accessed for a long period of time, since a client unaware of the migration might be referencing an unmigrated file system and not renewing the lease associated with the migrated file system.

#### 17.13.1. Client Transition Notifications

When there is a change in the network access path that a client is to use to access a file system, there are a number of related status indications with which clients need to deal:

- \* If an attempt is made to use or return a filehandle within a file system that is no longer accessible at the address previously used to access it, the error NFS4ERR\_MOVED is returned.

Exceptions are made to allow such filehandles to be used when interrogating a file system location attribute. This enables a client to determine a new replica's location or a new network access path.

This condition continues on subsequent attempts to access the file system in question. The only way the client can avoid the error is to cease accessing the file system in question at its old server location and access it instead using a different address at which it is now available.

- \* Whenever a client sends a SEQUENCE operation to a server that generated state held on that client and associated with a file system no longer accessible on that server, the response will contain the status bit SEQ4\_STATUS\_LEASE\_MOVED, indicating that there has been a lease migration.

This condition continues until the client acknowledges the notification by fetching a file system location attribute for the file system whose network access path is being changed. When there are multiple such file systems, a location attribute for each such file system needs to be fetched. The location attribute for all migrated file systems needs to be fetched in order to clear the condition. Even after the condition is cleared, the client needs to respond by using the location information to access the file system at its new location to ensure that leases are not needlessly expired.

Unlike NFSv4.0, in which the corresponding conditions are both errors and thus mutually exclusive, in NFSv4.1 the client can, and often will, receive both indications on the same request. As a result, implementations need to address the question of how to coordinate the necessary recovery actions when both indications arrive in the response to the same request. It should be noted that when processing an NFSv4 COMPOUND, the server will normally decide whether SEQ4\_STATUS\_LEASE\_MOVED is to be set before it determines which file system will be referenced or whether NFS4ERR\_MOVED is to be returned.

Since these indications are not mutually exclusive in NFSv4.1, the following combinations are possible results when a COMPOUND is issued:

- \* The COMPOUND status is NFS4ERR\_MOVED, and SEQ4\_STATUS\_LEASE\_MOVED is asserted.

In this case, transition recovery is required. While it is possible that migration discovery is needed in addition, it is likely that only the accessed file system has transitioned. In any case, because addressing NFS4ERR\_MOVED is necessary to allow the rejected requests to be processed on the target, dealing with it will typically have priority over migration discovery.



- \* The COMPOUND status is NFS4ERR\_MOVED, and SEQ4\_STATUS\_LEASE\_MOVED is clear.

In this case, transition recovery is also required. It is clear that migration discovery is not needed to find file systems that have been migrated other than the one returning NFS4ERR\_MOVED. Cases in which this result can arise include a referral or a migration for which there is no associated locking state. This can also arise in cases in which an access path transition other than migration occurs within the same server. In such a case, there is no need to set SEQ4\_STATUS\_LEASE\_MOVED, since the lease remains associated with the current server even though the access path has changed.

- \* The COMPOUND status is not NFS4ERR\_MOVED, and SEQ4\_STATUS\_LEASE\_MOVED is asserted.

In this case, no transition recovery activity is required on the file system(s) accessed by the request. However, to prevent avoidable lease expiration, migration discovery needs to be done.

- \* The COMPOUND status is not NFS4ERR\_MOVED, and SEQ4\_STATUS\_LEASE\_MOVED is clear.

In this case, neither transition-related activity nor migration discovery is required.

Note that the specified actions only need to be taken if they are not already going on. For example, when NFS4ERR\_MOVED is received while accessing a file system for which transition recovery is already occurring, the client merely waits for that recovery to be completed, while the receipt of the SEQ4\_STATUS\_LEASE\_MOVED indication only needs to initiate migration discovery for a server if such discovery is not already underway for that server.

The fact that a lease-migrated condition does not result in an error in NFSv4.1 has a number of important consequences. In addition to the fact that the two indications are not mutually exclusive, as discussed above, there are number of issues that are important in considering implementation of migration discovery, as discussed in Section 17.13.2.

Because SEQ4\_STATUS\_LEASE\_MOVED is not an error condition, it is possible for file systems whose access paths have not changed to be successfully accessed on a given server even though recovery is necessary for other file systems on the same server. As a result, access can take place while:

- \* The migration discovery process is happening for that server.
- \* The transition recovery process is happening for other file systems connected to that server.

#### 17.13.2. Performing Migration Discovery

Migration discovery can be performed in the same context as transition recovery, allowing recovery for each migrated file system to be invoked as it is discovered. Alternatively, it may be done in a separate migration discovery thread, allowing migration discovery to be done in parallel with one or more instances of transition recovery.

In either case, because the lease-migrated indication does not result in an error, other access to file systems on the server can proceed normally, with the possibility that further such indications will be received, raising the issue of how such indications are to be dealt with. In general:

- \* No action needs to be taken for such indications received by any threads performing migration discovery, since continuation of that work will address the issue.
- \* In other cases in which migration discovery is currently being performed, nothing further needs to be done to respond to such lease migration indications, as long as one can be certain that the migration discovery process would deal with those indications. See below for details.
- \* For such indications received in all other contexts, the appropriate response is to initiate or otherwise provide for the execution of migration discovery for file systems associated with the server IP address returning the indication.

This leaves a potential difficulty in situations in which the migration discovery process is near to completion but is still operating. One should not ignore a SEQ4\_STATUS\_LEASE\_MOVED indication if the migration discovery process is not able to respond to the discovery of additional migrating file systems without additional aid. A further complexity relevant in addressing such situations is that a lease-migrated indication may reflect the server's state at the time the SEQUENCE operation was processed, which may be different from that in effect at the time the response is received. Because new migration events may occur at any time, and because a SEQ4\_STATUS\_LEASE\_MOVED indication may reflect the situation in effect a considerable time before the indication is received, special care needs to be taken to ensure that SEQ4\_STATUS\_LEASE\_MOVED indications are not inappropriately ignored.

A useful approach to this issue involves the use of separate externally-visible migration discovery states for each server. Separate values could represent the various possible states for the migration discovery process for a server:

- \* Non-operation, in which migration discovery is not being performed.
- \* Normal operation, in which there is an ongoing scan for migrated file systems.
- \* Completion/verification of migration discovery processing, in which the possible completion of migration discovery processing needs to be verified.

Given that framework, migration discovery processing would proceed as follows:

- \* While in the normal-operation state, the thread performing discovery would fetch, for successive file systems known to the client on the server being worked on, a file system location attribute plus the fs\_status attribute.
- \* If the fs\_status attribute indicates that the file system is a migrated one (i.e., fss\_absent is true, and fss\_type != STATUS4\_REFERRAL), then a migrated file system has been found. In this situation, it is likely that the fetch of the file system location attribute has cleared one of the file systems contributing to the lease-migrated indication.

- \* In cases in which that happened, the thread cannot know whether the lease-migrated indication has been cleared, and so it enters the completion/verification state and proceeds to issue a COMPOUND to see if the SEQ4\_STATUS\_LEASE\_MOVED indication has been cleared.
- \* When the discovery process is in the completion/verification state, if other requests get a lease-migrated indication, they note that it was received. Later, the existence of such indications is used when the request completes, as described below.

When the request used in the completion/verification state completes:

- \* If a lease-migrated indication is returned, the discovery continues normally. Note that this is so even if all file systems have been traversed, since new migrations could have occurred while the process was going on.
- \* Otherwise, if there is any record that other requests saw a lease-migrated indication while the request was occurring, that record is cleared, and the verification request is retried. The discovery process remains in the completion/verification state.
- \* If there have been no lease-migrated indications, the work of migration discovery is considered completed, and it enters the non-operating state. Once it enters this state, subsequent lease-migrated indications will trigger a new migration discovery process.

It should be noted that the process described above is not guaranteed to terminate, as a long series of new migration events might continually delay the clearing of the SEQ4\_STATUS\_LEASE\_MOVED indication. To prevent unnecessary lease expiration, it is appropriate for clients to use the discovery of migrations to effect lease renewal immediately, rather than waiting for the clearing of the SEQ4\_STATUS\_LEASE\_MOVED indication when the complete set of migrations is available.

Lease discovery needs to be provided as described above. This ensures that the client discovers file system migrations soon enough to renew its leases on each destination server before they expire. Non-renewal of leases can lead to loss of locking state. While the consequences of such loss can be ameliorated through implementations of courtesy locks, servers are under no obligation to do so, and a conflicting lock request may mean that a lock is revoked unexpectedly. Clients should be aware of this possibility.

### 17.13.3. Overview of Client Response to NFS4ERR\_MOVED

This section outlines a way in which a client that receives NFS4ERR\_MOVED can effect transition recovery by using a new server or server endpoint if one is available. As part of that process, it will determine:

- \* Whether the NFS4ERR\_MOVED indicates migration has occurred, or whether it indicates another sort of file system access transition as discussed in Section 17.10 above.
- \* In the case of migration, whether Transparent State Migration has occurred.
- \* Whether any state has been lost during the process of Transparent State Migration.
- \* Whether sessions have been transferred as part of Transparent State Migration.

During the first phase of this process, the client proceeds to examine file system location entries to find the initial network address it will use to continue access to the file system or its replacement. For each location entry that the client examines, the process consists of five steps:

1. Performing an EXCHANGE\_ID directed at the location address. This operation is used to register the client owner (in the form of a client\_owner4) with the server, to obtain a client ID to be used subsequently to communicate with it, to obtain that client ID's confirmation status, and to determine server\_owner4 and scope for the purpose of determining if the entry is trunkable with the address previously being used to access the file system (i.e., that it represents another network access path to the same file system and can share locking state with it).
2. Making an initial determination of whether migration has occurred. The initial determination will be based on whether the EXCHANGE\_ID results indicate that the current location element is server-trunkable with that used to access the file system when access was terminated by receiving NFS4ERR\_MOVED. If it is, then migration has not occurred. In that case, the transition is dealt with, at least initially, as one involving continued access to the same file system on the same server through a new network address.

3. Obtaining access to existing session state or creating new sessions. How this is done depends on the initial determination of whether migration has occurred and can be done as described in Section 17.13.4 below in the case of migration or as described in Section 17.13.5 below in the case of a network address transfer without migration.
4. Verifying the trunking relationship assumed in step 2 as discussed in Section 7.5.1. Although this step will generally confirm the initial determination, it is possible for verification to invalidate the initial determination of network address shift (without migration) and instead determine that migration had occurred. There is no need to redo step 3 above, since it will be possible to continue use of the session established already.
5. Obtaining access to existing locking state and/or re-obtaining it. How this is done depends on the final determination of whether migration has occurred and can be done as described below in Section 17.13.4 in the case of migration or as described in Section 17.13.5 in the case of a network address transfer without migration.

Once the initial address has been determined, clients are free to apply an abbreviated process to find additional addresses trunkable with it (clients may seek session-trunkable or server-trunkable addresses depending on whether they support client ID trunking). During this later phase of the process, further location entries are examined using the abbreviated procedure specified below:

- A: Before the EXCHANGE\_ID, the fs name of the location entry is examined, and if it does not match that currently being used, the entry is ignored. Otherwise, one proceeds as specified by step 1 above.
- B: In the case that the network address is session-trunkable with one used previously, a BIND\_CONN\_TO\_SESSION is used to access that session using the new network address. Otherwise, or if the bind operation fails, a CREATE\_SESSION is done.
- C: The verification procedure referred to in step 4 above is used. However, if it fails, the entry is ignored and the next available entry is used.

#### 17.13.4. Obtaining Access to Sessions and State after Migration

In the event that migration has occurred, migration recovery will involve determining whether Transparent State Migration has occurred. This decision is made based on the client ID returned by the EXCHANGE\_ID and the reported confirmation status.

- \* If the client ID is an unconfirmed client ID not previously known to the client, then Transparent State Migration has not occurred.
- \* If the client ID is a confirmed client ID previously known to the client, then any transferred state would have been merged with an existing client ID representing the client to the destination server. In this state merger case, Transparent State Migration might or might not have occurred, and a determination as to whether it has occurred is deferred until sessions are established and the client is ready to begin state recovery.
- \* If the client ID is a confirmed client ID not previously known to the client, then the client can conclude that the client ID was transferred as part of Transparent State Migration. In this transferred client ID case, Transparent State Migration has occurred, although some state might have been lost.

Once the client ID has been obtained, it is necessary to obtain access to sessions to continue communication with the new server. In any of the cases in which Transparent State Migration has occurred, it is possible that a session was transferred as well. To deal with that possibility, clients can, after doing the EXCHANGE\_ID, issue a BIND\_CONN\_TO\_SESSION to connect the transferred session to a connection to the new server. If that fails, it is an indication that the session was not transferred and that a new session needs to be created to take its place.

In some situations, it is possible for a BIND\_CONN\_TO\_SESSION to succeed without session migration having occurred. If state merger has taken place, then the associated client ID may have already had a set of existing sessions, with it being possible that the session ID of a given session is the same as one that might have been migrated. In that event, a BIND\_CONN\_TO\_SESSION might succeed, even though there could have been no migration of the session with that session ID. In such cases, the client will receive sequence errors when the slot sequence values used are not appropriate on the new session. When this occurs, the client can create a new a session and cease using the existing one.

Once the client has determined the initial migration status, and determined that there was a shift to a new server, it needs to re-establish its locking state, if possible. To enable this to happen without loss of the guarantees normally provided by locking, the destination server needs to implement a per-fs grace period in all cases in which lock state was lost, including those in which Transparent State Migration was not implemented. Each client for which there was a transfer of locking state to the new server will have the duration of the grace period to reclaim its locks, from the time its locks were transferred.

Clients need to deal with the following cases:

- \* In the state merger case, it is possible that the server has not attempted Transparent State Migration, in which case state may have been lost without it being reflected in the SEQ4\_STATUS bits. To determine whether this has happened, the client can use TEST\_STATEID to check whether the stateids created on the source server are still accessible on the destination server. Once a single stateid is found to have been successfully transferred, the client can conclude that Transparent State Migration was begun, and any failure to transport all of the stateids will be reflected in the SEQ4\_STATUS bits. Otherwise, Transparent State Migration has not occurred.
- \* In a case in which Transparent State Migration has not occurred, the client can use the per-fs grace period provided by the destination server to reclaim locks that were held on the source server.
- \* In a case in which Transparent State Migration has occurred, and no lock state was lost (as shown by SEQ4\_STATUS flags), no lock reclaim is necessary.
- \* In a case in which Transparent State Migration has occurred, and some lock state was lost (as shown by SEQ4\_STATUS flags), existing stateids need to be checked for validity using TEST\_STATEID, and reclaim used to re-establish any that were not transferred.

For all of the cases above, RECLAIM\_COMPLETE with an rca\_one\_fs value of TRUE needs to be done before normal use of the file system, including obtaining new locks for the file system. This applies even if no locks were lost and there was no need for any to be reclaimed.



#### 17.13.5. Obtaining Access to Sessions and State after Network Address Transfer

The case in which there is a transfer to a new network address without migration is similar to that described in Section 17.13.4 above in that there is a need to obtain access to needed sessions and locking state. However, the details are simpler and will vary depending on the type of trunking between the address receiving NFS4ERR\_MOVED and that to which the transfer is to be made.

To make a session available for use, a BIND\_CONN\_TO\_SESSION should be used to obtain access to the session previously in use. Only if this fails, should a CREATE\_SESSION be done. While this procedure mirrors that in Section 17.13.4 above, there is an important difference in that preservation of the session is not purely optional but depends on the type of trunking.

Access to appropriate locking state will generally need no actions beyond access to the session. However, the SEQ4\_STATUS bits need to be checked for lost locking state, including the need to reclaim locks after a server reboot, since there is always a possibility of locking state being lost.

#### 17.14. Server Responsibilities Upon Migration

In the event of file system migration, when the client connects to the destination server, that server needs to be able to provide the client continued access to the files it had open on the source server. There are two ways to provide this:

- \* By provision of an fs-specific grace period, allowing the client the ability to reclaim its locks, in a fashion similar to what would have been done in the case of recovery from a server restart. See Section 17.14.1 for a more complete discussion.
- \* By implementing Transparent State Migration possibly in connection with session migration, the server can provide the client immediate access to the state built up on the source server on the destination server.

These features are discussed separately in Sections 17.14.2 and 17.14.3, which discuss Transparent State Migration and session migration, respectively.

All the features described above can involve transfer of lock-related information between source and destination servers. In some cases, this transfer is a necessary part of the implementation, while in other cases, it is a helpful implementation aid, which servers might

or might not use. The subsections below discuss the information that would be transferred but do not define the specifics of the transfer protocol. This is left as an implementation choice, although standards in this area could be developed at a later time.

#### 17.14.1. Server Responsibilities in Effecting State Reclaim after Migration

In this case, the destination server needs no knowledge of the locks held on the source server. It relies on the clients to accurately report (via reclaim operations) the locks previously held, and does not allow new locks to be granted on migrated file systems until the grace period expires. Disallowing of new locks applies to all clients accessing these file systems, while grace period expiration occurs for each migrated client independently.

During this grace period, clients have the opportunity to use reclaim operations to obtain locks for file system objects within the migrated file system, in the same way that they do when recovering from server restart, and the servers typically rely on clients to accurately report their locks, although they have the option of subjecting these requests to verification. If the clients only reclaim locks held on the source server, no conflict can arise. Once the client has reclaimed its locks, it indicates the completion of lock reclamation by performing a RECLAIM\_COMPLETE specifying `rca_one_fs` as TRUE.

While it is not necessary for source and destination servers to cooperate to transfer information about locks, implementations are well advised to consider transferring the following useful information:

- \* If information about the set of clients that have locking state for the transferred file system is made available, the destination server will be able to terminate the grace period once all such clients have reclaimed their locks, allowing normal locking activity to resume earlier than it would have otherwise.
- \* Locking summary information for individual clients (at various possible levels of detail) can detect some instances in which clients do not accurately represent the locks held on the source server.

#### 17.14.2. Server Responsibilities in Effecting Transparent State Migration

The basic responsibility of the source server in effecting Transparent State Migration is to make available to the destination server a description of each piece of locking state associated with the file system being migrated. In addition to client id string and verifier, the source server needs to provide for each stateid:

- \* The stateid including the current sequence value.
- \* The associated client ID.
- \* The handle of the associated file.
- \* The type of the lock, such as open, byte-range lock, delegation, or layout.
- \* For locks such as opens and byte-range locks, there will be information about the owner(s) of the lock.
- \* For recallable/revocable lock types, the current recall status needs to be included.
- \* For each lock type, there will be associated type-specific information. For opens, this will include share and deny mode while for byte-range locks and layouts, there will be a type and a byte-range.

Such information will most probably be organized by client id string on the destination server so that it can be used to provide appropriate context to each client when it makes itself known to the client. Issues connected with a client impersonating another by presenting another client's client id string can be addressed using NFSv4.1 state protection features, as described in Section 28.

A further server responsibility concerns locks that are revoked or otherwise lost during the process of file system migration. Because locks that appear to be lost during the process of migration will be reclaimed by the client, the servers have to take steps to ensure that locks revoked soon before or soon after migration are not inadvertently allowed to be reclaimed in situations in which the continuity of lock possession cannot be assured.

- \* For locks lost on the source but whose loss has not yet been acknowledged by the client (by using `FREE_STATEID`), the destination must be aware of this loss so that it can deny a request to reclaim them.

- \* For locks lost on the destination after the state transfer but before the client's RECLAIM\_COMPLETE is done, the destination server should note these and not allow them to be reclaimed.

An additional responsibility of the cooperating servers concerns situations in which a stateid cannot be transferred transparently because it conflicts with an existing stateid held by the client and associated with a different file system. In this case, there are two valid choices:

- \* Treat the transfer, as in NFSv4.0, as one without Transparent State Migration. In this case, conflicting locks cannot be granted until the client does a RECLAIM\_COMPLETE, after reclaiming the locks it had, with the exception of reclaims denied because they were attempts to reclaim locks that had been lost.
- \* Implement Transparent State Migration, except for the lock with the conflicting stateid. In this case, the client will be aware of a lost lock (through the SEQ4\_STATUS flags) and be allowed to reclaim it.

When transferring state between the source and destination, the issues discussed in Section 7.2 of [RFC7931] must still be attended to. In this case, the use of NFS4ERR\_DELAY may still be necessary in NFSv4.1, as it was in NFSv4.0, to prevent locking state changing while it is being transferred. See Section 22.1.1.3 for information about appropriate client retry approaches in the event that NFS4ERR\_DELAY is returned.

There are a number of important differences in the NFS4.1 context:

- \* The absence of RELEASE\_LOCKOWNER means that the one case in which an operation could not be deferred by use of NFS4ERR\_DELAY no longer exists.
- \* Sequencing of operations is no longer done using owner-based operation sequences numbers. Instead, sequencing is session-based.

As a result, when sessions are not transferred, the techniques discussed in Section 7.2 of [RFC7931] are adequate and will not be further discussed.

### 17.14.3. Server Responsibilities in Effecting Session Transfer

The basic responsibility of the source server in effecting session transfer is to make available to the destination server a description of the current state of each slot with the session, including the following:

- \* The last sequence value received for that slot.
- \* Whether there is cached reply data for the last request executed and, if so, the cached reply.

When sessions are transferred, there are a number of issues that pose challenges in terms of making the transferred state unmodifiable during the period it is gathered up and transferred to the destination server:

- \* A single session may be used to access multiple file systems, not all of which are being transferred.
- \* Requests made on a session may, even if rejected, affect the state of the session by advancing the sequence number associated with the slot used.

As a result, when the file system state might otherwise be considered unmodifiable, the client might have any number of in-flight requests, each of which is capable of changing session state, which may be of a number of types:

1. Those requests that were processed on the migrating file system before migration began.
2. Those requests that received the error NFS4ERR\_DELAY because the file system being accessed was in the process of being migrated.
3. Those requests that received the error NFS4ERR\_MOVED because the file system being accessed had been migrated.
4. Those requests that accessed the migrating file system in order to obtain location or status information.
5. Those requests that did not reference the migrating file system.

It should be noted that the history of any particular slot is likely to include a number of these request classes. In the case in which a session that is migrated is used by file systems other than the one migrated, requests of class 5 may be common and may be the last request processed for many slots.

Since session state can change even after the locking state has been fixed as part of the migration process, the session state known to the client could be different from that on the destination server, which necessarily reflects the session state on the source server at an earlier time. In deciding how to deal with this situation, it is helpful to distinguish between two sorts of behavioral consequences of the choice of initial sequence ID values:

- \* The error NFS4ERR\_SEQ\_MISORDERED is returned when the sequence ID in a request is neither equal to the last one seen for the current slot nor the next greater one.

In view of the difficulty of arriving at a mutually acceptable value for the correct last sequence value at the point of migration, it may be necessary for the server to show some degree of forbearance when the sequence ID is one that would be considered unacceptable if session migration were not involved.

- \* Returning the cached reply for a previously executed request when the sequence ID in the request matches the last value recorded for the slot.

In the cases in which an error is returned and there is no possibility of any non-idempotent operation having been executed, it may not be necessary to adhere to this as strictly as might be proper if session migration were not involved. For example, the fact that the error NFS4ERR\_DELAY was returned may not assist the client in any material way, while the fact that NFS4ERR\_MOVED was returned by the source server may not be relevant when the request was reissued and directed to the destination server.

An important issue is that the specification needs to take note of all potential COMPOUNDS, even if they might be unlikely in practice. For example, a COMPOUND is allowed to access multiple file systems and might perform non-idempotent operations in some of them before accessing a file system being migrated. Also, a COMPOUND may return considerable data in the response before being rejected with NFS4ERR\_DELAY or NFS4ERR\_MOVED, and may in addition be marked as sa\_cachethis. However, note that if the client and server adhere to rules in Section 22.1.1.3, there is no possibility of non-idempotent operations being spuriously reissued after receiving NFS4ERR\_DELAY response.

To address these issues, a destination server MAY do any of the following when implementing session transfer:

- \* Avoid enforcing any sequencing semantics for a particular slot until the client has established the starting sequence for that slot on the destination server.
- \* For each slot, avoid returning a cached reply returning NFS4ERR\_DELAY or NFS4ERR\_MOVED until the client has established the starting sequence for that slot on the destination server.
- \* Until the client has established the starting sequence for a particular slot on the destination server, avoid reporting NFS4ERR\_SEQ\_MISORDERED or returning a cached reply that contains either NFS4ERR\_DELAY or NFS4ERR\_MOVED and consists solely of a series of operations where the response is NFS4\_OK until the final error.

Because of the considerations mentioned above, including the rules for the handling of NFS4ERR\_DELAY included in Section 22.1.1.3, the destination server can respond appropriately to SEQUENCE operations received from the client by adopting the three policies listed below:

- \* Not responding with NFS4ERR\_SEQ\_MISORDERED for the initial request on a slot within a transferred session because the destination server cannot be aware of requests made by the client after the server handoff but before the client became aware of the shift. In cases in which NFS4ERR\_SEQ\_MISORDERED would normally have been reported, the request is to be processed normally as a new request.
- \* Replying as it would for a retry whenever the sequence matches that transferred by the source server, even though this would not provide retry handling for requests issued after the server handoff, under the assumption that, when such requests are issued, they will never be responded to in a state-changing fashion, making retry support for them unnecessary.
- \* Once a non-retry SEQUENCE is received for a given slot, using that as the basis for further sequence checking, with no further reference to the sequence value transferred by the source server.

#### 17.15. Effecting File System Referrals

Referrals are effected when an absent file system is encountered and one or more alternate locations are made available by the fs\_locations or fs\_locations\_info attributes. The client will typically get an NFS4ERR\_MOVED error, fetch the appropriate location information, and proceed to access the file system on a different server, even though it retains its logical position within the original namespace. Referrals differ from migration events in that

they happen only when the client has not previously referenced the file system in question (so there is nothing to transition). Referrals can only come into effect when an absent file system is encountered at its root.

The examples given in the sections below are somewhat artificial in that an actual client will not typically do a multi-component look up, but will have cached information regarding the upper levels of the name hierarchy. However, these examples are chosen to make the required behavior clear and easy to put within the scope of a small number of requests, without getting into a discussion of the details of how specific clients might choose to cache things.

#### 17.15.1. Referral Example (LOOKUP)

Let us suppose that the following COMPOUND is sent in an environment in which /this/is/the/path is absent from the target server. This may be for a number of reasons. It may be that the file system has moved, or it may be that the target server is functioning mainly, or solely, to refer clients to the servers on which various file systems are located.

```
* PUTROOTFH
* LOOKUP "this"
* LOOKUP "is"
* LOOKUP "the"
* LOOKUP "path"
* GETFH
* GETATTR (fsid, fileid, size, time_modify)
```

Under the given circumstances, the following will be the result.

```
* PUTROOTFH --> NFS_OK.  The current fh is now the root of the
  pseudo-fs.
* LOOKUP "this" --> NFS_OK.  The current fh is for /this and is
  within the pseudo-fs.
* LOOKUP "is" --> NFS_OK.  The current fh is for /this/is and is
  within the pseudo-fs.
```



- \* LOOKUP "the" --> NFS\_OK. The current fh is for /this/is/the and is within the pseudo-fs.
- \* LOOKUP "path" --> NFS\_OK. The current fh is for /this/is/the/path and is within a new, absent file system, but ... the client will never see the value of that fh.
- \* GETFH --> NFS4ERR\_MOVED. Fails because current fh is in an absent file system at the start of the operation, and the specification makes no exception for GETFH.
- \* GETATTR (fsid, fileid, size, time\_modify). Not executed because the failure of the GETFH stops processing of the COMPOUND.

Given the failure of the GETFH, the client has the job of determining the root of the absent file system and where to find that file system, i.e., the server and path relative to that server's root fh. Note that in this example, the client did not obtain filehandles and attribute information (e.g., fsid) for the intermediate directories, so that it would not be sure where the absent file system starts. It could be the case, for example, that /this/is/the is the root of the moved file system and that the reason that the look up of "path" succeeded is that the file system was not absent on that operation but was moved between the last LOOKUP and the GETFH (since COMPOUND is not atomic). Even if we had the fsids for all of the intermediate directories, we could have no way of knowing that /this/is/the/path was the root of a new file system, since we don't yet have its fsid.

In order to get the necessary information, let us re-send the chain of LOOKUPS with GETFHs and GETATTRs to at least get the fsids so we can be sure where the appropriate file system boundaries are. The client could choose to get fs\_locations\_info at the same time but in most cases the client will have a good guess as to where file system boundaries are (because of where NFS4ERR\_MOVED was, and was not, received) making fetching of fs\_locations\_info unnecessary.

OP01: PUTROOTFH --> NFS\_OK

- \* Current fh is root of pseudo-fs.

OP02: GETATTR(fsid) --> NFS\_OK

- \* Just for completeness. Normally, clients will know the fsid of the pseudo-fs as soon as they establish communication with a server.

OP03: LOOKUP "this" --> NFS\_OK

OP04: GETATTR(fsid) --> NFS\_OK

- \* Get current fsid to see where file system boundaries are. The fsid will be that for the pseudo-fs in this example, so no boundary.

OP05: GETFH --> NFS\_OK

- \* Current fh is for /this and is within pseudo-fs.

OP06: LOOKUP "is" --> NFS\_OK

- \* Current fh is for /this/is and is within pseudo-fs.

OP07: GETATTR(fsid) --> NFS\_OK

- \* Get current fsid to see where file system boundaries are. The fsid will be that for the pseudo-fs in this example, so no boundary.

OP08: GETFH --> NFS\_OK

- \* Current fh is for /this/is and is within pseudo-fs.

OP09: LOOKUP "the" --> NFS\_OK

- \* Current fh is for /this/is/the and is within pseudo-fs.

OP10: GETATTR(fsid) --> NFS\_OK

- \* Get current fsid to see where file system boundaries are. The fsid will be that for the pseudo-fs in this example, so no boundary.

OP11: GETFH --> NFS\_OK

- \* Current fh is for /this/is/the and is within pseudo-fs.

OP12: LOOKUP "path" --> NFS\_OK

- \* Current fh is for /this/is/the/path and is within a new, absent file system, but ...
- \* The client will never see the value of that fh.

OP13: GETATTR(fsid, fs\_locations\_info) --> NFS\_OK

- \* We are getting the fsid to know where the file system boundaries are. In this operation, the fsid will be different than that of the parent directory (which in turn was retrieved in OP10). Note that the fsid we are given will not necessarily be preserved at the new location. That fsid might be different, and in fact the fsid we have for this file system might be a valid fsid of a different file system on that new server.
- \* In this particular case, we are pretty sure anyway that what has moved is /this/is/the/path rather than /this/is/the since we have the fsid of the latter and it is that of the pseudo-fs, which presumably cannot move. However, in other examples, we might not have this kind of information to rely on (e.g., /this/is/the might be a non-pseudo file system separate from /this/is/the/path), so we need to have other reliable source information on the boundary of the file system that is moved. If, for example, the file system /this/is had moved, we would have a case of migration rather than referral, and once the boundaries of the migrated file system was clear we could fetch fs\_locations\_info.
- \* We are fetching fs\_locations\_info because the fact that we got an NFS4ERR\_MOVED at this point means that it is most likely that this is a referral and we need the destination. Even if it is the case that /this/is/the is a file system that has migrated, we will still need the location information for that file system.

OP14: GETFH --> NFS4ERR\_MOVED

- \* Fails because current fh is in an absent file system at the start of the operation, and the specification makes no exception for GETFH. Note that this means the server will never send the client a filehandle from within an absent file system.

Given the above, the client knows where the root of the absent file system is (/this/is/the/path) by noting where the change of fsid occurred (between "the" and "path"). The fs\_locations\_info attribute also gives the client the actual location of the absent file system, so that the referral can proceed. The server gives the client the bare minimum of information about the absent file system so that there will be very little scope for problems of conflict between information sent by the referring server and information of the file system's home. No filehandles and very few attributes are present on the referring server, and the client can treat those it receives as transient information with the function of enabling the referral.

## 17.15.2. Referral Example (READDIR)

Another context in which a client may encounter referrals is when it does a READDIR on a directory in which some of the sub-directories are the roots of absent file systems.

Suppose such a directory is read as follows:

```
* PUTROOTFH
* LOOKUP "this"
* LOOKUP "is"
* LOOKUP "the"
* READDIR (fsid, size, time_modify, mounted_on_fileid)
```

In this case, because `rdattr_error` is not requested, `fs_locations_info` is not requested, and some of the attributes cannot be provided, the result will be an `NFS4ERR_MOVED` error on the READDIR, with the detailed results as follows:

```
* PUTROOTFH --> NFS_OK.  The current fh is at the root of the
  pseudo-fs.
* LOOKUP "this" --> NFS_OK.  The current fh is for /this and is
  within the pseudo-fs.
* LOOKUP "is" --> NFS_OK.  The current fh is for /this/is and is
  within the pseudo-fs.
* LOOKUP "the" --> NFS_OK.  The current fh is for /this/is/the and
  is within the pseudo-fs.
* READDIR (fsid, size, time_modify, mounted_on_fileid) -->
  NFS4ERR_MOVED.  Note that the same error would have been returned
  if /this/is/the had migrated, but it is returned because the
  directory contains the root of an absent file system.
```

So now suppose that we re-send with `rdattr_error`:

```
* PUTROOTFH
* LOOKUP "this"
* LOOKUP "is"
```

- \* LOOKUP "the"
- \* READDIR (rdattr\_error, fsid, size, time\_modify, mounted\_on\_fileid)

The results will be:

- \* PUTROOTFH --> NFS\_OK. The current fh is at the root of the pseudo-fs.
- \* LOOKUP "this" --> NFS\_OK. The current fh is for /this and is within the pseudo-fs.
- \* LOOKUP "is" --> NFS\_OK. The current fh is for /this/is and is within the pseudo-fs.
- \* LOOKUP "the" --> NFS\_OK. The current fh is for /this/is/the and is within the pseudo-fs.
- \* READDIR (rdattr\_error, fsid, size, time\_modify, mounted\_on\_fileid) --> NFS\_OK. The attributes for directory entry with the component named "path" will only contain rdattr\_error with the value NFS4ERR\_MOVED, together with an fsid value and a value for mounted\_on\_fileid.

Suppose we do another READDIR to get fs\_locations\_info (although we could have used a GETATTR directly, as in Section 17.15.1).

- \* PUTROOTFH
- \* LOOKUP "this"
- \* LOOKUP "is"
- \* LOOKUP "the"
- \* READDIR (rdattr\_error, fs\_locations\_info, mounted\_on\_fileid, fsid, size, time\_modify)

The results would be:

- \* PUTROOTFH --> NFS\_OK. The current fh is at the root of the pseudo-fs.
- \* LOOKUP "this" --> NFS\_OK. The current fh is for /this and is within the pseudo-fs.
- \* LOOKUP "is" --> NFS\_OK. The current fh is for /this/is and is within the pseudo-fs.

- \* LOOKUP "the" --> NFS\_OK. The current fh is for /this/is/the and is within the pseudo-fs.
- \* REaddir (rdattr\_error, fs\_locations\_info, mounted\_on\_fileid, fsid, size, time\_modify) --> NFS\_OK. The attributes will be as shown below.

The attributes for the directory entry with the component named "path" will only contain:

- \* rdattr\_error (value: NFS\_OK)
- \* fs\_locations\_info
- \* mounted\_on\_fileid (value: unique fileid within referring file system)
- \* fsid (value: unique value within referring server)

The attributes for entry "path" will not contain size or time\_modify because these attributes are not available within an absent file system.

#### 17.16. The Attribute fs\_locations

The fs\_locations attribute is structured in the following way:

```
struct fs_location4 {
    utf8str_mixed  server<>;
    pathname4      rootpath;
};

struct fs_locations4 {
    pathname4      fs_root;
    fs_location4   locations<>;
};
```

The fs\_location4 data type is used to represent the location of a file system by providing a server name and the path to the root of the file system within that server's namespace. When a set of servers have corresponding file systems at the same path within their namespaces, an array of server names may be provided. An entry in the server array is a UTF-8 string and represents one of a traditional DNS host name, IPv4 address, IPv6 address, or a zero-length string. An IPv4 or IPv6 address is represented as a universal address (See Section 9.3.9 and [RFC5665]), minus the netid, and either with or without the trailing ".p1.p2" suffix that represents the port number. If the suffix is omitted, then the default port,

2049, is to be assumed. A zero-length string can be used to indicate the current address being used for the RPC call. It is not a requirement that all servers that share the same rootpath be listed in one `fs_location4` instance. The array of server names is provided for convenience. Servers that share the same rootpath may also be listed in separate `fs_location4` entries in the `fs_locations` attribute.

The `fs_locations4` data type and the `fs_locations` attribute each contain an array of such locations. Since the namespace of each server may be constructed differently, the "`fs_root`" field is provided. The path represented by `fs_root` represents the location of the file system in the current server's namespace, i.e., that of the server from which the `fs_locations` attribute was obtained. The `fs_root` path is meant to aid the client by clearly referencing the root of the file system whose locations are being reported, no matter what object within the current file system the current filehandle designates. The `fs_root` is simply the pathname the client used to reach the object on the current server (i.e., the object to which the `fs_locations` attribute applies).

When the `fs_locations` attribute is interrogated and there are no alternate file system locations, the server returns a zero-length array of `fs_location4` structures, together with a valid `fs_root`.

As an example, suppose there is a replicated file system located at two servers (`servA` and `servB`). At `servA`, the file system is located at path `/a/b/c`. At, `servB` the file system is located at path `/x/y/z`. If the client were to obtain the `fs_locations` value for the directory at `/a/b/c/d`, it might not necessarily know that the file system's root is located in `servA`'s namespace at `/a/b/c`. When the client switches to `servB`, it will need to determine that the directory it first referenced at `servA` is now represented by the path `/x/y/z/d` on `servB`. To facilitate this, the `fs_locations` attribute provided by `servA` would have an `fs_root` value of `/a/b/c` and two entries in `fs_locations`. One entry in `fs_locations` will be for itself (`servA`) and the other will be for `servB` with a path of `/x/y/z`. With this information, the client is able to substitute `/x/y/z` for the `/a/b/c` at the beginning of its access path and construct `/x/y/z/d` to use for the new server.

Note that there is no requirement that the number of components in each rootpath be the same; there is no relation between the number of components in rootpath or fs\_root, and none of the components in a rootpath and fs\_root have to be the same. In the above example, we could have had a third element in the locations array, with server equal to "servC" and rootpath equal to "/I/II", and a fourth element in locations with server equal to "servD" and rootpath equal to "/aleph/beth/gimel/dalet/he".

The relationship between fs\_root to a rootpath is that the client replaces the pathname indicated in fs\_root for the current server for the substitute indicated in rootpath for the new server.

For an example of a referred or migrated file system, suppose there is a file system located at serv1. At serv1, the file system is located at /az/buky/vedi/glagoli. The client finds that object at glagoli has migrated (or is a referral). The client gets the fs\_locations attribute, which contains an fs\_root of /az/buky/vedi/glagoli, and one element in the locations array, with server equal to serv2, and rootpath equal to /izhitsa/fita. The client replaces /az/buky/vedi/glagoli with /izhitsa/fita, and uses the latter pathname on serv2.

Thus, the server MUST return an fs\_root that is equal to the path the client used to reach the object to which the fs\_locations attribute applies. Otherwise, the client cannot determine the new path to use on the new server.

Since the fs\_locations attribute lacks information defining various attributes of the various file system choices presented, it should only be interrogated and used when fs\_locations\_info is not available. When fs\_locations is used, information about the specific locations should be assumed based on the following rules.

The following rules are general and apply irrespective of the context.

- \* All listed file system instances should be considered as of the same handle class, if and only if, the current fh\_expire\_type attribute does not include the FH4\_VOL\_MIGRATION bit. Note that in the case of referral, filehandle issues do not apply since there can be no filehandles known within the current file system, nor is there any access to the fh\_expire\_type attribute on the referring (absent) file system.
- \* All listed file system instances should be considered as of the same fileid class if and only if the fh\_expire\_type attribute indicates persistent filehandles and does not include the



FH4\_VOL\_MIGRATION bit. Note that in the case of referral, fileid issues do not apply since there can be no fileids known within the referring (absent) file system, nor is there any access to the fh\_expire\_type attribute.

- \* All file system instances servers should be considered as of different change classes.

For other class assignments, handling of file system transitions depends on the reasons for the transition:

- \* When the transition is due to migration, that is, the client was directed to a new file system after receiving an NFS4ERR\_MOVED error, the target should be treated as being of the same write-verifier class as the source.
- \* When the transition is due to failover to another replica, that is, the client selected another replica without receiving an NFS4ERR\_MOVED error, the target should be treated as being of a different write-verifier class from the source.

The specific choices reflect typical implementation patterns for failover and controlled migration, respectively. Since other choices are possible and useful, this information is better obtained by using fs\_locations\_info. When a server implementation needs to communicate other choices, it MUST support the fs\_locations\_info attribute.

See Section 28 for a discussion on the recommendations for the security flavor to be used by any GETATTR operation that requests the fs\_locations attribute.

#### 17.17. The Attribute fs\_locations\_info

The fs\_locations\_info attribute is intended as a more functional replacement for the fs\_locations attribute, which will continue to exist and be supported. Clients can use it to get a more complete set of data about alternative file system locations, including additional network paths to access replicas in use and additional replicas. When the server does not support fs\_locations\_info, fs\_locations can be used to get a subset of the data. A server that supports fs\_locations\_info MUST support fs\_locations as well.

There is additional data present in fs\_locations\_info that is not available in fs\_locations:

- \* Attribute continuity information. This information will allow a client to select a replica that meets the transparency requirements of the applications accessing the data and to

leverage optimizations due to the server guarantees of attribute continuity (e.g., if the change attribute of a file of the file system is continuous between multiple replicas, the client does not have to invalidate the file's cache when switching to a different replica).

- \* File system identity information that indicates when multiple replicas, from the client's point of view, correspond to the same target file system, allowing them to be used interchangeably, without disruption, as distinct synchronized replicas of the same file data.

Note that having two replicas with common identity information is distinct from the case of two (trunked) paths to the same replica.

- \* Information that will bear on the suitability of various replicas, depending on the use that the client intends. For example, many applications need an absolutely up-to-date copy (e.g., those that write), while others may only need access to the most up-to-date copy reasonably available.
- \* Server-derived preference information for replicas, which can be used to implement load-balancing while giving the client the entire file system list to be used in case the primary fails.

The `fs_locations_info` attribute is structured similarly to the `fs_locations` attribute. A top-level structure (`fs_locations_info4`) contains the entire attribute including the root pathname of the file system and an array of lower-level structures that define replicas that share a common rootpath on their respective servers. The lower-level structure in turn (`fs_locations_item4`) contains a specific pathname and information on one or more individual network access paths. For that last, lowest level, `fs_locations_info` has an `fs_locations_server4` structure that contains per-server-replica information in addition to the file system location entry. This per-server-replica information includes a nominally opaque array, `fls_info`, within which specific pieces of information are located at the specific indices listed below.

Two `fs_location_server4` entries that are within different `fs_location_item4` structures are never trunkable, while two entries within in the same `fs_location_item4` structure might or might not be trunkable. Two entries that are trunkable will have identical identity information, although, as noted above, the converse is not the case.

The attribute will always contain at least a single `fs_locations_server` entry. Typically, there will be an entry with the `FS4LIGF_CUR_REQ` flag set, although in the case of a referral there will be no entry with that flag set.

It should be noted that `fs_locations_info` attributes returned by servers for various replicas may differ for various reasons. One server may know about a set of replicas that are not known to other servers. Further, compatibility attributes may differ. Filehandles might be of the same class going from replica A to replica B but not going in the reverse direction. This might happen because the filehandles are the same, but replica B's server implementation might not have provision to note and report that equivalence.

The `fs_locations_info` attribute consists of a root pathname (`fli_fs_root`, just like `fs_root` in the `fs_locations` attribute), together with an array of `fs_location_item4` structures. The `fs_location_item4` structures in turn consist of a root pathname (`fli_rootpath`) together with an array (`fli_entries`) of elements of data type `fs_locations_server4`, all defined as follows.

```
/*
 * Defines an individual server access path
 */
struct fs_locations_server4 {
    int32_t          fls_currency;
    opaque           fls_info<>;
    utf8str_mixed    fls_server;
};

/*
 * Byte indices of items within
 * fls_info: flag fields, class numbers,
 * bytes indicating ranks and orders.
 */
const FSLI4BX_GFLAGS          = 0;
const FSLI4BX_TFLAGS          = 1;

const FSLI4BX_CLSIMUL         = 2;
const FSLI4BX_CLHANDLE        = 3;
const FSLI4BX_CLFILEID        = 4;
const FSLI4BX_CLWRITEVER      = 5;
const FSLI4BX_CLCHANGE        = 6;
const FSLI4BX_CLREADDIR       = 7;

const FSLI4BX_READRANK         = 8;
const FSLI4BX_WRITERANK        = 9;
const FSLI4BX_READORDER        = 10;
```

```
const FSLI4BX_WRITEORDER      = 11;

/*
 * Bits defined within the general flag byte.
 */
const FSLI4GF_WRITABLE        = 0x01;
const FSLI4GF_CUR_REQ         = 0x02;
const FSLI4GF_ABSENT          = 0x04;
const FSLI4GF_GOING           = 0x08;
const FSLI4GF_SPLIT           = 0x10;

/*
 * Bits defined within the transport flag byte.
 */
const FSLI4TF_RDMA             = 0x01;

/*
 * Defines a set of replicas sharing
 * a common value of the rootpath
 * within the corresponding
 * single-server namespaces.
 */
struct fs_locations_item4 {
    fs_locations_server4    fli_entries<>;
    pathname4               fli_rootpath;
};

/*
 * Defines the overall structure of
 * the fs_locations_info attribute.
 */
struct fs_locations_info4 {
    uint32_t                fli_flags;
    int32_t                 fli_valid_for;
    pathname4               fli_fs_root;
    fs_locations_item4      fli_items<>;
};

/*
 * Flag bits in fli_flags.
 */
const FSLI4IF_VAR_SUB        = 0x00000001;

typedef fs_locations_info4 fattr4_fs_locations_info;
```

As noted above, the `fs_locations_info` attribute, when supported, may be requested of absent file systems without causing `NFS4ERR_MOVED` to be returned. It is generally expected that it will be available for

both present and absent file systems even if only a single `fs_locations_server4` entry is present, designating the current (present) file system, or two `fs_locations_server4` entries designating the previous location of an absent file system (the one just referenced) and its successor location. Servers are strongly urged to support this attribute on all file systems if they support it on any file system.

The data presented in the `fs_locations_info` attribute may be obtained by the server in any number of ways, including specification by the administrator or by current protocols for transferring data among replicas and protocols not yet developed. NFSv4.1 only defines how this information is presented by the server to the client.

#### 17.17.1. The `fs_locations_server4` Structure

The `fs_locations_server4` structure consists of the following items in addition to the `fls_server` field, which specifies a network address or set of addresses to be used to access the specified file system. Note that both of these items (i.e., `fls_currency` and `fls_info`) specify attributes of the file system replica and should not be different when there are multiple `fs_locations_server4` structures, each specifying a network path to the chosen replica, for the same replica.

When these values are different in two `fs_locations_server4` structures, a client has no basis for choosing one over the other and is best off simply ignoring both entries, whether these entries apply to migration replication or referral. When there are more than two such entries, majority voting can be used to exclude a single erroneous entry from consideration. In the case in which trunking information is provided for a replica currently being accessed, the additional trunked addresses can be ignored while access continues on the address currently being used, even if the entry corresponding to that path might be considered invalid.

- \* An indication of how up-to-date the file system is (`fls_currency`) in seconds. This value is relative to the master copy. A negative value indicates that the server is unable to give any reasonably useful value here. A value of zero indicates that the file system is the actual writable data or a reliably coherent and fully up-to-date copy. Positive values indicate how out-of-date this copy can normally be before it is considered for update. Such a value is not a guarantee that such updates will always be performed on the required schedule but instead serves as a hint about how far the copy of the data would be expected to be behind the most up-to-date copy.

- \* A counted array of one-byte values (`fls_info`) containing information about the particular file system instance. This data includes general flags, transport capability flags, file system equivalence class information, and selection priority information. The encoding will be discussed below.
- \* The server string (`fls_server`). For the case of the replica currently being accessed (via `GETATTR`), a zero-length string MAY be used to indicate the current address being used for the RPC call. The `fls_server` field can also be an IPv4 or IPv6 address, formatted the same way as an IPv4 or IPv6 address in the "server" field of the `fs_location4` data type (see Section 17.16).

With the exception of the transport-flag field (at offset `FSLI4BX_TFLAGS` with the `fls_info` array), all of this data defined in this specification applies to the replica specified by the entry, rather than the specific network path used to access it. The classification of data in extensions to this data is discussed below.

Data within the `fls_info` array is in the form of 8-bit data items with constants giving the offsets within the array of various values describing this particular file system instance. This style of definition was chosen, in preference to explicit XDR structure definitions for these values, for a number of reasons.

- \* The kinds of data in the `fls_info` array, representing flags, file system classes, and priorities among sets of file systems representing the same data, are such that 8 bits provide a quite acceptable range of values. Even where there might be more than 256 such file system instances, having more than 256 distinct classes or priorities is unlikely.
- \* Explicit definition of the various specific data items within XDR would limit expandability in that any extension within would require yet another attribute, leading to specification and implementation clumsiness. In the context of the NFSv4 extension model in effect at the time `fs_locations_info` was designed (i.e., that which is described in [RFC5661]), this would necessitate a new minor version to effect any Standards Track extension to the data in `fls_info`.

The set of `fls_info` data is subject to expansion in a future minor version or in a Standards Track RFC within the context of a single minor version. The server should not send and the client MUST NOT use indices within the `fls_info` array or flag bits that are not defined in Standards Track RFCs.

In light of the new extension model defined in [RFC8178] and the fact that the individual items within `fls_info` are not explicitly referenced in the XDR, the following practices should be followed when extending or otherwise changing the structure of the data returned in `fls_info` within the scope of a single minor version:

- \* All extensions need to be described by Standards Track documents. There is no need for such documents to be marked as updating [RFC5661], [RFC8881], or this document.
- \* It needs to be made clear whether the information in any added data items applies to the replica specified by the entry or to the specific network paths specified in the entry.
- \* There needs to be a reliable way defined to determine whether the server is aware of the extension. This may be based on the length field of the `fls_info` array, but it is more flexible to provide fs-scope or server-scope attributes to indicate what extensions are provided.

This encoding scheme can be adapted to the specification of multi-byte numeric values, even though none are currently defined. If extensions are made via Standards Track RFCs, multi-byte quantities will be encoded as a range of bytes with a range of indices, with the byte interpreted in big-endian byte order. Further, any such index assignments will be constrained by the need for the relevant quantities not to cross XDR word boundaries.

The `fls_info` array currently contains:

- \* Two 8-bit flag fields, one devoted to general file-system characteristics and a second reserved for transport-related capabilities.
- \* Six 8-bit class values that define various file system equivalence classes as explained below.
- \* Four 8-bit priority values that govern file system selection as explained below.

The general file system characteristics flag (at byte index `FSLI4BX_GFLAGS`) has the following bits defined within it:

- \* `FSLI4GF_WRITABLE` indicates that this file system target is writable, allowing it to be selected by clients that may need to write on this file system. When the current file system instance is writable and is defined as of the same simultaneous use class (as specified by the value at index `FSLI4BX_CLSIMUL`) to which the

client was previously writing, then it must incorporate within its data any committed write made on the source file system instance. See Section 17.11.6, which discusses the write-verifier class. While there is no harm in not setting this flag for a file system that turns out to be writable, turning the flag on for a read-only file system can cause problems for clients that select a migration or replication target based on the flag and then find themselves unable to write.

- \* FSLI4GF\_CUR\_REQ indicates that this replica is the one on which the request is being made. Only a single server entry may have this flag set and, in the case of a referral, no entry will have it set. Note that this flag might be set even if the request was made on a network access path different from any of those specified in the current entry.
- \* FSLI4GF\_ABSENT indicates that this entry corresponds to an absent file system replica. It can only be set if FSLI4GF\_CUR\_REQ is set. When both such bits are set, it indicates that a file system instance is not usable but that the information in the entry can be used to determine the sorts of continuity available when switching from this replica to other possible replicas. Since this bit can only be true if FSLI4GF\_CUR\_REQ is true, the value could be determined using the fs\_status attribute, but the information is also made available here for the convenience of the client. An entry with this bit, since it represents a true file system (albeit absent), does not appear in the event of a referral, but only when a file system has been accessed at this location and has subsequently been migrated.
- \* FSLI4GF\_GOING indicates that a replica, while still available, should not be used further. The client, if using it, should make an orderly transfer to another file system instance as expeditiously as possible. It is expected that file systems going out of service will be announced as FSLI4GF\_GOING some time before the actual loss of service. It is also expected that the fli\_valid\_for value will be sufficiently small to allow clients to detect and act on scheduled events, while large enough that the cost of the requests to fetch the fs\_locations\_info values will not be excessive. Values on the order of ten minutes seem reasonable.



When this flag is seen as part of a transition into a new file system, a client might choose to transfer immediately to another replica, or it may reference the current file system and only transition when a migration event occurs. Similarly, when this flag appears as a replica in the referral, clients would likely avoid being referred to this instance whenever there is another choice.

This flag, like the other items within `fls_info`, applies to the replica rather than to a particular path to that replica. When it appears, a transition to a new replica, rather than to a different path to the same replica, is indicated.

- \* `FSLI4GF_SPLIT` indicates that when a transition occurs from the current file system instance to this one, the replacement may consist of multiple file systems. In this case, the client has to be prepared for the possibility that objects on the same file system before migration will be on different ones after. Note that `FSLI4GF_SPLIT` is not incompatible with the file systems belonging to the same fileid class since, if one has a set of fileids that are unique within a file system, each subset assigned to a smaller file system after migration would not have any conflicts internal to that file system.

A client, in the case of a split file system, will interrogate existing files with which it has continuing connection (it is free to simply forget cached filehandles). If the client remembers the directory filehandle associated with each open file, it may proceed upward using `LOOKUPP` to find the new file system boundaries. Note that in the event of a referral, there will not be any such files and so these actions will not be performed. Instead, a reference to a portion of the original file system now split off into other file systems will encounter an `fsid` change and possibly a further referral.

Once the client recognizes that one file system has been split into two, it can prevent the disruption of running applications by presenting the two file systems as a single one until a convenient point to recognize the transition, such as a restart. This would require a mapping from the server's fsids to fsids as seen by the client, but this is already necessary for other reasons. As noted above, existing fileids within the two descendant file systems will not conflict. Providing non-conflicting fileids for newly created files on the split file systems is the responsibility of the server (or servers working in concert). The server can encode filehandles such that filehandles generated before the split event can be discerned from those generated after the split, allowing the server to determine when the need for emulating two file systems as one is over.

Although it is possible for this flag to be present in the event of referral, it would generally be of little interest to the client, since the client is not expected to have information regarding the current contents of the absent file system.

The transport-flag field (at byte index FSLI4BX\_TFLAGS) contains the following bits related to the transport capabilities of the specific network path(s) specified by the entry:

- \* FSLI4TF\_RDMA indicates that any specified network paths provide NFSv4.1 clients access using an RDMA-capable transport.

Attribute continuity and file system identity information are expressed by defining equivalence relations on the sets of file systems presented to the client. Each such relation is expressed as a set of file system equivalence classes. For each relation, a file system has an 8-bit class number. Two file systems belong to the same class if both have identical non-zero class numbers. Zero is treated as non-matching. Most often, the relevant question for the client will be whether a given replica is identical to / continuous with the current one in a given respect, but the information should be available also as to whether two other replicas match in that respect as well.

The following fields specify the file system's class numbers for the equivalence relations used in determining the nature of file system transitions. See Sections 17.9 through 17.14 and their various subsections for details about how this information is to be used. Servers may assign these values as they wish, so long as file system instances that share the same value have the specified relationship to one another; conversely, file systems that have the specified relationship to one another share a common class value. As each instance entry is added, the relationships of this instance to

previously entered instances can be consulted, and if one is found that bears the specified relationship, that entry's class value can be copied to the new entry. When no such previous entry exists, a new value for that byte index (not previously used) can be selected, most likely by incrementing the value of the last class value assigned for that index.

- \* The field with byte index FSLI4BX\_CLSIMUL defines the simultaneous-use class for the file system.
- \* The field with byte index FSLI4BX\_CLHANDLE defines the handle class for the file system.
- \* The field with byte index FSLI4BX\_CLFILEID defines the fileid class for the file system.
- \* The field with byte index FSLI4BX\_CLWRITEVER defines the write-verifier class for the file system.
- \* The field with byte index FSLI4BX\_CLCHANGE defines the change class for the file system.
- \* The field with byte index FSLI4BX\_CLREADDIR defines the readdir class for the file system.

Server-specified preference information is also provided via 8-bit values within the `fls_info` array. The values provide a rank and an order (See below) to be used with separate values specifiable for the cases of read-only and writable file systems. These values are compared for different file systems to establish the server-specified preference, with lower values indicating "more preferred".

Rank is used to express a strict server-imposed ordering on clients, with lower values indicating "more preferred". Clients should attempt to use all replicas with a given rank before they use one with a higher rank. Only if all of those file systems are unavailable should the client proceed to those of a higher rank. Because specifying a rank will override client preferences, servers should be conservative about using this mechanism, particularly when the environment is one in which client communication characteristics are neither tightly controlled nor visible to the server.

Within a rank, the order value is used to specify the server's preference to guide the client's selection when the client's own preferences are not controlling, with lower values of order indicating "more preferred". If replicas are approximately equal in all respects, clients should defer to the order specified by the server. When clients look at server latency as part of their

selection, they are free to use this criterion, but it is suggested that when latency differences are not significant, the server-specified order should guide selection.

- \* The field at byte index FSLI4BX\_READRANK gives the rank value to be used for read-only access.
- \* The field at byte index FSLI4BX\_READORDER gives the order value to be used for read-only access.
- \* The field at byte index FSLI4BX\_WRITERANK gives the rank value to be used for writable access.
- \* The field at byte index FSLI4BX\_WRITEORDER gives the order value to be used for writable access.

Depending on the potential need for write access by a given client, one of the pairs of rank and order values is used. The read rank and order should only be used if the client knows that only reading will ever be done or if it is prepared to switch to a different replica in the event that any write access capability is required in the future.

#### 17.17.2. The fs\_locations\_info4 Structure

The fs\_locations\_info4 structure, encoding the fs\_locations\_info attribute, contains the following:

- \* The fli\_flags field, which contains general flags that affect the interpretation of this fs\_locations\_info4 structure and all fs\_locations\_item4 structures within it. The only flag currently defined is FSLI4IF\_VAR\_SUB. All bits in the fli\_flags field that are not defined should always be returned as zero.
- \* The fli\_fs\_root field, which contains the pathname of the root of the current file system on the current server, just as it does in the fs\_locations4 structure.
- \* An array called fli\_items of fs\_locations4\_item structures, which contain information about replicas of the current file system. Where the current file system is actually present, or has been present, i.e., this is not a referral situation, one of the fs\_locations\_item4 structures will contain an fs\_locations\_server4 for the current server. This structure will have FSLI4GF\_ABSENT set if the current file system is absent, i.e., normal access to it will return NFS4ERR\_MOVED.

- \* The `fli_valid_for` field specifies a time in seconds for which it is reasonable for a client to use the `fs_locations_info` attribute without refetch. The `fli_valid_for` value does not provide a guarantee of validity since servers can unexpectedly go out of service or become inaccessible for any number of reasons. Clients are well-advised to refetch this information for an actively accessed file system at every `fli_valid_for` seconds. This is particularly important when file system replicas may go out of service in a controlled way using the `FSLI4GF_GOING` flag to communicate an ongoing change. The server should set `fli_valid_for` to a value that allows well-behaved clients to notice the `FSLI4GF_GOING` flag and make an orderly switch before the loss of service becomes effective. If this value is zero, then no refetch interval is appropriate and the client need not refetch this data on any particular schedule. In the event of a transition to a new file system instance, a new value of the `fs_locations_info` attribute will be fetched at the destination. It is to be expected that this may have a different `fli_valid_for` value, which the client should then use in the same fashion as the previous value. Because a refetch of the attribute causes information from all component entries to be refetched, the server will typically provide a low value for this field if any of the replicas are likely to go out of service in a short time frame. Note that, because of the ability of the server to return `NFS4ERR_MOVED` to trigger the use of different paths, when alternate trunked paths are available, there is generally no need to use low values of `fli_valid_for` in connection with the management of alternate paths to the same replica.

The `FSLI4IF_VAR_SUB` flag within `fli_flags` controls whether variable substitution is to be enabled. See Section 17.17.3 for an explanation of variable substitution.

#### 17.17.3. The `fs_locations_item4` Structure

The `fs_locations_item4` structure contains a pathname (in the field `fli_rootpath`) that encodes the path of the target file system replicas on the set of servers designated by the included `fs_locations_server4` entries. The precise manner in which this target location is specified depends on the value of the `FSLI4IF_VAR_SUB` flag within the associated `fs_locations_info4` structure.

If this flag is not set, then `fli_rootpath` simply designates the location of the target file system within each server's single-server namespace just as it does for the rootpath within the `fs_location4` structure. When this bit is set, however, component entries of a certain form are subject to client-specific variable substitution so

as to allow a degree of namespace non-uniformity in order to accommodate the selection of client-specific file system targets to adapt to different client architectures or other characteristics.

When such substitution is in effect, a variable beginning with the string "\${" and ending with the string "}" and containing a colon is to be replaced by the client-specific value associated with that variable. The string "unknown" should be used by the client when it has no value for such a variable. The pathname resulting from such substitutions is used to designate the target file system, so that different clients may have different file systems, corresponding to that location in the multi-server namespace.

As mentioned above, such substituted pathname variables contain a colon. The part before the colon is to be a DNS domain name, and the part after is to be a case-insensitive alphanumeric string.

Where the domain is "ietf.org", only variable names defined in this document or subsequent Standards Track RFCs are subject to such substitution. Organizations are free to use their domain names to create their own sets of client-specific variables, to be subject to such substitution. In cases where such variables are intended to be used more broadly than a single organization, publication of an Informational RFC defining such variables is RECOMMENDED.

The variable `${ietf.org:CPU_ARCH}` is used to denote that the CPU architecture object files are compiled. This specification does not limit the acceptable values (except that they must be valid UTF-8 strings), but such values as "x86", "x86\_64", and "sparc" would be expected to be used in line with industry practice.

The variable `${ietf.org:OS_TYPE}` is used to denote the operating system, and thus the kernel and library APIs, for which code might be compiled. This specification does not limit the acceptable values (except that they must be valid UTF-8 strings), but such values as "linux" and "freebsd" would be expected to be used in line with industry practice.

The variable `${ietf.org:OS_VERSION}` is used to denote the operating system version, and thus the specific details of versioned interfaces, for which code might be compiled. This specification does not limit the acceptable values (except that they must be valid UTF-8 strings). However, combinations of numbers and letters with interspersed dots would be expected to be used in line with industry practice, with the details of the version format depending on the specific value of the variable `${ietf.org:OS_TYPE}` with which it is used.

Use of these variables could result in the direction of different clients to different file systems on the same server, as appropriate to particular clients. In cases in which the target file systems are located on different servers, a single server could serve as a referral point so that each valid combination of variable values would designate a referral hosted on a single server, with the targets of those referrals on a number of different servers.

Because namespace administration is affected by the values selected to substitute for various variables, clients should provide convenient means of determining what variable substitutions a client will implement, as well as, where appropriate, providing means to control the substitutions to be used. The exact means by which this will be done is outside the scope of this specification.

Although variable substitution is most suitable for use in the context of referrals, it may be used in the context of replication and migration. If it is used in these contexts, the server must ensure that no matter what values the client presents for the substituted variables, the result is always a valid successor file system instance to that from which a transition is occurring, i.e., that the data is identical or represents a later image of a writable file system.

Note that when `fli_rootpath` is a null pathname (that is, one with zero components), the file system designated is at the root of the specified server, whether or not the `FSLI4IF_VAR_SUB` flag within the associated `fs_locations_info4` structure is set.

#### 17.18. The Attribute `fs_status`

In an environment in which multiple copies of the same basic set of data are available, information regarding the particular source of such data and the relationships among different copies can be very helpful in providing consistent data to applications.

```
enum fs4_status_type {
    STATUS4_FIXED = 1,
    STATUS4_UPDATED = 2,
    STATUS4_VERSIONED = 3,
    STATUS4_WRITABLE = 4,
    STATUS4_REFERRAL = 5
};

struct fs4_status {
    bool                fss_absent;
    fs4_status_type     fss_type;
    utf8str_cs          fss_source;
    utf8str_cs          fss_current;
    int32_t             fss_age;
    nfstime4            fss_version;
};
```

The boolean `fss_absent` indicates whether the file system is currently absent. This value will be set if the file system was previously present and becomes absent, or if the file system has never been present and the type is `STATUS4_REFERRAL`. When this boolean is set and the type is not `STATUS4_REFERRAL`, the remaining information in the `fs4_status` reflects that last valid when the file system was present.

The `fss_type` field indicates the kind of file system image represented. This is of particular importance when using the version values to determine appropriate succession of file system images. When `fss_absent` is set, and the file system was previously present, the value of `fss_type` reflected is that when the file was last present. Five values are distinguished:

- \* `STATUS4_FIXED`, which indicates a read-only image in the sense that it will never change. The possibility is allowed that, as a result of migration or switch to a different image, changed data can be accessed, but within the confines of this instance, no change is allowed. The client can use this fact to cache aggressively.
- \* `STATUS4_VERSIONED`, which indicates that the image, like the `STATUS4_UPDATED` case, is updated externally, but it provides a guarantee that the server will carefully update an associated version value so that the client can protect itself from a situation in which it reads data from one version of the file system and then later reads data from an earlier version of the same file system. See below for a discussion of how this can be done.



- \* STATUS4\_UPDATED, which indicates an image that cannot be updated by the user writing to it but that may be changed externally, typically because it is a periodically updated copy of another writable file system somewhere else. In this case, version information is not provided, and the client does not have the responsibility of making sure that this version only advances upon a file system instance transition. In this case, it is the responsibility of the server to make sure that the data presented after a file system instance transition is a proper successor image and includes all changes seen by the client and any change made before all such changes.
- \* STATUS4\_WRITABLE, which indicates that the file system is an actual writable one. The client need not, of course, actually write to the file system, but once it does, it should not accept a transition to anything other than a writable instance of that same file system.
- \* STATUS4\_REFERRAL, which indicates that the file system in question is absent and has never been present on this server.

Note that in the STATUS4\_UPDATED and STATUS4\_VERSIONED cases, the server is responsible for the appropriate handling of locks that are inconsistent with external changes to delegations. If a server gives out delegations, they need to be recalled before an inconsistent change is made to the data, and MUST be revoked if this is not possible. Similarly, if an OPEN is inconsistent with data that is changed (the OPEN has OPEN4\_SHARE\_DENY\_WRITE/OPEN4\_SHARE\_DENY\_BOTH and the data is changed), that OPEN should be considered administratively revoked.

The opaque strings fss\_source and fss\_current provide a way of presenting information about the source of the file system image being present. It is not intended that the client do anything with this information other than make it available to administrative tools. It is intended that this information be helpful when researching possible problems with a file system image that might arise when it is unclear if the correct image is being accessed and, if not, how that image came to be made. This kind of diagnostic information will be helpful, if, as seems likely, copies of file systems are made in many different ways (e.g., simple user-level copies, file-system-level point-in-time copies, clones of the underlying storage), under a variety of administrative arrangements. In such environments, determining how a given set of data was constructed can be very helpful in resolving problems.

The opaque string `fss_source` is used to indicate the source of a given file system with the expectation that tools capable of creating a file system image propagate this information, when possible. It is understood that this may not always be possible since a user-level copy may be thought of as creating a new data set and the tools used may have no mechanism to propagate this data. When a file system is initially created, it is desirable to associate with it data regarding how the file system was created, where it was created, who created it, etc. Making this information available in this attribute in a human-readable string will be helpful for applications and system administrators and will also serve to make it available when the original file system is used to make subsequent copies.

The opaque string `fss_current` should provide whatever information is available about the source of the current copy. Such information includes the tool creating it, any relevant parameters to that tool, the time at which the copy was done, the user making the change, the server on which the change was made, etc. All information should be in a human-readable string.

The field `fss_age` provides an indication of how out-of-date the file system currently is with respect to its ultimate data source (in case of cascading data updates). This complements the `fls_currency` field of `fs_locations_server4` (See Section 17.17) in the following way: the information in `fls_currency` gives a bound for how out of date the data in a file system might typically get, while the value in `fss_age` gives a bound on how out-of-date that data actually is. Negative values imply that no information is available. A zero means that this data is known to be current. A positive value means that this data is known to be no older than that number of seconds with respect to the ultimate data source. Using this value, the client may be able to decide that a data copy is too old, so that it may search for a newer version to use.

The `fss_version` field provides a version identification, in the form of a time value, such that successive versions always have later time values. When the `fs_type` is anything other than `STATUS4_VERSIONED`, the server may provide such a value, but there is no guarantee as to its validity and clients will not use it except to provide additional information to add to `fss_source` and `fss_current`.

When `fss_type` is `STATUS4_VERSIONED`, servers should provide a value of `fss_version` that progresses monotonically whenever any new version of the data is established. This allows the client, if reliable image progression is important to it, to fetch this attribute as part of each `COMPOUND` where data or metadata from the file system is used.

When it is important to the client to make sure that only valid successor images are accepted, it must make sure that it does not read data or metadata from the file system without updating its sense of the current state of the image. This is to avoid the possibility that the `fs_status` that the client holds will be one for an earlier image, which would cause the client to accept a new file system instance that is later than that but still earlier than the updated data read by the client.

In order to accept valid images reliably, the client must do a `GETATTR` of the `fs_status` attribute that follows any interrogation of data or metadata within the file system in question. Often this is most conveniently done by appending such a `GETATTR` after all other operations that reference a given file system. When errors occur between reading file system data and performing such a `GETATTR`, care must be exercised to make sure that the data in question is not used before obtaining the proper `fs_status` value. In this connection, when an `OPEN` is done within such a versioned file system and the associated `GETATTR` of `fs_status` is not successfully completed, the open file in question must not be accessed until that `fs_status` is fetched.

The procedure above will ensure that before using any data from the file system the client has in hand a newly-fetched current version of the file system image. Multiple values for multiple requests in flight can be resolved by assembling them into the required partial order (and the elements should form a total order within the partial order) and using the last. The client may then, when switching among file system instances, decline to use an instance that does not have an `fss_type` of `STATUS4_VERSIONED` or whose `fss_version` field is earlier than the last one obtained from the predecessor file system instance.

## 18. Parallel NFS (pNFS)

### 18.1. Introduction

pNFS is an OPTIONAL feature within NFSv4.1. The pNFS feature provides for the use of multiple layout types within a common framework. Within this framework, each layout type specification is obliged to define how the general features of pNFS are to be addressed within that layout type. The requirements that needed to be addressed by the specification for each layout type are described in Section 19.1. The normative character of these obligations is discussed in Section 18.3

The description of the pNFS-related features are divided multiple documents and multiple top-level sections in this document:

- \* Common aspects of PNFS that apply to all layout types are described in Section 18.
- \* The specifics that might vary among different layout types are discussed in Section 19.
- \* The files layout type is described in Section 20.

Because the data access protocol for this layout type is a variant of NFSv4.1, it is described in this document. However, it should not be assumed that this placement makes this layout type more important or makes it somehow not optional.

- \* A number of other layout types made part of NFSv4.1 are described in separate layout type specification documents (i.e. [RFC5663], [RFC5664]).
- \* Additional layout types are part of NFSv4.2 and are described in their own layout type specification documents. The addition of further such layout types is provided for in Section 29.5.

Some Layout types allow direct access to file data's location including the possibility of unmediated access to the storage devices containing file data. Others allow the direction of IO requests to the appropriate data server (considered as a type of data storage device).

When file data for a single NFSv4 server (often referred to as the metadata server) is stored on multiple and/or higher-throughput storage devices (compared to the primary server's throughput capability), these approaches can provide significantly better file access performance. The relationship among multiple clients, a single server (assuming the role of a metadata server), and multiple file storage devices for pNFS (some of which act as data servers by implementing support for file access protocols), is shown in Figure 1.

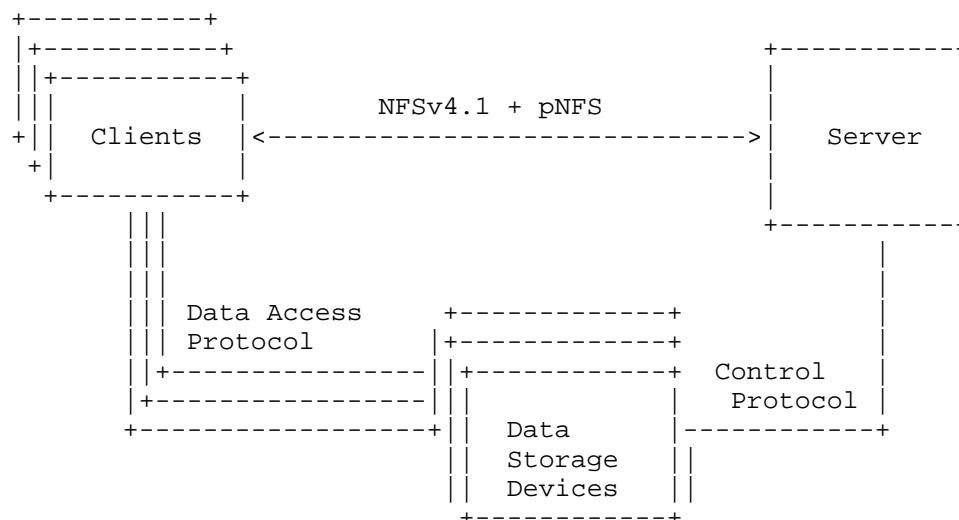


Figure 1

In this approach, the clients, the metadata server, and data storage devices work together to provide file data access and to deny it to those not appropriately authorized. The specifics of the responsibility of the participants to contribute to that authorization are defined in the layout type specification for the layouts used to perform IO operations.

This approach is in contrast to NFSv4 without pNFS, where this is primarily the server's responsibility while some of this responsibility may be delegated to the client under strictly specified conditions.

See Section 18.4.5 for a discussion of the Data Access Protocols. See Section 18.4.6 for a discussion of Control Protocols which provide the mechanisms used to coordinate the metadata server and the data storage devices.

pNFS involves OPTIONAL operations that manage protocol objects called 'layouts' (Section 18.2) that contain a byte-range and location information. Layouts are managed in a fashion similar to NFSv4.1 data delegations. For example, layouts are leased, recallable, and revocable. Note that layouts are distinct abstractions and are manipulated with new operations. When a client holds a layout, it is granted the ability to directly access the byte-range at the storage location designated in the layout.

There are interactions between layouts and other NFSv4.1 abstractions such as data delegations and byte-range locking. It is the responsibility of the layout type specification to address these issues as described in Section 19.1.10

## 18.2. Layout Types

A layout describes the mapping of a file's data to the file data devices that hold the data. Each layout is of a particular type (data type layouttype4, see Section 9.3.13). The existence of multiple layout type allows for multiple pNFS-based features to handle different data access protocols, such as those associated with the various supported layout types. A metadata server that supports pNFS MUST support at least one layout type. A private sub-range of the layout type namespace is also defined. Values from the private layout type range MAY be used for internal testing or experimentation (See Section 9.3.13).

Requests for pNFS-related operations will often specify a layout type. Examples of such operations are GETDEVICEINFO and LAYOUTGET. The response for these operations will include structures such as a device\_addr4 or a layout4, each of which includes a layout type within it. The layout type sent by the server MUST always be the same one requested by the client. When a server sends a response that includes a different layout type, the client should ignore the response and behave as if the server had returned an error response.

## 18.3. Normative Layout Specification Terminology

As described in Section 1.1, the keywords defined by [RFC2119] and [RFC8174] have special meanings that this document intends to adhere to. However, due to the nature of Sections Section 19 format="counter"/> and 29.5.3 together with some special circumstances resulting from the fact that complementations are described indirectly, there are some complexities that it is important to take note of:

- \* Where this document does not directly specify implementation requirements, use of these capitalized terms is often not appropriate since the guidance given in this document does not directly affect interoperability.
- \* In this document, what authors of RFCs defining new layout types need to do is stated without these specialized terms.

Although it is necessary to follow this guidance to provide successful pNFS layout type specifications, that sort of necessity is not of the sort defined as applicable to the use of the keywords defined in [RFC2119] [RFC8174].

The fact that these capitalized terms are not used should not be interpreted as indicating that this guidance does not need to be followed or is somehow not important.

Another way of stating this is that the material in Section 19.1 is "normative", as that word is normally used. This is so even though it is often assumed that text not using these terms is non-normative and that the implementer/author is free to disregard it.

- \* When one of these upper-case keywords defined in [RFC2119] [RFC8174] is used in Section 19.1, it is in the context of a rule directed to an implementer of client using a new layout type or a metadata server or data storage device implementing a new layout type, or in a quotation, sometimes indirect, from another document.

Overall, there are three types of normative statements used in Section 19.1:

- \* Requirements that applies to implementations.

These are stated using RFC2119-defined keywords.

- \* Requirements that apply to those defining layout types.

These are stated without using RFC2119-defined keywords. Nevertheless, they establish expectations that layout type specification need to satisfy (i.e norms) and are therefore described as "normative", the common use of that word in other IETF documents notwithstanding.

- \* Requirements which are general in that implementations need to conform to them, with the specific means by which the requirement is to be met left to layout type specification.

In this case, it is the responsibility of the layout type specification to clearly provide a description of the means to be used, the provision of which is obligatory despite the absence of RFC2119-defined terms.

#### 18.4. pNFS Definitions

NFSv4.1's pNFS feature provides parallel data access to a file system that stripes or otherwise divides its content across multiple storage servers. pNFS separates the file system protocol processing into two parts: metadata processing and data processing. Data consist of the contents of regular files that may be divided across storage servers. Division of data occurs in at least two ways: on a file-by-file basis and, within sufficiently large files, on a block-by-block basis. In contrast, divided access to metadata by pNFS clients is not provided in NFSv4.1, even though the file system back end of a pNFS server might stripe metadata. Metadata consist of everything other than file data, including the contents of non-regular files (e.g., directories); see Section 18.4.1. The metadata functionality is implemented by an NFSv4.1 server that supports pNFS and the operations described in Section 25; such a server is called a metadata server (Section 18.4.2).

The data functionality is implemented by one or more data storage devices, each of which is accessed by the client via a file access protocol. Definition of this protocol is the responsibility of the layout type specification as discussed in Section 19.1.

New terms are introduced to the NFSv4.1 nomenclature and existing terms are clarified to allow for the description of the pNFS feature suite.

##### 18.4.1. Metadata

Information about a file system object, such as its name, location within the namespace, owner, ACL, and other attributes. Metadata may also include storage location information, and this will with the layout type used based on the underlying storage mechanism that is used.

While many attributes are only known to the metadata server there are some necessary to appropriately process IO operations and others affected by executing IO operations. It is the job of the layout type specification to describe the necessary coordination. See Section 19.1 for details.



#### 18.4.2. Metadata Server

An NFSv4.1 server that supports the pNFS feature. A variety of architectural choices exist for the metadata server and its use of file system information held at the server. Some servers may contain metadata only for file objects residing at the metadata server, while the file data resides on associated storage devices. Other metadata servers may hold both metadata together with some portion of file data.

#### 18.4.3. pNFS Client

An NFSv4.1 client that supports pNFS operations and supports at least one data access protocol for performing I/O to data storage devices.

#### 18.4.4. Data Storage Devices

A data storage device stores a regular file's data, but while metadata management is done by the metadata server. A data storage device could be another NFSv4.1 server, an object-based storage device (OSD), a block device accessed over a System Area Network (SAN, e.g., either FiberChannel or iSCSI SAN), or some other entity.

#### 18.4.5. Data Access Protocol

As noted in Figure 1, the data access protocol is provided as a method used by the client to store and retrieve located on data storage devices.

The NFSv4.1 pNFS feature has been structured to allow a variety of data access protocols to be defined and used. The one actually used depends on the type of layout and is specified by the layout type specification as described in Section 19.1,

To use a particular data access protocol, both the metadata server and the client must have support for a layout type that uses that protocol as its data access protocol.

#### 18.4.6. Control Protocol

As illustrated in Figure 1, a control protocol is used between the metadata server and data storage devices. Specification of such protocols is outside the scope of the NFSv4.1 protocol. Such control protocols would be used to control activities such as the allocation and deallocation of storage, the management of state required by the data storage devices to perform client access control, and, depending on the file access protocol, the enforcement of authentication and authorization so that restrictions that would be enforced by the

metadata server could also be enforced by the data storage devices.

A particular control protocol is not REQUIRED by NFSv4.1 but requirements are placed on the control protocol for maintaining attributes such as modify time, the change attribute, and the end-of-file (EOF) position. Note that if pNFS is layered over a clustered, parallel file system (e.g., PVFS [PVFS]), the mechanisms that enable clustering and parallelism in that file system serve the same role as a control protocol. The differences in communication approaches allow these mechanisms to be treated as if they were structured as a control protocol.

In the flexible files layout [RFC8435], it is often stated that, when used in the "loose" coupling mode, there is no control protocol, which is at variance with the way control protocols are treated in this document, which requires that certain activities described be done and considers the means that are used to perform them as constituting a control protocol. In fact, these activities are performed using the same base protocol as the data access protocol, albeit in a different mode with greater privileges. In this document, we describe such arrangements as having no "separate control protocol."

#### 18.4.7. Layout

A layout defines how a file's data is organized on one or more data storage devices. There are many potential layout types; each of the layout types is differentiated by the data access protocol used to access data and by the aggregation scheme that lays out the file data on the various data storage devices. A layout is precisely identified by the tuple <client ID, filehandle, layout type, iomode, range>, where filehandle refers to the filehandle of the file on the metadata server.

It is important to define when layouts overlap and/or conflict with each other. For two layouts with overlapping byte-ranges to actually conflict, both layouts must be of the same layout type, correspond to the same filehandle, and meet layout-specific conditions specified by the layout type specification as discussed in Section 19.1.

#### 18.4.8. Layout Iomode

The layout iomode (data type layoutiomode4, see Section 9.3.20) indicates to the metadata server the client's intent to perform either just READ operations or a mixture containing READ and WRITE operations. For certain layout types, as defined by the layout type specification, it is appropriate for a client to specify this intent at the time it sends LAYOUTGET (Section 25.43). For example, for

block/volume-based protocols, block allocation could occur when a LAYOUTIOMODE4\_RW iomode is specified. A special LAYOUTIOMODE4\_ANY iomode is defined and can only be used for LAYOUTRETURN and CB\_LAYOUTRECALL, not for LAYOUTGET. It specifies that layouts pertaining to both LAYOUTIOMODE4\_READ and LAYOUTIOMODE4\_RW iomodes are being returned or recalled, respectively.

A data storage device can validate I/O for consistency with the iomode. Whether this happens depends on the layout type specification. In cases in which this is to be done, if the client's layout iomode is inconsistent with the I/O being performed, the data storage device may reject the client's I/O with an error indicating that a new layout with the correct iomode should be obtained via LAYOUTGET. For example, if a client gets a layout with a LAYOUTIOMODE4\_READ iomode and performs a WRITE to a data storage device, the device is allowed to reject that WRITE.

The use of the layout iomode does not conflict with OPEN share modes or byte-range LOCK operations; open share mode and byte-range lock conflicts are enforced as they are without the use of pNFS and are logically separate from issues related to pNFS layouts. Open share modes and byte-range locks are the preferred method for restricting user access to data files. For example, an OPEN of OPEN4\_SHARE\_ACCESS\_WRITE does not conflict with a LAYOUTGET containing an iomode of LAYOUTIOMODE4\_RW performed by another client. Applications that depend on writing into the same file concurrently may use byte-range locking to serialize their accesses.

#### 18.4.9. Device IDs

The device ID (data type deviceid4, see Section 9.3.14) identifies a group of storage devices. The scope of a device ID is the pair <client ID, layout type>. In practice, a significant amount of information may be required to fully address a storage device. Rather than embedding all such information in a layout, layouts embed device IDs. The NFSv4.1 operation GETDEVICEINFO (Section 25.40) is used to retrieve the complete address information (including all device addresses for the device ID) regarding the storage device according to its layout type and device ID. For example, the address of an NFSv4.1 data server or of an object-based storage device could be an IP address and port. The address of a block storage device could be a volume label.

Clients cannot expect the mapping between a device ID and its storage device address(es) to persist across metadata server restart. See Section 18.9.4 for a description of how recovery works in that situation.

A device ID is established by referencing it in the result of a GETDEVICELIST or LAYOUTGET operation and can be deleted by the server as soon as there are no layouts referring to the device ID.

If the client requested notifications for device ID mappings, the server should send CB\_NOTIFY\_DEVICEID notifications for device ID deletions or changes to the device-ID-to-device-address mappings to any client which has used the device-ID in question at least once, irrespective of whether the client has any layouts currently referring to it. If the server does not support or the client does not request notifications for device ID mappings, the client should periodically retired unused device IDs.

Given that GETDEVICELIST does not support requesting notifications a server that implements GETDEVICELIST MUST NOT advertise support for NOTIFY\_DEVICEID4\_CHANGE notification in GETDEVICEINFO, and client using GETDEVICELIST can not rely on NOTIFY\_DEVICEID4\_CHANGE or NOTIFY\_DEVICEID4\_DELETE notifications to work reliably.

Once a device ID is deleted by the server, the server MUST NOT reuse the device ID for the same layout type and client ID again. This requirement is feasible because the device ID is 16 bytes long, leaving sufficient room to store a generation number if the server's implementation requires most of the rest of the device ID's content to be reused. This requirement is necessary because otherwise the race conditions between asynchronous notification of device ID addition and deletion would be too difficult to sort out.

Device ID to device address mappings are not leased, and can be changed at any time. (Note that while device ID to device address mappings are likely to change after the metadata server restarts, the server is not required to change the mappings.) A server has two choices for changing mappings. It can recall all layouts referring to the device ID or it can use a notification mechanism.

The NFSv4.1 protocol has no optimal way to recall all layouts that referred to a particular device ID (unless the server associates a single device ID with a single fsid or a single client ID; in which case, CB\_LAYOUTRECALL has options for recalling all layouts associated with the fsid, client ID pair, or just the client ID).

Via a notification mechanism (See Section 27.12), device ID to device address mappings can change over the duration of server operation without recalling or revoking the layouts that refer to device ID. The notification mechanism can also delete a device ID, but only if the client has no layouts referring to the device ID. A notification of a change to a device ID to device address mapping will immediately or eventually invalidate some or all of the device ID's mappings.

The server MUST support notifications and the client must request them before they can be used. For further information about the notification types, see Section 27.12.

### 18.5. pNFS Operations

NFSv4.1 has several operations that are needed for pNFS servers, regardless of layout type or storage protocol. These operations are all sent to a metadata server and summarized here. While pNFS is an OPTIONAL feature, if pNFS is implemented, some operations are REQUIRED in order to comply with pNFS. See Section 24.

These are the fore channel pNFS operations:

GETDEVICEINFO (Section 25.40), as noted previously (Section 18.4.9), returns the mapping of device ID to storage device address.

GETDEVICELIST (Section 25.41) allows clients to fetch all device IDs for a specific file system.

LAYOUTGET (Section 25.43) is used by a client to get a layout for a file.

LAYOUTCOMMIT (Section 25.42) is used to:

- \* Inform the metadata server of the client's intent to commit data that has been written to the storage device (the storage device as originally indicated in the return value of LAYOUTGET), in situations in which this information needs to be provided to the metadata server,, as specified by the file layout type.
- \* Provide for the updating of metadata server's view of file attributes that can modified by WRITES to incorporate WRITES effected by using a data storage device. These attributes include size, modified\_time, and change.

LAYOUTRETURN (Section 25.44) is used to return layouts for a file, a file system ID (FSID), or a client ID.

These are the backchannel pNFS operations:

CB\_LAYOUTRECALL (Section 27.3) recalls a layout, all layouts belonging to a file system, or all layouts belonging to a client ID.

CB\_RECALL\_ANY (Section 27.6) tells a client that it needs to return

some number of recallable objects, including layouts, to the metadata server.

CB\_RECALLABLE\_OBJ\_AVAIL (Section 27.7) tells a client that a recallable object that it was denied (in case of pNFS, a layout denied by LAYOUTGET) due to resource exhaustion is now available.

CB\_NOTIFY\_DEVICEID (Section 27.12) notifies the client of changes to device IDs.

#### 18.6. pNFS Attributes

A number of attributes specific to pNFS are listed and described in Section 11.16.

#### 18.7. Layout Semantics

##### 18.7.1. Guarantees Provided by Layouts

Layouts grant to the client the ability to access data located at a data storage device using the associated data access protocol. The client is guaranteed the layout will be recalled when one of two things occur: either a conflicting layout is requested or the state encapsulated by the layout becomes invalid (this can happen when an event directly or indirectly modifies the layout). When a layout is recalled and returned by the client, the client continues with the ability to access file data with normal NFSv4.1 operations through the metadata server. Only the ability to access the file using the data storage device is affected.

The requirement of NFSv4.1 that all user access rights MUST be obtained through the appropriate OPEN, LOCK, and ACCESS operations is not modified with the existence of layouts. Layouts are provided to NFSv4.1 clients, and user access still follows the rules of the protocol as if they did not exist. It is a requirement that for a client to access a data storage device, a layout must be held by the client. If a device receives an I/O request for a byte-range for which the client does not hold a layout, the storage device MUST reject that I/O request. Note that the act of modifying a file for which a layout is held does not necessarily conflict with the holding of the layout that describes the file being modified. Therefore, it is the requirement of the data access protocol or layout type that determines the necessary behavior. For example, block/volume layout types require that the layout's iomode agree with the type of I/O being performed.

Depending upon the layout type and data access protocol in use, device-specific access permissions may be granted by LAYOUTGET and may be encoded within the type-specific layout. For an example of storage device access permissions, see an object-based protocol such as [OSD-T10]. If access permissions are encoded within the layout, the metadata server needs to recall the layout when those permissions become invalid for any reason -- for example, when a file becomes unwritable or inaccessible to a client. Note, clients are still required to perform the appropriate OPEN, LOCK, and ACCESS operations as described above. The degree to which it is possible for the client to circumvent these operations and the consequences of doing so must be clearly specified by the individual layout type specifications. In addition, these specifications must be clear about the requirements and non-requirements for the checking performed by the server.

In the presence of pNFS functionality, mandatory byte-range locks MUST behave as they would without pNFS. Therefore, if mandatory file locks and layouts are provided simultaneously, the storage device MUST be able to enforce the mandatory byte-range locks. For example, if one client obtains a mandatory byte-range lock and a second client accesses the on the data storage device, the device MUST appropriately restrict I/O for the range of the mandatory byte-range lock. If the device is incapable of providing this check in the presence of mandatory byte-range locks, then the metadata server MUST NOT grant potentially overlapping layouts and mandatory byte-range locks simultaneously.

#### 18.7.2. Getting a Layout

A client obtains a layout using the LAYOUTGET operation. The metadata server will grant layouts of a particular type (e.g., block/volume, object, or file). The client selects an appropriate layout type that the server supports and the client is prepared to use. The layout returned to the client might not exactly match the requested byte-range as described in Section 25.43.3. As needed, a client may send multiple LAYOUTGET operations. These might result in multiple overlapping, non-conflicting layouts (see Section 18.4.7).

In order to get a layout, the client must first have opened the file via the OPEN operation. When a client has no layout on a file, it MUST present an open stateid, a delegation stateid, or a byte-range lock stateid in the loga\_stateid argument. A successful LAYOUTGET result includes a layout stateid. The first successful LAYOUTGET processed by the server using a non-layout stateid as an argument MUST have the "seqid" field of the layout stateid in the response set to one. Thereafter, the client MUST use a layout stateid (See Section 18.7.3) on future invocations of LAYOUTGET on the file, and

the "seqid" MUST NOT be set to zero. The client MUST serialize LAYOUTGET operations using a non-layout stateid with any other operation affecting the layout state on the file, including CB\_LAYOUTRECALL, to allow consistent initialization of the layout state. Once the layout has been retrieved, it can be held across multiple OPEN and CLOSE sequences. Therefore, a client may hold a layout for a file that is not currently open by any user on the client. This allows for the caching of layouts beyond CLOSE.

The storage protocol used by the client to access the data on the storage device is determined by the layout's type. The client is responsible for matching the layout type with an available method to interpret and use the layout. The method for this layout type selection is outside the scope of the pNFS functionality.

Although the metadata server is in control of the layout for a file, the pNFS client can provide hints to the server when a file is opened or created about the preferred layout type and aggregation schemes. pNFS introduces a layout\_hint attribute (Section 11.16.4) that the client can set at file creation time to provide a hint to the server for new files. Setting this attribute separately, after the file has been created might make it difficult, or impossible, for the server implementation to comply.

Because the EXCLUSIVE4 createmode4 does not allow the setting of attributes at file creation time, NFSv4.1 introduces the EXCLUSIVE4\_1 createmode4, which does allow attributes to be set at file creation time. In addition, if the session is created with persistent reply caches, EXCLUSIVE4\_1 is neither necessary nor allowed. Instead, GUARDED4 both works better and is prescribed. Table 18 in Section 25.16.3 summarizes how a client is allowed to send an exclusive create.

#### 18.7.3. Layout Stateid

As with all other stateids, the layout stateid consists of a "seqid" and "other" field. Once a layout stateid is established, the "other" field will stay constant unless the stateid is revoked or the client returns all layouts on the file and the server disposes of the stateid. The "seqid" field is initially set to one, and is never zero on any NFSv4.1 operation that uses layout stateids, whether it is a fore channel or backchannel operation. After the layout stateid is established, the server increments by one the value of the "seqid" in each subsequent LAYOUTGET and LAYOUTRETURN response, and in each CB\_LAYOUTRECALL request.



Given the design goal of pNFS to provide parallelism, the layout stateid differs from other stateid types in that the client is expected to send LAYOUTGET and LAYOUTRETURN operations in parallel. The "seqid" value is used by the client to properly sort responses to LAYOUTGET and LAYOUTRETURN. The "seqid" is also used to prevent race conditions between LAYOUTGET and CB\_LAYOUTRECALL. Given that the processing rules differ from layout stateids and other stateid types, only the pNFS sections of this document should be considered to determine proper layout stateid handling.

Once the client receives a layout stateid, it MUST use the correct "seqid" for subsequent LAYOUTGET or LAYOUTRETURN operations. The correct "seqid" is defined as the highest "seqid" value from responses of fully processed LAYOUTGET or LAYOUTRETURN operations or arguments of a fully processed CB\_LAYOUTRECALL operation. Since the server is incrementing the "seqid" value on each layout operation, the client may determine the order of operation processing by inspecting the "seqid" value. In the case of overlapping layout ranges, the ordering information will provide the client the knowledge of which layout ranges are held. Note that overlapping layout ranges may occur because of the client's specific requests or because the server is allowed to expand the range of a requested layout and notify the client in the LAYOUTRETURN results. Additional layout stateid sequencing requirements are provided in Section 18.7.5.2.

The client's receipt of a "seqid" is not sufficient for it to be passed back to the metadata server. The client needs to fully process the operation in which the new seqid is seen before using it in further communication with the metadata server. For LAYOUTGET results, if the client is recording details of layout ranges received (See (Section 18.7.3.1 for information about specifics), it MUST first update its record of what ranges of the file's layout it has before using the seqid in this way. For LAYOUTRETURN results, the client MUST eliminate the range from its record of what ranges of the file's layout it had before using the seqid. For CB\_LAYOUTRECALL arguments, the client MUST send a response to the recall before using the seqid in a message to the metadata server. The fundamental basis of these requirements regarding client processing is that the "seqid" is used to define the order of processing. LAYOUTGET results may be processed in parallel. LAYOUTRETURN results may be processed in parallel. LAYOUTGET and LAYOUTRETURN responses may be processed in parallel as long as their ranges do not overlap. CB\_LAYOUTRECALL request processing MUST be processed in "seqid" order at all times.

Once a client has no more layouts on a file, the layout stateid is no longer valid and MUST NOT be used. Any attempt to use such a layout stateid will result in NFS4ERR\_BAD\_STATEID.

A client MAY always forget its layout state and associated layout stateid at any time (See also Section 18.7.3.1). When this happens, the client MUST use a non-layout stateid on a subsequent LAYOUTGET operation. Doing so will inform the server that the client has no more layouts of that layout type on the file and that its respective layout state can be released before issuing a new layout in response to LAYOUTGET.

#### 18.7.3.1. Requirements Regarding Retention of Layout Information.

It might reasonably be assumed that pNFS client state (layout ranges and iomode) for a file would exactly matches that of the pNFS metadata server for that file. This assumption would imply that any callback of a layout results in a LAYOUTRETURN or set of LAYOUTRETURNS that exactly match the range specified in the callback, since the client and metadata server would necessarily agree about the details of the state being maintained.

It is important to understand that the above assumption is not a protocol requirement and that the protocol does allow the client and metadata server to drop much of this information, in order to limit the need for additional storage and internode communication that maintaining the abovementioned assumption would require.

The obligations of the client and metadata server regarding the retention of layout information are discussed below. It is important to understand that data servers often need to be aware of the details of layouts gotten and that individual layout types might impose requirements regarding their retention by the client and how they might be used by the specific storage protocol used by the layout type.

The relevant protocol requirements discussed here are limited to those necessary to effect cancellation of layouts to prevent multiple clients using conflicting layouts and to prevent use of layouts when changes initiated by the metadata server make their further use incorrect or otherwise inadvisable. The relevant requirements and non-requirements can be summarized as follows:

- \* Within the scope of the base pNFS feature, and the use of LAYOUTRETURN and CB\_LAYOUTRECALL, both the client and the metadata server are free to drop detailed information about layout ranges and IO modes at any time.

Any restrictions on the client in this regard derive from the layout type used and its possible need for the client to have such information in order to effect IO operations.

The client MAY do so at any time as long as it does not use these dropped ranges subsequently or drop them while IO operations based on them are still being processed.

Even with the requirement above applying to the client, it is possible that I/O requests may be presented to a storage device no longer allowed to perform them. Since the server does not have strict control as to when client will return a recalled layout, the server needs to be able to unilaterally terminate the client's access to the storage devices as described by the layout. In terminating access, the server needs to deal with the possibility of lingering I/O requests, i.e., I/O requests that are still in flight to storage devices identified the cancelled layout. All layout type specifications need to specify whether such unilateral layout revocation by the metadata server is supported; if it is, the specification needs to also describe how lingering writes are dealt with. For example, storage devices identified by the canceled layout could be fenced off from the client that held the layout or the data server could be instructed to suppress access that was previously allowed by the cancelled layout. For details, see Section 19.1.9.

Similarly, the metadata server MAY also discard detailed information about layout ranges and IO modes. However, in doing so, it MUST NOT forget that it could have outstanding layouts for any part of the file so that it is unaware of layouts that might be retained by the clients. When that knowledge is lost, the layout MUST be returned or revoked so that it cannot be used subsequently.

- \* The freedom to forget layout details referred to above can be extended to allow the client to eliminate its knowledge of the existence of a layout for a particular file. In doing so, it must ensure that no IOs will use that layout subsequently
- \* Because of clients and server are allowed to independently discard layouts they were previously jointly aware of, the processing of layout recalls is more complicated than it would otherwise be. See Section 18.7.5.1 for further discussion.

The remainder of this section discusses how clients and metadata servers might deal with cases in which the other party chooses not to maintain detailed knowledge of layouts that it is aware of and might recall. For example,

- \* In situations in which conflicts that require callbacks are very rare, a server can use a multi-file callback to recover per-client resources (e.g., via an FSID recall or a multi-file recall within

a single CB\_COMPOUND). In this case, the result may be significantly less client-server pNFS traffic, although layouts might be recalled even if there might be no need to do so.

- \* It might also be useful for servers to maintain information about what ranges are held by a client on a coarse-grained basis, leading to the server's layout ranges being beyond those actually held by the client.
- \* In one possible extreme case, a server could manage conflicts on a per-file basis, only sending whole-file callbacks even though clients may request and be granted sub-file ranges.

#### 18.7.4. Committing a Layout

The pNFS protocol does not require the metadata server and storage devices to maintain a consistent view of file attributes that can be modified as a result of IO operations sent to data storage device. How each layout type provides for these changes to be reflected in the metadata server's view is described by the layout type specification, as described in Section 19.1.8

Some layout types also need to coordinate allocation-related data, including data location mapping which refers to aspects such as which offsets store data as opposed to storing holes (See Section 20.8.4 for a discussion). Related issues arise for storage protocols where a layout may hold provisionally allocated blocks where the allocation of those blocks does not survive a complete restart of both the client and server. Because of the potential for inconsistency between the client and the metadata server, it is necessary, in general, to resynchronize the client with the metadata server and its storage devices and make any potential changes available to other clients. This is accomplished by use of the LAYOUTCOMMIT operation. See the requirements described in Section 19.1.8 for the layout specification's obligations to describe and provide for the coordination of such items.

The LAYOUTCOMMIT operation is responsible for making globally visible the effects of layout-based file modifications as they affect the metadata server. The data should be written and committed to the appropriate storage devices before the LAYOUTCOMMIT occurs. The scope of data committed by a LAYOUTCOMMIT operation is specific to the type of layout because that scope depends on the storage protocol in use. It is important to note that the level of synchronization is from the point of view of the client that sent the LAYOUTCOMMIT. The updated state on the metadata server need only reflect the state as of the client's last operation previous to the LAYOUTCOMMIT. The metadata server is need not maintain a global view that accounts for other clients' I/O that may have occurred within the same time frame.

#### 18.7.5. Recalling a Layout

Since a layout protects a client's access to a file via a direct client-storage-device path, a layout need only be recalled when it is semantically unable to serve this function. Typically, this occurs when the layout no longer encapsulates the true location of the file over the byte-range it represents. Any operation or action, such as server-driven restriping or load balancing, that changes the layout will result in a recall of the layout. A layout is recalled by the CB\_LAYOUTRECALL callback operation (See Section 27.3) and returned with LAYOUTRETURN (See Section 25.44). The CB\_LAYOUTRECALL operation may recall a layout identified by a byte-range, all layouts associated with a file system ID (FSID), or all layouts associated with a client ID. Section 18.7.5.2 discusses sequencing issues surrounding the getting, returning, and recalling of layouts.

An iomode is also specified when recalling a layout. Generally, the iomode in the recall request must match the layout being returned; for example, a recall with an iomode of LAYOUTIOMODE4\_RW should cause the client to only return LAYOUTIOMODE4\_RW layouts and not LAYOUTIOMODE4\_READ layouts. However, a special LAYOUTIOMODE4\_ANY enumeration is defined to enable recalling a layout of any iomode; in other words, the client must return both LAYOUTIOMODE4\_READ and LAYOUTIOMODE4\_RW layouts.

A REMOVE operation need to make sure that existing layouts cannot be used access a non-existent or to reclaim. How this is to be done us defined by the layout type specification for the existing layout.

Some examples follow:

- \* In many cases, metadata server MUST recall the layout to prevent its subsequent use. Since a REMOVE may be delayed until the last close of the file has occurred, the recall may also be delayed until this time. After the last reference on the file has been released and the file has been removed, the client should no longer be able to perform I/O using the layout.
- \* In the case of a file-based layout, the data server returns NFS4ERR\_STALE in response to any operation on the removed file, making prompt recall of the layout unnecessary.

Later recall of the layout might be done later, as part server cleanup to object unlikely to be referenced.

Once a layout has been returned, the client MUST NOT send IO requests to the storage devices for the file, byte-range, and iomode represented by the returned layout. If a client does send an I/O to a storage device for which it does not hold a layout, the storage device will reject the I/O. The ability of and requirements for storage devices

#### 18.7.5.1. Layout recall/return interaction

As noted in Section 18.7.3.1, It may be useful for clients to "forget" details about what layouts and ranges the client actually has, leading to the server's layout ranges being beyond those that the client "thinks" it has.

In light of the above, it is useful for a server to be able to send callbacks for layout ranges it has not granted to a client, and for a client to return ranges it does not hold. A pNFS client MUST always return layouts that comprise the full range specified by the recall. Note, the full recalled layout range need not be returned as part of a single operation, but may be returned in portions. This allows the client to properly stage the flushing of dirty data and commits and returns of layouts. Also, it indicates to the metadata server that the client is making progress.

As long as the client does not assume it has layouts that are beyond what the server has granted, this is a safe practice. When a client forgets what ranges and layouts it has, and it receives a CB\_LAYOUTRECALL operation, the client follows up with a LAYOUTRETURN for what the server recalled (even though what was held does not match what being recalled) or alternatively return the NFS4ERR\_NOMATCHING\_LAYOUT error if it has no layout to return in the recalled range.

In order to avoid errors, it is vital that a client not assign itself layout permissions beyond what the server has granted, and that the server not forget layout permissions that have been granted. On the other hand, if a server believes that a client holds a layout that the client does not know about, it is useful for the client to cleanly indicate completion of the requested recall either by sending a LAYOUTRETURN operation for the entire requested range (even though it is not actually being "returned") or by returning an NFS4ERR\_NOMATCHING\_LAYOUT error to the CB\_LAYOUTRECALL.

In order to ensure client/server convergence with regard to layout state, the final LAYOUTRETURN operation in a sequence of LAYOUTRETURN operations for a particular recall MUST specify the entire range being recalled, echoing the recalled layout type, iomode, recall/return type (FILE, FSID, or ALL), and byte-range, even if layouts pertaining to partial ranges were previously returned. In addition, if the client holds no layouts that overlap the range being recalled, the client should return the NFS4ERR\_NOMATCHING\_LAYOUT error code to CB\_LAYOUTRECALL. This allows the server to update its view of the client's layout state.

Note that, in the case in which FSID or ALL are specified, the client needs to do a LAYOUTRETURN for all files for which it holds a layout stateid but need not do so for files for which it does not, even though the metadata server might be aware of these stateids.

#### 18.7.5.2. Sequencing of Layout Operations

As with other stateful operations, pNFS requires the correct sequencing of layout operations. pNFS uses the "seqid" in the layout stateid to provide the correct sequencing between regular operations and callbacks. It is the server's responsibility to avoid inconsistencies regarding the layouts provided and the client's responsibility to properly serialize its layout requests and layout returns.

#### 18.7.5.3. Layout Recall and Return Sequencing

One critical issue with regard to layout operations sequencing concerns callbacks. The protocol must defend against races between the reply to a LAYOUTGET or LAYOUTRETURN operation and a subsequent CB\_LAYOUTRECALL. A client MUST NOT process a CB\_LAYOUTRECALL that implies one or more outstanding LAYOUTGET or LAYOUTRETURN operations to which the client has not yet received a reply. The client detects such a CB\_LAYOUTRECALL by examining the "seqid" field of the recall's layout stateid. If the "seqid" is not exactly one higher than what the client currently has recorded, and the client has at least one LAYOUTGET and/or LAYOUTRETURN operation outstanding, or if the client

has a outstanding LAYOUTGET with a non-layout stateid, the client knows the server sent the CB\_LAYOUTRECALL after sending a response to an outstanding LAYOUTGET or LAYOUTRETURN. The client MUST wait before processing such a CB\_LAYOUTRECALL until it processes all replies for outstanding LAYOUTGET and LAYOUTRETURN operations for the corresponding file with seqid less than the seqid given by CB\_LAYOUTRECALL (lor\_stateid; see Section 27.3.)

In addition to the seqid-based mechanism, Section 7.6.3 describes the sessions mechanism for allowing the client to detect callback race conditions and delay processing such a CB\_LAYOUTRECALL. The server MAY reference conflicting operations in the CB\_SEQUENCE that precedes the CB\_LAYOUTRECALL. Because the server has already sent replies for these operations before sending the callback, the replies may race with the CB\_LAYOUTRECALL. The client MUST wait for all the referenced calls to complete and update its view of the layout state before processing the CB\_LAYOUTRECALL.

#### 18.7.5.3.1. Get/Return Sequencing

The protocol allows the client to send concurrent LAYOUTGET and LAYOUTRETURN operations to the server. The protocol does not provide any means for the server to process the requests in the same order in which they were created. However, through the use of the "seqid" field in the layout stateid, the client can determine the order in which parallel outstanding operations were processed by the server. Thus, when a layout retrieved by an outstanding LAYOUTGET operation intersects with a layout returned by an outstanding LAYOUTRETURN on the same file, the order in which the two conflicting operations are processed determines the final state of the overlapping layout. The order is determined by the "seqid" returned in each operation: the operation with the higher seqid was executed later.

It is permissible for the client to send multiple parallel LAYOUTGET operations for the same file or multiple parallel LAYOUTRETURN operations for the same file or a mix of both. It is permissible for the client to send multiple parallel LAYOUTGET operations for the same file using the layout stateid or multiple parallel LAYOUTRETURN operations for the same file or a mix of both.

It is permissible for the client to use the current stateid (see Section 23.2.3.1.2) for LAYOUTGET operations, for example, when compounding LAYOUTGETs or compounding OPEN and LAYOUTGETs. It is also permissible to use the current stateid when compounding LAYOUTRETURNs.



It is permissible for the client to use the current stateid when combining LAYOUTRETURN and LAYOUTGET operations for the same file in the same COMPOUND request since the server MUST process these in order. However, if a client does send such COMPOUND requests, it MUST NOT have more than one outstanding for the same file at the same time, and it MUST NOT have other LAYOUTGET or LAYOUTRETURN operations outstanding at the same time for that same file.

#### 18.7.5.3.2. Client Considerations

Consider a pNFS client that has sent a LAYOUTGET, and before it receives the reply to LAYOUTGET, it receives a CB\_LAYOUTRECALL for the same file with an overlapping range. There are two possibilities, which the client can distinguish via the layout stateid in the recall.

1. The server processed the LAYOUTGET before sending the recall, so the LAYOUTGET must be waited for because it may be carrying layout information that will need to be returned to deal with the CB\_LAYOUTRECALL.
2. The server sent the callback before receiving the LAYOUTGET. The server will not respond to the LAYOUTGET until the CB\_LAYOUTRECALL is processed.

If these possibilities cannot be distinguished, a deadlock could result, as the client must wait for the LAYOUTGET response before processing the recall in the first case, but that response will not arrive until after the recall is processed in the second case. Note that in the first case, the "seqid" in the layout stateid of the recall is two greater than what the client has recorded, or the client has an outstanding LAYOUTGET using a non-layout stateid; in the second case, the "seqid" is one greater than what the client has recorded. This allows the client to disambiguate between the two cases. The client thus knows precisely which possibility applies.

In case 1, the client knows it needs to wait for the LAYOUTGET response before processing the recall (or the client can return NFS4ERR\_DELAY).

In case 2, the client will not wait for the LAYOUTGET response before processing the recall because waiting would cause deadlock. Therefore, the action at the client will only require waiting in the case that the client has not yet seen the server's earlier responses to the LAYOUTGET operation(s).

The recall process can be considered completed when the final LAYOUTRETURN operation for the recalled range is completed. The LAYOUTRETURN uses the layout stateid (with seqid) specified in CB\_LAYOUTRECALL. If the client uses multiple LAYOUTRETURNS in processing the recall, the first LAYOUTRETURN will use the layout stateid as specified in CB\_LAYOUTRECALL. Subsequent LAYOUTRETURNS will use the highest seqid as is the usual case.

#### 18.7.5.3.3. Server Considerations

Consider a race from the metadata server's point of view. The metadata server has sent a CB\_LAYOUTRECALL and receives an overlapping LAYOUTGET for the same file before the LAYOUTRETURN(s) that respond to the CB\_LAYOUTRECALL. There are three cases:

1. The client sent the LAYOUTGET before processing the CB\_LAYOUTRECALL. The "seqid" in the layout stateid of the arguments of LAYOUTGET is one less than the "seqid" in CB\_LAYOUTRECALL. The server returns NFS4ERR\_RECALLCONFLICT to the client, which indicates to the client that there is a pending recall.
2. The client sent the LAYOUTGET after processing the CB\_LAYOUTRECALL, but the LAYOUTGET arrived before the LAYOUTRETURN and the response to CB\_LAYOUTRECALL that completed that processing. The "seqid" in the layout stateid of LAYOUTGET is equal to or greater than that of the "seqid" in CB\_LAYOUTRECALL. The server has not received a response to the CB\_LAYOUTRECALL, so it returns NFS4ERR\_RECALLCONFLICT.
3. The client sent the LAYOUTGET after processing the CB\_LAYOUTRECALL; the server received the CB\_LAYOUTRECALL response, but the LAYOUTGET arrived before the LAYOUTRETURN that completed that processing. The "seqid" in the layout stateid of LAYOUTGET is equal to that of the "seqid" in CB\_LAYOUTRECALL. The server has received a response to the CB\_LAYOUTRECALL, so it returns NFS4ERR\_RETURNCONFLICT.

#### 18.7.5.3.4. Wraparound and Validation of Seqid

The rules for layout stateid processing differ from other stateids in the protocol because the "seqid" value cannot be zero and the stateid's "seqid" value changes in a CB\_LAYOUTRECALL operation. The non-zero requirement combined with the inherent parallelism of layout operations means that a set of LAYOUTGET and LAYOUTRETURN operations may contain the same value for "seqid". The server uses a slightly modified version of the modulo arithmetic as described in Section 7.6.1 when incrementing the layout stateid's "seqid". The

difference is that zero is not a valid value for "seqid"; when the value of a "seqid" is 0xFFFFFFFF, the next valid value will be 0x00000001. The modulo arithmetic is also used for the comparisons of "seqid" values in the processing of CB\_LAYOUTRECALL events as described above in Section 18.7.5.3.3.

Just as the server validates the "seqid" in the event of CB\_LAYOUTRECALL usage, as described in Section 18.7.5.3.3, the server also validates the "seqid" value to ensure that it is within an appropriate range. This range represents the degree of parallelism the server supports for layout stateids. If the client is sending multiple layout operations to the server in parallel, by definition, the "seqid" value in the supplied stateid will not be the current "seqid" as held by the server. The range of parallelism spans from the highest or current "seqid" to a "seqid" value in the past. To assist in the discussion, the server's current "seqid" value for a layout stateid is defined as SERVER\_CURRENT\_SEQID. The lowest "seqid" value that is acceptable to the server is represented by PAST\_SEQID. And the value for the range of valid "seqid"s or range of parallelism is VALID\_SEQID\_RANGE. Therefore, the following holds:  $\text{VALID\_SEQID\_RANGE} = \text{SERVER\_CURRENT\_SEQID} - \text{PAST\_SEQID}$ . In the following, all arithmetic is the modulo arithmetic as described above.

The server MUST support a minimum VALID\_SEQID\_RANGE. The minimum is defined as:  $\text{VALID\_SEQID\_RANGE} = \text{summation over } 1..N \text{ of } (\text{ca\_maxoperations}(i) - 1)$ , where N is the number of session fore channels and ca\_maxoperations(i) is the value of the ca\_maxoperations returned from CREATE\_SESSION of the i'th session. The reason for "-1" is to allow for the required SEQUENCE operation. The server MAY support a VALID\_SEQID\_RANGE value larger than the minimum. The maximum VALID\_SEQID\_RANGE is  $(2^{32} - 2)$  (accounting for zero not being a valid "seqid" value).

If the server finds the "seqid" is zero, the NFS4ERR\_BAD\_STATEID error is returned to the client. The server further validates the "seqid" to ensure it is within the range of parallelism, VALID\_SEQID\_RANGE. If the "seqid" value is outside of that range, the error NFS4ERR\_OLD\_STATEID is returned to the client. Upon receipt of NFS4ERR\_OLD\_STATEID, the client updates the stateid in the layout request based on processing of other layout requests and re-sends the operation to the server.

## 18.7.5.3.5. Bulk Recall and Return

pNFS supports recalling and returning all layouts that are for files belonging to a particular fsid (LAYOUTRECALL4\_FSID, LAYOUTRETURN4\_FSID) or client ID (LAYOUTRECALL4\_ALL, LAYOUTRETURN4\_ALL). There are no "bulk" stateids, so detection of races via the seqid is not possible. The server MUST NOT initiate bulk recall while another recall is in progress, or the corresponding LAYOUTRETURN is in progress or pending. In the event the server sends a bulk recall while the client has a pending or in-progress LAYOUTRETURN, CB\_LAYOUTRECALL, or LAYOUTGET, the client returns NFS4ERR\_DELAY. In the event the client sends a LAYOUTGET or LAYOUTRETURN while a bulk recall is in progress, the server returns NFS4ERR\_RECALLCONFLICT. If the client sends a LAYOUTGET or LAYOUTRETURN after the server receives NFS4ERR\_DELAY from a bulk recall, then to ensure forward progress, the server MAY return NFS4ERR\_RECALLCONFLICT.

Once a CB\_LAYOUTRECALL of LAYOUTRECALL4\_ALL is sent, the server MUST NOT allow the client to use any layout stateid except for LAYOUTCOMMIT operations. Once the client receives a CB\_LAYOUTRECALL of LAYOUTRECALL4\_FSID, it MUST NOT use any layout stateid except for LAYOUTCOMMIT operations. Once a LAYOUTRETURN of LAYOUTRETURN4\_ALL is sent, all layout stateids granted to the client ID are freed. The client MUST NOT use the layout stateids again. It MUST use LAYOUTGET to obtain new layout stateids.

Once a CB\_LAYOUTRECALL of LAYOUTRECALL4\_FSID is sent, the server MUST NOT allow the client to use any layout stateid that refers to a file with the specified fsid except for LAYOUTCOMMIT operations. Once the client receives a CB\_LAYOUTRECALL of LAYOUTRECALL4\_ALL, it MUST NOT use any layout stateid that refers to a file with the specified fsid except for LAYOUTCOMMIT operations. Once a LAYOUTRETURN of LAYOUTRETURN4\_FSID is sent, all layout stateids granted to the referenced fsid are freed. The client MUST NOT use those freed layout stateids for files with the referenced fsid again. Subsequently, for any file with the referenced fsid, to use a layout, the client MUST first send a LAYOUTGET operation in order to obtain a new layout stateid for that file.

If the server has sent a bulk CB\_LAYOUTRECALL and receives a LAYOUTGET, or a LAYOUTRETURN with a stateid, the server MUST return NFS4ERR\_RECALLCONFLICT. If the server has sent a bulk CB\_LAYOUTRECALL and receives a LAYOUTRETURN with an lr\_returntype that is not equal to the lor\_recalltype of the CB\_LAYOUTRECALL, the server MUST return NFS4ERR\_RECALLCONFLICT.

#### 18.7.6. Revoking Layouts

Parallel NFS allow servers to revoke layouts from clients that fail to respond to recalls and/or fail to renew their lease in time. Depending on the layout type, the layout type specification might allow the server to revoke the layout in other situations and describe the actions to be taken with respect to the client's I/O to data servers.

#### 18.7.7. Metadata Server Write Propagation

Asynchronous writes written through the metadata server may be propagated lazily to the storage devices. For data written asynchronously through the metadata server, a client performing a read at the appropriate storage device is not guaranteed to see the newly written data until a COMMIT occurs at the metadata server. While the write is pending, reads to the storage device may give out either the old data, the new data, or a mixture of new and old. Upon completion of a synchronous WRITE or COMMIT (for asynchronously written data), the metadata server MUST ensure that storage devices give out the new data and that the data has been written to stable storage. If the server implements its storage in any way such that it cannot obey these constraints, then it MUST recall the layouts to prevent reads being done that cannot be handled correctly. Note that the layouts MUST be recalled prior to the server responding to the associated WRITE operations.

#### 18.8. pNFS Mechanics

This section describes the operations flow taken by a pNFS client to a metadata server and storage device.

When a pNFS client encounters a new FSID, it sends a GETATTR to the NFSv4.1 server for the `fs_layout_type` (Section 11.16.1) attribute. If the attribute returns at least one layout type, and the layout types returned are among the set supported by the client, the client knows that pNFS is a possibility for the file system. If, from the server that returned the new FSID, the client does not have a client ID that came from an EXCHANGE\_ID result that returned `EXCHGID4_FLAG_USE_PNFS_MDS`, it MUST send an EXCHANGE\_ID to the server with the `EXCHGID4_FLAG_USE_PNFS_MDS` bit set. If the server's response does not have `EXCHGID4_FLAG_USE_PNFS_MDS`, then contrary to what the `fs_layout_type` attribute said, the server does not support pNFS, and the client will not be able use pNFS to that server; in this case, the server MUST return NFS4ERR\_NOTSUPP in response to any pNFS operation.

The client then creates a session, requesting a persistent session, so that exclusive creates can be done with single round trip via the createmode4 of GUARDED4. If the session ends up not being persistent, the client will use EXCLUSIVE4\_1 for exclusive creates.

If a file is to be created on a pNFS-enabled file system, the client uses the OPEN operation. With the normal set of attributes that may be provided upon OPEN used for creation, there is an OPTIONAL layout\_hint attribute. The client's use of layout\_hint allows the client to express its preference for a layout type and its associated layout details. The use of a createmode4 of UNCHECKED4, GUARDED4, or EXCLUSIVE4\_1 will allow the client to provide the layout\_hint attribute at create time. The client MUST NOT use EXCLUSIVE4 (See Table 18). The client is RECOMMENDED to combine a GETATTR operation after the OPEN within the same COMPOUND. The GETATTR may then retrieve the layout\_type attribute for the newly created file. The client will then know what layout type the server has chosen for the file and therefore what storage protocol the client must use.

If the client wants to open an existing file, then it also includes a GETATTR to determine what layout type the file supports.

The GETATTR in either the file creation or plain file open case can also include the layout\_blksize and layout\_alignment attributes so that the client can determine optimal offsets and lengths for I/O on the file.

Assuming the client supports the layout type returned by GETATTR and it chooses to use pNFS for data access, it then sends LAYOUTGET using the filehandle and stateid returned by OPEN, specifying the range it wants to do I/O on. The response is a layout, which may be a subset of the range for which the client asked. It also includes device IDs and a description of how data is organized (or in the case of writing, how data is to be organized) across the devices. The device IDs and data description are encoded in a format that is specific to the layout type, but the client is expected to understand.

When the client wants to send an I/O, it determines to which device ID it needs to send the I/O command by examining the data description in the layout. It then sends a GETDEVICEINFO to find the device address(es) of the device ID. The client then sends the I/O request to one of device ID's device addresses, using the storage protocol defined for the layout type. Note that if a client has multiple I/Os to send, these I/O requests may be done in parallel.

If the I/O is a WRITE, then at some point the client will need to use LAYOUTCOMMIT to commit the modification time and the new size of the file (if it believes it could have extended the file size) to the metadata server and, depending on the layout type, the modified data to the file system.

#### 18.9. Recovery

Recovery is complicated by the distributed nature of the pNFS protocol. In general, crash recovery for layouts is similar to crash recovery for delegations in the base NFSv4.1 protocol. However, the client's ability to perform I/O without contacting the metadata server introduces subtleties that must be handled correctly if the possibility of file system corruption is to be avoided.

##### 18.9.1. Recovery from Client Restart

Client recovery for layouts is similar to client recovery for other lock and delegation state. When a pNFS client restarts, it will lose all information about the layouts that it previously owned. There are two methods by which the server can reclaim these resources and allow otherwise conflicting layouts to be provided to other clients.

The first is through the expiry of the client's lease. If the client recovery time is longer than the lease period, the client's lease will expire and the server will know that state may be released. For layouts, the server may release the state immediately upon lease expiry or it may allow the layout to persist, awaiting possible lease revival, as long as no other layout conflicts.

The second is through the client restarting in less time than it takes for the lease period to expire. In such a case, the client will contact the server through the standard EXCHANGE\_ID protocol. The server will find that the client's co\_ownerid matches the co\_ownerid of the previous client invocation, but that the verifier is different. The server uses this as a signal to release all layout state associated with the client's previous invocation. In this scenario, the data written by the client but not covered by a successful LAYOUTCOMMIT is in an undefined state; it may have been written or it may now be lost. This is acceptable behavior, making it necessary, depending on the layout type, for the client to use LAYOUTCOMMIT or other means to achieve the desired level of stability.

### 18.9.2. Dealing with Lease Expiration on the Client

If a client has reason to believe its lease has expired, it MUST NOT send I/O to the storage device until it has validated its lease. The client can send a SEQUENCE operation to the metadata server. If the SEQUENCE operation is successful, but `sr_status_flag` has `SEQ4_STATUS_EXPIRED_ALL_STATE_REVOKED`, `SEQ4_STATUS_EXPIRED_SOME_STATE_REVOKED`, or `SEQ4_STATUS_ADMIN_STATE_REVOKED` set, the client MUST NOT use currently held layouts. The client has two choices to recover from the lease expiration. First, for all modified but uncommitted data, the client writes it to the metadata server using the `FILE_SYNC4` flag for the WRITES, or WRITE and COMMIT. Second, the client re-establishes a client ID and session with the server and obtains new layouts and device-ID-to-device-address mappings for the modified data ranges and then writes the data to the storage devices with the newly obtained layouts.

If `sr_status_flags` from the metadata server has `SEQ4_STATUS_RESTART_RECLAIM_NEEDED` set (or SEQUENCE returns `NFS4ERR_BAD_SESSION` and `CREATE_SESSION` returns `NFS4ERR_STALE_CLIENTID`), then the metadata server has restarted, and the client should recover using the methods described in Section 18.9.4.

If `sr_status_flags` from the metadata server has `SEQ4_STATUS_LEASE_MOVED` set, then the client recovers by following the procedure described in Section 17.11.9.2. After that, the client may get an indication that the layout state was not moved with the file system. The client recovers as in the other applicable situations discussed in the first two paragraphs of this section.

If `sr_status_flags` reports no loss of state, then the lease for the layouts that the client has are valid and renewed, and the client can once again send I/O requests to the storage devices.

While clients should not send I/Os to storage devices that may extend past the lease expiration time period, enforcing this restriction is not always possible. For example, an extended network partition that starts after the I/O is sent and does not heal until the I/O request is received by the storage device. Thus, the metadata server and/or storage devices are responsible for protecting themselves from I/Os that are both sent before the lease expires and arrive after the lease expires. How this is to be done needs to be addressed by the layout type specification. See Section 18.9.3 for further discussion of the layout type specification's obligation in this regard



### 18.9.3. Dealing with Loss of Layout State on the Metadata Server

This is a description of the case where all of the following are true:

- \* the metadata server has not restarted
- \* a pNFS client's layouts have been discarded (usually because the client's lease expired) and are invalid
- \* an I/O from the pNFS client arrives at the storage device

The metadata server and its storage devices need to address this situation by fencing the client as described by the layout type specification. In other words, they **MUST** prevent the execution of I/O operations from the client to the storage devices after layout state loss.

The details of how fencing is done are specific to the layout type and need to be described in the layout specification, as required by Section 19.1.9. In some cases the control protocol can be used to effectively revoke the data server's record of the layout allowing it to be rejected. In other cases, the means available are not limited to specific layouts and have a force majeure character. However, it is always necessary in practice to prevent further prohibited access by the storage device, even it appears unresponsive. The means by which this is need not explicitly discussed but it should be clear that drastic action might be necessary to provide the requisite data integrity.

### 18.9.4. Recovery from Metadata Server Restart

The pNFS client will discover that the metadata server has restarted via the methods described in Section 13.4.2 and discussed in a pNFS-specific context in Section 18.9.2, Paragraph 2. The client **MUST** stop using layouts and delete the device ID to device address mappings it previously received from the metadata server. Having done that, if the client wrote data to the storage device without committing the layouts via LAYOUTCOMMIT, then the client has additional work to do in order to have the client, metadata server, and storage device(s) all synchronized on the state of the data.

- \* If the client has data still modified and unwritten in the client's memory, the client has only two choices.
  1. The client can obtain a layout via LAYOUTGET after the server's grace period and write the data to the storage devices.

2. The client can WRITE that data through the metadata server using the WRITE (Section 25.32) operation, and then obtain layouts as desired.

- \* If the client asynchronously wrote data to the storage device, but still has a copy of the data in its memory, then it has available to it the recovery options listed above in the previous bullet point. If the metadata server is also in its grace period, the client has available to it the options below in the next bullet point.
- \* The client does not have a copy of the data in its memory and the metadata server is still in its grace period. The client cannot use LAYOUTGET (within or outside the grace period) to reclaim a layout because the contents of the response from LAYOUTGET may not match what it had previously. The range might be different or the client might get the same range but the content of the layout might be different. Even if the content of the layout appears to be the same, the device IDs may map to different device addresses, and even if the device addresses are the same, the device addresses could have been assigned to a different storage device. The option of retrieving the data from the storage device and writing it to the metadata server per the recovery scenario described above is not available because, again, the mappings of range to device ID, device ID to device address, and device address to physical device are stale, and new mappings via new LAYOUTGET do not solve the problem.

The only recovery option for this scenario is to send a LAYOUTCOMMIT in reclaim mode, which the metadata server will accept as long as it is in its grace period. The use of LAYOUTCOMMIT in reclaim mode informs the metadata server that the layout has changed. It is critical that the metadata server receive this information before its grace period ends, and thus before it starts allowing updates to the file system.

To send LAYOUTCOMMIT in reclaim mode, the client sets the `loca_reclaim` field of the operation's arguments (Section 25.42.1) to TRUE. During the metadata server's recovery grace period (and only during the recovery grace period) the metadata server is prepared to accept LAYOUTCOMMIT requests with the `loca_reclaim` field set to TRUE.

When `loca_reclaim` is TRUE, the client is attempting to commit changes to the layout that occurred prior to the restart of the metadata server. The metadata server applies some consistency checks on the `loca_layoutupdate` field of the arguments to determine whether the client can commit the data written to the

storage device to the file system. The `loca_layoutupdate` field is of data type `layoutupdate4` and contains layout-type-specific content (in the `lou_body` field of `loca_layoutupdate`). The layout-type-specific information that `loca_layoutupdate` might have is defined in the layout type specification as described in Section 19.1.2. If the metadata server's consistency checks on `loca_layoutupdate` succeed, then the metadata server MUST commit the changed data that was written to the storage device within the scope of the `LAYOUTCOMMIT` operation. If the metadata server's consistency checks on `loca_layoutupdate` fail, the metadata server rejects the `LAYOUTCOMMIT` operation and makes no changes to the file system. However, any time `LAYOUTCOMMIT` with `loca_reclaim TRUE` fails, the pNFS client may have lost all uncommitted data within the scope of the failed `LAYOUTCOMMIT` operation. A client can defend against this risk by caching all data, whether written synchronously or asynchronously in its memory, and by not releasing the cached data until a successful `LAYOUTCOMMIT`. This condition does not hold true for all layout types; for example, file-based storage devices need not suffer from this limitation.

- \* The client does not have a copy of the data in its memory and the metadata server is no longer in its grace period; i.e., the metadata server returns `NFS4ERR_NO_GRACE`. As with the scenario in the above bullet point, the failure of `LAYOUTCOMMIT` means the data in the scope of that `LAYOUTCOMMIT` may have been lost. The defense against the risk is the same -- cache all written data on the client until a successful `LAYOUTCOMMIT`

#### 18.9.5. Operations during Metadata Server Grace Period

Some of the recovery scenarios thus far noted that some operations (namely, `WRITE` and `LAYOUTGET`) might be permitted during the metadata server's grace period. The metadata server may allow these operations during its grace period. For `LAYOUTGET`, the metadata server must reliably determine that servicing such a request will not conflict with an impending `LAYOUTCOMMIT` reclaim request. For `WRITE`, the metadata server must reliably determine that servicing the request will not conflict with an impending `OPEN` or with a `LOCK` where the file has mandatory byte-range locking enabled.

As mentioned previously, for expediency, the metadata server might reject some operations (namely, `WRITE` and `LAYOUTGET`) during its grace period, because the simplest correct approach is to reject all non-reclaim pNFS requests and `WRITE` operations by returning the `NFS4ERR_GRACE` error. However, depending on the storage protocol (which is specific to the layout type) and metadata server implementation, the metadata server may be able to determine that a particular request is safe. For example, a metadata server may save

provisional allocation mappings for each file to stable storage, as well as information about potentially conflicting OPEN share modes and mandatory byte-range locks that might have been in effect at the time of restart, and the metadata server may use this information during the recovery grace period to determine that a WRITE request is safe.

#### 18.9.6. Storage Device Recovery

Recovery from storage device restart dependent upon the layout type in use. Details need to be provided as described in Section 19.1.7.

Despite the above, there are a few general techniques a client can use if it discovers a storage device has crashed while holding modified, uncommitted data that was asynchronously written. First and foremost, it is important to realize that the client is the only one that has the information necessary to recover non-committed data since it holds the modified data and probably nothing else does. Second, the best solution is for the client to err on the side of caution and attempt to rewrite the modified data through another path.

In consequence, the client needs to immediately WRITE the data to the metadata server, with the stable field in the WRITE4args set to FILE\_SYNC4. Once it does this, there is no need to wait for recovery of access to the original storage device.

#### 18.10. Metadata and Storage Device Roles

If the same physical hardware is used to implement both a metadata server and storage device, then the same hardware entity is to be understood to be implementing two distinct roles and it is important that it be clearly understood on behalf of which role the hardware is executing at any given time.

Two sub-cases can be distinguished.

1. The storage device uses NFSv4.1 as the storage protocol, i.e., the same physical hardware is used to implement both a metadata and data server. See Section 20.5 for a description of how multiple roles are handled.

2. The storage device does not use NFSv4.1 as the storage protocol, and the same physical hardware is used to implement both a metadata and storage device. Whether distinct network addresses are used to access the metadata server and storage device is immaterial. This is because it is always clear to the pNFS client and server, from the upper-layer protocol being used (NFSv4.1 or non-NFSv4.1), to which role the request to the common server network address is directed.

#### 18.11. Security Issues for pNFS

pNFS separates file system metadata and data and provides access to both. For security (and other) purposes, data within the NFSv4.1 protocol can be divided as follows:

- \* Non-pNFS-related metadata, typically accessed and modified by using non-pNFS operations directed to the primary server aka the metadata server.
- \* pNFS-related metadata which provides metadata used to access file data, potentially on another device or server.

This information, in the form of layouts, is accessed and modified using special pNFS-related operations directed at the metadata server.

- \* File data, typically accessed using a data access protocol, which might or might not be an NFS protocol, using requests directed at what are called data servers or data storage devices, depending on the layout type.

In other cases, data can be accessed just as it would be on a non-pNFS server, by making READ and WRITE requests directed to the metadata server.

The combination of components in a pNFS system (See Figure 1) is required to preserve the security properties of NFSv4.1 with respect to an entity that is accessing file data from a client, regardless of whether this access is directed to the metadata server or to data elsewhere using layouts provided by the primary server.

Despite this important commonality, the ways in which this is done, depends on the layout type. The layout type affects the data access protocol used and the way that the activities of the metadata server and those providing file data access are coordinated

- \* Security issues for data access protocols not layered on RPC is discussed in Section 18.11.1.

- \* Security issues for data access protocols which are minor versions of NFSv4 and are supported by a separate control protocol are discussed in Section 18.11.2.
- \* Security for data access protocols for versions of NFS without the support of a separate control protocol is discussed in Section 18.11.3.

Given the multiple types of entities whose coordinated effort is required to implement pNFS access, there are a number of inter-entity communications paths which need to be provided with sufficient security to resist attack.

Often, this will require authentication of the requester in order to properly make authorization decisions. When authentication of the principal making the request is not possible, authentication of network peers need to be combined with a trust relationship between the connected peers.

There are a number of possible types of communication paths whose use is possible in various pNFS configurations, depending on the layout type and the specific entities involved.

- \* When an RPC-based communication path is used, the same sorts of techniques described in the NFSv4-wide security document (expected to be derived from [I-D.dnoveck-nfsv4-security]), are adequate to provide the necessary confidentiality and protection against the execution of unauthorized requests. The details may differ depending on the specific protocol used.
- \* In other cases, freedom from unauthorized access can be effected by physical isolation of the communication path between the two entities. However, this physical isolation needs to be supplemented in order to make sure that hostile entities cannot gain access to either of these endpoint.

One common and effective way of blocking such hostile access, is by making sure that the entity is configured so as to not be accessible for general services that can be compromised by external actors. This is often done if the entity is not implemented within a general-purpose operating system or is configured to not to be responsive to general internet traffic.

Where it is not possible to totally block such access, external authentication of principals is necessary.

Given the three types of entities whose coordinated effort is required to implement pNFS access, there are a number of inter-entity communications path which need to be provided with sufficient security to resist attack.

- \* For communication between the client and the metadata server, the RPC-based security described in the NFSv4-wide security document (expected to be derived from [I-D.dnoveck-nfsv4-security]), is to be used as it is when pNFS is not involved. Note that the associated threat analysis is to be found in that same document.
- \* For communication between the client and the data storage devices, there are multiple possibilities to consider, based on the type of file access protocol used.

For data access protocols not using RPC, in general it is not possible to determine whether particular request are appropriately authorized, since there might not be sufficient data present in the request to authenticate the sender or even identify it.

In such cases, the clients and data storage devices require mutual authentication and they need to trust one another.

- \* For communication between the metadata server and stat storage device, there needs to be authentication of both peers and a trust relationship between them.
- \* When there is communication between multiple data storage devices, it is generally best to rely on the metadata servers authentication and its trust of each devices.

#### 18.11.1. Security-related Handling for non-RPC Storage Protocols

These include the blocks layout [RFC5663], the SCSI layout [RFC8154], and the objects layout [RFC5664].

Because these storage protocols do not use RPC, the storage device has no way of determining the specific user making the request. Similarly the storage device has no way of determining the specific open with which a given IO request is associated. As a result, for these storage protocols, the client has the major responsibility for making sure that only valid requests are executed, by implementing the checking of requests to the storage device at the point of issue.

The important role of the client in enforcing these constraints makes authentication of the client peer (e.g. by the use of tls authentication) of critical importance. This is true not only in the case in which AUTH\_SYS is used, but also in the RPCSEC\_GSS case. In

that latter case, the credentials assure that we know the user making the request but we have no knowledge of the client implementation or any reason to view it as trustworthy in enforcing pNFS rules.

Given the reliance on the client, such storage protocols need to have some way of dealing with an unresponsive client when layouts need to be recalled. Typically, this involves some way for the metadata server to contact storage device directly (e.g. by effecting a persistent reservation) and locking out the unresponsive client (and possibly others).

#### 18.11.2. Security-related Handling for RPC Storage Protocols that are NFSv4 Minor Versions Assisted by Control Protocols

These include not only the use of NFSv4.1 as a storage protocol, as described in Section 20 but also the use of all NFSv4 minor versions as data access protocols as described in [RFC8435], in the "tight" coupling mode.

The use of NFSv4 (with RPC) eliminates the difficulties that apply to Section 18.11.1:

- \* Because of the use of RPC, the data server is aware of the user of whose behalf the request is made so the same sort of checking made by the metadata server can be done by the data server. When RPCSEC\_GSS is in use, authentication of that user is unproblematic. On the other hand, when AUTH\_SYS is used to access the data server, authentication of this identity is a serious concern, especially when client host authentication is not available. Generally, the use of AUTH\_SYS without client host authentication is to be avoided when accessing the data server, since you are trusting an unauthenticated client to "authenticate" a user's identity.
- \* Because of the use of NFSv4, the data server is aware of the specific open with which each IO request is associated.

Because, in this sort of environment, the client is not responsible for checking the validity of IO requests, situations in which a layout becomes invalid are dealt with the ordinary recall mechanism used for other recallable locking objects. However, to deal with the possibility of an unresponsive client, the metadata server will typically have the option of contacting the data server directly. Because of the possibility of communication issues, the control protocol will often use a lease-like mechanism so that, in the absence of communication, layouts are cancelled, rather than being kept indefinitely.



[Author Aside]: Need TH review for the flexible-files case.

### 18.11.3. Security-related Handling for RPC Storage Protocols using NFS Versions together with Control Protocol Assistance

These include the use of NFSv3, NFSv4.0, and NFSv4.1 as storage protocols, as described in [RFC8435], in the "loose" coupling mode.

With regard to the difficulties discussed in Section 18.11.1, specifics may vary based on the particular storage protocol used and the possible use of client-host authentication.

- \* Since loose authentication is only described for the AUTH\_SYS case, we have to assume that RPCSEC\_GSS will not be used. As a result, it is RECOMMENDED that client host authentication be available to prevent attackers acting as if they were unauthenticated clients and presenting their requests for execution. The authentication needed to prevent this is the authentication of the clients to the data server. When it is not used, serious security difficulties arise because the clients are not authenticated to the data server which is executing their requests.

The authentication of clients to the metadata server can ameliorate the problem, but the main value of doing so is the encryption of traffic between the client and the metadata server, provided by rpc-tls.

When this problem exists, use of RPCSEC\_GSS by the clients in accessing the metadata server, is not, by itself, helpful. The confidentiality of traffic between the client and the metadata server is necessary, whether that is provided by RPCSEC\_GSS privacy services or by rpc-tls encryption.

- \* In the case of NFSv3 as a storage protocol, there is no way for the data server to determine the open with which each request is assigned. While it might be possible to make this determination in the case in which NFSv4 minor versions are used as storage protocol, it appears that the loose coupling option makes no provision for this either.

As a result, the situation is quite like that described in Section 18.11.1 even though the storage protocol is RPC-based, with clients, rather than data servers responsible for checking request validity. As a result, similar issues arise when clients do not respond properly when layouts are recalled. The metadata server has the ability to make such layouts unusable by changing ownership of the data files involved.

[Author Aside]: TH Needs to review this section.

## 19. Specification of Layout Types

This section describes the required elements of pNFS layout type specifications. These layout types include the following:

- \* The files layout type, specified in Section 20.
- \* The blocks, objects, scsi and flexible files layout types, described in [RFC5663], [RFC5664], [RFC8154], and [RFC8435] respectively.
- \* Future layout types, expected to be specified in standards-track documents including those defining new minor versions and extensions to existing minor versions as provided for in [RFC8178].

What needs to be specified to define such new layout types is discussed in Section 19.1.

### 19.1. Layout Specification Needs

Layout type specifications need to provide:

- \* An understanding of the interoperability model associated with the layout type.

These models are discussed in Section 19.1.1.

- \* Description/definition used to the storage access protocol used to access the storage devices including specifying how information provided with the layout affects the IO operations performed.
- \* The types that need to be known by both the metadata server and client in order to obtain, return, and commit layouts. The relevant requests are generally defined containing a nominally opaque item whose realization by particular layout types is tailored to the specific needs of that layout type.

These types, whose use needs to explained in layout specification documents, are summarized in Section 19.1.2.

- \* Information regarding how pNFS requirements in Sections 19.1.3 through 19.1.5 are met needs to be provided by layout type specifications.

- \* In addition, the information described in Sections 19.1.6 through 19.1.10 needs to be provided.

#### 19.1.1.1. Layout Type Interoperability Models

One component of the required interoperability needs implementation on the metadata server and client, acting as requester and responder, although with a different polarity based on the nature of the requests involved. This applies to all layout types. These consists of:

- \* The implementation described in Section 18.
- \* The functions described Sections 25.40 through 25.44.

All of these need to be supported by the metadata server in the role of responder. It is up to the layout type specification to define the circumstance under which the client needs to use these in the role of requester.

- \* The function described in Section 27.3.

This needs to be supported by the client in the role of responder. It is up to the layout type specification to define the circumstance under which the metadata server is to use these in the role of requester. As part of this it needs to clearly define how the absence of a functioning reverse-direction request path.

Note that many of the requests discussed above contain per-layout-type data types, as discussed in Section 19.1.2. Discussing how these are used needs to be explained by the layout type specification.

In addition to the above, within the NFSv4.1 protocol proper, there is a set of communication requirements involving at least two of three components involved in performing IO using pNFS. As a result there are two interoperability choices for layout types to choose:

- \* Compatibility models requiring the development of a specific control protocol for use between the metadata server and one or more data servers.

Such models typically involve the development of file-system-specific control mechanisms with the metadata server and data server implementations being developed together.

In many cases, these control mechanisms take the form of a bespoke control protocol (in the strict sense). Often such protocols cannot be published as standards since the implementation might involve file-system-specific details that raise serious intellectual property concerns.

- \* Compatibility in which an existing protocol, such as a storage protocol is co-opted as a means to effect the required control mechanisms.

In this case, it is up to the layout type specification to explain how the existing protocol can be used to how the storage protocol can be used to ensure that pNFS requirements in Sections 19.1.3 through 19.1.5 are met.

#### 19.1.2. Layout-Type-Specific Data Types

The XDR definitions of the requests and responses for pNFS operations and callbacks often contain fields best describe as "nominally opaque". This term is appropriate because:

- \* These fields are defined within the protocol as variable-length opaque arrays.
- \* The actual data within these arrays is not opaque but is defined by overlay type whose content is established outside the DR framework.

While this approach differs from the spirit of XDR use, it has been used within NFSv4 for a long time, first being used in these protocol as part of the new NFSv4 attribute model.

The nominally opaque fields whose overlay type need to be provided by the layout type specification include the following:

- \* The `da_addr_body` field of the `device_addr4` data type.
- \* The `loh_body` field of the `layouthint4` data type.
- \* The `loc_body` field of `layout_content4` data type (which in turn is the `lo_content` field of the `layout4` data type)
- \* The `lou_body` field of the `layoutupdate4` data type.
- \* The `lrf_body` field of the `layoutreturn_file4` structure.

### 19.1.3. Control Protocol Layout Management Functions"

The following functions need to be addressed:

- \* Whether revocation of layouts is supported.
- \* When a layout is revoked, the metadata server MUST effect the suppression of future access using the revoked layout and ensure that there are no "lingering" WRITTEs left active after the revocation completes.

Some ways control protocols might provide support involve requests made by the metadata server to the data server to eliminate its record of previously-granted layouts.

- \* Layout conflicts and other situations that might result in layout recall or revocation being needed need to be described.

### 19.1.4. Other Control Protocol Functions"

Whatever control protocol is provided, it needs to support the implementation of the functions listed below. In this list a requirement incumbent upon implementation is presented first, together with examples of how control protocols might provide support to allow the requirements to be met. With regard to the potential normative character of these examples, discussed in Section 19.1.9, these examples are not normative while it the obligation of the specification to provide some possible means of satisfying the need to be addressed in defining the layout type that is normative. Those evaluating the suitability of proposed new layout types are urged to make sure this obligation is satisfied so that the requirements imposed on the implementer can be satisfied.

- \* Because the client MAY still send I/O requests to the metadata server, even though the metadata server has successfully given the client a layout, the metadata server needs to be able to perform IO requests directed at the metadata server, even though the data is present on a data storage device such as a data server. In doing so, it MUST perform the same operation as would be required by the NFSv4.1 protocol, if the operation were directed to an NFSv4 server when the pNFS feature is not involved.

Some ways control protocols might provide support involve ways for the metadata server to request execution of the IO request by the data server on behalf of the client making the original request so that the results can be returned to the client.

In some configurations, no inter-server request is necessary, since the data storage devices might be directly accessed by the metadata server, in the same way they would be accessed by the data access protocol.

- \* IOs sent to data storage devices must be considered authorized only if the authorization checking is the same as would be done when they were sent directly to the metadata server or a non-PNFS data server. When this authorization checking finds the request not authorized, it MUST NOT be performed.

Generally and as described in [I-D.dnoveck-nfsv4-security], explicit authorization checking at IO time is not required, since most IO operations are performed by the principal that opened the file. However, in the case in which other principals are performing IO, the layout type specification needs to indicate how the necessary checking is done.

This can involve use of the control protocol, to do the checking on the metadata server, or depending on the client to do the necessary checking. The layout type specification needs to make clear how this checking is done so that the result is the same as would have resulted for IOs directed at metadata server. In this regard, the uncertainty about the authorization semantics associated with the NFSv4 ACL model needs to be accommodated.

- \* The values reported by the metadata server as to the value of attributes such as modify time, the change attribute, and the EOF position (i.e. the size attribute) MUST reflect the effect of IO operations directed to data storage devices, except where the discrepancies indicated below are tolerated for a limited time:
  - While the size of a file can be affected by WRITE requests directed at the metadata server and by WRITE operations directed to data storage devices, it is not always clear whether or how any particular operation affects the file size, when files are divided into multiple pieces (using striping or other mechanisms) resident on multiple data storage devices.

In such cases the metadata server may not always have the definitive file size value but it MUST return the correct value when the attribute is interrogated by GETATTR/(N)VERIFY after a LAYOUTCOMMIT is done.

It is the obligation of the layout type specification to explain how that correct value is to be obtained. This can involve the knowledge that the metadata server might have as to the location of the server holding the last component, and

details of the control program whereby the data server keeps the metadata server aware if extension to the size of the last component that might need to be propagated to the metadata server in the event of a LAYOUTCOMMIT.

- Attributes such as `modified_time` and `change` are affected by WRITES performed at the data storage device. Such changes need to be reflected at the metadata server as soon as execution of LAYOUTCOMMIT for the layout associated with the WRITES. Situation in which these are propagated without an explicit LAYOUTCOMMIT need to be documented in the layout type specification document.
- A LAYOUTCOMMIT requires that changes in attributes resulting from operations on the storage device using the layout committed need to be reflected in the attributes available on the metadata server before the LAYOUTCOMMIT is responded to.

It is the obligation of the layout type specification to make it clear how this requirement is to be adhered to.

- \* Operations directed at the metadata server often need to involve storage allocation, whether the operation involves creating, deleting, extending, or truncating files.

The storage allocation MUST be the same as would have been done if the corresponding operation was performed on an NFSv4.1 server for which pNFS was not active (or active but not used).

The layout type specification needs to explain how this is to be done. The following approaches are common for existing layout types:

- The storage allocation is done by the file system component, whether that resides with the metadata server or another data server.

In such cases, the client is unaware of the details of storage allocation since it sees only the file abstraction, which is designed to hide these details.

- The storage allocation is done by the metadata server with the client informed of the detail so that it can use the storage assigned as the source or destination for IO requests

- The storage allocation is done by a distinct file system component and the client needs to be aware of the assigned location as seen by the metadata server and the data storage device despite the fact that the latter are different from the physical addresses used by the data server.

#### 19.1.5. IO Checking Requirements

The requirements listed below were introduced by Section 18.7.1 and do not always require control protocol support. Nevertheless, they are discussed here because corresponding requirements for the data storage device might exist for some layout types and these would require control protocol support.

The client has these obligations when making I/O requests to the storage devices:

- \* Clients MUST NOT perform I/O to the storage device if they do not have layouts for the files in question.
- \* Clients MUST NOT perform I/O operations outside of the specified ranges in the layout segment.
- \* Clients MUST NOT perform I/O operations that would be inconsistent with the iomode specified in the layout segments it holds.
- \* Clients MUST NOT perform I/O operations that would not be authorized if performed by the metadata server.

It is important to understand that the need to make these checks might make it inappropriate to discard layout information as discussed in Section 18.7.3.1

The ability of the data storage device to verify that these requirements are being adhered to varies with the layout type, as does the potentially normative requirements for the data server to do this verification and for the layout type specification to describe the means by which the control protocol provides the data storage device the means information necessary to effect such verification.

It is not inherently necessary for a layout type specification to address the question of storage device normative verification requirement. Nevertheless, if there is no such requirement, the security consequences need to be addressed. Furthermore, if this requirement exists, it is incumbent on the layout type to explain how the data storage device can obtain the information necessary to do the required verification.



#### 19.1.6. Security-related Requirements"

A security considerations section needs to be provided by each layout type specification,. This section needs to explain how the NFSv4.1 authentication and authorization functions are preserved. That is, if a metadata server would prevent processing of a READ or WRITE operation, how would a corresponding IO issued using the layout type being described be prevented.

#### 19.1.7. Recovery Requirements"

The specification needs to describe the methods of recovery from the following situations:

- \* Failure and restart for client, metadata server, storage device.
- \* Lease expiration from perspective of the active client, metadata server, and possibly the storage device when applicable.
- \* Loss of layout state resulting in fencing of client access to storage devices (for an example, see Section 18.9.3).

#### 19.1.8. Requirements Regarding Committing Layouts

While the XDR provided by Section 25.42 provides extensive flexibility as to how this function is to be provided, it is not clear how, for particular layout types, flexibility is to be appropriately used.

As a result, layout type specifications need to provide the following information:

- \* Whether the size value presented is always used if it extends the file or whether there are circumstances in which the metadata server is not to do so.

This information is necessary because the statement "may use this information" makes it extremely unclear how a global view of the file size might be maintained if it is not so used.

- \* Whether, and under what circumstances, the metadata server might use a modified time suggestion provided with LAYOUTCOMMIT.

When this is allowed, there needs to be an explanation about how lack of time synchronization is to be dealt with and how it deals with necessary checking of suggested values.

Also needed is a discussion of how acceptance of such suggestions affects the change attribute.

- \* While the XDR defining the format of the `lou_body` field is made available as part of defining overlays for nominally opaque types, explanation of the semantics and uses of this information needs to be provided to make it clear how LAYOUTCOMMIT is to be used.
- \* The circumstances under which LAYOUTCOMMIT is necessary. Given the general assumption that this is necessary iff WRITE operations are done using layouts, this can be addressed by combining specification of other situations in which LAYOUTCOMMIT is required and discussion of situations which WRITES are done and LAYOUTCOMMIT is NOT required.

#### 19.1.9. Requirements Regarding Layout Termination

Layout termination involves the return, recall, and revocation of layouts, and participants' discarding of layout information.

As a result, layout type specifications need to provide the following information:

- \* In the case of layout return, there is not expected to be any restrictions on clients' ability to return layouts, other than requiring the client to wait for pending IO's to complete. In the event such restrictions exist, the layout type specification needs to describe them.

For layout types in which IO validation is done on data storage devices, requirements regarding making the device aware of the termination of the layout (and consequent refusal of further IO and draining of current ones) need to be clearly stated.

- \* The case of layout recall is similar in that restrictions, if they exist, need to be described by the layout type specification.

In cases in which recall needs to be done only for some layout types (e.g. REMOVE), the requirements need to be clearly stated.

- \* While layout revocation is always possible due to an unacceptably delayed layout recall or a failure-like situation such as an expired lease, some layout types might allow it in other situations. These need to be described clearly.

Where these situations are described using the term "conflicting layouts", the meaning of this term needs to be explained clearly. This is particularly important because [RFC8881] defined this in a

way that does not apply to the files layout when server-multipathing is in effect. Also, it need to be made clear whether non-layout-based IO operations can have the same effect.

Descriptions of how these situations are to be dealt with need to describe how IO operations using the revoked layouts are prevented and existing ones drained. In this context, where referring to "fencing" is stated in a requirement, it needs to be supplemented by an explanation of how the requirement is dealt with in all situations. For example, informing a data server of a layout's termination needs to be supplemented by consideration of the situation in which the data server cannot be contacted.

- \* As discussed in Section 18.7.3.1, clients and metadata servers may discard layout information without worry about the inconsistency between the MDA and clients thereby created. However, there can be restrictions on such discarding based on the layout type. These need to be clearly described.

#### 19.1.10. Requirements Regarding Feature Interactions"

The layout type specifications need to provide information about OPTIONAL features that are not supported (or for which special handling for support is required) for particular layout types.

These features can include those defined after publication of the layout type specification. While it would be appropriate for the specification of relevant gaps to be the responsibility of the extension defining the feature, it is often the case that these issues were not recognized for some new features.

This situation can result in restrictions that are hard to explain when the features involved are part of Minor Version Two. Mention of such features is not fully appropriate in this document, but the prospect for bis documents for the optional feature or the existing layout types seems too far off to properly advise potential implementors. For this reason, we will mention such difficulties in order to provide helpful implementation information in the layout type specification or in subsections of Section 19.

#### 19.2. Addressing Requirements for Existing Layout Types

Each of the existing layout types is expected to address the requirements discussed in previous subsections of Section 19. However, because some of the definitions of these layout types were published before these requirements were laid out in detail, it might be that the existing definition might not make it sufficiently clear how these requirements were addressed.

In addition, it might be important to note potential difficulties with OPTIONAL added in a late minor version.

Descriptions of how existing layout types address these requirements are to be found as described below:

- \* The definitions of the blocks, objects, and scsi layout types, were published before [RFC8434] was published and the need for the layout type requirements was fully recognized. For a discussion of how the layout type requirements are addressed, see Sections 19.2.1, 19.2.3 and, 19.2.2 respectively.
- \* The definition of the flexible files layout type was published after [RFC8434] but before some the necessity for some of the clarifications provided in this document were recognized. For a discussion of how the layout type requirements were addressed, see Section 19.2.4,
- \* Although the original definition of the files layout type was published before [RFC8434], the revisions made as part of the NFSv4.1 respecification effort can reasonably be expected to make it clear how the layout type requirements were addressed. In any case, for a discussion how those revision might affect existing implementation that relied on earlier specifications, see Section 20.26 for details.

#### 19.2.1. Blocks Layout Type and Layout Type Requirements

This layout type was described in [RFC5663], before the obligations of layout type specifications first discussed in [RFC8434] were understood. While it is has proved possible for these gaps to be dealt with by implementers involved in the pNFS effort, the need for wider understanding of the pNFS feature at this stage of protocol development makes it necessary to fill in these gaps in this section.

One important class of gaps concerns the need to deal with matters that need to be explained, as provided for in Section 19.1.4, Although, the answers could be considered obvious, it might not be so to all readers and it would be helpful if the explanations below were taken into account when reading [RFC5663]

- \* It is not made sufficiently clear how IO request authorization is provided for.

It needs to be clearly stated that this layout type's approach to authorization checking relies on the checking being done at OPEN time, and that cases in which this checking is not adequate are problematic.

Of particular concern are IO requests to be performed by a principal different from the one that the opened file. In order to provide the same authorization results that would be provided for an IO request sent to the MDS would require the client to use ACCESS on each such IO but that needs to be stated clearly if that is to be done.

One likely approach to this issue and some discussed below (e.g. mandatory byte-range locks) could be addressed as is done in Section 3 of [RFC8154].

- \* The role of the control protocol with respect to storage allocation needs to be clarified.

While [RFC8434] suggests that this is somehow a control protocol function and [RFC5663] is silent on the matter, it needs to be made clear that, for this layout type, the control protocol is not involved in storage allocation and that allocation are performed by the file system as they are when pNFS is not involved. The only exception, which is mentioned in [RFC5663], concerns allocation which occurs as part of granting writable layouts and their possible role in implementing copy-on-write.

Another important class of gaps concern the need for discussion of OPTIONAL features which cannot be supported or require special effort to support when using this layout type.

- \* Mandatory byte-range locks are not supportable together with layouts of the layout type.

Servers which do support such locking need to avoid granting layouts of this type while such locks are active.

Granting of layout of this type cannot be done when such locks are in effect for the target file.

- \* When a server supports NFSv4 ACLs and the server supports use of distinct permissions for extending file and for overwriting existing bytes, proper authorization cannot be guaranteed, making it impossible to grant layout on file with ACLs that sometime allow one of these actions and not the other.

#### 19.2.2. Scsi Layout Type and Layout Type Requirements

This layout type was described in [RFC8154], before the obligations of layout type specifications first discussed in [RFC8434] were fully described.

Nevertheless, these issues appear have been better understood when this document was written since this document does address many of the issues noted in Section 19.2.1 in its Section 3 which clarifies authorization and locking semantics conformance. While it appears that, given these, improvements, it might be better if [RFC8154] obsoleted [RFC5663], as things currently stand, [RFC5663], has been neither updated or obsoleted by [RFC8154], For a discussion of how this issues is to be addressed, see Appendix D.2.15.

Despite this large improvement, there remain a number of difficulties with [RFC8154] lack of attention to the specification's requirements of [RFC8434] and Section 19, including the following.

- \* The issue with regard to the clarification of storage allocation noted in Section 19.2.1 applies equally to this layout type.
- \* The issue with regard to handling of NFSv4 ACLs noted in Section 19.2.1 applies equally to this layout type.

#### 19.2.3. Object Layout Type and Layout Type Requirements

This layout type was described in [RFC5664], before the obligations of layout type specifications first discussed in [RFC8434] were understood. As a result it will be necessary to fill in these gaps in this section.

The following issues, noted for other similar layout types, need to be noted:

- \* This document, like [RFC5663] does not adequately address many of the issues noted in Section 19.2.1 regarding authorization and locking semantics.

As in the case of the blocks layout type, the material in Section 3 of [RFC8154] is helpful.

- \* The discussion of allocation while it is not explicit as it might be, is adequate since it clear that, in this case storage allocation is a control protocol function, since the control protocol and the storage protocol are the same.
- \* The issue with regard to handling of NFSv4 ACLs noted in Section 19.2.1 applies equally to this layout type and needs to be addressed similarly to this document.

#### 19.2.4. Flexible Files Layout Type and Layout Type Requirements

Issues relating to the suitability of this for description of a file layouts with tight coupling are discussed in Section 19.2.4.1.

The following issues, including a number noted for other layout types, need to be noted:

- \* Issues related to locking semantics appear to have been adequately dealt with in Section 2.3 of [RFC8434].
- \* Issues with regard to authorization, including dealing with IOs issued by non-opening principals and other matters relating to the use of NFSv4 ACLs do not appear to be addressed adequately. This appears to be due primarily to the fact that the working group was constrained, in dealing with the authorization-related issues dealt with in [I-D.dnoveck-nfsv4-security] and [I-D.ietf-nfsv4-acls-update] to allow a wide range of previously allowed behavior, invalidating apparently reasonable assumptions made by the authors of [RFC8435].

Because the existing text relating to authorization is different for the tight and loose cases, this matter will be discussed in detail in Section 19.2.4.1.

##### 19.2.4.1. Tight Coupling Using the Flex File Layout Type

As this document is currently written, it attempts to describe support for multiple storage protocols and multiple approaches to the handling of locking and authorization. While the description of the tight coupling model is helpful, it cannot, for reasons explained below, serve simultaneously as the layout type specification for what are essentially two different ways of providing access to data storage devices.

To illustrate the problem, let us look at the statement "The requirements for a control protocol are specified in [RFC5661] and clarified in [RFC8434]" appearing in the document Introduction. Part of the clarification provided in [RFC8434] and Section 19 includes the requirement that certain matters need to be addressed in a layout-type-specific fashion with the responsibility for explaining the particulars of many matters left to the layout type specification for that layout type. This creates the following difficulties for the approach currently taken in [RFC8435]

- \* The differences between tight and loose coupling are so great that they cannot be addressed in this framework using the same layout type for both, as many of the layout-type-specific matters are likely to be different for these two approaches.

When the description of these two models are interspersed, as they are now, it is difficult to deal simultaneously with situations in which request verification is the responsibility of two different parties and fencing has such wide differences in granularity and disruptiveness.

A related practical difficulty concerns the inability of clients to know before, receiving a layout, the type of coupling and the nature of the storage protocol. Given that the client can only specify the layout type, it has no way of making sure before requesting a layout, that it is one that it can use or of helping the MDS to provide one it can use.

- \* Layout-type-specific matters are addressed in Section 20, where layout-type-specific descriptions are provided for the file layout type, since this section is layout specification for the files layout type. There appears to be no specification of the tight flexible file layout type even if a value were assigned for this purpose.

If it is intended that these all be dealt with identically, there needs to be document stating this clearly. As it is, we have a set of remarks, not always correct, about how these issues might be addressed in various circumstances.

Of particular importance is the very different approaches, to request validation and fencing taken in the files layout type and the loose-coupling case. As a result, it needs to be clearly stated what approach is taken in the case of tight-coupling of flex-files layouts.

- \* The handling of request authorization appear to be different for the two types of coupling with both different from the handling specified in Section 20. All of the cases discussed below need to deal with the cases that have proved difficult to deal with for other layout types, including IOs issued by principals other than the opener, potential finer-grained write permissions for NFSv4 ACLs, and the existing underspecification of NFSv4 ACLs.
- \* Authorization for loose coupling appears to be described by the third paragraph of the Security Considerations section of [RFC8434] which is quoted (with interspersed comments) below:



The metadata server enforces the file access control policy at LAYOUTGET time. The client should use RPC authorization credentials for getting the layout for the requested iomode ((LAYOUTIOMODE4\_READ or LAYOUTIOMODE4\_RW), and the server verifies the permissions and ACL for these credentials, possibly returning NFS4ERR\_ACCESS if the client is not allowed the requested iomode.

Although following the above probably harmless, it considerably confuses the matter since layouts, like delegations, are of client scope and can be used by multiple principals doing IO operations on multiple OPENS of the same file.

If the LAYOUTGET operation succeeds, the client receives, as part of the layout, a set of credentials allowing it I/O access to the specified data files corresponding to the requested iomode.

This is the essence of the loose file authentication approach.

When the client acts on I/O operations on behalf of its local users, it MUST authenticate and authorize the user by issuing respective OPEN and ACCESS calls to the metadata server, similar to having NFSv4 data delegations.

Needs further clarification although it is correct except for the confusing mention of authentication by ACCESS and OPEN. Suggest something like the following:

When the client issues IO requests on behalf of local users, it use the credential provided by LAYOUTGET. However, it need the authorization provided by a previous OPEN to ensure that the opening principal can validly perform the operation. In cases in which IO operations are performed by principals other than the opener, ACCESS MUST be used to verify the authorization to perform the IO request.

- \* The authorization described for the tight coupling case seems to assume that if the MDS copies the authorization-related attributes from the MDS to the data server, the authorization will be the same. While it would be nice if this were the case, the underspecification of ACL authorization semantics and the liberal approach to existing variances from POSIX authorization makes it hard to rely on. This is probably not a practical problem, if, as seems to be the case, the discussion of tight coupling is basically exploratory, making it likely that it can be changed before being used in a standards-track environment

- \* For authorization handling in the files layout type, the material in Section 20.14.2 provides helpful suggestions.

It should be consulted as possibilities for addressing authorization issues in the context of tight coupling are considered.

#### 19.2.5. Needed Updates to Layout Specification Documents

As described in Sections 19.2.2 and 19.2.3, there need to be changes made to [RFC8154] and [RFC5663]

- \* [RFC5664] needs to incorporate the material in Section 3 of [RFC8154] and the update regarding ACL handling currently presented in Section 19.2.1.
- \* [RFC8154] needs to the update regarding ACL handling currently presented in Section 19.2.1.

As a result, this draft and subsequent ones will be marked as updating [RFC8154] and [RFC5664].

There need to be updates to [RFC5664] and [RFC8435]. However, because of the issues noted below, these will only affect the status of these documents in later drafts.

- \* Because of uncertainty about the continuing need for a non-SCSI block layout type and how that need might be best provided for, the working group needs to discuss the issues in Appendix D.2.15 and decide on a way forward involving obsoleting or updating [RFC5663].
- \* Because of the interaction between authorization issues and the question of how and when to address questions related to tight coupling, it is would be best if the working group addresses the need for changes following the recommendations of the authors of [RFC8435].

The working group, in deciding on a way to address issues in this document, will need to accommodate the authors' plans for future development of this layout type.

## 20. The pNFS File Layout Type

This section describes the structure and semantics of the file-based layout type or pNFS. These file-based layouts use the layout type LAYOUT4\_NFSV4\_1\_FILES and use NFSv4.1 as a Data Access Protocol

The information required of all layout type specifications by Section 19 are provided, for the file layout type, in this section and included subsections.

There a number of noteworthy characteristics of this layout type that affect many of the specific characteristics that are to be provided by layout type specification documents. How these issues are dealt with in the case of the files layout type are effected by the following salient design characteristics:

- \* Because the data storage protocol is identical to that used when pNFS is not present, there is no potential performance benefit, as there might be with other layout types, from the use of simpler, more direct means of performing IO operations.

However, there is a corresponding benefit from directing IO requests to the servers able to support them directly. This IO direction function can improve performance by eliminating the overhead associated with request forwarding with clustered servers,

- \* Because the data storage device is a data server containing, at the least, facilities to persistently record, the physical location and extent of written data, there is no need for this information to be part of the interaction with the data storage device, allowing them to use a file-like access paradigm.

As a result, facilities such as the use of LAYOUTCOMMIT can be dispensed with in many cases and will never require client estimates of file attributes to be provided to the metadata server.

- \* Because data server devices are expected to have significant intracluster communication facilities that can be used as part of a control protocol, information can be passed between the metadata sever and data storage devices, allowing IO checking to be performed by the data storage device.

This allows extensive changes to the client to be avoided, thereby eliminating dependence on IO checking being done by the client with its negative effect on security.

### 20.1. Recent Changes to the File Layout Type

In addition to the re-organization made necessary by the incorporation of the new Section 19, there are a number of changes made to fill in gaps left in previous descriptions of this layout type and correct misunderstandings where previous descriptions of pNFS did not pay adequate attention to important aspects of this layout type and the need to deal with the possibility of server-multipathing (See Section 20.9).

The following issues should be noted:

- \* Although there were previously references to the need for layout specifications to specify whether those layout types had the ability to revoke layouts or to do so "unilaterally", there was no explicit statement either accepting or forbidding this for the files layout type.

Given this lack, it is reasonable to suppose this is allowed since the files layout type, unlike others, has facilities for fine-grained non-disruptive fencing.

We have adopted the approach of allowing this because there is no basis to deny it and because of the need to correctly deal with WRITES in the server-multipathing case.

- \* There are some statements regarding the possibility of layout conflict which seem to be based on experience with and the needs of block layouts and only apply to file layouts if the possibility of server-multipathing is not considered.

This has no practical effect, since the MDS is always free to recall layouts. Nevertheless, to make things clearer, we needed to explain such conflicts and how they can arise when server multipathing is possible.

In order to address this issue we limited the core of recognition of conflicts to cases of range overlaps and made the rest the responsibility of the layout type (for the file layout type, this is discussed in Section 20.18.)

- \* There are some confusing statements allowing clients to not check IOs for proper values of layout iomode before issuing WRITES.

It appears that, for the file layout type, which relies on the data server to do IO checking, iomode is checked for WRITES as it should be.

Making the check obligatory or using "SHOULD" raises compatibility issues so we are not changing this, since it workable where server-multipathing is not an issue.

As modified, we state that checking the iomode is not a general protocol need but that there are situations (e.g. server-multipathing) that make it advisable.

Section 13.5 of [RFC5661] and [RFC8881] contains the following statement which appears to be an early and ultimately not very helpful discussion of what we have come to call server-multipathing.

It is not always necessary for the two data server addresses to designate the same server with trunking being used. For example, the data could be read-only, and the data consist of exact replicas.

As it wov, about sixteen years later appears necessary to deal, within in the protocol, transitions to and from situations in which two areas of a file are read-only replicas, the revised discussion of the files layout type in Section 20. In so doing, the changes already discussed to revocation, conflicts and the handling of layout iomode are pretty much essential, even though they would be justifiable on other grounds.

## 20.2. File Layout Type Motivation

The compatibility model (see Section 19.1.1 is one which requires the development of file-system-specific control mechanisms to meet the requirements laid out in Section 19 and the performance goals of the layout type, as described below:

- \* To provide access to the greater IO throughput made available by a server cluster, without requiring the user to be aware of detailed file location information,
- \* To avoid unnecessary request forwarding by directing IO requests to servers prepared to execute them directly. This IO direction function often makes it preferable to request IO operations using file layout rater than using the metadata server as a point of access.
- \* To provide the ability to provide greater throughput while accessing a single file by including facilities for striping data across multiple NFSv4.1 data servers.

### 20.3. Protocol Feature Support for the Files Layout Type

In general, the addition of post-NFSv4.1 OPTIONAL features are unlikely to cause difficulties for existing implementations of the files layout type, although some level of control protocol extension is likely to be necessary. For example, features that involve storage allocation changes such those providing hole-punching can generally be accommodated using facilities for remote request execution on the associated data server(s).

There is one important case in which support for server-multipathing (See Section 20.9) was not present for previous implementations of this layout type and the lack of recognition of this as an important case in [RFC8881] could lead to lack of support for this case, if, as suggested in previous specification, conformance to iomode was ignored. For this reason, those who need support for server multipathing should be aware of the importance of checking for iomode validity in order to provide support for server multipathing.

### 20.4. Storage Protocol for the Files Layout Type

The storage protocol used to perform IO to the storage devices, which are data servers in this case, is identical in form to that used to perform IO to the metadata server, or to an NFSv4.1 server operating outside the PNFS framework.

Despite this general arrangement, the following discussion of instances of different treatments and for adherence to uniform treatment are to be adhered to:

- \* The checking of authorization done on the data server MUST be identical to that which would be performed if the IO request were directed to the metadata server.

While, for the most part, the server performing the IO only needs to explicitly check authorization if the principal performing the IO is different from the opener, there can be occasions in which to provide the identical handling required in previous paragraph, by adapting to the consequences of some unfortunate gaps in the specification of NFSv4.1 authorization semantics, the data server might need to adhere to what is now non-standard authorization checking, matching that of the metadata server.

Because previous specifications (i.e. [RFC3530], [RFC7530], [RFC5661], [RFC8881]) did not discuss authorization semantics, it was necessary to fill this gap in [I-D.dnoveck-nfsv4-security] but, in so doing, previously acceptable behavior could not be invalidated. In addition, because of ACL underspecification in

previous documents, [I-D.ietf-nfsv4-acls-update] had to be written with similar difficulties because of existence of current implementations with a wide range of semantics implemented.

As a result. because of the possibility of various undocumented authorization schemes existing, it might not be possible to appropriately document the necessary control protocol activity in Sections 20.14.2 and 20.15 while clearly explaining the most common case, following POSIX, with limited modification to support that part of NFSv4 ACLs that fits within the POSIX authorization model. As a result, implementors need to focus on meeting the requirement regarding identical authorization and are responsible for effecting it in the control protocol provided for coordination among servers

- \* The possibility, within NFSv4 ACLs, of separate ACE mask bits separately controlling authorization for extending and for overwriting existing file data, might make it impossible to fully ascertain write authorization at open time and limit checking to the case in which principals other than the opener issue IO.
- \* When stateids are sent to the data server, they are obtained from the metadata server rather than from the same server used to do the IO, as is done when pNFS is not involved.

Despite the above, there are restrictions described in Section 20.10 that limit the stateids that can be used. Only stateids with a zero seqid are to be used in performing IO operations using the file layout type and special stateids cannot be used.

- \* Since stateids are interpreted based on the associated clientid, it is necessary that IO operations be done as part of sessions that are associated with the same clientid associated with the clientid used to obtain the stateid.

For a discussion of how such a clientid is to be obtained, see Section 20.5.

This situation, in which multiple sessions share a clientid, is similar to that that in effect when clientid trunking is in effect. However, implementation is simpler since in the pNFS files case since only one session can cause changes to the locking state, unless clientid trunking is supported as well.

- \* The file handles used to designate files on which IO is to be done are likely to be different since the handle established as part of the layout is to be used.

- \* Although the principals are often the same and need to be compared, their existence in different forms (e.g. name@domain vs. numeric uids) will require reliable mapping between multiple forms as is often necessary at the protocol-filesystem interface.

## 20.5. Client ID and Session Considerations

Sessions are a REQUIRED feature of NFSv4.1, and this requirement for use applies to both the metadata server and file-based (NFSv4.1-based) data servers.

The role a server plays in pNFS is determined by the result it returns from EXCHANGE\_ID. The roles are:

- \* Metadata server (EXCHGID4\_FLAG\_USE\_PNFS\_MDS is set in the result `eir_flags`).
- \* Data server (EXCHGID4\_FLAG\_USE\_PNFS\_DS).
- \* Non-metadata server (EXCHGID4\_FLAG\_USE\_NON\_PNFS). This is an NFSv4.1 server that does not support operations (e.g., LAYOUTGET) or attributes that pertain to pNFS.

The client MAY request zero or more of EXCHGID4\_FLAG\_USE\_NON\_PNFS, EXCHGID4\_FLAG\_USE\_PNFS\_DS, or EXCHGID4\_FLAG\_USE\_PNFS\_MDS, even though some combinations (e.g., EXCHGID4\_FLAG\_USE\_NON\_PNFS | EXCHGID4\_FLAG\_USE\_PNFS\_MDS) are contradictory. However, the server MUST only return the following acceptable combinations:

+=====+	
Acceptable Results from EXCHANGE_ID	
+=====+	
EXCHGID4_FLAG_USE_PNFS_MDS	
+-----+	
EXCHGID4_FLAG_USE_PNFS_MDS   EXCHGID4_FLAG_USE_PNFS_DS	
+-----+	
EXCHGID4_FLAG_USE_PNFS_DS	
+-----+	
EXCHGID4_FLAG_USE_NON_PNFS	
+-----+	
EXCHGID4_FLAG_USE_PNFS_DS   EXCHGID4_FLAG_USE_NON_PNFS	
+-----+	

Table 7

As the above table implies, a server can have one or two roles. A server can be both a metadata server and a data server, or it can be both a data server and non-metadata server. In addition to returning



two roles in the EXCHANGE\_ID's results, and thus serving both roles via a common client ID, a server can serve two roles by returning a unique client ID and server owner for each role in each of two EXCHANGE\_ID results, with each result indicating each role.

In the case of a server with concurrent pNFS roles that are served by a common client ID, if the EXCHANGE\_ID request from the client has zero or a combination of the bits set in eia\_flags, the server result should set bits that represent the higher of the acceptable combination of the server roles, with a preference to match the roles requested by the client. Thus, if a client request has (EXCHGID4\_FLAG\_USE\_NON\_PNFS | EXCHGID4\_FLAG\_USE\_PNFS\_MDS | EXCHGID4\_FLAG\_USE\_PNFS\_DS) flags set, and the server is both a metadata server and a data server, serving both the roles by a common client ID, the server MUST return with (EXCHGID4\_FLAG\_USE\_PNFS\_MDS | EXCHGID4\_FLAG\_USE\_PNFS\_DS) set.

In the case of a server that has multiple concurrent pNFS roles, each role served by a unique client ID, if the client specifies zero or a combination of roles in the request, the server is to return only one of the roles from the combination specified by the client request. If the role specified by the server result does not match the intended use by the client, the client should send the EXCHANGE\_ID specifying just the interested pNFS role.

If a pNFS metadata client gets a layout that refers it to an NFSv4.1 data server, it needs a client ID on that data server. If it does not yet have a client ID from the server that had the EXCHGID4\_FLAG\_USE\_PNFS\_DS flag set in the EXCHANGE\_ID results, then the client needs to send an EXCHANGE\_ID to the data server, using the same co\_ownerid as it sent to the metadata server, with the EXCHGID4\_FLAG\_USE\_PNFS\_DS flag set in the arguments. If the server's EXCHANGE\_ID results have EXCHGID4\_FLAG\_USE\_PNFS\_DS set, then the client may use the client ID to create sessions that will exchange pNFS data operations. The client ID returned by the data server has no relationship with the client ID returned by a metadata server unless the client IDs are equal, and the server owners and server scopes of the data server and metadata server are equal.

In NFSv4.1, the session ID in the SEQUENCE operation implies the client ID, which in turn might be used by the server to map the stateid to the right client/server pair. However, when a data server is presented with a READ or WRITE operation with a stateid, because the stateid is associated with a client ID on a metadata server, and because the session ID in the preceding SEQUENCE operation is tied to the client ID of the data server, the data server has no obvious way to determine the metadata server from the COMPOUND procedure, and thus has no way to validate the stateid. One RECOMMENDED approach is for pNFS servers to encode metadata server routing and/or identity information in the data server filehandles as returned in the layout.

If metadata server routing and/or identity information is encoded in data server filehandles, when the metadata server identity or location changes, the data server filehandles it gave out will become invalid (stale), and so the metadata server MUST first recall the layouts. Invalidating a data server filehandle does not render the NFS client's data cache invalid. The client's cache should map a data server filehandle to a metadata server filehandle, and a metadata server filehandle to cached data.

If a server is both a metadata server and a data server, the server might need to distinguish operations on files that are directed to the metadata server from those that are directed to the data server. It is RECOMMENDED that the values of the filehandles returned by the LAYOUTGET operation be different than the value of the filehandle returned by the OPEN of the same file.

Another scenario is for the metadata server and the storage device to be distinct from one client's point of view, and the roles reversed from another client's point of view. For example, in the cluster file system model, a metadata server to one client might be a data server to another client. If NFSv4.1 is being used as the storage protocol, then pNFS servers need to encode the values of filehandles according to their specific roles.

#### 20.5.1. Sessions Considerations for Data Servers

Section 7.11.2 states that a client has to keep its lease renewed in order to prevent a session from being deleted by the server. If the reply to EXCHANGE\_ID has just the EXCHGID4\_FLAG\_USE\_PNFS\_DS role set, then (as noted in Section 20.10) the client will not be able to determine the data server's lease\_time attribute because GETATTR will not be permitted. Instead, the rule is that any time a client receives a layout referring it to a data server that returns just the EXCHGID4\_FLAG\_USE\_PNFS\_DS role, the client can validly assume that the lease\_time attribute from the metadata server that returned the layout applies to the data server. Thus, the data server MUST be

aware of the values of all `lease_time` attributes of all metadata servers for which it is providing I/O, and it MUST use the maximum of all such `lease_time` values as the lease interval for all client IDs and sessions established on it.

For example, if one metadata server has a `lease_time` attribute of 20 seconds, and a second metadata server has a `lease_time` attribute of 10 seconds, then if both servers return layouts that refer to an `EXCHGID4_FLAG_USE_PNFS_DS`-only data server, the data server MUST renew a client's lease if the interval between two `SEQUENCE` operations on different `COMPOUND` requests is less than 20 seconds.

## 20.6. File Layout Definitions

The following definitions apply to the `LAYOUT4_NFSV4_1_FILES` layout type and might be applicable to other layout types.

**Unit.** A unit is a fixed-size quantity of data written to a data server.

**Pattern.** A pattern is a method of distributing one or more equal sized units across a set of data servers. A pattern is iterated one or more times.

**Stripe.** A stripe is a set of data distributed across a set of data servers in a pattern before that pattern repeats.

**Stripe Count.** A stripe count is the number of units in a pattern.

**Stripe Width.** A stripe width is the size of a stripe in bytes. The stripe width = the stripe count \* the size of the stripe unit.

Hereafter, this document will refer to a unit that is written in a pattern as a "stripe unit".

A pattern may have more stripe units than data servers. If so, some data servers will have more than one stripe unit per stripe. A data server that has multiple stripe units per stripe MAY store each unit in a different data file (and depending on the implementation, will possibly assign a unique data filehandle to each data file).

## 20.7. File Layout Data Types

This section defines the specific types overlaying the nominally opaque types listed in Section 19.1.2

The high level NFSv4.1 layout types are:

- \* `nfsv4_1_file_layouthint4` overlaying the `loh_body` field of the `layouthint4` data type. This type is described in detail in Section 20.7.1.
- \* `nfsv4_1_file_layout_ds_addr4` overlaying the `da_addr_body` field of the `deice_sddr4` data type. This type is described in detail in Section 20.7.2.
- \* `nfsv4_1_file_layout4` overlaying the `loc_body` field of the `layout_content4` data type. This type is described in detail in Section 20.7.2.
- \* Within the file layout type, there is no layout-type-specific overlay for the `lou_body` field of the `layouytupdate4` data type or for the `lrf_body` field of the `layoutreturn_file4` structure.

As a result, this field can have any value with a zero-length opaque string being most convenient.

#### 20.7.1. File Layout Hint-related Data Types

The SETATTR operation supports a layout hint attribute (Section 11.16.4). When the client sets a layout hint (data type `layouthint4`) with a layout type of `LAYOUT4_NFSV4_1_FILES` (the `loh_type` field), the opaque `loh_body` field is overlaid by a value of data type `nfsv4_1_file_layouthint4`.

```
const NFL4_UFLG_MASK          = 0x0000003F;
const NFL4_UFLG_DENSE        = 0x00000001;
const NFL4_UFLG_COMMIT_THRU_MDS = 0x00000002;
const NFL4_UFLG_STRIPE_UNIT_SIZE_MASK
                              = 0xFFFFF0;

typedef uint32_t nfl_util4;

enum filelayout_hint_care4 {
    NFLH4_CARE_DENSE          = NFL4_UFLG_DENSE,

    NFLH4_CARE_COMMIT_THRU_MDS
                              = NFL4_UFLG_COMMIT_THRU_MDS,

    NFLH4_CARE_STRIPE_UNIT_SIZE
                              = 0x00000040,

    NFLH4_CARE_STRIPE_COUNT = 0x00000080;
};

/* Encoded in the loh_body field of data type layouthint4: */

struct nfsv4_1_file_layouthint4 {
    uint32_t      nflh_care;
    nfl_util4     nflh_util;
    count4        nflh_stripe_count;
};
```

The generic layout hint structure is described in Section 9.3.19. The client uses the layout hint in the layout\_hint (Section 11.16.4) attribute to indicate the preferred type of layout to be used for a newly created file. The layout-type-specific contents for the files layout in the layout hint is composed of three fields. The first field, nflh\_care, is a set of flags indicating which values of the hint the client cares about. If the NFLH4\_CARE\_DENSE flag is set, then the client indicates in the second field, nflh\_util, a preference for how the data file is packed (Section 20.8.4), which is controlled by the value of the expression nflh\_util & NFL4\_UFLG\_DENSE ("&" represents the bitwise AND operator). If the NFLH4\_CARE\_COMMIT\_THRU\_MDS flag is set, then the client indicates a preference for whether the client should send COMMIT operations to the metadata server or data server (Section 20.11), which is controlled by the value of nflh\_util & NFL4\_UFLG\_COMMIT\_THRU\_MDS. If the NFLH4\_CARE\_STRIPE\_UNIT\_SIZE flag is set, the client indicates its preferred stripe unit size, which is indicated in nflh\_util & NFL4\_UFLG\_STRIPE\_UNIT\_SIZE\_MASK (thus, the stripe unit size MUST be a multiple of 64 bytes). The minimum stripe unit size is 64 bytes. If the NFLH4\_CARE\_STRIPE\_COUNT flag is set, the client indicates in the third field, nflh\_stripe\_count, the stripe count. The stripe count multiplied by the stripe unit size is the stripe width.

#### 20.7.2. File Layout Content-related Data Types

When LAYOUTGET returns a files layout (indicated in the loc\_type field of the lo\_content field), the loc\_body field of the lo\_content field contains a value of data type nfsv4\_1\_file\_layout4. Among other content, nfsv4\_1\_file\_layout4 has a storage device ID (field nfl\_deviceid) of data type deviceid4. The GETDEVICEINFO operation maps a device ID to a storage device address (type device\_addr4). When GETDEVICEINFO returns a device address with a layout type of LAYOUT4\_NFSV4\_1\_FILES (the da\_layout\_type field), the da\_addr\_body field contains a value of data type nfsv4\_1\_file\_layout\_ds\_addr4.

```
typedef netaddr4 multipath_list4<>;

/*
 * Encoded in the da_addr_body field of
 * data type device_addr4:
 */
struct nfsv4_1_file_layout_ds_addr4 {
    uint32_t      nfl_da_stripe_indices<>;
    multipath_list4 nfl_da_multipath_ds_list<>;
};
```

The nfsv4\_1\_file\_layout\_ds\_addr4 data type represents the device address. It is composed of two fields:

1. `nflda_multipath_ds_list`: An array of lists of data servers, where each list can be one or more elements, and each element represents a data server address that may serve equally as the target of I/O operations (see Section 20.9). The length of this array might be different than the stripe count.
2. `nflda_stripe_indices`: An array of indices used to index into `nflda_multipath_ds_list`. The value of each element of `nflda_stripe_indices` MUST be less than the number of elements in `nflda_multipath_ds_list`. Each element of `nflda_multipath_ds_list` should be referred to by one or more elements of `nflda_stripe_indices`. The number of elements in `nflda_stripe_indices` is always equal to the stripe count.

```
/*  
 * Encoded in the loc_body field of  
 * data type layout_content4:  
 */  
struct nfsv4_1_file_layout4 {  
    deviceid4      nfl_deviceid;  
    nfl_util4      nfl_util;  
    uint32_t       nfl_first_stripe_index;  
    offset4        nfl_pattern_offset;  
    nfs_fh4        nfl_fh_list<>;  
};
```

The `nfsv4_1_file_layout4` data type represents the layout. It is composed of the following fields:

1. `nfl_deviceid`: The device ID that maps to a value of type `nfsv4_1_file_layout_ds_addr4`.
2. `nfl_util`: Like the `nflh_util` field of data type `nfsv4_1_file_layouthint4`, a compact representation of how the data on a file on each data server is packed, whether the client should send COMMIT operations to the metadata server or data server, and the stripe unit size. If a server returns two or more overlapping layouts, each stripe unit size in each overlapping layout MUST be the same.
3. `nfl_first_stripe_index`: The index into the first element of the `nflda_stripe_indices` array to use.

4. `nfl_pattern_offset`: This field is the logical offset into the file where the striping pattern starts. It is required for converting the client's logical I/O offset (e.g., the current offset in a POSIX file descriptor before the `read()` or `write()` system call is sent) into the stripe unit number (see Section 20.8.1).

If dense packing is used, then `nfl_pattern_offset` is also needed to convert the client's logical I/O offset to an offset on the file on the data server corresponding to the stripe unit number (See Section 20.8.4).

Note that `nfl_pattern_offset` is not always the same as `lo_offset`. For example, via the `LAYOUTGET` operation, a client might request a layout starting at offset 1000 of a file that has its striping pattern start at offset zero.

5. `nfl_fh_list`: An array of data server filehandles for each list of data servers in each element of the `nfl_da_multipath_ds_list` array. The number of elements in `nfl_fh_list` depends on whether sparse or dense packing is being used.

- \* If sparse packing is being used, the number of elements in `nfl_fh_list` MUST be one of three values:
  - Zero. This means that filehandles used for each data server are the same as the filehandle returned by the `OPEN` operation from the metadata server.
  - One. This means that every data server uses the same filehandle: what is specified in `nfl_fh_list[0]`.
  - The same number of elements in `nfl_da_multipath_ds_list`. Thus, in this case, when sending an I/O operation to any data server in `nfl_da_multipath_ds_list[X]`, the filehandle in `nfl_fh_list[X]` MUST be used.

See the discussion on sparse packing in Section 20.8.4.

- \* If dense packing is being used, the number of elements in `nfl_fh_list` MUST be the same as the number of elements in `nfl_da_stripe_indices`. Thus, when sending an I/O operation to any data server in `nfl_da_multipath_ds_list[nfl_da_stripe_indices[Y]]`, the filehandle in `nfl_fh_list[Y]` MUST be used. In addition, any time there exists `i` and `j`, (`i != j`), such that the intersection of `nfl_da_multipath_ds_list[nfl_da_stripe_indices[i]]` and



`nflda_multipath_ds_list[nflda_stripe_indices[j]]` is not empty, then `nfl_fh_list[i]` MUST NOT equal `nfl_fh_list[j]`. In other words, when dense packing is being used, if a data server appears in two or more units of a striping pattern, each reference to the data server MUST use a different filehandle.

Indeed, if there are multiple striping patterns, as indicated by the presence of multiple objects of data type `layout4` (either returned in one or multiple `LAYOUTGET` operations), and a data server is the target of a unit of one pattern and another unit of another pattern, then each reference to each data server MUST use a different filehandle.

See the discussion on dense packing in Section 20.8.4.

The details on the interpretation of the layout are in Section 20.8.

## 20.8. Interpreting the File Layout

### 20.8.1. Determining the Stripe Unit Number

To find the stripe unit number that corresponds to the client's logical file offset, the pattern offset will also be used. The *i*'th stripe unit (`SUi`) is:

```
relative_offset = file_offset - nfl_pattern_offset;  
SUi = floor(relative_offset / stripe_unit_size);
```

### 20.8.2. Interpreting the File Layout Using Sparse Packing

When sparse packing is used, the algorithm for determining the filehandle and set of data-server network addresses to write stripe unit *i* (`SUi`) to is:

```

stripe_count = number of elements in nflda_stripe_indices;

j = (SUi + nfl_first_stripe_index) % stripe_count;

idx = nflda_stripe_indices[j];

fh_count = number of elements in nfl_fh_list;
ds_count = number of elements in nflda_multipath_ds_list;

switch (fh_count) {
    case ds_count:
        fh = nfl_fh_list[idx];
        break;

    case 1:
        fh = nfl_fh_list[0];
        break;

    case 0:
        fh = filehandle returned by OPEN;
        break;

    default:
        throw a fatal exception;
        break;
}

address_list = nflda_multipath_ds_list[idx];

```

The client would then select a data server from `address_list`, and send a READ or WRITE operation using the filehandle specified in `fh`.

Consider the following example:

Suppose we have a device address consisting of seven data servers, arranged in three equivalence (Section 20.9) classes:

```
{ A, B, C, D }, { E }, { F, G }
```

where A through G are network addresses.

Then

```
nflda_multipath_ds_list<> = { A, B, C, D }, { E }, { F, G }
```

i.e.,

```
nflda_multipath_ds_list[0] = { A, B, C, D }
```

```
nfl_da_multipath_ds_list[1] = { E }
```

```
nfl_da_multipath_ds_list[2] = { F, G }
```

Suppose the striping index array is:

```
nfl_da_stripe_indices<> = { 2, 0, 1, 0 }
```

Now suppose the client gets a layout that has a device ID that maps to the above device address. The initial index contains

```
nfl_first_stripe_index = 2,
```

and the filehandle list is

```
nfl_fh_list = { 0x36, 0x87, 0x67 }.
```

If the client wants to write to SU0, the set of valid { network address, filehandle } combinations for SUi are determined by:

```
nfl_first_stripe_index = 2
```

So

```
idx = nfl_da_stripe_indices[(0 + 2) % 4]
    = nfl_da_stripe_indices[2]
    = 1
```

So

```
nfl_da_multipath_ds_list[1] = { E }
```

and

```
nfl_fh_list[1] = { 0x87 }
```

The client can thus write SU0 to { 0x87, { E } }.

The destinations of the first 13 stripe units are:

SU <sub>i</sub>	filehandle	data servers
0	87	E
1	36	A,B,C,D
2	67	F,G
3	36	A,B,C,D
4	87	E
5	36	A,B,C,D
6	67	F,G
7	36	A,B,C,D
8	87	E
9	36	A,B,C,D
10	67	F,G
11	36	A,B,C,D
12	87	E

Table 8

### 20.8.3. Interpreting the File Layout Using Dense Packing

When dense packing is used, the algorithm for determining the filehandle and set of data server network addresses to write stripe unit *i* (SU<sub>*i*</sub>) to is:

```

stripe_count = number of elements in nflda_stripe_indices;

j = (SUi + nfl_first_stripe_index) % stripe_count;

idx = nflda_stripe_indices[j];

fh_count = number of elements in nfl_fh_list;
ds_count = number of elements in nflda_multipath_ds_list;

switch (fh_count) {
  case stripe_count:
    fh = nfl_fh_list[j];
    break;

  default:
    throw a fatal exception;
    break;
}

address_list = nflda_multipath_ds_list[idx];

```

The client would then select a data server from `address_list`, and send a READ or WRITE operation using the filehandle specified in `fh`.

Consider the following example (which is the same as the sparse packing example, except for the filehandle list):

Suppose we have a device address consisting of seven data servers, arranged in three equivalence (Section 20.9) classes:

```
{ A, B, C, D }, { E }, { F, G }
```

where A through G are network addresses.

Then

```
nflda_multipath_ds_list<> = { A, B, C, D }, { E }, { F, G }
```

i.e.,

```
nflda_multipath_ds_list[0] = { A, B, C, D }
```

```
nflda_multipath_ds_list[1] = { E }
```

```
nflda_multipath_ds_list[2] = { F, G }
```

Suppose the striping index array is:

```
nfl_da_stripe_indices<> = { 2, 0, 1, 0 }
```

Now suppose the client gets a layout that has a device ID that maps to the above device address. The initial index contains

```
nfl_first_stripe_index = 2,
```

and

```
nfl_fh_list = { 0x67, 0x37, 0x87, 0x36 }.
```

The interesting examples for dense packing are SU1 and SU3 because each stripe unit refers to the same data server list, yet each stripe unit MUST use a different filehandle. If the client wants to write to SU1, the set of valid { network address, filehandle } combinations for SUi are determined by:

```
nfl_first_stripe_index = 2
```

So

```
j = (1 + 2) % 4 = 3
```

```
idx = nfl_da_stripe_indices[j]
```

```
= nfl_da_stripe_indices[3]
```

```
= 0
```

So

```
nfl_da_multipath_ds_list[0] = { A, B, C, D }
```

and

```
nfl_fh_list[3] = { 0x36 }
```

The client can thus write SU1 to { 0x36, { A, B, C, D } }.

For SU3,  $j = (3 + 2) \% 4 = 1$ , and `nfl_da_stripe_indices[1] = 0`. Then `nfl_da_multipath_ds_list[0] = { A, B, C, D }`, and `nfl_fh_list[1] = 0x37`. The client can thus write SU3 to { 0x37, { A, B, C, D } }.

The destinations of the first 13 stripe units are:

SUi	filehandle	data servers
0	87	E
1	36	A,B,C,D
2	67	F,G
3	37	A,B,C,D
4	87	E
5	36	A,B,C,D
6	67	F,G
7	37	A,B,C,D
8	87	E
9	36	A,B,C,D
10	67	F,G
11	37	A,B,C,D
12	87	E

Table 9

#### 20.8.4. Sparse and Dense Stripe Unit Packing

The flag NFL4\_UFLG\_DENSE of the nfl\_util4 data type (field nflh\_util of the data type nfsv4\_1\_file\_layouthint4 and field nfl\_util of data type nfsv4\_1\_file\_layout) specifies how the data is packed within the data file on a data server. It allows for two different data packings: sparse and dense. The packing type determines the calculation that will be made to map the client-visible file offset to the offset within the data file located on the data server.

If nfl\_util & NFL4\_UFLG\_DENSE is zero, this means that sparse packing is being used. Hence, the logical offsets of the file as viewed by a client sending READs and WRITEs directly to the metadata server are

the same offsets each data server uses when storing a stripe unit. The effect then, for striping patterns consisting of at least two stripe units, is for each data server file to be sparse or "holey". So for example, suppose there is a pattern with three stripe units, the stripe unit size is 4096 bytes, and there are three data servers in the pattern. Then, the file in data server 1 will have stripe units 0, 3, 6, 9, ... filled; data server 2's file will have stripe units 1, 4, 7, 10, ... filled; and data server 3's file will have stripe units 2, 5, 8, 11, ... filled. The unfilled stripe units of each file will be holes; hence, the files in each data server are sparse.

If sparse packing is being used and a client attempts I/O to one of the holes, then an error MUST be returned by the data server. Using the above example, if data server 3 received a READ or WRITE operation for block 4, the data server would return NFS4ERR\_PNFS\_IO\_HOLE. Thus, data servers need to understand the striping pattern in order to support sparse packing.

If nfl\_util & NFL4\_UFLG\_DENSE is one, this means that dense packing is being used, and the data server files have no holes. Dense packing might be selected because the data server does not (efficiently) support holey files or because the data server cannot recognize read-ahead unless there are no holes. If dense packing is indicated in the layout, the data files will be packed. Using the same striping pattern and stripe unit size that were used for the sparse packing example, the corresponding dense packing example would have all stripe units of all data files filled as follows:

- \* Logical stripe units 0, 3, 6, ... of the file would live on stripe units 0, 1, 2, ... of the file of data server 1.
- \* Logical stripe units 1, 4, 7, ... of the file would live on stripe units 0, 1, 2, ... of the file of data server 2.
- \* Logical stripe units 2, 5, 8, ... of the file would live on stripe units 0, 1, 2, ... of the file of data server 3.

Because dense packing does not leave holes on the data servers, the pNFS client is allowed to write to any offset of any data file of any data server in the stripe. Thus, the data servers need not know the file's striping pattern.

The calculation to determine the byte offset within the data file for dense data server layouts is:



```
stripe_width = stripe_unit_size * N;
  where N = number of elements in nflda_stripe_indices.

relative_offset = file_offset - nfl_pattern_offset;

data_file_offset = floor(relative_offset / stripe_width)
  * stripe_unit_size
  + relative_offset % stripe_unit_size
```

If dense packing is being used, and a data server appears more than once in a striping pattern, then to distinguish one stripe unit from another, the data server MUST use a different filehandle. Let's suppose there are two data servers. Logical stripe units 0, 3, 6 are served by data server 1; logical stripe units 1, 4, 7 are served by data server 2; and logical stripe units 2, 5, 8 are also served by data server 2. Unless data server 2 has two filehandles (each referring to a different data file), then, for example, a write to logical stripe unit 1 overwrites the write to logical stripe unit 2 because both logical stripe units are located in the same stripe unit (0) of data server 2.

#### 20.9. Multipathing to Data Servers

The NFSv4.1 file layout supports multipathing to multiple data server addresses. This can take the form of network multipathing or server multipathing:

- \* Network multipathing is used for bandwidth scaling via trunking (Section 7.5) and for higher availability of use in the event of a network discontinuity.
- \* Server Multipathing allows the client to switch to other data server addresses which may be that of another data server that is providing access to the same data, without having to contact the metadata server for a new layout.

Note however, that care needs to be taken when there are WRITES involved as users of the two data servers for multiple IOs can result in data corruption because this situation could result in formerly pieces of data becoming non-identical, making server-multipathing inappropriate until the data sets become identical once again. As a result, layouts based on server-multipathing can be recalled/revoked due to conflicts arising from writable layouts or WRITES performed by the MDS. For details see Section 20.18.

To support multipathing, each element of the `nfla_multipath_ds_list` contains an array of one more data server network addresses. This array (data type `multipath_list4`) represents a list of data servers (each identified by a network address), with the possibility that some data servers will appear in the list multiple times.

The client is free to use any of the network addresses as a destination to send data server requests. If some network addresses are less optimal paths to the data than others, then the MDS should not include those network addresses in an element of `nfla_multipath_ds_list`. If less optimal network addresses exist to provide failover, the RECOMMENDED method to offer the addresses is to provide them in a replacement device-ID-to-device-address mapping, or a replacement device ID. When a client finds that no data server in an element of `nfla_multipath_ds_list` responds, it should send a `GETDEVICEINFO` to attempt to replace the existing device-ID-to-device-address mappings. If the MDS detects that all data servers represented by an element of `nfla_multipath_ds_list` are unavailable, the MDS should send a `CB_NOTIFY_DEVICEID` (if the client has indicated it wants device ID notifications for changed device IDs) to change the device-ID-to-device-address mappings to the available data servers. If the device ID itself will be replaced, the MDS should recall all layouts with the device ID, and thus force the client to get new layouts and device ID mappings via `LAYOUTGET` and `GETDEVICEINFO`.

Typically, if two network addresses appear in an element of `nfla_multipath_ds_list`, they will designate the same data server, and the two data server addresses will support the implementation of client ID or session trunking (the latter is RECOMMENDED) as defined in Section 7.5. The two data server addresses will share the same server owner or major ID of the server owner.

It is not always necessary for the two data server addresses to designate the same server with trunking being used. One example, consists of situations in which the data could be read-only, and the data consist of exact replicas. In cases in which these replicas are subject to change, the occurrence of `WRITES` constitute a layout conflict (see Section 20.18 even if a layout is not used to perform the `WRITE`, leading to layout recall and/or revocation (see 20.17 and 20.18)).

## 20.10. Operations Sent to NFSv4.1 Data Servers

Clients accessing data on an NFSv4.1 data server MUST send only the NULL procedure and COMPOUND procedures whose operations are taken only from two restricted subsets of the operations defined as valid NFSv4.1 operations. Clients MUST use the filehandle specified by the layout when accessing data on NFSv4.1 data servers.

The first of these operation subsets consists of management operations. This subset consists of the BACKCHANNEL\_CTL, BIND\_CONN\_TO\_SESSION, CREATE\_SESSION, DESTROY\_CLIENTID, DESTROY\_SESSION, EXCHANGE\_ID, SECINFO\_NO\_NAME, SET\_SSV, and SEQUENCE operations. The client may use these operations in order to set up and maintain the appropriate client IDs, sessions, and security contexts involved in communication with the data server. Henceforth, these will be referred to as data-server housekeeping operations.

The second subset consists of COMMIT, READ, WRITE, and PUTFH. These operations MUST be used with a current filehandle specified by the layout. In the case of PUTFH, the new current filehandle MUST be one taken from the layout. Henceforth, these will be referred to as data-server I/O operations. As described in Section 18.7.1, a client MUST NOT send an I/O to a data server for which it does not hold a valid layout; the data server MUST reject such an I/O.

Unless the server has a concurrent non-data-server personality -- i.e., EXCHANGE\_ID results returned (EXCHGID4\_FLAG\_USE\_PNFS\_DS | EXCHGID4\_FLAG\_USE\_PNFS\_MDS) or (EXCHGID4\_FLAG\_USE\_PNFS\_DS | EXCHGID4\_FLAG\_USE\_NON\_PNFS) see Section 20.5 -- any attempted use of operations against a data server other than those specified in the two subsets above MUST return NFS4ERR\_NOTSUPP to the client.

When the server has concurrent data-server and non-data-server personalities, each COMPOUND sent by the client MUST be constructed so that it is appropriate to one of the two personalities, and it MUST NOT contain operations directed to a mix of those personalities. The server MUST enforce this. To understand the constraints, operations within a COMPOUND are divided into the following three classes:

1. An operation that is ambiguous regarding its personality assignment. This includes all of the data-server housekeeping operations. Additionally, if the server has assigned filehandles so that the ones defined by the layout are the same as those used by the metadata server, all operations using such filehandles are within this class, with the following exception. The exception is that if the operation uses a stateid that is incompatible with a data-server personality (e.g., a special stateid or the stateid

has a non-zero "seqid" field, see Section 20.14.1), the operation is in class 3, as described below. A COMPOUND containing multiple class 1 operations (and operations of no other class) MAY be sent to a server with multiple concurrent data server and non-data-server personalities.

2. An operation that is unambiguously referable to the data-server personality. This includes data-server I/O operations where the filehandle is one that can only be validly directed to the data-server personality.
3. An operation that is unambiguously referable to the non-data-server personality. This includes all COMPOUND operations that are neither data-server housekeeping nor data-server I/O operations, plus data-server I/O operations where the current fh (or the one to be made the current fh in the case of PUTFH) is only valid on the metadata server or where a stateid is used that is incompatible with the data server, i.e., is a special stateid or has a non-zero seqid value.

When a COMPOUND first executes an operation from class 3 above, it acts as a normal COMPOUND on any other server, and the data-server personality ceases to be relevant. There are no special restrictions on the operations in the COMPOUND to limit them to those for a data server. When a PUTFH is done, filehandles derived from the layout are not valid. If their format is not normally acceptable, then NFS4ERR\_BADHANDLE MUST result. Similarly, current filehandles for other operations do not accept filehandles derived from layouts and are not normally usable on the metadata server. Using these will result in NFS4ERR\_STALE.

When a COMPOUND first executes an operation from class 2, which would be PUTFH where the filehandle is one from a layout, the COMPOUND henceforth is interpreted with respect to the data-server personality. Operations outside the two classes discussed above MUST result in NFS4ERR\_NOTSUPP. Filehandles are validated using the rules of the data server, resulting in NFS4ERR\_BADHANDLE and/or NFS4ERR\_STALE even when they would not normally do so when addressed to the non-data-server personality. Stateids must obey the rules of the data server in that any use of special stateids or stateids with non-zero seqid values must result in NFS4ERR\_BAD\_STATEID.

Until the server first executes an operation from class 2 or class 3, the client MUST NOT depend on the operation being executed by either the data-server or the non-data-server personality. The server MUST pick one personality consistently for a given COMPOUND, with the only possible transition being a single one when the first operation from class 2 or class 3 is executed.

Because of the complexity induced by assigning filehandles so they can be used on both a data server and a metadata server, it is RECOMMENDED that where the same server can have both personalities, the server assign separate unique filehandles to both personalities. This makes it unambiguous for which server a given request is intended.

GETATTR and SETATTR MUST be directed to the metadata server. In the case of a SETATTR of the size attribute, the control protocol is responsible for propagating size updates/truncations to the data servers. In the case of extending WRITES to the data servers, the new size must be visible on the metadata server once a LAYOUTCOMMIT has completed and in some other circumstances (See Section 18.7.4). Section 20.16 which describes the mechanisms by which the client is to handle data-server files that do not reflect the metadata server's size).

#### 20.11. COMMIT through Metadata Server

The file layout provides two alternate means of providing for the commit of data written through data servers. The flag `NFL4_UFLG_COMMIT_THRU_MDS` in the field `nfl_util` of the file layout (data type `nfsv4_1_file_layout4`) is an indication from the metadata server to the client of the REQUIRED way of performing COMMIT, either by sending the COMMIT to the data server or the metadata server. These two methods of dealing with the issue correspond to broad styles of implementation for a pNFS server supporting the file layout type.

- \* When the flag is FALSE, COMMIT operations MUST to be sent to the data server to which the corresponding WRITE operations were sent. This approach is sometimes useful when file striping is implemented within the pNFS server (instead of the file system), with the individual data servers each implementing their own file systems.
- \* When the flag is TRUE, COMMIT operations MUST be sent to the metadata server, rather than to the individual data servers. This approach is sometimes useful when file striping is implemented within the clustered file system that is the backend to the pNFS server. In such an implementation, each COMMIT to each data server might result in repeated writes of metadata blocks to the detriment of write performance. Sending a single COMMIT to the metadata server can be more efficient when there exists a clustered file system capable of implementing such a coordinated COMMIT.

If `nfl_util` & `NFL4_UFLG_COMMIT_THRU_MDS` is TRUE, then in order to maintain the current NFSv4.1 commit and recovery model, the data servers MUST return a common `writeverf` verifier in all WRITE responses for a given file layout, and the metadata server's COMMIT implementation must return the same `writeverf`. The value of the `writeverf` verifier MUST be changed at the metadata server or any data server that is referenced in the layout, whenever there is a server event that can possibly lead to loss of uncommitted data. The scope of the verifier can be for a file or for the entire pNFS server. It might be more difficult for the server to maintain the verifier at the file level, but the benefit is that only events that impact a given file will require recovery action.

Note that if the layout specified dense packing, then the offset used to a COMMIT to the MDS may differ than that of an offset used to a COMMIT to the data server.

The single COMMIT to the metadata server will return a verifier, and the client should compare it to all the verifiers from the WRITES and fail the COMMIT if there are any mismatched verifiers. If COMMIT to the metadata server fails, the client should re-send WRITES for all the modified data in the file. The client should treat modified data with a mismatched verifier as a WRITE failure and try to recover by resending the WRITES to the original data server or using another path to that data if the layout has not been recalled. Alternatively, the client can obtain a new layout or it could rewrite the data directly to the metadata server. If `nfl_util` & `NFL4_UFLG_COMMIT_THRU_MDS` is FALSE, sending a COMMIT to the metadata server might have no effect. If `nfl_util` & `NFL4_UFLG_COMMIT_THRU_MDS` is FALSE, a COMMIT sent to the metadata server should be used only to commit data that was written to the metadata server. See Section 18.9.6 for recovery options.

## 20.12. The Layout Iomode

The layout iomode need not be used in all cases by the metadata server when servicing NFSv4.1 file-based layouts, although in some circumstances its use might be necessary. For example, if the server implementation supports reading from read-only replicas or mirrors, it is necessary for the server to return a layout enabling the client to do so. To support cases in which WRITES need to be prevented to ensure that data provided by multiple servers is the same (e.g. to support server-multipathing) the client needs to set the iomode based on its intent to read or write the data. The iomode need not be checked by all data servers when clients perform I/O. However, the data servers MUST still validate that the client holds a valid layout and return an error if the client does not.

### 20.13. LAYOUTCOMMIT on file layouts

As discussed in detail below, requirement for the use of LAYOUTCOMMIT are quite limited, since the control protocol has the ability to coordinate attribute changes arising from WRITE operations and make the result visible to the metadata server and clients. This information, together with the details below, provides the appropriate file-layout-related information required in Section 19.1.4.

- \* There is no need for the client to provide the file size as part of LAYOUTCOMMIT.
- \* There is no need for the client to provide a modified time as part of LAYOUTCOMMIT.
- \* When LAYOUTCOMMIT is done the resulting change attribute should reflect either the change value derived from the modified time or a value treating the LAYOUTCOMMIT as if it were a WRITE.
- \* There is no provision for additional information to be provided as part of LAYOUTCOMMIT, as it is for other layout types.

For file layouts, WRITES to a Data Server that return a stable\_how4 value of FILE\_SYNC4 guarantee that data and file system metadata are on stable storage. This implies that a LAYOUTCOMMIT is not needed in order to make the data and metadata visible to the metadata server and other clients.

For file layouts, when WRITE to the data server returns UNSTABLE4 or DATA\_SYNC4 and the NFL4\_UFLG\_COMMIT\_THRU\_MDS flag is set, the client MUST send the COMMIT to the metadata server. A successful COMMIT to the metadata server guarantees that data and file system metadata are on stable storage. As a result, any time that NFS4\_UFLG\_COMMIT\_THRU\_MDS is set, a LAYOUTCOMMIT (of the byte range specified by the layout) is not needed.

For file layouts, when NFL4\_UFLG\_COMMIT\_THRU\_MDS flag is not set, and WRITE or COMMIT to the data server return DATA\_SYNC4, the client MUST send the LAYOUTCOMMIT to the metadata server in order to synchronize file metadata.

The following table summarizes the conditions under which a LAYOUTCOMMIT is needed, and the effects of a COMMIT to a data server and metadata server.

NFL4_UFLG_ COMMIT_ THRU_MDS	WRITE to DS returns	Meaning of COMMIT to DS	Meaning of COMMIT to DS	LAYOUT COMMIT reequired
Not Set	UNSTABLE	DATA_SYNC4	Nothing	Yes
Not Set	DATA_SYNC4	Nothing	Nothing	Yes
Not Set	FILE_SYNC4	Nothing	Nothing	NO
Set	UNSTABLE	Nothing	FILE_SYNC4	NO
Set	DATA_SYNC4	Nothing	FILE_SYNC4	NO
Not Set	FILE_SYNC4	Nothing	Nothing	NO

Table 10

Note that a client can always demand FILE\_SYNC4 or DATA\_SYNC4 via WRITE arguments. Also note that specifying these stability levels might adversely impact performance.

If a LAYOUTCOMMIT is required, it should be sent before CLOSE to maintain close-to-open semantics. If required, it should be sent before LOCKU, OPEN\_DOWNGRADE, LAYOUTRETURN, and when the application issues fsync() [fsync]. Again, if LAYOUTCOMMIT is required, it should be sent periodically to keep the file size and modification time approximately up-to-date. This allows the use of commands such as "tail -f" which copies its input file to the standard output and updates the output as new lines become available in the input file. It is the client implementation's choice to determine how frequently LAYOUTCOMMIT is issued. Possible policies include every N'th COMMIT to a data server, every N'th unit of time, or after writing a stripe to a set of data servers.

Even if a required LAYOUTCOMMIT is not issued by the client, the data server and metadata servers have a set of responsibilities necessary to provide data consistency:

- 1) Data servers MUST commit data and synchronize modification and size attributes with the metadata server before a layout revocation is completed as described in section 12.5.4.



- 2) Data servers should commit data and synchronize modification and size attributes with the metadata server after the metadata server reboots. In theory the client should commit the data, but this avoids the problem where both the client and metadata server crash at the same time.
- 3) The metadata server MAY periodically poll data servers to synchronize modification and size attributes.

#### 20.14. File Layout Type Control Protocol

Although the specifics of the control protocol that might be used are not specified in this document, implementations of the file layout type need to provide control mechanisms, often in the form of a control protocol to provide the functions described in Sections 19.1.3 through 19.1.5.

One important function is the coordination of stateid- related information as described in 20.14.1 and 20.14.2

The functions described by Section 19.1.3 are provided as discussed below:

- \* Revocation needs to be supported. One common way of doing this is integrating the cancellation of the sharing of layout stateids with sharing mechanisms described in Section 20.14.1.
- \* Similarly, when layouts are returned, the metadata server needs to be made aware of their deletion.

This involves the layout management functions described in Section 20.14.3.

The functions described by Section 19.1.4 are provided as discussed below:

- \* Execution of I/O requests directed to the metadata server is discussed in Section 20.19.
- \* Appropriate authorization of IO requests sent to data servers are discussed in Section 20.14.2.
- \* Information relating to the committing of layouts, used to update the metadata server's view of attributes that can be changed as a consequence of IO operations, is provided in Section 20.13.

- \* Information relating to the use of the control protocol in providing needed storage allocation is provided in Section 20.14.4.

#### 20.14.1. Control Protocol State Coordination

As described in Section 20.4 it is often necessary for the metadata server and some set of data server to share knowledge about the lock-related objects denoted by stateids that may be used in IO operations (i.e. open stateids, lock stateids, and delegation stateids) and layout stateids. In each case the unique portion of the stateid has an associated type and a file with which it is associated. Depending on the type of stateids used with IO operations, the following additional information needs to be shared if the file in question has a layout involving one or more data servers.

- \* For open stateids, the open mode and the opening principal need to be provided.

Because previous specifications required open-upgrade to be done even when the principals were different, following this guidance might make it impossible to correctly provide the principal as described above. Although Section 14.9 corrects this issue, implementers need to consider whether doing so (currently specified as "SHOULD NOT") is still possible in this context even though it might be possible in non-pNFS contexts.

When these change due to use of OPEN\_DOWNGRADE or CLOSEs and revocations of open state (treated as reducing the open mode to zero), the client and associated data servers need to be made aware of the changes of state.

Deny modes do not need to be provided to the data server since IO operations outside of OPENS are not allowed to be sent data servers when the file layout type is used. In this situation, the open mode is always checked and there is no reason for the deny mode to be accessed during checking.

- \* For delegation stateids, the type of delegation (read or write) needs to be provided to the data server.
- \* For byte-range lock stateids, unless mandatory byte-range locking is supported, sharing of this information is unnecessary.

Such locks do not affect the processing of IO operation.

When mandatory byte-range locking is supported, the mode, bounds and lockowner needs to be made available, unless, as discussed in Section 20.20, data server support for this feature can be avoided.

- \* For layout stateids, the range and IO mode needs to be provided.

If the metadata server chooses not to provide this information for locking state created when there are no layouts forcing it to be provided, then it is obliged to provide this information together with the layout that makes its sharing necessary.

Given the coordination discussed above, the following rules are to be used in determining the stateids to be used when client send IO requests to data servers based of file type layouts. They differ from those in Section 13.2.5.

- \* If the entity corresponding to the lock-owner (e.g., a process) sending the I/O has a mandatory byte-range lock stateid for the associated open file, then the byte-range lock stateid for that lock-owner and open file should be used.
- \* If there is no appropriate mandatory byte-range lock stateid, then the OPEN stateid for the open file in question should be used.

Although clients who hold delegations can do the equivalent of opening files on their own, in such cases there is no MDS assigned open stateid and the delegation is to be used instead.

- \* If there is neither an open stateid nor a mandatory byte-range lock stateid, a delegation stateid, if it exists, should be used.
- \* If none of the above are available or there is no associated layout, the IO MUST be sent to the MDS. See Section 20.19.

There are three basic approaches to sharing that can be used for the above and for information that needs to be shared for reason discussed in Section 20.14.2. Control protocols can use one of these or a workable combination due to workloads and performance requirements.

- \* Synchronous Propagation of changes is often used although its contribution to latency often makes it desirable to find an alternative.
- \* Asynchronous Propagation with Verification, when available, can be used to reduce latency by providing a way for a subsequent operation to determine whether the necessary change has occurred.

- \* Synchronous Interrogation can be used to make it simpler to obtain the current data directly from the metadata server using the control protocol.

Because valid IO execution required a layout and a state id to be sent with the IO request (i.e. an OPEN, delegation, or lock stateid), the data server will need these only when stateids of both types are available. This circumstance can be dealt with in a number of ways:

- \* Both the layouts and IO-request-associated stateids are propagated independently to the data server. As a result, the data server will honor an IO request when it has both a layout (including the associated filehandle sent as part of the layout) and a valid stateid.

In this case, the propagation of the of the IO stateids need to be done synchronously while asynchronous propagation with verification can be used for layout propagation.

Propagating stateid information even when layouts are not present simplifies the handling of IO requests sent to the MDS. See Section 20.19 for further discussion.

- \* Propagation can be limited to cases in which both necessary elements are present. In this approach, stateids are not sent to data servers unless a layout exists. As a result, when a layout is generated, stateid information for a number of stateids will need to be sent to the data server.
- \* A synchronous interrogation model can be used to avoid the need to propagate stateid information other than layouts. Instead, the MDS, who has the necessary information is interrogate using the control protocol. Caching of successful interrogation can be used but it will need to be supplemented layout protocol messages advising the data server of a change in the associated stateid set.

In the case of IO directed to the MDS, the same interrogation paradigm is used locally, on the MDS.

In either of the cases described above, data servers need to be made of the creation and deletion of layouts. This propagation is necessary for the files layout type because state changes need to be shared among three nodes. The consequences are discussed in Section 20.14.3.

#### 20.14.2. Control Protocol Authorization Coordination

As provided for in Section 20.14.1, the need to apply the IO authorization arrived at via OPEN, in the context of use of file layouts can be addressed by providing data servers that receive layout information with additional authorization-related data about the stateids that can be used with those layouts. This information will generally include identification of the principal that opened and information about whether the open is for READ, WRITE or both.

While that addresses the common case in which the IO is only requested by the OPENing principal, there is a need to provide valid authorization when other principals perform layout-based IO operations. There are multiple facilities the control protocol can provide that can serve as part of a suitable facility to check for valid authorization:

- \* Provide the ability within the control protocol for the data server to request that an ACCESS-like operation be provided by the MDS.
- \* Provide notification to LAYOUT-handling data servers of changes in authorization-related attributes. This could used to allow caching of ACCESS results with flushing of the cache on attribute changes,

In the case in which IO operations are done using a delegation stateid, it is the responsibility of the client performing OPENS locally to check for authorization for the Opening principal locally. Once this is done, I/O operations are authorized as discussed above in that I/O performed by the opener are authorized based on the authorization done at OPEN time, while those done by other principals on the same client need to be authorized using the control protocol facilities discussed above.

#### 20.14.3. Control Protocol Layout Management

Due to the structure of implementations of the file layout type the role of the data server in layout management is limited to responding to propagation initiated by the MDS. This includes case of layout creation and layout termination.

#### 20.14.4. Control Protocol Role in Storage Allocation for File Layout Type

Within the NFSv4.1 protocol there is relatively little need for the control protocol to concern itself with matters related to storage allocation, as explained below:

- \* Allocation of new storage (and possible deallocation of existing storage) as a result of WRITES happen to the data server just as they would if the WRITE was forwarded by the MDS or performed by the data server without pNFS involvement.
- \* In case of file removal or truncation, the necessary storage deallocation is the responsible of the data server executing a control program requests issued by the MDS.

#### 20.14.5. IO Validation for the Files Layout Type

IO validation, which needs to be done on the data server consists of the following elements:

- \* Check for layout validity including the existence of a layout for the file accessed by the client and the range of bytes being within those specified

The validity of the layout iomode need not always be checked but support for some features (e.g. server multipathing) are not supported without checking this, so, for configurations in which such support is needed, WRITES to read-only layouts cannot be processed and might cause layout revocation as described in Section 20.18.

- \* Checking the validity of the stateid.

This generally requires control protocol support as described in Section 20.14.1. In addition, if mandatory byte-range locks are supported without pNFS, they need to be supported when the files layout type is used as discussed in Section 20.20.

- \* Checking for proper IO authorization.

This is not necessary the stateid is for an open file and the principal that opened the file (made available as described in Section 20.14.1) is performing the IO operation.

#### 20.15. File Attributes

Since the SETATTR operation has the ability to modify state that is visible on both the metadata and data servers (e.g., the size), care must be taken to ensure that the resultant state across the set of data servers is consistent, especially when truncating or growing the file.

As described earlier, the LAYOUTCOMMIT operation is used to ensure that the metadata is synchronized with changes made to the data servers. For the NFSv4.1-based data storage protocol, it may be necessary to re-synchronize state such as the size attribute, and the setting of mtime/change/ctime. See Section 18.7.4 for a full description of the semantics regarding LAYOUTCOMMIT and attribute synchronization. It should be noted that by using an NFSv4.1-based layout type, it is possible to synchronize this state before LAYOUTCOMMIT occurs. For example, the control protocol can be used to query the attributes present on the data servers.

The OPEN operation (Section 25.16.4) does not impose any requirement that I/O operations on an open file have the same credentials as the OPEN itself (unless EXCHGID4\_FLAG\_BIND\_PRINC\_STATEID is set when EXCHANGE\_ID creates the client ID), and so it requires the server's READ and WRITE operations to perform appropriate access checking, when the principal making the IO request is not the same as that doing the OPEN. As a result, changes to ACLs and other authorization-related attributes (i.e. mode, owner, group\_owner, or similar attributes added later) may require effective sharing of the new attribute values to deal with IO being done by principals other than the opener, as discussed in Section 20.14.2.

#### 20.16. Data Server Component File Size

A problem can arise when a component data file on a particular data server has grown past EOF. This problem can occur for both dense and sparse layouts.

Consider the following scenario: a client creates a new file (size == 0) and writes to byte 131072; the client then seeks to the beginning of the file and reads byte 100. The client expects to receive zeroes back as a result of the READ. However, if the striping pattern directs the client to send the READ to a data server other than the one that received the client's original WRITE, the data server servicing the READ might believe that the file's size is still 0 bytes. In that event, the data server's READ response would contain zero bytes and an indication of EOF.

The data server is only able to return zeroes if it knows that the file's size has been extended. While it might be possible to arrange the control protocol to provide the requisite knowledge, no such requirement has dealt with this issue in the past, making it possible that data server could return EOF as discussed above.

For this reason, when the offset in the arguments to READ is less than the client's view of the file size and if the READ response indicates EOF and/or contains fewer bytes than requested, the client will need to interpret such a response as a hole in the file, and the NFS client would substitute zeroes for the data.

Once the file is closed and a subsequent OPEN of the file is done, the change attribute can be inspected for a difference from a cached value for the change attribute. In cases such as those described above, this implies that, if LAYOUTCOMMIT is done at close whenever necessary (See Section 20.13), subsequent READs cannot encounter the issue.

#### 20.17. Revocation and Fencing for the Files Layout Type

This section provides information regarding the possibility of layout revocation, as provided for by Section 19.1.3

There are a number of situations in which the implementations of the file layout type, through its associated control protocol, need to cancel existing layout and make it impossible for data servers to use those layouts to execute further I/O operations using those layouts. When this process is applied to all layouts associated with a client for a file, the client is said to be "fenced" from access to the file in question.

The following situations can result in this sort of layout cancellation:

- \* Lease expiration for a client's lease to the MDS. In this case all data server IO for the client is prevented.
- \* A Metadata sever restart results in all layouts issued by the previous MDS instance being cancelled.
- \* A client fails to respond a CB\_LAYOUTRECALL, making it necessary for the MDs to cancel the layout without the client's assent.
- \* administrative intervention.
- \* A conflict such as those described in Section 20.18. The include the issue of conflicting layouts or other IO operations that invalidate the assumptions under which layout were previously granted.

Fencing works as follows. As described in Section 20.5, in COMPOUND procedure requests to the data server, the data filehandle provided by the PUTFH operation and the stateid in the READ or WRITE operation



are used to ensure that the client has a valid layout for the I/O being performed. If it does not have such a layout, the I/O is rejected with NFS4ERR\_PNFS\_NO\_LAYOUT. The server can simply check the stateid and, additionally, make the data filehandle stale if the layout specified a data filehandle that is different from the metadata server's filehandle for the file (See the nfl\_fh\_list description in Section 20.7.2).

The procedure described above eliminate the possibility of execution of "lingering" WRITES, since requests that have not reached the data server cannot be executed. However, it is possible that IO requests that have passed this check might not complete in a reasonable time. In this case, the client has the option of ceasing to wait and considering the request failed or lease expiration might force it to be discarded. In either case, there is no possibility of the request being reissued since it will fail the check due to revocation of the layout.

Before the metadata server takes any action to revoke layout state given out by a previous instance, it must make sure that all layout state from that previous instance are invalidated at the data servers. This has the following implications.

- \* The metadata server must not restripe a file until it has contacted all of the data servers to invalidate the layouts from the previous instance.
- \* The metadata server must not give out mandatory locks that conflict with layouts from the previous instance without either doing a specific layout invalidation (as it would have to do anyway) or doing a global data server invalidation.

#### 20.18. Layout/IO Conflicts for the Files Layout Type

When the files layout type is used, conflicts between overlapping layouts can arise in some circumstances when at least one of the layouts allows WRITES to be performed.

Such circumstances can arise when server-multipathing (as opposed to network-multipathing) is in effect. (See Sections 20.9 and 20.12 for details).

When server-multipathing is in effect, two distinct regions of storage (or the same regions as mediated by two distinct data server need to be kept identical. The presence of WRITES to one or more of those paths can invalidate this assumption in the ways listed below. This can happen even when the iomodes are different.

- \* If two distinct storage areas are involved an update to one will result in incorrect data being returned by some READs
- \* If access to the same storage area is mediated by two different data servers, caching done by either server will result in incorrect results. When persistent write caching is used to acknowledge stable WRITES before WRITES to the ultimate storage location are performed, the problem is exacerbated
- \* When WRITES are done to multiple paths, allocation might be done separately on separate data servers, resulting in data corruption, unless the two servers coordinate allocation, which they are rarely in a position to do.

When WRITES are performed by the MDS, the same sorts of problems arise. As a result, when server-multipathing is being used to effect read-sharing, WRITES, whether issued to the MDS or data servers, need to cause the revocation of WRITE-capable layouts so that we can establish a state of affairs in which either multiple read-only layouts or a single READ-WRITE layout is in effect.

#### 20.19. File Layout MDS-directed IO

When IO requests are directed at a metadata server and the data is located on a different server, the request is typically sent to the appropriate data server even if there is no layout in effect that would support the request being sent to that data server directly.

Despite the existing distribution arrangement which might arise in the context of a distributed clustered NFSv4 server, there a number of issues that require special attention.

- \* Validation of the stateid is often difficult to do on the data server since stateid creation is an MDS-specific activity.
- \* Authorization checking, for the case in which the principal performing the IO is not the opener.

#### 20.20. File Layout Handling of Mandatory Byte-Range Locks

Although mandatory byte range locking is an OPTIONAL feature that is not commonly used, there will be situations in which it is present and needs to be accommodated in an environment in which pNFS using the files layout type needs to be supported. There are multiple approaches to doing so. These include:

- \* Those that exclude layouts and byte-range locks on the same files.

This can be accomplished by recalling layouts associated with a file before granting a mandatory byte-range lock. As a result IO operations on such files will be formed by the MDS as describes in Section 20.19.

Although this approach is not mentioned in [RFC8881] as a way to support the requirement to honor these locks, it is, in many environments, a way to take advantage of pNFS while avoiding its use where it would hurt performance.

- \* Those that try to accommodate these features on the same file.

We are not considering the possibility of the data server requesting IO clearance from MDS, since it similar but poorer-performing than sending the request directly to the MDS.

This would require the propagation of byte-range locks and unlocks to data servers (usually one but sometimes more) who might perform IOs overlapping the range specified range. The fact that it is hard to do this propagation asynchronously makes it a poor choice in terms of performance when such locking is done.

- \* Those try to exclude read-write layouts and exclusive mandatory byte range locks on the same file. This requires recall of write-capable layouts before granting a request for an exclusive mandatory byte-range lock.

This allows READs (more common than WRITES) and shared byte range locks to avoid the additional overhead resulting from byte-range locking.

In other cases, the IO can be sent to the MDS or the locks can be propagated as in the case described above.

The handling (and potential mishandling) of read-only layout needs to be considered before using this approach. While this and previous specifications clearly give the responsibility of enforcing these to the client, the fact that the data server had previously been given latitude in this area is concerning,, since, if the client enforced these, server laxity as to enforcement would never be needed. As a result, those anticipating this strategy need to be sure that read-only layouts are enforced in fact for this to work, with the specification's view that this is possibly non-compliant not improving things.

## 20.21. Layout Recall Issues for File Layout Type

There are some noteworthy peculiarities regarding layout recalls that result from file deletion. These include cases in which a file is deleted because another file is renamed to have the existing name deleting the original holder of that name.

In such case other layout types will need to arrange for the reclamation of storage and the recall of associated layouts to prevent further access to the deleted file. The files layout type is different in that:

- \* Although the storage occupied by the file does need to be freed, this is more appropriately considered apart from storage allocation facilities such as those used to support file truncation. It is better to think of this as a remote file deletion request effected by the control protocol.
- \* Because of the proceeding item, there is no possibility of the existing layout being used to access the deleted file. As a result, immediate layout recall is not needed.
- \* Since the layout in question cannot be used, it will need to be cleaned up eventually and this can occur without coordination between the MDS and the client who can each drop the layout independently.

## 20.22. Restrictions on Layout Information Discard for File Layout Type

As discussed in Section 18.7.3.1, clients are allowed to discard records of existing layouts, it is the obligation of the layout type's specification to clearly state any restrictions on the ability of the client and the MDS to do so. Specifically:

- \* The client cannot use layouts that have been discarded because it is required to validate IOs sent to the data server.

As a result, it needs to avoid discarding information that is likely to be needed again.

- \* Whether the MDS can discard layout information depends on whether the control protocol supports retrieving it from the data server.

In the case in which it can, it is free to discard such information, only concerning itself with the likely cost of retrieving it when necessary.

- \* Data servers MUST NOT discard layout information since doing so would cause them to reject valid IO requests. Such information can only be dropped in response to layout return or revocation.

#### 20.23. Dealing with Client and Server Failure

This treatment is considerably more flexible than that appearing in previous NFSv4.1 specifications. This expanded scope derives from two important considerations:

- \* The appearance in Section 18.43.3 of [RFC5661] and [RFC8881] of the clause "If the metadata server is in a grace period, and does not persist layouts and device ID to device address mappings" suggesting that the metadata might validly do these things and respond substantively during the grace period.

Despite the fact that this suggestion is not followed up in descriptions of pNFS recovery or of the file layout type, we are exploring the implications of this choice here, referring to the metadata server making this choice as being "layout-persistent".

- \* The inclusion, in changes made in [RFC8881], of the possibility of clientids and their associated state surviving a disruptive event such as a node restart or inter-node transfer.

We refer to such occasions as being "state-persistent" and assume, for our purposes, that layouts would be included among the persistent stateids.

- \* Because of the change of descriptive approach prompted by [RFC8434], it is now clearer that the control protocol has the option of making a new instance of a data server after restart aware of layouts and associated stateids held by the previous data server instance.

In the current descriptive framework, this might seem a necessary corollary to the control protocols role in providing sharing for such information, but within the previous descriptive framework which focused on specific instances of state propagation, the absence of specific mention could have been interpreted as foreclosing this possibility.

We refer to the state of affairs in which layout and their associated stateids are known to a new data server instance as manifesting "layout continuity", with the understanding that layout continuity can arise as a result of either layout persistence or be provided by the control protocol together with the MDS' uninterrupted existence.

The OPTIONAL character of these potential choices makes it important that to deal with clients unaware of their existence so it needs to be made clear:

- \* How these choices have dependencies or where they are independent.
- \* How the client can determine whether they are in effect.
- \* How a client unaware of these choices could successfully recover, even if it does so in a more complicated and time-consuming manner than it would if it were aware of these choices.

We deal with issues related to recovery from MDS, data server, and client failure in Sections 20.23.1, 20.23.2, and 20.23.3 respectively.

#### 20.23.1. Dealing with MDS Failure

Because no provision is made is made for reclaiming layouts, when a client sees an MDS restart, it normally needs to do the actions below. While we will discuss possible checks for state-persistence or layout-persistence, the discussion of how these situations are dealt with is deferred to later in the section.

These actions need to performed, independently of the possibility that the failed node was acting as a data server as well a metadata sever. In such cases, once recovery from the MDS failure, the recovery from the data server failure needs to be done, as described in Section 20.23.2.

Recovery needs to proceed as follows:

- \* It is appropriate to consider the possibility of state-persistence at this point.

State-persistence can arise if the client-MDS session is a persistent one or if, in establishing a new session, it is found that the original client-MDS clientid is still in existence.

In either case, we deal with the situation as described later in this section, whether layout-persistence is present or not.

- \* A new clientid and session is obtained using a sequence of an EXCHANGE\_ID operation followed by a CREATE\_SESSION operation using that client ID (eir\_clientid as returned from EXCHANGE\_ID) is required to establish and confirm the client ID on the server. This clientid serves as the replacement for the clientid terminated by the MDS restart.

- \* Locks other than layouts need to be reclaimed as they for other (non-pNFS) occasions of server restart. Whether any of these exist or not, the determination of whether layout-persistence still needs to be done although it is up to the client when in the process this happens.

The presence of layout persistence can best be determined by attempting to use an existing layout to perform a short READ operation which only succeeds if layout-persistence is in effect. It is easier to interpret failures if the stateid under which the READ is to be done is already reclaimed. Once a successful READ is done, layout persistence is established and stateids, once reclaimed, can be used together with pre-existing layouts to continue work delayed by the MDS restart.

- \* At that point, recovery is complete even though layouts cannot be used. At this point, RECAIM\_COMPLETE can be done.

Once the last such RECAIM\_COMPLETE is done, the grace period is over and LAYOUTGET can be done to obtain new layouts, completing the recovery process. It is the responsibility of the control protocol, as described in Section 20.14.1, to provide stateid information to the data server for files for which a layout is obtained.

At this point, the basic recovery process, without state-persistence, is complete.

The rest of this section will discuss the case in which state persistence is in effect and its presence can allow the recovery process to be abbreviated. In this case, further handling depends on the possible presence of layout-persistence which can be determined by using an existing layout to do a READ using pNFS. If it succeeds, recovery is done and normal processing can continue.

In the case, layout-persistence is not in effect, the client needs to re-establish needed layouts once state recovery is complete. Because other clients may have obtained state from the restarting server, even if the current client has not, such use can result in failure if layout-persistence is not in effect until the grace period is over. Once the grace period is over, replacement layouts can be obtained and used.

### 20.23.2. Dealing with Data Server Failure

There is potential uncertainty about the actions necessary for recovery from a data server restart because of the way that the description of the files layout type has been forced to evolve in response to the creation of [RFC8434] and its incorporation in the NFSv4.1 respecification effort:

In essence, because earlier specification efforts did not specify how recovery was to be done, we are forced to speculate based on the existing specifications regarding the control protocol about how these might be used to satisfy recovery needs in the case of data server restart. This creates a number of questions that need to be addressed.

- \* In [RFC5661] and [RFC8881] the general pNFS description indicated this (restart of a storage device) was a layout-type specific matter but the section on the files layout type did not address the issue.

This left considerable uncertainty about whether the restarting data server would have access to lock stateids and layouts established by the previous instance (possibly provided by the control protocol).

While one might expect that the control protocol responsible for propagating this information in other cases could do upon restart, there is no explicit statement to this effect. Furthermore, the appearance of specific directives to propagate information in other situations could be interpreted as allowing the information not to be propagated in this case.

- \* The publication of [RFC8434] and its subsequent incorporation in rfc8881lbis changes things in that it focused on the responsibilities of the control protocol to provide effective sharing and avoided specifics about how that sharing was to be done (via information propagation or later cross-node interrogation).

Even though this new descriptive approach made it natural for the control protocol to provide the necessary sharing for simpler recovery, there is no evidence for the existence of implementations that do so.

- \* Given the nature of structure of stateids and clientid mandated by the files layout type, improving recovery in this way is not expected to be difficult.



As a result, we have the following possibilities, for the situations that a client might see after data server restart:

- (A) The control protocol, as part of sharing layout information including information about associated stateids has made this information available to the new data server instance, once it is aware of a new instance of the original server owner.

We refer to this situation as manifesting full state-continuity.

- (B) As part of implementing an approximation to a "global" stateid space, making stateids that are part of a given clientid makes these available (in read-only fashion) to the new instance of the server s soon as it establishes a connection to the existing (originally MDS-based) clientid.

Although layout stateids might be propagated, we are assuming, in this case, that, in the absence of full control protocol involvement, functional layouts are not immediately available.

We refer to this situation as manifesting limited state-continuity.

- (C) The new data server initializes itself with an empty stateid set, despite the use of a common clientid with the data server.

This inconsistency might be expected to be troublesome, but it could exist and it is possible that it resolves itself once new layouts are gotten, triggering the transfer of associated stateid information.

We refer to this situation as manifesting state-discontinuity.

- (D) The new data server initializes itself with an empty stateid set and the MDS maintains synchrony between the two nodes by resetting the shared clientid to have an empty stateid set.

We refer to this situation as manifesting state-destruction.

As a result of limited testing of available client and server implementations, we have found instances of (D) but no bases of (A), (B), or (C). While reasonable arguments could be made that these behaviors are valid and acceptable as defined by [RFC8881], the current draft will assume (D) that servers will implement that and that clients can validly assume that server will behave in that way until and unless the working group decides differently. See Appendix D.2.14 for further information.

At this point, we are assuming that (D) is in effect and the necessary recovery is as described below"

We need to reobtain any locks that we held. It is not clear that RECLAIM is possible, since the MDS is not in a grace period. As a result, it is possible the locks might be unavailable due to the fact that a conflicting lock might be obtained. In either case, we need to reobtain layouts in order to re-establish pNFS access. If these are not obtainable, IO can be directed to the MDS, as described in Section 20.19.

#### 20.23.3. Dealing with Client Failure

There appears nothing layout-type-specific to add to the discussion in Section 18.9.1

#### 20.24. Dealing with Lease Expiration

Although there are leases associated with the a client connection to MDS and to data servers, we need to treat each of the following cases separately:

- \* There is a lease loss between the client and a node acting as the MDS.

In this case, the handling described in Section 18.9.2 can be followed

- \* There is a lease loss between the client a single node acting as both MDS and data server.

In this case as well, the handling described in Section 18.9.2 can be followed since the client has no locks owned by the data server.

- \* There is a lease loss between the client and a node acting as a data server.

In this case, there should be no lease because there are no lock held on the data server. However, it is possible that a data sever might treat the lack of session use over the lease period as amounting to a lease loss. If so, since there are no locks to be lost, recovery should be trivial.

## 20.25. Security Issues for the File Layout Type

All IO performed using file type layouts MUST adhere to the security considerations outlined in Section 18.11. To effect this requirement, NFSv4.1 data servers MUST perform the required access checks for each READ or WRITE operation where the IO has not been not authorized by a preceding OPEN allowing that IO type.

Because of lax specification practices with regard to authorization, in previous NFSv4.1 specifications, additional checks, matching those that would be done in the absence of pNFS might be needed. If the metadata server would deny a READ or WRITE operation on a file due to its ACL, mode attribute, open access mode, mandatory byte-range lock state, or any other attributes and state, the data server MUST also deny the READ or WRITE operation. This impacts the control protocol and the propagation of state from the metadata server to the data servers; see Sections 20.14.1, 20.14.2, and 20.20 for more details, including explanations of why the deny mode of each Open does not need to be known by data servers.

The methods for authentication, integrity, and privacy for data servers based on the LAYOUT4\_NFSV4\_1\_FILES layout type are the same as those used by metadata servers. Metadata and data servers use ONC RPC security flavors to authenticate, and SECINFO and SECINFO\_NO\_NAME to negotiate the security mechanism and services to be used. Thus, when using the LAYOUT4\_NFSV4\_1\_FILES layout type, the impact on the RPC-based security model due to pNFS (as alluded to in Sections 2.7 and 2.8.2) is zero.

For a given file object, a metadata server MAY require different security parameters (secinfo4 value) than the data server. For a given file object with multiple data servers, the secinfo4 value should be the same across all data servers. If the secinfo4 values across a metadata server and its data servers differ for a specific file, the mapping of the principal to the server's internal user identifier MUST be the same in order for the access-control checks based on ACL, mode, open and deny mode, and mandatory locking to be consistent across on the pNFS server.

If an NFSv4.1 implementation supports pNFS and supports NFSv4.1 file layouts, then the implementation MUST support the SECINFO\_NO\_NAME operation on both the metadata and data servers.

## 20.26. Necessary Changes to the Files Layout Type to Meet Layout Type Requirements

This section describes the changes that were made from the treatment in [RFC8881] to adequately describe the file layout type within the descriptive framework established in Section 19. These include:

The following changes were needed to provide a layout type description that meets the requirements specified in Section 19.1. Although that section was derived from [RFC8434], it is not precisely the same.

- \* Changes in the treatment of authorization to deal with previous inadequacies in the handling of OPEN.
- \* Changes in the treatment of authorization that ignored possible difficulties caused by unclear elements of the underspecified NFSv4 ACL model.
- \* A previous lack of clarity regarding the possible need for revocation.
- \* A lack of attention to the consequences of server-multipathing (previously alluded to but not recognized as distinct from network-multipathing) and the potential for WRITES to result in conflicts requiring layout revocation.
- \* The uncertainty and likely confusion arising from non-normative statements indicating that there is no need to check IO conformance with the iomode together with the fact that iomode is often a necessary element in determining whether a conflict exists

This gap and the possible ensuing confusion are of major importance when there is a possibility of conflicts arising as a result of server- multipathing.

The needed changes can be summarized as follows:

- \* We try to clearly distinguish server-multipathing from network-multipathing so that we can make clear the issues raised by the former.

This replaces a situation in which the possibility of server-multipathing was mentioned if the two data sets were identical but there was no consideration of the issues that would arise if a WRITE caused them to cease to be identical.

- \* We try to make it clear that there are circumstances that make it necessary to check iomode conformance, even though there might be existing implementations, with certain functionality limits, that work properly despite a lack of such checking
- \* All the material regarding stateid sharing with the data server makes clear the important role of authorization of the opening principal and the lack of need to respond to changes in authorization-related attributes.
- \* We mention the troublesome aspects of the NFSv4 ACL model including the provision of separate bits to control authorization of extending the file or overwriting existing bytes.
- \* We make it clear that, given the uncertainty within older specification of authorization within the NFSv4 protocol, the focus often needs to be on whether the handling on the data server matches that that the MDS would provide, rather than the conformance with previous specification of authorization semantics.
- \* We have adapted the treatment of recovery to incorporate the persistence options made available in [RFC8881] and to a more general approach to the handling of state sharing by the control protocol.

#### 20.27. The File Layout Type and Layout Type Requirements

This section surveys how the requirements for layout type specifications identified in Section 19.1 are addressed, for the files layout type.

The overall framework for understanding the file layout type is provided as follows:

- \* Information regarding the interoperability model (see Section 19.1.1) is presented in Section 20.2.
- \* Information regarding layout-specific data types overlaying the nominally opaque types specified as part of the pNFS feature (See Section 19.1.2) is presented in Section 20.7.
- \* The handling of various OPTIONAL features are discussed later in this section. Currently the, only such troublesome feature is mandatory byte-range locking.

Features added in NFSv4.2 need to be checked for how they might be implemented together with the files layout type.

- \* The storage protocol for the files layout type (a variant of NFSv4.1 or subsequent minor versions is discussed in Section 20.4.

The function of the control protocol whose requirements are specified in Sections 19.1.3 through 19.1.5 are addressed as follows:

- \* Control program functions related to layout management (discussed in Section 19.1.3) are addressed in Sections 20.17, 20.18, and 20.14.1.
- \* The functions of the control protocol related to storage allocation whose requirements are specified Section 19.1.4 are discussed in Section 20.14.4
- \* The functions of the control protocol related to the performance of IO requests directed to the MDS that are specified in Section 20.19.

For important background, see Section 20.14.1.

- \* The functions of the control protocol related to IO request authorization that are specified in Section 19.1.4 are discussed in Section 20.14.2.
- \* The functions of the control protocol related to the maintenance of attribute values that can be affected by WRITES that are specified in Section 19.1.4 are discussed in Section 20.13.
- \* The functions described by Section 19.1.5 are provided as described by Section 20.14.5

The layout specific information regarding the termination of layouts discussed in Section 19.1.9 is provided as follows:

- \* Although the sharing of information regarding the return of layouts can be integrated with the sharing of other state information as described in Section 20.14.1, it can also be provided independently, as described in Section 20.14.3.

In either case, it is important that client, data server and MDS all know that a returned layout no longer exists,' although it is possible that the MDS and client might not be aware of some existing layouts.

- \* Information regarding layout-type-specific issues with regard to layout recall is discussed in Section 20.21.

- \* Information regarding layout-type-specific issues with regard to layout recall is discussed in Section 20.17.
- \* Information regarding layout-type-specific issues with regard to layout recall is discussed in Section 20.18.
- \* Information regard potential layout-type-specific restrictions regarding the discarding of layout-related information are discussed in Section 20.22.

The requirements described by Section 19.1.7 are addressed in Sections 20.23 and 20.24

The requirements described by Section 19.1.6 are addressed, for the files layout type, in Section 20.25. To provide some guidance in satisfying these requirements, the following section are helpful:

- \* Section 20.14.1 provide information about sharing of state necessary to verify the appropriate locking state (and indirectly authorization for the opener) to perform IO.
- \* Section 20.14.2, provides information about IO requests made by principals other than the opener and IOs using a delegation stateid.
- \* Section 20.20 provides information about options that can be use to support the case of mandatory byte range locks.

## 21. Internationalization

Internationalization for NFSv4.1 is described in [I-D.ietf-nfsv4-internationalization], just as it is for other minor versions. The only NFSv4.1-specific element, the `fs_charset_cap` attribute is described in Section 21.1 below.

### 21.1. UTF-8 Capabilities

```
const FSCHARSET_CAP4_CONTAINS_NON_UTF8 = 0x1;
const FSCHARSET_CAP4_ALLOWS_ONLY_UTF8  = 0x2;

typedef uint32_t      fs_charset_cap4;
```

This attribute provides a simple way of determining whether a particular file system behaves as a UTF-8-only server and rejects file names which are not valid Unicode strings encoded using UTF-8. When this attribute is supported and the value returned has the `FSCHARSET_CAP4_ALLOWS_ONLY_UTF8` flag set, the error `NFS4ERR_INVAL` MUST be returned if any file name argument contains a string which is not a valid UTF-8-encoded string.

When this attribute is supported and the value returned has the `FSCHARSET_CAP4_ALLOWS_ONLY_UTF8` flag clear, the error `NFS4ERR_INVAL` will not be returned based on adherence to the rules of UTF-8. While such file systems are generally UTF-8-unaware, this cannot be assumed, since server are allowed (in some circumstances; it is a "SHOULD NOT") to accept non-UTF-8 names while being aware of the structure of UTF-8-conforming names, for the purposes of determining canonical equivalence, for example.

With regard to the flag `FSCHARSET_CAP4_CONTAINS_NON_UTF8`, it has proved impossible to determine, from existing treatments of this attribute, any value that might be helpful here. As a result, we are forced to assume that this flag is always a complement of `FSCHARSET_CAP4_ALLOWS_ONLY_UTF8` and that any result in which it is not is to be ignored, with the appropriate handling being the same as would apply if the attribute were not supported.

When this attribute is not supported, the client can perform a `LOOKUP` using a name not conforming to the rules of UTF-8 and use the error returned to determine whether names which not UTF-8-encoded Unicode are accepted.

## 22. Error Values

NFS error numbers are assigned to failed operations within a Compound (`COMPOUND` or `CB_COMPOUND`) request. A Compound request contains a number of NFS operations that have their results encoded in sequence in a Compound reply. The results of successful operations will consist of an `NFS4_OK` status followed by the encoded results of the operation. If an NFS operation fails, an error status will be entered in the reply and the Compound request will be terminated.

### 22.1. Error Definitions

Error	Number	Description
<code>NFS4_OK</code>	0	Section 22.1.3.1
<code>NFS4ERR_ACCESS</code>	13	Section 22.1.6.1



NFS4ERR_ACLST_POORFIT	20000	REFTBD22
NFS4ERR_ACLST_NOFIT	20001	REFTBD22
NFS4ERR_ACLST_TOOBIG	20002	REFTBD22
NFS4ERR_ACLFT_TOOBIG	20003	REFTBD22
NFS4ERR_ATTRNOTSUPP	10032	Section 22.1.15.1
NFS4ERR_ADMIN_REVOKED	10047	Section 22.1.5.1
NFS4ERR_BACK_CHAN_BUSY	10057	Section 22.1.12.1
NFS4ERR_BADCHAR	10040	Section 22.1.7.1
NFS4ERR_BADHANDLE	10001	Section 22.1.2.1
NFS4ERR_BADIOMODE	10049	Section 22.1.10.1
NFS4ERR_BADLAYOUT	10050	Section 22.1.10.2
NFS4ERR_BADNAME	10041	Section 22.1.7.2
NFS4ERR_BADOWNER	10039	Section 22.1.15.2
NFS4ERR_BADSESSION	10052	Section 22.1.11.1
NFS4ERR_BADSLOT	10053	Section 22.1.11.2
NFS4ERR_BADTYPE	10007	Section 22.1.4.1
NFS4ERR_BADXDR	10036	Section 22.1.1.1
NFS4ERR_BAD_COOKIE	10003	Section 22.1.1.2
NFS4ERR_BAD_HIGH_SLOT	10077	Section 22.1.11.3
NFS4ERR_BAD_RANGE	10042	Section 22.1.8.1
NFS4ERR_BAD_SEQID	10026	Section 22.1.16.1
NFS4ERR_BAD_SESSION_DIGEST	10051	Section 22.1.12.2
NFS4ERR_BAD_STATEID	10025	Section 22.1.5.2
NFS4ERR_CB_PATH_DOWN	10048	Section 22.1.11.4

NFS4ERR_CLID_INUSE	10017	Section 22.1.13.2
NFS4ERR_CLIENTID_BUSY	10074	Section 22.1.13.1
NFS4ERR_COMPLETE_ALREADY	10054	Section 22.1.9.1
NFS4ERR_CONN_NOT_BOUND_TO_SESSION	10055	Section 22.1.11.6
NFS4ERR_DEADLOCK	10045	Section 22.1.8.2
NFS4ERR_DEADSESSION	10078	Section 22.1.11.5
NFS4ERR_DELAY	10008	Section 22.1.1.3
NFS4ERR_DELEG_ALREADY_WANTED	10056	Section 22.1.14.1
NFS4ERR_DELEG_REVOKED	10087	Section 22.1.5.3
NFS4ERR_DENIED	10010	Section 22.1.8.3
NFS4ERR_DIRDELEG_UNAVAIL	10084	Section 22.1.14.2
NFS4ERR_DQUOT	69	Section 22.1.4.2
NFS4ERR_ENCR_ALG_UNSUPP	10079	Section 22.1.13.3
NFS4ERR_EXIST	17	Section 22.1.4.3
NFS4ERR_EXPIRED	10011	Section 22.1.5.4
NFS4ERR_FBIG	27	Section 22.1.4.4
NFS4ERR_FHEXPIRED	10014	Section 22.1.2.2
NFS4ERR_FILE_OPEN	10046	Section 22.1.4.5
NFS4ERR_GRACE	10013	Section 22.1.9.2
NFS4ERR_HASH_ALG_UNSUPP	10072	Section 22.1.13.4
NFS4ERR_INVALID	22	Section 22.1.1.4
NFS4ERR_IO	5	Section 22.1.4.6
NFS4ERR_ISDIR	21	Section 22.1.2.3
NFS4ERR_LAYOUTTRYLATER	10058	Section 22.1.10.3

NFS4ERR_LAYOUTUNAVAILABLE	10059	Section 22.1.10.4
NFS4ERR_LEASE_MOVED	10031	Section 22.1.16.2
NFS4ERR_LOCKED	10012	Section 22.1.8.4
NFS4ERR_LOCKS_HELD	10037	Section 22.1.8.5
NFS4ERR_LOCK_NOTSUPP	10043	Section 22.1.8.6
NFS4ERR_LOCK_RANGE	10028	Section 22.1.8.7
NFS4ERR_MINOR_VERS_MISMATCH	10021	Section 22.1.3.2
NFS4ERR_MLINK	31	Section 22.1.4.7
NFS4ERR_MOVED	10019	Section 22.1.2.4
NFS4ERR_NAMETOOLONG	63	Section 22.1.7.3
NFS4ERR_NOENT	2	Section 22.1.4.8
NFS4ERR_NOFILEHANDLE	10020	Section 22.1.2.5
NFS4ERR_NOMATCHING_LAYOUT	10060	Section 22.1.10.5
NFS4ERR_NOSPC	28	Section 22.1.4.9
NFS4ERR_NOTDIR	20	Section 22.1.2.6
NFS4ERR_NOTEMPTY	66	Section 22.1.4.10
NFS4ERR_NOTSUPP	10004	Section 22.1.1.5
NFS4ERR_NOT_ONLY_OP	10081	Section 22.1.3.3
NFS4ERR_NOT_SAME	10027	Section 22.1.15.3
NFS4ERR_NO_GRACE	10033	Section 22.1.9.3
NFS4ERR_NXIO	6	Section 22.1.16.3
NFS4ERR_OLD_STATEID	10024	Section 22.1.5.5
NFS4ERR_OPENMODE	10038	Section 22.1.8.8
NFS4ERR_OP_ILLEGAL	10044	Section 22.1.3.4

NFS4ERR_OP_NOT_IN_SESSION	10071	Section 22.1.3.5
NFS4ERR_PERM	1	Section 22.1.6.2
NFS4ERR_PNFS_IO_HOLE	10075	Section 22.1.10.6
NFS4ERR_PNFS_NO_LAYOUT	10080	Section 22.1.10.7
NFS4ERR_RECALLCONFLICT	10061	Section 22.1.14.3
NFS4ERR_RECLAIM_BAD	10034	Section 22.1.9.4
NFS4ERR_RECLAIM_CONFLICT	10035	Section 22.1.9.5
NFS4ERR_REJECT_DELEG	10085	Section 22.1.14.4
NFS4ERR_REP_TOO_BIG	10066	Section 22.1.3.6
NFS4ERR_REP_TOO_BIG_TO_CACHE	10067	Section 22.1.3.7
NFS4ERR_REQ_TOO_BIG	10065	Section 22.1.3.8
NFS4ERR_RESOURCE	10018	Section 22.1.16.4
NFS4ERR_RESTOREFH	10030	Section 22.1.16.5
NFS4ERR_RETRY_UNCACHED_REP	10068	Section 22.1.3.9
NFS4ERR_RETURNCONFLICT	10086	Section 22.1.10.8
NFS4ERR_ROFS	30	Section 22.1.4.11
NFS4ERR_SAME	10009	Section 22.1.15.4
NFS4ERR_SHARE_DENIED	10015	Section 22.1.8.9
NFS4ERR_SEQUENCE_POS	10064	Section 22.1.3.10
NFS4ERR_SEQ_FALSE_RETRY	10076	Section 22.1.11.7
NFS4ERR_SEQ_MISORDERED	10063	Section 22.1.11.8
NFS4ERR_SERVERFAULT	10006	Section 22.1.1.6
NFS4ERR_STALE	70	Section 22.1.2.7
NFS4ERR_STALE_CLIENTID	10022	Section 22.1.13.5

NFS4ERR_STALE_STATEID	10023	Section 22.1.16.6
NFS4ERR_SYMLINK	10029	Section 22.1.2.8
NFS4ERR_TOOSMALL	10005	Section 22.1.1.7
NFS4ERR_TOO_MANY_OPS	10070	Section 22.1.3.11
NFS4ERR_UNKNOWN_LAYOUTTYPE	10062	Section 22.1.10.9
NFS4ERR_UNSAFE_COMPOUND	10069	Section 22.1.3.12
NFS4ERR_WRONGSEC	10016	Section 22.1.6.3
NFS4ERR_WRONG_CRED	10082	Section 22.1.6.4
NFS4ERR_WRONG_TYPE	10083	Section 22.1.2.9
NFS4ERR_XDEV	18	Section 22.1.4.12

Table 11: Protocol Error Definitions

## 22.1.1. General Errors

This section deals with errors that are applicable to a broad set of different purposes.

## 22.1.1.1. NFS4ERR\_BADXDR (Error Code 10036)

The arguments for this operation do not match those specified in the XDR definition. This includes situations in which the request ends before all the arguments have been seen. Note that this error applies when fixed enumerations (these include booleans) have a value within the input stream that is not valid for the enum. A replier may pre-parse all operations for a Compound procedure before doing any operation execution and return RPC-level XDR errors in that case.

## 22.1.1.2. NFS4ERR\_BAD\_COOKIE (Error Code 10003)

Used for operations that provide a set of information indexed by some quantity provided by the client or cookie sent by the server for an earlier invocation. Where the value cannot be used for its intended purpose, this error results.

#### 22.1.1.3. NFS4ERR\_DELAY (Error Code 10008)

For any of a number of reasons, the replier could not process this operation in what was deemed a reasonable time. The requester should wait and then try the request with a new slot and sequence value.

Some examples of situations that might lead to this error being returned:

- \* A server that supports hierarchical storage receives a request to process a file that had been migrated.
- \* An operation requires a delegation recall to proceed, but the need to wait for this delegation to be recalled and returned makes processing this request in a timely fashion impossible.
- \* A request is being performed on a session being migrated from another server as described in Section 17.14.3, and the lack of full information about the state of the session on the source makes it impossible to process the request immediately.

In such cases, returning the error NFS4ERR\_DELAY allows necessary preparatory operations to proceed without holding up requester resources such as a session slot. After delaying for period of time, the requester can then re-send the operation in question, often as part of a nearly identical request. Because of the need to avoid spurious reissues of non-idempotent operations and to avoid acting in response to NFS4ERR\_DELAY errors returned on responses returned from the replier's reply cache, integration with the session-provided reply cache is necessary. There are a number of cases to deal with, each of which requires different sorts of handling by the requester and replier:

- \* If NFS4ERR\_DELAY is returned on a SEQUENCE operation, the request is retried in full with the SEQUENCE operation containing the same slot and sequence values. In this case, the replier MUST avoid returning a response containing NFS4ERR\_DELAY as the response to SEQUENCE solely because an earlier instance of the same request returned that error and it was stored in the reply cache. If the replier did this, the retries would not be effective as there would be no opportunity for the replier to see whether the condition that generated the NFS4ERR\_DELAY had been rectified during the time between the original request and the retry.
- \* If NFS4ERR\_DELAY is returned on an operation other than SEQUENCE that validly appears as the first operation of a request, the handling is similar. The request can be retried in full without modification. In this case as well, the replier MUST avoid

returning a response containing NFS4ERR\_DELAY as the response to an initial operation of a request solely on the basis of its presence in the reply cache. If the replier did this, the retries would not be effective as there would be no opportunity for the replier to see whether the condition that generated the NFS4ERR\_DELAY had been rectified during the interim between the original request and the retry.

- \* If NFS4ERR\_DELAY is returned on an operation other than the first in the request, the request when retried MUST contain a SEQUENCE operation that is different than the original one, with either the slot ID or the sequence value different from that in the original request. Because requesters do this, there is no need for the replier to take special care to avoid returning an NFS4ERR\_DELAY error obtained from the reply cache. When no non-idempotent operations have been processed before the NFS4ERR\_DELAY was returned, the requester should retry the request in full, with the only difference from the original request being the modification to the slot ID or sequence value in the reissued SEQUENCE operation.
- \* When NFS4ERR\_DELAY is returned on an operation other than the first within a request and there has been a non-idempotent operation processed before the NFS4ERR\_DELAY was returned, reissuing the request as is normally done would incorrectly cause the re-execution of the non-idempotent operation.

To avoid this situation, the requester should reissue the request without the non-idempotent operation. The request still must use a SEQUENCE operation with either a different slot ID or sequence value from the SEQUENCE in the original request. Because this is done, there is no way the replier could avoid spuriously re-executing the non-idempotent operation since the different SEQUENCE parameters prevent the requester from recognizing that the non-idempotent operation is being retried.

Note that without the ability to return NFS4ERR\_DELAY and the requester's willingness to re-send when receiving it, deadlock might result. For example, if a recall is done, and if the delegation return or operations preparatory to delegation return are held up by other operations that need the delegation to be returned, session slots might not be available. The result could be deadlock.

#### 22.1.1.4. NFS4ERR\_INVALID (Error Code 22)

The arguments for this operation are not valid for some reason, even though they do match those specified in the XDR definition for the request.

#### 22.1.1.5. NFS4ERR\_NOTSUPP (Error Code 10004)

Operation not supported because the operation is either of the following:

- \* an OPTIONAL one and is not supported by this server or the file system on which it is issued.
- \* an operation which MUST NOT be implemented in the current minor version.

In addition, this error may be returned in certain unsupported instances of the LINK operation.

#### 22.1.1.6. NFS4ERR\_SERVERFAULT (Error Code 10006)

An error occurred on the server that does not map to any of the specific legal NFSv4.1 protocol error values. The client should translate this into an appropriate error. UNIX clients may choose to translate this to EIO.

#### 22.1.1.7. NFS4ERR\_TOOSMALL (Error Code 10005)

Used where an operation returns a variable amount of data, with a limit specified by the client. Where the data returned cannot be fit within the limit specified by the client, this error results.

### 22.1.2. Filehandle Errors

These errors deal with the situation in which the current or saved filehandle, or the filehandle passed to PUTFH intended to become the current filehandle, is invalid in some way. This includes situations in which the filehandle is a valid filehandle in general but is not of the appropriate object type for the current operation.

Where the error description indicates a problem with the current or saved filehandle, it is to be understood that filehandles are only checked for the condition if they are implicit arguments of the operation in question.

#### 22.1.2.1. NFS4ERR\_BADHANDLE (Error Code 10001)

Illegal NFS filehandle for the current server. The current filehandle failed internal consistency checks. Once accepted as valid (by PUTFH), no subsequent status change can cause the filehandle to generate this error.



#### 22.1.2.2. NFS4ERR\_FHEXPIRED (Error Code 10014)

A current or saved filehandle that is an argument to the current operation is volatile and has expired at the server.

#### 22.1.2.3. NFS4ERR\_ISDIR (Error Code 21)

The current or saved filehandle designates a directory when the current operation does not allow a directory to be accepted as the target of this operation.

#### 22.1.2.4. NFS4ERR\_MOVED (Error Code 10019)

The file system that contains the current filehandle object is not present at the server or is not accessible with the network address used. It may have been made accessible on a different set of network addresses, relocated or migrated to another server, or it may have never been present. The client may obtain the new file system location by obtaining the `fs_locations` or `fs_locations_info` attribute for the current filehandle. For further discussion, refer to Section 17.3.

As with the case of `NFS4ERR_DELAY`, it is possible that one or more non-idempotent operations may have been successfully executed within a COMPOUND before `NFS4ERR_MOVED` is returned. Because of this, once the new location is determined, the original request that received the `NFS4ERR_MOVED` should not be re-executed in full. Instead, the client should send a new COMPOUND with any successfully executed non-idempotent operations removed. When the client uses the same session for the new COMPOUND, its `SEQUENCE` operation should use a different slot ID or sequence.

#### 22.1.2.5. NFS4ERR\_NOFILEHANDLE (Error Code 10020)

The logical current or saved filehandle value is required by the current operation and is not set. This may be a result of a malformed COMPOUND operation (i.e., no `PUTFH` or `PUTROOTFH` before an operation that requires the current filehandle be set).

#### 22.1.2.6. NFS4ERR\_NOTDIR (Error Code 20)

The current (or saved) filehandle designates an object that is not a directory for an operation in which a directory is required.

#### 22.1.2.7. NFS4ERR\_STALE (Error Code 70)

The current or saved filehandle value designating an argument to the current operation is invalid. The file referred to by that filehandle no longer exists or access to it has been revoked.

#### 22.1.2.8. NFS4ERR\_SYMLINK (Error Code 10029)

The current filehandle designates a symbolic link when the current operation does not allow a symbolic link as the target.

#### 22.1.2.9. NFS4ERR\_WRONG\_TYPE (Error Code 10083)

The current (or saved) filehandle designates an object that is of an invalid type for the current operation, and there is no more specific error (such as NFS4ERR\_ISDIR or NFS4ERR\_SYMLINK) that applies. Note that in NFSv4.0, such situations generally resulted in the less-specific error NFS4ERR\_INVALID.

### 22.1.3. Compound Structure Errors

This section deals with errors that relate to the overall structure of a Compound request (by which we mean to include both COMPOUND and CB\_COMPOUND), rather than to particular operations.

There are a number of basic constraints on the operations that may appear in a Compound request. Sessions add to these basic constraints by requiring a Sequence operation (either SEQUENCE or CB\_SEQUENCE) at the start of the Compound.

#### 22.1.3.1. NFS\_OK (Error code 0)

Indicates the operation completed successfully, in that all of the constituent operations completed without error.

#### 22.1.3.2. NFS4ERR\_MINOR\_VERS\_MISMATCH (Error code 10021)

The minor version specified is not one that the current listener supports. This value is returned in the overall status for the Compound but is not associated with a specific operation since the results will specify a result count of zero.

#### 22.1.3.3. NFS4ERR\_NOT\_ONLY\_OP (Error Code 10081)

Certain operations, which are allowed to be executed outside of a session, MUST be the only operation within a Compound whenever the Compound does not start with a Sequence operation. This error results when that constraint is not met.

#### 22.1.3.4. NFS4ERR\_OP\_ILLEGAL (Error Code 10044)

The operation code is not a valid one for the current Compound procedure. The opcode in the result stream matched with this error is the ILLEGAL value, although the value that appears in the request stream may be different. Where an illegal value appears and the replier pre-parses all operations for a Compound procedure before doing any operation execution, an RPC-level XDR error may be returned.

#### 22.1.3.5. NFS4ERR\_OP\_NOT\_IN\_SESSION (Error Code 10071)

Most forward operations and all callback operations are only valid within the context of a session, so that the Compound request in question MUST begin with a Sequence operation. If an attempt is made to execute these operations outside the context of session, this error results.

#### 22.1.3.6. NFS4ERR\_REP\_TOO\_BIG (Error Code 10066)

The reply to a Compound would exceed the channel's negotiated maximum response size.

#### 22.1.3.7. NFS4ERR\_REP\_TOO\_BIG\_TO\_CACHE (Error Code 10067)

The reply to a Compound would exceed the channel's negotiated maximum size for replies cached in the reply cache when the Sequence for the current request specifies that this request is to be cached.

#### 22.1.3.8. NFS4ERR\_REQ\_TOO\_BIG (Error Code 10065)

The Compound request exceeds the channel's negotiated maximum size for requests.

#### 22.1.3.9. NFS4ERR\_RETRY\_UNCACHED\_REP (Error Code 10068)

The requester has attempted a retry of a Compound that it previously requested not be placed in the reply cache.

#### 22.1.3.10. NFS4ERR\_SEQUENCE\_POS (Error Code 10064)

A Sequence operation appeared in a position other than the first operation of a Compound request.

#### 22.1.3.11. NFS4ERR\_TOO\_MANY\_OPS (Error Code 10070)

The Compound request has too many operations, exceeding the count negotiated when the session was created.

#### 22.1.3.12. NFS4ERR\_UNSAFE\_COMPOUND (Error Code 10068)

The client has sent a COMPOUND request with an unsafe mix of operations -- specifically, with a non-idempotent operation that changes the current filehandle and that is not followed by a GETFH.

#### 22.1.4. File System Errors

These errors describe situations that occurred in the underlying file system implementation rather than in the protocol or any NFSv4.x feature.

##### 22.1.4.1. NFS4ERR\_BADTYPE (Error Code 10007)

An attempt was made to create an object with an inappropriate type specified to CREATE. This may be because the type is undefined, because the type is not supported by the server, or because the type is not intended to be created by CREATE (such as a regular file or named attribute, for which OPEN is used to do the file creation).

##### 22.1.4.2. NFS4ERR\_DQUOT (Error Code 69)

Resource (quota) hard limit exceeded. The user's resource limit on the server has been exceeded.

##### 22.1.4.3. NFS4ERR\_EXIST (Error Code 17)

A file of the specified target name (when creating, renaming, or linking) already exists.

##### 22.1.4.4. NFS4ERR\_FBIG (Error Code 27)

The file is too large. The operation would have caused the file to grow beyond the server's limit.

##### 22.1.4.5. NFS4ERR\_FILE\_OPEN (Error Code 10046)

The operation is not allowed because a file involved in the operation is currently open. Servers may, but are not required to, disallow linking-to, removing, or renaming open files.

##### 22.1.4.6. NFS4ERR\_IO (Error Code 5)

Indicates that an I/O error occurred for which the file system was unable to provide recovery.

#### 22.1.4.7. NFS4ERR\_MLINK (Error Code 31)

The request would have caused the server's limit for the number of hard links a file may have to be exceeded.

#### 22.1.4.8. NFS4ERR\_NOENT (Error Code 2)

Indicates no such file or directory. The file or directory name specified does not exist.

#### 22.1.4.9. NFS4ERR\_NOSPC (Error Code 28)

Indicates there is no space left on the device. The operation would have caused the server's file system to exceed its limit.

#### 22.1.4.10. NFS4ERR\_NOTEMPTY (Error Code 66)

An attempt was made to remove a directory that was not empty.

#### 22.1.4.11. NFS4ERR\_ROFS (Error Code 30)

Indicates a read-only file system. A modifying operation was attempted on a read-only file system.

#### 22.1.4.12. NFS4ERR\_XDEV (Error Code 18)

Indicates an attempt to do an operation, such as linking, that inappropriately crosses a boundary. This may be due to such boundaries as:

- \* that between file systems (where the fsids are different).
- \* that between different named attribute directories or between a named attribute directory and an ordinary directory.
- \* that between byte-ranges of a file system that the file system implementation treats as separate (for example, for space accounting purposes), and where cross-connection between the byte-ranges are not allowed.

#### 22.1.5. State Management Errors

These errors indicate problems with the stateid (or one of the stateids) passed to a given operation. This includes situations in which the stateid is invalid as well as situations in which the stateid is valid but designates locking state that has been revoked. Depending on the operation, the stateid when valid may designate opens, byte-range locks, file or directory delegations, layouts, or

device maps.

#### 22.1.5.1. NFS4ERR\_ADMIN\_REVOKED (Error Code 10047)

A stateid designates locking state of any type that has been revoked due to administrative interaction, possibly while the lease is valid.

#### 22.1.5.2. NFS4ERR\_BAD\_STATEID (Error Code 10026)

A stateid does not properly designate any valid state. See Sections 13.2.4 and 13.2.3 for a discussion of how stateids are validated.

#### 22.1.5.3. NFS4ERR\_DELEG\_REVOKED (Error Code 10087)

A stateid designates recallable locking state of any type (delegation or layout) that has been revoked due to the failure of the client to return the lock when it was recalled.

#### 22.1.5.4. NFS4ERR\_EXPIRED (Error Code 10011)

A stateid designates locking state of any type that has been revoked due to expiration of the client's lease, either immediately upon lease expiration, or following a later request for a conflicting lock.

#### 22.1.5.5. NFS4ERR\_OLD\_STATEID (Error Code 10024)

A stateid with a non-zero seqid value is not the most current seqid for the state.

#### 22.1.6. Security Errors

These are the various permission-related errors in NFSv4.1.

##### 22.1.6.1. NFS4ERR\_ACCESS (Error Code 13)

Indicates permission denied. The caller does not have the correct permission to perform the requested operation. Contrast this with NFS4ERR\_PERM (Section 22.1.6.2), which restricts itself to owner or privileged-user permission failures, and NFS4ERR\_WRONG\_CRED (Section 22.1.6.4), which deals with appropriate permission to delete or modify transient objects based on the credentials of the user that created them.

#### 22.1.6.2. NFS4ERR\_PERM (Error Code 1)

Indicates requester is not the owner. The operation was not allowed because the caller is neither a privileged user (root) nor the owner of the target of the operation.

#### 22.1.6.3. NFS4ERR\_WRONGSEC (Error Code 10016)

Indicates that the security mechanism being used by the client for the operation does not match the server's security policy. The client should change the security mechanism being used and re-send the operation (but not with the same slot ID and sequence ID; one or both MUST be different on the re-send). SECINFO and SECINFO\_NO\_NAME can be used to determine the appropriate mechanism.

#### 22.1.6.4. NFS4ERR\_WRONG\_CRED (Error Code 10082)

An operation that manipulates state was attempted by a principal that was not allowed to modify that piece of state.

#### 22.1.7. Name Errors

Names in NFSv4 are typically UTF-8 strings, although it is possible for servers, when accessing certain file systems, to support other encodings to support internationalization. When the strings are of length zero or are not valid UTF-8 encoding in a file system that only supports UTF-8 encodings, the error NFS4ERR\_INVALID results. Besides this, there are a number of other errors to indicate specific problems with names.

##### 22.1.7.1. NFS4ERR\_BADCHAR (Error Code 10040)

A string contains a character that is not supported by the server in the context in which it being used.

##### 22.1.7.2. NFS4ERR\_BADNAME (Error Code 10041)

A name string in a request consisted of valid characters supported by the server, but the name is not supported by the server as a valid name for the current operation. An example might be creating a file or directory named ".." on a server whose file system uses that name for links to parent directories.

##### 22.1.7.3. NFS4ERR\_NAMETOOLONG (Error Code 63)

Returned when the filename in an operation exceeds the server's implementation limit.

### 22.1.8. Locking Errors

This section deals with errors related to locking, both as to share reservations and byte-range locking. It does not deal with errors specific to the process of reclaiming locks. Those are dealt with in Section 22.1.9.

#### 22.1.8.1. NFS4ERR\_BAD\_RANGE (Error Code 10042)

The byte-range of a LOCK, LOCKT, or LOCKU operation is not allowed by the server. For example, this error results when a server that only supports 32-bit ranges receives a range that cannot be handled by that server. (See Section 25.10.3.)

#### 22.1.8.2. NFS4ERR\_DEADLOCK (Error Code 10045)

The server has been able to determine a byte-range locking deadlock condition for a READW\_LT or WRITEW\_LT LOCK operation.

#### 22.1.8.3. NFS4ERR\_DENIED (Error Code 10010)

An attempt to lock a file is denied. Since this may be a temporary condition, the client is encouraged to re-send the lock request (but not with the same slot ID and sequence ID; one or both MUST be different on the re-send) until the lock is accepted. See Section 14.6 for a discussion of the re-send.

#### 22.1.8.4. NFS4ERR\_LOCKED (Error Code 10012)

A READ or WRITE operation was attempted on a file where there was a conflict between the I/O and an existing lock:

- \* There is a share reservation inconsistent with the I/O being done.
- \* The range to be read or written intersects an existing mandatory byte-range lock.

#### 22.1.8.5. NFS4ERR\_LOCKS\_HELD (Error Code 10037)

An operation was prevented by the unexpected presence of locks.

#### 22.1.8.6. NFS4ERR\_LOCK\_NOTSUPP (Error Code 10043)

A LOCK operation was attempted that would require the upgrade or downgrade of a byte-range lock range already held by the owner, and the server does not support atomic upgrade or downgrade of locks.



#### 22.1.8.7. NFS4ERR\_LOCK\_RANGE (Error Code 10028)

A LOCK operation is operating on a range that overlaps in part a currently held byte-range lock for the current lock-owner and does not precisely match a single such byte-range lock where the server does not support this type of request, and thus does not implement POSIX locking semantics [fcntl]. See Sections 25.10.4, 25.11.4, and 25.12.4 for a discussion of how this applies to LOCK, LOCKT, and LOCKU respectively.

#### 22.1.8.8. NFS4ERR\_OPENMODE (Error Code 10038)

The client attempted a READ, WRITE, LOCK, or other operation not sanctioned by the stateid passed (e.g., writing to a file opened for read-only access).

#### 22.1.8.9. NFS4ERR\_SHARE\_DENIED (Error Code 10015)

An attempt to OPEN a file with a share reservation has failed because of a share conflict.

#### 22.1.9. Reclaim Errors

These errors relate to the process of reclaiming locks after a server restart.

##### 22.1.9.1. NFS4ERR\_COMPLETE\_ALREADY (Error Code 10054)

The client previously sent a successful RECLAIM\_COMPLETE operation specifying the same scope, whether that scope is global or for the same file system in the case of a per-fs RECLAIM\_COMPLETE. An additional RECLAIM\_COMPLETE operation is not necessary and results in this error.

##### 22.1.9.2. NFS4ERR\_GRACE (Error Code 10013)

This error is returned when the server is in its grace period with regard to the file system object for which the lock was requested. In this situation, a non-reclaim locking request cannot be granted. This can occur because either:

- \* The server does not have sufficient information about locks that might be potentially reclaimed to determine whether the lock could be granted.
- \* The request is made by a client responsible for reclaiming its locks that has not yet done the appropriate RECLAIM\_COMPLETE operation, allowing it to proceed to obtain new locks.

In the case of a per-fs grace period, there may be clients (i.e., those currently using the destination file system) who might be unaware of the circumstances resulting in the initiation of the grace period. Such clients need to periodically retry the request until the grace period is over, just as other clients do.

#### 22.1.9.3. NFS4ERR\_NO\_GRACE (Error Code 10033)

A reclaim of client state was attempted in circumstances in which the server cannot guarantee that conflicting state has not been provided to another client. This occurs in any of the following situations:

- \* There is no active grace period applying to the file system object for which the request was made.
- \* The client making the request has no current role in reclaiming locks.
- \* Previous operations have created a situation in which the server is not able to determine that a reclaim-interfering edge condition does not exist.

#### 22.1.9.4. NFS4ERR\_RECLAIM\_BAD (Error Code 10034)

The server has determined that a reclaim attempted by the client is not valid, i.e., the lock specified as being reclaimed could not possibly have existed before the server restart or file system migration event. A server is not obliged to make this determination and will typically rely on the client to only reclaim locks that the client was granted prior to restart. However, when a server does have reliable information to enable it to make this determination, this error indicates that the reclaim has been rejected as invalid. This is as opposed to the error NFS4ERR\_RECLAIM\_CONFLICT (See Section 22.1.9.5) where the server can only determine that there has been an invalid reclaim, but cannot determine which request is invalid.

#### 22.1.9.5. NFS4ERR\_RECLAIM\_CONFLICT (Error Code 10035)

The reclaim attempted by the client has encountered a conflict and cannot be satisfied. This potentially indicates a misbehaving client, although not necessarily the one receiving the error. The misbehavior might be on the part of the client that established the lock with which this client conflicted. See also Section 22.1.9.4 for the related error, NFS4ERR\_RECLAIM\_BAD.

#### 22.1.10. pNFS Errors

This section deals with pNFS-related errors including those that are associated with using NFSv4.1 to communicate with a data server.

##### 22.1.10.1. NFS4ERR\_BADIOMODE (Error Code 10049)

An invalid or inappropriate layout iomode was specified. For example an inappropriate layout iomode, suppose a client's LAYOUTGT operation specified an iomode of LAYOUTIOMODE4\_RW, and the server is neither able nor willing to let the client send write requests to data servers; the server can reply with NFS4ERR\_BADIOMODE. The client would then send another LAYOUTGET with an iomode of LAYOUTIOMODE4\_READ.

##### 22.1.10.2. NFS4ERR\_BADLAYOUT (Error Code 10050)

The layout specified is invalid in some way. For LAYOUTCOMMIT, this indicates that the specified layout is not held by the client or is not of mode LAYOUTIOMODE4\_RW. For LAYOUTGET, it indicates that a layout matching the client's specification as to minimum length cannot be granted.

##### 22.1.10.3. NFS4ERR\_LAYOUTTRYLATER (Error Code 10058)

Layouts are temporarily unavailable for the file. The client should re-send later (but not with the same slot ID and sequence ID; one or both MUST be different on the re-send).

##### 22.1.10.4. NFS4ERR\_LAYOUTUNAVAILABLE (Error Code 10059)

Returned when layouts are not available for the current file system or the particular specified file.

##### 22.1.10.5. NFS4ERR\_NOMATCHING\_LAYOUT (Error Code 10060)

Returned when layouts are recalled and the client has no layouts matching the specification of the layouts being recalled.

##### 22.1.10.6. NFS4ERR\_PNFS\_IO\_HOLE (Error Code 10075)

The pNFS client has attempted to read from or write to an illegal hole of a file of a data server that is using sparse packing. See Section 20.8.4.

#### 22.1.10.7. NFS4ERR\_PNFS\_NO\_LAYOUT (Error Code 10080)

The pNFS client has attempted to read from or write to a file (using a request to a data server) without holding a valid layout. This includes the case where the client had a layout, but the iomode does not allow a WRITE.

#### 22.1.10.8. NFS4ERR\_RETURNCONFLICT (Error Code 10086)

A layout is unavailable due to an attempt to perform the LAYOUTGET before a pending LAYOUTRETURN on the file has been received. See Section 18.7.5.3.3.

#### 22.1.10.9. NFS4ERR\_UNKNOWN\_LAYOUTTYPE (Error Code 10062)

The client has specified a layout type that is not supported by the server.

#### 22.1.11. Session Use Errors

This section deals with errors encountered when using sessions, that is, errors encountered when a request uses a Sequence (i.e., either SEQUENCE or CB\_SEQUENCE) operation.

##### 22.1.11.1. NFS4ERR\_BADSESSION (Error Code 10052)

The specified session ID is unknown to the server to which the operation is addressed.

##### 22.1.11.2. NFS4ERR\_BADSLOT (Error Code 10053)

The requester sent a Sequence operation that attempted to use a slot the replier does not have in its slot table. It is possible the slot may have been retired.

##### 22.1.11.3. NFS4ERR\_BAD\_HIGH\_SLOT (Error Code 10077)

The highest\_slot argument in a Sequence operation exceeds the replier's enforced highest\_slotid. Also returned when the rsa\_target\_highest\_slotid argument in a CB\_RECALL\_SLOT operation exceeds maximum enforced slot ID of the session's fore channel.

##### 22.1.11.4. NFS4ERR\_CB\_PATH\_DOWN (Error Code 10048)

There is a problem contacting the client via the callback path. The function of this error has been mostly superseded by the use of status flags in the reply to the SEQUENCE operation (See Section 25.46).

#### 22.1.11.5. NFS4ERR\_DEADSESSION (Error Code 10078)

The specified session is a persistent session that is dead and does not accept new requests or perform new operations on existing requests (in the case in which a request was partially executed before server restart).

#### 22.1.11.6. NFS4ERR\_CONN\_NOT\_BOUND\_TO\_SESSION (Error Code 10055)

A Sequence operation was sent on a connection that has not been associated with the specified session, where the client specified that connection association was to be enforced with SP4\_MACH\_CRED or SP4\_SSV state protection.

#### 22.1.11.7. NFS4ERR\_SEQ\_FALSE\_RETRY (Error Code 10076)

The requester sent a Sequence operation with a slot ID and sequence ID that are in the reply cache, but the replier has detected that the retried request is not the same as the original request. See Section 7.6.1.3.1.

#### 22.1.11.8. NFS4ERR\_SEQ\_MISORDERED (Error Code 10063)

The requester sent a Sequence operation with an invalid sequence ID.

#### 22.1.12. Session Management Errors

This section deals with errors associated with requests used in session management.

##### 22.1.12.1. NFS4ERR\_BACK\_CHAN\_BUSY (Error Code 10057)

An attempt was made to destroy a session when the session cannot be destroyed because the server has callback requests outstanding.

##### 22.1.12.2. NFS4ERR\_BAD\_SESSION\_DIGEST (Error Code 10051)

The digest used in a SET\_SSV request is not valid.

#### 22.1.13. Client Management Errors

This section deals with errors associated with requests used to create and manage client IDs.

##### 22.1.13.1. NFS4ERR\_CLIENTID\_BUSY (Error Code 10074)

The DESTROY\_CLIENTID operation has found there are sessions and/or unexpired state associated with the client ID to be destroyed.

#### 22.1.13.2. NFS4ERR\_CLID\_INUSE (Error Code 10017)

While processing an EXCHANGE\_ID operation, the server was presented with a co\_ownerid field that matches an existing client with valid leased state, but the principal sending the EXCHANGE\_ID operation differs from the principal that established the existing client. This indicates a collision (most likely due to chance) between clients. The client should recover by changing the co\_ownerid and re-sending EXCHANGE\_ID (but not with the same slot ID and sequence ID; one or both MUST be different on the re-send).

#### 22.1.13.3. NFS4ERR\_ENCR\_ALG\_UNSUPP (Error Code 10079)

An EXCHANGE\_ID was sent that specified state protection via SSV, and where the set of encryption algorithms presented by the client did not include any supported by the server.

#### 22.1.13.4. NFS4ERR\_HASH\_ALG\_UNSUPP (Error Code 10072)

An EXCHANGE\_ID was sent that specified state protection via SSV, and where the set of hashing algorithms presented by the client did not include any supported by the server.

#### 22.1.13.5. NFS4ERR\_STALE\_CLIENTID (Error Code 10022)

A client ID not recognized by the server was passed to an operation. Note that unlike the case of NFSv4.0, client IDs are not passed explicitly to the server in ordinary locking operations and cannot result in this error. Instead, when there is a server restart, it is first manifested through an error on the associated session, and the staleness of the client ID is detected when trying to associate a client ID with a new session.

#### 22.1.14. Delegation Errors

This section deals with errors associated with requesting and returning delegations.

##### 22.1.14.1. NFS4ERR\_DELEG\_ALREADY\_WANTED (Error Code 10056)

The client has requested a delegation when it had already registered that it wants that same delegation.

##### 22.1.14.2. NFS4ERR\_DIRDELEG\_UNAVAIL (Error Code 10084)

This error is returned when the server is unable or unwilling to provide a requested directory delegation.

#### 22.1.14.3. NFS4ERR\_RECALLCONFLICT (Error Code 10061)

A recallable object (i.e., a layout or delegation) is unavailable due to a conflicting recall operation that is currently in progress for that object.

#### 22.1.14.4. NFS4ERR\_REJECT\_DELEG (Error Code 10085)

The callback operation invoked to deal with a new delegation has rejected it.

#### 22.1.15. Attribute Handling Errors

This section deals with errors specific to attribute handling within NFSv4.

##### 22.1.15.1. NFS4ERR\_ATTRNOTSUPP (Error Code 10032)

An attribute specified is not supported by the server. This error MUST NOT be returned by the GETATTR operation.

##### 22.1.15.2. NFS4ERR\_BADOWNER (Error Code 10039)

This error is returned when an owner or owner\_group attribute value or the who field of an ACE within an ACL attribute value cannot be translated to a local representation.

##### 22.1.15.3. NFS4ERR\_NOT\_SAME (Error Code 10027)

This error is returned by the VERIFY operation to signify that the attributes compared were not the same as those provided in the client's request.

##### 22.1.15.4. NFS4ERR\_SAME (Error Code 10009)

This error is returned by the NVERIFY operation to signify that the attributes compared were the same as those provided in the client's request.

#### 22.1.16. Obsoleted Errors

These errors MUST NOT be generated by any NFSv4.1 operation. This can be for a number of reasons.

- \* The function provided by the error has been superseded by one of the status bits returned by the SEQUENCE operation.

- \* The new session structure and associated change in locking have made the error unnecessary.
- \* There has been a restructuring of some errors for NFSv4.1 that resulted in the elimination of certain errors.

#### 22.1.16.1. NFS4ERR\_BAD\_SEQID (Error Code 10026)

The sequence number (seqid) in a locking request is neither the next expected number or the last number processed. These seqids are ignored in NFSv4.1.

#### 22.1.16.2. NFS4ERR\_LEASE\_MOVED (Error Code 10031)

A lease being renewed is associated with a file system that has been migrated to a new server. The error has been superseded by the SEQ4\_STATUS\_LEASE\_MOVED status bit (See Section 25.46).

#### 22.1.16.3. NFS4ERR\_NXIO (Error Code 6)

I/O error. No such device or address. This error is for errors involving block and character device access, but because NFSv4.1 is not a device-access protocol, this error is not applicable.

#### 22.1.16.4. NFS4ERR\_RESOURCE (Error Code 10018)

For the processing of the COMPOUND procedure, the server may exhaust available resources and cannot continue processing operations within the COMPOUND procedure. This error will be returned from the server in those instances of resource exhaustion related to the processing of the COMPOUND procedure.

In NFSv4.1, the need for this general error has been eliminated because explicit limits on compound sizes are established when the session is created.

#### 22.1.16.5. NFS4ERR\_RESTOREFH (Error Code 10030)

The RESTOREFH operation does not have a saved filehandle (identified by SAVEFH) to operate upon. In NFSv4.1, this error has been superseded by NFS4ERR\_NOFILEHANDLE.



## 22.1.16.6. NFS4ERR\_STALE\_STATEID (Error Code 10023)

A stateid generated by an earlier server instance was used. This error is moot in NFSv4.1 because all operations that take a stateid MUST be preceded by the SEQUENCE operation, and the earlier server instance is detected by the session infrastructure that supports SEQUENCE.

## 22.2. Operations and Their Valid Errors

This section contains a table that gives the valid error returns for each protocol operation. The error code NFS4\_OK (indicating no error) is not listed but should be understood to be returnable by all operations with two important exceptions:

- \* The operations that MUST NOT be implemented: OPEN\_CONFIRM, RELEASE\_LOCKOWNER, RENEW, SETCLIENTID, and SETCLIENTID\_CONFIRM.
- \* All illegal (i.e. undefined) operations.

Operation	Errors
Illegal Ops	NFS4ERR_BADXDR, NFS4ERR_OP_ILLEGAL
ACCESS	NFS4ERR_ACCESS, NFS4ERR_BADXDR, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_FHEXPIRED, NFS4ERR_INVAL, NFS4ERR_IO, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOO_MANY_OPS
BACKCHANNEL_CTL	NFS4ERR_BADXDR, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_INVAL, NFS4ERR_NOENT, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_TOO_MANY_OPS

BIND_CONN_TO_SESSION	NFS4ERR_BADSESSION, NFS4ERR_BADXDR, NFS4ERR_BAD_SESSION_DIGEST, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_INVAL, NFS4ERR_NOT_ONLY_OP, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_TOO_MANY_OPS
CLOSE	NFS4ERR_ADMIN_REVOKED, NFS4ERR_BADXDR, NFS4ERR_BAD_STATEID, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_EXPIRED, NFS4ERR_FHEXPIRED, NFS4ERR_LOCKS_HELD, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_OLD_STATEID, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONG_CRED
COMMIT	NFS4ERR_ACCESS, NFS4ERR_BADXDR, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_FHEXPIRED, NFS4ERR_IO, NFS4ERR_ISDIR, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_SYMLINK, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONG_TYPE
CREATE	NFS4ERR_ACCESS, NFS4ERR_ATTRNOTSUPP, NFS4ERR_BADCHAR, NFS4ERR_BADNAME, NFS4ERR_BADOWNER, NFS4ERR_BADTYPE, NFS4ERR_BADXDR, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_DQUOT, NFS4ERR_EXIST, NFS4ERR_FHEXPIRED, NFS4ERR_INVAL, NFS4ERR_IO,

	NFS4ERR_MLINK, NFS4ERR_MOVED, NFS4ERR_NAMETOOLONG, NFS4ERR_NOFILEHANDLE, NFS4ERR_NOSPC, NFS4ERR_NOTDIR, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_PERM, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_ROFS, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOO_MANY_OPS, NFS4ERR_UNSAFE_COMPOUND
CREATE_SESSION	NFS4ERR_BADXDR, NFS4ERR_CLID_INUSE, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_INVAL, NFS4ERR_NOENT, NFS4ERR_NOT_ONLY_OP, NFS4ERR_NOSPC, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SEQ_MISORDERED, NFS4ERR_SERVERFAULT, NFS4ERR_STALE_CLIENTID, NFS4ERR_TOOSMALL, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONG_CRED
DELEGPURGE	NFS4ERR_BADXDR, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_NOTSUPP, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONG_CRED
DELEGRETURN	NFS4ERR_ADMIN_REVOKED, NFS4ERR_BADXDR, NFS4ERR_BAD_STATEID, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_DELEG_REVOKED, NFS4ERR_EXPIRED, NFS4ERR_FHEXPIRED, NFS4ERR_INVAL, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_NOTSUPP, NFS4ERR_OLD_STATEID, NFS4ERR_OP_NOT_IN_SESSION,

	NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONG_CRED
DESTROY_CLIENTID	NFS4ERR_BADXDR, NFS4ERR_CLIENTID_BUSY, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_NOT_ONLY_OP, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE_CLIENTID, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONG_CRED
DESTROY_SESSION	NFS4ERR_BACK_CHAN_BUSY, NFS4ERR_BADSESSION, NFS4ERR_BADXDR, NFS4ERR_CB_PATH_DOWN, NFS4ERR_CONN_NOT_BOUND_TO_SESSION, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_NOT_ONLY_OP, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE_CLIENTID, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONG_CRED
EXCHANGE_ID	NFS4ERR_BADCHAR, NFS4ERR_BADXDR, NFS4ERR_CLID_INUSE, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_ENCR_ALG_UNSUPP, NFS4ERR_HASH_ALG_UNSUPP, NFS4ERR_INVAL, NFS4ERR_NOENT, NFS4ERR_NOT_ONLY_OP, NFS4ERR_NOT_SAME, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_TOO_MANY_OPS

FREE_STATEID	NFS4ERR_BADXDR, NFS4ERR_BAD_STATEID, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_LOCKS_HELD, NFS4ERR_OLD_STATEID, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONG_CRED
GET_DIR_DELEGATION	NFS4ERR_ACCESS, NFS4ERR_BADXDR, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_DIRDELEG_UNAVAIL, NFS4ERR_FHEXPIRED, NFS4ERR_GRACE, NFS4ERR_INVAL, NFS4ERR_IO, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_NOTDIR, NFS4ERR_NOTSUPP, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOO_MANY_OPS
GETATTR	NFS4ERR_ACCESS, NFS4ERR_ACLFT_TOOBIG, NFS4ERR_BADXDR, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_FHEXPIRED, NFS4ERR_GRACE, NFS4ERR_INVAL, NFS4ERR_IO, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONG_TYPE
GETDEVICEINFO	NFS4ERR_BADXDR, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_INVAL, NFS4ERR_NOENT, NFS4ERR_NOTSUPP, NFS4ERR_OP_NOT_IN_SESSION,

	NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_TOOSMALL, NFS4ERR_TOO_MANY_OPS, NFS4ERR_UNKNOWN_LAYOUTTYPE
GETDEVICELIST	NFS4ERR_BADXDR, NFS4ERR_BAD_COOKIE, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_FHEXPIRED, NFS4ERR_INVAL, NFS4ERR_IO, NFS4ERR_NOFILEHANDLE, NFS4ERR_NOTSUPP, NFS4ERR_NOT_SAME, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_TOO_MANY_OPS, NFS4ERR_UNKNOWN_LAYOUTTYPE
GETFH	NFS4ERR_FHEXPIRED, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_STALE
LAYOUTCOMMIT	NFS4ERR_ACCESS, NFS4ERR_ADMIN_REVOKED, NFS4ERR_ATTRNOTSUPP, NFS4ERR_BADIOMODE, NFS4ERR_BADLAYOUT, NFS4ERR_BADXDR, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_DELEG_REVOKED, NFS4ERR_EXPIRED, NFS4ERR_FBIG, NFS4ERR_FHEXPIRED, NFS4ERR_GRACE, NFS4ERR_INVAL, NFS4ERR_IO, NFS4ERR_ISDIR NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_NOTSUPP, NFS4ERR_NO_GRACE, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_RECLAIM_BAD, NFS4ERR_RECLAIM_CONFLICT, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_SYMLINK, NFS4ERR_TOO_MANY_OPS, NFS4ERR_UNKNOWN_LAYOUTTYPE,

	NFS4ERR_WRONG_CRED
LAYOUTGET	NFS4ERR_ACCESS, NFS4ERR_ADMIN_REVOKED, NFS4ERR_BADIOMODE, NFS4ERR_BADLAYOUT, NFS4ERR_BADXDR, NFS4ERR_BAD_STATEID, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_DELEG_REVOKED, NFS4ERR_DQUOT, NFS4ERR_EXPIRED, NFS4ERR_FHEXPIRED, NFS4ERR_GRACE, NFS4ERR_INVALID, NFS4ERR_IO, NFS4ERR_LAYOUTTRYLATER, NFS4ERR_LAYOUTUNAVAILABLE, NFS4ERR_LOCKED, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_NOSPC, NFS4ERR_NOTSUPP, NFS4ERR_OLD_STATEID, NFS4ERR_OPENMODE, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_RECALLCONFLICT, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOOSMALL, NFS4ERR_TOO_MANY_OPS, NFS4ERR_UNKNOWN_LAYOUTTYPE, NFS4ERR_WRONG_TYPE
LAYOUTRETURN	NFS4ERR_ADMIN_REVOKED, NFS4ERR_BADXDR, NFS4ERR_BAD_STATEID, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_DELEG_REVOKED, NFS4ERR_EXPIRED, NFS4ERR_FHEXPIRED, NFS4ERR_GRACE, NFS4ERR_INVALID, NFS4ERR_ISDIR, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_NOTSUPP, NFS4ERR_NO_GRACE, NFS4ERR_OLD_STATEID, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOO_MANY_OPS, NFS4ERR_UNKNOWN_LAYOUTTYPE, NFS4ERR_WRONG_CRED, NFS4ERR_WRONG_TYPE
LINK	NFS4ERR_ACCESS, NFS4ERR_BADCHAR, NFS4ERR_BADNAME, NFS4ERR_BADXDR,

	NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_DQUOT, NFS4ERR_EXIST, NFS4ERR_FHEXPIRED, NFS4ERR_FILE_OPEN, NFS4ERR_GRACE, NFS4ERR_INVAL, NFS4ERR_ISDIR, NFS4ERR_IO, NFS4ERR_MLINK, NFS4ERR_MOVED, NFS4ERR_NAMETOOLONG, NFS4ERR_NOFILEHANDLE, NFS4ERR_NOSPC, NFS4ERR_NOTDIR, NFS4ERR_NOTSUPP, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_ROFS, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_SYMLINK, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONGSEC, NFS4ERR_WRONG_TYPE, NFS4ERR_XDEV
LOCK	NFS4ERR_ACCESS, NFS4ERR_ADMIN_REVOKED, NFS4ERR_BADXDR, NFS4ERR_BAD_RANGE, NFS4ERR_BAD_STATEID, NFS4ERR_DEADLOCK, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_DENIED, NFS4ERR_EXPIRED, NFS4ERR_FHEXPIRED, NFS4ERR_GRACE, NFS4ERR_INVAL, NFS4ERR_ISDIR, NFS4ERR_LOCK_NOTSUPP, NFS4ERR_LOCK_RANGE, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_NO_GRACE, NFS4ERR_OLD_STATEID, NFS4ERR_OPENMODE, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_RECLAIM_BAD, NFS4ERR_RECLAIM_CONFLICT, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_ROFS, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_SYMLINK, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONG_CRED, NFS4ERR_WRONG_TYPE
LOCKT	NFS4ERR_ACCESS, NFS4ERR_BADXDR, NFS4ERR_BAD_RANGE, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_DENIED, NFS4ERR_FHEXPIRED,



	NFS4ERR_GRACE, NFS4ERR_INVAL, NFS4ERR_ISDIR, NFS4ERR_LOCK_RANGE, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_ROFS, NFS4ERR_STALE, NFS4ERR_SYMLINK, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONG_CRED, NFS4ERR_WRONG_TYPE
LOCKU	NFS4ERR_ACCESS, NFS4ERR_ADMIN_REVOKED, NFS4ERR_BADXDR, NFS4ERR_BAD_RANGE, NFS4ERR_BAD_STATEID, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_EXPIRED, NFS4ERR_FHEXPIRED, NFS4ERR_INVAL, NFS4ERR_LOCK_RANGE, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_OLD_STATEID, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONG_CRED
LOOKUP	NFS4ERR_ACCESS, NFS4ERR_BADCHAR, NFS4ERR_BADNAME, NFS4ERR_BADXDR, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_FHEXPIRED, NFS4ERR_INVAL, NFS4ERR_IO, NFS4ERR_MOVED, NFS4ERR_NAMETOOLONG, NFS4ERR_NOENT, NFS4ERR_NOFILEHANDLE, NFS4ERR_NOTDIR, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_SYMLINK, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONGSEC
LOOKUPP	NFS4ERR_ACCESS, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_FHEXPIRED, NFS4ERR_IO, NFS4ERR_MOVED,

	NFS4ERR_NOENT, NFS4ERR_NOFILEHANDLE, NFS4ERR_NOTDIR, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_SYMLINK, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONGSEC
NVERIFY	NFS4ERR_ACCESS, NFS4ERR_ATTRNOTSUPP, NFS4ERR_BADCHAR, NFS4ERR_BADXDR, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_FHEXPIRED, NFS4ERR_GRACE, NFS4ERR_INVAL, NFS4ERR_IO, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SAME, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOO_MANY_OPS, NFS4ERR_UNKNOWN_LAYOUTTYPE, NFS4ERR_WRONG_TYPE
OPEN	NFS4ERR_ACCESS, NFS4ERR_ADMIN_REVOKED, NFS4ERR_ATTRNOTSUPP, NFS4ERR_BADCHAR, NFS4ERR_BADNAME, NFS4ERR_BADOWNER, NFS4ERR_BADXDR, NFS4ERR_BAD_STATEID, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_DELEG_ALREADY_WANTED, NFS4ERR_DELEG_REVOKED, NFS4ERR_DQUOT, NFS4ERR_EXIST, NFS4ERR_EXPIRED, NFS4ERR_FBIG, NFS4ERR_FHEXPIRED, NFS4ERR_GRACE, NFS4ERR_INVAL, NFS4ERR_ISDIR, NFS4ERR_IO, NFS4ERR_MOVED, NFS4ERR_NAME_TOO_LONG, NFS4ERR_NOENT, NFS4ERR_NOFILEHANDLE, NFS4ERR_NOSPC, NFS4ERR_NOTDIR, NFS4ERR_NO_GRACE, NFS4ERR_OLD_STATEID, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_PERM, NFS4ERR_RECLAIM_BAD, NFS4ERR_RECLAIM_CONFLICT, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG,

	NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_ROFS, NFS4ERR_SERVERFAULT, NFS4ERR_SHARE_DENIED, NFS4ERR_STALE, NFS4ERR_SYMLINK, NFS4ERR_TOO_MANY_OPS, NFS4ERR_UNSAFE_COMPOUND, NFS4ERR_WRONGSEC, NFS4ERR_WRONG_TYPE
OPEN_CONFIRM	NFS4ERR_NOTSUPP
OPEN_DOWNGRADE	NFS4ERR_ADMIN_REVOKED, NFS4ERR_BADXDR, NFS4ERR_BAD_STATEID, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_EXPIRED, NFS4ERR_FHEXPIRED, NFS4ERR_INVAL, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_OLD_STATEID, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_ROFS, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONG_CRED
OPENATTR	NFS4ERR_ACCESS, NFS4ERR_BADXDR, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_DQUOT, NFS4ERR_FHEXPIRED, NFS4ERR_IO, NFS4ERR_MOVED, NFS4ERR_NOENT, NFS4ERR_NOFILEHANDLE, NFS4ERR_NOSPC, NFS4ERR_NOTSUPP, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_ROFS, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOO_MANY_OPS, NFS4ERR_UNSAFE_COMPOUND, NFS4ERR_WRONG_TYPE
PUTFH	NFS4ERR_BADHANDLE, NFS4ERR_BADXDR, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_MOVED, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG,

	NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONGSEC
PUTPUBFH	NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONGSEC
PUTROOTFH	NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONGSEC
READ	NFS4ERR_ACCESS, NFS4ERR_ADMIN_REVOKED, NFS4ERR_BADXDR, NFS4ERR_BAD_STATEID, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_DELEG_REVOKED, NFS4ERR_EXPIRED, NFS4ERR_FHEXPIRED, NFS4ERR_GRACE, NFS4ERR_INVAL, NFS4ERR_ISDIR, NFS4ERR_IO, NFS4ERR_LOCKED, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_OLD_STATEID, NFS4ERR_OPENMODE, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_PNFS_IO_HOLE, NFS4ERR_PNFS_NO_LAYOUT, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_SYMLINK, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONG_TYPE
READDIR	NFS4ERR_ACCESS, NFS4ERR_BADXDR, NFS4ERR_BAD_COOKIE, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_FHEXPIRED, NFS4ERR_INVAL, NFS4ERR_IO, NFS4ERR_MOVED,

	NFS4ERR_NOFILEHANDLE, NFS4ERR_NOTDIR, NFS4ERR_NOT_SAME, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOOSMALL, NFS4ERR_TOO_MANY_OPS
READLINK	NFS4ERR_ACCESS, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_FHEXPIRED, NFS4ERR_INVAL, NFS4ERR_IO, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONG_TYPE
RECLAIM_COMPLETE	NFS4ERR_BADXDR, NFS4ERR_COMPLETE_ALREADY, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_FHEXPIRED, NFS4ERR_INVAL, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONG_CRED, NFS4ERR_WRONG_TYPE
RELEASE_LOCKOWNER	NFS4ERR_NOTSUPP
REMOVE	NFS4ERR_ACCESS, NFS4ERR_BADCHAR, NFS4ERR_BADNAME, NFS4ERR_BADXDR, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_FHEXPIRED, NFS4ERR_FILE_OPEN, NFS4ERR_GRACE, NFS4ERR_INVAL, NFS4ERR_IO, NFS4ERR_MOVED, NFS4ERR_NAMETOOLONG, NFS4ERR_NOENT, NFS4ERR_NOFILEHANDLE, NFS4ERR_NOTDIR, NFS4ERR_NOTEMPTY,

	NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_ROFS, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOO_MANY_OPS
RENAME	NFS4ERR_ACCESS, NFS4ERR_BADCHAR, NFS4ERR_BADNAME, NFS4ERR_BADXDR, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_DQUOT, NFS4ERR_EXIST, NFS4ERR_FHEXPIRED, NFS4ERR_FILE_OPEN, NFS4ERR_GRACE, NFS4ERR_INVAL, NFS4ERR_IO, NFS4ERR_MLINK, NFS4ERR_MOVED, NFS4ERR_NAMETOOLONG, NFS4ERR_NOENT, NFS4ERR_NOFILEHANDLE, NFS4ERR_NOSPC, NFS4ERR_NOTDIR, NFS4ERR_NOTEMPTY, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_ROFS, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONGSEC, NFS4ERR_XDEV
RENEW	NFS4ERR_NOTSUPP
RESTOREFH	NFS4ERR_DEADSESSION, NFS4ERR_FHEXPIRED, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONGSEC
SAVEFH	NFS4ERR_DEADSESSION, NFS4ERR_FHEXPIRED, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG,

	NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOO_MANY_OPS
SECINFO	NFS4ERR_ACCESS, NFS4ERR_BADCHAR, NFS4ERR_BADNAME, NFS4ERR_BADXDR, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_FHEXPIRED, NFS4ERR_INVAL, NFS4ERR_MOVED, NFS4ERR_NAMETOOLONG, NFS4ERR_NOENT, NFS4ERR_NOFILEHANDLE, NFS4ERR_NOTDIR, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOO_MANY_OPS
SECINFO_NO_NAME	NFS4ERR_ACCESS, NFS4ERR_BADXDR, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_FHEXPIRED, NFS4ERR_INVAL, NFS4ERR_MOVED, NFS4ERR_NOENT, NFS4ERR_NOFILEHANDLE, NFS4ERR_NOTDIR, NFS4ERR_NOTSUPP, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOO_MANY_OPS
SEQUENCE	NFS4ERR_BADSESSION, NFS4ERR_BADSLLOT, NFS4ERR_BADXDR, NFS4ERR_BAD_HIGH_SLOT, NFS4ERR_CONN_NOT_BOUND_TO_SESSION, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SEQUENCE_POS, NFS4ERR_SEQ_FALSE_RETRY, NFS4ERR_SEQ_MISORDERED, NFS4ERR_TOO_MANY_OPS
SET_SSV	NFS4ERR_BADXDR, NFS4ERR_BAD_SESSION_DIGEST,

	NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_INVAL, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_TOO_MANY_OPS
SETATTR	NFS4ERR_ACCESS, NFS4ERR_ACLST_POORFIT, NFS4ERR_ACLST_NOFIT, NFS4ERR_ACLST_TOOBIG, NFS4ERR_ADMIN_REVOKED, NFS4ERR_ATTRNOTSUPP, NFS4ERR_BADCHAR, NFS4ERR_BADOWNER, NFS4ERR_BADXDR, NFS4ERR_BAD_STATEID, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_DELEG_REVOKED, NFS4ERR_DQUOT, NFS4ERR_EXPIRED, NFS4ERR_FBIG, NFS4ERR_FHEXPIRED, NFS4ERR_GRACE, NFS4ERR_INVAL, NFS4ERR_IO, NFS4ERR_LOCKED, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_NOSPC, NFS4ERR_OLD_STATEID, NFS4ERR_OPENMODE, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_PERM, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_ROFS, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOO_MANY_OPS, NFS4ERR_UNKNOWN_LAYOUTTYPE, NFS4ERR_WRONG_TYPE
SETCLIENTID	NFS4ERR_NOTSUPP
SETCLIENTID_CONFIRM	NFS4ERR_NOTSUPP
TEST_STATEID	NFS4ERR_BADXDR, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_TOO_MANY_OPS



VERIFY	NFS4ERR_ACCESS, NFS4ERR_ATTRNOTSUPP, NFS4ERR_BADCHAR, NFS4ERR_BADXDR, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_FHEXPIRED, NFS4ERR_GRACE, NFS4ERR_INVAL, NFS4ERR_IO, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_NOT_SAME, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOO_MANY_OPS, NFS4ERR_UNKNOWN_LAYOUTTYPE, NFS4ERR_WRONG_TYPE
WANT_DELEGATION	NFS4ERR_BADXDR, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_DELEG_ALREADY_WANTED, NFS4ERR_FHEXPIRED, NFS4ERR_GRACE, NFS4ERR_INVAL, NFS4ERR_IO, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_NOTSUPP, NFS4ERR_NO_GRACE, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_RECALLCONFLICT, NFS4ERR_RECLAIM_BAD, NFS4ERR_RECLAIM_CONFLICT, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONG_TYPE
WRITE	NFS4ERR_ACCESS, NFS4ERR_ADMIN_REVOKED, NFS4ERR_BADXDR, NFS4ERR_BAD_STATEID, NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_DELEG_REVOKED, NFS4ERR_DQUOT, NFS4ERR_EXPIRED, NFS4ERR_FBIG, NFS4ERR_FHEXPIRED, NFS4ERR_GRACE, NFS4ERR_INVAL, NFS4ERR_IO, NFS4ERR_ISDIR, NFS4ERR_LOCKED, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE, NFS4ERR_NOSPC, NFS4ERR_OLD_STATEID, NFS4ERR_OPENMODE, NFS4ERR_OP_NOT_IN_SESSION,

NFS4ERR_PNFS_IO_HOLE,
NFS4ERR_PNFS_NO_LAYOUT,
NFS4ERR_REP_TOO_BIG,
NFS4ERR_REP_TOO_BIG_TO_CACHE,
NFS4ERR_REQ_TOO_BIG,
NFS4ERR_RETRY_UNCACHED_REP,
NFS4ERR_ROFS, NFS4ERR_SERVERFAULT,
NFS4ERR_STALE, NFS4ERR_SYMLINK,
NFS4ERR_TOO_MANY_OPS,
NFS4ERR_WRONG_TYPE

Table 12: Valid Error Returns for Each Protocol Operation

### 22.3. Callback Operations and Their Valid Errors

This section contains a table that gives the valid error returns for each callback operation. The error code NFS4\_OK (indicating no error) is not listed but should be understood to be returnable by all callback operations with the exception of CB\_ILLEGAL.

Callback Operation	Errors
CB_GETATTR	NFS4ERR_BADHANDLE, NFS4ERR_BADXDR, NFS4ERR_DELAY, NFS4ERR_INVAL, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_TOO_MANY_OPS,
CB_ILLEGAL	NFS4ERR_BADXDR, NFS4ERR_OP_ILLEGAL
CB_LAYOUTRECALL	NFS4ERR_BADHANDLE, NFS4ERR_BADIOMODE, NFS4ERR_BADXDR, NFS4ERR_BAD_STATEID, NFS4ERR_DELAY, NFS4ERR_INVAL, NFS4ERR_NOMATCHING_LAYOUT, NFS4ERR_NOTSUPP, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_TOO_MANY_OPS, NFS4ERR_UNKNOWN_LAYOUTTYPE,

	NFS4ERR_WRONG_TYPE
CB_NOTIFY	NFS4ERR_BADHANDLE, NFS4ERR_BADXDR, NFS4ERR_BAD_STATEID, NFS4ERR_DELAY, NFS4ERR_INVAL, NFS4ERR_NOTSUPP, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_TOO_MANY_OPS
CB_NOTIFY_DEVICEID	NFS4ERR_BADXDR, NFS4ERR_DELAY, NFS4ERR_INVAL, NFS4ERR_NOTSUPP, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_TOO_MANY_OPS
CB_NOTIFY_LOCK	NFS4ERR_BADHANDLE, NFS4ERR_BADXDR, NFS4ERR_BAD_STATEID, NFS4ERR_DELAY, NFS4ERR_NOTSUPP, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_TOO_MANY_OPS
CB_PUSH_DELEG	NFS4ERR_BADHANDLE, NFS4ERR_BADXDR, NFS4ERR_DELAY, NFS4ERR_INVAL, NFS4ERR_NOTSUPP, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REJECT_DELEG, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_TOO_MANY_OPS, NFS4ERR_WRONG_TYPE
CB_RECALL	NFS4ERR_BADHANDLE, NFS4ERR_BADXDR,

	NFS4ERR_BAD_STATEID, NFS4ERR_DELAY, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_TOO_MANY_OPS
CB_RECALL_ANY	NFS4ERR_BADXDR, NFS4ERR_DELAY, NFS4ERR_INVAL, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_TOO_MANY_OPS
CB_RECALLABLE_OBJ_AVAIL	NFS4ERR_BADXDR, NFS4ERR_DELAY, NFS4ERR_INVAL, NFS4ERR_NOTSUPP, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_TOO_MANY_OPS
CB_RECALL_SLOT	NFS4ERR_BADXDR, NFS4ERR_BAD_HIGH_SLOT, NFS4ERR_DELAY, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_TOO_MANY_OPS
CB_SEQUENCE	NFS4ERR_BADSESSION, NFS4ERR_BADSLOT, NFS4ERR_BADXDR, NFS4ERR_BAD_HIGH_SLOT, NFS4ERR_CONN_NOT_BOUND_TO_SESSION, NFS4ERR_DELAY, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SEQUENCE_POS, NFS4ERR_SEQ_FALSE_RETRY, NFS4ERR_SEQ_MISORDERED,

	NFS4ERR_TOO_MANY_OPS
CB_WANTS_CANCELLED	NFS4ERR_BADXDR, NFS4ERR_DELAY, NFS4ERR_NOTSUPP, NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG, NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG, NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_TOO_MANY_OPS

Table 13: Valid Error Returns for Each Protocol Callback Operation

## 22.4. Errors and the Operations That Use Them

Error	Operations
NFS4ERR_ACCESS	ACCESS, COMMIT, CREATE, GETATTR, GET_DIR_DELEGATION, LAYOUTCOMMIT, LAYOUTGET, LINK, LOCK, LOCKT, LOCKU, LOOKUP, LOOKUPP, NVERIFY, OPEN, OPENATTR, READ, READDIR, READLINK, REMOVE, RENAME, SECINFO, SECINFO_NO_NAME, SETATTR, VERIFY, WRITE
NFS4ERR_ACLFT_TOOBIG	GETATTR
NFS4ERR_ACLST_POORFIT	SETATTR
NFS4ERR_ACLST_NOFIT	SETATTR
NFS4ERR_ACLST_TOOBIG	SETATTR
NFS4ERR_ADMIN_REVOKED	CLOSE, DELEGRETURN, LAYOUTCOMMIT, LAYOUTGET, LAYOUTRETURN, LOCK, LOCKU, OPEN, OPEN_DOWNGRADE, READ, SETATTR, WRITE
NFS4ERR_ATTRNOTSUPP	CREATE, LAYOUTCOMMIT, NVERIFY, OPEN, SETATTR, VERIFY

NFS4ERR_BACK_CHAN_BUSY	DESTROY_SESSION
NFS4ERR_BADCHAR	CREATE, EXCHANGE_ID, LINK, LOOKUP, NVERIFY, OPEN, REMOVE, RENAME, SECINFO, SETATTR, VERIFY
NFS4ERR_BADHANDLE	CB_GETATTR, CB_LAYOUTRECALL, CB_NOTIFY, CB_NOTIFY_LOCK, CB_PUSH_DELEG, CB_RECALL, PUTFH
NFS4ERR_BADIOMODE	CB_LAYOUTRECALL, LAYOUTCOMMIT, LAYOUTGET
NFS4ERR_BADLAYOUT	LAYOUTCOMMIT, LAYOUTGET
NFS4ERR_BADNAME	CREATE, LINK, LOOKUP, OPEN, REMOVE, RENAME, SECINFO
NFS4ERR_BADOWNER	CREATE, OPEN, SETATTR
NFS4ERR_BADSESSION	BIND_CONN_TO_SESSION, CB_SEQUENCE, DESTROY_SESSION, SEQUENCE
NFS4ERR_BADSLOT	CB_SEQUENCE, SEQUENCE
NFS4ERR_BADTYPE	CREATE
NFS4ERR_BADXDR	ACCESS, BACKCHANNEL_CTL, BIND_CONN_TO_SESSION, CB_GETATTR, CB_ILLEGAL, CB_LAYOUTRECALL, CB_NOTIFY, CB_NOTIFY_DEVICEID, CB_NOTIFY_LOCK, CB_PUSH_DELEG, CB_RECALL, CB_RECALLABLE_OBJ_AVAIL, CB_RECALL_ANY, CB_RECALL_SLOT, CB_SEQUENCE, CB_WANTS_CANCELLED, CLOSE, COMMIT, CREATE, CREATE_SESSION, DELEGPURGE, DELEGRETURN, DESTROY_CLIENTID, DESTROY_SESSION, EXCHANGE_ID, FREE_STATEID,

	GETATTR, GETDEVICEINFO, GETDEVICELIST, GET_DIR_DELEGATION, ILLEGAL, LAYOUTCOMMIT, LAYOUTGET, LAYOUTRETURN, LINK, LOCK, LOCKT, LOCKU, LOOKUP, NVERIFY, OPEN, OPENATTR, OPEN_DOWNGRADE, PUTFH, READ, READDIR, RECLAIM_COMPLETE, REMOVE, RENAME, SECINFO, SECINFO_NO_NAME, SEQUENCE, SETATTR, SET_SSV, TEST_STATEID, VERIFY, WANT_DELEGATION, WRITE
NFS4ERR_BAD_COOKIE	GETDEVICELIST, READDIR
NFS4ERR_BAD_HIGH_SLOT	CB_RECALL_SLOT, CB_SEQUENCE, SEQUENCE
NFS4ERR_BAD_RANGE	LOCK, LOCKT, LOCKU
NFS4ERR_BAD_SESSION_DIGEST	BIND_CONN_TO_SESSION, SET_SSV
NFS4ERR_BAD_STATEID	CB_LAYOUTRECALL, CB_NOTIFY, CB_NOTIFY_LOCK, CB_RECALL, CLOSE, DELEGRETURN, FREE_STATEID, LAYOUTGET, LAYOUTRETURN, LOCK, LOCKU, OPEN, OPEN_DOWNGRADE, READ, SETATTR, WRITE
NFS4ERR_CB_PATH_DOWN	DESTROY_SESSION
NFS4ERR_CLID_INUSE	CREATE_SESSION, EXCHANGE_ID
NFS4ERR_CLIENTID_BUSY	DESTROY_CLIENTID
NFS4ERR_COMPLETE_ALREADY	RECLAIM_COMPLETE
NFS4ERR_CONN_NOT_BOUND_TO_SESSION	CB_SEQUENCE, DESTROY_SESSION, SEQUENCE
NFS4ERR_DEADLOCK	LOCK
NFS4ERR_DEADSESSION	ACCESS, BACKCHANNEL_CTL, BIND_CONN_TO_SESSION, CLOSE,

	COMMIT, CREATE, CREATE_SESSION, DELEGPURGE, DELEGRETURN, DESTROY_CLIENTID, DESTROY_SESSION, EXCHANGE_ID, FREE_STATEID, GETATTR, GETDEVICEINFO, GETDEVICELIST, GET_DIR_DELEGATION, LAYOUTCOMMIT, LAYOUTGET, LAYOUTRETURN, LINK, LOCK, LOCKT, LOCKU, LOOKUP, LOOKUPP, NVERIFY, OPEN, OPENATTR, OPEN_DOWNGRADE, PUTFH, PUTPUBFH, PUTROOTFH, READ, READDIR, READLINK, RECLAIM_COMPLETE, REMOVE, RENAME, RESTOREFH, SAVEFH, SECINFO, SECINFO_NO_NAME, SEQUENCE, SETATTR, SET_SSV, TEST_STATEID, VERIFY, WANT_DELEGATION, WRITE
NFS4ERR_DELAY	ACCESS, BACKCHANNEL_CTL, BIND_CONN_TO_SESSION, CB_GETATTR, CB_LAYOUTRECALL, CB_NOTIFY, CB_NOTIFY_DEVICEID, CB_NOTIFY_LOCK, CB_PUSH_DELEG, CB_RECALL, CB_RECALLABLE_OBJ_AVAIL, CB_RECALL_ANY, CB_RECALL_SLOT, CB_SEQUENCE, CB_WANTS_CANCELLED, CLOSE, COMMIT, CREATE, CREATE_SESSION, DELEGPURGE, DELEGRETURN, DESTROY_CLIENTID, DESTROY_SESSION, EXCHANGE_ID, FREE_STATEID, GETATTR, GETDEVICEINFO, GETDEVICELIST, GET_DIR_DELEGATION, LAYOUTCOMMIT, LAYOUTGET, LAYOUTRETURN, LINK, LOCK, LOCKT, LOCKU, LOOKUP, LOOKUPP, NVERIFY, OPEN, OPENATTR, OPEN_DOWNGRADE,



	PUTFH, PUTPUBFH, PUTROOTFH, READ, READDIR, READLINK, RECLAIM_COMPLETE, REMOVE, RENAME, SECINFO, SECINFO_NO_NAME, SEQUENCE, SETATTR, SET_SSV, TEST_STATEID, VERIFY, WANT_DELEGATION, WRITE
NFS4ERR_DELEG_ALREADY_WANTED	OPEN, WANT_DELEGATION
NFS4ERR_DELEG_REVOKED	DELEGRETURN, LAYOUTCOMMIT, LAYOUTGET, LAYOUTRETURN, OPEN, READ, SETATTR, WRITE
NFS4ERR_DENIED	LOCK, LOCKT
NFS4ERR_DIRDELEG_UNAVAIL	GET_DIR_DELEGATION
NFS4ERR_DQUOT	CREATE, LAYOUTGET, LINK, OPEN, OPENATTR, RENAME, SETATTR, WRITE
NFS4ERR_ENCR_ALG_UNSUPP	EXCHANGE_ID
NFS4ERR_EXIST	CREATE, LINK, OPEN, RENAME
NFS4ERR_EXPIRED	CLOSE, DELEGRETURN, LAYOUTCOMMIT, LAYOUTGET, LAYOUTRETURN, LOCK, LOCKU, OPEN, OPEN_DOWNGRADE, READ, SETATTR, WRITE
NFS4ERR_FBIG	LAYOUTCOMMIT, OPEN, SETATTR, WRITE
NFS4ERR_FHEXPIRED	ACCESS, CLOSE, COMMIT, CREATE, DELEGRETURN, GETATTR, GETDEVICELIST, GETFH, GET_DIR_DELEGATION, LAYOUTCOMMIT, LAYOUTGET, LAYOUTRETURN, LINK, LOCK, LOCKT, LOCKU, LOOKUP, LOOKUP, NVERIFY, OPEN, OPENATTR, OPEN_DOWNGRADE, READ, READDIR, READLINK, RECLAIM_COMPLETE, REMOVE, RENAME, RESTOREFH, SAVEFH,

	SECINFO, SECINFO_NO_NAME, SETATTR, VERIFY, WANT_DELEGATION, WRITE
NFS4ERR_FILE_OPEN	LINK, REMOVE, RENAME
NFS4ERR_GRACE	GETATTR, GET_DIR_DELEGATION, LAYOUTCOMMIT, LAYOUTGET, LAYOUTRETURN, LINK, LOCK, LOCKT, NVERIFY, OPEN, READ, REMOVE, RENAME, SETATTR, VERIFY, WANT_DELEGATION, WRITE
NFS4ERR_HASH_ALG_UNSUPP	EXCHANGE_ID
NFS4ERR_INVAL	ACCESS, BACKCHANNEL_CTL, BIND_CONN_TO_SESSION, CB_GETATTR, CB_LAYOUTRECALL, CB_NOTIFY, CB_NOTIFY_DEVICEID, CB_PUSH_DELEG, CB_RECALLABLE_OBJ_AVAIL, CB_RECALL_ANY, CREATE, CREATE_SESSION, DELEGRETURN, EXCHANGE_ID, GETATTR, GETDEVICEINFO, GETDEVICELIST, GET_DIR_DELEGATION, LAYOUTCOMMIT, LAYOUTGET, LAYOUTRETURN, LINK, LOCK, LOCKT, LOCKU, LOOKUP, NVERIFY, OPEN, OPEN_DOWNGRADE, READ, READDIR, READLINK, RECLAIM_COMPLETE, REMOVE, RENAME, SECINFO, SECINFO_NO_NAME, SETATTR, SET_SSV, VERIFY, WANT_DELEGATION, WRITE
NFS4ERR_IO	ACCESS, COMMIT, CREATE, GETATTR, GETDEVICELIST, GET_DIR_DELEGATION, LAYOUTCOMMIT, LAYOUTGET, LINK, LOOKUP, LOOKUPP, NVERIFY, OPEN, OPENATTR, READ, READDIR, READLINK,

	REMOVE, RENAME, SETATTR, VERIFY, WANT_DELEGATION, WRITE
NFS4ERR_ISDIR	COMMIT, LAYOUTCOMMIT, LAYOUTRETURN, LINK, LOCK, LOCKT, OPEN, READ, WRITE
NFS4ERR_LAYOUTTRYLATER	LAYOUTGET
NFS4ERR_LAYOUTUNAVAILABLE	LAYOUTGET
NFS4ERR_LOCKED	LAYOUTGET, READ, SETATTR, WRITE
NFS4ERR_LOCKS_HELD	CLOSE, FREE_STATEID
NFS4ERR_LOCK_NOTSUPP	LOCK
NFS4ERR_LOCK_RANGE	LOCK, LOCKT, LOCKU
NFS4ERR_MLINK	CREATE, LINK, RENAME
NFS4ERR_MOVED	ACCESS, CLOSE, COMMIT, CREATE, DELEGRETURN, GETATTR, GETFH, GET_DIR_DELEGATION, LAYOUTCOMMIT, LAYOUTGET, LAYOUTRETURN, LINK, LOCK, LOCKT, LOCKU, LOOKUP, LOOKUPP, NVERIFY, OPEN, OPENATTR, OPEN_DOWNGRADE, PUTFH, READ, READDIR, READLINK, RECLAIM_COMPLETE, REMOVE, RENAME, RESTOREFH, SAVEFH, SECINFO, SECINFO_NO_NAME, SETATTR, VERIFY, WANT_DELEGATION, WRITE
NFS4ERR_NAMETOOLONG	CREATE, LINK, LOOKUP, OPEN, REMOVE, RENAME, SECINFO
NFS4ERR_NOENT	BACKCHANNEL_CTL, CREATE_SESSION, EXCHANGE_ID, GETDEVICEINFO, LOOKUP, LOOKUPP, OPEN, OPENATTR, REMOVE, RENAME, SECINFO,

	SECINFO_NO_NAME
NFS4ERR_NOFILEHANDLE	ACCESS, CLOSE, COMMIT, CREATE, DELEGRETURN, GETATTR, GETDEVICELIST, GETFH, GET_DIR_DELEGATION, LAYOUTCOMMIT, LAYOUTGET, LAYOUTRETURN, LINK, LOCK, LOCKT, LOCKU, LOOKUP, LOOKUPP, NVERIFY, OPEN, OPENATTR, OPEN_DOWNGRADE, READ, READDIR, READLINK, RECLAIM_COMPLETE, REMOVE, RENAME, RESTOREFH, SAVEFH, SECINFO, SECINFO_NO_NAME, SETATTR, VERIFY, WANT_DELEGATION, WRITE
NFS4ERR_NOMATCHING_LAYOUT	CB_LAYOUTRECALL
NFS4ERR_NOSPC	CREATE, CREATE_SESSION, LAYOUTGET, LINK, OPEN, OPENATTR, RENAME, SETATTR, WRITE
NFS4ERR_NOTDIR	CREATE, GET_DIR_DELEGATION, LINK, LOOKUP, LOOKUPP, OPEN, READDIR, REMOVE, RENAME, SECINFO, SECINFO_NO_NAME
NFS4ERR_NOTEMPTY	REMOVE, RENAME
NFS4ERR_NOTSUPP	CB_LAYOUTRECALL, CB_NOTIFY, CB_NOTIFY_DEVICEID, CB_NOTIFY_LOCK, CB_PUSH_DELEG, CB_RECALLABLE_OBJ_AVAIL, CB_WANTS_CANCELLED, DELEGPURGE, DELEGRETURN, GETDEVICEINFO, GETDEVICELIST, GET_DIR_DELEGATION, LAYOUTCOMMIT, LAYOUTGET, LAYOUTRETURN, LINK, OPENATTR, OPEN_CONFIRM, RELEASE_LOCKOWNER, RENEW, SECINFO_NO_NAME, SETCLIENTID,

	SETCLIENTID_CONFIRM, WANT_DELEGATION
NFS4ERR_NOT_ONLY_OP	BIND_CONN_TO_SESSION, CREATE_SESSION, DESTROY_CLIENTID, DESTROY_SESSION, EXCHANGE_ID
NFS4ERR_NOT_SAME	EXCHANGE_ID, GETDEVICELIST, READDIR, VERIFY
NFS4ERR_NO_GRACE	LAYOUTCOMMIT, LAYOUTRETURN, LOCK, OPEN, WANT_DELEGATION
NFS4ERR_OLD_STATEID	CLOSE, DELEGRETURN, FREE_STATEID, LAYOUTGET, LAYOUTRETURN, LOCK, LOCKU, OPEN, OPEN_DOWNGRADE, READ, SETATTR, WRITE
NFS4ERR_OPENMODE	LAYOUTGET, LOCK, READ, SETATTR, WRITE
NFS4ERR_OP_ILLEGAL	CB_ILLEGAL, ILLEGAL
NFS4ERR_OP_NOT_IN_SESSION	ACCESS, BACKCHANNEL_CTL, CB_GETATTR, CB_LAYOUTRECALL, CB_NOTIFY, CB_NOTIFY_DEVICEID, CB_NOTIFY_LOCK, CB_PUSH_DELEG, CB_RECALL, CB_RECALLABLE_OBJ_AVAIL, CB_RECALL_ANY, CB_RECALL_SLOT, CB_WANTS_CANCELLED, CLOSE, COMMIT, CREATE, DELEGPURGE, DELEGRETURN, FREE_STATEID, GETATTR, GETDEVICEINFO, GETDEVICELIST, GETFH, GET_DIR_DELEGATION, LAYOUTCOMMIT, LAYOUTGET, LAYOUTRETURN, LINK, LOCK, LOCKT, LOCKU, LOOKUP, LOOKUPP, NVERIFY, OPEN, OPENATTR, OPEN_DOWNGRADE, PUTFH, PUTPUBFH, PUTROOTFH, READ, READDIR, READLINK, RECLAIM_COMPLETE, REMOVE,

	RENAME, RESTOREFH, SAVEFH, SECINFO, SECINFO_NO_NAME, SETATTR, SET_SSV, TEST_STATEID, VERIFY, WANT_DELEGATION, WRITE
NFS4ERR_PERM	CREATE, OPEN, SETATTR
NFS4ERR_PNFS_IO_HOLE	READ, WRITE
NFS4ERR_PNFS_NO_LAYOUT	READ, WRITE
NFS4ERR_RECALLCONFLICT	LAYOUTGET, WANT_DELEGATION
NFS4ERR_RECLAIM_BAD	LAYOUTCOMMIT, LOCK, OPEN, WANT_DELEGATION
NFS4ERR_RECLAIM_CONFLICT	LAYOUTCOMMIT, LOCK, OPEN, WANT_DELEGATION
NFS4ERR_REJECT_DELEG	CB_PUSH_DELEG
NFS4ERR_REP_TOO_BIG	ACCESS, BACKCHANNEL_CTL, BIND_CONN_TO_SESSION, CB_GETATTR, CB_LAYOUTRECALL, CB_NOTIFY, CB_NOTIFY_DEVICEID, CB_NOTIFY_LOCK, CB_PUSH_DELEG, CB_RECALL, CB_RECALLABLE_OBJ_AVAIL, CB_RECALL_ANY, CB_RECALL_SLOT, CB_SEQUENCE, CB_WANTS_CANCELLED, CLOSE, COMMIT, CREATE, CREATE_SESSION, DELEGPURGE, DELEGRETURN, DESTROY_CLIENTID, DESTROY_SESSION, EXCHANGE_ID, FREE_STATEID, GETATTR, GETDEVICEINFO, GETDEVICELIST, GET_DIR_DELEGATION, LAYOUTCOMMIT, LAYOUTGET, LAYOUTRETURN, LINK, LOCK, LOCKT, LOCKU, LOOKUP, LOOKUPP, NVERIFY, OPEN, OPENATTR, OPEN_DOWNGRADE, PUTFH, PUTPUBFH, PUTROOTFH,

	READ, READDIR, READLINK, RECLAIM_COMPLETE, REMOVE, RENAME, RESTOREFH, SAVEFH, SECINFO, SECINFO_NO_NAME, SEQUENCE, SETATTR, SET_SSV, TEST_STATEID, VERIFY, WANT_DELEGATION, WRITE
NFS4ERR_REP_TOO_BIG_TO_CACHE	ACCESS, BACKCHANNEL_CTL, BIND_CONN_TO_SESSION, CB_GETATTR, CB_LAYOUTRECALL, CB_NOTIFY, CB_NOTIFY_DEVICEID, CB_NOTIFY_LOCK, CB_PUSH_DELEG, CB_RECALL, CB_RECALLABLE_OBJ_AVAIL, CB_RECALL_ANY, CB_RECALL_SLOT, CB_SEQUENCE, CB_WANTS_CANCELLED, CLOSE, COMMIT, CREATE, CREATE_SESSION, DELEGPURGE, DELEGRETURN, DESTROY_CLIENTID, DESTROY_SESSION, EXCHANGE_ID, FREE_STATEID, GETATTR, GETDEVICEINFO, GETDEVICELIST, GET_DIR_DELEGATION, LAYOUTCOMMIT, LAYOUTGET, LAYOUTRETURN, LINK, LOCK, LOCKT, LOCKU, LOOKUP, LOOKUPP, NVERIFY, OPEN, OPENATTR, OPEN_DOWNGRADE, PUTFH, PUTPUBFH, PUTROOTFH, READ, READDIR, READLINK, RECLAIM_COMPLETE, REMOVE, RENAME, RESTOREFH, SAVEFH, SECINFO, SECINFO_NO_NAME, SEQUENCE, SETATTR, SET_SSV, TEST_STATEID, VERIFY, WANT_DELEGATION, WRITE
NFS4ERR_REQ_TOO_BIG	ACCESS, BACKCHANNEL_CTL, BIND_CONN_TO_SESSION, CB_GETATTR, CB_LAYOUTRECALL, CB_NOTIFY, CB_NOTIFY_DEVICEID, CB_NOTIFY_LOCK,

	CB_PUSH_DELEG, CB_RECALL, CB_RECALLABLE_OBJ_AVAIL, CB_RECALL_ANY, CB_RECALL_SLOT, CB_SEQUENCE, CB_WANTS_CANCELLED, CLOSE, COMMIT, CREATE, CREATE_SESSION, DELEGPURGE, DELEGRETURN, DESTROY_CLIENTID, DESTROY_SESSION, EXCHANGE_ID, FREE_STATEID, GETATTR, GETDEVICEINFO, GETDEVICELIST, GET_DIR_DELEGATION, LAYOUTCOMMIT, LAYOUTGET, LAYOUTRETURN, LINK, LOCK, LOCKT, LOCKU, LOOKUP, LOOKUPP, NVERIFY, OPEN, OPENATTR, OPEN_DOWNGRADE, PUTFH, PUTPUBFH, PUTROOTFH, READ, READDIR, READLINK, RECLAIM_COMPLETE, REMOVE, RENAME, RESTOREFH, SAVEFH, SECINFO, SECINFO_NO_NAME, SEQUENCE, SETATTR, SET_SSV, TEST_STATEID, VERIFY, WANT_DELEGATION, WRITE
NFS4ERR_RETRY_UNCACHED_REP	ACCESS, BACKCHANNEL_CTL, BIND_CONN_TO_SESSION, CB_GETATTR, CB_LAYOUTRECALL, CB_NOTIFY, CB_NOTIFY_DEVICEID, CB_NOTIFY_LOCK, CB_PUSH_DELEG, CB_RECALL, CB_RECALLABLE_OBJ_AVAIL, CB_RECALL_ANY, CB_RECALL_SLOT, CB_SEQUENCE, CB_WANTS_CANCELLED, CLOSE, COMMIT, CREATE, CREATE_SESSION, DELEGPURGE, DELEGRETURN, DESTROY_CLIENTID, DESTROY_SESSION, EXCHANGE_ID, FREE_STATEID, GETATTR, GETDEVICEINFO, GETDEVICELIST, GET_DIR_DELEGATION,



	LAYOUTCOMMIT, LAYOUTGET, LAYOUTRETURN, LINK, LOCK, LOCKT, LOCKU, LOOKUP, LOOKUPP, NVERIFY, OPEN, OPENATTR, OPEN_DOWNGRADE, PUTFH, PUTPUBFH, PUTROOTFH, READ, READDIR, READLINK, RECLAIM_COMPLETE, REMOVE, RENAME, RESTOREFH, SAVEFH, SECINFO, SECINFO_NO_NAME, SEQUENCE, SETATTR, SET_SSV, TEST_STATEID, VERIFY, WANT_DELEGATION, WRITE
NFS4ERR_ROFS	CREATE, LINK, LOCK, LOCKT, OPEN, OPENATTR, OPEN_DOWNGRADE, REMOVE, RENAME, SETATTR, WRITE
NFS4ERR_SAME	NVERIFY
NFS4ERR_SEQUENCE_POS	CB_SEQUENCE, SEQUENCE
NFS4ERR_SEQ_FALSE_RETRY	CB_SEQUENCE, SEQUENCE
NFS4ERR_SEQ_MISORDERED	CB_SEQUENCE, CREATE_SESSION, SEQUENCE
NFS4ERR_SERVERFAULT	ACCESS, BIND_CONN_TO_SESSION, CB_GETATTR, CB_NOTIFY, CB_NOTIFY_DEVICEID, CB_NOTIFY_LOCK, CB_PUSH_DELEG, CB_RECALL, CB_RECALLABLE_OBJ_AVAIL, CB_WANTS_CANCELLED, CLOSE, COMMIT, CREATE, CREATE_SESSION, DELEGPURGE, DELEGRETURN, DESTROY_CLIENTID, DESTROY_SESSION, EXCHANGE_ID, FREE_STATEID, GETATTR, GETDEVICEINFO, GETDEVICELIST, GET_DIR_DELEGATION, LAYOUTCOMMIT, LAYOUTGET, LAYOUTRETURN, LINK, LOCK, LOCKU, LOOKUP, LOOKUPP,

	NVERIFY, OPEN, OPENATTR, OPEN_DOWNGRADE, PUTFH, PUTPUBFH, PUTROOTFH, READ, READDIR, READLINK, RECLAIM_COMPLETE, REMOVE, RENAME, RESTOREFH, SAVEFH, SECINFO, SECINFO_NO_NAME, SETATTR, TEST_STATEID, VERIFY, WANT_DELEGATION, WRITE
NFS4ERR_SHARE_DENIED	OPEN
NFS4ERR_STALE	ACCESS, CLOSE, COMMIT, CREATE, DELEGRETURN, GETATTR, GETFH, GET_DIR_DELEGATION, LAYOUTCOMMIT, LAYOUTGET, LAYOUTRETURN, LINK, LOCK, LOCKT, LOCKU, LOOKUP, LOOKUPP, NVERIFY, OPEN, OPENATTR, OPEN_DOWNGRADE, PUTFH, READ, READDIR, READLINK, RECLAIM_COMPLETE, REMOVE, RENAME, RESTOREFH, SAVEFH, SECINFO, SECINFO_NO_NAME, SETATTR, VERIFY, WANT_DELEGATION, WRITE
NFS4ERR_STALE_CLIENTID	CREATE_SESSION, DESTROY_CLIENTID, DESTROY_SESSION
NFS4ERR_SYMLINK	COMMIT, LAYOUTCOMMIT, LINK, LOCK, LOCKT, LOOKUP, LOOKUPP, OPEN, READ, WRITE
NFS4ERR_TOOSMALL	CREATE_SESSION, GETDEVICEINFO, LAYOUTGET, READDIR
NFS4ERR_TOO_MANY_OPS	ACCESS, BACKCHANNEL_CTL, BIND_CONN_TO_SESSION, CB_GETATTR, CB_LAYOUTRECALL, CB_NOTIFY, CB_NOTIFY_DEVICEID, CB_NOTIFY_LOCK,

	CB_PUSH_DELEG, CB_RECALL, CB_RECALLABLE_OBJ_AVAIL, CB_RECALL_ANY, CB_RECALL_SLOT, CB_SEQUENCE, CB_WANTS_CANCELLED, CLOSE, COMMIT, CREATE, CREATE_SESSION, DELEGPURGE, DELEGRETURN, DESTROY_CLIENTID, DESTROY_SESSION, EXCHANGE_ID, FREE_STATEID, GETATTR, GETDEVICEINFO, GETDEVICELIST, GET_DIR_DELEGATION, LAYOUTCOMMIT, LAYOUTGET, LAYOUTRETURN, LINK, LOCK, LOCKT, LOCKU, LOOKUP, LOOKUPP, NVERIFY, OPEN, OPENATTR, OPEN_DOWNGRADE, PUTFH, PUTPUBFH, PUTROOTFH, READ, READDIR, READLINK, RECLAIM_COMPLETE, REMOVE, RENAME, RESTOREFH, SAVEFH, SECINFO, SECINFO_NO_NAME, SEQUENCE, SETATTR, SET_SSV, TEST_STATEID, VERIFY, WANT_DELEGATION, WRITE
NFS4ERR_UNKNOWN_LAYOUTTYPE	CB_LAYOUTRECALL, GETDEVICEINFO, GETDEVICELIST, LAYOUTCOMMIT, LAYOUTGET, LAYOUTRETURN, NVERIFY, SETATTR, VERIFY
NFS4ERR_UNSAFE_COMPOUND	CREATE, OPEN, OPENATTR
NFS4ERR_WRONGSEC	LINK, LOOKUP, LOOKUPP, OPEN, PUTFH, PUTPUBFH, PUTROOTFH, RENAME, RESTOREFH
NFS4ERR_WRONG_CRED	CLOSE, CREATE_SESSION, DELEGPURGE, DELEGRETURN, DESTROY_CLIENTID, DESTROY_SESSION, FREE_STATEID, LAYOUTCOMMIT, LAYOUTRETURN, LOCK, LOCKT, LOCKU, OPEN_DOWNGRADE, RECLAIM_COMPLETE

NFS4ERR_WRONG_TYPE	CB_LAYOUTRECALL, CB_PUSH_DELEG, COMMIT, GETATTR, LAYOUTGET, LAYOUTRETURN, LINK, LOCK, LOCKT, NVERIFY, OPEN, OPENATTR, READ, READLINK, RECLAIM_COMPLETE, SETATTR, VERIFY, WANT_DELEGATION, WRITE
NFS4ERR_XDEV	LINK, RENAME

Table 14: Errors and the Operations That Use Them

## 23. NFSv4.1 Procedures

Both procedures, NULL and COMPOUND, MUST be implemented.

## 23.1. Procedure 0: NULL - No Operation

## 23.1.1. ARGUMENTS

void;

## 23.1.2. RESULTS

void;

## 23.1.3. DESCRIPTION

This is the standard NULL procedure with the standard void argument and void response. This procedure has no functionality associated with it. Because of this, it is sometimes used to measure the overhead of processing a service request. Therefore, the server should ensure that no unnecessary work is done in servicing this procedure.

## 23.1.4. ERRORS

None.

## 23.2. Procedure 1: COMPOUND - Compound Operations

## 23.2.1. ARGUMENTS

```
enum nfs_opnum4 {
    OP_ACCESS           = 3,
    OP_CLOSE            = 4,
    OP_COMMIT           = 5,
    OP_CREATE           = 6,
    OP_DELEGPURGE       = 7,
    OP_DELEGRETURN      = 8,
    OP_GETATTR          = 9,
    OP_GETFH            = 10,
    OP_LINK              = 11,
    OP_LOCK              = 12,
    OP_LOCKT            = 13,
    OP_LOCKU            = 14,
    OP_LOOKUP            = 15,
    OP_LOOKUPP          = 16,
    OP_NVERIFY          = 17,
    OP_OPEN              = 18,
    OP_OPENATTR         = 19,
    OP_OPEN_CONFIRM     = 20, /* Mandatory not-to-implement */
    OP_OPEN_DOWNGRADE   = 21,
    OP_PUTFH            = 22,
    OP_PUTPUBFH         = 23,
    OP_PUTROOTFH        = 24,
    OP_READ              = 25,
    OP_READDIR          = 26,
    OP_READLINK         = 27,
    OP_REMOVE           = 28,
    OP_RENAME           = 29,
    OP_RENEW            = 30, /* Mandatory not-to-implement */
    OP_RESTOREFH        = 31,
    OP_SAVEFH           = 32,
    OP_SECINFO          = 33,
    OP_SETATTR          = 34,
    OP_SETCLIENTID      = 35, /* Mandatory not-to-implement */
    OP_SETCLIENTID_CONFIRM = 36, /* Mandatory not-to-implement */
    OP_VERIFY           = 37,
    OP_WRITE            = 38,
    OP_RELEASE_LOCKOWNER = 39, /* Mandatory not-to-implement */

    /* new operations for NFSv4.1 */

    OP_BACKCHANNEL_CTL  = 40,
    OP_BIND_CONN_TO_SESSION = 41,
    OP_EXCHANGE_ID      = 42,
    OP_CREATE_SESSION   = 43,
    OP_DESTROY_SESSION  = 44,
    OP_FREE_STATEID     = 45,
    OP_GET_DIR_DELEGATION = 46,
```

```
OP_GETDEVICEINFO      = 47,
OP_GETDEVICELIST      = 48,
OP_LAYOUTCOMMIT       = 49,
OP_LAYOUTGET         = 50,
OP_LAYOUTRETURN       = 51,
OP_SECINFO_NO_NAME    = 52,
OP_SEQUENCE           = 53,
OP_SET_SSV            = 54,
OP_TEST_STATEID       = 55,
OP_WANT_DELEGATION    = 56,
OP_DESTROY_CLIENTID   = 57,
OP_RECLAIM_COMPLETE   = 58,
OP_ILLEGAL            = 10044
};

union nfs_argop4 switch (nfs_opnum4 argop) {
  case OP_ACCESS:      ACCESS4args opaccess;
  case OP_CLOSE:       CLOSE4args opclose;
  case OP_COMMIT:      COMMIT4args opcommit;
  case OP_CREATE:      CREATE4args opcreate;
  case OP_DELEGPURGE:  DELEGPURGE4args opdeleGPurge;
  case OP_DELEGRETURN: DELEGRETURN4args opdelegreturn;
  case OP_GETATTR:     GETATTR4args opgetattr;
  case OP_GETFH:       void;
  case OP_LINK:        LINK4args oplink;
  case OP_LOCK:        LOCK4args oplock;
  case OP_LOCKT:       LOCKT4args oplockt;
  case OP_LOCKU:       LOCKU4args oplocku;
  case OP_LOOKUP:      LOOKUP4args oplookup;
  case OP_LOOKUPP:     void;
  case OP_NVERIFY:     NVERIFY4args opnverify;
  case OP_OPEN:        OPEN4args opopen;
  case OP_OPENATTR:    OPENATTR4args opopenattr;

  /* Not for NFSv4.1 */
  case OP_OPEN_CONFIRM: OPEN_CONFIRM4args opopen_confirm;

  case OP_OPEN_DOWNGRADE:
    OPEN_DOWNGRADE4args opopen_downgrade;

  case OP_PUTFH:       PUTFH4args opputfh;
  case OP_PUTPUBFH:    void;
  case OP_PUTROOTFH:   void;
  case OP_READ:        READ4args opread;
  case OP_READDIR:     READDIR4args opreaddir;
  case OP_READLINK:    void;
  case OP_REMOVE:      REMOVE4args opremove;
  case OP_RENAME:      RENAME4args oprename;
```

```
/* Not for NFSv4.1 */
case OP_RENEW:          RENEW4args oprenew;

case OP_RESTOREFH:      void;
case OP_SAVEFH:         void;
case OP_SECINFO:        SECINFO4args opsecinfo;
case OP_SETATTR:        SETATTR4args opsetattr;

/* Not for NFSv4.1 */
case OP_SETCLIENTID:    SETCLIENTID4args opsetclientid;

/* Not for NFSv4.1 */
case OP_SETCLIENTID_CONFIRM: SETCLIENTID_CONFIRM4args
                                opsetclientid_confirm;
case OP_VERIFY:         VERIFY4args opverify;
case OP_WRITE:          WRITE4args opwrite;

/* Not for NFSv4.1 */
case OP_RELEASE_LOCKOWNER:
                        RELEASE_LOCKOWNER4args
                        oprelease_lockowner;

/* Operations new to NFSv4.1 */
case OP_BACKCHANNEL_CTL:
                        BACKCHANNEL_CTL4args opbackchannel_ctl;

case OP_BIND_CONN_TO_SESSION:
                        BIND_CONN_TO_SESSION4args
                        opbind_conn_to_session;

case OP_EXCHANGE_ID:    EXCHANGE_ID4args opexchange_id;

case OP_CREATE_SESSION:
                        CREATE_SESSION4args opcreate_session;

case OP_DESTROY_SESSION:
                        DESTROY_SESSION4args opdestroy_session;

case OP_FREE_STATEID:   FREE_STATEID4args opfree_stateid;

case OP_GET_DIR_DELEGATION:
                        GET_DIR_DELEGATION4args
                        opget_dir_delegation;

case OP_GETDEVICEINFO:  GETDEVICEINFO4args opgetdeviceinfo;
case OP_GETDEVICELIST:  GETDEVICELIST4args opgetdevicelist;
case OP_LAYOUTCOMMIT:   LAYOUTCOMMIT4args oplayoutcommit;
case OP_LAYOUTGET:      LAYOUTGET4args oplayoutget;
```

```

case OP_LAYOUTRETURN:  LAYOUTRETURN4args oplayoutreturn;

case OP_SECINFO_NO_NAME:
                        SECINFO_NO_NAME4args opsecinfo_no_name;

case OP_SEQUENCE:      SEQUENCE4args opsequence;
case OP_SET_SSV:       SET_SSV4args opset_ssv;
case OP_TEST_STATEID:  TEST_STATEID4args optest_stateid;

case OP_WANT_DELEGATION:
                        WANT_DELEGATION4args opwant_delegation;

case OP_DESTROY_CLIENTID:
                        DESTROY_CLIENTID4args
                        opdestroy_clientid;

case OP_RECLAIM_COMPLETE:
                        RECLAIM_COMPLETE4args
                        opreclaim_complete;

/* Operations not new to NFSv4.1 */
case OP_ILLEGAL:       void;
};

struct COMPOUND4args {
    utf8str_cs         tag;
    uint32_t            minorversion;
    nfs_argop4         argarray<>;
};

```

#### 23.2.2. RESULTS

```

union nfs_resop4 switch (nfs_opnum4 resop) {
case OP_ACCESS:        ACCESS4res opaccess;
case OP_CLOSE:         CLOSE4res opclose;
case OP_COMMIT:        COMMIT4res opcommit;
case OP_CREATE:        CREATE4res opcreate;
case OP_DELEGPURGE:    DELEGPURGE4res opdeleGPurge;
case OP_DELEGRETURN:   DELEGRETURN4res opdelegreturn;
case OP_GETATTR:       GETATTR4res opgetattr;
case OP_GETFH:         GETFH4res opgetfh;
case OP_LINK:          LINK4res oplink;
case OP_LOCK:          LOCK4res oplock;
case OP_LOCKT:         LOCKT4res oplockt;
case OP_LOCKU:         LOCKU4res oplocku;
case OP_LOOKUP:        LOOKUP4res oplookup;
case OP_LOOKUPP:       LOOKUPP4res oplookupp;
case OP_NVERIFY:       NVERIFY4res opnverify;

```



```
case OP_OPEN:          OPEN4res opopen;
case OP_OPENATTR:      OPENATTR4res opopenattr;
/* Not for NFSv4.1 */
case OP_OPEN_CONFIRM:  OPEN_CONFIRM4res opopen_confirm;

case OP_OPEN_DOWNGRADE:
                        OPEN_DOWNGRADE4res
                        opopen_downgrade;

case OP_PUTFH:         PUTFH4res opputfh;
case OP_PUTPUBFH:      PUTPUBFH4res opputpubfh;
case OP_PUTROOTFH:     PUTROOTFH4res opputrootfh;
case OP_READ:          READ4res opread;
case OP_READDIR:       READDIR4res opreaddir;
case OP_READLINK:      READLINK4res opreadlink;
case OP_REMOVE:        REMOVE4res opremove;
case OP_RENAME:        RENAME4res oprename;
/* Not for NFSv4.1 */
case OP_RENEW:         RENEW4res oprenew;
case OP_RESTOREFH:     RESTOREFH4res oprestorefh;
case OP_SAVEFH:        SAVEFH4res opsavefh;
case OP_SECINFO:        SECINFO4res opsecinfo;
case OP_SETATTR:        SETATTR4res opsetattr;
/* Not for NFSv4.1 */
case OP_SETCLIENTID:  SETCLIENTID4res opsetclientid;

/* Not for NFSv4.1 */
case OP_SETCLIENTID_CONFIRM:
                        SETCLIENTID_CONFIRM4res
                        opsetclientid_confirm;

case OP_VERIFY:        VERIFY4res opverify;
case OP_WRITE:         WRITE4res opwrite;

/* Not for NFSv4.1 */
case OP_RELEASE_LOCKOWNER:
                        RELEASE_LOCKOWNER4res
                        oprelease_lockowner;

/* Operations new to NFSv4.1 */
case OP_BACKCHANNEL_CTL:
                        BACKCHANNEL_CTL4res
                        opbackchannel_ctl;

case OP_BIND_CONN_TO_SESSION:
                        BIND_CONN_TO_SESSION4res
                        opbind_conn_to_session;

case OP_EXCHANGE_ID:   EXCHANGE_ID4res opexchange_id;
```

```
case OP_CREATE_SESSION:
    CREATE_SESSION4res
    opcreate_session;

case OP_DESTROY_SESSION:
    DESTROY_SESSION4res
    opdestroy_session;

case OP_FREE_STATEID:  FREE_STATEID4res
    opfree_stateid;

case OP_GET_DIR_DELEGATION:
    GET_DIR_DELEGATION4res
    opget_dir_delegation;

case OP_GETDEVICEINFO: GETDEVICEINFO4res
    opgetdeviceinfo;

case OP_GETDEVICELIST: GETDEVICELIST4res
    opgetdevicelist;

case OP_LAYOUTCOMMIT:  LAYOUTCOMMIT4res oplayoutcommit;
case OP_LAYOUTGET:     LAYOUTGET4res  oplayoutget;
case OP_LAYOUTRETURN:  LAYOUTRETURN4res oplayoutreturn;

case OP_SECINFO_NO_NAME:
    SECINFO_NO_NAME4res
    opsecinfo_no_name;

case OP_SEQUENCE:      SEQUENCE4res opsequence;
case OP_SET_SSV:       SET_SSV4res  opset_ssv;
case OP_TEST_STATEID:  TEST_STATEID4res optest_stateid;

case OP_WANT_DELEGATION:
    WANT_DELEGATION4res
    opwant_delegation;

case OP_DESTROY_CLIENTID:
    DESTROY_CLIENTID4res
    opdestroy_clientid;

case OP_RECLAIM_COMPLETE:
    RECLAIM_COMPLETE4res
    opreclaim_complete;

/* Operations not new to NFSv4.1 */
case OP_ILLEGAL:      ILLEGAL4res  opillegal;
};
```

```
struct COMPOUND4res {  
    nfsstat4      status;  
    utf8str_cs    tag;  
    nfs_resop4    resarray<>;  
};
```

### 23.2.3. DESCRIPTION

The COMPOUND procedure is used to combine one or more NFSv4 operations into a single RPC request. The server interprets each of the operations in turn. If an operation is executed by the server and the status of that operation is NFS4\_OK, then the next operation in the COMPOUND procedure is executed. The server continues this process until there are no more operations to be executed or until one of the operations has a status value other than NFS4\_OK.

In the processing of the COMPOUND procedure, the server may find that it does not have the available resources to execute any or all of the operations within the COMPOUND sequence. See Section 7.6.4 for a more detailed discussion.

The server will generally choose between two methods of decoding the client's request. The first would be the traditional one-pass XDR decode. If there is an XDR decoding error in this case, the RPC XDR decode error would be returned. The second method would be to make an initial pass to decode the basic COMPOUND request and then to XDR decode the individual operations; the most interesting is the decode of attributes. In this case, the server may encounter an XDR decode error during the second pass. If it does, the server would return the error NFS4ERR\_BADXDR to signify the decode error.

The COMPOUND arguments contain a "minorversion" field. For NFSv4.1, the value for this field is 1. If the server receives a COMPOUND procedure with a minorversion field value that it does not support, the server MUST return an error of NFS4ERR\_MINOR\_VERS\_MISMATCH and a zero-length resultdata array.

Contained within the COMPOUND results is a "status" field. If the results array length is non-zero, this status must be equivalent to the status of the last operation that was executed within the COMPOUND procedure. Therefore, if an operation incurred an error then the "status" value will be the same error value as is being returned for the operation that failed.

Note that operations zero and one are not defined for the COMPOUND procedure. Operation 2 is not defined and is reserved for future definition and use with minor versioning. If the server receives an operation array that contains operation 2 and the minorversion field

has a value of zero, an error of NFS4ERR\_OP\_ILLEGAL, as described in the next paragraph, is returned to the client. If an operation array contains an operation 2 and the minorversion field is non-zero and the server does not support the minor version, the server returns an error of NFS4ERR\_MINOR\_VERS\_MISMATCH. Therefore, the NFS4ERR\_MINOR\_VERS\_MISMATCH error takes precedence over all other errors.

It is possible that the server receives a request that contains an operation that is less than the first legal operation (OP\_ACCESS) or greater than the last legal operation (OP\_RECLAIM\_COMPLETE for NFSv4.1). In this case, the server's response will encode the opcode OP\_ILLEGAL rather than the illegal opcode of the request. The status field in the ILLEGAL return results will be set to NFS4ERR\_OP\_ILLEGAL. The COMPOUND procedure's return results will also be NFS4ERR\_OP\_ILLEGAL. Note that for future minor versions, the last legal operation might be different but that illegal operations are dealt with similarly.

The definition of the "tag" in the request is left to the implementer. It may be used to summarize the content of the Compound request for the benefit of packet-sniffers and engineers debugging implementations. However, the value of "tag" in the response needs to be the same as the value provided in the request. This applies to the tag field of the CB\_COMPOUND procedure as well.

#### 23.2.3.1. Current Filehandle and Stateid

The COMPOUND procedure offers a simple environment for the execution of the operations specified by the client. The first two relate to the filehandle while the second two relate to the current stateid.

##### 23.2.3.1.1. Current Filehandle

The current and saved filehandles are used throughout the protocol. Most operations implicitly use the current filehandle as an argument, and many set the current filehandle as part of the results. The combination of client-specified sequences of operations and current and saved filehandle arguments and results allows for greater protocol flexibility. The best or easiest example of current filehandle usage is a sequence like the following:

PUTFH fh1	{fh1}
LOOKUP "compA"	{fh2}
GETATTR	{fh2}
LOOKUP "compB"	{fh3}
GETATTR	{fh3}
LOOKUP "compC"	{fh4}
GETATTR	{fh4}
GETFH	

Figure 2

In this example, the PUTFH (Section 25.19) operation explicitly sets the current filehandle value while the result of each LOOKUP operation sets the current filehandle value to the resultant file system object. Also, the client is able to insert GETATTR operations using the current filehandle as an argument.

The PUTROOTFH (Section 25.21) and PUTPUBFH (Section 25.20) operations also set the current filehandle. The above example would replace "PUTFH fh1" with PUTROOTFH or PUTPUBFH with no filehandle argument in order to achieve the same effect (on the assumption that "compA" is directly below the root of the namespace).

Along with the current filehandle, there is a saved filehandle. While the current filehandle is set as the result of operations like LOOKUP, the saved filehandle must be set directly with the use of the SAVEFH operation. The SAVEFH operation copies the current filehandle value to the saved value. The saved filehandle value is used in combination with the current filehandle value for the LINK and RENAME operations. The RESTOREFH operation will copy the saved filehandle value to the current filehandle value; as a result, the saved filehandle value may be used a sort of "scratch" area for the client's series of operations.

#### 23.2.3.1.2. Current Stateid

With NFSv4.1, additions of a current stateid and a saved stateid have been made to the COMPOUND processing environment; this allows for the passing of stateids between operations. There are no changes to the syntax of the protocol, only changes to the semantics of a few operations.

A "current stateid" is the stateid that is associated with the current filehandle. The current stateid may only be changed by an operation that modifies the current filehandle or returns a stateid. If an operation returns a stateid, it MUST set the current stateid to the returned value. If an operation sets the current filehandle but does not return a stateid, the current stateid MUST be set to the

all-zeros special stateid, i.e., (seqid, other) = (0, 0). If an operation uses a stateid as an argument but does not return a stateid, the current stateid MUST NOT be changed. For example, PUTFH, PUTROOTFH, and PUTPUBFH will change the current server state from {ocfh, (osid)} to {cfh, (0, 0)}, while LOCK will change the current state from {cfh, (osid)} to {cfh, (nsid)}. Operations like LOOKUP that transform a current filehandle and component name into a new current filehandle will also change the current state to {0, 0}. The SAVEFH and RESTOREFH operations will save and restore both the current filehandle and the current stateid as a set.

The following example is the common case of a simple READ operation with a normal stateid showing that the PUTFH initializes the current stateid to (0, 0). The subsequent READ with stateid (sid1) leaves the current stateid unchanged.

```
PUTFH fh1                                - -> {fh1, (0, 0)}
READ (sid1), 0, 1024    {fh1, (0, 0)} -> {fh1, (0, 0)}
```

Figure 3

This next example performs an OPEN with the root filehandle and, as a result, generates stateid (sid1). The next operation specifies the READ with the argument stateid set such that (seqid, other) are equal to (1, 0), but the current stateid set by the previous operation is actually used when the operation is evaluated. This allows correct interaction with any existing, potentially conflicting, locks.

```
PUTROOTFH                                - -> {fh1, (0, 0)}
OPEN "compA"    {fh1, (0, 0)} -> {fh2, (sid1)}
READ (1, 0), 0, 1024    {fh2, (sid1)} -> {fh2, (sid1)}
CLOSE (1, 0)          {fh2, (sid1)} -> {fh2, (sid2)}
```

Figure 4

This next example is similar to the second in how it passes the stateid sid2 generated by the LOCK operation to the next READ operation. This allows the client to explicitly surround a single I/O operation with a lock and its appropriate stateid to guarantee correctness with other client locks. The example also shows how SAVEFH and RESTOREFH can save and later reuse a filehandle and stateid, passing them as the current filehandle and stateid to a READ operation.

```

PUTFH fh1                                - -> {fh1, (0, 0)}
LOCK 0, 1024, (sid1)                     {fh1, (sid1)} -> {fh1, (sid2)}
READ (1, 0), 0, 1024                     {fh1, (sid2)} -> {fh1, (sid2)}
LOCKU 0, 1024, (1, 0)                    {fh1, (sid2)} -> {fh1, (sid3)}
SAVEFH                                    {fh1, (sid3)} -> {fh1, (sid3)}

PUTFH fh2                                {fh1, (sid3)} -> {fh2, (0, 0)}
WRITE (1, 0), 0, 1024                    {fh2, (0, 0)} -> {fh2, (0, 0)}

RESTOREFH                                {fh2, (0, 0)} -> {fh1, (sid3)}
READ (1, 0), 1024, 1024                  {fh1, (sid3)} -> {fh1, (sid3)}

```

Figure 5

The final example shows a disallowed use of the current stateid. The client is attempting to implicitly pass an anonymous special stateid, (0,0), to the READ operation. The server MUST return NFS4ERR\_BAD\_STATEID in the reply to the READ operation.

```

PUTFH fh1                                - -> {fh1, (0, 0)}
READ (1, 0), 0, 1024                     {fh1, (0, 0)} -> NFS4ERR_BAD_STATEID

```

Figure 6

#### 23.2.4. ERRORS

COMPOUND will of course return every error that each operation on the fore channel can return (See Table 12). However, if COMPOUND returns zero operations, obviously the error returned by COMPOUND has nothing to do with an error returned by an operation. The list of errors COMPOUND will return if it processes zero operations include:

Error	Notes
NFS4ERR_BADCHAR	The tag argument has a character the replier does not support.
NFS4ERR_BADXDR	
NFS4ERR_DELAY	
NFS4ERR_INVAL	The tag argument is not in UTF-8 encoding.
NFS4ERR_MINOR_VERS_MISMATCH	
NFS4ERR_SERVERFAULT	
NFS4ERR_TOO_MANY_OPS	
NFS4ERR_REP_TOO_BIG	
NFS4ERR_REP_TOO_BIG_TO_CACHE	
NFS4ERR_REQ_TOO_BIG	

Table 15: COMPOUND Error Returns

#### 24. Operations: REQUIRED, RECOMMENDED, or OPTIONAL

The following tables summarize the operations of the NFSv4.1 protocol and the corresponding designation of REQUIRED, RECOMMENDED, and OPTIONAL to implement or MUST NOT implement. The designation of MUST NOT implement is reserved for those operations that were defined in NFSv4.0 and MUST NOT be implemented in NFSv4.1.

For the most part, the REQUIRED, RECOMMENDED, or OPTIONAL designation for operations sent by the client is for the server implementation. The client is generally required to implement the operations needed for the operating environment for which it serves. For example, a read-only NFSv4.1 client would have no need to implement the WRITE operation and is not required to do so.



The REQUIRED or OPTIONAL designation for callback operations sent by the server is for both the client and server. Generally, the client has the option of creating the backchannel and sending the operations on the fore channel that will be a catalyst for the server sending callback operations. A partial exception is CB\_RECALL\_SLOT; the only way the client can avoid supporting this operation is by not creating a backchannel.

Since this is a summary of the operations and their designation, there are subtleties that are not presented here. Therefore, if there is a question of the requirements of implementation, the operation descriptions themselves must be consulted along with other relevant explanatory text within this specification.

The abbreviations used in the second and third columns of the table are defined as follows.

REQ REQUIRED to implement

REC RECOMMEND to implement

OPT OPTIONAL to implement

MNI MUST NOT implement

For the NFSv4.1 features that are OPTIONAL, the operations that support those features are OPTIONAL, and the server would return NFS4ERR\_NOTSUPP in response to the client's use of those operations. If an OPTIONAL feature is supported, it is possible that a set of operations related to the feature become REQUIRED to implement. The third column of the table designates the feature(s) and if the operation is REQUIRED or OPTIONAL in the presence of support for the feature.

The OPTIONAL features identified and their abbreviations are as follows:

pNFS Parallel NFS

FDELG File Delegations

DDELG Directory Delegations

Operation	REQ, REC, OPT, or MNI	Feature (REQ, REC, or OPT)	Definition
ACCESS	REQ		Section 25.1
BACKCHANNEL_CTL	REQ		Section 25.33
BIND_CONN_TO_SESSION	REQ		Section 25.34
CLOSE	REQ		Section 25.2
COMMIT	REQ		Section 25.3
CREATE	REQ		Section 25.4
CREATE_SESSION	REQ		Section 25.36
DELEGPURGE	OPT	FDELG (REQ)	Section 25.5
DELEGRETURN	OPT	FDELG, DDELG, pNFS (REQ)	Section 25.6
DESTROY_CLIENTID	REQ		Section 25.50
DESTROY_SESSION	REQ		Section 25.37
EXCHANGE_ID	REQ		Section 25.35
FREE_STATEID	REQ		Section 25.38
GETATTR	REQ		Section 25.7
GETDEVICEINFO	OPT	pNFS (REQ)	Section 25.40
GETDEVICELIST	OPT	pNFS (OPT)	Section 25.41
GETFH	REQ		Section 25.8
GET_DIR_DELEGATION	OPT	DDELG (REQ)	Section 25.39
LAYOUTCOMMIT	OPT	pNFS (REQ)	Section 25.42
LAYOUTGET	OPT	pNFS (REQ)	Section 25.43

LAYOUTRETURN	OPT	pNFS (REQ)	Section 25.44
LINK	OPT		Section 25.9
LOCK	REQ		Section 25.10
LOCKT	REQ		Section 25.11
LOCKU	REQ		Section 25.12
LOOKUP	REQ		Section 25.13
LOOKUPP	REQ		Section 25.14
NVERIFY	REQ		Section 25.15
OPEN	REQ		Section 25.16
OPENATTR	OPT		Section 25.17
OPEN_CONFIRM	MNI		N/A
OPEN_DOWNGRADE	REQ		Section 25.18
PUTFH	REQ		Section 25.19
PUTPUBFH	REQ		Section 25.20
PUTROOTFH	REQ		Section 25.21
READ	REQ		Section 25.22
REaddir	REQ		Section 25.23
READLINK	OPT		Section 25.24
RECLAIM_COMPLETE	REQ		Section 25.51
RELEASE_LOCKOWNER	MNI		N/A
REMOVE	REQ		Section 25.25
RENAME	REQ		Section 25.26
RENEW	MNI		N/A
RESTOREFH	REQ		Section 25.27

SAVEFH	REQ		Section 25.28	
SECINFO	REQ		Section 25.29	
SECINFO_NO_NAME	REC	pNFS file layout (REQ)	Section 25.45, Section 20.25	
SEQUENCE	REQ		Section 25.46	
SETATTR	REQ		Section 25.30	
SETCLIENTID	MNI		N/A	
SETCLIENTID_CONFIRM	MNI		N/A	
SET_SSV	REQ		Section 25.47	
TEST_STATEID	REQ		Section 25.48	
VERIFY	REQ		Section 25.31	
WANT_DELEGATION	OPT	FDELG (OPT)	Section 25.49	
WRITE	REQ		Section 25.32	

Table 16: Operations

Operation	REQ, REC, OPT, or MNI	Feature (REQ, REC, or OPT)	Definition
CB_GETATTR	OPT	FDELG (REQ)	Section 27.1
CB_LAYOUTRECALL	OPT	pNFS (REQ)	Section 27.3
CB_NOTIFY	OPT	DDELG (REQ)	Section 27.4
CB_NOTIFY_DEVICEID	OPT	pNFS (OPT)	Section 27.12
CB_NOTIFY_LOCK	OPT		Section 27.11
CB_PUSH_DELEG	OPT	FDELG (OPT)	Section 27.5
CB_RECALL	OPT	FDELG, DDELG, pNFS (REQ)	Section 27.2
CB_RECALL_ANY	OPT	FDELG, DDELG, pNFS (REQ)	Section 27.6
CB_RECALL_SLOT	REQ		Section 27.8
CB_RECALLABLE_OBJ_AVAIL	OPT	DDELG, pNFS (REQ)	Section 27.7
CB_SEQUENCE	REQ		Section 27.9
CB_WANTS_CANCELLED	OPT	FDELG, DDELG, pNFS (REQ)	Section 27.10

Table 17: Callback Operations

## 25. NFSv4.1 Operations

### 25.1. Operation 3: ACCESS - Check Access Rights

#### 25.1.1. ARGUMENTS

```
const ACCESS4_READ      = 0x00000001;
const ACCESS4_LOOKUP    = 0x00000002;
const ACCESS4_MODIFY    = 0x00000004;
const ACCESS4_EXTEND    = 0x00000008;
const ACCESS4_DELETE    = 0x00000010;
const ACCESS4_EXECUTE   = 0x00000020;

struct ACCESS4args {
    /* CURRENT_FH: object */
    uint32_t      access;
};
```

#### 25.1.2. RESULTS

```
struct ACCESS4resok {
    uint32_t      supported;
    uint32_t      access;
};

union ACCESS4res switch (nfsstat4 status) {
    case NFS4_OK:
        ACCESS4resok  resok4;
    default:
        void;
};
```

#### 25.1.3. DESCRIPTION

ACCESS determines the access rights that a user, as identified by the credentials in the RPC request, has with respect to the file system object specified by the current filehandle. The client encodes the set of access rights that are to be checked in the bit mask "access". The server checks the permissions encoded in the bit mask. If a status of NFS4\_OK is returned, two bit masks are included in the response. The first, "supported", represents the access rights for which the server can verify reliably. The second, "access", represents the access rights available to the user for the filehandle provided. On success, the current filehandle retains its value.

Note that the reply's supported and access fields MUST NOT contain more values than originally set in the request's access field. For example, if the client sends an ACCESS operation with just the

ACCESS4\_READ value set and the server supports this value, the server MUST NOT set more than ACCESS4\_READ in the supported field even if it could have reliably checked other values.

The reply's access field MUST NOT contain more values than the supported field.

The results of this operation are necessarily advisory in nature. A return status of NFS4\_OK and the appropriate bit set in the bit mask do not imply that such access will be allowed to the file system object in the future. This is because access rights can be revoked by the server at any time.

The following access permissions may be requested:

ACCESS4\_READ Read data from file or read a directory.

ACCESS4\_LOOKUP Look up a name in a directory (no meaning for non-directory objects).

ACCESS4\_MODIFY Rewrite existing file data or modify existing directory entries.

ACCESS4\_EXTEND Write new data to a file as controlled by the ACE mask bit for appending data or add directory entries (for either files or sub-directories).

ACCESS4\_DELETE Delete an existing directory entry. Often, when this is done the link count of the referenced object is decremented to zero and the object itself is deleted when no longer linked- to or open.

ACCESS4\_EXECUTE Execute a regular file (no meaning for a directory).

On success, the current filehandle retains its value.

ACCESS4\_EXECUTE is a challenging semantic to implement because NFS provides remote file access, not remote execution. This leads to the following:

- \* Whether or not a regular file is executable ought to be the responsibility of the NFS client and not the server. And yet the ACCESS operation is specified to seemingly require a server to own that responsibility.
- \* When a client executes a regular file, it has to read the file from the server. Strictly speaking, the server should not allow the client to read a file being executed unless the user has read

permissions on the file. Requiring explicit read permissions on executable files in order to access them over NFS is not going to be acceptable to some users and storage administrators. Historically, NFS servers have allowed a user to READ a file if the user has execute access to the file.

As a practical example, the UNIX specification [access\_api] states that an implementation claiming conformance to UNIX may indicate in the access() programming interface's result that a privileged user has execute rights, even if no execute permission bits are set on the regular file's attributes. It is possible to claim conformance to the UNIX specification and instead not indicate execute rights in that situation, which is true for some operating environments. Suppose the operating environments of the client and server are implementing the access() semantics for privileged users differently, and the ACCESS operation implementations of the client and server follow their respective access() semantics. This can cause undesired behavior:

- \* Suppose the client's access() interface returns X\_OK if the user is privileged and no execute permission bits are set on the regular file's attribute, and the server's access() interface does not return X\_OK in that situation. Then the client will be unable to execute files stored on the NFS server that could be executed if stored on a non-NFS file system.
- \* Suppose the client's access() interface does not return X\_OK if the user is privileged, and no execute permission bits are set on the regular file's attribute, and the server's access() interface does return X\_OK in that situation. Then:
  - The client will be able to execute files stored on the NFS server that could be executed if stored on a non-NFS file system, unless the client's execution subsystem also checks for execute permission bits.
  - Even if the execution subsystem is checking for execute permission bits, there are more potential issues. For example, suppose the client is invoking access() to build a "path search table" of all executable files in the user's "search path", where the path is a list of directories each containing executable files. Suppose there are two files each in separate directories of the search path, such that files have the same component name. In the first directory the file has no execute permission bits set, and in the second directory the file has execute bits set. The path search table will indicate that the first directory has the executable file, but the execute subsystem will fail to execute it. The command shell might



fail to try the second file in the second directory. And even if it did, this is a potential performance issue. Clearly, the desired outcome for the client is for the path search table to not contain the first file.

To deal with the problems described above, the "smart client, stupid server" principle is used. The client owns overall responsibility for determining execute access and relies on the server to parse the execution permissions within the file's mode, acl, and dacl attributes. The rules for the client and server follow:

- \* If the client is sending ACCESS in order to determine if the user can read the file, the client will set ACCESS4\_READ in the request's access field.
- \* If the client's operating environment only grants execution to the user if the user has execute access according to the execute permissions in the mode, acl, and dacl attributes, then if the client wants to determine execute access, the client will send an ACCESS request with ACCESS4\_EXECUTE bit set in the request's access field.
- \* If the client's operating environment grants execution to the user even if the user does not have execute access according to the execute permissions in the mode, acl, and dacl attributes, then if the client wants to determine execute access, it would send an ACCESS request with both the ACCESS4\_EXECUTE and ACCESS4\_READ bits set in the request's access field. This way, if any read or execute permission grants the user read or execute access (or if the server interprets the user as privileged), as indicated by the presence of ACCESS4\_EXECUTE and/or ACCESS4\_READ in the reply's access field, the client will be able to grant the user execute access to the file.
- \* If the server supports execute permission bits, or some other method for denoting executability (e.g., the suffix of the name of the file might indicate execute), it MUST check only execute permissions, not read permissions, when determining whether or not the reply will have ACCESS4\_EXECUTE set in the access field. The server MUST NOT also examine read permission bits when determining whether or not the reply will have ACCESS4\_EXECUTE set in the access field. Even if the server's operating environment would grant execute access to the user (e.g., the user is privileged), the server MUST NOT reply with ACCESS4\_EXECUTE set in reply's access field unless there is at least one execute permission bit set in the mode, acl, or dacl attributes. In the case of acl and dacl, the "one execute permission bit" MUST be an ACE4\_EXECUTE bit set in an ALLOW ACE.

- \* If the server does not support execute permission bits or some other method for denoting executability, it MUST NOT set ACCESS4\_EXECUTE in the reply's supported and access fields. If the client set ACCESS4\_EXECUTE in the ACCESS request's access field, and ACCESS4\_EXECUTE is not set in the reply's supported field, then the client will have to send an ACCESS request with the ACCESS4\_READ bit set in the request's access field.
- \* If the server supports read permission bits, it MUST only check for read permissions in the mode, acl, and dacl attributes when it receives an ACCESS request with ACCESS4\_READ set in the access field. The server MUST NOT also examine execute permission bits when determining whether the reply will have ACCESS4\_READ set in the access field or not.

Note that if the ACCESS reply has ACCESS4\_READ or ACCESS\_EXECUTE set, then the user also has permissions to OPEN (Section 25.16) or READ (Section 25.22) the file. In other words, if the client sends an ACCESS request with the ACCESS4\_READ and ACCESS\_EXECUTE set in the access field (or two separate requests, one with ACCESS4\_READ set and the other with ACCESS4\_EXECUTE set), and the reply has just ACCESS4\_EXECUTE set in the access field (or just one reply has ACCESS4\_EXECUTE set), then the user has authorization to OPEN or READ the file.

#### 25.1.4. IMPLEMENTATION

In general, it is not sufficient for the client to attempt to deduce access permissions by inspecting the uid, gid, and mode fields in the file attributes or by attempting to interpret the contents of the ACL attribute. This is because the server may perform uid or gid mapping or enforce additional access-control restrictions. It is also possible that the server may not be in the same ID space as the client. In these cases (and perhaps others), the client cannot reliably perform an access check with only current file attributes.

In the NFSv2 protocol, the only reliable way to determine whether an operation was allowed was to try it and see if it succeeded or failed. Using the ACCESS operation in the NFSv4.1 protocol, the client can ask the server to indicate whether or not one or more classes of operations are permitted. The ACCESS operation is provided to allow clients to check before doing a series of operations that will result in an access failure. The OPEN operation provides a point where the server can verify access to the file object and a method to return that information to the client. The ACCESS operation is still useful for directory operations or for use in the case that the UNIX interface access() is used on the client.

The information returned by the server in response to an ACCESS call is not permanent. It was correct at the exact time that the server performed the checks, but not necessarily afterwards. The server can revoke access permission at any time.

The client should use the effective credentials of the user to build the authentication information in the ACCESS request used to determine access rights. It is the effective user and group credentials that are used in subsequent READ and WRITE operations.

Many implementations do not directly support the ACCESS4\_DELETE permission. Operating systems like UNIX will ignore the ACCESS4\_DELETE bit if set on an access request on a non-directory object. In these systems, delete permission on a file is determined by the access permissions on the directory in which the file resides, instead of being determined by the permissions of the file itself. Therefore, the mask returned enumerating which access rights can be determined will have the ACCESS4\_DELETE value set to 0. This indicates to the client that the server was unable to check that particular access right. The ACCESS4\_DELETE bit in the access mask returned will then be ignored by the client.

## 25.2. Operation 4: CLOSE - Close File

### 25.2.1. ARGUMENTS

```
struct CLOSE4args {
    /* CURRENT_FH: object */
    seqid4      seqid;
    stateid4    open_stateid;
};
```

### 25.2.2. RESULTS

```
union CLOSE4res switch (nfsstat4 status) {
    case NFS4_OK:
        stateid4      open_stateid;
    default:
        void;
};
```

### 25.2.3. DESCRIPTION

The CLOSE operation releases share reservations for the regular or named attribute file as specified by the current filehandle. The share reservations and other state information released at the server as a result of this CLOSE are only those associated with the supplied stateid. State associated with other OPENS is not affected.

If byte-range locks are held, the client should release all locks before sending a CLOSE. The server MAY free all outstanding locks on CLOSE, but some servers might not support the CLOSE of a file that still has byte-range locks held. The server MUST return failure if any locks would exist after the CLOSE.

The argument seqid MAY have any value, and the server MUST ignore seqid.

On success, the current filehandle retains its value.

The server MAY require that the combination of principal, security flavor, and, if applicable, GSS mechanism that sent the OPEN request also be the one to CLOSE the file. This might not be possible if credentials for the principal are no longer available. The server MAY allow the machine credential or SSV credential (See Section 25.35) to send CLOSE.

#### 25.2.4. IMPLEMENTATION

Even though CLOSE returns a stateid, this stateid is not useful to the client and should be treated as deprecated. CLOSE "shuts down" the state associated with all OPENS for the file by a single open-owner. As noted above, CLOSE will either release all file-locking state or return an error. Therefore, the stateid returned by CLOSE is not useful for operations that follow. To help find any uses of this stateid by clients, the server should return the invalid special stateid (the "other" value is zero and the "seqid" field is NFS4\_UINT32\_MAX, see Section 13.2.3).

A CLOSE operation may make delegations grantable where they were not previously. Servers may choose to respond immediately if there are pending delegation want requests or may respond to the situation at a later time.

### 25.3. Operation 5: COMMIT - Commit Cached Data

#### 25.3.1. ARGUMENTS

```
struct COMMIT4args {
    /* CURRENT_FH: file */
    offset4      offset;
    count4      count;
};
```

#### 25.3.2. RESULTS

```
struct COMMIT4resok {
    verifier4    writeverf;
};

union COMMIT4res switch (nfsstat4 status) {
    case NFS4_OK:
        COMMIT4resok    resok4;
    default:
        void;
};
```

#### 25.3.3. DESCRIPTION

The COMMIT operation forces or flushes uncommitted, modified data to stable storage for the file specified by the current filehandle. The flushed data is that which was previously written with one or more WRITE operations that had the "committed" field of their results field set to UNSTABLE4.

The offset specifies the position within the file where the flush is to begin. An offset value of zero means to flush data starting at the beginning of the file. The count specifies the number of bytes of data to flush. If the count is zero, a flush from the offset to the end of the file is done.

The server returns a write verifier upon successful completion of the COMMIT. The write verifier is used by the client to determine if the server has restarted between the initial WRITE operations and the COMMIT. The client does this by comparing the write verifier returned from the initial WRITE operations and the verifier returned by the COMMIT operation. The server must vary the value of the write verifier at each server event or instantiation that may lead to a loss of uncommitted data. Most commonly this occurs when the server is restarted; however, other events at the server may result in uncommitted data loss as well.

On success, the current filehandle retains its value.

#### 25.3.4. IMPLEMENTATION

The COMMIT operation is similar in operation and semantics to the POSIX fsync() [fsync] system interface that synchronizes a file's state with the disk (file data and metadata is flushed to disk or stable storage). COMMIT performs the same operation for a client, flushing any unsynchronized data and metadata on the server to the server's disk or stable storage for the specified file. When using pNFS, if a WRITE returned UNSTABLE4 and NFL4\_UFLG\_COMMIT\_THRU\_MDS is not set, then the client MUST COMMIT to the data server. The COMMIT

may result in flushing the data but not the metadata. In this case, the metadata MUST be flushed with a subsequent LAYOUTCOMMIT to the metadata server. A complete set of pNFS rules for flushing data and metadata is described in Section 20.13. As in the case of fsync(), it may be that there is some modified data or no modified data to synchronize. The data may have been synchronized by the server's normal periodic buffer synchronization activity. COMMIT should return NFS4\_OK, unless there has been an unexpected error.

COMMIT differs from fsync() in that it is possible for the client to flush a range of the file (most likely triggered by a buffer-reclamation scheme on the client before the file has been completely written).

The server implementation of COMMIT is reasonably simple. If the server receives a full file COMMIT request, that is, starting at offset zero and count zero, it should do the equivalent of applying fsync() to the entire file. Otherwise, it should arrange to have the modified data in the range specified by offset and count to be flushed to stable storage. In both cases, any metadata associated with the file must be flushed to stable storage before returning. It is not an error for there to be nothing to flush on the server. This means that the data and metadata that needed to be flushed have already been flushed or lost during the last server failure.

The client implementation of COMMIT is a little more complex. There are two reasons for wanting to commit a client buffer to stable storage. The first is that the client wants to reuse a buffer. In this case, the offset and count of the buffer are sent to the server in the COMMIT request. The server then flushes any modified data based on the offset and count, and flushes any modified metadata associated with the file. It then returns the status of the flush and the write verifier. The second reason for the client to generate a COMMIT is for a full file flush, such as may be done at close. In this case, the client would gather all of the buffers for this file that contain uncommitted data, do the COMMIT operation with an offset of zero and count of zero, and then free all of those buffers. Any other dirty buffers would be sent to the server in the normal fashion.

After a buffer is written (via the WRITE operation) by the client with the "committed" field in the result of WRITE set to UNSTABLE4, the buffer must be considered as modified by the client until the buffer has either been flushed via a COMMIT operation or written via a WRITE operation with the "committed" field in the result set to FILE\_SYNC4 or DATA\_SYNC4. This is done to prevent the buffer from being freed and reused before the data can be flushed to stable storage on the server.

When a response is returned from either a WRITE or a COMMIT operation and it contains a write verifier that differs from that previously returned by the server, the client will need to retransmit all of the buffers containing uncommitted data to the server. How this is to be done is up to the implementer. If there is only one buffer of interest, then it should be sent in a WRITE request with the FILE\_SYNC4 stable parameter. If there is more than one buffer, it might be worthwhile retransmitting all of the buffers in WRITE operations with the stable parameter set to UNSTABLE4 and then retransmitting the COMMIT operation to flush all of the data on the server to stable storage. However, if the server repeatably returns from COMMIT a verifier that differs from that returned by WRITE, the only way to ensure progress is to retransmit all of the buffers with WRITE requests with the FILE\_SYNC4 stable parameter.

The above description applies to page-cache-based systems as well as buffer-cache-based systems. In the former systems, the virtual memory systems, the virtual memory system will need to be modified instead of the buffer cache.

#### 25.4. Operation 6: CREATE - Create a Non-Regular File Object

##### 25.4.1. ARGUMENTS

```
union createtype4 switch (nfs_ftype4 type) {
    case NF4LNK:
        linktext4 linkdata;
    case NF4BLK:
    case NF4CHR:
        specdata4 devdata;
    case NF4SOCK:
    case NF4FIFO:
    case NF4DIR:
        void;
    default:
        void; /* server should return NFS4ERR_BADTYPE */
};

struct CREATE4args {
    /* CURRENT_FH: directory for creation */
    createtype4      objtype;
    component4       objname;
    fattr4           createattrs;
};
```

##### 25.4.2. RESULTS

```
struct CREATE4resok {
    change_info4    cinfo;
    bitmap4         attrset;      /* attributes set */
};

union CREATE4res switch (nfsstat4 status) {
    case NFS4_OK:
        /* new CURRENTFH: created object */
        CREATE4resok resok4;
    default:
        void;
};
```

#### 25.4.3. DESCRIPTION

The CREATE operation creates a file object other than an ordinary file in a directory with a given name. The OPEN operation MUST be used to create a regular file or a named attribute.

The current filehandle must be a directory: an object of type NF4DIR. If the current filehandle is an attribute directory (type NF4ATTRDIR), the error NFS4ERR\_WRONG\_TYPE is returned. If the current filehandle designates any other type of object, the error NFS4ERR\_NOTDIR results.

The objname specifies the name for the new object. The objtype determines the type of object to be created: directory, symlink, etc. If the object type specified is that of an ordinary file, a named attribute, or a named attribute directory, the error NFS4ERR\_BADTYPE results.

If an object of the same name already exists in the directory, the server will return the error NFS4ERR\_EXIST.

For the directory where the new file object was created, the server returns change\_info4 information in cinfo. With the atomic field of the change\_info4 data type, the server will indicate if the before and after change attributes were obtained atomically with respect to the file object creation.

If the objname has a length of zero, or if objname does not obey the UTF-8 definition, the error NFS4ERR\_INVALID will be returned.

The current filehandle is replaced by that of the new object.



The `createattrs` specifies the initial set of attributes for the object. The set of attributes may include any writable attribute valid for the object type. When the operation is successful, the server will return to the client an attribute mask signifying which attributes were successfully set for the object.

If `createattrs` includes neither the owner attribute nor an ACL with an ACE for the owner, and if the server's file system both supports and requires an owner attribute (or an owner ACE), then the server MUST derive the owner (or the owner ACE). This would typically be from the principal indicated in the RPC credentials of the call, but the server's operating environment or file system semantics may dictate other methods of derivation. Similarly, if `createattrs` includes neither the group attribute nor a group ACE, and if the server's file system both supports and requires the notion of a group attribute (or group ACE), the server MUST derive the group attribute (or the corresponding owner ACE) for the file. This could be from the RPC call's credentials, such as the group principal if the credentials include it (such as with `AUTH_SYS`), from the group identifier associated with the principal in the credentials (e.g., POSIX systems have a user database `[passwd]` that has a group identifier for every user identifier), inherited from the directory in which the object is created, or whatever else the server's operating environment or file system semantics dictate. This applies to the `OPEN` operation too.

Conversely, it is possible that the client will specify in `createattrs` an owner attribute, group attribute, or ACL that the principal indicated the RPC call's credentials does not have permissions to create files for. The error to be returned in this instance is `NFS4ERR_PERM`. This applies to the `OPEN` operation too.

If the current filehandle designates a directory for which another client holds a directory delegation, then, unless the delegation is such that the situation can be resolved by sending a notification, the delegation MUST be recalled, and the `CREATE` operation MUST NOT proceed until the delegation is returned or revoked. Except where this happens very quickly, one or more `NFS4ERR_DELAY` errors will be returned to requests made while delegation remains outstanding.

When the current filehandle designates a directory for which one or more directory delegations exist, then, when those delegations request such notifications, `NOTIFY4_ADD_ENTRY` will be generated as a result of this operation.

#### 25.4.4. IMPLEMENTATION

If the client desires to set attribute values after the create, a SETATTR operation can be added to the COMPOUND request so that the appropriate attributes will be set.

#### 25.5. Operation 7: DELEGPURGE - Purge Delegations Awaiting Recovery

##### 25.5.1. ARGUMENTS

```
struct DELEGPURGE4args {  
    clientid4      clientid;  
};
```

##### 25.5.2. RESULTS

```
struct DELEGPURGE4res {  
    nfsstat4      status;  
};
```

##### 25.5.3. DESCRIPTION

This operation purges all of the delegations awaiting recovery for a given client. This is useful for clients that do not commit delegation information to stable storage to indicate that conflicting requests need not be delayed by the server awaiting recovery of delegation information.

The client is NOT specified by the clientid field of the request. The client should set the client field to zero, and the server MUST ignore the clientid field. Instead, the server MUST derive the client ID from the value of the session ID in the arguments of the SEQUENCE operation that precedes DELEGPURGE in the COMPOUND request.

The DELEGPURGE operation should be used by clients that record delegation information on stable storage on the client. In this case, after the client recovers all delegations it knows of, it should immediately send a DELEGPURGE operation. Doing so will notify the server that no additional delegations for the client will be recovered allowing it to free resources, and avoid delaying other clients which make requests that conflict with the unrecovered delegations. The set of delegations known to the server and the client might be different. The reason for this is that after sending a request that resulted in a delegation, the client might experience a failure before it both received the delegation and committed the delegation to the client's stable storage.

The server MAY support DELEGPURGE, but if it does not, it MUST NOT support CLAIM\_DELEGATE\_PREV and MUST NOT support CLAIM\_DELEG\_PREV\_FH.

## 25.6. Operation 8: DELEGRETURN - Return Delegation

### 25.6.1. ARGUMENTS

```
struct DELEGRETURN4args {  
    /* CURRENT_FH: delegated object */  
    stateid4      deleg_stateid;  
};
```

### 25.6.2. RESULTS

```
struct DELEGRETURN4res {  
    nfsstat4      status;  
};
```

### 25.6.3. DESCRIPTION

The DELEGRETURN operation returns the delegation represented by the current filehandle and stateid.

Delegations may be returned voluntarily (i.e., before the server has recalled them) or when recalled. In either case, the client must properly propagate state changed under the context of the delegation to the server before returning the delegation.

The server MAY require that the principal, security flavor, and if applicable, the GSS mechanism, combination that acquired the delegation also be the one to send DELEGRETURN on the file. This might not be possible if credentials for the principal are no longer available. The server MAY allow the machine credential or SSV credential (See Section 25.35) to send DELEGRETURN.

## 25.7. Operation 9: GETATTR - Get Attributes

### 25.7.1. ARGUMENTS

```
struct GETATTR4args {  
    /* CURRENT_FH: object */  
    bitmap4      attr_request;  
};
```

### 25.7.2. RESULTS

```
struct GETATTR4resok {
    fattr4      obj_attributes;
};

union GETATTR4res switch (nfsstat4 status) {
    case NFS4_OK:
        GETATTR4resok  resok4;
    default:
        void;
};
```

### 25.7.3. DESCRIPTION

The GETATTR operation will obtain attributes for the file system object specified by the current filehandle. The client sets a bit in the bitmap argument for each attribute value that it would like the server to return. The server returns an attribute bitmap that indicates the attribute values that it was able to return, which will include all attributes requested by the client that are attributes supported by the server for the target file system. This bitmap is followed by the attribute values ordered lowest attribute number first.

The server MUST return a value for each attribute that the client requests if the attribute is supported by the server for the target file system. If the server does not support a particular attribute on the target file system, then it MUST NOT return the attribute value and MUST NOT set the attribute bit in the result bitmap. The server MUST return an error if it supports an attribute on the target but cannot obtain its value. In that case, no attribute values will be returned.

File systems that are absent should be treated as having support for a very small set of attributes as described in Section 17.4.1, even if previously, when the file system was present, more attributes were supported.

All servers MUST support the REQUIRED attributes as specified in Section 11.10, for all file systems, with the exception of absent file systems.

On success, the current filehandle retains its value.

#### 25.7.4. IMPLEMENTATION

Suppose there is an OPEN\_DELEGATE\_WRITE delegation held by another client for the file in question and size and/or change are among the set of attributes being interrogated. The server has two choices. First, the server can obtain the actual current value of these attributes from the client holding the delegation by using the CB\_GETATTR callback. Second, the server, particularly when the delegated client is unresponsive, can recall the delegation in question. The GETATTR MUST NOT proceed until one of the following occurs:

- \* The requested attribute values are returned in the response to CB\_GETATTR.
- \* The OPEN\_DELEGATE\_WRITE delegation is returned.
- \* The OPEN\_DELEGATE\_WRITE delegation is revoked.

Unless one of the above happens very quickly, one or more NFS4ERR\_DELAY errors will be returned while a delegation is outstanding.

#### 25.8. Operation 10: GETFH - Get Current Filehandle

##### 25.8.1. ARGUMENTS

```
/* CURRENT_FH: */  
void;
```

##### 25.8.2. RESULTS

```
struct GETFH4resok {  
    nfs_fh4      object;  
};  
  
union GETFH4res switch (nfsstat4 status) {  
    case NFS4_OK:  
        GETFH4resok      resok4;  
    default:  
        void;  
};
```

##### 25.8.3. DESCRIPTION

This operation returns the current filehandle value.

On success, the current filehandle retains its value.

As described in Section 7.6.4, GETFH is REQUIRED or RECOMMENDED to immediately follow certain operations, and servers are free to reject such operations if the client fails to insert GETFH in the request as REQUIRED or RECOMMENDED. Section 25.16.4.1 provides additional justification for why GETFH MUST follow OPEN.

#### 25.8.4. IMPLEMENTATION

Operations that change the current filehandle like LOOKUP or CREATE do not automatically return the new filehandle as a result. For instance, if a client needs to look up a directory entry and obtain its filehandle, then the following request is needed.

PUTFH (directory filehandle)

LOOKUP (entry name)

GETFH

#### 25.9. Operation 11: LINK - Create Link to a File

##### 25.9.1. ARGUMENTS

```
struct LINK4args {  
    /* SAVED_FH: source object */  
    /* CURRENT_FH: target directory */  
    component4    newname;  
};
```

##### 25.9.2. RESULTS

```
struct LINK4resok {  
    change_info4    cinfo;  
};  
  
union LINK4res switch (nfsstat4 status) {  
    case NFS4_OK:  
        LINK4resok resok4;  
    default:  
        void;  
};
```

### 25.9.3. DESCRIPTION

The LINK operation creates an additional newname for the file represented by the saved filehandle, as set by the SAVEFH operation, in the directory represented by the current filehandle. The existing file and the target directory must reside within the same file system on the server. On success, the current filehandle will continue to be the target directory. If an object exists in the target directory with the same name as newname, the server must return NFS4ERR\_EXIST.

For the target directory, the server returns change\_info4 information in cinfo. With the atomic field of the change\_info4 data type, the server will indicate if the before and after change attributes were obtained atomically with respect to the link creation.

If the newname has a length of zero, or if newname does not obey the UTF-8 definition, the error NFS4ERR\_INVALID will be returned.

### 25.9.4. IMPLEMENTATION

The server MAY impose restrictions on the LINK operation such that LINK may not be done when the file is open or when that open is done by particular protocols, or with particular options or access modes. When LINK is rejected because of such restrictions, the error NFS4ERR\_FILE\_OPEN is returned.

If a server does implement such restrictions and those restrictions include cases of NFSv4 opens preventing successful execution of a link, the server needs to recall any delegations that could hide the existence of opens relevant to that decision. The reason is that when a client holds a delegation, the server might not have an accurate account of the opens for that client, since the client may execute OPENS and CLOSEs locally. The LINK operation must be delayed only until a definitive result can be obtained. For example, suppose there are multiple delegations and one of them establishes an open whose presence would prevent the link. Given the server's semantics, NFS4ERR\_FILE\_OPEN may be returned to the caller as soon as that delegation is returned without waiting for other delegations to be returned. Similarly, if such opens are not associated with delegations, NFS4ERR\_FILE\_OPEN can be returned immediately with no delegation recall being done.

If the current filehandle designates a directory for which another client holds a directory delegation, then, unless the delegation is such that the situation can be resolved by sending a notification, the delegation **MUST** be recalled, and the operation cannot be performed successfully until the delegation is returned or revoked. Except where this happens very quickly, one or more `NFS4ERR_DELAY` errors will be returned to requests made while delegation remains outstanding.

When the current filehandle designates a directory for which one or more directory delegations exist, then, when those delegations request such notifications, instead of a recall, `NOTIFY4_ADD_ENTRY` will be generated as a result of the `LINK` operation.

If the current file system supports the `numlinks` attribute, and other clients have delegations to the file being linked, then those delegations **MUST** be recalled and the `LINK` operation **MUST NOT** proceed until all delegations are returned or revoked. Except where this happens very quickly, one or more `NFS4ERR_DELAY` errors will be returned to requests made while delegation remains outstanding.

Changes to any property of the "hard" linked files are reflected in all of the linked files. When a link is made to a file, the attributes for the file should have a value for `numlinks` that is one greater than the value before the `LINK` operation.

The statement "file and the target directory must reside within the same file system on the server" means that the `fsid` fields in the attributes for the objects are the same. If they reside on different file systems, the error `NFS4ERR_XDEV` is returned. This error may be returned by some servers when there is an internal partitioning of a file system that the `LINK` operation would violate.

On some servers, `"."` and `".."` are illegal values for `newname` and the error `NFS4ERR_BADNAME` will be returned if they are specified.

When the current filehandle designates a named attribute directory and the object to be linked (the saved filehandle) is not a named attribute for the same object, the error `NFS4ERR_XDEV` **MUST** be returned. When the saved filehandle designates a named attribute and the current filehandle is not the appropriate named attribute directory, the error `NFS4ERR_XDEV` **MUST** also be returned.

When the current filehandle designates a named attribute directory and the object to be linked (the saved filehandle) is a named attribute within that directory, the server may return the error `NFS4ERR_NOTSUPP`.



In the case that newname is already linked to the file represented by the saved filehandle, the server will return NFS4ERR\_EXIST.

Note that symbolic links are created with the CREATE operation.

## 25.10. Operation 12: LOCK - Create Lock

### 25.10.1. ARGUMENTS

```
/*
 * For LOCK, transition from open_stateid and lock_owner
 * to a lock stateid.
 */
struct open_to_lock_owner4 {
    seqid4      open_seqid;
    stateid4    open_stateid;
    seqid4      lock_seqid;
    lock_owner4 lock_owner;
};

/*
 * For LOCK, existing lock stateid continues to request new
 * file lock for the same lock_owner and open_stateid.
 */
struct exist_lock_owner4 {
    stateid4    lock_stateid;
    seqid4      lock_seqid;
};

union locker4 switch (bool new_lock_owner) {
    case TRUE:
        open_to_lock_owner4    open_owner;
    case FALSE:
        exist_lock_owner4      lock_owner;
};

/*
 * LOCK/LOCKT/LOCKU: Record lock management
 */
struct LOCK4args {
    /* CURRENT_FH: file */
    nfs_lock_type4 locktype;
    bool           reclaim;
    offset4        offset;
    length4        length;
    locker4        locker;
};
```

## 25.10.2. RESULTS

```
struct LOCK4denied {
    offset4      offset;
    length4      length;
    nfs_lock_type4 locktype;
    lock_owner4  owner;
};

struct LOCK4resok {
    stateid4      lock_stateid;
};

union LOCK4res switch (nfsstat4 status) {
    case NFS4_OK:
        LOCK4resok      resok4;
    case NFS4ERR_DENIED:
        LOCK4denied      denied;
    default:
        void;
};
```

## 25.10.3. DESCRIPTION

The LOCK operation requests a byte-range lock for the byte-range specified by the offset and length parameters, and lock type specified in the locktype parameter. If this is a reclaim request, the reclaim parameter will be TRUE.

Bytes in a file may be locked even if those bytes are not currently allocated to the file. To lock the file from a specific offset through the end-of-file (no matter how long the file actually is) use a length field equal to NFS4\_UINT64\_MAX. The server MUST return NFS4ERR\_INVALID under the following combinations of length and offset:

- \* Length is equal to zero.
- \* Length is not equal to NFS4\_UINT64\_MAX, and the sum of length and offset exceeds NFS4\_UINT64\_MAX.

32-bit servers are servers that support locking for byte offsets that fit within 32 bits (i.e., less than or equal to NFS4\_UINT32\_MAX). If the client specifies a range that overlaps one or more bytes beyond offset NFS4\_UINT32\_MAX but does not end at offset NFS4\_UINT64\_MAX, then such a 32-bit server MUST return the error NFS4ERR\_BAD\_RANGE.

If the server returns NFS4ERR\_DENIED, the owner, offset, and length of a conflicting lock are returned.

The `locker` argument specifies the lock-owner that is associated with the LOCK operation. The `locker4` structure is a switched union that indicates whether the client has already created byte-range locking state associated with the current open file and lock-owner. In the case in which it has, the argument is just a `stateid` representing the set of locks associated with that open file and lock-owner, together with a `lock_seqid` value that MAY be any value and MUST be ignored by the server. In the case where no byte-range locking state has been established, or the client does not have the `stateid` available, the argument contains the `stateid` of the open file with which this lock is to be associated, together with the lock-owner with which the lock is to be associated. The `open_to_lock_owner` case covers the very first lock done by a lock-owner for a given open file and offers a method to use the established state of the `open_stateid` to transition to the use of a lock `stateid`.

The following fields of the `locker` parameter MAY be set to any value by the client and MUST be ignored by the server:

- \* The `clientid` field of the `lock_owner` field of the `open_owner` field (`locker.open_owner.lock_owner.clientid`). The reason the server MUST ignore the `clientid` field is that the server MUST derive the client ID from the session ID from the SEQUENCE operation of the COMPOUND request.
- \* The `open_seqid` and `lock_seqid` fields of the `open_owner` field (`locker.open_owner.open_seqid` and `locker.open_owner.lock_seqid`).
- \* The `lock_seqid` field of the `lock_owner` field (`locker.lock_owner.lock_seqid`).

Note that the client ID appearing in a `LOCK4denied` structure is the actual client associated with the conflicting lock, whether this is the client ID associated with the current session or a different one. Thus, if the server returns `NFS4ERR_DENIED`, it MUST set the `clientid` field of the owner field of the denied field.

If the current filehandle is not an ordinary file, an error will be returned to the client. In the case that the current filehandle represents an object of type `NF4DIR`, `NFS4ERR_ISDIR` is returned. If the current filehandle designates a symbolic link, `NFS4ERR_SYMLINK` is returned. In all other cases, `NFS4ERR_WRONG_TYPE` is returned.

On success, the current filehandle retains its value.

#### 25.10.4. IMPLEMENTATION

If the server is unable to determine the exact offset and length of the conflicting byte-range lock, the same offset and length that were provided in the arguments should be returned in the denied results.

LOCK operations are subject to permission checks and to checks against the access type of the associated file. However, the specific right and modes required for various types of locks reflect the semantics of the server-exported file system, and are not specified by the protocol. For example, Windows 2000 allows a write lock of a file open for read access, while a POSIX-compliant system does not.

When the client sends a LOCK operation that corresponds to a range that the lock-owner has locked already (with the same or different lock type), or to a sub-range of such a range, or to a byte-range that includes multiple locks already granted to that lock-owner, in whole or in part, and the server does not support such locking operations (i.e., does not support POSIX locking semantics), the server will return the error NFS4ERR\_LOCK\_RANGE. In that case, the client may return an error, or it may emulate the required operations, using only LOCK for ranges that do not include any bytes already locked by that lock-owner and LOCKU of locks held by that lock-owner (specifying an exactly matching range and type). Similarly, when the client sends a LOCK operation that amounts to upgrading (changing from a READ\_LT lock to a WRITE\_LT lock) or downgrading (changing from WRITE\_LT lock to a READ\_LT lock) an existing byte-range lock, and the server does not support such a lock, the server will return NFS4ERR\_LOCK\_NOTSUPP. Such operations may not perfectly reflect the required semantics in the face of conflicting LOCK operations from other clients.

When a client holds an OPEN\_DELEGATE\_WRITE delegation, the client holding that delegation is assured that there are no opens by other clients. Thus, there can be no conflicting LOCK operations from such clients. Therefore, the client may be handling locking requests locally, without doing LOCK operations on the server. If it does that, it must be prepared to update the lock status on the server, by sending appropriate LOCK and LOCKU operations before returning the delegation.

When one or more clients hold OPEN\_DELEGATE\_READ delegations, any LOCK operation where the server is implementing mandatory locking semantics MUST result in the recall of all such delegations. The LOCK operation may not be granted until all such delegations are returned or revoked. Except where this happens very quickly, one or more NFS4ERR\_DELAY errors will be returned to requests made while the delegation remains outstanding.

## 25.11. Operation 13: LOCKT - Test for Lock

### 25.11.1. ARGUMENTS

```
struct LOCKT4args {
    /* CURRENT_FH: file */
    nfs_lock_type4  locktype;
    offset4         offset;
    length4         length;
    lock_owner4     owner;
};
```

### 25.11.2. RESULTS

```
union LOCKT4res switch (nfsstat4 status) {
    case NFS4ERR_DENIED:
        LOCK4denied    denied;
    case NFS4_OK:
        void;
    default:
        void;
};
```

### 25.11.3. DESCRIPTION

The LOCKT operation tests the lock as specified in the arguments. If a conflicting lock exists, the owner, offset, length, and type of the conflicting lock are returned. The owner field in the results includes the client ID of the owner of the conflicting lock, whether this is the client ID associated with the current session or a different client ID. If no lock is held, nothing other than NFS4\_OK is returned. Lock types READ\_LT and READW\_LT are processed in the same way in that a conflicting lock test is done without regard to blocking or non-blocking. The same is true for WRITE\_LT and WRITEW\_LT.

The ranges are specified as for LOCK. The NFS4ERR\_INVAL and NFS4ERR\_BAD\_RANGE errors are returned under the same circumstances as for LOCK.

The `clientid` field of the owner MAY be set to any value by the client and MUST be ignored by the server. The reason the server MUST ignore the `clientid` field is that the server MUST derive the client ID from the session ID from the `SEQUENCE` operation of the `COMPOUND` request.

If the current filehandle is not an ordinary file, an error will be returned to the client. In the case that the current filehandle represents an object of type `NF4DIR`, `NFS4ERR_ISDIR` is returned. If the current filehandle designates a symbolic link, `NFS4ERR_SYMLINK` is returned. In all other cases, `NFS4ERR_WRONG_TYPE` is returned.

On success, the current filehandle retains its value.

#### 25.11.4. IMPLEMENTATION

If the server is unable to determine the exact offset and length of the conflicting lock, the same offset and length that were provided in the arguments should be returned in the denied results.

`LOCKT` uses a `lock_owner4` rather a `stateid4`, as is used in `LOCK` to identify the owner. This is because the client does not have to open the file to test for the existence of a lock, so a `stateid` might not be available.

As noted in Section 25.10.4, some servers may return `NFS4ERR_LOCK_RANGE` to certain (otherwise non-conflicting) `LOCK` operations that overlap ranges already granted to the current lock-owner.

The `LOCKT` operation's test for conflicting locks should exclude locks for the current lock-owner, and thus should return `NFS4_OK` in such cases. Note that this means that a server might return `NFS4_OK` to a `LOCKT` request even though a `LOCK` operation for the same range and lock-owner would fail with `NFS4ERR_LOCK_RANGE`.

When a client holds an `OPEN_DELEGATE_WRITE` delegation, it may choose (See Section 25.10.4) to handle `LOCK` requests locally. In such a case, `LOCKT` requests will similarly be handled locally.

### 25.12. Operation 14: `LOCKU` - Unlock File

#### 25.12.1. ARGUMENTS

```
struct LOCKU4args {  
    /* CURRENT_FH: file */  
    nfs_lock_type4  locktype;  
    seqid4          seqid;  
    stateid4        lock_stateid;  
    offset4         offset;  
    length4         length;  
};
```

#### 25.12.2. RESULTS

```
union LOCKU4res switch (nfsstat4 status) {  
    case NFS4_OK:  
        stateid4        lock_stateid;  
    default:  
        void;  
};
```

#### 25.12.3. DESCRIPTION

The LOCKU operation unlocks the byte-range lock specified by the parameters. The client may set the locktype field to any value that is legal for the nfs\_lock\_type4 enumerated type, and the server MUST accept any legal value for locktype. Any legal value for locktype has no effect on the success or failure of the LOCKU operation.

The ranges are specified as for LOCK. The NFS4ERR\_INVAL and NFS4ERR\_BAD\_RANGE errors are returned under the same circumstances as for LOCK.

The seqid parameter MAY be any value and the server MUST ignore it.

If the current filehandle is not an ordinary file, an error will be returned to the client. In the case that the current filehandle represents an object of type NF4DIR, NFS4ERR\_ISDIR is returned. If the current filehandle designates a symbolic link, NFS4ERR\_SYMLINK is returned. In all other cases, NFS4ERR\_WRONG\_TYPE is returned.

On success, the current filehandle retains its value.

The server MAY require that the principal, security flavor, and if applicable, the GSS mechanism, combination that sent a LOCK operation also be the one to send LOCKU on the file. This might not be possible if credentials for the principal are no longer available. The server MAY allow the machine credential or SSV credential (See Section 25.35) to send LOCKU.

#### 25.12.4. IMPLEMENTATION

If the area to be unlocked does not correspond exactly to a lock actually held by the lock-owner, the server may return the error NFS4ERR\_LOCK\_RANGE. This includes the case in which the area is not locked, where the area is a sub-range of the area locked, where it overlaps the area locked without matching exactly, or the area specified includes multiple locks held by the lock-owner. In all of these cases, allowed by POSIX locking [fcntl] semantics, a client receiving this error should, if it desires support for such operations, simulate the operation using LOCKU on ranges corresponding to locks it actually holds, possibly followed by LOCK operations for the sub-ranges not being unlocked.

When a client holds an OPEN\_DELEGATE\_WRITE delegation, it may choose (See Section 25.10.4) to handle LOCK requests locally. In such a case, LOCKU operations will similarly be handled locally.

#### 25.13. Operation 15: LOOKUP - Lookup Filename

##### 25.13.1. ARGUMENTS

```
struct LOOKUP4args {  
    /* CURRENT_FH: directory */  
    component4      objname;  
};
```

##### 25.13.2. RESULTS

```
struct LOOKUP4res {  
    /* New CURRENT_FH: object */  
    nfsstat4      status;  
};
```

##### 25.13.3. DESCRIPTION

The LOOKUP operation looks up or finds a file system object using the directory specified by the current filehandle. LOOKUP evaluates the component and if the object exists, the current filehandle is replaced with the component's filehandle.

If the component cannot be evaluated either because it does not exist or because the client does not have permission to evaluate the component, then an error will be returned and the current filehandle will be unchanged.

If the component is a zero-length string or if any component does not obey the UTF-8 definition, the error NFS4ERR\_INVALID will be returned.



#### 25.13.4. IMPLEMENTATION

If the client wants to achieve the effect of a multi-component look up, it may construct a COMPOUND request such as (and obtain each filehandle):

```
PUTFH (directory filehandle)
LOOKUP "pub"
GETFH
LOOKUP "foo"
GETFH
LOOKUP "bar"
GETFH
```

Unlike NFSv3, NFSv4.1 allows LOOKUP requests to cross mountpoints on the server. The client can detect a mountpoint crossing by comparing the fsid attribute of the directory with the fsid attribute of the directory looked up. If the fsids are different, then the new directory is a server mountpoint. UNIX clients that detect a mountpoint crossing will need to mount the server's file system. This needs to be done to maintain the file object identity checking mechanisms common to UNIX clients.

Servers that limit NFS access to "shared" or "exported" file systems should provide a pseudo file system into which the exported file systems can be integrated, so that clients can browse the server's namespace. The clients view of a pseudo file system will be limited to paths that lead to exported file systems.

Note: previous versions of the protocol assigned special semantics to the names "." and "..". NFSv4.1 assigns no special semantics to these names. The LOOKUPP operator must be used to look up a parent directory.

Note that this operation does not follow symbolic links. The client is responsible for all parsing of filenames including filenames that are modified by symbolic links encountered during the look up process.

If the current filehandle supplied is not a directory but a symbolic link, the error NFS4ERR\_SYMLINK is returned as the error. For all other non-directory file types, the error NFS4ERR\_NOTDIR is returned.

#### 25.14. Operation 16: LOOKUPP - Lookup Parent Directory

##### 25.14.1. ARGUMENTS

```
/* CURRENT_FH: object */  
void;
```

#### 25.14.2. RESULTS

```
struct LOOKUPP4res {  
    /* new CURRENT_FH: parent directory */  
    nfsstat4      status;  
};
```

#### 25.14.3. DESCRIPTION

The current filehandle is assumed to refer to a regular directory or a named attribute directory. LOOKUPP assigns the filehandle for its parent directory to be the current filehandle. If there is no parent directory, an NFS4ERR\_NOENT error must be returned. Therefore, NFS4ERR\_NOENT will be returned by the server when the current filehandle is at the root or top of the server's file tree.

As is the case with LOOKUP, LOOKUPP will also cross mountpoints.

If the current filehandle is not a directory or named attribute directory, the error NFS4ERR\_NOTDIR is returned.

If the requester's security flavor does not match that configured for the parent directory, then the server should return NFS4ERR\_WRONGSEC (a future minor revision of NFSv4 might require that error return in the LOOKUPP response). However, if the server does so, it MUST support the SECINFO\_NO\_NAME operation (Section 25.45), so that the client can gracefully determine the correct security flavor.

If the current filehandle is a named attribute directory that is associated with a file system object via OPENATTR (i.e., not a sub-directory of a named attribute directory), LOOKUPP is to return the filehandle of the associated file system object.

#### 25.14.4. IMPLEMENTATION

An issue to note is upward navigation from named attribute directories. The named attribute directories are essentially detached from the namespace, and this property should be safely represented in the client operating environment. LOOKUPP on a named attribute directory may return the filehandle of the associated file, and conveying this to applications might be unsafe as many applications expect the parent of an object to always be a directory. Therefore, the client may want to hide the parent of named attribute directories (represented as ".." in UNIX) or represent the named attribute directory as its own parent (as is typically done for the

file system root directory in UNIX).

## 25.15. Operation 17: NVERIFY - Verify Difference in Attributes

### 25.15.1. ARGUMENTS

```
struct NVERIFY4args {  
    /* CURRENT_FH: object */  
    fattr4          obj_attributes;  
};
```

### 25.15.2. RESULTS

```
struct NVERIFY4res {  
    nfsstat4          status;  
};
```

### 25.15.3. DESCRIPTION

This operation is used to prefix a sequence of operations to be performed if one or more attributes have changed on some file system object. If all the attributes match, then the error NFS4ERR\_SAME MUST be returned.

On success, the current filehandle retains its value.

### 25.15.4. IMPLEMENTATION

This operation is useful as a cache validation operator. If the object to which the attributes belong has changed, then the following operations may obtain new data associated with that object, for instance, to check if a file has been changed and obtain new data if it has:

```
SEQUENCE  
PUTFH fh  
NVERIFY attrbits attrs  
READ 0 32767
```

Contrast this with NFSv3, which would first send a GETATTR in one request/reply round trip, and then if attributes indicated that the client's cache was stale, then send a READ in another request/reply round trip.

In the case that an OPTIONAL attribute is specified in the NVERIFY operation and the server does not support that attribute for the file system object, the error NFS4ERR\_ATTRNOTSUPP is returned to the client.

When an attribute is a supported one but is one that is not supported for use in NVERIFY (e.g., rdattrib\_error or any set-only attribute (such as time\_modify\_set)) is specified, the error NFS4ERR\_INVAL is returned to the client.

## 25.16. Operation 18: OPEN - Open a Regular File

### 25.16.1. ARGUMENTS

```

/*
 * Various definitions for OPEN
 */
enum createmode4 {
    UNCHECKED4      = 0,
    GUARDED4        = 1,
    /* Deprecated in NFSv4.1. */
    EXCLUSIVE4       = 2,
    /*
     * New to NFSv4.1. If session is persistent,
     * GUARDED4 MUST be used. Otherwise, use
     * EXCLUSIVE4_1 instead of EXCLUSIVE4.
     */
    EXCLUSIVE4_1     = 3
};

struct creatverfattr {
    verifier4      cva_verf;
    fattr4         cva_attrs;
};

union createhow4 switch (createmode4 mode) {
    case UNCHECKED4:
    case GUARDED4:
        fattr4      createattrs;
    case EXCLUSIVE4:
        verifier4    createverf;
    case EXCLUSIVE4_1:
        creatverfattr ch_createboth;
};

enum opentype4 {
    OPEN4_NOCREATE = 0,
    OPEN4_CREATE   = 1
};

union openflag4 switch (opentype4 opentype) {
    case OPEN4_CREATE:
        createhow4    how;
};

```

```
default:
    void;
};

/* Next definitions used for OPEN delegation */
enum limit_by4 {
    NFS_LIMIT_SIZE          = 1,
    NFS_LIMIT_BLOCKS        = 2
    /* others as needed */
};

struct nfs_modified_limit4 {
    uint32_t    num_blocks;
    uint32_t    bytes_per_block;
};

union nfs_space_limit4 switch (limit_by4 limitby) {
    /* limit specified as file size */
    case NFS_LIMIT_SIZE:
        uint64_t    filesize;
    /* limit specified by number of blocks */
    case NFS_LIMIT_BLOCKS:
        nfs_modified_limit4    mod_blocks;
} ;

/*
 * Share Access and Deny constants for open argument
 */
const OPEN4_SHARE_ACCESS_READ    = 0x00000001;
const OPEN4_SHARE_ACCESS_WRITE   = 0x00000002;
const OPEN4_SHARE_ACCESS_BOTH    = 0x00000003;

const OPEN4_SHARE_DENY_NONE      = 0x00000000;
const OPEN4_SHARE_DENY_READ      = 0x00000001;
const OPEN4_SHARE_DENY_WRITE     = 0x00000002;
const OPEN4_SHARE_DENY_BOTH      = 0x00000003;

/* new flags for share_access field of OPEN4args */
const OPEN4_SHARE_ACCESS_WANT_DELEG_MASK    = 0xFF00;
const OPEN4_SHARE_ACCESS_WANT_NO_PREFERENCE = 0x0000;
const OPEN4_SHARE_ACCESS_WANT_READ_DELEG     = 0x0100;
const OPEN4_SHARE_ACCESS_WANT_WRITE_DELEG    = 0x0200;
const OPEN4_SHARE_ACCESS_WANT_ANY_DELEG      = 0x0300;
const OPEN4_SHARE_ACCESS_WANT_NO_DELEG       = 0x0400;
const OPEN4_SHARE_ACCESS_WANT_CANCEL         = 0x0500;

const
```

```
OPEN4_SHARE_ACCESS_WANT_SIGNAL_DELEG_WHEN_RESRC_AVAIL
= 0x10000;

const
OPEN4_SHARE_ACCESS_WANT_PUSH_DELEG_WHEN_UNCONTENDED
= 0x20000;

enum open_delegation_type4 {
    OPEN_DELEGATE_NONE      = 0,
    OPEN_DELEGATE_READ      = 1,
    OPEN_DELEGATE_WRITE     = 2,
    OPEN_DELEGATE_NONE_EXT  = 3 /* new to v4.1 */
};

/*
 * Includes multiple types of operations:
 *
 * - Non-reclaim operations valid independent of grace period
 *   status.
 * - Reclaim operations only used during a grace period.
 * - Reclaim operations only used during a special delegation
 *   recovery period.
 */

enum open_claim_type4 {
    CLAIM_NULL              = 0,      /* Non-reclaim operation, */
    CLAIM_PREVIOUS          = 1,      /* Reclaim operation --
                                     grace period only. */
    CLAIM_DELEGATE_CUR      = 2,      /* Non-reclaim operation. */
    CLAIM_DELEGATE_PREV     = 3,      /* Reclaim operation --
                                     special delegation
                                     recovery period only. */

    /*
     * Beyond this point, all values are new to v4.1.
     */

    /*
     * Like CLAIM_NULL, but object identified
     * by the current filehandle.
     */
    CLAIM_FH                = 4,      /* Non-reclaim operation. */

    /*
     * Like CLAIM_DELEGATE_CUR, but object identified
     * by current filehandle.
     */
    CLAIM_DELEG_CUR_FH      = 5,      /* Non-reclaim operation. */
}
```

```
/*
 * Like CLAIM_DELEGATE_PREV, but object identified
 * by current filehandle.
 */
CLAIM_DELEG_PREV_FH      = 6      /* Reclaim operation --
                                     special delegation
                                     recovery period
                                     only. */
};

struct open_claim_delegate_cur4 {
    stateid4      delegate_stateid;
    component4    file;
};

union open_claim4 switch (open_claim_type4 claim) {
/*
 * No special rights to file.
 * Ordinary OPEN of the specified file.
 */
case CLAIM_NULL:
    /* CURRENT_FH: directory */
    component4    file;
/*
 * Right to the file established by an
 * open previous to server reboot. File
 * identified by filehandle obtained at
 * that time rather than by name.
 */
case CLAIM_PREVIOUS:
    /* CURRENT_FH: file being reclaimed */
    open_delegation_type4  delegate_type;

/*
 * Right to file based on a delegation
 * granted by the server. File is
 * specified by name.
 */
case CLAIM_DELEGATE_CUR:
    /* CURRENT_FH: directory */
    open_claim_delegate_cur4      delegate_cur_info;

/*
 * Right to file based on a delegation
 * granted to a previous boot instance
 * of the client. File is specified by name.
 */
case CLAIM_DELEGATE_PREV:
```

```

        /* CURRENT_FH: directory */
        component4      file_delegate_prev;

/*
 * Like CLAIM_NULL.  No special rights
 * to file.  Ordinary OPEN of the
 * specified file by current filehandle.
 */
case CLAIM_FH: /* new to v4.1 */
    /* CURRENT_FH: regular file to open */
    void;

/*
 * Like CLAIM_DELEGATE_PREV.  Right to file based on a
 * delegation granted to a previous boot
 * instance of the client.  File is identified
 * by filehandle.
 */
case CLAIM_DELEG_PREV_FH: /* new to v4.1 */
    /* CURRENT_FH: file being opened */
    void;

/*
 * Like CLAIM_DELEGATE_CUR.  Right to file based on
 * a delegation granted by the server.
 * File is identified by filehandle.
 */
case CLAIM_DELEG_CUR_FH: /* new to v4.1 */
    /* CURRENT_FH: file being opened */
    stateid4      oc_delegate_stateid;

};

/*
 * OPEN: Open a file, potentially receiving an OPEN delegation
 */
struct OPEN4args {
    seqid4      seqid;
    uint32_t    share_access;
    uint32_t    share_deny;
    open_owner4 owner;
    openflag4   openhow;
    open_claim4 claim;
};

```

#### 25.16.2. RESULTS



```
struct open_read_delegation4 {
    stateid4 stateid; /* Stateid for delegation*/
    bool      recall; /* Pre-recalled flag for
                      delegations obtained
                      by reclaim (CLAIM_PREVIOUS) */

    nfsace4 permissions; /* Defines users who don't
                          need an ACCESS call to
                          open for read */
};

struct open_write_delegation4 {
    stateid4 stateid; /* Stateid for delegation */
    bool      recall; /* Pre-recalled flag for
                      delegations obtained
                      by reclaim
                      (CLAIM_PREVIOUS) */

    nfs_space_limit4
        space_limit; /* Defines condition that
                      the client must check to
                      determine whether the
                      file needs to be flushed
                      to the server on close. */

    nfsace4 permissions; /* Defines users who don't
                          need an ACCESS call as
                          part of a delegated
                          open. */
};

enum why_no_delegation4 { /* new to v4.1 */
    WND4_NOT_WANTED          = 0,
    WND4_CONTENTION          = 1,
    WND4_RESOURCE            = 2,
    WND4_NOT_SUPP_FTYPE      = 3,
    WND4_WRITE_DELEG_NOT_SUPP_FTYPE = 4,
    WND4_NOT_SUPP_UPGRADE    = 5,
    WND4_NOT_SUPP_DOWNGRADE  = 6,
    WND4_CANCELLED           = 7,
    WND4_IS_DIR              = 8
};

union open_none_delegation4 /* new to v4.1 */
switch (why_no_delegation4 ond_why) {
    case WND4_CONTENTION:
        bool ond_server_will_push_deleg;
```

```

        case WND4_RESOURCE:
            bool ond_server_will_signal_avail;
        default:
            void;
    };

union open_delegation4
switch (open_delegation_type4 delegation_type) {
    case OPEN_DELEGATE_NONE:
        void;
    case OPEN_DELEGATE_READ:
        open_read_delegation4 read;
    case OPEN_DELEGATE_WRITE:
        open_write_delegation4 write;
    case OPEN_DELEGATE_NONE_EXT: /* new to v4.1 */
        open_none_delegation4 od_whynone;
};

/*
 * Result flags
 */

/* Client must confirm open */
const OPEN4_RESULT_CONFIRM = 0x00000002;
/* Type of file locking behavior at the server */
const OPEN4_RESULT_LOCKTYPE_POSIX = 0x00000004;
/* Server will preserve file if removed while open */
const OPEN4_RESULT_PRESERVE_UNLINKED = 0x00000008;

/*
 * Server may use CB_NOTIFY_LOCK on locks
 * derived from this open
 */
const OPEN4_RESULT_MAY_NOTIFY_LOCK = 0x00000020;

struct OPEN4resok {
    stateid4      stateid;      /* Stateid for open */
    change_info4  cinfo;        /* Directory Change Info */
    uint32_t      rflags;        /* Result flags */
    bitmap4       attrset;       /* attribute set for create*/
    open_delegation4 delegation; /* Info on any open
                                delegation */
};

union OPEN4res switch (nfsstat4 status) {
    case NFS4_OK:
        /* New CURRENT_FH: opened file */
        OPEN4resok      resok4;

```

```
default:
    void;
};
```

### 25.16.3. DESCRIPTION

The OPEN operation opens a regular file in a directory with the provided name or filehandle. OPEN can also create a file if a name is provided, and the client specifies it wants to create a file. Specification of whether or not a file is to be created, and the method of creation is via the openhow parameter. The openhow parameter consists of a switched union (data type opengflag4), which switches on the value of opentype (OPEN4\_NOCREATE or OPEN4\_CREATE). If OPEN4\_CREATE is specified, this leads to another switched union (data type createhow4) that supports four cases of creation methods: UNCHECKED4, GUARDED4, EXCLUSIVE4, or EXCLUSIVE4\_1. If opentype is OPEN4\_CREATE, then the claim field of the claim field MUST be one of CLAIM\_NULL, CLAIM\_DELEGATE\_CUR, or CLAIM\_DELEGATE\_PREV, because these claim methods include a component of a file name.

Upon success (which might entail creation of a new file), the current filehandle is replaced by that of the created or existing object.

If the current filehandle is a named attribute directory, OPEN will then create or open a named attribute file. Note that exclusive create of a named attribute is not supported. If the createmode is EXCLUSIVE4 or EXCLUSIVE4\_1 and the current filehandle is a named attribute directory, the server will return EINVAL.

UNCHECKED4 means that the file should be created if a file of that name does not exist and encountering an existing regular file of that name is not an error. For this type of create, createattrs specifies the initial set of attributes for the file. The set of attributes may include any writable attribute valid for regular files. When an UNCHECKED4 create encounters an existing file, the attributes specified by createattrs are not used, except that when createattrs specifies the size attribute with a size of zero, the existing file is truncated.

If GUARDED4 is specified, the server checks for the presence of a duplicate object by name before performing the create. If a duplicate exists, NFS4ERR\_EXIST is returned. If the object does not exist, the request is performed as described for UNCHECKED4.

For the UNCHECKED4 and GUARDED4 cases, where the operation is successful, the server will return to the client an attribute mask signifying which attributes were successfully set for the object.

EXCLUSIVE4\_1 and EXCLUSIVE4 specify that the server is to follow exclusive creation semantics, using the verifier to ensure exclusive creation of the target. The server should check for the presence of a duplicate object by name. If the object does not exist, the server creates the object and stores the verifier with the object. If the object does exist and the stored verifier matches the client provided verifier, the server uses the existing object as the newly created object. If the stored verifier does not match, then an error of NFS4ERR\_EXIST is returned.

If using EXCLUSIVE4, and if the server uses attributes to store the exclusive create verifier, the server will signify which attributes it used by setting the appropriate bits in the attribute mask that is returned in the results. Unlike UNCHECKED4, GUARDED4, and EXCLUSIVE4\_1, EXCLUSIVE4 does not support the setting of attributes at file creation, and after a successful OPEN via EXCLUSIVE4, the client MUST send a SETATTR to set attributes to a known state.

In NFSv4.1, EXCLUSIVE4 has been deprecated in favor of EXCLUSIVE4\_1. Unlike EXCLUSIVE4, attributes may be provided in the EXCLUSIVE4\_1 case, but because the server may use attributes of the target object to store the verifier, the set of allowable attributes may be fewer than the set of attributes SETATTR allows. The allowable attributes for EXCLUSIVE4\_1 are indicated in the `suppattr_exclcreat` (Section 11.12.1.14) attribute. If the client attempts to set in `cva_attrs` an attribute that is not in `suppattr_exclcreat`, the server MUST return NFS4ERR\_INVALID. The response field, `attrset`, indicates both which attributes the server set from `cva_attrs` and which attributes the server used to store the verifier. As described in Section 25.16.4, the client can compare `cva_attrs.attrmask` with `attrset` to determine which attributes were used to store the verifier.

With the addition of persistent sessions and pNFS, under some conditions EXCLUSIVE4 MUST NOT be used by the client or supported by the server. The following table summarizes the appropriate and mandated exclusive create methods for implementations of NFSv4.1:

Persistent Reply Cache Enabled	Server Supports pNFS	Server REQUIRED	Client Allowed
no	no	EXCLUSIVE4_1 and EXCLUSIVE4	EXCLUSIVE4_1 (should) or EXCLUSIVE4 (should not)
no	yes	EXCLUSIVE4_1	EXCLUSIVE4_1
yes	no	GUARDED4	GUARDED4
yes	yes	GUARDED4	GUARDED4

Table 18: Required Methods for Exclusive Create

If `CREATE_SESSION4_FLAG_PERSIST` is set in the results of `CREATE_SESSION`, the reply cache is persistent (See Section 25.36). If the `EXCHGID4_FLAG_USE_PNFS_MDS` flag is set in the results from `EXCHANGE_ID`, the server is a pNFS server (See Section 25.35). If the client attempts to use `EXCLUSIVE4` on a persistent session, or a session derived from an `EXCHGID4_FLAG_USE_PNFS_MDS` client ID, the server MUST return `NFS4ERR_INVAL`.

With persistent sessions, exclusive create semantics are fully achievable via `GUARDED4`, and so `EXCLUSIVE4` or `EXCLUSIVE4_1` MUST NOT be used. When pNFS is being used, the `layout_hint` attribute might not be supported after the file is created. Only the `EXCLUSIVE4_1` and `GUARDED` methods of exclusive file creation allow the atomic setting of attributes.

For the target directory, the server returns `change_info4` information in `cinfo`. With the atomic field of the `change_info4` data type, the server will indicate if the before and after change attributes were obtained atomically with respect to the link creation.

The `OPEN` operation provides for Windows share reservation capability with the use of the `share_access` and `share_deny` fields of the `OPEN` arguments. The client specifies at `OPEN` the required `share_access` and `share_deny` modes. For clients that do not directly support `SHAREs` (i.e., `UNIX`), the expected deny value is `OPEN4_SHARE_DENY_NONE`. In the case that there is an existing `SHARE` reservation that conflicts with the `OPEN` request, the server returns the error `NFS4ERR_SHARE_DENIED`. For additional discussion of `SHARE` semantics, see Section 14.7.

For each OPEN, the client provides a value for the owner field of the OPEN argument. The owner field is of data type open\_owner4, and contains a field called clientid and a field called owner. The client can set the clientid field to any value and the server MUST ignore it. Instead, the server MUST derive the client ID from the session ID of the SEQUENCE operation of the COMPOUND request.

The "seqid" field of the request is not used in NFSv4.1, but it MAY be any value and the server MUST ignore it.

In the case that the client is recovering state from a server failure, the claim field of the OPEN argument is used to signify that the request is meant to reclaim state previously held.

The "claim" field of the OPEN argument is used to specify the file to be opened and the state information that the client claims to possess. There are seven claim types as follows:

open type	description
CLAIM_NULL, CLAIM_FH	For the client, this is a new OPEN request and there is no previous state associated with the file for the client. With CLAIM_NULL, the file is identified by the current filehandle and the specified component name. With CLAIM_FH (new to NFSv4.1), the file is identified by just the current filehandle.
CLAIM_PREVIOUS	The client is claiming basic OPEN state for a file that was held previous to a server restart. Generally used when a server is returning persistent filehandles; the client may not have the file name to reclaim the OPEN.
CLAIM_DELEGATE_CUR, CLAIM_DELEG_CUR_FH	The client is claiming a delegation for OPEN as granted by the server. Generally, this is done as part of recalling a delegation. With CLAIM_DELEGATE_CUR, the file is identified by the current filehandle and the specified component name. With CLAIM_DELEG_CUR_FH (new to NFSv4.1), the file is identified by just the current filehandle.
CLAIM_DELEGATE_PREV, CLAIM_DELEG_PREV_FH	The client is claiming a delegation granted to a previous client instance; used after the client restarts. The server MAY support CLAIM_DELEGATE_PREV and/or CLAIM_DELEG_PREV_FH (new to NFSv4.1). If it does support either claim type, CREATE_SESSION MUST NOT remove the client's delegation state, and the server MUST support the DELEGPURGE operation.

Table 19

For OPEN requests that reach the server during the grace period, the server returns an error of NFS4ERR\_GRACE. The following claim types are exceptions:

- \* OPEN requests specifying the claim type CLAIM\_PREVIOUS are devoted to reclaiming opens after a server restart and are typically only valid during the grace period.
- \* OPEN requests specifying the claim types CLAIM\_DELEGATE\_CUR and CLAIM\_DELEG\_CUR\_FH are valid both during and after the grace period. Since the granting of the delegation that they are subordinate to assures that there is no conflict with locks to be reclaimed by other clients, the server need not return NFS4ERR\_GRACE when these are received during the grace period.
- \* OPEN requests specifying the claim types CLAIM\_DELEGATE\_PREV and CLAIM\_DELEG\_PREV\_FH are valid both during and after the grace period. They must be done during the special delegation recovery period which can overlap a grace period.

For any OPEN request, the server may return an OPEN delegation, which allows further opens and closes to be handled locally on the client as described in Section 15.4. Note that delegation is up to the server to decide. The client should never assume that delegation will or will not be granted in a particular instance. It should always be prepared for either case. A partial exception is the reclaim (CLAIM\_PREVIOUS) case, in which a delegation type is claimed. In this case, delegation will always be granted, although the server may specify an immediate recall in the delegation structure.

The rflags returned by a successful OPEN allow the server to return information governing how the open file is to be handled.

- \* OPEN4\_RESULT\_CONFIRM is deprecated and MUST NOT be returned by an NFSv4.1 server.
- \* OPEN4\_RESULT\_LOCKTYPE\_POSIX indicates that the server's byte-range locking behavior supports the complete set of POSIX locking techniques [fcntl]. From this, the client can choose to manage byte-range locking state in a way to handle a mismatch of byte-range locking management.



- \* OPEN4\_RESULT\_PRESERVE\_UNLINKED indicates that the NFSv4.1 server will preserve the open file if the client (or any other client) removes the file as long as it remains open. Since the server cannot be aware of files opened using NFSv3 and the client has no information regarding this bit for NFSv4.0 opens, the client, in situations which such opens might exist could find it necessary to do a silly-rename even when this bit is set for all NFSv4.1 OPENS.

In addition to the basic guarantee above, the server, by returning this bit, promises to preserve the file through any necessary grace period after server restart or file system migration, thereby giving the client the opportunity to reclaim its open.

In cases in which a client knows that this flag is returned for all NFSv4.1 opens of a particular file, it can avoid the need for a possible "silly rename" of the file to assure its preservation. It should be noted the possibility of files being open by NFSv3 and NFSv4.0 clients may make the use of silly-rename if necessary. See Section 25.25.4 for further details.

- \* OPEN4\_RESULT\_MAY\_NOTIFY\_LOCK indicates that the server may attempt CB\_NOTIFY\_LOCK callbacks for locks on this file. This flag is a hint only, and may be safely ignored by the client.

If the component is of zero length, NFS4ERR\_INVALID will be returned. The component may also be subject to UTF-8, character support, or other name validity checks. See Section 22.1.7 for further discussion.

When an OPEN is done and the specified open-owner already has the resulting filehandle open, the result is to "OR" together the new share and deny status together with the existing status. In this case, only a single CLOSE need be done, even though multiple OPENS were completed. When such an OPEN is done, checking of share reservations for the new OPEN proceeds normally, with no exception for the existing OPEN held by the same open-owner. In this case, the stateid returned as an "other" field that matches that of the previous open while the "seqid" field is incremented to reflect the change status due to the new open.

If the underlying file system at the server is only accessible in a read-only mode and the OPEN request has specified ACCESS\_WRITE or ACCESS\_BOTH, the server will return NFS4ERR\_ROFS to indicate a read-only file system.

As with the CREATE operation, the server MUST derive the owner, owner ACE, group, or group ACE if any of the four attributes are required and supported by the server's file system. For an OPEN with the

EXCLUSIVE4 createmode, the server has no choice, since such OPEN calls do not include the createattrs field. Conversely, if createattrs (UNCHECKED4 or GUARDED4) or cva\_attrs (EXCLUSIVE4\_1) is specified, and includes an owner, owner\_group, or ACE that the principal in the RPC call's credentials does not have authorization to create files for, then the server may return NFS4ERR\_PERM.

In the case of an OPEN that specifies a size of zero (e.g., truncation) and the file has named attributes, the named attributes are left as is and are not removed.

NFSv4.1 gives more precise control to clients over acquisition of delegations via the following new flags for the share\_access field of OPEN4args:

OPEN4\_SHARE\_ACCESS\_WANT\_READ\_DELEG

OPEN4\_SHARE\_ACCESS\_WANT\_WRITE\_DELEG

OPEN4\_SHARE\_ACCESS\_WANT\_ANY\_DELEG

OPEN4\_SHARE\_ACCESS\_WANT\_NO\_DELEG

OPEN4\_SHARE\_ACCESS\_WANT\_CANCEL

OPEN4\_SHARE\_ACCESS\_WANT\_SIGNAL\_DELEG\_WHEN\_RESRC\_AVAIL

OPEN4\_SHARE\_ACCESS\_WANT\_PUSH\_DELEG\_WHEN\_UNCONTENDED

If (share\_access & OPEN4\_SHARE\_ACCESS\_WANT\_DELEG\_MASK) is not zero, then the client will have specified one and only one of:

OPEN4\_SHARE\_ACCESS\_WANT\_READ\_DELEG

OPEN4\_SHARE\_ACCESS\_WANT\_WRITE\_DELEG

OPEN4\_SHARE\_ACCESS\_WANT\_ANY\_DELEG

OPEN4\_SHARE\_ACCESS\_WANT\_NO\_DELEG

OPEN4\_SHARE\_ACCESS\_WANT\_CANCEL

Otherwise, the client is neither indicating a desire nor a non-desire for a delegation, and the server MAY or MAY not return a delegation in the OPEN response.

If the server supports the new `_WANT_` flags and the client sends one or more of the new flags, then in the event the server does not return a delegation, it **MUST** return a delegation type of `OPEN_DELEGATE_NONE_EXT`. The field `ond_why` in the reply indicates why no delegation was returned and will be one of:

`WND4_NOT_WANTED`

The client specified `OPEN4_SHARE_ACCESS_WANT_NO_DELEG`.

`WND4_CONTENTION`

There is a conflicting delegation or open on the file.

`WND4_RESOURCE`

Resource limitations prevent the server from granting a delegation.

`WND4_NOT_SUPP_FTYPE`

The server does not support delegations on this file type.

`WND4_WRITE_DELEG_NOT_SUPP_FTYPE`

The server does not support `OPEN_DELEGATE_WRITE` delegations on this file type.

`WND4_NOT_SUPP_UPGRADE`

The server does not support atomic upgrade of an `OPEN_DELEGATE_READ` delegation to an `OPEN_DELEGATE_WRITE` delegation.

`WND4_NOT_SUPP_DOWNGRADE`

The server does not support atomic downgrade of an `OPEN_DELEGATE_WRITE` delegation to an `OPEN_DELEGATE_READ` delegation.

`WND4_CANCELED`

The client specified `OPEN4_SHARE_ACCESS_WANT_CANCEL` and now any "want" for this file object is cancelled.

`WND4_IS_DIR`

The specified file object is a directory, and the operation is `OPEN` or `WANT_DELEGATION`, which do not support delegations on directories.

`OPEN4_SHARE_ACCESS_WANT_READ_DELEG`,  
`OPEN_SHARE_ACCESS_WANT_WRITE_DELEG`, or  
`OPEN_SHARE_ACCESS_WANT_ANY_DELEG` mean, respectively, the client wants an `OPEN_DELEGATE_READ`, `OPEN_DELEGATE_WRITE`, or any delegation regardless which of `OPEN4_SHARE_ACCESS_READ`, `OPEN4_SHARE_ACCESS_WRITE`, or `OPEN4_SHARE_ACCESS_BOTH` is set. If the

client has an `OPEN_DELEGATE_READ` delegation on a file and requests an `OPEN_DELEGATE_WRITE` delegation, then the client is requesting atomic upgrade of its `OPEN_DELEGATE_READ` delegation to an `OPEN_DELEGATE_WRITE` delegation. If the client has an `OPEN_DELEGATE_WRITE` delegation on a file and requests an `OPEN_DELEGATE_READ` delegation, then the client is requesting atomic downgrade to an `OPEN_DELEGATE_READ` delegation. A server MAY support atomic upgrade or downgrade. If it does, then the returned `delegation_type` of `OPEN_DELEGATE_READ` or `OPEN_DELEGATE_WRITE` that is different from the delegation type the client currently has, indicates successful upgrade or downgrade. If the server does not support atomic delegation upgrade or downgrade, then `ond_why` will be set to `WND4_NOT_SUPP_UPGRADE` or `WND4_NOT_SUPP_DOWNGRADE`.

`OPEN4_SHARE_ACCESS_WANT_NO_DELEG` means that the client wants no delegation.

`OPEN4_SHARE_ACCESS_WANT_CANCEL` means that the client wants no delegation and wants to cancel any previously registered "want" for a delegation.

The client may set one or both of `OPEN4_SHARE_ACCESS_WANT_SIGNAL_DELEG_WHEN_RESRC_AVAIL` and `OPEN4_SHARE_ACCESS_WANT_PUSH_DELEG_WHEN_UNCONTENDED`. However, they will have no effect unless one of following is set:

- \* `OPEN4_SHARE_ACCESS_WANT_READ_DELEG`
- \* `OPEN4_SHARE_ACCESS_WANT_WRITE_DELEG`
- \* `OPEN4_SHARE_ACCESS_WANT_ANY_DELEG`

If the client specifies `OPEN4_SHARE_ACCESS_WANT_SIGNAL_DELEG_WHEN_RESRC_AVAIL`, then it wishes to register a "want" for a delegation, in the event the `OPEN` results do not include a delegation. If so and the server denies the delegation due to insufficient resources, the server MAY later inform the client, via the `CB_RECALLABLE_OBJ_AVAIL` operation, that the resource limitation condition has eased. The server will tell the client that it intends to send a future `CB_RECALLABLE_OBJ_AVAIL` operation by setting `delegation_type` in the results to `OPEN_DELEGATE_NONE_EXT`, `ond_why` to `WND4_RESOURCE`, and `ond_server_will_signal_avail` set to `TRUE`. If `ond_server_will_signal_avail` is set to `TRUE`, the server MUST later send a `CB_RECALLABLE_OBJ_AVAIL` operation.

If the client specifies

OPEN4\_SHARE\_ACCESS\_WANT\_SIGNAL\_DELEG\_WHEN\_UNCONTENDED, then it wishes to register a "want" for a delegation, in the event the OPEN results do not include a delegation. If so and the server denies the delegation due to contention, the server MAY later inform the client, via the CB\_PUSH\_DELEG operation, that the contention condition has eased. The server will tell the client that it intends to send a future CB\_PUSH\_DELEG operation by setting delegation\_type in the results to OPEN\_DELEGATE\_NONE\_EXT, ond\_why to WND4\_CONTENTION, and ond\_server\_will\_push\_deleg to TRUE. If ond\_server\_will\_push\_deleg is TRUE, the server MUST later send a CB\_PUSH\_DELEG operation.

If the client has previously registered a want for a delegation on a file, and then sends a request to register a want for a delegation on the same file, the server MUST return a new error:

NFS4ERR\_DELEG\_ALREADY\_WANTED. If the client wishes to register a different type of delegation want for the same file, it MUST cancel the existing delegation WANT.

#### 25.16.4. IMPLEMENTATION

In absence of a persistent session, the client invokes exclusive create by setting the how parameter to EXCLUSIVE4 or EXCLUSIVE4\_1. In these cases, the client provides a verifier that can reasonably be expected to be unique. A combination of a client identifier, perhaps the client network address, and a unique number generated by the client, perhaps the RPC transaction identifier, may be appropriate.

If the object does not exist, the server creates the object and stores the verifier in stable storage. For file systems that do not provide a mechanism for the storage of arbitrary file attributes, the server may use one or more elements of the object's metadata to store the verifier. The verifier MUST be stored in stable storage to prevent erroneous failure on retransmission of the request. It is assumed that an exclusive create is being performed because exclusive semantics are critical to the application. Because of the expected usage, exclusive CREATE does not rely solely on the server's reply cache for storage of the verifier. A nonpersistent reply cache does not survive a crash and the session and reply cache may be deleted after a network partition that exceeds the lease time, thus opening failure windows.

An NFSv4.1 server should not store the verifier in any of the file's OPTIONAL or REQUIRED attributes. If it does, the server should use time\_modify\_set or time\_access\_set to store the verifier. The server should not store the verifier in the following attributes:

acl (it is desirable for access control to be established at creation),

dacl (ditto),

mode (ditto),

owner (ditto),

owner\_group (ditto),

retentevt\_set (it may be desired to establish retention at creation)

retention\_hold (ditto),

retention\_set (ditto),

sacl (it is desirable for auditing control to be established at creation),

size (on some servers, size may have a limited range of values),

mode\_set\_masked (as with mode),

and

time\_creation (a meaningful file creation should be set when the file is created).

Another alternative for the server is to use a named attribute to store the verifier.

Because the EXCLUSIVE4 create method does not specify initial attributes when processing an EXCLUSIVE4 create, the server

- \* MUST set the owner of the file to that corresponding to the credential of request's RPC header.
- \* MUST NOT allow file access to anyone other than the owner of the file.

If the server cannot support exclusive create semantics, possibly because of the requirement to commit the verifier to stable storage, it should fail the OPEN request with the error NFS4ERR\_NOTSUPP.

During an exclusive CREATE request, if the object already exists, the server reconstructs the object's verifier and compares it with the verifier in the request. If they match, the server treats the request as a success. The request is presumed to be a duplicate of an earlier, successful request for which the reply was lost and that the server duplicate request cache mechanism did not detect. If the verifiers do not match, the request is rejected with the status NFS4ERR\_EXIST.

After the client has performed a successful exclusive create, the attrset response indicates which attributes were used to store the verifier. If EXCLUSIVE4 was used, the attributes set in attrset were used for the verifier. If EXCLUSIVE4\_1 was used, the client determines the attributes used for the verifier by comparing attrset with cva\_attrs.attrmask; any bits set in the former but not the latter identify the attributes used to store the verifier. The client MUST immediately send a SETATTR to set attributes used to store the verifier. Until it does so, the attributes used to store the verifier cannot be relied upon. The subsequent SETATTR MUST NOT occur in the same COMPOUND request as the OPEN.

Unless a persistent session is used, use of the GUARDED4 attribute does not provide exactly once semantics. In particular, if a reply is lost and the server does not detect the retransmission of the request, the operation can fail with NFS4ERR\_EXIST, even though the create was performed successfully. The client would use this behavior in the case that the application has not requested an exclusive create but has asked to have the file truncated when the file is opened. In the case of the client timing out and retransmitting the create request, the client can use GUARDED4 to prevent against a sequence like create, write, create (retransmitted) from occurring.

For SHARE reservations, the value of the expression (share\_access & ~OPEN4\_SHARE\_ACCESS\_WANT\_DELEG\_MASK) MUST be one of OPEN4\_SHARE\_ACCESS\_READ, OPEN4\_SHARE\_ACCESS\_WRITE, or OPEN4\_SHARE\_ACCESS\_BOTH. If not, the server MUST return NFS4ERR\_INVALID. The value of share\_deny MUST be one of OPEN4\_SHARE\_DENY\_NONE, OPEN4\_SHARE\_DENY\_READ, OPEN4\_SHARE\_DENY\_WRITE, or OPEN4\_SHARE\_DENY\_BOTH. If not, the server MUST return NFS4ERR\_INVALID.

Based on the share\_access value (OPEN4\_SHARE\_ACCESS\_READ, OPEN4\_SHARE\_ACCESS\_WRITE, or OPEN4\_SHARE\_ACCESS\_BOTH), the client should check that the requester has the proper access rights to perform the specified operation. This would generally be the results of applying the ACL access rules to the file for the current requester. However, just as with the ACCESS operation, the client

should not attempt to second-guess the server's decisions, as access rights may change and may be subject to server administrative controls outside the ACL framework. If the requester's READ or WRITE operation is not authorized (depending on the share\_access value), the server MUST return NFS4ERR\_ACCESS.

Note that if the client ID was not created with the EXCHGID4\_FLAG\_BIND\_PRINC\_STATEID capability set in the reply to EXCHANGE\_ID, then the server MUST NOT impose any requirement that READs and WRITEs sent for an open file have the same credentials as the OPEN itself, and the server is REQUIRED to perform access checking on the READs and WRITEs themselves. Otherwise, if the reply to EXCHANGE\_ID did have EXCHGID4\_FLAG\_BIND\_PRINC\_STATEID set, then with one exception, the credentials used in the OPEN request MUST match those used in the READs and WRITEs, and the stateids in the READs and WRITEs MUST match, or be derived from the stateid from the reply to OPEN. The exception is if SP4\_SSV or SP4\_MACH\_CRED state protection is used, and the spo\_must\_allow result of EXCHANGE\_ID includes the READ and/or WRITE operations. In that case, the machine or SSV credential will be allowed to send READ and/or WRITE. See Section 25.35.

If the component provided to OPEN is a symbolic link, the error NFS4ERR\_SYMLINK will be returned to the client, while if it is a directory the error NFS4ERR\_ISDIR will be returned. If the component is neither of those but not an ordinary file, the error NFS4ERR\_WRONG\_TYPE is returned. If the current filehandle is not a directory, the error NFS4ERR\_NOTDIR will be returned.

The use of the OPEN4\_RESULT\_PRESERVE\_UNLINKED result flag allows a client to avoid the common implementation practice of renaming an open file to ".nfs<unique value>" instead of removing the file. After the server returns OPEN4\_RESULT\_PRESERVE\_UNLINKED for all NFsv4.1 OPENS and there are no OPENS issued by other protocols to deal with, if a client sends a REMOVE operation that would reduce the file's link count to zero, the server will report a value of zero for the numlinks attribute on the file, when GETATTR can be done on this file.



If another client has a delegation of the file being opened that conflicts with open being done (sometimes depending on the `share_access` or `share_deny` value specified), the delegation(s) MUST be recalled, and the operation cannot proceed until each such delegation is returned or revoked. Except where this happens very quickly, one or more `NFS4ERR_DELAY` errors will be returned to requests made while delegation remains outstanding. In the case of an `OPEN_DELEGATE_WRITE` delegation, any open by a different client will conflict, while for an `OPEN_DELEGATE_READ` delegation, only opens with one of the following characteristics will be considered conflicting:

- \* The value of `share_access` includes the bit `OPEN4_SHARE_ACCESS_WRITE`.
- \* The value of `share_deny` specifies `OPEN4_SHARE_DENY_READ` or `OPEN4_SHARE_DENY_BOTH`.
- \* `OPEN4_CREATE` is specified together with `UNCHECKED4`, the size attribute is specified as zero (for truncation), and an existing file is truncated.

If `OPEN4_CREATE` is specified and the file does not exist and the current filehandle designates a directory for which another client holds a directory delegation, then, unless the delegation is such that the situation can be resolved by sending a notification, the delegation MUST be recalled, and the operation cannot proceed until the delegation is returned or revoked. Except where this happens very quickly, one or more `NFS4ERR_DELAY` errors will be returned to requests made while delegation remains outstanding.

If `OPEN4_CREATE` is specified and the file does not exist and the current filehandle designates a directory for which one or more directory delegations exist, then, when those delegations request such notifications, `NOTIFY4_ADD_ENTRY` will be generated as a result of this operation.

#### 25.16.4.1. Warning to Client implementers

`OPEN` resembles `LOOKUP` in that it generates a filehandle for the client to use. Unlike `LOOKUP` though, `OPEN` creates server state on the filehandle. In normal circumstances, the client can only release this state with a `CLOSE` operation. `CLOSE` uses the current filehandle to determine which file to close. Therefore, the client MUST follow every `OPEN` operation with a `GETFH` operation in the same `COMPOUND` procedure. This will supply the client with the filehandle such that `CLOSE` can be used appropriately.

Simply waiting for the lease on the file to expire is insufficient because the server may maintain the state indefinitely as long as another client does not attempt to make a conflicting access to the same file.

See also Section 7.6.4.

## 25.17. Operation 19: OPENATTR - Open Named Attribute Directory

### 25.17.1. ARGUMENTS

```
struct OPENATTR4args {
    /* CURRENT_FH: object */
    bool    createdir;
};
```

### 25.17.2. RESULTS

```
struct OPENATTR4res {
    /*
     * If status is NFS4_OK,
     *   new CURRENT_FH: named attribute
     *                   directory
     */
    nfsstat4    status;
};
```

### 25.17.3. DESCRIPTION

The OPENATTR operation is used to obtain the filehandle of the named attribute directory associated with the current filehandle. The result of the OPENATTR will be a filehandle to an object of type NF4ATTRDIR. From this filehandle, READDIR and LOOKUP operations can be used to obtain filehandles for the various named attributes associated with the original file system object. Filehandles returned within the named attribute directory will designate objects of type of NF4NAMEDATTR.

The createdir argument allows the client to signify if a named attribute directory should be created as a result of the OPENATTR operation. Some clients may use the OPENATTR operation with a value of FALSE for createdir to determine if any named attributes exist for the object. If none exist, then NFS4ERR\_NOENT will be returned. If createdir has a value of TRUE and no named attribute directory exists, one is created and its filehandle becomes the current filehandle. On the other hand, if createdir has a value of TRUE and the named attribute directory already exists, no error results and the filehandle of the existing directory becomes the current

filehandle. The creation of a named attribute directory assumes that the server has implemented named attribute support in this fashion and is not required to do so by this definition.

If the current filehandle designates an object of type NF4NAMEDATTR (a named attribute) or NF4ATTRDIR (a named attribute directory), an error of NFS4ERR\_WRONG\_TYPE is returned to the client. Named attributes or a named attribute directory MUST NOT have their own named attributes.

#### 25.17.4. IMPLEMENTATION

If the server does not support named attributes for file system objects on the file system associated with the current filehandle, an error of NFS4ERR\_NOTSUPP will be returned to the client.

#### 25.18. Operation 21: OPEN\_DOWNGRADE - Reduce Open File Access

##### 25.18.1. ARGUMENTS

```
struct OPEN_DOWNGRADE4args {
    /* CURRENT_FH: opened file */
    stateid4      open_stateid;
    seqid4        seqid;
    uint32_t      share_access;
    uint32_t      share_deny;
};
```

##### 25.18.2. RESULTS

```
struct OPEN_DOWNGRADE4resok {
    stateid4      open_stateid;
};

union OPEN_DOWNGRADE4res switch(nfsstat4 status) {
    case NFS4_OK:
        OPEN_DOWNGRADE4resok    resok4;
    default:
        void;
};
```

## 25.18.3. DESCRIPTION

This operation is used to adjust the access and deny states for a given open. This is necessary when a given open-owner opens the same file multiple times with different access and deny values. In this situation, a close of one of the opens may change the appropriate share\_access and share\_deny flags to remove bits associated with opens no longer in effect.

Valid values for the expression (share\_access & ~OPEN4\_SHARE\_ACCESS\_WANT\_DELEG\_MASK) are OPEN4\_SHARE\_ACCESS\_READ, OPEN4\_SHARE\_ACCESS\_WRITE, or OPEN4\_SHARE\_ACCESS\_BOTH. If the client specifies other values, the server MUST reply with NFS4ERR\_INVALID.

Valid values for the share\_deny field are OPEN4\_SHARE\_DENY\_NONE, OPEN4\_SHARE\_DENY\_READ, OPEN4\_SHARE\_DENY\_WRITE, or OPEN4\_SHARE\_DENY\_BOTH. If the client specifies other values, the server MUST reply with NFS4ERR\_INVALID.

After checking for valid values of share\_access and share\_deny, the server replaces the current access and deny modes on the file with share\_access and share\_deny subject to the following constraints:

- \* The bits in share\_access should equal the union of the share\_access bits (not including OPEN4\_SHARE\_WANT\_\* bits) specified for some subset of the OPENS in effect for the current open-owner on the current file.
- \* The bits in share\_deny should equal the union of the share\_deny bits specified for some subset of the OPENS in effect for the current open-owner on the current file.

If the above constraints are not respected, the server should return the error NFS4ERR\_INVALID. Since share\_access and share\_deny bits should be subsets of those already granted, short of a defect in the client or server implementation, it is not possible for the OPEN\_DOWNGRADE request to be denied because of conflicting share reservations.

The seqid argument is not used in NFSv4.1, MAY be any value, and MUST be ignored by the server.

On success, the current filehandle retains its value.

#### 25.18.4. IMPLEMENTATION

An OPEN\_DOWNGRADE operation may make OPEN\_DELEGATE\_READ delegations grantable where they were not previously. Servers may choose to respond immediately if there are pending delegation want requests or may respond to the situation at a later time.

#### 25.19. Operation 22: PUTFH - Set Current Filehandle

##### 25.19.1. ARGUMENTS

```
struct PUTFH4args {  
    nfs_fh4      object;  
};
```

##### 25.19.2. RESULTS

```
struct PUTFH4res {  
    /*  
     * If status is NFS4_OK,  
     *   new CURRENT_FH: argument to PUTFH  
     */  
    nfsstat4      status;  
};
```

##### 25.19.3. DESCRIPTION

This operation replaces the current filehandle with the filehandle provided as an argument. It clears the current stateid.

If the security mechanism used by the requester does not meet the requirements of the filehandle provided to this operation, the server MUST return NFS4ERR\_WRONGSEC.

See Section 23.2.3.1.1 for more details on the current filehandle.

See Section 23.2.3.1.2 for more details on the current stateid.

##### 25.19.4. IMPLEMENTATION

This operation is used in an NFS request to set the context for file accessing operations that follow in the same COMPOUND request.

## 25.20. Operation 23: PUTPUBFH - Set Public Filehandle

### 25.20.1. ARGUMENT

```
void;
```

### 25.20.2. RESULT

```
struct PUTPUBFH4res {  
    /*  
     * If status is NFS4_OK,  
     *   new CURRENT_FH: public fh  
     */  
    nfsstat4      status;  
};
```

### 25.20.3. DESCRIPTION

This operation replaces the current filehandle with the filehandle that represents the public filehandle of the server's namespace. This filehandle may be different from the "root" filehandle that may be associated with some other directory on the server.

PUTPUBFH also clears the current stateid.

The public filehandle represents the concepts embodied in [RFC2054], [RFC2055], and [RFC2224]. The intent for NFSv4.1 is that the public filehandle (represented by the PUTPUBFH operation) be used as a method of providing WebNFS server compatibility with NFSv3.

The public filehandle and the root filehandle (represented by the PUTROOTFH operation) should be equivalent. If the public and root filehandles are not equivalent, then the directory corresponding to the public filehandle MUST be a descendant of the directory corresponding to the root filehandle.

See Section 23.2.3.1.1 for more details on the current filehandle.

See Section 23.2.3.1.2 for more details on the current stateid.

### 25.20.4. IMPLEMENTATION

This operation is used in an NFS request to set the context for file accessing operations that follow in the same COMPOUND request.

With the NFSv3 public filehandle, the client is able to specify whether the pathname provided in the LOOKUP should be evaluated as either an absolute path relative to the server's root or relative to

the public filehandle. [RFC2224] contains further discussion of the functionality. With NFSv4.1, that type of specification is not directly available in the LOOKUP operation. The reason for this is because the component separators needed to specify absolute vs. relative are not allowed in NFSv4. Therefore, the client is responsible for constructing its request such that the use of either PUTROOTFH or PUTPUBFH signifies absolute or relative evaluation of an NFS URL, respectively.

Note that there are warnings mentioned in [RFC2224] with respect to the use of absolute evaluation and the restrictions the server may place on that evaluation with respect to how much of its namespace has been made available. These same warnings apply to NFSv4.1. It is likely, therefore, that because of server implementation details, an NFSv3 absolute public filehandle look up may behave differently than an NFSv4.1 absolute resolution.

There is a form of security negotiation as described in [RFC2755] that uses the public filehandle and an overloading of the pathname. This method is not available with NFSv4.1 as filehandles are not overloaded with special meaning and therefore do not provide the same framework as NFSv3. Clients should therefore use the security negotiation mechanisms described in Section 12 [To be Updated] of the NFSv4-wide security document, currently

## 25.21. Operation 24: PUTROOTFH - Set Root Filehandle

### 25.21.1. ARGUMENTS

void;

### 25.21.2. RESULTS

```
struct PUTROOTFH4res {
    /*
     * If status is NFS4_OK,
     *   new CURRENT_FH: root fh
     */
    nfsstat4      status;
};
```

### 25.21.3. DESCRIPTION

This operation replaces the current filehandle with the filehandle that represents the root of the server's namespace. From this filehandle, a LOOKUP operation can locate any other filehandle on the server. This filehandle may be different from the "public" filehandle that may be associated with some other directory on the server.

PUTROOTFH also clears the current stateid.

See Section 23.2.3.1.1 for more details on the current filehandle.

See Section 23.2.3.1.2 for more details on the current stateid.

### 25.21.4. IMPLEMENTATION

This operation is used in an NFS request to set the context for file accessing operations that follow in the same COMPOUND request.

## 25.22. Operation 25: READ - Read from File

### 25.22.1. ARGUMENTS

```
struct READ4args {
    /* CURRENT_FH: file */
    stateid4      stateid;
    offset4       offset;
    count4        count;
};
```

### 25.22.2. RESULTS

```
struct READ4resok {
    bool          eof;
    opaque        data<>;
};

union READ4res switch (nfsstat4 status) {
    case NFS4_OK:
        READ4resok      resok4;
    default:
        void;
};
```



### 25.22.3. DESCRIPTION

The READ operation reads data from the regular file identified by the current filehandle.

The client provides an offset of where the READ is to start and a count of how many bytes are to be read. An offset of zero means to read data starting at the beginning of the file. If offset is greater than or equal to the size of the file, the status NFS4\_OK is returned with a data length set to zero and eof is set to TRUE. The READ is subject to access permissions checking.

If the client specifies a count value of zero, the READ succeeds and returns zero bytes of data again subject to access permissions checking. The server may choose to return fewer bytes than specified by the client. The client needs to check for this condition and handle the condition appropriately.

Except when special stateids are used, the stateid value for a READ request represents a value returned from a previous byte-range lock or share reservation request or the stateid associated with a delegation. The stateid identifies the associated owners if any and is used by the server to verify that the associated locks are still valid (e.g., have not been revoked).

If the read ended at the end-of-file (formally, in a correctly formed READ operation, if offset + count is equal to the size of the file), or the READ operation extends beyond the size of the file (if offset + count is greater than the size of the file), eof is returned as TRUE; otherwise, it is FALSE. A successful READ of an empty file will always return eof as TRUE.

If the current filehandle is not an ordinary file, an error will be returned to the client. In the case that the current filehandle represents an object of type NF4DIR, NFS4ERR\_ISDIR is returned. If the current filehandle designates a symbolic link, NFS4ERR\_SYMLINK is returned. In all other cases, NFS4ERR\_WRONG\_TYPE is returned.

For a READ with a stateid value of all bits equal to zero, the server MAY allow the READ to be serviced subject to mandatory byte-range locks or the current share deny modes for the file. For a READ with a stateid value of all bits equal to one, the server MAY allow READ operations to bypass locking checks at the server.

On success, the current filehandle retains its value.

#### 25.22.4. IMPLEMENTATION

If the server returns a "short read" (i.e., fewer data than requested and eof is set to FALSE), the client should send another READ to get the remaining data. A server may return less data than requested under several circumstances. The file may have been truncated by another client or perhaps on the server itself, changing the file size from what the requesting client believes to be the case. This would reduce the actual amount of data available to the client. It is possible that the server reduce the transfer size and so return a short read result. Server resource exhaustion may also occur in a short read.

If mandatory byte-range locking is in effect for the file, and if the byte-range corresponding to the data to be read from the file is WRITE\_LT locked by an owner not associated with the stateid, the server will return the NFS4ERR\_LOCKED error. The client should try to get the appropriate READ\_LT via the LOCK operation before re-attempting the READ. When the READ completes, the client should release the byte-range lock via LOCKU.

If another client has an OPEN\_DELEGATE\_WRITE delegation for the file being read, the delegation must be recalled, and the operation cannot proceed until that delegation is returned or revoked. Except where this happens very quickly, one or more NFS4ERR\_DELAY errors will be returned to requests made while the delegation remains outstanding. Normally, delegations will not be recalled as a result of a READ operation since the recall will occur as a result of an earlier OPEN. However, since it is possible for a READ to be done with a special stateid, the server needs to check for this case even though the client should have done an OPEN previously.

#### 25.23. Operation 26: READDIR - Read Directory

##### 25.23.1. ARGUMENTS

```
struct READDIR4args {
    /* CURRENT_FH: directory */
    nfs_cookie4      cookie;
    verifier4        cookieverf;
    count4           dircount;
    count4           maxcount;
    bitmap4          attr_request;
};
```

##### 25.23.2. RESULTS

```
struct entry4 {
    nfs_cookie4    cookie;
    component4     name;
    fattr4         attrs;
    entry4         *nextentry;
};

struct dirlist4 {
    entry4         *entries;
    bool           eof;
};

struct READDIR4resok {
    verifier4      cookieverf;
    dirlist4       reply;
};

union READDIR4res switch (nfsstat4 status) {
    case NFS4_OK:
        READDIR4resok  resok4;
    default:
        void;
};
```

### 25.23.3. DESCRIPTION

The READDIR operation retrieves a variable number of entries from a file system directory and returns client-requested attributes for each entry along with information to allow the client to request additional directory entries in a subsequent READDIR.

The arguments contain a cookie value that represents where the READDIR should start within the directory. A value of zero for the cookie is used to start reading at the beginning of the directory. For subsequent READDIR requests, the client specifies a cookie value that is provided by the server on a previous READDIR request.

The request's cookieverf field should be set to 0 (zero) when the request's cookie field is zero (first read of the directory). On subsequent requests, the cookieverf field must match the cookieverf returned by the READDIR in which the cookie was acquired. If the server determines that the cookieverf is no longer valid for the directory, the error NFS4ERR\_NOT\_SAME must be returned.

The dircount field of the request is a hint of the maximum number of bytes of directory information that should be returned. This value represents the total length of the names of the directory entries and

the cookie value for these entries. This length represents the XDR encoding of the data (names and cookies) and not the length in the native format of the server.

The maxcount field of the request represents the maximum total size of all of the data being returned within the READDIR4resok structure and includes the XDR overhead. The server MAY return less data. If the server is unable to return a single directory entry within the maxcount limit, the error NFS4ERR\_TOOSMALL MUST be returned to the client.

Finally, the request's attr\_request field represents the set of attributes to be returned for each directory entry supplied by the server. As in the case of GETATTR, if this set includes unsupported attributes, they are not included in the returned data and no error results. Because of the possibility of mount points within a directory, different sets of attributes might be supported for different entries since they might be parts of distinct file systems.

A successful reply consists of a list of directory entries. Each of these entries contains the name of the directory entry, a cookie value for that entry, and the associated attributes as requested. The "eof" flag has a value of TRUE if there are no more entries in the directory.

The cookie value is only meaningful to the server and is used as a cursor for the directory entry. As mentioned, this cookie is used by the client for subsequent READDIR operations so that it may continue reading a directory. The cookie is similar in concept to a READ offset but MUST NOT be interpreted as such by the client. Ideally, the cookie value should not change if the directory is modified since the client may be caching these values.

In some cases, the server may encounter an error while obtaining the attributes for a directory entry. Instead of returning an error for the entire READDIR operation, the server can instead return the attribute rgetattr\_error (Section 11.12.1.12). With this, the server is able to communicate the failure to the client and not fail the entire operation in the instance of what might be a transient failure. Obviously, the client must request the fattr4\_rgetattr\_error attribute for this method to work properly. If the client does not request the attribute, the server has no choice but to return failure for the entire READDIR operation.

For some file system environments, the directory entries "." and ".." have special meaning, and in other environments, they do not. If the server supports these special entries within a directory, they should not be returned to the client as part of the READDIR response. To

enable some client environments, the cookie values of zero, 1, and 2 are to be considered reserved. Note that the UNIX client will use these values when combining the server's response and local representations to enable a fully formed UNIX directory presentation to the application.

For READDIR arguments, cookie values of one and two should not be used, and for READDIR results, cookie values of zero, one, and two should not be returned.

On success, the current filehandle retains its value.

#### 25.23.4. IMPLEMENTATION

The server's file system directory representations can differ greatly. A client's programming interfaces may also be bound to the local operating environment in a way that does not translate well into the NFS protocol. Therefore, the use of the `dircount` and `maxcount` fields are provided to enable the client to provide hints to the server. If the client is aggressive about attribute collection during a READDIR, the server has an idea of how to limit the encoded response.

If `dircount` is zero, the server bounds the reply's size based on the request's `maxcount` field.

The `cookieverf` may be used by the server to help manage cookie values that may become stale. It should be a rare occurrence that a server is unable to continue properly reading a directory with the provided cookie/`cookieverf` pair. The server needs to make every effort to avoid this condition since the application at the client might be unable to properly handle this type of failure.

The use of the `cookieverf` will also protect the client from using READDIR cookie values that might be stale. For example, if the file system has been migrated, the server might or might not be able to use the same cookie values to service READDIR as the previous server used. With the client providing the `cookieverf`, the server is able to provide the appropriate response to the client. This prevents the case where the server accepts a cookie value but the underlying directory has changed and the response is invalid from the client's context of its previous READDIR.

Since some servers will not be returning "." and ".." entries as has been done with previous versions of the NFS protocol, the client that requires these entries be present in READDIR responses must fabricate them.

## 25.24. Operation 27: READLINK - Read Symbolic Link

### 25.24.1. ARGUMENTS

```
/* CURRENT_FH: symlink */  
void;
```

### 25.24.2. RESULTS

```
struct READLINK4resok {  
    linktext4    link;  
};  
  
union READLINK4res switch (nfsstat4 status) {  
    case NFS4_OK:  
        READLINK4resok resok4;  
    default:  
        void;  
};
```

### 25.24.3. DESCRIPTION

READLINK reads the data associated with a symbolic link. Depending on the value of the UTF-8 capability attribute (Section 21.1), the data is encoded in UTF-8. Whether created by an NFS client or created locally on the server, the data in a symbolic link is not interpreted (except possibly to check for proper UTF-8 encoding) when created, but is simply stored.

On success, the current filehandle retains its value.

### 25.24.4. IMPLEMENTATION

A symbolic link is nominally a pointer to another file. The data is not necessarily interpreted by the server, just stored in the file. It is possible for a client implementation to store a pathname that is not meaningful to the server operating system in a symbolic link. A READLINK operation returns the data to the client for interpretation. If different implementations want to share access to symbolic links, then they must agree on the interpretation of the data in the symbolic link.

The READLINK operation is only allowed on objects of type NF4LNK. The server should return the error NFS4ERR\_WRONG\_TYPE if the object is not of type NF4LNK.

## 25.25. Operation 28: REMOVE - Remove File System Object

## 25.25.1. ARGUMENTS

```
struct REMOVE4args {  
    /* CURRENT_FH: directory */  
    component4      target;  
};
```

## 25.25.2. RESULTS

```
struct REMOVE4resok {  
    change_info4    cinfo;  
};  
  
union REMOVE4res switch (nfsstat4 status) {  
    case NFS4_OK:  
        REMOVE4resok    resok4;  
    default:  
        void;  
};
```

## 25.25.3. DESCRIPTION

The REMOVE operation removes (deletes) a directory entry which names a file system object from the directory corresponding to the current filehandle. If the entry in the directory was the last reference to the (i.e., there are no other links to that object), the specified object may be destroyed. In addition, as discussed below, the destruction of the object can be delayed by its use as an open file. The directory may be either of type NF4DIR or NF4ATTRDIR.

For the directory where the filename was removed, the server returns change\_info4 information in cinfo. With the atomic field of the change\_info4 data type, the server will indicate if the before and after change attributes were obtained atomically with respect to the removal.

If the target has a length of zero, or if the target does not obey the UTF-8 definition (and the server is enforcing UTF-8 encoding; see Section 21.1), the error NFS4ERR\_INVALID will be returned.

On success, the current filehandle retains its value.

## 25.25.4. IMPLEMENTATION

In two important respects, the REMOVE operation within NFSv4.1 differs from remove operations for earlier versions of NFS:

- \* NFSv3 required a different operator RMDIR for directory removal and together with REMOVE for non-directory removal. This allowed clients to skip checking the file type when being passed a non-directory delete system call (e.g., unlink() [unlink] in POSIX) to remove a directory, as well as the converse (e.g., a rmdir() on a non-directory) because they knew the server would check the file type. NFSv4.1 REMOVE can be used to delete any directory entry irrespective of its file type.

The implementer of an NFSv4.1 client's entry points from the unlink() and rmdir() system calls should first check the file type against the types the system call is allowed to remove before sending a REMOVE operation. Alternatively, the implementer can produce a COMPOUND call that includes a LOOKUP/VERIFY sequence of operations to verify the file type before a REMOVE operation in the same COMPOUND call.

- \* In order to deal with removal of open files in a manner consistent with local file system semantics, the server has the option of returning the flag OPEN4\_RESULT\_PRESERVE\_UNLINKED, to indicate to the client that the file will be preserved as long as it has an outstanding NFSv4.1 open (See Section 25.16) that returned this flag.

Regardless of the state of OPEN4\_RESULT\_PRESERVE\_UNLINKED, which controls the continued existence of the object to be deleted, it is unwise for the client to rely on the availability of disk space due to the REMOVE. This is because server file space allocation policies may differ.

If there is an OPEN preventing writing the file (i.e one specifying OPEN4\_SHARE\_DENY\_WRITE or OPEN4\_SHARE\_DENY\_BOTH) the server MUST reject the removal operation (i.e a REMOVE or a removal that occurs as part of RENAME when a file is renamed-over).

In addition, the server MAY implement its own restrictions on removal of a file while it is open. The server might disallow such a removal operation. The conditions that influence the restrictions on removal of a file while it is still open include:

- \* Whether certain access protocols (i.e., those other than NFS) are holding the file open.
- \* Whether particular options or policies on the server are enabled.



If a file has an outstanding OPEN and this prevents the removal of the file's directory entry, the error NFS4ERR\_FILE\_OPEN is returned. This applies to both rejection because of an OPEN preventing writing the file and to that due to the server's own policies.

The case of a removal operation being done by a client that holds a write delegation presents important issues that allows the server to proceed to process the remove without the delegation being recalled, since the client is presumed aware of the existence of OPENS that were sent by it previously or never sent to the server because the delegation holder can process OPENS on its own. Although the server has considerable freedom in determining the OPENS that might prevent a successful removal operation, the server may omit the recall in this case, unless the server is aware of OPENS performed by other access protocols or because options or policies are in effect that might prevent the removal operation independent of the existence of OPENS with deny modes specified.

To deal efficiently with the common case of servers whose policies are such that only the potential existence of OPENS denying READ and/or WRITE, such situations can be dealt with as follows:

- \* The client doing the removal operation and holding the delegation needs to make the server aware of the existence of such OPENS without necessarily returning the delegations.
- \* The server, using the knowledge that he has about any such OPENS proceeds to make its decision using only the OPENS it is aware of, without recalling the delegation.

In cases in which the recall can be dispensed with:

- \* If the removal operation succeeds, the server may treat the associate write delegation as effectively canceled just as if it had recalled and returned.
- \* If the removal operation returns successfully, the client can simply forget about the delegation, as if it had been returned.
- \* If the removal operation returns successfully, the client can client can avoid flushing pending data to be written to the file being removed.
- \* If the removal operation returns successfully, the client can simply forget about opens the server is unaware of/ These open cannot have any active deny modes specified.

Where the determination above cannot be made definitively because delegations are being held by other clients they MUST be recalled to allow processing of the removal operation to continue. When such a delegation is held, the server has no reliable knowledge of the status of OPENS for that client, so unless there are files opened with the particular deny modes by clients without delegations other than the client doing the removal operation, the determination cannot be made until all such delegations are recalled, and the operation cannot proceed until each sufficient delegation has been returned or revoked to allow the server to make a correct determination.

If the current filehandle designates a directory for which another client holds a directory delegation, then, unless the situation can be resolved by sending a notification, the directory delegation MUST be recalled, and the operation MUST NOT proceed until the delegation is returned or revoked. Except where this happens very quickly, one or more NFS4ERR\_DELAY errors will be returned to requests made while delegation remains outstanding.

When the current filehandle designates a directory for which one or more directory delegations exist, then, when those delegations request such notifications, NOTIFY4\_REMOVE\_ENTRY will be generated as a result of this operation.

Note that when a remove occurs as a result of a RENAME, NOTIFY4\_REMOVE\_ENTRY will only be generated if the removal happens as a separate operation. In the case in which the removal is integrated and atomic with RENAME, the notification of the removal is integrated with notification for the RENAME. See the discussion of the NOTIFY4\_RENAME\_ENTRY notification in Section 27.4.

In all cases in which delegations are recalled, the server is likely to return one or more NFS4ERR\_DELAY errors while delegations remain outstanding.

The handling of the REMOVE operation, when it not rejected as described above involves the following activities:

- A): The elimination of the directory entry identified by the REMOVE parameters.
- B): The reduction of the link count associated with the file being removed.
- C): The deletion of the file data and its eventual re-use to accommodate newly-written data.

This last item need only be done when the link count decremented as part of item B reaches the value zero. In addition, this last state transition, will, under certain circumstances. be deferred as long as the file is known by the server to be open. See below for details.

Items A and B are done in sequence but there is no atomicity requirement so that other request may see A having been done without B occurring. The execution of C is often delayed until the link count reaches zero and, in many cases, until the file is no longer open.

- \* If the reply from the OPEN had the flag OPEN4\_RESULT\_PRESERVE\_UNLINKED set, the server is obligated to maintain access to the removed object (using a filehandle) until the last OPEN is closed.

This obligation continues across reboots and grace periods, so the file is preserved through the grace period and only considered closed, if it is not reclaimed during the grace period.

When all of the directory entries within a directory are deleted, it is subject to deletion itself, despite the fact that there still might be files actively used by their filehandles, even though they were once referred by directory entries since removed.

- \* If a client does not have support for the OPEN4\_RESULT\_PRESERVE\_UNLINKED flag, it will ignore the value and behave as if it were not set.
- \* If the reply from the OPEN did not have the flag OPEN4\_RESULT\_PRESERVE\_UNLINKED set, the client has the option, as it did in NFSv3 and NFSv4.0, of renaming the file instead of removing it (referred to as "silly rename")

## 25.26. Operation 29: RENAME - Rename Directory Entry

### 25.26.1. ARGUMENTS

```
struct RENAME4args {  
    /* SAVED_FH: source directory */  
    component4      oldname;  
    /* CURRENT_FH: target directory */  
    component4      newname;  
};
```

### 25.26.2. RESULTS

```
struct RENAME4resok {
    change_info4    source_cinfo;
    change_info4    target_cinfo;
};

union RENAME4res switch (nfsstat4 status) {
    case NFS4_OK:
        RENAME4resok    resok4;
    default:
        void;
};
```

### 25.26.3. DESCRIPTION

The RENAME operation renames the object identified by oldname in the source directory corresponding to the saved filehandle, as set by the SAVEFH operation, to newname in the target directory corresponding to the current filehandle. The operation is required to be atomic to the client. Source and target directories MUST reside on the same file system on the server. On success, the current filehandle will continue to be the target directory.

If the target directory already contains an entry with the name newname, the source object MUST be compatible with the target: either both are non-directories or both are directories and the target MUST be empty. If compatible, the existing target is removed before the rename occurs or, preferably, the target is removed atomically as part of the rename. See Section 25.25.4 for client and server actions whenever a target is removed. Note however that when the removal is performed atomically with the rename, certain parts of the removal described there are integrated with the rename. For example, notification of the removal will not be via a NOTIFY4\_REMOVE\_ENTRY but will be indicated as part of the NOTIFY4\_ADD\_ENTRY or NOTIFY4\_RENAME\_ENTRY generated by the rename.

If the source object and the target are not compatible or if the target is a directory but not empty, the server will return the error NFS4ERR\_EXIST.

If oldname and newname both refer to the same file (e.g., they might be hard links of each other), then, unless the file is open (see Section 23.26.4), RENAME MUST perform no action and return NFS4\_OK.

For both directories involved in the RENAME, the server returns change\_info4 information. With the atomic field of the change\_info4 data type, the server will indicate if the before and after change attributes were obtained atomically with respect to the rename.

If oldname refers to a named attribute and the saved and current filehandles refer to different file system objects, the server will return NFS4ERR\_XDEV just as if the saved and current filehandles represented directories on different file systems.

If oldname or newname has a length of zero, or if oldname or newname does not obey the UTF-8 definition in the case of a Unicode-aware file system, the error NFS4ERR\_INVALID will be returned.

#### 25.26.4. IMPLEMENTATION

The server MAY impose restrictions on the RENAME operation such that RENAME may not be done when the file being renamed is open or when that open is done by particular protocols, or with particular options or access modes. Similar restrictions may be applied when a file exists with the target name and is open. When RENAME is rejected because of either of the above restrictions, the error NFS4ERR\_FILE\_OPEN is returned.

When oldname and rename refer to the same file and that file is open in a fashion such that RENAME would normally be rejected with NFS4ERR\_FILE\_OPEN if oldname and newname were different files, then RENAME is to be rejected with NFS4ERR\_FILE\_OPEN.

When the restrictions regarding open files apply to an open file with the target name and the client has a write delegation for that file, delegation recalls can be avoided in the same fashion as described in Section 25.25.4.

If a server does implement such restrictions and those restrictions include cases of NFSv4 opens preventing successful execution of a rename, the server needs to recall any delegations that could hide the existence of opens relevant to that decision. This is because when a client holds a delegation, the server might not have an accurate account of the opens for that client, since the client may execute OPENS and CLOSEs locally. The RENAME operation need only be delayed until a definitive result can be obtained. For example, if there are multiple delegations and one of them establishes an open whose presence would prevent the rename, given the server's semantics, NFS4ERR\_FILE\_OPEN may be returned to the caller as soon as that delegation is returned without waiting for other delegations to be returned. Similarly, if such opens are not associated with delegations, NFS4ERR\_FILE\_OPEN can be returned immediately with no delegation recall being done.

If the current filehandle or the saved filehandle designates a directory for which another client holds a directory delegation, then, unless the situation can be resolved by sending a notification,

the delegation MUST be recalled, and the operation cannot proceed until the delegation is returned or revoked. Except where this happens very quickly, one or more NFS4ERR\_DELAY errors will be returned to requests made while delegation remains outstanding.

When the current and saved filehandles are the same and they designate a directory for which one or more directory delegations exist, then, when those delegations request such notifications, a notification of type NOTIFY4\_RENAME\_ENTRY will be generated as a result of this operation. When oldname and rename refer to the same file, no notification is generated (because, as Section 25.26.3 states, the server MUST take no action). When a file is removed because it has the same name as the target, if that removal is done atomically with the rename, a NOTIFY4\_REMOVE\_ENTRY notification will not be generated. Instead, the deletion of the file will be reported as part of the NOTIFY4\_RENAME\_ENTRY notification.

When the current and saved filehandles are not the same:

- \* If the current filehandle designates a directory for which one or more directory delegations exist, then, when those delegations request such notifications, NOTIFY4\_ADD\_ENTRY will be generated as a result of this operation. When a file is removed because it has the same name as the target, if that removal is done atomically with the rename, a NOTIFY4\_REMOVE\_ENTRY notification will not be generated. Instead, the deletion of the file will be reported as part of the NOTIFY4\_ADD\_ENTRY notification.
- \* If the saved filehandle designates a directory for which one or more directory delegations exist, then, when those delegations request such notifications, NOTIFY4\_REMOVE\_ENTRY will be generated as a result of this operation.

If the object being renamed has file delegations held by clients other than the one doing the RENAME, the delegations MUST be recalled, and the operation cannot proceed until each such delegation is returned or revoked. Note that in the case of multiply linked files, the delegation recall requirement applies even if the delegation was obtained through a different name than the one being renamed. In all cases in which delegations are recalled, the server is likely to return one or more NFS4ERR\_DELAY errors while the delegation(s) remains outstanding, although it might not do that if the delegations are returned quickly.

The direct modification resulting from the RENAME operation must be atomic to the client. However, any necessary changes to the link count attributes for the file involved and the eventual deletion of the file's data in the case of a file deleted because it is renamed

over need not be atomic. In addition, the delay of deletion until last close functions in the renamed over case behaves as indicated in Section 25.25.4.

The statement "source and target directories MUST reside on the same file system on the server" means that the fsid fields in the attributes for the directories are the same. If they reside on different file systems, the error NFS4ERR\_XDEV is returned.

Based on the value of the fh\_expire\_type attribute for the object, the filehandle may or may not expire on a RENAME. However, server implementers are strongly encouraged to attempt to keep filehandles from expiring in this fashion.

On some servers, the file names "." and ".." are illegal as either oldname or newname, and will result in the error NFS4ERR\_BADNAME. In addition, on many servers the case of oldname or newname being an alias for the source directory will be checked for. Such servers will return the error NFS4ERR\_INVALID in these cases.

If either of the source or target filehandles are not directories, the server will return NFS4ERR\_NOTDIR.

## 25.27. Operation 31: RESTOREFH - Restore Saved Filehandle

### 25.27.1. ARGUMENTS

```
/* SAVED_FH: */  
void;
```

### 25.27.2. RESULTS

```
struct RESTOREFH4res {  
    /*  
     * If status is NFS4_OK,  
     *    new CURRENT_FH: value of saved fh  
     */  
    nfsstat4      status;  
};
```

### 25.27.3. DESCRIPTION

The RESTOREFH operation sets the current filehandle and stateid to the values in the saved filehandle and stateid. If there is no saved filehandle, then the server will return the error NFS4ERR\_NOFILEHANDLE.

See Section 23.2.3.1.1 for more details on the current filehandle.

See Section 23.2.3.1.2 for more details on the current stateid.

#### 25.27.4. IMPLEMENTATION

Operations like OPEN and LOOKUP use the current filehandle to represent a directory and replace it with a new filehandle. Assuming that the previous filehandle was saved with a SAVEFH operator, the previous filehandle can be restored as the current filehandle. This is commonly used to obtain post-operation attributes for the directory, e.g.,

```
PUTFH (directory filehandle)
SAVEFH
GETATTR attrbits      (pre-op dir attrs)
CREATE optbits "foo" attrs
GETATTR attrbits      (file attributes)
RESTOREFH
GETATTR attrbits      (post-op dir attrs)
```

#### 25.28. Operation 32: SAVEFH - Save Current Filehandle

##### 25.28.1. ARGUMENTS

```
/* CURRENT_FH: */
void;
```

##### 25.28.2. RESULTS

```
struct SAVEFH4res {
    /*
     * If status is NFS4_OK,
     *   new SAVED_FH: value of current fh
     */
    nfsstat4      status;
};
```

##### 25.28.3. DESCRIPTION

The SAVEFH operation saves the current filehandle and stateid. If a previous filehandle was saved, then it is no longer accessible. The saved filehandle can be restored as the current filehandle with the RESTOREFH operator.

On success, the current filehandle retains its value.

See Section 23.2.3.1.1 for more details on the current filehandle.

See Section 23.2.3.1.2 for more details on the current stateid.



## 25.28.4. IMPLEMENTATION

## 25.29. Operation 33: SECINFO - Obtain Available Security

Although this is an existing NFSv4.1 operation and appropriately described in this document, much of the detail regarding the values returned and their role in security negotiation is described in Section 16 of the NFSv4-wide security document, currently [I-D.dnoveck-nfsv4-security]. This adaptation has been necessary since connection characteristics are now an appropriate subject of negotiation, where previously negotiation only concerned the choice of appropriate auth flavors on existing connection.

## 25.29.1. ARGUMENTS

```
struct SECINFO4args {  
    /* CURRENT_FH: directory */  
    component4      name;  
};
```

## 25.29.2. RESULTS

```

/*
 * From RFC 2203
 */
enum rpc_gss_svc_t {
    RPC_GSS_SVC_NONE          = 1,
    RPC_GSS_SVC_INTEGRITY     = 2,
    RPC_GSS_SVC_PRIVACY       = 3
};

struct rpcsec_gss_info {
    sec_oid4      oid;
    qop4          qop;
    rpc_gss_svc_t service;
};

/* RPCSEC_GSS has a value of '6' - See RFC 2203 */
union secinfo4 switch (uint32_t flavor) {
    case RPCSEC_GSS:
        rpcsec_gss_info      flavor_info;
    default:
        void;
};

typedef secinfo4 SECINFO4resok<>;

union SECINFO4res switch (nfsstat4 status) {
    case NFS4_OK:
        /* CURRENTFH: consumed */
        SECINFO4resok resok4;
    default:
        void;
};

```

### 25.29.3. DESCRIPTION

The SECINFO operation is used by the client to obtain a list of valid RPC authentication flavors for a specific directory filehandle, file name pair. SECINFO should apply the same access methodology used for LOOKUP when evaluating the name. Therefore, if the requester does not have the appropriate access to LOOKUP the name, then SECINFO MUST behave the same way and return NFS4ERR\_ACCESS.

The result will contain an array that represents the security mechanisms available, with an order corresponding to the server's preferences, the most preferred being first in the array. The client is free to pick whatever security mechanism it both desires and supports, or to pick in the server's preference order the first one

it supports. The array entries are represented by the `secinfo4` structure. The field 'flavor' will contain a value of `AUTH_NONE`, `AUTH_SYS` (as defined in RFC 5531 [RFC5531]), or `RPCSEC_GSS` (as defined in RFC 2203 [RFC2203]). The field flavor can also be any other security flavor registered with IANA.

For the flavors `AUTH_NONE` and `AUTH_SYS`, no additional security information is returned. The same is true of many (if not most) other security flavors, including `AUTH_DH`. For a return value of `RPCSEC_GSS`, a security triple is returned that contains the mechanism object identifier (OID, as defined in RFC 2743 [RFC2743]), the quality of protection (as defined in RFC 2743 [RFC2743]), and the service type (as defined in RFC 2203 [RFC2203]). It is possible for `SECINFO` to return multiple entries with flavor equal to `RPCSEC_GSS` with different security triple values.

On success, the current filehandle is consumed (see Section 6.2.1.8), and if the next operation after `SECINFO` tries to use the current filehandle, that operation will fail with the status `NFS4ERR_NOFILEHANDLE`.

If the name has a length of zero, or if the name does not obey the UTF-8 definition (assuming UTF-8 capabilities are enabled; see Section 21.1), the error `NFS4ERR_INVALID` will be returned.

See Section 16 of [I-D.dnoveck-nfsv4-security] for additional information on the use of `SECINFO`.

#### 25.29.4. IMPLEMENTATION

The `SECINFO` operation is expected to be used by the NFS client when the error value of `NFS4ERR_WRONGSEC` is returned from another NFS operation. This signifies to the client that the server's security policy is different from what the client is currently using. At this point, the client is expected to obtain a list of possible security flavors and choose what best suits its policies.

As mentioned, the server's security policies will determine when a client request receives `NFS4ERR_WRONGSEC`. See Table 14 for a list of operations that can return `NFS4ERR_WRONGSEC`. In addition, when `READDIR` returns attributes, the `rdattr_error` (Section 11.12.1.12) can contain `NFS4ERR_WRONGSEC`. Note that `CREATE` and `REMOVE` MUST NOT return `NFS4ERR_WRONGSEC`. The rationale for `CREATE` is that unless the target name exists, it cannot have a separate security policy from the parent directory, and the security policy of the parent was checked when its filehandle was injected into the `COMPOUND` request's operations stream (for similar reasons, an `OPEN` operation that creates the target MUST NOT return `NFS4ERR_WRONGSEC`). If the target

name exists, while it might have a separate security policy, that is irrelevant because CREATE MUST return NFS4ERR\_EXIST. The rationale for REMOVE is that while that target might have a separate security policy, the target is going to be removed, and so the security policy of the parent trumps that of the object being removed. RENAME and LINK MAY return NFS4ERR\_WRONGSEC, but the NFS4ERR\_WRONGSEC error applies only to the saved filehandle (See Section 6.2.2). Any NFS4ERR\_WRONGSEC error on the current filehandle used by LINK and RENAME MUST be returned by the PUTFH, PUTPUBFH, PUTROOTFH, or RESTOREFH operation that injected the current filehandle.

With the exception of LINK and RENAME, the set of operations that can return NFS4ERR\_WRONGSEC represents the point at which the client can inject a filehandle into the "current filehandle" at the server. The filehandle is either provided by the client (PUTFH, PUTPUBFH, PUTROOTFH), generated as a result of a name-to-filehandle translation (LOOKUP and OPEN), or generated from the saved filehandle via RESTOREFH. As Section 6.2.1.1 states, a put filehandle operation followed by SAVEFH MUST NOT return NFS4ERR\_WRONGSEC. Thus, the RESTOREFH operation, under certain conditions (See Section 6.2.1), is permitted to return NFS4ERR\_WRONGSEC so that security policies can be honored.

The READDIR operation will not directly return the NFS4ERR\_WRONGSEC error. However, if the READDIR request included a request for attributes, it is possible that the READDIR request's security triple did not match that of a directory entry. If this is the case and the client has requested the rdattrib\_error attribute, the server will return the NFS4ERR\_WRONGSEC error in rdattrib\_error for the entry.

To resolve an error return of NFS4ERR\_WRONGSEC, the client does the following:

- \* For LOOKUP and OPEN, the client will use SECINFO with the same current filehandle and name as provided in the original LOOKUP or OPEN to enumerate the available security triples.
- \* For the rdattrib\_error, the client will use SECINFO with the same current filehandle as provided in the original READDIR. The name passed to SECINFO will be that of the directory entry (as returned from READDIR) that had the NFS4ERR\_WRONGSEC error in the rdattrib\_error attribute.

- \* For PUTFH, PUTROOTFH, PUTPUBFH, RESTOREFH, LINK, and RENAME, the client will use SECINFO\_NO\_NAME { style = SECINFO\_STYLE4\_CURRENT\_FH }. The client will prefix the SECINFO\_NO\_NAME operation with the appropriate PUTFH, PUTPUBFH, or PUTROOTFH operation that provides the filehandle originally provided by the PUTFH, PUTPUBFH, PUTROOTFH, or RESTOREFH operation.

NOTE: In NFSv4.0, the client was required to use SECINFO, and had to reconstruct the parent of the original filehandle and the component name of the original filehandle. The introduction in NFSv4.1 of SECINFO\_NO\_NAME obviates the need for reconstruction.

- \* For LOOKUPP, the client will use SECINFO\_NO\_NAME { style = SECINFO\_STYLE4\_PARENT } and provide the filehandle that equals the filehandle originally provided to LOOKUPP.

See Section 28 for a discussion on the recommendations for the security flavor used by SECINFO and SECINFO\_NO\_NAME.

## 25.30. Operation 34: SETATTR - Set Attributes

### 25.30.1. ARGUMENTS

```
struct SETATTR4args {  
    /* CURRENT_FH: target object */  
    stateid4      stateid;  
    fattr4        obj_attributes;  
};
```

### 25.30.2. RESULTS

```
struct SETATTR4res {  
    nfsstat4      status;  
    bitmap4       attrset;  
};
```

### 25.30.3. DESCRIPTION

The SETATTR operation changes one or more of the attributes of a file system object. The new attributes are specified with a bitmap and the attributes that follow the bitmap in bit order.

The `stateid` argument for `SETATTR` is used to provide locking context that is necessary for `SETATTR` requests that set the size attribute. Since setting the size attribute modifies the file's data, it has the same locking requirements as a corresponding `WRITE`. The area between the old end-of-file and the new end-of-file is considered to be modified just as would have been the case had the area in question been specified as a result:

- \* Any `SETATTR` that sets the size attribute is incompatible with a share reservation that specifies `OPEN4_SHARE_DENY_WRITE`.
- \* The target of `WRITE`, for the purpose of checking conflicts with mandatory byte-range locks, when a server is implementing mandatory byte-range locking.

A valid `stateid` needs to be specified when the size attributed is said so that the server can determine, when writing is being denied, whether the size modification is allowed under an open allowing writing. When the file size attribute is not set, the special `stateid` consisting of all bits equal to zero MAY be passed.

On either success or failure of the operation, the server will return the `attrsset` bitmask to represent what (if any) attributes were successfully set. The `attrsset` in the response is a subset of the `attrmask` field of the `obj_attributes` field in the argument.

On success, the current filehandle retains its value.

#### 25.30.4. IMPLEMENTATION

If the request specifies the owner attribute to be set, the server should allow the operation to succeed if the current owner of the object matches the value specified in the request. Some servers may be implemented in a way that prohibits the setting of the owner attribute unless the requester has privilege to do so. If the server is lenient in this one case of matching owner values, the client implementation can be simplified in cases of creation of an object (e.g., using an exclusive create via `OPEN`) followed by a `SETATTR`.

The file size attribute is used to request changes to the size of a file. A value of zero causes the file to be truncated, a value less than the current size of the file causes data from new size to the end of the file to be discarded, and a size greater than the current size of the file causes logically zeroed data bytes to be added to the end of the file. Servers are free to implement this using unallocated bytes (holes) or allocated data bytes set to zero. Clients should not make any assumptions regarding a server's implementation of this feature, beyond that the bytes in the affected byte-range returned by READ will be zeroed. Servers MUST support extending the file size via SETATTR.

SETATTR is not guaranteed to be atomic. A failed SETATTR may partially change a file's attributes, hence the reason why the reply always includes the status and the list of attributes that were set.

If the object whose attributes are being changed has a file delegation that is held by a client other than the one doing the SETATTR, the delegation(s) must be recalled, and the operation cannot proceed to actually change an attribute until each such delegation is returned or revoked. In all cases in which delegations are recalled, the server is likely to return one or more NFS4ERR\_DELAY errors while the delegation(s) remains outstanding, although it might not do that if the delegations are returned quickly.

If the object whose attributes are being set is a directory and another client holds a directory delegation for that directory, then if enabled, asynchronous notifications will be generated when the set of attributes changed has a non-null intersection with the set of attributes for which notification is requested. Notifications of type NOTIFY4\_CHANGE\_DIR\_ATTRS will be sent to the appropriate client(s), but the SETATTR is not delayed by waiting for these notifications to be sent.

If the object whose attributes are being set is a member of the directory for which another client holds a directory delegation, then asynchronous notifications will be generated when the set of attributes changed has a non-null intersection with the set of attributes for which notification is requested. Notifications of type NOTIFY4\_CHANGE\_CHILD\_ATTRS will be sent to the appropriate clients, but the SETATTR is not delayed by waiting for these notifications to be sent.

Changing the size of a file with SETATTR indirectly changes the time\_modify and change attributes. A client must account for this as size changes can result in data deletion.

The attributes `time_access_set` and `time_modify_set` are write-only attributes constructed as a switched union so the client can direct the server in setting the time values. If the switched union specifies `SET_TO_CLIENT_TIME4`, the client has provided an `nfstime4` to be used for the operation. If the switch union does not specify `SET_TO_CLIENT_TIME4`, the server is to use its current time for the `SETATTR` operation.

If server and client times differ, programs that compare client time to file times can break. A time synchronization protocol should be used to limit client/server time skew.

Use of a COMPOUND containing a `VERIFY` operation specifying only the change attribute, immediately followed by a `SETATTR`, provides a means whereby a client may specify a request that emulates the functionality of the `SETATTR` guard mechanism of NFSv3. Since the function of the guard mechanism is to avoid changes to the file attributes based on stale information, delays between checking of the guard condition and the setting of the attributes have the potential to compromise this function, as would the corresponding delay in the NFSv4 emulation. Therefore, NFSv4.1 servers should take care to avoid such delays, to the degree possible, when executing such a request.

If the server does not support an attribute as requested by the client, the server MUST return `NFS4ERR_ATTRNOTSUPP`.

A mask of the attributes actually set is returned by `SETATTR` in all cases. That mask MUST NOT include attribute bits not requested to be set by the client. If the attribute masks in the request and reply are equal, the status field in the reply MUST be `NFS4_OK`.

## 25.31. Operation 37: `VERIFY` - Verify Same Attributes

### 25.31.1. ARGUMENTS

```
struct VERIFY4args {  
    /* CURRENT_FH: object */  
    fattr4          obj_attributes;  
};
```

### 25.31.2. RESULTS

```
struct VERIFY4res {  
    nfsstat4        status;  
};
```



### 25.31.3. DESCRIPTION

The VERIFY operation is used to verify that attributes have the value assumed by the client before proceeding with the following operations in the COMPOUND request. If any of the attributes do not match, then the error NFS4ERR\_NOT\_SAME must be returned. The current filehandle retains its value after successful completion of the operation.

### 25.31.4. IMPLEMENTATION

One possible use of the VERIFY operation is the following series of operations. With this, the client is attempting to verify that the file being removed will match what the client expects to be removed. This series can help prevent the unintended deletion of a file.

```
PUTFH (directory filehandle)
LOOKUP (file name)
VERIFY (filehandle == fh)
PUTFH (directory filehandle)
REMOVE (file name)
```

This series does not prevent a second client from removing and creating a new file in the middle of this sequence, but it does help avoid the unintended result.

In the case that a OPTIONAL attribute is specified in the VERIFY operation and the server does not support that attribute for the file system object, the error NFS4ERR\_ATTRNOTSUPP is returned to the client.

When the attribute specified is supported but is one for which use of VERIFY is inappropriate (e.g. rgetattr\_error or any set-only attribute (such as time\_modify\_set)), the error NFS4ERR\_INVALID is returned to the client.

## 25.32. Operation 38: WRITE - Write to File

### 25.32.1. ARGUMENTS

```

enum stable_how4 {
    UNSTABLE4      = 0,
    DATA_SYNC4    = 1,
    FILE_SYNC4     = 2
};

struct WRITE4args {
    /* CURRENT_FH: file */
    stateid4        stateid;
    offset4         offset;
    stable_how4     stable;
    opaque          data<>;
};

```

### 25.32.2. RESULTS

```

struct WRITE4resok {
    count4          count;
    stable_how4     committed;
    verifier4       writeverf;
};

union WRITE4res switch (nfsstat4 status) {
    case NFS4_OK:
        WRITE4resok   resok4;
    default:
        void;
};

```

### 25.32.3. DESCRIPTION

The WRITE operation is used to write data to a regular file. The target file is specified by the current filehandle. The offset specifies the offset where the data should be written. An offset of zero specifies that the write should start at the beginning of the file. The count, as encoded as part of the opaque data parameter, represents the number of bytes of data that are to be written. If the count is zero, the WRITE will succeed and return a count of zero subject to permissions checking. The server MAY write fewer bytes than requested by the client.

The client specifies with the stable parameter the method of how the data is to be processed by the server. If stable is FILE\_SYNC4, the server MUST commit the data written plus all file system metadata to stable storage before returning results. This corresponds to the NFSv2 protocol semantics. Any other behavior constitutes a protocol violation. If stable is DATA\_SYNC4, then the server MUST commit all of the data to stable storage and enough of the metadata to retrieve

the data before returning. The server implementer is free to implement DATA\_SYNC4 in the same fashion as FILE\_SYNC4, but with a possible performance drop. If stable is UNSTABLE4, the server is free to commit any part of the data and the metadata to stable storage, including all or none, before returning a reply to the client. There is no guarantee whether or when any uncommitted data will subsequently be committed to stable storage. The only guarantees made by the server are that it will not destroy any data without changing the value of writeverf and that it will not commit the data and metadata at a level less than that requested by the client.

Except when special stateids are used, the stateid value for a WRITE request represents a value returned from a previous byte-range LOCK or OPEN request or the stateid associated with a delegation. The stateid identifies the associated owners if any and is used by the server to verify that the associated locks are still valid (e.g., have not been revoked).

Upon successful completion, the following results are returned. The count result is the number of bytes of data written to the file. The server may write fewer bytes than requested. If so, the actual number of bytes written starting at location, offset, is returned.

The server also returns an indication of the level of commitment of the data and metadata via committed. Per Table 20,

- \* The server MAY commit the data at a stronger level than requested.
- \* The server MUST commit the data at a level at least as strong as that requested.

stable	committed
UNSTABLE4	FILE_SYNC4, DATA_SYNC4, UNSTABLE4
DATA_SYNC4	FILE_SYNC4, DATA_SYNC4
FILE_SYNC4	FILE_SYNC4

Table 20: Valid Combinations of the Fields  
Stable in the Request and Committed in the  
Reply

The final portion of the result is the field `writeverf`. This field is the write verifier and is a cookie that the client can use to determine whether a server has changed instance state (e.g., server restart) between a call to `WRITE` and a subsequent call to either `WRITE` or `COMMIT`. This cookie **MUST** be unchanged during a single instance of the NFSv4.1 server and **MUST** be unique between instances of the NFSv4.1 server. If the cookie changes, then the client **MUST** assume that any data written with an `UNSTABLE4` value for committed and an old `writeverf` in the reply has been lost and will need to be recovered.

If a client writes data to the server with the `stable` argument set to `UNSTABLE4` and the reply yields a committed response of `DATA_SYNC4` or `UNSTABLE4`, the client will follow up some time in the future with a `COMMIT` operation to synchronize outstanding asynchronous data and metadata with the server's stable storage, barring client error. It is possible that due to client crash or other error that a subsequent `COMMIT` will not be received by the server.

For a `WRITE` with a `stateid` value of all bits equal to zero, the server **MAY** allow the `WRITE` to be serviced subject to mandatory byte-range locks or the current share deny modes for the file. For a `WRITE` with a `stateid` value of all bits equal to 1, the server **MUST** **NOT** allow the `WRITE` operation to bypass locking checks at the server and otherwise is treated as if a `stateid` of all bits equal to zero were used.

On success, the current filehandle retains its value.

#### 25.32.4. IMPLEMENTATION

It is possible for the server to write fewer bytes of data than requested by the client. In this case, the server should not return an error unless no data was written at all. If the server writes less than the number of bytes specified, the client will need to send another `WRITE` to write the remaining data.

It is assumed that the act of writing data to a file will cause the `time_modified` and `change` attributes of the file to be updated. However, these attributes are not to be changed unless the contents of the file are changed. Thus, a `WRITE` request with `count` set to zero will not cause the `time_modified` and `change` attributes of the file to be updated.

Stable storage is persistent storage that survives:

1. Repeated power failures.

2. Hardware failures (of any board, power supply, etc.).
3. Repeated software crashes and restarts.

This definition does not address failure of the stable storage module itself.

The verifier is defined to allow a client to detect different instances of an NFSv4.1 protocol server over which cached, uncommitted data may be lost. In the most likely case, the verifier allows the client to detect server restarts. This information is required so that the client can safely determine whether the server could have lost cached data. If the server fails unexpectedly and the client has uncommitted data from previous WRITE requests (done with the stable argument set to UNSTABLE4 and in which the result committed was returned as UNSTABLE4 as well), the server might not have flushed cached data to stable storage. The burden of recovery is on the client, and the client will need to retransmit the data to the server.

A suggested verifier would be to use the time that the server was last started (if restarting the server results in lost buffers).

The reply's committed field allows the client to do more effective caching. If the server is committing all WRITE requests to stable storage, then it will return with committed set to FILE\_SYNC4, regardless of the value of the stable field in the arguments. A server that uses an NVRAM accelerator may choose to implement this policy. The client can use this to increase the effectiveness of the cache by discarding cached data that has already been committed on the server.

Some implementations may return NFS4ERR\_NOSPC instead of NFS4ERR\_DQUOT when a user's quota is exceeded.

In the case that the current filehandle is of type NF4DIR, the server will return NFS4ERR\_ISDIR. If the current file is a symbolic link, the error NFS4ERR\_SYMLINK will be returned. Otherwise, if the current filehandle does not designate an ordinary file, the server will return NFS4ERR\_WRONG\_TYPE.

If mandatory byte-range locking is in effect for the file, and the corresponding byte-range of the data to be written to the file is READ\_LT or WRITE\_LT locked by an owner that is not associated with the stateid, the server MUST return NFS4ERR\_LOCKED. If so, the client MUST check if the owner corresponding to the stateid used with the WRITE operation has a conflicting READ\_LT lock that overlaps with the byte-range that was to be written. If the stateid's owner has no

conflicting READ\_LT lock, then the client should try to get the appropriate write byte-range lock via the LOCK operation before re-attempting the WRITE. When the WRITE completes, the client can release the byte-range lock via LOCKU.

If the stateid's owner had a conflicting READ\_LT lock, then the client has no choice but to return an error to the application that attempted the WRITE. The reason is that since the stateid's owner had a READ\_LT lock, either the server attempted to temporarily effectively upgrade this READ\_LT lock to a WRITE\_LT lock or the server has no upgrade capability. If the server attempted to upgrade the READ\_LT lock and failed, it is pointless for the client to re-attempt the upgrade via the LOCK operation, because there might be another client also trying to upgrade. If two clients are blocked trying to upgrade the same lock, the clients deadlock. If the server has no upgrade capability, then it is pointless to try a LOCK operation to upgrade.

If one or more other clients have delegations for the file being written, those delegations MUST be recalled, and the operation cannot proceed until those delegations are returned or revoked. Except where this happens very quickly, one or more NFS4ERR\_DELAY errors will be returned to requests made while the delegation remains outstanding. Normally, delegations will not be recalled as a result of a WRITE operation since the recall will occur as a result of an earlier OPEN. However, since it is possible for a WRITE to be done with a special stateid, the server needs to check for this case even though the client should have done an OPEN previously.

### 25.33. Operation 40: BACKCHANNEL\_CTL - Backchannel Control

#### 25.33.1. ARGUMENT

```

typedef opaque gsshandle4_t<>;

struct gss_cb_handles4 {
    rpc_gss_service_t      gcbp_service; /* RFC 2203 */
    gsshandle4_t           gcbp_handle_from_server;
    gsshandle4_t           gcbp_handle_from_client;
};

union callback_sec_parms4 switch (uint32_t cb_secflavor) {
case AUTH_NONE:
    void;
case AUTH_SYS:
    authsys_parms    cbsp_sys_cred; /* RFC 5531 */
case RPCSEC_GSS:
    gss_cb_handles4 cbsp_gss_handles;
};

struct BACKCHANNEL_CTL4args {
    uint32_t          bca_cb_program;
    callback_sec_parms4 bca_sec_parms<>;
};

```

#### 25.33.2. RESULT

```

struct BACKCHANNEL_CTL4res {
    nfsstat4          bcr_status;
};

```

#### 25.33.3. DESCRIPTION

The BACKCHANNEL\_CTL operation replaces the backchannel's callback program number and adds (not replaces) RPCSEC\_GSS handles for use by the backchannel.

The arguments of the BACKCHANNEL\_CTL call are a subset of the CREATE\_SESSION parameters. In the arguments of BACKCHANNEL\_CTL, the bca\_cb\_program field and bca\_sec\_parms fields correspond respectively to the csa\_cb\_program and csa\_sec\_parms fields of the arguments of CREATE\_SESSION (Section 25.36).

BACKCHANNEL\_CTL MUST appear in a COMPOUND that starts with SEQUENCE.

If the RPCSEC\_GSS handle identified by gcbp\_handle\_from\_server does not exist on the server, the server MUST return NFS4ERR\_NOENT.

If an RPCSEC\_GSS handle is using the SSV context (See Section 7.9), then because each SSV RPCSEC\_GSS handle shares a common SSV GSS context, there are security considerations specific to this situation discussed in Section 7.10.

#### 25.34. Operation 41: BIND\_CONN\_TO\_SESSION - Associate Connection with Session

##### 25.34.1. ARGUMENT

```
enum channel_dir_from_client4 {
    CDFC4_FORE          = 0x1,
    CDFC4_BACK          = 0x2,
    CDFC4_FORE_OR_BOTH  = 0x3,
    CDFC4_BACK_OR_BOTH  = 0x7
};

struct BIND_CONN_TO_SESSION4args {
    sessionid4      bctsa_sessid;

    channel_dir_from_client4
                    bctsa_dir;

    bool            bctsa_use_conn_in_rdma_mode;
};
```

##### 25.34.2. RESULT



```
enum channel_dir_from_server4 {
    CDFS4_FORE      = 0x1,
    CDFS4_BACK      = 0x2,
    CDFS4_BOTH      = 0x3
};

struct BIND_CONN_TO_SESSION4resok {
    sessionid4      bctsr_sessid;

    channel_dir_from_server4
                    bctsr_dir;

    bool            bctsr_use_conn_in_rdma_mode;
};

union BIND_CONN_TO_SESSION4res
    switch (nfsstat4 bctsr_status) {

    case NFS4_OK:
        BIND_CONN_TO_SESSION4resok
            bctsr_resok4;

    default:
        void;
    };
```

#### 25.34.3. DESCRIPTION

BIND\_CONN\_TO\_SESSION is used to associate additional connections with a session. It MUST be used on the connection being associated with the session. It MUST be the only operation in the COMPOUND procedure. If SP4\_NONE (Section 25.35) state protection is used, any principal, security flavor, or RPCSEC\_GSS context MAY be used to invoke the operation. If SP4\_MACH\_CRED is used, RPCSEC\_GSS MUST be used with the integrity or privacy services, using the principal that created the client ID. If SP4\_SSV is used, RPCSEC\_GSS with the SSV GSS mechanism (Section 7.9) and integrity or privacy MUST be used.

If, when the client ID was created, the client opted for SP4\_NONE state protection, the client is not required to use BIND\_CONN\_TO\_SESSION to associate the connection with the session, unless the client wishes to associate the connection with the backchannel. When SP4\_NONE protection is used, simply sending a COMPOUND request with a SEQUENCE operation is sufficient to associate the connection with the session specified in SEQUENCE.

The field bctsa\_dir indicates whether the client wants to associate the connection with the fore channel or the backchannel or both channels. The value CDFS4\_FORE\_OR\_BOTH indicates that the client

wants to associate the connection with both the fore channel and backchannel, but will accept the connection being associated to just the fore channel. The value CDFC4\_BACK\_OR\_BOTH indicates that the client wants to associate with both the fore channel and backchannel, but will accept the connection being associated with just the backchannel. The server replies in bctsr\_dir which channel(s) the connection is associated with. If the client specified CDFC4\_FORE, the server MUST return CDFS4\_FORE. If the client specified CDFC4\_BACK, the server MUST return CDFS4\_BACK. If the client specified CDFC4\_FORE\_OR\_BOTH, the server MUST return CDFS4\_FORE or CDFS4\_BOTH. If the client specified CDFC4\_BACK\_OR\_BOTH, the server MUST return CDFS4\_BACK or CDFS4\_BOTH.

See the CREATE\_SESSION operation (Section 25.36), and the description of the argument csa\_use\_conn\_in\_rdma\_mode to understand bctsa\_use\_conn\_in\_rdma\_mode, and the description of csr\_use\_conn\_in\_rdma\_mode to understand bctsr\_use\_conn\_in\_rdma\_mode.

Invoking BIND\_CONN\_TO\_SESSION on a connection already associated with the specified session has no effect, and the server MUST respond with NFS4\_OK, unless the client is demanding changes to the set of channels the connection is associated with. If so, the server MUST return NFS4ERR\_INVALID.

#### 25.34.4. IMPLEMENTATION

If a session's channel loses all connections, depending on the client ID's state protection and type of channel, the client might need to use BIND\_CONN\_TO\_SESSION to associate a new connection. If the server restarted and does not keep the reply cache in stable storage, the server will not recognize the session ID. The client will ultimately have to invoke EXCHANGE\_ID to create a new client ID and session.

Suppose SP4\_SSV state protection is being used, and BIND\_CONN\_TO\_SESSION is among the operations included in the spo\_must\_enforce set when the client ID was created (Section 25.35). If so, there is an issue if SET\_SSV is sent, no response is returned, and the last connection associated with the client ID drops. The client, per the sessions model, MUST retry the SET\_SSV. But it needs a new connection to do so, and MUST associate that connection with the session via a BIND\_CONN\_TO\_SESSION authenticated with the SSV GSS mechanism. The problem is that the RPCSEC\_GSS message integrity codes use a subkey derived from the SSV as the key and the SSV may have changed. While there are multiple recovery strategies, a single, general strategy is described here.

- \* The client reconnects.

- \* The client assumes that the SET\_SSV was executed, and so sends BIND\_CONN\_TO\_SESSION with the subkey (derived from the new SSV, i.e., what SET\_SSV would have set the SSV to) used as the key for the RPCSEC\_GSS credential message integrity codes.
- \* If the request succeeds, this means that the original attempted SET\_SSV did execute successfully. The client re-sends the original SET\_SSV, which the server will reply to via the reply cache.
- \* If the server returns an RPC authentication error, this means that the server's current SSV was not changed (and the SET\_SSV was likely not executed). The client then tries BIND\_CONN\_TO\_SESSION with the subkey derived from the old SSV as the key for the RPCSEC\_GSS message integrity codes.
- \* The attempted BIND\_CONN\_TO\_SESSION with the old SSV should succeed. If so, the client re-sends the original SET\_SSV. If the original SET\_SSV was not executed, then the server executes it. If the original SET\_SSV was executed but failed, the server will return the SET\_SSV from the reply cache.

#### 25.35. Operation 42: EXCHANGE\_ID - Instantiate Client ID

The EXCHANGE\_ID operation exchanges long-hand client and server identifiers (owners) and provides access to a client ID, creating one if necessary. This client ID becomes associated with the connection on which the operation is done, so that it is available when a CREATE\_SESSION is done or when the connection is used to issue a request on an existing session associated with the current client.

##### 25.35.1. ARGUMENT

```

const EXCHGID4_FLAG_SUPP_MOVED_REFER    = 0x00000001;
const EXCHGID4_FLAG_SUPP_MOVED_MIGR     = 0x00000002;

const EXCHGID4_FLAG_BIND_PRINC_STATEID  = 0x00000100;

const EXCHGID4_FLAG_USE_NON_PNFS        = 0x00010000;
const EXCHGID4_FLAG_USE_PNFS_MDS        = 0x00020000;
const EXCHGID4_FLAG_USE_PNFS_DS         = 0x00040000;

const EXCHGID4_FLAG_MASK_PNFS           = 0x00070000;

const EXCHGID4_FLAG_UPD_CONFIRMED_REC_A = 0x40000000;
const EXCHGID4_FLAG_CONFIRMED_R        = 0x80000000;

struct state_protect_ops4 {

```

```

        bitmap4 spo_must_enforce;
        bitmap4 spo_must_allow;
};

struct ssv_sp_parms4 {
    state_protect_ops4    ssp_ops;
    sec_oid4              ssp_hash_algs<>;
    sec_oid4              ssp_encr_algs<>;
    uint32_t              ssp_window;
    uint32_t              ssp_num_gss_handles;
};

enum state_protect_how4 {
    SP4_NONE = 0,
    SP4_MACH_CRED = 1,
    SP4_SSV = 2
};

union state_protect4_a switch(state_protect_how4 spa_how) {
    case SP4_NONE:
        void;
    case SP4_MACH_CRED:
        state_protect_ops4    spa_mach_ops;
    case SP4_SSV:
        ssv_sp_parms4         spa_ssv_parms;
};

struct EXCHANGE_ID4args {
    client_owner4          eia_clientowner;
    uint32_t               eia_flags;
    state_protect4_a       eia_state_protect;
    nfs_impl_id4           eia_client_impl_id<1>;
};

```

#### 25.35.2. RESULT

```

struct ssv_prot_info4 {
    state_protect_ops4    spi_ops;
    uint32_t              spi_hash_alg;
    uint32_t              spi_encr_alg;
    uint32_t              spi_ssv_len;
    uint32_t              spi_window;
    gsshandle4_t          spi_handles<>;
};

union state_protect4_r switch(state_protect_how4 spr_how) {
    case SP4_NONE:
        void;
    case SP4_MACH_CRED:
        state_protect_ops4    spr_mach_ops;
    case SP4_SSV:
        ssv_prot_info4        spr_ssv_info;
};

struct EXCHANGE_ID4resok {
    clientid4              eir_clientid;
    sequenceid4            eir_sequenceid;
    uint32_t               eir_flags;
    state_protect4_r        eir_state_protect;
    server_owner4           eir_server_owner;
    opaque                  eir_server_scope<NFS4_OPAQUE_LIMIT>;
    nfs_impl_id4            eir_server_impl_id<1>;
};

union EXCHANGE_ID4res switch (nfsstat4 eir_status) {
    case NFS4_OK:
        EXCHANGE_ID4resok    eir_resok4;

    default:
        void;
};

```

### 25.35.3. DESCRIPTION

The client uses the EXCHANGE\_ID operation to register a particular instance of that client with the server, as represented by a client\_owner4. However, when the client\_owner4 has already been registered by other means (e.g., Transparent State Migration), the client may still use EXCHANGE\_ID to obtain the client ID assigned previously.

The client ID returned from this operation will be associated with the connection on which the EXCHANGE\_ID is received and will serve as a parent object for sessions created by the client on this connection

or to which the connection is bound. As a result of using those sessions to make requests involving the creation of state, that state will become associated with the client ID returned.

In situations in which the registration of the client\_owner has not occurred previously, the client ID must first be used, along with the returned eir\_sequenceid, in creating an associated session using CREATE\_SESSION.

If the flag EXCHGID4\_FLAG\_CONFIRMED\_R is set in the result, eir\_flags, then it is an indication that the registration of the client\_owner has already occurred and that a further CREATE\_SESSION is not needed to confirm it. Of course, subsequent CREATE\_SESSION operations may be needed for other reasons.

The value eir\_sequenceid is used to establish an initial sequence value associated with the client ID returned. In cases in which a CREATE\_SESSION has already been done, there is no need for this value, since sequencing of such request has already been established, and the client has no need for this value and will ignore it.

EXCHANGE\_ID MAY be sent in a COMPOUND procedure that starts with SEQUENCE. However, when a client communicates with a server for the first time, it will not have a session, so using SEQUENCE will not be possible. If EXCHANGE\_ID is sent without a preceding SEQUENCE, then it MUST be the only operation in the COMPOUND procedure's request. If it is not, the server MUST return NFS4ERR\_NOT\_ONLY\_OP.

The eia\_clientowner field is composed of a co\_verifier field and a co\_ownerid string. As noted in Section 5.5, the co\_ownerid identifies the client, and the co\_verifier specifies a particular incarnation of that client. An EXCHANGE\_ID sent with a new incarnation of the client will lead to the server removing lock state of the old incarnation. On the other hand, when an EXCHANGE\_ID sent with the current incarnation and co\_ownerid does not result in an unrelated error, it will potentially update an existing client ID's properties or simply return information about the existing client\_id. The latter would happen when this operation is done to the same server using different network addresses as part of creating trunked connections.

A server MUST NOT provide the same client ID to two different incarnations of an eia\_clientowner.

In addition to the client ID and sequence ID, the server returns a server owner (eir\_server\_owner) and server scope (eir\_server\_scope). The former field is used in connection with network trunking as described in Section 7.5. The latter field is used to allow clients

to determine when client IDs sent by one server may be recognized by another in the event of file system migration (See Section 17.11.9 of the current document).

The client ID returned by EXCHANGE\_ID is only unique relative to the combination of `eir_server_owner.so_major_id` and `eir_server_scope`. Thus, if two servers return the same client ID, the onus is on the client to distinguish the client IDs on the basis of `eir_server_owner.so_major_id` and `eir_server_scope`. In the event two different servers claim matching `server_owner.so_major_id` and `eir_server_scope`, the client can use the verification techniques discussed in Section 7.5.1 to determine if the servers are distinct. If they are distinct, then the client will need to note the destination network addresses of the connections used with each server and use the network address as the final discriminator.

The server, as defined by the unique identity expressed in the `so_major_id` of the server owner and the server scope, needs to track several properties of each client ID it hands out. The properties apply to the client ID and all sessions associated with the client ID. The properties are derived from the arguments and results of EXCHANGE\_ID. The client ID properties include:

- \* The capabilities expressed by the following bits, which come from the results of EXCHANGE\_ID:
  - EXCHGID4\_FLAG\_SUPP\_MOVED\_REFERER
  - EXCHGID4\_FLAG\_SUPP\_MOVED\_MIGR
  - EXCHGID4\_FLAG\_BIND\_PRINC\_STATEID
  - EXCHGID4\_FLAG\_USE\_NON\_PNFS
  - EXCHGID4\_FLAG\_USE\_PNFS\_MDS
  - EXCHGID4\_FLAG\_USE\_PNFS\_DS

These properties may be updated by subsequent EXCHANGE\_ID operations on confirmed client IDs though the server MAY refuse to change them.

- \* The state protection method used, one of SP4\_NONE, SP4\_MACH\_CRED, or SP4\_SSV, as set by the `spa_how` field of the arguments to EXCHANGE\_ID. Once the client ID is confirmed, this property cannot be updated by subsequent EXCHANGE\_ID operations.
- \* For SP4\_MACH\_CRED or SP4\_SSV state protection:

- The list of operations (`spo_must_enforce`) that MUST use the specified state protection. This list comes from the results of `EXCHANGE_ID`.
- The list of operations (`spo_must_allow`) that MAY use the specified state protection. This list comes from the results of `EXCHANGE_ID`.

Once the client ID is confirmed, these properties cannot be updated by subsequent `EXCHANGE_ID` requests.

\* For `SP4_SSV` protection:

- The OID of the hash algorithm. This property is represented by one of the algorithms in the `ssp_hash_algs` field of the `EXCHANGE_ID` arguments. Once the client ID is confirmed, this property cannot be updated by subsequent `EXCHANGE_ID` requests.
- The OID of the encryption algorithm. This property is represented by one of the algorithms in the `ssp_encr_algs` field of the `EXCHANGE_ID` arguments. Once the client ID is confirmed, this property cannot be updated by subsequent `EXCHANGE_ID` requests.
- The length of the SSV. This property is represented by the `spi_ssv_len` field in the `EXCHANGE_ID` results. Once the client ID is confirmed, this property cannot be updated by subsequent `EXCHANGE_ID` operations.

There are REQUIRED and RECOMMENDED relationships among the length of the key of the encryption algorithm ("key length"), the length of the output of hash algorithm ("hash length"), and the length of the SSV ("SSV length").

- o key length MUST be  $\leq$  hash length. This is because the keys used for the encryption algorithm are actually subkeys derived from the SSV, and the derivation is via the hash algorithm. The selection of an encryption algorithm with a key length that exceeded the length of the output of the hash algorithm would require padding, and thus weaken the use of the encryption algorithm.
- o hash length should be  $\leq$  SSV length. This is because the SSV is a key used to derive subkeys via an HMAC, and it is recommended that the key used as input to an HMAC be at least as long as the length of the HMAC's hash algorithm's output (See Section 3 of [RFC2104]).



- o key length should be  $\leq$  SSV length. This is a transitive result of the above two invariants.
- o key length should be  $\geq$  hash length / 2. This is because the subkey derivation is via an HMAC and it is recommended that if the HMAC has to be truncated, it should not be truncated to less than half the hash length (See Section 4 of [RFC2104]).
- Number of concurrent versions of the SSV the client and server will support (See Section 7.9). This property is represented by `spi_window` in the `EXCHANGE_ID` results. The property may be updated by subsequent `EXCHANGE_ID` operations.
- \* The client's implementation ID as represented by the `eia_client_impl_id` field of the arguments. The property may be updated by subsequent `EXCHANGE_ID` requests.
- \* The server's implementation ID as represented by the `eir_server_impl_id` field of the reply. The property may be updated by replies to subsequent `EXCHANGE_ID` requests.

The `eia_flags` passed as part of the arguments and the `eir_flags` results allow the client and server to inform each other of their capabilities as well as indicate how the client ID will be used. Whether a bit is set or cleared on the arguments' flags does not force the server to set or clear the same bit on the results' side. Bits not defined above cannot be set in the `eia_flags` field. If they are, the server MUST reject the operation with `NFS4ERR_INVAL`.

The `EXCHGID4_FLAG_UPD_CONFIRMED_REC_A` bit can only be set in `eia_flags`; it is always off in `eir_flags`. The `EXCHGID4_FLAG_CONFIRMED_R` bit can only be set in `eir_flags`; it is always off in `eia_flags`. If the server recognizes the `co_ownerid` and `co_verifier` as mapping to a confirmed client ID, it sets `EXCHGID4_FLAG_CONFIRMED_R` in `eir_flags`. The `EXCHGID4_FLAG_CONFIRMED_R` flag allows a client to tell if the client ID it is trying to create already exists and is confirmed.

If `EXCHGID4_FLAG_UPD_CONFIRMED_REC_A` is set in `eia_flags`, this means that the client is attempting to update properties of an existing confirmed client ID (if the client wants to update properties of an unconfirmed client ID, it MUST NOT set `EXCHGID4_FLAG_UPD_CONFIRMED_REC_A`). If so, it is RECOMMENDED that the client send the update `EXCHANGE_ID` operation in the same COMPOUND as a SEQUENCE so that the `EXCHANGE_ID` is executed exactly once. Whether the client can update the properties of client ID depends on the state protection it selected when the client ID was created, and

the principal and security flavor it used when sending the EXCHANGE\_ID operation. The situations described in items 6, 7, 8, or 9 of the second numbered list of Section 25.35.4 below will apply. Note that if the operation succeeds and returns a client ID that is already confirmed, the server MUST set the EXCHGID4\_FLAG\_CONFIRMED\_R bit in eir\_flags.

If EXCHGID4\_FLAG\_UPD\_CONFIRMED\_REC\_A is not set in eia\_flags, this means that the client is trying to establish a new client ID; it is attempting to trunk data communication to the server (See Section 7.5); or it is attempting to update properties of an unconfirmed client ID. The situations described in items 1, 2, 3, 4, or 5 of the second numbered list of Section 25.35.4 below will apply. Note that if the operation succeeds and returns a client ID that was previously confirmed, the server MUST set the EXCHGID4\_FLAG\_CONFIRMED\_R bit in eir\_flags.

When the EXCHGID4\_FLAG\_SUPP\_MOVED\_REFER flag bit is set, the client indicates that it is capable of dealing with an NFS4ERR\_MOVED error as part of a referral sequence. When this bit is not set, it is still legal for the server to perform a referral sequence. However, a server may use the fact that the client is incapable of correctly responding to a referral, by avoiding it for that particular client. It may, for instance, act as a proxy for that particular file system, at some cost in performance, although it is not obligated to do so. If the server will potentially perform a referral, it MUST set EXCHGID4\_FLAG\_SUPP\_MOVED\_REFER in eir\_flags.

When the EXCHGID4\_FLAG\_SUPP\_MOVED\_MIGR is set, the client indicates that it is capable of dealing with an NFS4ERR\_MOVED error as part of a file system migration sequence. When this bit is not set, it is still legal for the server to indicate that a file system has moved, when this in fact happens. However, a server may use the fact that the client is incapable of correctly responding to a migration in its scheduling of file systems to migrate so as to avoid migration of file systems being actively used. It may also hide actual migrations from clients unable to deal with them by acting as a proxy for a migrated file system for particular clients, at some cost in performance, although it is not obligated to do so. If the server will potentially perform a migration, it MUST set EXCHGID4\_FLAG\_SUPP\_MOVED\_MIGR in eir\_flags.

When EXCHGID4\_FLAG\_BIND\_PRINC\_STATEID is set, the client indicates that it wants the server to bind the stateid to the principal. This means that when a principal creates a stateid, it has to be the one to use the stateid. If the server will perform binding, it will return EXCHGID4\_FLAG\_BIND\_PRINC\_STATEID. The server MAY return EXCHGID4\_FLAG\_BIND\_PRINC\_STATEID even if the client does not request

it. If an update to the client ID changes the value of EXCHGID4\_FLAG\_BIND\_PRINC\_STATEID's client ID property, the effect applies only to new stateids. Existing stateids (and all stateids with the same "other" field) that were created with stateid to principal binding in force will continue to have binding in force. Existing stateids (and all stateids with the same "other" field) that were created with stateid to principal not in force will continue to have binding not in force.

The EXCHGID4\_FLAG\_USE\_NON\_PNFS, EXCHGID4\_FLAG\_USE\_PNFS\_MDS, and EXCHGID4\_FLAG\_USE\_PNFS\_DS bits are described in Section 20.5 and convey roles the client ID is to be used for in a pNFS environment. The server MUST set one of the acceptable combinations of these bits (roles) in `eir_flags`, as specified in that section. Note that the same client owner/server owner pair can have multiple roles. Multiple roles can be associated with the same client ID or with different client IDs. Thus, if a client sends EXCHANGE\_ID from the same client owner to the same server owner multiple times, but specifies different pNFS roles each time, the server might return different client IDs. Given that different pNFS roles might have different client IDs, the client may ask for different properties for each role/client ID.

The `spa_how` field of the `eia_state_protect` field specifies how the client wants to protect its client, locking, and session states from unauthorized changes (Section 7.8.3):

- \* SP4\_NONE. The client does not request the NFSv4.1 server to enforce state protection. The NFSv4.1 server MUST NOT enforce state protection for the returned client ID.
- \* SP4\_MACH\_CRED. If `spa_how` is SP4\_MACH\_CRED, then the client MUST send the EXCHANGE\_ID operation with RPCSEC\_GSS as the security flavor, and with a service of RPC\_GSS\_SVC\_INTEGRITY or RPC\_GSS\_SVC\_PRIVACY. If SP4\_MACH\_CRED is specified, then the client wants to use an RPCSEC\_GSS-based machine credential to protect its state. The server MUST note the principal the EXCHANGE\_ID operation was sent with, and the GSS mechanism used. These notes collectively comprise the machine credential.

After the client ID is confirmed, as long as the lease associated with the client ID is unexpired, a subsequent EXCHANGE\_ID operation that uses the same `eia_clientowner.co_owner` as the first EXCHANGE\_ID MUST also use the same machine credential as the first EXCHANGE\_ID. The server returns the same client ID for the subsequent EXCHANGE\_ID as that returned from the first EXCHANGE\_ID.

- \* SP4\_SSV. If spa\_how is SP4\_SSV, then the client MUST send the EXCHANGE\_ID operation with RPCSEC\_GSS as the security flavor, and with a service of RPC\_GSS\_SVC\_INTEGRITY or RPC\_GSS\_SVC\_PRIVACY. If SP4\_SSV is specified, then the client wants to use the SSV to protect its state. The server records the credential used in the request as the machine credential (as defined above) for the eia\_clientowner.co\_owner. The CREATE\_SESSION operation that confirms the client ID MUST use the same machine credential.

When a client specifies SP4\_MACH\_CRED or SP4\_SSV, it also provides two lists of operations (each expressed as a bitmap). The first list is spo\_must\_enforce and consists of those operations the client MUST send (subject to the server confirming the list of operations in the result of EXCHANGE\_ID) with the machine credential (if SP4\_MACH\_CRED protection is specified) or the SSV-based credential (if SP4\_SSV protection is used). The client MUST send the operations with RPCSEC\_GSS credentials that specify the RPC\_GSS\_SVC\_INTEGRITY or RPC\_GSS\_SVC\_PRIVACY security service. Typically, the first list of operations includes EXCHANGE\_ID, CREATE\_SESSION, DELEGPURGE, DESTROY\_SESSION, BIND\_CONN\_TO\_SESSION, and DESTROY\_CLIENTID. The client should not specify in this list any operations that require a filehandle because the server's access policies MAY conflict with the client's choice, and thus the client would then be unable to access a subset of the server's namespace.

Note that if SP4\_SSV protection is specified, and the client indicates that CREATE\_SESSION must be protected with SP4\_SSV, because the SSV cannot exist without a confirmed client ID, the first CREATE\_SESSION MUST instead be sent using the machine credential, and the server MUST accept the machine credential.

There is a corresponding result, also called spo\_must\_enforce, of the operations for which the server will require SP4\_MACH\_CRED or SP4\_SSV protection. Normally, the server's result equals the client's argument, but the result MAY be different. If the client requests one or more operations in the set { EXCHANGE\_ID, CREATE\_SESSION, DELEGPURGE, DESTROY\_SESSION, BIND\_CONN\_TO\_SESSION, DESTROY\_CLIENTID }, then the result spo\_must\_enforce MUST include the operations the client requested from that set.

If spo\_must\_enforce in the results has BIND\_CONN\_TO\_SESSION set, then connection binding enforcement is enabled, and the client MUST use the machine (if SP4\_MACH\_CRED protection is used) or SSV (if SP4\_SSV protection is used) credential on calls to BIND\_CONN\_TO\_SESSION.

The second list is `spo_must_allow` and consists of those operations the client wants to have the option of sending with the machine credential or the SSV-based credential, even if the object the operations are performed on is not owned by the machine or SSV credential.

The corresponding result, also called `spo_must_allow`, consists of the operations the server will allow the client to use `SP4_SSV` or `SP4_MACH_CRED` credentials with. Normally, the server's result equals the client's argument, but the result MAY be different.

The purpose of `spo_must_allow` is to allow clients to solve the following conundrum. Suppose the client ID is confirmed with `EXCHGID4_FLAG_BIND_PRINC_STATEID`, and it calls `OPEN` with the `RPCSEC_GSS` credentials of a normal user. Now suppose the user's credentials expire, and cannot be renewed (e.g., a Kerberos ticket granting ticket expires, and the user has logged off and will not be acquiring a new ticket granting ticket). The client will be unable to send `CLOSE` without the user's credentials, which is to say the client has to either leave the state on the server or re-send `EXCHANGE_ID` with a new verifier to clear all state, that is, unless the client includes `CLOSE` on the list of operations in `spo_must_allow` and the server agrees.

The `SP4_SSV` protection parameters also have:

`ssp_hash_algs`:

This is the set of algorithms the client supports for the purpose of computing the digests needed for the internal SSV GSS mechanism and for the `SET_SSV` operation. Each algorithm is specified as an object identifier (OID). The REQUIRED algorithms for a server are `id-sha1`, `id-sha224`, `id-sha256`, `id-sha384`, and `id-sha512` [RFC4055].

Due to known weaknesses in `id-sha1`, it is RECOMMENDED that the client specify at least one algorithm within `ssp_hash_algs` other than `id-sha1`.

The algorithm the server selects among the set is indicated in `spi_hash_alg`, a field of `spr_ssv_prot_info`. The field `spi_hash_alg` is an index into the array `ssp_hash_algs`. Because of known the weaknesses in `id-sha1`, it is RECOMMENDED that it not be selected by the server as long as `ssp_hash_algs` contains any other supported algorithm.

If the server does not support any of the offered algorithms, it returns `NFS4ERR_HASH_ALG_UNSUPP`. If `ssp_hash_algs` is empty, the server MUST return `NFS4ERR_INVALID`.

**ssp\_encr\_algs:**

This is the set of algorithms the client supports for the purpose of providing privacy protection for the internal SSV GSS mechanism. Each algorithm is specified as an OID. The REQUIRED algorithm for a server is id-aes256-CBC. The RECOMMENDED algorithms are id-aes192-CBC and id-aes128-CBC [CSOR\_AES]. The selected algorithm is returned in spi\_encr\_alg, an index into ssp\_encr\_algs. If the server does not support any of the offered algorithms, it returns NFS4ERR\_ENCR\_ALG\_UNSUPP. If ssp\_encr\_algs is empty, the server MUST return NFS4ERR\_INVALID. Note that due to previously stated requirements and recommendations on the relationships between key length and hash length, some combinations of RECOMMENDED and REQUIRED encryption algorithm and hash algorithm either should not or MUST NOT be used. Table 21 summarizes the illegal and discouraged combinations.

**ssp\_window:**

This is the number of SSV versions the client wants the server to maintain (i.e., each successful call to SET\_SSV produces a new version of the SSV). If ssp\_window is zero, the server MUST return NFS4ERR\_INVALID. The server responds with spi\_window, which MUST NOT exceed ssp\_window and MUST be at least one. Any requests on the backchannel or fore channel that are using a version of the SSV that is outside the window will fail with an ONC RPC authentication error, and the requester will have to retry them with the same slot ID and sequence ID.

**ssp\_num\_gss\_handles:**

This is the number of RPCSEC\_GSS handles the server should create that are based on the GSS SSV mechanism (See Section 7.9). It is not the total number of RPCSEC\_GSS handles for the client ID. Indeed, subsequent calls to EXCHANGE\_ID will add RPCSEC\_GSS handles. The server responds with a list of handles in spi\_handles. If the client asks for at least one handle and the server cannot create it, the server MUST return an error. The handles in spi\_handles are not available for use until the client ID is confirmed, which could be immediately if EXCHANGE\_ID returns EXCHGID4\_FLAG\_CONFIRMED\_R, or upon successful confirmation from CREATE\_SESSION.

While a client ID can span all the connections that are connected to a server sharing the same eir\_server\_owner.so\_major\_id, the RPCSEC\_GSS handles returned in spi\_handles can only be used on connections connected to a server that returns the same the eir\_server\_owner.so\_major\_id and eir\_server\_owner.so\_minor\_id on each connection. It is permissible for the client to set ssp\_num\_gss\_handles to zero; the client can create more handles with another EXCHANGE\_ID call.

Because each SSV RPCSEC\_GSS handle shares a common SSV GSS context, there are security considerations specific to this situation discussed in Section 7.10.

The seq\_window (See Section 5.2.3.1 of [RFC2203]) of each RPCSEC\_GSS handle in spi\_handle MUST be the same as the seq\_window of the RPCSEC\_GSS handle used for the credential of the RPC request of which the EXCHANGE\_ID operation was sent as a part.

Encryption Algorithm	MUST NOT be combined with	should not be combined with
id-aes128-CBC		id-sha384, id-sha512
id-aes192-CBC	id-sha1	id-sha512
id-aes256-CBC	id-sha1, id-sha224	

Table 21

The arguments include an array of up to one element in length called eia\_client\_impl\_id. If eia\_client\_impl\_id is present, it contains the information identifying the implementation of the client. Similarly, the results include an array of up to one element in length called eir\_server\_impl\_id that identifies the implementation of the server. Servers MUST accept a zero-length eia\_client\_impl\_id array, and clients MUST accept a zero-length eir\_server\_impl\_id array.

A possible use for implementation identifiers would be in diagnostic software that extracts this information in an attempt to identify interoperability problems, performance workload behaviors, or general usage statistics. Since the intent of having access to this information is for planning or general diagnosis only, the client and server MUST NOT interpret this implementation identity information in a way that affects how the implementation interacts with its peer. The client and server are not allowed to depend on the peer's manifesting a particular allowed behavior based on an implementation identifier but are required to interoperate as specified elsewhere in the protocol specification.

Because it is possible that some implementations might violate the protocol specification and interpret the identity information, implementations MUST provide facilities to allow the NFSv4 client and server to be configured to set the contents of the `nfs_impl_id` structures sent to any specified value.

#### 25.35.4. IMPLEMENTATION

A server's client record is a 5-tuple:

1. `co_ownerid`:

The client identifier string, from the `eia_clientowner` structure of the `EXCHANGE_ID4args` structure.

2. `co_verifier`:

A client-specific value used to indicate incarnations (where a client restart represents a new incarnation), from the `eia_clientowner` structure of the `EXCHANGE_ID4args` structure.

3. `principal`:

The principal that was defined in the RPC header's credential and/or verifier at the time the client record was established.

4. `client ID`:

The shorthand client identifier, generated by the server and returned via the `eir_clientid` field in the `EXCHANGE_ID4resok` structure.

5. `confirmed`:

A private field on the server indicating whether or not a client record has been confirmed. A client record is confirmed if there has been a successful `CREATE_SESSION` operation to confirm it. Otherwise, it is unconfirmed. An unconfirmed record is established by an `EXCHANGE_ID` call. Any unconfirmed record that is not confirmed within a lease period should be removed.

The following identifiers represent special values for the fields in the records.

`ownerid_arg`:

The value of the `eia_clientowner.co_ownerid` subfield of the `EXCHANGE_ID4args` structure of the current request.



**verifier\_arg:**

The value of the eia\_clientowner.co\_verifier subfield of the EXCHANGE\_ID4args structure of the current request.

**old\_verifier\_arg:**

A value of the eia\_clientowner.co\_verifier field of a client record received in a previous request; this is distinct from verifier\_arg.

**principal\_arg:**

The value of the RPCSEC\_GSS principal for the current request.

**old\_principal\_arg:**

A value of the principal of a client record as defined by the RPC header's credential or verifier of a previous request. This is distinct from principal\_arg.

**clientid\_ret:**

The value of the eir\_clientid field the server will return in the EXCHANGE\_ID4resok structure for the current request.

**old\_clientid\_ret:**

The value of the eir\_clientid field the server returned in the EXCHANGE\_ID4resok structure for a previous request. This is distinct from clientid\_ret.

**confirmed:**

The client ID has been confirmed.

**unconfirmed:**

The client ID has not been confirmed.

Since EXCHANGE\_ID is a non-idempotent operation, we must consider the possibility that retries occur as a result of a client restart, network partition, malfunctioning router, etc. Retries are identified by the value of the eia\_clientowner field of EXCHANGE\_ID4args, and the method for dealing with them is outlined in the scenarios below.

The scenarios are described in terms of the client record(s) a server has for a given co\_ownerid. Note that if the client ID was created specifying SP4\_SSV state protection and EXCHANGE\_ID as the one of the operations in spo\_must\_allow, then the server MUST authorize EXCHANGE\_IDs with the SSV principal in addition to the principal that created the client ID.

**1. New Owner ID**

If the server has no client records with `eia_clientowner.co_ownerid` matching `ownerid_arg`, and `EXCHGID4_FLAG_UPD_CONFIRMED_REC_A` is not set in the `EXCHANGE_ID`, then a new shorthand client ID (let us call it `clientid_ret`) is generated, and the following unconfirmed record is added to the server's state.

```
{ ownerid_arg, verifier_arg, principal_arg, clientid_ret,
  unconfirmed }
```

Subsequently, the server returns `clientid_ret`.

## 2. Non-Update on Existing Client ID

If the server has the following confirmed record, and the request does not have `EXCHGID4_FLAG_UPD_CONFIRMED_REC_A` set, then the request is the result of a retried request due to a faulty router or lost connection, or the client is trying to determine if it can perform trunking.

```
{ ownerid_arg, verifier_arg, principal_arg, clientid_ret,
  confirmed }
```

Since the record has been confirmed, the client must have received the server's reply from the initial `EXCHANGE_ID` request. Since the server has a confirmed record, and since `EXCHGID4_FLAG_UPD_CONFIRMED_REC_A` is not set, with the possible exception of `eir_server_owner.so_minor_id`, the server returns the same result it did when the client ID's properties were last updated (or if never updated, the result when the client ID was created). The confirmed record is unchanged.

## 3. Client Collision

If `EXCHGID4_FLAG_UPD_CONFIRMED_REC_A` is not set, and if the server has the following confirmed record, then this request is likely the result of a chance collision between the values of the `eia_clientowner.co_ownerid` subfield of `EXCHANGE_ID4args` for two different clients.

```
{ ownerid_arg, *, old_principal_arg, old_clientid_ret, confirmed
}
```

If there is currently no state associated with `old_clientid_ret`, or if there is state but the lease has expired, then this case is effectively equivalent to the New Owner ID case of Section 25.35.4, Paragraph 7, Item 1. The confirmed record is deleted, the `old_clientid_ret` and its lock state are deleted, a new shorthand client ID is generated, and the following unconfirmed record is added to the server's state.

```
{ ownerid_arg, verifier_arg, principal_arg, clientid_ret,
  unconfirmed }
```

Subsequently, the server returns `clientid_ret`.

If `old_clientid_ret` has an unexpired lease with state, then no state of `old_clientid_ret` is changed or deleted. The server returns `NFS4ERR_CLID_INUSE` to indicate that the client should retry with a different value for the `eia_clientowner.co_ownerid` subfield of `EXCHANGE_ID4args`. The client record is not changed.

#### 4. Replacement of Unconfirmed Record

If the `EXCHGID4_FLAG_UPD_CONFIRMED_REC_A` flag is not set, and the server has the following unconfirmed record, then the client is attempting `EXCHANGE_ID` again on an unconfirmed client ID, perhaps due to a retry, a client restart before client ID confirmation (i.e., before `CREATE_SESSION` was called), or some other reason.

```
{ ownerid_arg, *, *, old_clientid_ret, unconfirmed }
```

It is possible that the properties of `old_clientid_ret` are different than those specified in the current `EXCHANGE_ID`. Whether or not the properties are being updated, to eliminate ambiguity, the server deletes the unconfirmed record, generates a new client ID (`clientid_ret`), and establishes the following unconfirmed record:

```
{ ownerid_arg, verifier_arg, principal_arg, clientid_ret,
  unconfirmed }
```

#### 5. Client Restart

If `EXCHGID4_FLAG_UPD_CONFIRMED_REC_A` is not set, and if the server has the following confirmed client record, then this request is likely from a previously confirmed client that has restarted.

```
{ ownerid_arg, old_verifier_arg, principal_arg, old_clientid_ret,
  confirmed }
```

Since the previous incarnation of the same client will no longer be making requests, once the new client ID is confirmed by `CREATE_SESSION`, byte-range locks and share reservations should be released immediately rather than forcing the new incarnation to wait for the lease time on the previous incarnation to expire. Furthermore, session state should be removed since if the client had maintained that information across restart, this request would not have been sent. If the server supports neither the `CLAIM_DELEGATE_PREV` nor `CLAIM_DELEG_PREV_FH` claim types, associated delegations should be purged as well; otherwise, delegations are retained and recovery proceeds according to Section 15.2.1.

After processing, `clientid_ret` is returned to the client and this client record is added:

```
{ ownerid_arg, verifier_arg, principal_arg, clientid_ret,
  unconfirmed }
```

The previously described confirmed record continues to exist, and thus the same `ownerid_arg` exists in both a confirmed and unconfirmed state at the same time. The number of states can collapse to one once the server receives an applicable `CREATE_SESSION` or `EXCHANGE_ID`.

- \* If the server subsequently receives a successful `CREATE_SESSION` that confirms `clientid_ret`, then the server atomically destroys the confirmed record and makes the unconfirmed record confirmed as described in Section 25.36.3.
- \* If the server instead subsequently receives an `EXCHANGE_ID` with the client owner equal to `ownerid_arg`, one strategy is to simply delete the unconfirmed record, and process the `EXCHANGE_ID` as described in the entirety of Section 25.35.4.

## 6. Update

If `EXCHGID4_FLAG_UPD_CONFIRMED_REC_A` is set, and the server has the following confirmed record, then this request is an attempt at an update.

```
{ ownerid_arg, verifier_arg, principal_arg, clientid_ret,
  confirmed }
```

Since the record has been confirmed, the client must have received the server's reply from the initial `EXCHANGE_ID` request. The server allows the update, and the client record is left intact.

#### 7. Update but No Confirmed Record

If EXCHGID4\_FLAG\_UPD\_CONFIRMED\_REC\_A is set, and the server has no confirmed record corresponding ownerid\_arg, then the server returns NFS4ERR\_NOENT and leaves any unconfirmed record intact.

#### 8. Update but Wrong Verifier

If EXCHGID4\_FLAG\_UPD\_CONFIRMED\_REC\_A is set, and the server has the following confirmed record, then this request is an illegal attempt at an update, perhaps because of a retry from a previous client incarnation.

```
{ ownerid_arg, old_verifier_arg, *, clientid_ret, confirmed }
```

The server returns NFS4ERR\_NOT\_SAME and leaves the client record intact.

#### 9. Update but Wrong Principal

If EXCHGID4\_FLAG\_UPD\_CONFIRMED\_REC\_A is set, and the server has the following confirmed record, then this request is an illegal attempt at an update by an unauthorized principal.

```
{ ownerid_arg, verifier_arg, old_principal_arg, clientid_ret, confirmed }
```

The server returns NFS4ERR\_PERM and leaves the client record intact.

### 25.36. Operation 43: CREATE\_SESSION - Create New Session and Confirm Client ID

#### 25.36.1. ARGUMENT

```

struct channel_attrs4 {
    count4          ca_headerpadsize;
    count4          ca_maxrequestsize;
    count4          ca_maxresponsesize;
    count4          ca_maxresponsesize_cached;
    count4          ca_maxoperations;
    count4          ca_maxrequests;
    uint32_t        ca_rdma_ird<1>;
};

const CREATE_SESSION4_FLAG_PERSIST          = 0x00000001;
const CREATE_SESSION4_FLAG_CONN_BACK_CHAN   = 0x00000002;
const CREATE_SESSION4_FLAG_CONN_RDMA        = 0x00000004;

struct CREATE_SESSION4args {
    clientid4      csa_clientid;
    sequenceid4    csa_sequence;

    uint32_t        csa_flags;

    channel_attrs4 csa_fore_chan_attrs;
    channel_attrs4 csa_back_chan_attrs;

    uint32_t        csa_cb_program;
    callback_sec_parms4 csa_sec_parms<>;
};

```

#### 25.36.2. RESULT

```

struct CREATE_SESSION4resok {
    sessionid4      csr_sessionid;
    sequenceid4     csr_sequence;

    uint32_t        csr_flags;

    channel_attrs4  csr_fore_chan_attrs;
    channel_attrs4  csr_back_chan_attrs;
};

union CREATE_SESSION4res switch (nfsstat4 csr_status) {
case NFS4_OK:
    CREATE_SESSION4resok  csr_resok4;
default:
    void;
};

```

## 25.36.3. DESCRIPTION

This operation is used by the client to create new session objects on the server.

CREATE\_SESSION can be sent with or without a preceding SEQUENCE operation in the same COMPOUND procedure. If CREATE\_SESSION is sent with a preceding SEQUENCE operation, any session created by CREATE\_SESSION has no direct relation to the session specified in the SEQUENCE operation, although the two sessions might be associated with the same client ID. If CREATE\_SESSION is sent without a preceding SEQUENCE, then it MUST be the only operation in the COMPOUND procedure's request. If it is not, the server MUST return NFS4ERR\_NOT\_ONLY\_OP.

In addition to creating a session, CREATE\_SESSION has the following effects:

- \* The first session created with a new client ID serves to confirm the creation of that client's state on the server. The server returns the parameter values for the new session.
- \* The connection CREATE\_SESSION that is sent over is associated with the session's fore channel.

The arguments and results of CREATE\_SESSION are described as follows:

`csa_clientid`: This is the client ID with which the new session will be associated. The corresponding result is `csr_sessionid`, the session ID of the new session.

`csa_sequence`: Each client ID serializes CREATE\_SESSION via a per-client ID sequence number (see Section 25.36.4). The corresponding result is `csr_sequence`, which MUST be equal to `csa_sequence`.

In the next three arguments, the client offers a value that is to be a property of the session. Except where stated otherwise, it is RECOMMENDED that the server accept the value. If it is not acceptable, the server MAY use a different value. Regardless, the server MUST return the value the session will use (which will be either what the client offered, or what the server is insisting on) to the client.

`csa_flags`: The `csa_flags` field contains a list of the following flag bits:

CREATE\_SESSION4\_FLAG\_PERSIST:

If `CREATE_SESSION4_FLAG_PERSIST` is set, the client wants the server to provide a persistent reply cache. For sessions in which only idempotent operations will be used (e.g., a read-only session), clients have no reason to set `CREATE_SESSION4_FLAG_PERSIST`. If the server does not or cannot provide a persistent reply cache, the server **MUST NOT** set `CREATE_SESSION4_FLAG_PERSIST` in the field `csr_flags`.

If the server is a pNFS metadata server, for reasons described in Section 18.7.2 it should support `CREATE_SESSION4_FLAG_PERSIST` if it supports the `layout_hint` (Section 11.16.4) attribute.

**CREATE\_SESSION4\_FLAG\_CONN\_BACK\_CHAN:**

If `CREATE_SESSION4_FLAG_CONN_BACK_CHAN` is set in `csa_flags`, the client is requesting that the connection over which the `CREATE_SESSION` operation arrived be associated with the session's backchannel in addition to its fore channel. If the server agrees, it sets `CREATE_SESSION4_FLAG_CONN_BACK_CHAN` in the result field `csr_flags`. If `CREATE_SESSION4_FLAG_CONN_BACK_CHAN` is not set in `csa_flags`, then `CREATE_SESSION4_FLAG_CONN_BACK_CHAN` **MUST NOT** be set in `csr_flags`.

**CREATE\_SESSION4\_FLAG\_CONN\_RDMA:**

If `CREATE_SESSION4_FLAG_CONN_RDMA` is set in `csa_flags`, and if the connection over which the `CREATE_SESSION` operation arrived is currently in non-RDMA mode but has the capability to operate in RDMA mode, then the client is requesting that the server "step up" to RDMA mode on the connection. If the server agrees, it sets `CREATE_SESSION4_FLAG_CONN_RDMA` in the result field `csr_flags`. If `CREATE_SESSION4_FLAG_CONN_RDMA` is not set in `csa_flags`, then `CREATE_SESSION4_FLAG_CONN_RDMA` **MUST NOT** be set in `csr_flags`. Note that once the server agrees to step up, it and the client **MUST** exchange all future traffic on the connection with RPC RDMA framing and not Record Marking ([RFC8166]).

**csa\_fore\_chan\_attrs, csa\_back\_chan\_attrs:** The `csa_fore_chan_attrs` and `csa_back_chan_attrs` fields apply to attributes of the fore channel (which conveys requests originating from the client to the server), and the backchannel (the channel that conveys callback requests originating from the server to the client), respectively. The results are in corresponding structures called `csr_fore_chan_attrs` and `csr_back_chan_attrs`. The results establish attributes for each channel, and on all subsequent use of each channel of the session. Each structure has the following fields:



**ca\_headerpadsize:**

The maximum amount of padding the requester is willing to apply to ensure that write payloads are aligned on some boundary at the replier. For each channel, the server

- \* will reply in ca\_headerpadsize with its preferred value, or zero if padding is not in use, and
- \* MAY decrease this value but MUST NOT increase it.

**ca\_maxrequestsize:**

The maximum size of a COMPOUND or CB\_COMPOUND request that will be sent. This size represents the XDR encoded size of the request, including the RPC headers (including security flavor credentials and verifiers) but excludes any RPC transport framing headers. Imagine a request coming over a non-RDMA TCP/IP connection, and that it has a single Record Marking header preceding it. The maximum allowable count encoded in the header will be ca\_maxrequestsize. If a requester sends a request that exceeds ca\_maxrequestsize, the error NFS4ERR\_REQ\_TOO\_BIG will be returned per the description in Section 7.6.4. For each channel, the server MAY decrease this value but MUST NOT increase it.

**ca\_maxresponsesize:**

The maximum size of a COMPOUND or CB\_COMPOUND reply that the requester will accept from the replier including RPC headers (see the ca\_maxrequestsize definition). For each channel, the server MAY decrease this value, but MUST NOT increase it. However, if the client selects a value for ca\_maxresponsesize such that a replier on a channel could never send a response, the server is to return NFS4ERR\_TOOSMALL in the CREATE\_SESSION reply. After the session is created, if a requester sends a request for which the size of the reply would exceed this value, the replier will return NFS4ERR\_REP\_TOO\_BIG, per the description in Section 7.6.4.

**ca\_maxresponsesize\_cached:**

Like ca\_maxresponsesize, but the maximum size of a reply that will be stored in the reply cache (Section 7.6.1). For each channel, the server MAY decrease this value, but MUST NOT increase it. If, in the reply to CREATE\_SESSION, the value of ca\_maxresponsesize\_cached of a channel is less than the value of ca\_maxresponsesize of the same channel, then this is an indication to the requester that it needs to be selective about which replies it directs the replier to cache; for example, large replies from non-idempotent operations (e.g., COMPOUND requests with a READ operation) should not be cached. The

requester decides which replies to cache via an argument to the SEQUENCE (the `sa_cachethis` field, see Section 25.46) or CB\_SEQUENCE (the `csa_cachethis` field, see Section 27.9) operations. After the session is created, if a requester sends a request for which the size of the reply would exceed `ca_maxresponsesize_cached`, the replier will return NFS4ERR\_REP\_TOO\_BIG\_TO\_CACHE, per the description in Section 7.6.4.

**ca\_maxoperations:**

The maximum number of operations the replier will accept in a COMPOUND or CB\_COMPOUND. For the backchannel, the server MUST NOT change the value the client offers. For the fore channel, the server MAY change the requested value. After the session is created, if a requester sends a COMPOUND or CB\_COMPOUND with more operations than `ca_maxoperations`, the replier MUST return NFS4ERR\_TOO\_MANY\_OPS.

**ca\_maxrequests:**

The maximum number of concurrent COMPOUND or CB\_COMPOUND requests the requester will send on the session. Subsequent requests will each be assigned a slot identifier by the requester within the range zero to `ca_maxrequests - 1` inclusive. For the backchannel, the server MUST NOT change the value the client offers. For the fore channel, the server MAY change the requested value.

**ca\_rdma\_ird:**

This array has a maximum of one element. If this array has one element, then the element contains the inbound RDMA read queue depth (IRD). For each channel, the server MAY decrease this value, but MUST NOT increase it.

**csa\_cb\_program** This is the ONC RPC program number the server MUST use in any callbacks sent through the backchannel to the client. The server MUST specify an ONC RPC program number equal to `csa_cb_program` and an ONC RPC version number equal to 4 in callbacks sent to the client. If a CB\_COMPOUND is sent to the client, the server MUST use a minor version number of 1. There is no corresponding result.

**csa\_sec\_parms** The field `csa_sec_parms` is an array of acceptable security credentials the server can use on the session's backchannel. Three security flavors are supported: AUTH\_NONE, AUTH\_SYS, and RPCSEC\_GSS. If AUTH\_NONE is specified for a credential, then this says the client is authorizing the server to use AUTH\_NONE on all callbacks for the session. If AUTH\_SYS is specified, then the client is authorizing the server to use

AUTH\_SYS on all callbacks, using the credential specified cbsp\_sys\_cred. If RPCSEC\_GSS is specified, then the server is allowed to use the RPCSEC\_GSS context specified in cbsp\_gss\_parms as the RPCSEC\_GSS context in the credential of the RPC header of callbacks to the client. There is no corresponding result.

The RPCSEC\_GSS context for the backchannel is specified via a pair of values of data type gsshandle4\_t. The data type gsshandle4\_t represents an RPCSEC\_GSS handle, and is precisely the same as the data type of the "handle" field of the rpc\_gss\_init\_res data type defined in "Context Creation Response - Successful Acceptance", Section 5.2.3.1 of [RFC2203].

The first RPCSEC\_GSS handle, gcbp\_handle\_from\_server, is the fore handle the server returned to the client (either in the handle field of data type rpc\_gss\_init\_res or as one of the elements of the spi\_handles field returned in the reply to EXCHANGE\_ID) when the RPCSEC\_GSS context was created on the server. The second handle, gcbp\_handle\_from\_client, is the back handle to which the client will map the RPCSEC\_GSS context. The server can immediately use the value of gcbp\_handle\_from\_client in the RPCSEC\_GSS credential in callback RPCs. That is, the value in gcbp\_handle\_from\_client can be used as the value of the field "handle" in data type rpc\_gss\_cred\_t (See "Elements of the RPCSEC\_GSS Security Protocol", Section 5 of [RFC2203]) in callback RPCs. The server MUST use the RPCSEC\_GSS security service specified in gcbp\_service, i.e., it MUST set the "service" field of the rpc\_gss\_cred\_t data type in RPCSEC\_GSS credential to the value of gcbp\_service (see "RPC Request Header", Section 5.3.1 of [RFC2203]).

If the RPCSEC\_GSS handle identified by gcbp\_handle\_from\_server does not exist on the server, the server will return NFS4ERR\_NOENT.

Within each element of csa\_sec\_parms, the fore and back RPCSEC\_GSS contexts MUST share the same GSS context and MUST have the same seq\_window (See Section 5.2.3.1 of RFC 2203 [RFC2203]). The fore and back RPCSEC\_GSS context state are independent of each other as far as the RPCSEC\_GSS sequence number (See the seq\_num field in the rpc\_gss\_cred\_t data type of Sections 5 and 5.3.1 of [RFC2203]).

If an RPCSEC\_GSS handle is using the SSV context (See Section 7.9), then because each SSV RPCSEC\_GSS handle shares a common SSV GSS context, there are security considerations specific to this situation discussed in Section 7.10.

Once the session is created, the first SEQUENCE or CB\_SEQUENCE received on a slot MUST have a sequence ID equal to 1; if not, the replier MUST return NFS4ERR\_SEQ\_MISORDERED.

#### 25.36.4. IMPLEMENTATION

To describe a possible implementation, the same notation for client records introduced in the description of EXCHANGE\_ID is used with the following addition:

`clientid_arg`: The value of the `csa_clientid` field of the `CREATE_SESSION4args` structure of the current request.

Since `CREATE_SESSION` is a non-idempotent operation, we need to consider the possibility that retries may occur as a result of a client restart, network partition, malfunctioning router, etc. For each client ID created by `EXCHANGE_ID`, the server maintains a separate reply cache (called the `CREATE_SESSION` reply cache) similar to the session reply cache used for `SEQUENCE` operations, with two distinctions.

- \* First, this is a reply cache just for detecting and processing `CREATE_SESSION` requests for a given client ID.
- \* Second, the size of the client ID reply cache is of one slot (and as a result, the `CREATE_SESSION` request does not carry a slot number). This means that at most one `CREATE_SESSION` request for a given client ID can be outstanding.

As previously stated, `CREATE_SESSION` can be sent with or without a preceding `SEQUENCE` operation. Even if a `SEQUENCE` precedes `CREATE_SESSION`, the server MUST maintain the `CREATE_SESSION` reply cache, which is separate from the reply cache for the session associated with a `SEQUENCE`. If `CREATE_SESSION` was originally sent by itself, the client MAY send a retry of the `CREATE_SESSION` operation within a `COMPOUND` preceded by a `SEQUENCE`. If `CREATE_SESSION` was originally sent in a `COMPOUND` that started with a `SEQUENCE`, then the client should send a retry in a `COMPOUND` that starts with a `SEQUENCE` that has the same session ID as the `SEQUENCE` of the original request. However, the client MAY send a retry in a `COMPOUND` that either has no preceding `SEQUENCE`, or has a preceding `SEQUENCE` that refers to a different session than the original `CREATE_SESSION`. This might be necessary if the client sends a `CREATE_SESSION` in a `COMPOUND` preceded by a `SEQUENCE` with session ID X, and session X no longer exists. Regardless, any retry of `CREATE_SESSION`, with or without a preceding `SEQUENCE`, MUST use the same value of `csa_sequence` as the original.

After the client received a reply to an EXCHANGE\_ID operation that contains a new, unconfirmed client ID, the server expects the client to follow with a CREATE\_SESSION operation to confirm the client ID. The server expects value of csa\_sequenceid in the arguments to that CREATE\_SESSION to be equal to the value of the field eir\_sequenceid that was returned in results of the EXCHANGE\_ID that returned the unconfirmed client ID. Before the server replies to that EXCHANGE\_ID operation, it initializes the client ID slot to be equal to eir\_sequenceid - 1 (accounting for underflow), and records a contrived CREATE\_SESSION result with a "cached" result of NFS4ERR\_SEQ\_MISORDERED. With the client ID slot thus initialized, the processing of the CREATE\_SESSION operation is divided into four phases:

1. Client record look up. The server looks up the client ID in its client record table. If the server contains no records with client ID equal to clientid\_arg, then most likely the client's state has been purged during a period of inactivity, possibly due to a loss of connectivity. NFS4ERR\_STALE\_CLIENTID is returned, and no changes are made to any client records on the server. Otherwise, the server goes to phase 2.
2. Sequence ID processing. If csa\_sequenceid is equal to the sequence ID in the client ID's slot, then this is a replay of the previous CREATE\_SESSION request, and the server returns the cached result. If csa\_sequenceid is not equal to the sequence ID in the slot, and is more than one greater (accounting for wraparound), then the server returns the error NFS4ERR\_SEQ\_MISORDERED, and does not change the slot. If csa\_sequenceid is equal to the slot's sequence ID + 1 (accounting for wraparound), then the slot's sequence ID is set to csa\_sequenceid, and the CREATE\_SESSION processing goes to the next phase. A subsequent new CREATE\_SESSION call over the same client ID MUST use a csa\_sequenceid that is one greater than the sequence ID in the slot.
3. Client ID confirmation. If this would be the first session for the client ID, the CREATE\_SESSION operation serves to confirm the client ID. Otherwise, the client ID confirmation phase is skipped and only the session creation phase occurs. Any case in which there is more than one record with identical values for client ID represents a server implementation error. Operation in the potential valid cases is summarized as follows.

\* Successful Confirmation

If the server has the following unconfirmed record, then this is the expected confirmation of an unconfirmed record.

```
{ ownerid, verifier, principal_arg, clientid_arg,
  unconfirmed }
```

As noted in Section 25.35.4, the server might also have the following confirmed record.

```
{ ownerid, old_verifier, principal_arg, old_clientid,
  confirmed }
```

The server schedules the replacement of both records with:

```
{ ownerid, verifier, principal_arg, clientid_arg, confirmed
}
```

The processing of CREATE\_SESSION continues on to session creation. Once the session is successfully created, the scheduled client record replacement is committed. If the session is not successfully created, then no changes are made to any client records on the server.

\* Unsuccessful Confirmation

If the server has the following record, then the client has changed principals after the previous EXCHANGE\_ID request, or there has been a chance collision between shorthand client identifiers.

```
{ *, *, old_principal_arg, clientid_arg, * }
```

Neither of these cases is permissible. Processing stops and NFS4ERR\_CLID\_INUSE is returned to the client. No changes are made to any client records on the server.

4. Session creation. The server confirmed the client ID, either in this CREATE\_SESSION operation, or a previous CREATE\_SESSION operation. The server examines the remaining fields of the arguments.

The server creates the session by recording the parameter values used (including whether the CREATE\_SESSION4\_FLAG\_PERSIST flag is set and has been accepted by the server) and allocating space for the session reply cache (if there is not enough space, the server returns NFS4ERR\_NOSPC). For each slot in the reply cache, the server sets the sequence ID to zero, and records an entry containing a COMPOUND reply with zero operations and the error NFS4ERR\_SEQ\_MISORDERED. This way, if the first SEQUENCE request sent has a sequence ID equal to zero, the server can simply return what is in the reply cache: NFS4ERR\_SEQ\_MISORDERED. The

client initializes its reply cache for receiving callbacks in the same way, and similarly, the first CB\_SEQUENCE operation on a slot after session creation MUST have a sequence ID of one.

If the session state is created successfully, the server associates the session with the client ID provided by the client.

When a request that had CREATE\_SESSION4\_FLAG\_CONN\_RDMA set needs to be retried, the retry MUST be done on a new connection that is in non-RDMA mode. If properties of the new connection are different enough that the arguments to CREATE\_SESSION need to change, then a non-retry MUST be sent. The server will eventually dispose of any session that was created on the original connection.

On the backchannel, the client and server might wish to have many slots, in some cases perhaps more than the fore channel, in order to deal with the situations where the network link has high latency and is the primary bottleneck for response to recalls. If so, and if the client provides too few slots to the backchannel, the server might limit the number of recallable objects it gives to the client.

Implementing RPCSEC\_GSS callback support requires changes to both the client and server implementations of RPCSEC\_GSS. One possible set of changes includes:

- \* Adding a data structure that wraps the GSS-API context with a reference count.
- \* New functions to increment and decrement the reference count. If the reference count is decremented to zero, the wrapper data structure and the GSS-API context it refers to would be freed.
- \* Change RPCSEC\_GSS to create the wrapper data structure upon receiving GSS-API context from gss\_accept\_sec\_context() and gss\_init\_sec\_context(). The reference count would be initialized to 1.
- \* Adding a function to map an existing RPCSEC\_GSS handle to a pointer to the wrapper data structure. The reference count would be incremented.
- \* Adding a function to create a new RPCSEC\_GSS handle from a pointer to the wrapper data structure. The reference count would be incremented.

- \* Replacing calls from RPCSEC\_GSS that free GSS-API contexts, with calls to decrement the reference count on the wrapper data structure.

## 25.37. Operation 44: DESTROY\_SESSION - Destroy a Session

### 25.37.1. ARGUMENT

```
struct DESTROY_SESSION4args {  
    sessionid4      dsa_sessionid;  
};
```

### 25.37.2. RESULT

```
struct DESTROY_SESSION4res {  
    nfsstat4      dsr_status;  
};
```

### 25.37.3. DESCRIPTION

The DESTROY\_SESSION operation closes the session and discards the session's reply cache, if any. Any remaining connections associated with the session are immediately disassociated. If the connection has no remaining associated sessions, the connection MAY be closed by the server. Locks, delegations, layouts, wants, and the lease, which are all tied to the client ID, are not affected by DESTROY\_SESSION.

DESTROY\_SESSION MUST be invoked on a connection that is associated with the session being destroyed. In addition, if SP4\_MACH\_CRED state protection was specified when the client ID was created, the RPCSEC\_GSS principal that created the session MUST be the one that destroys the session, using RPCSEC\_GSS privacy or integrity. If SP4\_SSV state protection was specified when the client ID was created, RPCSEC\_GSS using the SSV mechanism (Section 7.9) MUST be used, with integrity or privacy.

If the COMPOUND request starts with SEQUENCE, and if the sessionids specified in SEQUENCE and DESTROY\_SESSION are the same, then

- \* DESTROY\_SESSION MUST be the final operation in the COMPOUND request.
- \* It is advisable to avoid placing DESTROY\_SESSION in a COMPOUND request with other state-modifying operations, because the DESTROY\_SESSION will destroy the reply cache.



- \* Because the session and its reply cache are destroyed, a client that retries the request may receive an error in reply to the retry, even though the original request was successful.

If the COMPOUND request starts with SEQUENCE, and if the sessionids specified in SEQUENCE and DESTROY\_SESSION are different, then DESTROY\_SESSION can appear in any position of the COMPOUND request (except for the first position). The two sessionids can belong to different client IDs.

If the COMPOUND request does not start with SEQUENCE, and if DESTROY\_SESSION is not the sole operation, then server MUST return NFS4ERR\_NOT\_ONLY\_OP.

If there is a backchannel on the session and the server has outstanding CB\_COMPOUND operations for the session which have not been replied to, then the server MAY refuse to destroy the session and return an error. If so, then in the event the backchannel is down, the server should return NFS4ERR\_CB\_PATH\_DOWN to inform the client that the backchannel needs to be repaired before the server will allow the session to be destroyed. Otherwise, the error NFSERR\_BACK\_CHAN\_BUSY will be returned to indicate that there are CB\_COMPOUNDS that need to be replied to. The client should reply to all outstanding CB\_COMPOUNDS before re-sending DESTROY\_SESSION.

## 25.38. Operation 45: FREE\_STATEID - Free Stateid with No Locks

### 25.38.1. ARGUMENT

```
struct FREE_STATEID4args {  
    stateid4      fsa_stateid;  
};
```

### 25.38.2. RESULT

```
struct FREE_STATEID4res {  
    nfsstat4      fsr_status;  
};
```

### 25.38.3. DESCRIPTION

The FREE\_STATEID operation is used to free a stateid that no longer has any associated locks (including opens, byte-range locks, delegations, and layouts). This may be because of client LOCKU operations or because of server revocation. If there are valid locks (of any kind) associated with the stateid in question, the error NFS4ERR\_LOCKS\_HELD will be returned, and the associated stateid will not be freed.

When a stateid is freed that had been associated with revoked locks, by sending the FREE\_STATEID operation, the client acknowledges the loss of those locks. This allows the server, once all such revoked state is acknowledged, to allow that client again to reclaim locks, without encountering the edge conditions discussed in Section 13.4.2.

Once a successful FREE\_STATEID is done for a given stateid, any subsequent use of that stateid will result in an NFS4ERR\_BAD\_STATEID error.

## 25.39. Operation 46: GET\_DIR\_DELEGATION - Get a Directory Delegation

### 25.39.1. ARGUMENT

```
typedef nfstime4 attr_notice4;

struct GET_DIR_DELEGATION4args {
    /* CURRENT_FH: delegated directory */
    bool          gdda_signal_deleg_avail;
    bitmap4       gdda_notification_types;
    attr_notice4  gdda_child_attr_delay;
    attr_notice4  gdda_dir_attr_delay;
    bitmap4       gdda_child_attributes;
    bitmap4       gdda_dir_attributes;
};
```

The values to be set for the following arguments need to set as described below:

- \* gdda\_child\_attr\_delay indicate an acceptable delay between an child attribute change and the sending of a corresponding notification.

If it is zero, prompt notification is requested. It is likely that this request will be denied because of the difficulty of providing such notification in the presence of multiply-linked files.

If the value is below the value returned in the dirent\_notif\_delay attribute the request is ignored and no NOTIFY4\_CHILD\_ATTR notifications will be generated.

- \* gdda\_dir\_attr\_delay indicate an acceptable delay between a directory attribute change and the sending of a corresponding notification.

If it is zero, prompt notification is requested. Not that prompt notification is always used for attributes tied to directory content notification. See the description of `gdda_dir_attributes` below for details.

- \* `gdda_child_attributes` indicates the set of attributes for which notifications (mostly in the form of `NOTIFY4_CHILD_ATTR` notifications) are requested.

Note that a specific set of attributes which are not subject to change are handled separately. These are `creation_time`, `fileid`, `fsid`, and `file_handle`. These attributes, which are not subject to change, are presented as part of the `attributes` field of `notify4_entry` structure reporting on new and deleted files rather than in a separate notification as a result of change.

- \* `gdda_dir_attributes` indicates the set of attributes for which notifications (in the form of `NOTIFY4_DIR_ATTR` notifications) are requested.

Attributes whose update is inherently tied to content modifications are treated specially. These include `modified_time` and `change`. These are always delivered promptly, regardless of the value of `gdda_dir_attr_delay`.

The generation of notifications for these special attributes is not controlled by `NOTIFY4_DIR_ATTR` bit in `gdda_notification_types`. Instead they are generated (or not) based on the content notification to which they would be appended.

#### 25.39.2. RESULT

```

struct GET_DIR_DELEGATION4resok {
    verifier4      gddr_cookieverf;
    /* Stateid for get_dir_delegation */
    stateid4       gddr_stateid;
    /* Which notifications can the server support */
    bitmap4        gddr_notification;
    bitmap4        gddr_child_attributes;
    bitmap4        gddr_dir_attributes;
};

enum gddrnf4_status {
    GDD4_OK          = 0,
    GDD4_UNAVAIL     = 1
};

union GET_DIR_DELEGATION4res_non_fatal
switch (gddrnf4_status gddrnf_status) {
case GDD4_OK:
    GET_DIR_DELEGATION4resok      gddrnf_resok4;
case GDD4_UNAVAIL:
    bool                          gddrnf_will_signal_deleg_avail;
};

union GET_DIR_DELEGATION4res
switch (nfsstat4 gddr_status) {
case NFS4_OK:
    GET_DIR_DELEGATION4res_non_fatal      gddr_res_non_fatal4;
default:
    void;
};

```

### 25.39.3. DESCRIPTION

The GET\_DIR\_DELEGATION operation is used by a client to request a directory delegation. The directory is represented by the current filehandle. The client also specifies whether it wants the server to notify it when the directory changes in certain ways by setting one or more bits in a bitmap. The server may refuse to grant the delegation. In that case, the server will return NFS4ERR\_DIRDELEG\_UNAVAIL. If the server decides to hand out the delegation, it will return a cookie verifier for that directory. If the cookie verifier changes when the client is holding the delegation, the delegation will be recalled unless the client has asked for notification for this event.

The server will also return a directory delegation stateid, gddr\_stateid, as a result of the GET\_DIR\_DELEGATION operation. This stateid will appear in callback messages related to the delegation,

such as notifications and delegation recalls. The client will use this stateid to return the delegation voluntarily or upon recall. A delegation is returned by calling the DELEGRETURN operation.

The server might not be able to support notifications of certain events. If the client asks for such notifications, the server MUST inform the client of its inability to do so as part of the GET\_DIR\_DELEGATION reply by not setting the appropriate bits in the supported notifications bitmask, gddr\_notification, contained in the reply. The server MUST NOT add bits to gddr\_notification that the client did not request.

The GET\_DIR\_DELEGATION operation can be used for both normal and named attribute directories.

If client sets gdda\_signal\_deleg\_avail to TRUE, then it is registering with the client a "want" for a directory delegation. If the delegation is not available, and the server supports and will honor the "want", the results will have gddrnf\_will\_signal\_deleg\_avail set to TRUE and no error will be indicated on return. If so, the client should expect a future CB\_RECALLABLE\_OBJ\_AVAIL operation to indicate that a directory delegation is available. If the server does not wish to honor the "want" or is not able to do so, it returns the error NFS4ERR\_DIRDELEG\_UNAVAIL. If the delegation is immediately available, the server should return it with the response to the operation, rather than via a callback.

When a client makes a request for a directory delegation while it already holds a directory delegation for that directory (including the case where it has been recalled but not yet returned by the client or revoked by the server), the server MUST reply with the value of gddr\_status set to NFS4\_OK, the value of gddrnf\_status set to GDD4\_UNAVAIL, and the value of gddrnf\_will\_signal\_deleg\_avail set to FALSE. The delegation the client held before the request remains intact, and its state is unchanged. The current stateid is not changed (See Section 23.2.3.1.2 for a description of the current stateid).

#### 25.39.4. IMPLEMENTATION

Directory delegations provide the benefit of improving cache consistency of namespace information. This is done through synchronous callbacks. A server must support synchronous callbacks in order to support directory delegations. In addition to that, notifications, which can be either synchronous or asynchronous, provide a way to reduce network traffic as well as improve client performance under certain conditions.

The bitmap `gdda_notification_types` allows the client to request sending of particular notification types and to inform the server of other information relevant to the provision of notifications. For detailed description of the notification, see the appropriate subsection of Section 27.4. The bits, which are defined in Section 16.2.5, can be classified as follows:

- \* Content update notifications can be requested to allow the client to maintain directory information in accord with that on the server, despite ongoing changes on the server.

The client can ask for notifications on addition of entries to a directory (by setting the bit `NOTIFY4_ADD_ENTRY`), notifications on entry removal (`NOTIFY4_REMOVE_ENTRY`), and renames (`NOTIFY4_RENAME_ENTRY`).

If a client is interested in directory entry caching or negative name caching, it can set the `gdda_notification_types` appropriately to its particular need and the server will notify it of all changes that would otherwise invalidate its name cache. The kind of notification a client asks for may depend on the directory size, its rate of change, and the applications being used to access that directory. The enumeration of the conditions under which a client might ask for a notification is out of the scope of this specification.

In addition, the client can ask for notification of other sorts of directory change by setting `NOTIFY4_CHANGE_COOKIE_VERIFIER`. These include changes to cookie verifiers, cookies within the directory, or the order of directory entries.

- \* The client can ask for notification of attribute changes by setting either `NOTIFY4_CHANGE_DIR_ATTRIBUTE` (for changes to directory attributes) or `NOTIFY4_CHANGE_CHILD_ATTRIBUTE` (for change to attributes of objects associated with the entries within the directory)

For attribute notifications, the client will set bits in the `gdda_dir_attributes` bitmap to indicate which attributes it wants to be notified of. If the server does not support notifications for changes to a certain attribute, it should not set that attribute in the supported attribute bitmap specified in the reply (`gddr_dir_attributes`). The client will also set in the `gdda_child_attributes` bitmap the attributes of directory entries it wants to be notified of, and the server will indicate in `gddr_child_attributes` which attributes of directory entries it will notify the client of.

The client will also let the server know if it wants to get the notification as soon as the attribute change occurs or after a certain delay by setting a delay factor; `gdda_child_attr_delay` is for attribute changes to directory entries and `gdda_dir_attr_delay` is for attribute changes to the directory. If this delay factor is set to zero, that indicates to the server that the client wants to be notified of any attribute changes as soon as they occur. If the delay factor is set to N seconds, the server will make a best-effort guarantee that attribute updates are synchronized within N seconds. If the client asks for a delay factor that the server does not support or that may cause significant resource consumption on the server by causing the server to send a lot of notifications, the server should not commit to sending out notifications for attributes and therefore must not set the appropriate bit in the `gddr_child_attributes` and `gddr_dir_attributes` bitmaps in the response.

- \* Authorization notifications are used to inform the client of information useful to determine when the local equivalents of LOOKUP, READDIR, and GETATTR can be considered authorized with the use of ACCESS to check for authorization. These notifications are discussed in Section 16.2.13.

NOTIFY4\_CHANGE\_AUTH is used to inform the client of a necessary change in the handling of authorization for the local equivalents of LOOKUP and READDIR operations. The structure of this notification is described in Section 27.4.11.

NOTIFY4\_CHANGE\_AUTH is used to inform the client of a necessary change in the handling of for the local equivalents of GETATTR operations. The structure of this notification is described in Section 27.4.12.

- \* In addition to requesting particular types of notifications some of the bits in `gdda_notification_types` are used as flags to inform the server of notification-related choices that the client can make. These include NOTIFY4\_GFLAG\_EXTEND and NOTIFY4\_CFLAG\_ORDER.

The bitmap `gddr_notification_types` allows the server to indicate that particular notification types will be sent when necessary and to inform the client of other information useful in connection with the provision of notifications. The bits, which are defined in Section 16.2.5, can be classified as follows:

- \* For bits that have associated notifications, the bit is zero if that notification was not requested and only set to one if that notification was requested and the server undertook to send it when necessary.

These notifications are discussed in Sections 16.2.11 through 16.2.13.

The client MUST use security policy that the directory or its applicable ancestor (Section 6.2) is exported with. If not, the server MUST return NFS4ERR\_WRONGSEC to the operation that both precedes GET\_DIR\_DELEGATION and sets the current filehandle (See Section 6.2).

The directory delegation covers all the entries in the directory except the parent entry. That means if a directory and its parent both hold directory delegations, any changes to the parent will not cause a notification to be sent for the child even though the child's parent entry points to the parent directory.

#### 25.40. Operation 47: GETDEVICEINFO - Get Device Information

##### 25.40.1. ARGUMENT

```
struct GETDEVICEINFO4args {
    deviceid4      gdia_device_id;
    layouttype4    gdia_layout_type;
    count4         gdia_maxcount;
    bitmap4        gdia_notify_types;
};
```

##### 25.40.2. RESULT

```
struct GETDEVICEINFO4resok {
    device_addr4    gdir_device_addr;
    bitmap4         gdir_notification;
};

union GETDEVICEINFO4res switch (nfsstat4 gdir_status) {
case NFS4_OK:
    GETDEVICEINFO4resok      gdir_resok4;
case NFS4ERR_TOOSMALL:
    count4                   gdir_mincount;
default:
    void;
};
```



### 25.40.3. DESCRIPTION

The GETDEVICEINFO operation returns pNFS storage device address information for the specified device ID. The client identifies the device information to be returned by providing the `gdia_device_id` and `gdia_layout_type` that uniquely identify the device. The client provides `gdia_maxcount` to limit the number of bytes for the result. This maximum size represents all of the data being returned within the GETDEVICEINFO4resok structure and includes the XDR overhead. The server may return less data. If the server is unable to return any information within the `gdia_maxcount` limit, the error NFS4ERR\_TOOSMALL will be returned. However, if `gdia_maxcount` is zero, NFS4ERR\_TOOSMALL MUST NOT be returned.

The `da_layout_type` field of the `gdir_device_addr` returned by the server MUST be equal to the `gdia_layout_type` specified by the client. If it is not equal, the client is to ignore the response as invalid and behave as if the server returned an error, even if the client does have support for the layout type returned.

The client also provides a notification bitmap, `gdia_notify_types`, for the device ID mapping notification for which it is interested in receiving; the server must support device ID notifications for the notification request to have affect. The notification mask is composed in the same manner as the bitmap for file attributes (Section 9.3.7). The numbers of bit positions are listed in the `notify_device_type4` enumeration type (Section 27.12). Only two enumerated values of `notify_device_type4` currently apply to GETDEVICEINFO: NOTIFY\_DEVICEID4\_CHANGE and NOTIFY\_DEVICEID4\_DELETE (See Section 27.12).

The notification bitmap applies only to the specified device ID. If a client sends a GETDEVICEINFO operation on a deviceID multiple times, the last notification bitmap is used by the server for subsequent notifications. If the bitmap is zero or empty, then the device ID's notifications are turned off.

If the client wants to just update or turn off notifications, it MAY send a GETDEVICEINFO operation with `gdia_maxcount` set to zero. In that event, if the device ID is valid, the reply's `da_addr_body` field of the `gdir_device_addr` field will be of zero length.

If an unknown device ID is given in `gdia_device_id`, the server returns NFS4ERR\_NOENT. Otherwise, the device address information is returned in `gdir_device_addr`. Finally, if the server supports notifications for device ID mappings, the `gdir_notification` result will contain a bitmap of which notifications it will actually send to the client (via CB\_NOTIFY\_DEVICEID, see Section 27.12).

If NFS4ERR\_TOOSMALL is returned, the results also contain gdir\_mincount. The value of gdir\_mincount represents the minimum size necessary to obtain the device information.

#### 25.40.4. IMPLEMENTATION

Aside from updating or turning off notifications, another use case for gdia\_maxcount being set to zero is to validate a device ID.

The client should request a notification for changes or deletion of a device ID to device address mapping so that the server can allow the client gracefully use a new mapping, without having pending I/O fail abruptly, or force layouts using the device ID to be recalled or revoked.

It is possible that GETDEVICEINFO (and GETDEVICELIST) will race with CB\_NOTIFY\_DEVICEID, i.e., CB\_NOTIFY\_DEVICEID arrives before the client gets and processes the response to GETDEVICEINFO or GETDEVICELIST. The analysis of the race leverages the fact that the server MUST NOT delete a device ID that is referred to by a layout the client has.

- \* CB\_NOTIFY\_DEVICEID deletes a device ID. If the client believes it has layouts that refer to the device ID, then it is possible that layouts referring to the deleted device ID have been revoked. The client should send a TEST\_STATEID request using the stateid for each layout that might have been revoked. If TEST\_STATEID indicates that any layouts have been revoked, the client must recover from layout revocation as described in Section 18.7.6. If TEST\_STATEID indicates that at least one layout has not been revoked, the client should send a GETDEVICEINFO operation on the supposedly deleted device ID to verify that the device ID has been deleted.

If GETDEVICEINFO indicates that the device ID does not exist, then the client assumes the server is faulty and recovers by sending an EXCHANGE\_ID operation. If GETDEVICEINFO indicates that the device ID does exist, then while the server is faulty for sending an erroneous device ID deletion notification, the degree to which it is faulty does not require the client to create a new client ID.

If the client does not have layouts that refer to the device ID, no harm is done. The client should mark the device ID as deleted, and when GETDEVICEINFO or GETDEVICELIST results are received that indicate that the device ID has been in fact deleted, the device ID should be removed from the client's cache.

- \* CB\_NOTIFY\_DEVICEID indicates that a device ID's device addressing mappings have changed. The client should assume that the results from the in-progress GETDEVICEINFO will be stale for the device ID once received, and so it should send another GETDEVICEINFO on the device ID.

#### 25.41. Operation 48: GETDEVICELIST - Get All Device Mappings for a File System

##### 25.41.1. ARGUMENT

```
struct GETDEVICELIST4args {
    /* CURRENT_FH: object belonging to the file system */
    layouttype4      gdla_layout_type;

    /* number of deviceIDs to return */
    count4           gdla_maxdevices;

    nfs_cookie4      gdla_cookie;
    verifier4        gdla_cookieverf;
};
```

##### 25.41.2. RESULT

```
struct GETDEVICELIST4resok {
    nfs_cookie4      gdlr_cookie;
    verifier4        gdlr_cookieverf;
    deviceid4        gdlr_deviceid_list<>;
    bool             gdlr_eof;
};

union GETDEVICELIST4res switch (nfsstat4 gdlr_status) {
case NFS4_OK:
    GETDEVICELIST4resok      gdlr_resok4;
default:
    void;
};
```

##### 25.41.3. DESCRIPTION

This operation is used by the client to enumerate all of the device IDs that a server's file system uses.

The client provides a current filehandle of a file object that belongs to the file system (i.e., all file objects sharing the same fsid as that of the current filehandle) and the layout type in `gdla_layout_type`. Since this operation might require multiple calls to enumerate all the device IDs (and is thus similar to the READDIR

(Section 25.23) operation), the client also provides `gdia_cookie` and `gdia_cookieverf` to specify the current cursor position in the list. When the client wants to read from the beginning of the file system's device mappings, it sets `gdla_cookie` to zero. The field `gdla_cookieverf` MUST be ignored by the server when `gdla_cookie` is zero. The client provides `gdla_maxdevices` to limit the number of device IDs in the result. If `gdla_maxdevices` is zero, the server MUST return `NFS4ERR_INVALID`. The server MAY return fewer device IDs.

The successful response to the operation will contain the cookie, `gdlr_cookie`, and the cookie verifier, `gdlr_cookieverf`, to be used on the subsequent `GETDEVICELIST`. A `gdlr_eof` value of `TRUE` signifies that there are no remaining entries in the server's device list. Each element of `gdlr_deviceid_list` contains a device ID.

#### 25.41.4. IMPLEMENTATION

An example of the use of this operation is for pNFS clients and servers that use `LAYOUT4_BLOCK_VOLUME` layouts. In these environments it may be helpful for a client to determine device accessibility upon first file system access.

### 25.42. Operation 49: `LAYOUTCOMMIT` - Commit Writes Made Using a Layout

#### 25.42.1. ARGUMENT

```
union newtime4 switch (bool nt_timechanged) {
case TRUE:
    nfstime4          nt_time;
case FALSE:
    void;
};

union newoffset4 switch (bool no_newoffset) {
case TRUE:
    offset4           no_offset;
case FALSE:
    void;
};

struct LAYOUTCOMMIT4args {
    /* CURRENT_FH: file */
    offset4           loca_offset;    /* Unused */
    length4           loca_length;    /* Unused */
    bool              loca_reclaim;
    stateid4          loca_stateid;
    newoffset4        loca_last_write_offset;
    newtime4          loca_time_modify;
    layoutupdate4     loca_layoutupdate;
};
```

#### 25.42.2. RESULT

```
union newsize4 switch (bool ns_sizechanged) {
case TRUE:
    length4          ns_size;
case FALSE:
    void;
};

struct LAYOUTCOMMIT4resok {
    newsize4          locr_newsize;
};

union LAYOUTCOMMIT4res switch (nfsstat4 locr_status) {
case NFS4_OK:
    LAYOUTCOMMIT4resok    locr_resok4;
default:
    void;
};
```

### 25.42.3. DESCRIPTION

The LAYOUTCOMMIT operation commits changes made using the layout represented by the current filehandle, client ID (derived from the session ID in the preceding SEQUENCE operation), and stateid.

Because of important differences among layout types in handling of storage allocation and possible server- coordination facilities, this operation is inherently layout-type-specific, even though all layout types share a common function and a core set of parameters. Despite this shared core, many aspects of the functioning of LAYOUTCOMMIT, including the circumstances when it is required, are defined by the specification of the layout type. See Section 19.1.8 for details relating to places where the layout type specification is referred to below.

In general, LAYOUTCOMMIT commits the entire layout; layout type-specific data (loca\_layoutupdate) may specify a smaller scope of data that is to be committed. It is the responsibility of the specification of the layout type involved to specify the possible forms and effects of such limits, including any restrictions as to layouts that need to be held by the client, if these are affected by limits presented by the layout-type-specific data).

The loca\_offset and loca\_length arguments are no longer used. The client should set both loca\_offset and loca\_length to 0. The server is to ignore the loca\_offset and loca\_length arguments. The client MUST hold the layout designated by loca\_stateid, having it previously been granted via LAYOUTGET (Section 25.43), with an iomode of LAYOUTIOMODE4\_RW. For the case where the client does not hold the specified layout, the server MUST return the error NFS4ERR\_BAD\_LAYOUT. Otherwise, where the specified layout does not have an iomode of LAYOUTIOMODE4\_RW, the server MUST return the error NFS4ERR\_BAD\_IOMODE.

The LAYOUTCOMMIT operation indicates that the client has completed writes using a layout obtained by a previous LAYOUTGET. The client may have only written a subset of the data range it previously requested. For layout types for which provisional allocation is valid, as defined by the layout type specification, LAYOUTCOMMIT allows it to commit or discard provisionally allocated space and to update the server with a new end-of-file. The layout referenced by LAYOUTCOMMIT is still valid after the operation completes and can be continued to be referenced using the same the client ID, filehandle, byte-range, layout type, and stateid.

If the `loca_reclaim` field is set to `TRUE`, this indicates that the client is attempting to commit changes to a layout after the restart of the metadata server during the metadata server's recovery grace period (See Section 18.9.4). This type of request may be necessary when the client has uncommitted writes to provisionally allocated byte-ranges of a file that were sent to the storage devices before the restart of the metadata server. In this case, the layout provided by the client **MUST** be a subset of a writable layout that the client held immediately before the restart of the metadata server. The value of the field `loca_stateid` **MUST** be a value that the metadata server returned before it restarted. The metadata server is free to accept or reject this request based on its own internal metadata consistency checks. If the metadata server finds that the layout provided by the client does not pass its consistency checks, it **MUST** reject the request with the status `NFS4ERR_RECLAIM_BAD`. The successful completion of the `LAYOUTCOMMIT` request with `loca_reclaim` set to `TRUE` does **NOT** provide the client with a layout for the file. It simply commits the changes to the layout specified in the `loca_layoutupdate` field. To obtain a layout for the file, the client must send a `LAYOUTGET` request to the server after the server's grace period has expired. If the metadata server receives a `LAYOUTCOMMIT` request with `loca_reclaim` set to `TRUE` when the metadata server is not in its recovery grace period, it **MUST** reject the request with the status `NFS4ERR_NO_GRACE`.

Setting the `loca_reclaim` field to `TRUE` is required if and only if the committed layout was acquired before the metadata server restart. If the client is committing a layout that was acquired during the metadata server's grace period, it **MUST** set the "reclaim" field to `FALSE`.

The `loca_stateid` is a layout stateid value as returned by previously successful layout operations (See Section 18.7.3).

The `loca_last_write_offset` field specifies the offset of the last byte written by the client previous to the `LAYOUTCOMMIT`. Note that this value is never equal to the file's size (at most it is one byte less than the file's size) and **MUST** be less than or equal to `NFS4_MAXFILEOFF`. The metadata server uses this information to determine whether the file's size needs to be updated. If the metadata server updates the file's size as the result of the `LAYOUTCOMMIT` operation, it must return the new size (`loca_newsize.ns_size`) as part of the results.

The `loca_time_modify` field provide a way for the client to suggest a modification time it would like the metadata server to set, whether this facility is allowed is specified by the layout type specification. The metadata server can, subject to the layout type

specification, use the suggestion but if not it will use the time of the LAYOUTCOMMIT operation to set the modification time. If the metadata server uses the client-provided modification time, it need to ensure that time does not flow backwards, either immediately or as a consequence of accepting a time in the future resulting in time flowing backward once it is set to the time of the LAYOUTCOMMIT later. In any case, layout type specifications allowing this need to address issues related to the lack of synchronization of client and MDS clocks.

The handling of the change attribute is similar, For file systems that derive the change attribute from the modified time, the same procedure should be used for modified times deriving from LAYOUTCOMMIT. In other cases, the same approach used for WRITES to the MDS should be followed for LAYOUTCOMMITs.

If the client wants to force the metadata server to set an exact time, the client should use a SETATTR operation in a COMPOUND right after LAYOUTCOMMIT. See Section 18.7.4 for more details. If the client desires the resultant modification time or change attribute, it should construct the COMPOUND so that a GETATTR follows the LAYOUTCOMMIT.

The `loca_layoutupdate` argument to LAYOUTCOMMIT provides a mechanism for a client to provide layout-specific updates to the metadata server. For example, the layout update can describe what byte-ranges of the original layout have been used and what byte-ranges can be deallocated.

The layout information is more verbose for block devices than for objects and files because the latter two hide the details of block allocation behind their storage protocols. At the minimum, the client needs to communicate changes to the end-of-file location back to the server, and, if desired, its view of the file's modification time. For block/volume layouts, it needs to specify precisely which blocks have been used.

If the layout identified in the arguments does not exist, the error NFS4ERR\_BADLAYOUT is returned. The layout being committed will also be rejected if it does not correspond to an existing layout with an `iomode` of LAYOUTIOMODE4\_RW.

On success, the current filehandle retains its value and the current `stateid` retains its value.



#### 25.42.4. IMPLEMENTATION

The client MAY also use LAYOUTCOMMIT with the loca\_reclaim field set to TRUE to convey hints to modified file attributes or to report layout-type specific information such as I/O errors for object-based storage layouts, as normally done during normal operation. Doing so may help the metadata server to recover files more efficiently after restart. For example, some file system implementations may require expansive recovery of file system objects if the metadata server does not get a positive indication from all clients holding a LAYOUTIOMODE4\_RW layout that they have successfully completed all their writes. Sending a LAYOUTCOMMIT (if required) and then following with LAYOUTRETURN can provide such an indication and allow for graceful and efficient recovery.

If loca\_reclaim is TRUE, the metadata server is free to either examine or ignore the value in the field loca\_stateid. The metadata server implementation might or might not encode in its layout stateid information that allows the metadata server to perform a consistency check on the LAYOUTCOMMIT request.

#### 25.43. Operation 50: LAYOUTGET - Get Layout Information

##### 25.43.1. ARGUMENT

```
struct LAYOUTGET4args {  
    /* CURRENT_FH: file */  
    bool                loga_signal_layout_avail;  
    layouttype4         loga_layout_type;  
    layoutiomode4       loga_iomode;  
    offset4             loga_offset;  
    length4             loga_length;  
    length4             loga_minlength;  
    stateid4            loga_stateid;  
    count4              loga_maxcount;  
};
```

##### 25.43.2. RESULT

```
struct LAYOUTGET4resok {
    bool        logr_return_on_close;
    stateid4    logr_stateid;
    layout4     logr_layout<>;
};

union LAYOUTGET4res switch (nfsstat4 logr_status) {
case NFS4_OK:
    LAYOUTGET4resok    logr_resok4;
case NFS4ERR_LAYOUTTRYLATER:
    bool                logr_will_signal_layout_avail;
default:
    void;
};
```

### 25.43.3. DESCRIPTION

The LAYOUTGET operation requests a layout from the metadata server for reading or writing the file given by the filehandle at the byte-range specified by offset and length. Layouts are identified by the client ID (derived from the session ID in the preceding SEQUENCE operation), current filehandle, layout type (`loga_layout_type`), and the layout stateid (`loga_stateid`). The use of the `loga_iomode` field depends upon the layout type, but should reflect the client's data access intent.

If the metadata server is in a grace period, and does not persist layouts and device ID to device address mappings, then it MUST return NFS4ERR\_GRACE (See Section 13.4.2.1).

The LAYOUTGET operation returns layout information for the specified byte-range: a layout. The client actually specifies two ranges, both starting at the offset in the `loga_offset` field. The first range is between `loga_offset` and `loga_offset + loga_length - 1` inclusive. This range indicates the desired range the client wants the layout to cover. The second range is between `loga_offset` and `loga_offset + loga_minlength - 1` inclusive. This range indicates the required range the client needs the layout to cover. Thus, `loga_minlength` MUST be less than or equal to `loga_length`.

When a length field is set to NFS4\_UINT64\_MAX, this indicates a desire (when `loga_length` is NFS4\_UINT64\_MAX) or requirement (when `loga_minlength` is NFS4\_UINT64\_MAX) to get a layout from `loga_offset` through the end-of-file, regardless of the file's length.

The following rules govern the relationships among, and the minima of, `loga_length`, `loga_minlength`, and `loga_offset`.

- \* If `loga_length` is less than `loga_minlength`, the metadata server MUST return `NFS4ERR_INVAL`.
- \* If `loga_minlength` is zero, this is an indication to the metadata server that the client desires any layout at offset `loga_offset` or less that the metadata server has "readily available". Readily is subjective, and depends on the layout type and the pNFS server implementation. For example, some metadata servers might have to pre-allocate stable storage when they receive a request for a range of a file that goes beyond the file's current length. If `loga_minlength` is zero and `loga_length` is greater than zero, this tells the metadata server what range of the layout the client would prefer to have. If `loga_length` and `loga_minlength` are both zero, then the client is indicating that it desires a layout of any length with the ending offset of the range no less than the value specified `loga_offset`, and the starting offset at or below `loga_offset`. If the metadata server does not have a layout that is readily available, then it MUST return `NFS4ERR_LAYOUTTRYLATER`.
- \* If the sum of `loga_offset` and `loga_minlength` exceeds `NFS4_UINT64_MAX`, and `loga_minlength` is not `NFS4_UINT64_MAX`, the error `NFS4ERR_INVAL` MUST result.
- \* If the sum of `loga_offset` and `loga_length` exceeds `NFS4_UINT64_MAX`, and `loga_length` is not `NFS4_UINT64_MAX`, the error `NFS4ERR_INVAL` MUST result.

After the metadata server has performed the above checks on `loga_offset`, `loga_minlength`, and `loga_offset`, the metadata server MUST return a layout according to the rules in Table 22.

Acceptable layouts based on `loga_minlength`. Note: `u64m` = `NFS4_UINT64_MAX`; `a_off` = `loga_offset`; `a_minlen` = `loga_minlength`.

Layout iomode of request	Layout a_minlen of request	Layout iomode of reply	Layout offset of reply	Layout length of reply
_READ	u64m	MAY be _READ	MUST be <= a_off	MUST be >= file length - layout offset
_READ	u64m	MAY be _RW	MUST be <= a_off	MUST be u64m
_READ	> 0 and < u64m	MAY be _READ	MUST be <= a_off	MUST be >= MIN(file length, a_minlen + a_off) - layout offset
_READ	> 0 and < u64m	MAY be _RW	MUST be <= a_off	MUST be >= a_off - layout offset + a_minlen
_READ	0	MAY be _READ	MUST be <= a_off	MUST be > 0
_READ	0	MAY be _RW	MUST be <= a_off	MUST be > 0
_RW	u64m	MUST be _RW	MUST be <= a_off	MUST be u64m
_RW	> 0 and < u64m	MUST be _RW	MUST be <= a_off	MUST be >= a_off - layout offset + a_minlen
_RW	0	MUST be _RW	MUST be <= a_off	MUST be > 0

Table 22

If loga\_minlength is not zero and the metadata server cannot return a layout according to the rules in Table 22, then the metadata server MUST return the error NFS4ERR\_BADLAYOUT. If loga\_minlength is zero and the metadata server cannot or will not return a layout according to the rules in Table 22, then the metadata server MUST return the error NFS4ERR\_LAYOUTTRYLATER. Assuming that loga\_length is greater than loga\_minlength or equal to zero, the metadata server should return a layout according to the rules in Table 23.

Desired layouts based on `loga_length`. The rules of Table 22 MUST be applied first. Note: `u64m = NFS4_UINT64_MAX`; `a_off = loga_offset`; `a_len = loga_length`.

Layout iomode of request	Layout a_len of request	Layout iomode of reply	Layout offset of reply	Layout length of reply
<code>_READ</code>	<code>u64m</code>	MAY be <code>_READ</code>	MUST be <code>&lt;= a_off</code>	should be <code>u64m</code>
<code>_READ</code>	<code>u64m</code>	MAY be <code>_RW</code>	MUST be <code>&lt;= a_off</code>	should be <code>u64m</code>
<code>_READ</code>	<code>&gt; 0 and &lt; u64m</code>	MAY be <code>_READ</code>	MUST be <code>&lt;= a_off</code>	should be <code>&gt;= a_off - layout offset + a_len</code>
<code>_READ</code>	<code>&gt; 0 and &lt; u64m</code>	MAY be <code>_RW</code>	MUST be <code>&lt;= a_off</code>	should be <code>&gt;= a_off - layout offset + a_len</code>
<code>_READ</code>	<code>0</code>	MAY be <code>_READ</code>	MUST be <code>&lt;= a_off</code>	should be <code>&gt; a_off - layout offset</code>
<code>_READ</code>	<code>0</code>	MAY be <code>_READ</code>	MUST be <code>&lt;= a_off</code>	should be <code>&gt; a_off - layout offset</code>
<code>_RW</code>	<code>u64m</code>	MUST be <code>_RW</code>	MUST be <code>&lt;= a_off</code>	should be <code>u64m</code>
<code>_RW</code>	<code>&gt; 0 and &lt; u64m</code>	MUST be <code>_RW</code>	MUST be <code>&lt;= a_off</code>	should be <code>&gt;= a_off - layout offset + a_len</code>
<code>_RW</code>	<code>0</code>	MUST be <code>_RW</code>	MUST be <code>&lt;= a_off</code>	should be <code>&gt; a_off - layout offset</code>

Table 23

The `loga_stateid` field specifies a valid stateid. If a layout is not currently held by the client, the `loga_stateid` field represents a stateid reflecting the correspondingly valid open, byte-range lock,

or delegation stateid. Once a layout is held on the file by the client, the loga\_stateid field MUST be a stateid as returned from a previous LAYOUTGET or LAYOUTRETURN operation or provided by a CB\_LAYOUTRECALL operation (See Section 18.7.3).

The loga\_maxcount field specifies the maximum layout size (in bytes) that the client can handle. If the size of the layout structure exceeds the size specified by maxcount, the metadata server will return the NFS4ERR\_TOOSMALL error.

The returned layout is expressed as an array, logr\_layout, with each element of type layout4. If a file has a single striping pattern, then logr\_layout should contain just one entry. Otherwise, if the requested range overlaps more than one striping pattern, logr\_layout will contain the required number of entries. The elements of logr\_layout MUST be sorted in ascending order of the value of the lo\_offset field of each element. There MUST be no gaps or overlaps in the range between two successive elements of logr\_layout. The lo\_iomode field in each element of logr\_layout MUST be the same.

Table 22 and Table 23 both refer to a returned layout iomode, offset, and length. Because the returned layout is encoded in the logr\_layout array, more description is required.

**iomode** The value of the returned layout iomode listed in Table 22 and Table 23 is equal to the value of the lo\_iomode field in each element of logr\_layout. As shown in Table 22 and Table 23, the metadata server MAY return a layout with an lo\_iomode different from the requested iomode (field loga\_iomode of the request). If it does so, it MUST ensure that the lo\_iomode is more permissive than the loga\_iomode requested. For example, this behavior allows an implementation to upgrade LAYOUTIOMODE4\_READ requests to LAYOUTIOMODE4\_RW requests at its discretion, within the limits of the layout type specific protocol. A lo\_iomode of either LAYOUTIOMODE4\_READ or LAYOUTIOMODE4\_RW MUST be returned.

**offset** The value of the returned layout offset listed in Table 22 and Table 23 is always equal to the lo\_offset field of the first element logr\_layout.

**length** When setting the value of the returned layout length, the situation is complicated by the possibility that the special layout length value NFS4\_UINT64\_MAX is involved. For a logr\_layout array of N elements, the lo\_length field in the first N-1 elements MUST NOT be NFS4\_UINT64\_MAX. The lo\_length field of the last element of logr\_layout can be NFS4\_UINT64\_MAX under some conditions as described in the following list.

- \* If an applicable rule of Table 22 states that the metadata server MUST return a layout of length NFS4\_UINT64\_MAX, then the lo\_length field of the last element of logr\_layout MUST be NFS4\_UINT64\_MAX.
- \* If an applicable rule of Table 22 states that the metadata server MUST NOT return a layout of length NFS4\_UINT64\_MAX, then the lo\_length field of the last element of logr\_layout MUST NOT be NFS4\_UINT64\_MAX.
- \* If an applicable rule of Table 23 states that the metadata server should return a layout of length NFS4\_UINT64\_MAX, then the lo\_length field of the last element of logr\_layout should be NFS4\_UINT64\_MAX.
- \* When the value of the returned layout length of Table 22 and Table 23 is not NFS4\_UINT64\_MAX, then the returned layout length is equal to the sum of the lo\_length fields of each element of logr\_layout.

Once a LAYOUTGET operation returns with logr\_return\_on\_close set to TRUE for a given file, then all subsequent LAYOUTGET requests by that client for the same file and layout type, MUST reply with logr\_return\_on\_close set to TRUE until the client returns all its open state for that file using CLOSE and DELEGRETURN. Note that return\_on\_close also applies retroactively to all layout segments retrieved by the client for that file and layout type.

After the client has closed all open stateids and returned the delegation stateids for a file for which logr\_return\_on\_close was set to TRUE, the server MUST invalidate all layout segments that were issued to the client for that file. The client MUST NOT attempt to use that layout or the layout stateid.

If the server needs to revoke all open stateids and delegation stateids owned by the client for a file for which logr\_return\_on\_close was set to TRUE, then it MUST also revoke all layout segments of type loga\_layout\_type that were issued for that file to that client, and take action to fence the access to the DSes

The logr\_stateid stateid is returned to the client for use in subsequent layout related operations. See Sections 13.2, 18.7.3, and 18.7.5.2 for a further discussion and requirements.

The format of the returned layout (lo\_content) is specific to the layout type. The value of the layout type (lo\_content.loc\_type) for each of the elements of the array of layouts returned by the metadata server (logr\_layout) MUST be equal to the loga\_layout\_type specified

by the client. If it is not equal, the client should ignore the response as invalid and behave as if the metadata server returned an error, even if the client does have support for the layout type returned.

If neither the requested file nor its containing file system support layouts, the metadata server MUST return NFS4ERR\_LAYOUTUNAVAILABLE. If the layout type is not supported, the metadata server MUST return NFS4ERR\_UNKNOWN\_LAYOUTTYPE. If layouts are supported but no layout matches the client provided layout identification, the metadata server MUST return NFS4ERR\_BADLAYOUT. If an invalid loga\_iomode is specified, or a loga\_iomode of LAYOUTIOMODE4\_ANY is specified, the metadata server MUST return NFS4ERR\_BADIOMODE.

If the layout for the file is unavailable due to transient conditions, e.g., file sharing prohibits layouts, the metadata server MUST return NFS4ERR\_LAYOUTTRYLATER.

If the layout request is rejected due to an overlapping layout recall, the metadata server MUST return NFS4ERR\_RECALLCONFLICT. See Section 18.7.5.2 for details.

If the layout conflicts with a mandatory byte-range lock held on the file, and if the storage devices have no method of enforcing mandatory locks, other than through the restriction of layouts, the metadata server should return NFS4ERR\_LOCKED.

If client sets loga\_signal\_layout\_avail to TRUE, then it is registering with the client a "want" for a layout in the event the layout cannot be obtained due to resource exhaustion. If the metadata server supports and will honor the "want", the results will have logr\_will\_signal\_layout\_avail set to TRUE. If so, the client should expect a CB\_RECALLABLE\_OBJ\_AVAIL operation to indicate that a layout is available.

On success, the current filehandle retains its value and the current stateid is updated to match the value as returned in the results.

#### 25.43.4. IMPLEMENTATION

Typically, LAYOUTGET will be called as part of a COMPOUND request after an OPEN operation and results in the client having location information for the file. This requires that loga\_stateid be set to the special stateid that tells the metadata server to use the current stateid, which is set by OPEN (See Section 23.2.3.1.2). A client may also hold a layout across multiple OPENS. The client specifies a layout type that limits what kind of layout the metadata server will return. This prevents metadata servers from granting layouts that



are unusable by the client.

As indicated by Table 22 and Table 23, the specification of LAYOUTGET allows a pNFS client and server considerable flexibility. A pNFS client can take several strategies for sending LAYOUTGET. Some examples are as follows.

- \* If LAYOUTGET is preceded by OPEN in the same COMPOUND request and the OPEN requests OPEN4\_SHARE\_ACCESS\_READ access, the client might opt to request a \_READ layout with loga\_offset set to zero, loga\_minlength set to zero, and loga\_length set to NFS4\_UINT64\_MAX. If the file has space allocated to it, that space is striped over one or more storage devices, and there is either no conflicting layout or the concept of a conflicting layout does not apply to the pNFS server's layout type or implementation, then the metadata server might return a layout with a starting offset of zero, and a length equal to the length of the file, if not NFS4\_UINT64\_MAX. If the length of the file is not a multiple of the pNFS server's stripe width (See Section 20.6 for a formal definition), the metadata server might round up the returned layout's length.
- \* If LAYOUTGET is preceded by OPEN in the same COMPOUND request, and the OPEN requests OPEN4\_SHARE\_ACCESS\_WRITE access and does not truncate the file, the client might opt to request a \_RW layout with loga\_offset set to zero, loga\_minlength set to zero, and loga\_length set to the file's current length (if known), or NFS4\_UINT64\_MAX. As with the previous case, under some conditions the metadata server might return a layout that covers the entire length of the file or beyond.
- \* This strategy is as above, but the OPEN truncates the file. In this case, the client might anticipate it will be writing to the file from offset zero, and so loga\_offset and loga\_minlength are set to zero, and loga\_length is set to the value of threshold4\_write\_iosize. The metadata server might return a layout from offset zero with a length at least as long as threshold4\_write\_iosize.
- \* A process on the client invokes a request to read from offset 10000 for length 50000. The client is using buffered I/O, and has buffer sizes of 4096 bytes. The client intends to map the request of the process into a series of READ requests starting at offset 8192. The end offset needs to be higher than  $10000 + 50000 = 60000$ , and the next offset that is a multiple of 4096 is 61440. The difference between 61440 and that starting offset of the layout is 53248 (which is the product of 4096 and 15). The value of threshold4\_read\_iosize is less than 53248, so the client sends

a LAYOUTGET request with `loga_offset` set to 8192, `loga_minlength` set to 53248, and `loga_length` set to the file's length (if known) minus 8192 or `NFS4_UINT64_MAX` (if the file's length is not known). Since this LAYOUTGET request exceeds the metadata server's threshold, it grants the layout, possibly with an initial offset of zero, with an end offset of at least  $8192 + 53248 - 1 = 61439$ , but preferably a layout with an offset aligned on the stripe width and a length that is a multiple of the stripe width.

- \* This strategy is as above, but the client is not using buffered I/O, and instead all internal I/O requests are sent directly to the server. The LAYOUTGET request has `loga_offset` equal to 10000 and `loga_minlength` set to 50000. The value of `loga_length` is set to the length of the file. The metadata server is free to return a layout that fully overlaps the requested range, with a starting offset and length aligned on the stripe width.
- \* Again, a process on the client invokes a request to read from offset 10000 for length 50000 (i.e. a range with a starting offset of 10000 and an ending offset of 69999), and buffered I/O is in use. The client is expecting that the server might not be able to return the layout for the full I/O range. The client intends to map the request of the process into a series of thirteen READ requests starting at offset 8192, each with length 4096, with a total length of 53248 (which equals  $13 * 4096$ ), which fully contains the range that client's process wants to read. Because the value of `threshold4_read_iosize` is equal to 4096, it is practical and reasonable for the client to use several LAYOUTGET operations to complete the series of READs. The client sends a LAYOUTGET request with `loga_offset` set to 8192, `loga_minlength` set to 4096, and `loga_length` set to 53248 or higher. The server will grant a layout possibly with an initial offset of zero, with an end offset of at least  $8192 + 4096 - 1 = 12287$ , but preferably a layout with an offset aligned on the stripe width and a length that is a multiple of the stripe width. This will allow the client to make forward progress, possibly sending more LAYOUTGET operations for the remainder of the range.

- \* An NFS client detects a sequential read pattern, and so sends a LAYOUTGET operation that goes well beyond any current or pending read requests to the server. The server might likewise detect this pattern, and grant the LAYOUTGET request. Once the client reads from an offset of the file that represents 50% of the way through the range of the last layout it received, in order to avoid stalling I/O that would wait for a layout, the client sends more operations from an offset of the file that represents 50% of the way through the last layout it received. The client continues to request layouts with byte-ranges that are well in advance of the byte-ranges of recent and/or read requests of processes running on the client.
- \* This strategy is as above, but the client fails to detect the pattern, but the server does. The next time the metadata server gets a LAYOUTGET, it returns a layout with a length that is well beyond `loga_minlength`.
- \* A client is using buffered I/O, and has a long queue of write-behinds to process and also detects a sequential write pattern. It sends a LAYOUTGET for a layout that spans the range of the queued write-behinds and well beyond, including ranges beyond the filer's current length. The client continues to send LAYOUTGET operations once the write-behind queue reaches 50% of the maximum queue length.

Once the client has obtained a layout referring to a particular device ID, the metadata server MUST NOT delete the device ID until the layout is returned or revoked.

CB\_NOTIFY\_DEVICEID can race with LAYOUTGET. One race scenario is that LAYOUTGET returns a device ID for which the client does not have device address mappings, and the metadata server sends a CB\_NOTIFY\_DEVICEID to add the device ID to the client's awareness and meanwhile the client sends GETDEVICEINFO on the device ID. This scenario is discussed in Section 25.40.4. Another scenario is that the CB\_NOTIFY\_DEVICEID is processed by the client before it processes the results from LAYOUTGET. The client will send a GETDEVICEINFO on the device ID. If the results from GETDEVICEINFO are received before the client gets results from LAYOUTGET, then there is no longer a race. If the results from LAYOUTGET are received before the results from GETDEVICEINFO, the client can either wait for results of GETDEVICEINFO or send another one to get possibly more up-to-date device address mappings for the device ID.

## 25.44. Operation 51: LAYOUTRETURN - Release Layout Information

## 25.44.1. ARGUMENT

```
/* Constants used for LAYOUTRETURN and CB_LAYOUTRECALL */
const LAYOUT4_RET_REC_FILE      = 1;
const LAYOUT4_RET_REC_FSID     = 2;
const LAYOUT4_RET_REC_ALL      = 3;

enum layoutreturn_type4 {
    LAYOUTRETURN4_FILE = LAYOUT4_RET_REC_FILE,
    LAYOUTRETURN4_FSID = LAYOUT4_RET_REC_FSID,
    LAYOUTRETURN4_ALL  = LAYOUT4_RET_REC_ALL
};

struct layoutreturn_file4 {
    offset4      lrf_offset;
    length4      lrf_length;
    stateid4     lrf_stateid;
    /* layouttype4 specific data */
    opaque       lrf_body<>;
};

union layoutreturn4 switch(layoutreturn_type4 lr_returntype) {
    case LAYOUTRETURN4_FILE:
        layoutreturn_file4      lr_layout;
    default:
        void;
};

struct LAYOUTRETURN4args {
    /* CURRENT_FH: file */
    bool                lora_reclaim;
    layouttype4         lora_layout_type;
    layoutiomode4       lora_iomode;
    layoutreturn4       lora_layoutreturn;
};
```

## 25.44.2. RESULT

```

union layoutreturn_stateid switch (bool lrs_present) {
case TRUE:
    stateid4          lrs_stateid;
case FALSE:
    void;
};

union LAYOUTRETURN4res switch (nfsstat4 lorr_status) {
case NFS4_OK:
    layoutreturn_stateid    lorr_stateid;
default:
    void;
};

```

### 25.44.3. DESCRIPTION

This operation returns from the client to the server one or more layouts represented by the client ID (derived from the session ID in the preceding SEQUENCE operation), `lora_layout_type`, and `lora_iomode`. When `lr_returntype` is `LAYOUTRETURN4_FILE`, the returned layout is further identified by the current filehandle, `lrf_offset`, `lrf_length`, and `lrf_stateid`. If the `lrf_length` field is `NFS4_UINT64_MAX`, all bytes of the layout, starting at `lrf_offset`, are returned. When `lr_returntype` is `LAYOUTRETURN4_FSID`, the current filehandle is used to identify the file system and all layouts matching the client ID, the `fsid` of the file system, `lora_layout_type`, and `lora_iomode` are returned. When `lr_returntype` is `LAYOUTRETURN4_ALL`, all layouts matching the client ID, `lora_layout_type`, and `lora_iomode` are returned and the current filehandle is not used. After this call, the client MUST NOT use the returned layout(s) and the associated storage protocol to access the file data.

If the set of layouts designated in the case of `LAYOUTRETURN4_FSID` or `LAYOUTRETURN4_ALL` is empty, then no error results. In the case of `LAYOUTRETURN4_FILE`, the byte-range specified is returned even if it is a subdivision of a layout previously obtained with `LAYOUTGET`, a combination of multiple layouts previously obtained with `LAYOUTGET`, or a combination including some layouts previously obtained with `LAYOUTGET`, and one or more subdivisions of such layouts. When the byte-range does not designate any bytes for which a layout is held for the specified file, client ID, layout type and mode, no error results. See Section 18.7.5.3.5 for considerations with "bulk" return of layouts.

The layout being returned may be a subset or superset of a layout specified by `CB_LAYOUTRECALL`. However, if it is a subset, the recall is not complete until the full recalled scope has been returned. Recalled scope refers to the byte-range in the case of

LAYOUTRETURN4\_FILE, the use of LAYOUTRETURN4\_FSID, or the use of LAYOUTRETURN4\_ALL. There must be a LAYOUTRETURN with a matching scope to complete the return even if all current layout ranges have been previously individually returned.

For all lr\_returntype values, an iomode of LAYOUTIOMODE4\_ANY specifies that all layouts that match the other arguments to LAYOUTRETURN (i.e., client ID, lora\_layout\_type, and one of current filehandle and range; fsid derived from current filehandle; or LAYOUTRETURN4\_ALL) are being returned.

In the case that lr\_returntype is LAYOUTRETURN4\_FILE, the lrf\_stateid provided by the client is a layout stateid as returned from previous layout operations. Note that the "seqid" field of lrf\_stateid MUST NOT be zero. See Sections 13.2, 18.7.3, and 18.7.5.2 for a further discussion and requirements.

Return of a layout or all layouts does not invalidate the mapping of storage device ID to a storage device address. The mapping remains in effect until specifically changed or deleted via device ID notification callbacks. Of course if there are no remaining layouts that refer to a previously used device ID, the server is free to delete a device ID without a notification callback, which will be the case when notifications are not in effect.

If the lora\_reclaim field is set to TRUE, the client is attempting to return a layout that was acquired before the restart of the metadata server during the metadata server's grace period. When returning layouts that were acquired during the metadata server's grace period, the client MUST set the lora\_reclaim field to FALSE. The lora\_reclaim field MUST be set to FALSE also when lr\_layoutreturn is LAYOUTRETURN4\_FSID or LAYOUTRETURN4\_ALL. See LAYOUTCOMMIT (Section 25.42) for more details.

Layouts may be returned when recalled or voluntarily (i.e., before the server has recalled them). In either case, the client must properly propagate state changed under the context of the layout to the storage device(s) or to the metadata server before returning the layout.

If the client returns the layout in response to a CB\_LAYOUTRECALL where the lor\_recalltype field of the clora\_recall field was LAYOUTRECALL4\_FILE, the client should use the lor\_stateid value from CB\_LAYOUTRECALL as the value for lrf\_stateid. Otherwise, it should use logr\_stateid (from a previous LAYOUTGET result) or lorr\_stateid (from a previous LAYRETURN result). This is done to indicate the point in time (in terms of layout stateid transitions) when the recall was sent. The client uses the precise lora\_recallstateid

value and MUST NOT set the stateid's seqid to zero; otherwise, NFS4ERR\_BAD\_STATEID MUST be returned. NFS4ERR\_OLD\_STATEID can be returned if the client is using an old seqid, and the server knows the client should not be using the old seqid. For example, the client uses the seqid on slot 1 of the session, receives the response with the new seqid, and uses the slot to send another request with the old seqid.

If a client fails to return a layout in a timely manner, then the metadata server should use its control protocol with the storage devices to fence the client from accessing the data referenced by the layout. See Section 18.7.5 for more details.

If the LAYOUTRETURN request sets the lora\_reclaim field to TRUE after the metadata server's grace period, NFS4ERR\_NO\_GRACE is returned.

If the LAYOUTRETURN request sets the lora\_reclaim field to TRUE and lr\_returntype is set to LAYOUTRETURN4\_FSID or LAYOUTRETURN4\_ALL, NFS4ERR\_INVALID is returned.

If the client sets the lr\_returntype field to LAYOUTRETURN4\_FILE, then the lrs\_stateid field will represent the layout stateid as updated for this operation's processing; the current stateid will also be updated to match the returned value. If the last byte of any layout for the current file, client ID, and layout type is being returned and there are no remaining pending CB\_LAYOUTRECALL operations for which a LAYOUTRETURN operation must be done, lrs\_present MUST be FALSE, and no stateid will be returned. In addition, the COMPOUND request's current stateid will be set to the all-zeroes special stateid (See Section 23.2.3.1.2). The server MUST reject with NFS4ERR\_BAD\_STATEID any further use of the current stateid in that COMPOUND until the current stateid is re-established by a later stateid-returning operation.

On success, the current filehandle retains its value.

If the EXCHGID4\_FLAG\_BIND\_PRINC\_STATEID capability is set on the client ID (See Section 25.35), the server will require that the principal, security flavor, and if applicable, the GSS mechanism, combination that acquired the layout also be the one to send LAYOUTRETURN. This might not be possible if credentials for the principal are no longer available. The server will allow the machine credential or SSV credential (See Section 25.35) to send LAYOUTRETURN if LAYOUTRETURN's operation code was set in the spo\_must\_allow result of EXCHANGE\_ID.

## 25.44.4. IMPLEMENTATION

The final LAYOUTRETURN operation in response to a CB\_LAYOUTRECALL callback MUST be serialized with any outstanding, intersecting LAYOUTRETURN operations. Note that it is possible that while a client is returning the layout for some recalled range, the server may recall a superset of that range (e.g., LAYOUTRECALL4\_ALL); the final return operation for the latter must block until the former layout recall is done.

Returning all layouts in a file system using LAYOUTRETURN4\_FSID is typically done in response to a CB\_LAYOUTRECALL for that file system as the final return operation. Similarly, LAYOUTRETURN4\_ALL is used in response to a recall callback for all layouts. It is possible that the client already returned some outstanding layouts via individual LAYOUTRETURN calls and the call for LAYOUTRETURN4\_FSID or LAYOUTRETURN4\_ALL marks the end of the LAYOUTRETURN sequence. See Section 18.7.3.1 for more details.

Once the client has returned all layouts referring to a particular device ID, the server MAY delete the device ID.

## 25.45. Operation 52: SECINFO\_NO\_NAME - Get Security on Unnamed Object

Although this is a new NFSv4.1 operation and appropriately described in this document, much of the detail regarding the values returned and their role in security negotiation is described in Section 16 of the NFSv4-wide security document, currently [I-D.dnoveck-nfsv4-security]. This adaptation has been necessary since connection characteristics are now an appropriate subject of negotiation, where previously negotiation only concerned the choice of appropriate auth flavors on existing connection.

## 25.45.1. ARGUMENT

```
enum secinfo_style4 {
    SECINFO_STYLE4_CURRENT_FH      = 0,
    SECINFO_STYLE4_PARENT          = 1
};
```

```
/* CURRENT_FH: object or child directory */
typedef secinfo_style4 SECINFO_NO_NAME4args;
```

## 25.45.2. RESULT

```
/* CURRENTFH: consumed if status is NFS4_OK */
typedef SECINFO4res SECINFO_NO_NAME4res;
```



### 25.45.3. DESCRIPTION

Like the SECINFO operation, SECINFO\_NO\_NAME is used by the client to obtain a list of valid RPC authentication flavors and transport characteristics for a specific file object. Unlike SECINFO, SECINFO\_NO\_NAME only works with objects that are accessed by filehandle.

There are two styles of SECINFO\_NO\_NAME, as determined by the value of the secinfo\_style4 enumeration. If SECINFO\_STYLE4\_CURRENT\_FH is passed, then SECINFO\_NO\_NAME is querying for the required security for the current filehandle. If SECINFO\_STYLE4\_PARENT is passed, then SECINFO\_NO\_NAME is querying for the required security of the current filehandle's parent, where the current filehandle MUST be that of directory (an object of type NF4DIR). If the style selected is SECINFO\_STYLE4\_PARENT, then SECINFO should apply the same access methodology used for LOOKUPP when evaluating the traversal to the parent directory. Therefore, if the requester does not have the appropriate access to LOOKUPP the parent, then SECINFO\_NO\_NAME must behave the same way and return NFS4ERR\_ACCESS.

If PUTFH, PUTPUBFH, PUTROOTFH, or RESTOREFH returns NFS4ERR\_WRONGSEC, then the client resolves the situation by sending a COMPOUND request that consists of PUTFH, PUTPUBFH, or PUTROOTFH immediately followed by SECINFO\_NO\_NAME, style SECINFO\_STYLE4\_CURRENT\_FH. See Section 6.2 for instructions on dealing with NFS4ERR\_WRONGSEC error returns from PUTFH, PUTROOTFH, PUTPUBFH, or RESTOREFH.

If SECINFO\_STYLE4\_PARENT is specified and there is no parent directory, SECINFO\_NO\_NAME MUST return NFS4ERR\_NOENT.

On success, the current filehandle is consumed (See Section 6.2.1.8), and if the next operation after SECINFO\_NO\_NAME tries to use the current filehandle, that operation will fail with the status NFS4ERR\_NOFILEHANDLE.

Everything else about SECINFO\_NO\_NAME is the same as SECINFO. See the discussion of SECINFO in Section 12.5 of the NFSv4-wide security document.

### 25.45.4. IMPLEMENTATION

See the discussion on SECINFO in Section 12.5.4.2 of the NFSv4-wide security document, currently [I-D.dnoveck-nfsv4-security].

### 25.46. Operation 53: SEQUENCE - Supply Per-Procedure Sequencing and Control

## 25.46.1. ARGUMENT

```

struct SEQUENCE4args {
    sessionid4    sa_sessionid;
    sequenceid4   sa_sequenceid;
    slotid4       sa_slotid;
    slotid4       sa_highest_slotid;
    bool          sa_cachethis;
};

```

## 25.46.2. RESULT

```

const SEQ4_STATUS_CB_PATH_DOWN           = 0x00000001;
const SEQ4_STATUS_CB_GSS_CONTEXTS_EXPIRING = 0x00000002;
const SEQ4_STATUS_CB_GSS_CONTEXTS_EXPIRED  = 0x00000004;
const SEQ4_STATUS_EXPIRED_ALL_STATE_REVOKED = 0x00000008;
const SEQ4_STATUS_EXPIRED_SOME_STATE_REVOKED = 0x00000010;
const SEQ4_STATUS_ADMIN_STATE_REVOKED      = 0x00000020;
const SEQ4_STATUS_RECALLABLE_STATE_REVOKED = 0x00000040;
const SEQ4_STATUS_LEASE_MOVED              = 0x00000080;
const SEQ4_STATUS_RESTART_RECLAIM_NEEDED  = 0x00000100;
const SEQ4_STATUS_CB_PATH_DOWN_SESSION    = 0x00000200;
const SEQ4_STATUS_BACKCHANNEL_FAULT        = 0x00000400;
const SEQ4_STATUS_DEVID_CHANGED            = 0x00000800;
const SEQ4_STATUS_DEVID_DELETED           = 0x00001000;

```

```

struct SEQUENCE4resok {
    sessionid4    sr_sessionid;
    sequenceid4   sr_sequenceid;
    slotid4       sr_slotid;
    slotid4       sr_highest_slotid;
    slotid4       sr_target_highest_slotid;
    uint32_t      sr_status_flags;
};

```

```

union SEQUENCE4res switch (nfsstat4 sr_status) {
case NFS4_OK:
    SEQUENCE4resok  sr_resok4;
default:
    void;
};

```

## 25.46.3. DESCRIPTION

The SEQUENCE operation is used by the server to implement session request control and the reply cache semantics.

SEQUENCE MUST appear as the first operation of any COMPOUND in which it appears. The error NFS4ERR\_SEQUENCE\_POS will be returned when it is found in any position in a COMPOUND beyond the first. Operations other than SEQUENCE, BIND\_CONN\_TO\_SESSION, EXCHANGE\_ID, DESTROY\_CLIENTID, CREATE\_SESSION, and DESTROY\_SESSION, MUST NOT appear as the first operation in a COMPOUND. Such operations MUST yield the error NFS4ERR\_OP\_NOT\_IN\_SESSION if they do appear at the start of a COMPOUND.

If SEQUENCE is received on a connection not associated with the session via CREATE\_SESSION or BIND\_CONN\_TO\_SESSION, and connection association enforcement is enabled (See Section 25.35), then the server returns NFS4ERR\_CONN\_NOT\_BOUND\_TO\_SESSION.

The sa\_sessionid argument identifies the session to which this request applies. The sr\_sessionid result MUST equal sa\_sessionid.

The sa\_slotid argument is the index in the reply cache for the request. The sa\_sequenceid field is the sequence number of the request for the reply cache entry (slot). The sr\_slotid result MUST equal sa\_slotid. The sr\_sequenceid result MUST equal sa\_sequenceid.

The sa\_highest\_slotid argument is the highest slot ID for which the client has a request outstanding; it could be equal to sa\_slotid. The server returns two "highest\_slotid" values: sr\_highest\_slotid and sr\_target\_highest\_slotid. The former is the highest slot ID the server will accept in future SEQUENCE operation, and should not be less than the value of sa\_highest\_slotid (but see Section 7.6.1 for an exception). The latter is the highest slot ID the server would prefer the client use on a future SEQUENCE operation.

If sa\_cachethis is TRUE, then the client is requesting that the server cache the entire reply in the server's reply cache; therefore, the server MUST cache the reply (See Section 7.6.1.3). The server MAY cache the reply if sa\_cachethis is FALSE. If the server does not cache the entire reply, it MUST still record that it executed the request at the specified slot and sequence ID.

The response to the SEQUENCE operation contains a word of status flags (sr\_status\_flags) that can provide to the client information related to the status of the client's lock state and communications paths. Note that any status bits relating to lock state MAY be reset when lock state is lost due to a server restart (even if the session is persistent across restarts; session persistence does not imply lock state persistence) or the establishment of a new client instance.

**SEQ4\_STATUS\_CB\_PATH\_DOWN**

When set, indicates that the client has no operational backchannel path for any session associated with the client ID, making it necessary for the client to re-establish one. This bit remains set on all SEQUENCE responses on all sessions associated with the client ID until at least one backchannel is available on any session associated with the client ID. If the client fails to re-establish a backchannel for the client ID, it is subject to having recallable state revoked.

**SEQ4\_STATUS\_CB\_PATH\_DOWN\_SESSION**

When set, indicates that the session has no operational backchannel. There are two reasons why SEQ4\_STATUS\_CB\_PATH\_DOWN\_SESSION may be set and not SEQ4\_STATUS\_CB\_PATH\_DOWN. First is that a callback operation that applies specifically to the session (e.g., CB\_RECALL\_SLOT, see Section 27.8) needs to be sent. Second is that the server did send a callback operation, but the connection was lost before the reply. The server cannot be sure whether or not the client received the callback operation, and so, per rules on request retry, the server MUST retry the callback operation over the same session. The SEQ4\_STATUS\_CB\_PATH\_DOWN\_SESSION bit is the indication to the client that it needs to associate a connection to the session's backchannel. This bit remains set on all SEQUENCE responses of the session until a connection is associated with the session's a backchannel. If the client fails to re-establish a backchannel for the session, it is subject to having recallable state revoked.

**SEQ4\_STATUS\_CB\_GSS\_CONTEXTS\_EXPIRING**

When set, indicates that all GSS contexts or RPCSEC\_GSS handles assigned to the session's backchannel will expire within a period equal to the lease time. This bit remains set on all SEQUENCE replies until at least one of the following are true:

- \* All SSV RPCSEC\_GSS handles on the session's backchannel have been destroyed and all non-SSV GSS contexts have expired.
- \* At least one more SSV RPCSEC\_GSS handle has been added to the backchannel.
- \* The expiration time of at least one non-SSV GSS context of an RPCSEC\_GSS handle is beyond the lease period from the current time (relative to the time of when a SEQUENCE response was sent)

**SEQ4\_STATUS\_CB\_GSS\_CONTEXTS\_EXPIRED**

When set, indicates all non-SSV GSS contexts and all SSV RPCSEC\_GSS handles assigned to the session's backchannel have expired or have been destroyed. This bit remains set on all SEQUENCE replies until at least one non-expired non-SSV GSS context for the session's backchannel has been established or at least one SSV RPCSEC\_GSS handle has been assigned to the backchannel.

**SEQ4\_STATUS\_EXPIRED\_ALL\_STATE\_REVOKED**

When set, indicates that the lease has expired and as a result the server released all of the client's locking state. This status bit remains set on all SEQUENCE replies until the loss of all such locks has been acknowledged by use of FREE\_STATEID (See Section 25.38), or by establishing a new client instance by destroying all sessions (via DESTROY\_SESSION), the client ID (via DESTROY\_CLIENTID), and then invoking EXCHANGE\_ID and CREATE\_SESSION to establish a new client ID.

**SEQ4\_STATUS\_EXPIRED\_SOME\_STATE\_REVOKED**

When set, indicates that some subset of the client's locks have been revoked due to expiration of the lease period followed by another client's conflicting LOCK operation. This status bit remains set on all SEQUENCE replies until the loss of all such locks has been acknowledged by use of FREE\_STATEID.

**SEQ4\_STATUS\_ADMIN\_STATE\_REVOKED**

When set, indicates that one or more locks have been revoked without expiration of the lease period, due to administrative action. This status bit remains set on all SEQUENCE replies until the loss of all such locks has been acknowledged by use of FREE\_STATEID.

**SEQ4\_STATUS\_RECALLABLE\_STATE\_REVOKED**

When set, indicates that one or more recallable objects have been revoked without expiration of the lease period, due to the client's failure to return them when recalled, which may be a consequence of there being no working backchannel and the client failing to re-establish a backchannel per the SEQ4\_STATUS\_CB\_PATH\_DOWN, SEQ4\_STATUS\_CB\_PATH\_DOWN\_SESSION, or SEQ4\_STATUS\_CB\_GSS\_CONTEXTS\_EXPIRED status flags. This status bit remains set on all SEQUENCE replies until the loss of all such locks has been acknowledged by use of FREE\_STATEID.

**SEQ4\_STATUS\_LEASE\_MOVED**

When set, indicates that responsibility for lease renewal has been transferred to one or more new servers. This condition will continue until the client receives an NFS4ERR\_MOVED error and the

server receives the subsequent GETATTR for the `fs_locations` or `fs_locations_info` attribute for an access to each file system for which a lease has been moved to a new server. See Section 17.11.9.2.

#### SEQ4\_STATUS\_RESTART\_RECLAIM\_NEEDED

When set, indicates that due to server restart, the client must reclaim locking state. Until the client sends a global RECLAIM\_COMPLETE (Section 25.51), every SEQUENCE operation will return SEQ4\_STATUS\_RESTART\_RECLAIM\_NEEDED.

#### SEQ4\_STATUS\_BACKCHANNEL\_FAULT

The server has encountered an unrecoverable fault with the backchannel (e.g., it has lost track of the sequence ID for a slot in the backchannel). The client MUST stop sending more requests on the session's fore channel, wait for all outstanding requests to complete on the fore and back channel, and then destroy the session.

#### SEQ4\_STATUS\_DEVID\_CHANGED

The client is using device ID notifications and the server has changed a device ID mapping held by the client. This flag will stay present until the client has obtained the new mapping with GETDEVICEINFO.

#### SEQ4\_STATUS\_DEVID\_DELETED

The client is using device ID notifications and the server has deleted a device ID mapping held by the client. This flag will stay in effect until the client sends a GETDEVICEINFO on the device ID with a null value in the argument `gdi_notify_types`.

The value of the `sa_sequenceid` argument relative to the cached sequence ID on the slot falls into one of three cases.

- \* If the difference between `sa_sequenceid` and the server's cached sequence ID at the slot ID is two (2) or more, or if `sa_sequenceid` is less than the cached sequence ID (accounting for wraparound of the unsigned sequence ID value), then the server MUST return NFS4ERR\_SEQ\_MISORDERED.
- \* If `sa_sequenceid` and the cached sequence ID are the same, this is a retry, and the server replies with what is recorded in the reply cache. The lease is possibly renewed as described below.
- \* If `sa_sequenceid` is one greater (accounting for wraparound) than the cached sequence ID, then this is a new request, and the slot's sequence ID is incremented. The operations subsequent to SEQUENCE, if any, are processed. If there are no other

operations, the only other effects are to cache the SEQUENCE reply in the slot, maintain the session's activity, and possibly renew the lease.

If the client reuses a slot ID and sequence ID for a completely different request, the server MAY treat the request as if it is a retry of what it has already executed. The server MAY however detect the client's illegal reuse and return NFS4ERR\_SEQ\_FALSE\_RETRY.

If SEQUENCE returns an error, then the state of the slot (sequence ID, cached reply) MUST NOT change, and the associated lease MUST NOT be renewed.

If SEQUENCE returns NFS4\_OK, then the associated lease MUST be renewed (See Section 13.3), except if SEQ4\_STATUS\_EXPIRED\_ALL\_STATE\_REVOKED is returned in sr\_status\_flags.

#### 25.46.4. IMPLEMENTATION

The server MUST maintain a mapping of session ID to client ID in order to validate any operations that follow SEQUENCE that take a stateid as an argument and/or result.

If the client establishes a persistent session, then a SEQUENCE received after a server restart might encounter requests performed and recorded in a persistent reply cache before the server restart. In this case, SEQUENCE will be processed successfully, while requests that were not previously performed and recorded are rejected with NFS4ERR\_DEADSESSION.

Depending on which of the operations within the COMPOUND were successfully performed before the server restart, these operations will also have replies sent from the server reply cache. Note that when these operations establish locking state, it is locking state that applies to the previous server instance and to the previous client ID, even though the server restart, which logically happened after these operations, eliminated that state. In the case of a partially executed COMPOUND, processing may reach an operation not processed during the earlier server instance, making this operation a new one and not performable on the existing session. In this case, NFS4ERR\_DEADSESSION will be returned from that operation.

#### 25.47. Operation 54: SET\_SSV - Update SSV for a Client ID

##### 25.47.1. ARGUMENT

```
struct ssa_digest_input4 {
    SEQUENCE4args sdi_seqargs;
};

struct SET_SSV4args {
    opaque        ssa_ssv<>;
    opaque        ssa_digest<>;
};
```

#### 25.47.2. RESULT

```
struct ssr_digest_input4 {
    SEQUENCE4res sdi_seqres;
};

struct SET_SSV4resok {
    opaque        ssr_digest<>;
};

union SET_SSV4res switch (nfsstat4 ssr_status) {
case NFS4_OK:
    SET_SSV4resok    ssr_resok4;
default:
    void;
};
```

#### 25.47.3. DESCRIPTION

This operation is used to update the SSV for a client ID. Before SET\_SSV is called the first time on a client ID, the SSV is zero. The SSV is the key used for the SSV GSS mechanism (Section 7.9)

SET\_SSV MUST be preceded by a SEQUENCE operation in the same COMPOUND. It MUST NOT be used if the client did not opt for SP4\_SSV state protection when the client ID was created (See Section 25.35); the server returns NFS4ERR\_INVALID in that case.

The field ssa\_digest is computed as the output of the HMAC [RFC2104] using the subkey derived from the SSV4\_SUBKEY\_MIC\_I2T and current SSV as the key (See Section 7.9 for a description of subkeys), and an XDR encoded value of data type ssa\_digest\_input4. The field sdi\_seqargs is equal to the arguments of the SEQUENCE operation for the COMPOUND procedure that SET\_SSV is within.

The argument ssa\_ssv is XORed with the current SSV to produce the new SSV. The argument ssa\_ssv should be generated randomly.



In the response, `ssr_digest` is the output of the HMAC using the subkey derived from `SSV4_SUBKEY_MIC_T2I` and new SSV as the key, and an XDR encoded value of data type `ssr_digest_input4`. The field `sdi_seqres` is equal to the results of the SEQUENCE operation for the COMPOUND procedure that `SET_SSV` is within.

As noted in Section 25.35, the client and server can maintain multiple concurrent versions of the SSV. The client and server each MUST maintain an internal SSV version number, which is set to one the first time `SET_SSV` executes on the server and the client receives the first `SET_SSV` reply. Each subsequent `SET_SSV` increases the internal SSV version number by one. The value of this version number corresponds to the `smpt_ssv_seq`, `smt_ssv_seq`, `sspt_ssv_seq`, and `ssct_ssv_seq` fields of the SSV GSS mechanism tokens (See Section 7.9).

#### 25.47.4. IMPLEMENTATION

When the server receives `ssa_digest`, it MUST verify the digest by computing the digest the same way the client did and comparing it with `ssa_digest`. If the server gets a different result, this is an error, `NFS4ERR_BAD_SESSION_DIGEST`. This error might be the result of another `SET_SSV` from the same client ID changing the SSV. If so, the client recovers by sending a `SET_SSV` operation again with a recomputed digest based on the subkey of the new SSV. If the transport connection is dropped after the `SET_SSV` request is sent, but before the `SET_SSV` reply is received, then there are special considerations for recovery if the client has no more connections associated with sessions associated with the client ID of the SSV. See Section 25.34.4.

Clients should not send an `ssa_ssv` that is equal to a previous `ssa_ssv`, nor equal to a previous or current SSV (including an `ssa_ssv` equal to zero since the SSV is initialized to zero when the client ID is created).

Clients should not send `SET_SSV` with `RPCSEC_GSS` privacy. Servers MUST support `RPCSEC_GSS` with privacy for any COMPOUND that has { SEQUENCE, SET\_SSV }.

A client should not send `SET_SSV` with the SSV GSS mechanism's credential because the purpose of `SET_SSV` is to seed the SSV from non-SSV credentials. Instead, `SET_SSV` should be sent with the credential of a user that is accessing the client ID for the first time (Section 7.8.3). However, if the client does send `SET_SSV` with SSV credentials, the digest protecting the arguments uses the value of the SSV before `ssa_ssv` is XORed in, and the digest protecting the results uses the value of the SSV after the `ssa_ssv` is XORed in.

## 25.48. Operation 55: TEST\_STATEID - Test Stateids for Validity

## 25.48.1. ARGUMENT

```
struct TEST_STATEID4args {  
    stateid4      ts_stateids<>;  
};
```

## 25.48.2. RESULT

```
struct TEST_STATEID4resok {  
    nfsstat4      tsr_status_codes<>;  
};  
  
union TEST_STATEID4res switch (nfsstat4 tsr_status) {  
    case NFS4_OK:  
        TEST_STATEID4resok tsr_resok4;  
    default:  
        void;  
};
```

## 25.48.3. DESCRIPTION

The TEST\_STATEID operation is used to check the validity of a set of stateids. It can be used at any time, but the client should definitely use it when it receives an indication that one or more of its stateids have been invalidated due to lock revocation. This occurs when the SEQUENCE operation returns with one of the following sr\_status\_flags set:

- \* SEQ4\_STATUS\_EXPIRED\_SOME\_STATE\_REVOKED
- \* SEQ4\_STATUS\_EXPIRED\_ADMIN\_STATE\_REVOKED
- \* SEQ4\_STATUS\_EXPIRED\_RECALLABLE\_STATE\_REVOKED

The client can use TEST\_STATEID one or more times to test the validity of its stateids. Each use of TEST\_STATEID allows a large set of such stateids to be tested and avoids problems with earlier stateids in a COMPOUND request from interfering with the checking of subsequent stateids, as would happen if individual stateids were tested by a series of corresponding by operations in a COMPOUND request.

For each stateid, the server returns the status code that would be returned if that stateid were to be used in normal operation. Returning such a status indication is not an error and does not cause COMPOUND processing to terminate. Checks for the validity of the stateid proceed as they would for normal operations with a number of exceptions:

- \* There is no check for the type of stateid object, as would be the case for normal use of a stateid.
- \* There is no reference to the current filehandle.
- \* Special stateids are always considered invalid (they result in the error code NFS4ERR\_BAD\_STATEID).

All stateids are interpreted as being associated with the client for the current session. Any possible association with a previous instance of the client (as stale stateids) is not considered.

The valid status values in the returned status\_code array are NFS4ERR\_OK, NFS4ERR\_BAD\_STATEID, NFS4ERR\_OLD\_STATEID, NFS4ERR\_EXPIRED, NFS4ERR\_ADMIN\_REVOKED, and NFS4ERR\_DELEG\_REVOKED.

#### 25.48.4. IMPLEMENTATION

See Sections 13.2.2 and 13.2.4 for a discussion of stateid structure, lifetime, and validation.

#### 25.49. Operation 56: WANT\_DELEGATION - Request Delegation

##### 25.49.1. ARGUMENT

```

union deleg_claim4 switch (open_claim_type4 dc_claim) {
/*
 * No special rights to object. Ordinary delegation
 * request of the specified object. Object identified
 * by filehandle.
 */
case CLAIM_FH: /* new to v4.1 */
    /* CURRENT_FH: object being delegated */
    void;

/*
 * Right to file based on a delegation granted
 * to a previous boot instance of the client.
 * File is specified by filehandle.
 */
case CLAIM_DELEG_PREV_FH: /* new to v4.1 */
    /* CURRENT_FH: object being delegated */
    void;

/*
 * Right to the file established by an open previous
 * to server reboot. File identified by filehandle.
 * Used during server reclaim grace period.
 */
case CLAIM_PREVIOUS:
    /* CURRENT_FH: object being reclaimed */
    open_delegation_type4    dc_delegate_type;
};

struct WANT_DELEGATION4args {
    uint32_t        wda_want;
    deleg_claim4    wda_claim;
};

```

#### 25.49.2. RESULT

```

union WANT_DELEGATION4res switch (nfsstat4 wdr_status) {
case NFS4_OK:
    open_delegation4 wdr_resok4;
default:
    void;
};

```

#### 25.49.3. DESCRIPTION

Where this description mandates the return of a specific error code for a specific condition, and where multiple conditions apply, the server MAY return any of the mandated error codes.

This operation allows a client to:

- \* Get a delegation on all types of files except directories.
- \* Register a "want" for a delegation for the specified file object, and be notified via a callback when the delegation is available. The server MAY support notifications of availability via callbacks. If the server does not support registration of wants, it MUST NOT return an error to indicate that, and instead MUST return with `ond_why` set to `WND4_CONTENTION` or `WND4_RESOURCE` and `ond_server_will_push_deleg` or `ond_server_will_signal_avail` set to `FALSE`. When the server indicates that it will notify the client by means of a callback, it will either provide the delegation using a `CB_PUSH_DELEG` operation or cancel its promise by sending a `CB_WANTS_CANCELLED` operation.
- \* Cancel a want for a delegation.

The client should not set `OPEN4_SHARE_ACCESS_READ` and should not set `OPEN4_SHARE_ACCESS_WRITE` in `wda_want`. If it does, the server MUST ignore them.

The meanings of the following flags in `wda_want` are the same as they are in `OPEN`, except as noted below.

- \* `OPEN4_SHARE_ACCESS_WANT_READ_DELEG`
- \* `OPEN4_SHARE_ACCESS_WANT_WRITE_DELEG`
- \* `OPEN4_SHARE_ACCESS_WANT_ANY_DELEG`
- \* `OPEN4_SHARE_ACCESS_WANT_NO_DELEG`. Unlike the `OPEN` operation, this flag should not be set by the client in the arguments to `WANT_DELEGATION`, and MUST be ignored by the server.
- \* `OPEN4_SHARE_ACCESS_WANT_CANCEL`
- \* `OPEN4_SHARE_ACCESS_WANT_SIGNAL_DELEG_WHEN_RESRC_AVAIL`
- \* `OPEN4_SHARE_ACCESS_WANT_PUSH_DELEG_WHEN_UNCONTENDED`

The handling of the above flags in `WANT_DELEGATION` is the same as in `OPEN`. Information about the delegation and/or the promises the server is making regarding future callbacks are the same as those described in the `open_delegation4` structure.

The successful results of WANT\_DELEGATION are of data type `open_delegation4`, which is the same data type as the "delegation" field in the results of the OPEN operation (See Section 25.16.3). The server constructs `wdr_resok4` the same way it constructs OPEN's "delegation" with one difference: WANT\_DELEGATION MUST NOT return a delegation type of OPEN\_DELEGATE\_NONE.

If  $((\text{wda\_want} \ \& \ \text{OPEN4\_SHARE\_ACCESS\_WANT\_DELEG\_MASK}) \ \& \sim \text{OPEN4\_SHARE\_ACCESS\_WANT\_NO\_DELEG})$  is zero, then the client is indicating no explicit desire or non-desire for a delegation and the server MUST return NFS4ERR\_INVAL.

The client uses the OPEN4\_SHARE\_ACCESS\_WANT\_CANCEL flag in the WANT\_DELEGATION operation to cancel a previously requested want for a delegation. Note that if the server is in the process of sending the delegation (via CB\_PUSH\_DELEG) at the time the client sends a cancellation of the want, the delegation might still be pushed to the client.

If WANT\_DELEGATION fails to return a delegation, and the server returns NFS4\_OK, the server MUST set the delegation type to OPEN4\_DELEGATE\_NONE\_EXT, and set `od_whynone`, as described in Section 25.16. Write delegations are not available for file types that are not writable. This includes file objects of types NF4BLK, NF4CHR, NF4LNK, NF4SOCK, and NF4FIFO. If the client requests OPEN4\_SHARE\_ACCESS\_WANT\_WRITE\_DELEG without OPEN4\_SHARE\_ACCESS\_WANT\_READ\_DELEG on an object with one of the aforementioned file types, the server must set `wdr_resok4.od_whynone.ond_why` to WND4\_WRITE\_DELEG\_NOT\_SUPP\_FTYPE.

#### 25.49.4. IMPLEMENTATION

A request for a conflicting delegation is not normally intended to trigger the recall of the existing delegation. Servers may choose to treat some clients as having higher priority such that their wants will trigger recall of an existing delegation, although that is expected to be an unusual situation.

Servers will generally recall delegations assigned by WANT\_DELEGATION on the same basis as those assigned by OPEN. CB\_RECALL will generally be done only when other clients perform operations inconsistent with the delegation. The normal response to aging of delegations is to use CB\_RECALL\_ANY, in order to give the client the opportunity to keep the delegations most useful from its point of view.

#### 25.50. Operation 57: DESTROY\_CLIENTID - Destroy a Client ID

## 25.50.1. ARGUMENT

```
struct DESTROY_CLIENTID4args {  
    clientid4      dca_clientid;  
};
```

## 25.50.2. RESULT

```
struct DESTROY_CLIENTID4res {  
    nfsstat4      dcr_status;  
};
```

## 25.50.3. DESCRIPTION

The DESTROY\_CLIENTID operation destroys the client ID. If there are sessions (both idle and non-idle), opens, locks, delegations, and/or wants (Section 25.49) associated with the unexpired lease of the client ID, the server MUST return NFS4ERR\_CLIENTID\_BUSY. DESTROY\_CLIENTID MAY be preceded with a SEQUENCE operation as long as the client ID derived from the session ID of SEQUENCE is not the same as the client ID to be destroyed. If the client IDs are the same, then the server MUST return NFS4ERR\_CLIENTID\_BUSY.

If DESTROY\_CLIENTID is not prefixed by SEQUENCE, it MUST be the only operation in the COMPOUND request (otherwise, the server MUST return NFS4ERR\_NOT\_ONLY\_OP). If the operation is sent without a SEQUENCE preceding it, a client that retransmits the request may receive an error in response, because the original request might have been successfully executed.

## 25.50.4. IMPLEMENTATION

DESTROY\_CLIENTID allows a server to immediately reclaim the resources consumed by an unused client ID, and also to forget that it ever generated the client ID. By forgetting that it ever generated the client ID, the server can safely reuse the client ID on a future EXCHANGE\_ID operation.

## 25.51. Operation 58: RECLAIM\_COMPLETE - Indicates Reclaims Finished

## 25.51.1. ARGUMENT

```
struct RECLAIM_COMPLETE4args {  
    /*  
     * If rca_one_fs TRUE,  
     *  
     * CURRENT_FH: object in  
     * file system reclaim is  
     * complete for.  
     */  
    bool                rca_one_fs;  
};
```

#### 25.51.2. RESULTS

```
struct RECLAIM_COMPLETE4res {  
    nfsstat4            rcr_status;  
};
```

#### 25.51.3. DESCRIPTION

A RECLAIM\_COMPLETE operation is used to indicate that the client has reclaimed all of the locking state that it will recover using reclaim-type operation used to re-establish locking state during a server grace period. It is not used in connection with the special delegation recovery period used after client restart.

It does so when it is recovering state due to either a server restart or the migration of a file system to another server. There are two types of RECLAIM\_COMPLETE operations:

- \* When rca\_one\_fs is FALSE, a global RECLAIM\_COMPLETE is being done. This indicates that recovery of all locks that the client held on the previous server instance has been completed. The current filehandle need not be set in this case.
- \* When rca\_one\_fs is TRUE, a file system-specific RECLAIM\_COMPLETE is being done. This indicates that recovery of locks for a single fs (the one designated by the current filehandle) due to the migration of the file system has been completed. Presence of a current filehandle is required when rca\_one\_fs is set to TRUE. When the current filehandle designates a filehandle in a file system not in the process of migration, the operation returns NFS4\_OK and is otherwise ignored.

Once a RECLAIM\_COMPLETE is done, there can be no further reclaim operations for locks whose scope is defined as having completed recovery. Once the client sends RECLAIM\_COMPLETE, the server will not allow the client to do subsequent reclaims of locking state for that scope and, if these are attempted, will return NFS4ERR\_NO\_GRACE.



Whenever a client establishes a new client ID as a result of one of server restart and before it does the first non-reclaim operation that obtains a lock, it MUST send a RECLAIM\_COMPLETE with `rca_one_fs` set to FALSE, even if there are no locks to reclaim. If non-reclaim locking operations are done before the RECLAIM\_COMPLETE, an NFS4ERR\_GRACE error will be returned.

Similarly, when the client accesses a migrated file system on a new server, before it sends the first non-reclaim operation that obtains a lock on this new server, it MUST send a RECLAIM\_COMPLETE with `rca_one_fs` set to TRUE and current filehandle within that file system, even if there are no locks to reclaim. If non-reclaim locking operations are done on that file system before the RECLAIM\_COMPLETE, an NFS4ERR\_GRACE error will be returned.

It should be noted that there are situations in which a client needs to issue both forms of RECLAIM\_COMPLETE. An example is an instance of file system migration in which the file system is migrated to a server for which the client has no clientid. As a result, the client needs to obtain a clientid from the server (incurring the responsibility to do RECLAIM\_COMPLETE with `rca_one_fs` set to FALSE) as well as RECLAIM\_COMPLETE with `rca_one_fs` set to TRUE to complete the per-fs grace period associated with the file system migration. These two may be done in any order as long as all necessary lock reclaims have been done before issuing either of them.

Any locks not reclaimed at the point at which RECLAIM\_COMPLETE is done become non-reclaimable. The client MUST NOT attempt to reclaim them, either during the current server instance or in any subsequent server instance, or on another server to which responsibility for that file system is transferred. If the client were to do so, it would be violating the protocol by representing itself as owning locks that it does not own, and so has no right to reclaim. See Section 8.4.3 of [RFC5661] for a discussion of edge conditions related to lock reclaim.

By sending a RECLAIM\_COMPLETE, the client indicates readiness to proceed to do normal non-reclaim locking operations. The client should be aware that such operations may temporarily result in NFS4ERR\_GRACE errors until the server is ready to terminate its grace period.

#### 25.51.4. IMPLEMENTATION

Servers will typically use the information as to when reclaim activity is complete to reduce the length of the grace period. When the server maintains in persistent storage a list of clients that might have had locks, it is able to use the fact that all such clients have done a RECLAIM\_COMPLETE to terminate the grace period and begin normal operations (i.e., grant requests for new locks) sooner than it might otherwise.

Latency can be minimized by doing a RECLAIM\_COMPLETE as part of the COMPOUND request in which the last lock-reclaiming operation is done. When there are no reclaims to be done, RECLAIM\_COMPLETE should be done immediately in order to allow the grace period to end as soon as possible.

RECLAIM\_COMPLETE should only be done once for each server instance or occasion of the transition of a file system. If it is done a second time, the error NFS4ERR\_COMPLETE\_ALREADY will result. Note that because of the session feature's retry protection, retries of COMPOUND requests containing RECLAIM\_COMPLETE operation will not result in this error.

When a RECLAIM\_COMPLETE is sent, the client effectively acknowledges any locks not yet reclaimed as lost. This allows the server to re-enable the client to recover locks if the occurrence of edge conditions, as described in Section 13.4.3, had caused the server to disable the client's ability to recover locks.

Because previous descriptions of RECLAIM\_COMPLETE were not sufficiently explicit about the circumstances in which use of RECLAIM\_COMPLETE with rca\_one\_fs set to TRUE was appropriate, there have been cases in which it has been misused by clients who have issued RECLAIM\_COMPLETE with rca\_one\_fs set to TRUE when it should have not been. There have also been cases in which servers have, in various ways, not responded to such misuse as described above, either ignoring the rca\_one\_fs setting (treating the operation as a global RECLAIM\_COMPLETE) or ignoring the entire operation.

While clients should not misuse this feature, and servers are to respond to such misuse as described above, implementers need to be aware of the following considerations as they make necessary trade-offs between interoperability with existing implementations and proper support for facilities to allow lock recovery in the event of file system migration.

- \* When servers have no support for becoming the destination server of a file system subject to migration, there is no possibility of a per-fs RECLAIM\_COMPLETE being done legitimately, and occurrences of it should be ignored. However, the negative consequences of accepting such mistaken use are quite limited as long as the client does not issue it before all necessary reclaims are done.
- \* When a server might become the destination for a file system being migrated, inappropriate use of per-fs RECLAIM\_COMPLETE is more concerning. In the case in which the file system designated is not within a per-fs grace period, the per-fs RECLAIM\_COMPLETE is to be ignored, with the negative consequences of accepting it being limited, as in the case in which migration is not supported. However, if the server encounters a file system undergoing migration, the operation cannot be accepted as if it were a global RECLAIM\_COMPLETE without invalidating its intended use.

## 25.52. Operation 10044: ILLEGAL - Illegal Operation

### 25.52.1. ARGUMENTS

void;

### 25.52.2. RESULTS

```
struct ILLEGAL4res {  
    nfsstat4      status;  
};
```

### 25.52.3. DESCRIPTION

This operation is a placeholder for encoding a result to handle the case of the client sending an operation code within COMPOUND that is not supported. See the COMPOUND procedure description for more details.

The status field of ILLEGAL4res MUST be set to NFS4ERR\_OP\_ILLEGAL.

### 25.52.4. IMPLEMENTATION

A client will probably not send an operation with code OP\_ILLEGAL but if it does, the response will be ILLEGAL4res just as it would be with any other invalid operation code. Note that if the server gets an illegal operation code that is not OP\_ILLEGAL, and if the server checks for legal operation codes during the XDR decode phase, then the ILLEGAL4res would not be returned.

## 26. NFSv4.1 Callback Procedures

The procedures used for callbacks are defined in the following sections. In the interest of clarity, the terms "client" and "server" refer to NFS clients and servers, despite the fact that for an individual callback RPC, the sense of these terms would be precisely the opposite.

Both procedures, CB\_NULL and CB\_COMPOUND, MUST be implemented.

### 26.1. Procedure 0: CB\_NULL - No Operation

#### 26.1.1. ARGUMENTS

void;

#### 26.1.2. RESULTS

void;

#### 26.1.3. DESCRIPTION

CB\_NULL is the standard ONC RPC NULL procedure, with the standard void argument and void response. Even though there is no direct functionality associated with this procedure, the server will use CB\_NULL to confirm the existence of a path for RPCs from the server to client.

#### 26.1.4. ERRORS

None.

### 26.2. Procedure 1: CB\_COMPOUND - Compound Operations

#### 26.2.1. ARGUMENTS

```
enum nfs_cb_opnum4 {
    OP_CB_GETATTR          = 3,
    OP_CB_RECALL           = 4,
    /* Callback operations new to NFSv4.1 */
    OP_CB_LAYOUTRECALL     = 5,
    OP_CB_NOTIFY           = 6,
    OP_CB_PUSH_DELEG       = 7,
    OP_CB_RECALL_ANY       = 8,
    OP_CB_RECALLABLE_OBJ_AVAIL = 9,
    OP_CB_RECALL_SLOT      = 10,
    OP_CB_SEQUENCE         = 11,
    OP_CB_WANTS_CANCELLED  = 12,
```

```

        OP_CB_NOTIFY_LOCK          = 13,
        OP_CB_NOTIFY_DEVICEID      = 14,

        OP_CB_ILLEGAL              = 10044
};

union nfs_cb_argop4 switch (unsigned argop) {
    case OP_CB_GETATTR:
        CB_GETATTR4args            opcbgetattr;
    case OP_CB_RECALL:
        CB_RECALL4args             opcbrecall;
    case OP_CB_LAYOUTRECALL:
        CB_LAYOUTRECALL4args       opcblayoutrecall;
    case OP_CB_NOTIFY:
        CB_NOTIFY4args             opcbnotify;
    case OP_CB_PUSH_DELEG:
        CB_PUSH_DELEG4args         opcbpush_deleg;
    case OP_CB_RECALL_ANY:
        CB_RECALL_ANY4args         opcbrecall_any;
    case OP_CB_RECALLABLE_OBJ_AVAIL:
        CB_RECALLABLE_OBJ_AVAIL4args opcbrecallable_obj_avail;
    case OP_CB_RECALL_SLOT:
        CB_RECALL_SLOT4args        opcbrecall_slot;
    case OP_CB_SEQUENCE:
        CB_SEQUENCE4args           opcbsequence;
    case OP_CB_WANTS_CANCELLED:
        CB_WANTS_CANCELLED4args    opcbwants_cancelled;
    case OP_CB_NOTIFY_LOCK:
        CB_NOTIFY_LOCK4args        opcbnotify_lock;
    case OP_CB_NOTIFY_DEVICEID:
        CB_NOTIFY_DEVICEID4args    opcbnotify_deviceid;
    case OP_CB_ILLEGAL:
        void;
};

struct CB_COMPOUND4args {
    utf8str_cs    tag;
    uint32_t      minorversion;
    uint32_t      callback_ident;
    nfs_cb_argop4 argarray<>;
};

```

#### 26.2.2. RESULTS

```
union nfs_cb_resop4 switch (unsigned resop) {
  case OP_CB_GETATTR:      CB_GETATTR4res  opcbgetattr;
  case OP_CB_RECALL:      CB_RECALL4res   opcbrecall;

  /* new NFSv4.1 operations */
  case OP_CB_LAYOUTRECALL:
                        CB_LAYOUTRECALL4res
                        opcblayoutrecall;

  case OP_CB_NOTIFY:      CB_NOTIFY4res   opcbnotify;

  case OP_CB_PUSH_DELEG:  CB_PUSH_DELEG4res
                        opcbpush_deleg;

  case OP_CB_RECALL_ANY:  CB_RECALL_ANY4res
                        opcbrecall_any;

  case OP_CB_RECALLABLE_OBJ_AVAIL:
                        CB_RECALLABLE_OBJ_AVAIL4res
                        opcbrecallable_obj_avail;

  case OP_CB_RECALL_SLOT:
                        CB_RECALL_SLOT4res
                        opcbrecall_slot;

  case OP_CB_SEQUENCE:    CB_SEQUENCE4res opcbsequence;

  case OP_CB_WANTS_CANCELLED:
                        CB_WANTS_CANCELLED4res
                        opcbwants_cancelled;

  case OP_CB_NOTIFY_LOCK:
                        CB_NOTIFY_LOCK4res
                        opcbnotify_lock;

  case OP_CB_NOTIFY_DEVICEID:
                        CB_NOTIFY_DEVICEID4res
                        opcbnotify_deviceid;

  /* Not new operation */
  case OP_CB_ILLEGAL:      CB_ILLEGAL4res  opcbillegal;
};

struct CB_COMPOUND4res {
  nfsstat4 status;
  utf8str_cs  tag;
  nfs_cb_resop4  resarray<>;
};
```

### 26.2.3. DESCRIPTION

The CB\_COMPOUND procedure is used to combine one or more of the callback procedures into a single RPC request. The main callback RPC program has two main procedures: CB\_NULL and CB\_COMPOUND. All other operations use the CB\_COMPOUND procedure as a wrapper.

During the processing of the CB\_COMPOUND procedure, the client may find that it does not have the available resources to execute any or all of the operations within the CB\_COMPOUND sequence. Refer to Section 7.6.4 for details.

The minorversion field of the arguments MUST be the same as the minorversion of the COMPOUND procedure used to create the client ID and session. For NFSv4.1, minorversion MUST be set to 1.

Contained within the CB\_COMPOUND results is a "status" field. This status MUST be equal to the status of the last operation that was executed within the CB\_COMPOUND procedure. Therefore, if an operation incurred an error, then the "status" value will be the same error value as is being returned for the operation that failed.

The "tag" field is handled the same way as that of the COMPOUND procedure (See Section 23.2.3).

Illegal operation codes are handled in the same way as they are handled for the COMPOUND procedure.

### 26.2.4. IMPLEMENTATION

The CB\_COMPOUND procedure is used to combine individual operations into a single RPC request. The client interprets each of the operations in turn. If an operation is executed by the client and the status of that operation is NFS4\_OK, then the next operation in the CB\_COMPOUND procedure is executed. The client continues this process until there are no more operations to be executed or one of the operations has a status value other than NFS4\_OK.

### 26.2.5. ERRORS

CB\_COMPOUND will of course return every error that each operation on the backchannel can return (See Table 13). However, if CB\_COMPOUND returns zero operations, obviously the error returned by COMPOUND has nothing to do with an error returned by an operation. The list of errors CB\_COMPOUND will return if it processes zero operations includes:

Error	Notes
NFS4ERR_BADCHAR	The tag argument has a character the replier does not support.
NFS4ERR_BADXDR	
NFS4ERR_DELAY	
NFS4ERR_INVAL	The tag argument is not in UTF-8 encoding.
NFS4ERR_MINOR_VERS_MISMATCH	
NFS4ERR_SERVERFAULT	
NFS4ERR_TOO_MANY_OPS	
NFS4ERR_REP_TOO_BIG	
NFS4ERR_REP_TOO_BIG_TO_CACHE	
NFS4ERR_REQ_TOO_BIG	

Table 24: CB\_COMPOUND Error Returns

## 27. NFSv4.1 Callback Operations

## 27.1. Operation 3: CB\_GETATTR - Get Attributes

## 27.1.1. ARGUMENT

```

struct CB_GETATTR4args {
    nfs_fh4 fh;
    bitmap4 attr_request;
};

```

## 27.1.2. RESULT



```

struct CB_GETATTR4resok {
    fattr4  obj_attributes;
};

union CB_GETATTR4res switch (nfsstat4 status) {
    case NFS4_OK:
        CB_GETATTR4resok      resok4;
    default:
        void;
};

```

#### 27.1.3. DESCRIPTION

The CB\_GETATTR operation is used by the server to obtain the current modified state of a file that has been OPEN\_DELEGATE\_WRITE delegated. The size and change attributes are the only ones guaranteed to be serviced by the client. See Section 15.4.3 for a full description of how the client and server are to interact with the use of CB\_GETATTR.

If the filehandle specified is not one for which the client holds an OPEN\_DELEGATE\_WRITE delegation, an NFS4ERR\_BADHANDLE error is returned.

#### 27.1.4. IMPLEMENTATION

The client returns attrmask bits and the associated attribute values only for the change attribute, and attributes that it may change (time\_modify, and size).

### 27.2. Operation 4: CB\_RECALL - Recall a Delegation

#### 27.2.1. ARGUMENT

```

struct CB_RECALL4args {
    stateid4  stateid;
    bool      truncate;
    nfs_fh4   fh;
};

```

#### 27.2.2. RESULT

```

struct CB_RECALL4res {
    nfsstat4  status;
};

```

### 27.2.3. DESCRIPTION

The CB\_RECALL operation is used to begin the process of recalling a delegation and returning it to the server.

The truncate flag is used to optimize recall for a file object that is a regular file and is about to be truncated to zero. When it is TRUE, the client is freed of the obligation to propagate modified data for the file to the server, since this data is irrelevant.

If the handle specified is not one for which the client holds a delegation, an NFS4ERR\_BADHANDLE error is returned.

If the stateid specified is not one corresponding to an OPEN delegation for the file specified by the filehandle, an NFS4ERR\_BAD\_STATEID is returned.

### 27.2.4. IMPLEMENTATION

The client should reply to the callback immediately. Replying does not complete the recall except when the value of the reply's status field is neither NFS4ERR\_DELAY nor NFS4\_OK. The recall is not complete until the delegation is returned using a DELEGRETURN operation.

## 27.3. Operation 5: CB\_LAYOUTRECALL - Recall Layout from Client

### 27.3.1. ARGUMENT

```

/*
 * NFSv4.1 callback arguments and results
 */

enum layoutrecall_type4 {
    LAYOUTRECALL4_FILE = LAYOUT4_RET_REC_FILE,
    LAYOUTRECALL4_FSID = LAYOUT4_RET_REC_FSID,
    LAYOUTRECALL4_ALL = LAYOUT4_RET_REC_ALL
};

struct layoutrecall_file4 {
    nfs_fh4        lor_fh;
    offset4        lor_offset;
    length4        lor_length;
    stateid4       lor_stateid;
};

union layoutrecall4 switch(layoutrecall_type4 lor_recalltype) {
case LAYOUTRECALL4_FILE:
    layoutrecall_file4 lor_layout;
case LAYOUTRECALL4_FSID:
    fsid4             lor_fsid;
case LAYOUTRECALL4_ALL:
    void;
};

struct CB_LAYOUTRECALL4args {
    layouttype4      clora_type;
    layoutiomode4    clora_iomode;
    bool             clora_changed;
    layoutrecall4    clora_recall;
};

```

### 27.3.2. RESULT

```

struct CB_LAYOUTRECALL4res {
    nfsstat4        clorr_status;
};

```

### 27.3.3. DESCRIPTION

The CB\_LAYOUTRECALL operation is used by the server to recall layouts from the client; as a result, the client will begin the process of returning layouts via LAYOUTRETURN. The CB\_LAYOUTRECALL operation specifies one of three forms of recall processing with the value of layoutrecall\_type4. The recall is for one of the following: a specific layout of a specific file (LAYOUTRECALL4\_FILE), an entire file system ID (LAYOUTRECALL4\_FSID), or all file systems

(LAYOUTRECALL4\_ALL).

The behavior of the operation varies based on the value of the `layoutrecall_type4`. The value and behaviors are:

#### LAYOUTRECALL4\_FILE

For a layout to match the recall request, the values of the following fields must match those of the layout: `clora_type`, `clora_iomode`, `lor_fh`, and the byte-range specified by `lor_offset` and `lor_length`. The `clora_iomode` field may have a special value of `LAYOUTIOMODE4_ANY`. The special value `LAYOUTIOMODE4_ANY` will match any iomode originally returned in a layout; therefore, it acts as a wild card. The other special value used is for `lor_length`. If `lor_length` has a value of `NFS4_UINT64_MAX`, the `lor_length` field means the maximum possible file size. If a matching layout is found, it MUST be returned using the `LAYOUTRETURN` operation (See Section 25.44). An example of the field's special value use is if `clora_iomode` is `LAYOUTIOMODE4_ANY`, `lor_offset` is zero, and `lor_length` is `NFS4_UINT64_MAX`, then the entire layout is to be returned.

The `NFS4ERR_NOMATCHING_LAYOUT` error is only returned when the client does not hold layouts for the file or if the client does not have any overlapping layouts for the specification in the layout recall.

#### LAYOUTRECALL4\_FSID and LAYOUTRECALL4\_ALL

If `LAYOUTRECALL4_FSID` is specified, the `fsid` specifies the file system for which any outstanding layouts MUST be returned. If `LAYOUTRECALL4_ALL` is specified, all outstanding layouts MUST be returned. In addition, `LAYOUTRECALL4_FSID` and `LAYOUTRECALL4_ALL` specify that all the storage device ID to storage device address mappings in the affected file system(s) are also recalled. The respective `LAYOUTRETURN` with either `LAYOUTRETURN4_FSID` or `LAYOUTRETURN4_ALL` acknowledges to the server that the client invalidated the said device mappings. See Section 18.7.5.3.5 for considerations with "bulk" recall of layouts.

The `NFS4ERR_NOMATCHING_LAYOUT` error is only returned when the client does not hold layouts and does not have valid deviceid mappings.

In processing the layout recall request, the client also varies its behavior based on the value of the `clora_changed` field. This field is used by the server to provide additional context for the reason why the layout is being recalled. A `FALSE` value for `clora_changed` indicates that no change in the layout is expected and the client may write modified data to the storage devices involved; this must be

done prior to returning the layout via LAYOUTRETURN. A TRUE value for `clora_changed` indicates that the server is changing the layout. Examples of layout changes and reasons for a TRUE indication are the following: the metadata server is restriping the file or a permanent error has occurred on a storage device and the metadata server would like to provide a new layout for the file. Therefore, a `clora_changed` value of TRUE indicates some level of change for the layout and that the client should not write and commit modified data to the storage devices. In this case, the client writes and commits data through the metadata server.

See Section 18.7.3 for a description of how the `lor_stateid` field in the arguments is to be constructed. Note that the "seqid" field of `lor_stateid` MUST NOT be zero. See Sections 13.2, 18.7.3, and 18.7.5.2 for a further discussion and requirements.

#### 27.3.4. IMPLEMENTATION

The client's processing for CB\_LAYOUTRECALL is similar to CB\_RECALL (recall of file delegations) in that the client responds to the request before actually returning layouts via the LAYOUTRETURN operation. While the client responds to the CB\_LAYOUTRECALL immediately, the operation is not considered complete (i.e., considered pending) until all affected layouts are returned to the server via the LAYOUTRETURN operation.

Before returning the layout to the server via LAYOUTRETURN, the client should wait for the response from in-process or in-flight READ, WRITE, or COMMIT operations that use the recalled layout.

If the client is holding modified data that is affected by a recalled layout, the client has various options for writing the data to the server. As always, the client may write the data through the metadata server. In fact, the client may not have a choice other than writing to the metadata server when the `clora_changed` argument is TRUE and a new layout is unavailable from the server. However, the client may be able to write the modified data to the storage device if the `clora_changed` argument is FALSE; this needs to be done before returning the layout via LAYOUTRETURN. If the client were to obtain a new layout covering the modified data's byte-range, then writing to the storage devices is an available alternative. Note that before obtaining a new layout, the client must first return the original layout.

In the case of modified data being written while the layout is held, the client must use LAYOUTCOMMIT operations at the appropriate time; as required LAYOUTCOMMIT must be done before the LAYOUTRETURN. If a large amount of modified data is outstanding, the client may send

LAYOUTRETURNS for portions of the recalled layout; this allows the server to monitor the client's progress and adherence to the original recall request. However, the last LAYOUTRETURN in a sequence of returns MUST specify the full range being recalled (See Section 18.7.3.1 for details).

If a server needs to delete a device ID and there are layouts referring to the device ID, CB\_LAYOUTRECALL MUST be invoked to cause the client to return all layouts referring to the device ID before the server can delete the device ID. If the client does not return the affected layouts, the server MAY revoke the layouts.

#### 27.4. Operation 6: CB\_NOTIFY - Notify Client Using Directory Delegations

##### 27.4.1. ARGUMENT

```
/* Changed entry information. */
struct notify_entry4 {
    component4    ne_file;
    fattr4        ne_attrs;
};

/* Previous entry information */
struct prev_entry4 {
    notify_entry4 pe_prev_entry;
    /* what REaddir returned for this entry */
    nfs_cookie4   pe_prev_entry_cookie;
};

struct notify_remove4 {
    notify_entry4 nrm_old_entry;
    nfs_cookie4   nrm_old_entry_cookie;
};

/*
 * Objects of types defined below are encoded within
 * notifylist4, depending on the associated bit map,
 * in a fashion similar to the way that attributes
 * are presented within an attrlist4 in a fattr4.
 */
typedef opaque notifylist4<>;

struct notify4 {
    /* composed from notify_type4 or notify_deviceid_type4 */
    bitmap4      notify_mask;
    notifylist4  notify_vals;
};

struct CB_NOTIFY4args {
    stateid4     cna_stateid;
    nfs_fh4      cna_fh;
    notify4      cna_changes<>;
};

cna_stateid designates the associated directory delegation while
cna_fh designates the directory for which the delegation is held.

Each element of cna_changes provides a relevant notification with
type based on notify_mask and associated data, with associated
information in a format that depends on the type, within a nominally
opaque array. The elements are processed in sequence.
```

The bitmap `notify_mask` contains bits whose indices are derived from the enum `notify_type4` defined in Section 16.2.5. The following issues need to be noted:

- \* Many of bits defined in `notify_type4` are flags which do not have an associated notification message, as explained in Section 16.2.5.

If any such bits are set in notify mask, the callback is invalid and `NFS4ERR_INVALID` is to be returned.

- \* If the mask contains and bits in positions not defined as valid elements of the enum `notify_type4`, the callback is invalid and `NFS4ERR_INVALID` is to be returned.
- \* If the bitmask contains no bits set or more than one bit set, the callback is invalid and `NFS4ERR_INVALID` is to be returned.

When an element is found to be invalid, there is no processing of further elements.

#### 27.4.2. RESULT

```
struct CB_NOTIFY4res {  
    nfsstat4    cnr_status;  
};
```

#### 27.4.3. DESCRIPTION

The `CB_NOTIFY` operation is used by the server to send notifications to clients about events related to the maintenance of cached information regarding delegated directories that is used to locally satisfy needs for information normally provided by the use of `LOOKUP`, `REaddir`, and `GETATTR` requests.

The registration of notifications occurs when the delegation is established using `GET_DIR_DELEGATION`. As a result, these notifications are sent over the backchannel, when certain events occur that affect the directory, the files within it, or the delegation itself. Most notifications are sent asynchronously but the sending of some of these is initiated promptly, as part of operations making changes, while others are subject to delay and might be sent periodically. Those sent promptly need to be responded to before the generated operation proceeds but are unlike recalls in that the operation does not await any completing request from the client before proceeding. Those subject to delay can be sent some time after the motivating change occurs. This choice is shown in the Modes column in Table 25.



These notifications are used in providing the following functions:

- \* Notifications relating to the updating of directory contents, as discussed in Section 16.2.11.
- \* Notifications relating to the updating of attributes for directories and objects within them, as discussed in Section 16.2.12.
- \* Notifications relating to authorization for use of cached information in locally satisfying requests, as discussed in Section 16.2.13.

The notifications are sent as list of pairs of bitmaps and values with each bitmap consisting of a single bit selected from the enum `notify_type`, defined in Section 16.2.5 which identifies the specific type of notification being sent. Although the description in Section 9.3.7 is relevant, these bitmaps each have only a single bit set so that the contents of the accompanying opaque array is described by the notification structure associated with that notification type. The individual types are shown in Table 25 with the Modes used in that table defined as follows:

**Synch:**

Notifications are sent synchronously in the context of the operation causing the change that the client needs to be informed about.

When there is a situation in which the notification is to be sent but the client has not requested that type of notification, the delegation is recalled and needs to be returned or revoked before the operation proceeds.

**Ordered:**

Notification are sent promptly in the context of the operation causing the change that the client needs to be informed about.

There is a potential need to order such notifications since processing some notifications in an order different from that in which events occurred can confuse the client. Normally this ordering is provided by putting a number of notifications in the same `CB_NOTIFY` so that they are processed in order

In case in which a server has more notifications than can fit in a single CB\_COMPOUND request, enforcing appropriate ordering will involve serializing multiple CB\_COMPOUND requests. This can involve waiting for responses before sending new callbacks or sending all callbacks associated with a given delegation using the same slot of the session.

When there is a situation in which the notification is to be sent but the client has not requested that type of notification, the delegation is recalled but processing of the operation proceeds without waiting for a client response.

Prompt:

Requests sent promptly as in the case of Ordered notifications but without need for ordering support outside of the context of particular notification types.

The functions of such notifications are either inherently order-independent (e.g., two request to purge a cache are effectively the same as one, independent of the order) or where state is updated protected by an ascending sequence value to prevent difficulties with out-of-order updates.

When there is a situation in which the notification is to be sent but the client has not requested that type of notification, the delegation is recalled but processing of the operation proceeds without waiting for a client response.

Batched:

Sent outside the context of the change, with substantial delays and with no commitment to deliver changes in the order made. Such updates can be sent periodically with sufficient delays between updates to eliminate misordering issues.

When there is a situation in which the notification is to be sent but the client has not requested that type of notification, processing of the operation proceeds normally

Note that for many notifications normally sent batched and described that way in the table below, there are situations in which the client can ask for them be sent using the Ordered approach and the server can undertake to do so. This applies to attribute notifications when the delay of zero is chosen by the client and agreed to by the server,

Name	Function	Data	Mode	Desc.	Disc.
ADD_ENTRY	Add dir. entry	notify_ add4	Ordered	S. 27.4.4	S. 16.2.11
REMOVE_ENTRY	Remove dir. entry	notify4_ remove	Ordered	S. 27.4.5	S. 16.2.11
RENAME_ENTRY	Rename dir. entry	notify_ rename4	Ordered	S. 27.4.6	S. 16.2.11
CHANGE_CHILD_ATTR	Update dir entry attr.	notify4_ chattr	Batched	S. 27.4.8	S. 16.2.12
CHANGE_DIR_ATTR	Update dir attr.	notify4_ dattr	Batched	S. 27.4.7	S. 16.2.12
CHANGE_COOKIE _VERIFIER	Update dir. contents	notify_ verifier4	Prompt	S. 27.4.9	S. 16.2.11
CHANGE_AMASK	Update Attribute Masks	notify_ changeam	Prompt	S. 27.4.10	S. 16.2.12
CHANGE_AUTH	Update Authorization	notify_ changeu4	Prompt	S. 27.4.11	S. 16.2.13
CHANGE_GA	Update getattr processing	notify_ changea4	Prompt	S. 27.4.12	S. 16.2.13

Table 25: Notification Types

## 27.4.4. NOTIFY4\_ADD\_ENTRY

```
struct notify_add4 {
    /*
     * Information on object possibly renamed over.
     * zero-length array if none.
     */
    notify_remove4      nad_old_entry<1>;

    /*
     * Information on new object with requested values of
     * unmodifiable attributes.
     */
    notify_entry4       nad_new_entry;

    /*
     * What REaddir would return for this entry.
     *
     * Will be invalid (length zero) if client is
     * order-unaware.
     */
    nfs_cookie4        nad_new_entry_cookie<1>;

    /*
     * Ordering information.
     *
     * Invalid/ignored if client is order-unaware.
     *
     * zero-length nad_prev_entry if this is the
     * first entry.
     *
     * if nad_prev_entry is length one and contains
     * zero-length name prev_entry is not provided,
     * as happens in the order-unaware and cookie-derived
     * order cases.
     */
    prev_entry4        nad_prev_entry<1>;
    bool               nad_last_entry;
};
```

When this notification is sent, the associated data will be in the form of a notify\_add4, as defined above.

The server will send information about the new directory entry being created. If the client is known to be interested in the order of the entries, the cookie for that entry is also sent. The entry information (data type `notify_add4`) includes the component name of the entry and unmodifiable attributes. These attributes will be the ones specified by the client when the delegation is created, which are not modifiable (`type`, `fileid`, `fsid`, `creation_time`, and `file_handle`).

The server will send this type of entry when a file is actually being created, when an entry is being added to a directory as a result of a rename across directories (See below), and when a hard link is being created to an existing file.

If the client is known to be interested in the order of the entries, additional information to place the new entry as provided as described in the rest of this paragraph. If this entry is added to the end of the directory, the server will set the `nad_last_entry` flag to `TRUE`. If the file is added such that there is at least one entry before it, the server will also return the previous entry information (`nad_prev_entry`, a variable-length array of up to one element. If the array is of zero length, there is no previous entry), along with its cookie. This is to help clients find the right location in their file name caches and directory caches where this entry should be cached. If the new entry's cookie is available, it will be in the `nad_new_entry_cookie` (another variable-length array of up to one element) field.

If the addition of the entry causes another entry to be deleted (which can only happen in the rename case) atomically with the addition, then information on this entry is reported in `nad_old_entry`.

#### 27.4.5. NOTIFY4\_REMOVE\_ENTRY

```
typedef notify4_remove notify_remove4;
```

When this notification is sent, the associated data will be in the form of a `notify4_remove`, as defined in Section 27.4.1

The server will send information about the directory entry being deleted. If the client is order-aware, the server will send the cookie value as part of this.

#### 27.4.6. NOTIFY4\_RENAME\_ENTRY

```
struct notify_rename4 {  
    notify_remove4   nrn_old_entry;  
    notify_add4      nrn_new_entry;  
};
```

When this notification is sent, the associated data will be in the form of a `notify_rename4`, as defined above.

The server will send information about both the old entry and the new entry. This includes the name and attributes for each entry. In addition, if the rename causes the deletion of an entry (i.e., the case of a file renamed over), then this is reported in `nrn_new_new_entry.nad_old_entry`. This notification is only sent if both entries are in the same directory. If the rename is across directories, the server will send a remove notification to one directory and an add notification to the other directory, assuming both have a directory delegation.

#### 27.4.7. Directory Attribute Update Notifications

```
typedef fattr4 notify4_dattr;
```

The `fattr4` identifies the attributes being changed together with the new value for each such attribute.

For most attributes, this notification is sent in a batched or prompt fashion as requested by the delay specified when the delegation is created.

For the attributes `modified_time` and `change`, notification is always prompt and is part of the same `notiff4` as:

- \* An associated add, remove, or rename notification.
- \* A combination of add and remove notifications used to signal the results of a cross-directory RENAME which deletes an existing renamed-over file.
- \* An associated change-verifier notification.

#### 27.4.8. Child Attribute Update Notifications

```
struct notify4_chattr {  
    notify_entry4   na_changed_entry;  
};
```

When this notification is sent, the associated `notify_entry4` will contain two fields:

- \* `ne_file` identifies the entry whose attributes are being reported.
- \* `ne_attrs` provides the changes attribute values.

The client will specify an attribute mask to inform the server of attributes for which it wants to receive notifications. This change notification can be requested for changes the attributes of any file in the directory. The client cannot ask for change attribute notification for a specific file. One attribute mask covers all the files in the directory. Upon any attribute change, the server will send back the values of changed attributes. Notifications might not make sense for some file system-wide attributes, and it is up to the server to decide which subset it wants to support. The client can negotiate the frequency of attribute notifications by letting the server know how often it wants to be notified of an attribute change. The server will return supported notification frequencies or an indication that no notification is permitted for directory or child attributes by setting the `dir_notif_delay` and `dir_entry_notif_delay` attributes, respectively.

Certain attributes are unmodifiable during the life of the file system object. These include `creation_time`, `fileid`, `fsid`, and `file_handle`. Although these attributes cannot appear in a `notify4_chattr`, their inclusion in the mask will cause them to be provided as part of the `ne_attrs` for `notify_entry4` structure appearing in content modification notifications.

#### 27.4.9. NOTIFY4\_CHANGE\_COOKIE\_VERIFIER

```
struct notify_verifier4 {  
    verifier4      nv_old_cookieverf;  
    verifier4      nv_new_cookieverf;  
};
```

When this notification is sent, the associated data will be in the form of a `notify_verifier4`, as defined above.

The holder is informed via this notification of a number of potential events:

- \* When the cookie verifier changes, the client is informed of the new value.
- \* When there is any change in the cookie assigned to an existing directory entry, the client is informed of the change even if the verifier has remained the same.

This is necessary because servers are free to not change cookie verifiers in many cases in which a cookie is changed.

- \* If there is a change in the order of directory entries and the client has previously indicated concern with keeping its order in sync with that of the server by using the NOTIFY4\_CFLAG\_ORDER flag. The notification is sent even if there is no corresponding change in directory entry cookies.

In this case as well, the message can be sent without a verifier change.

Upon receiving this notification, the client can invalidate its cookies and re-send a READDIR to get the new set of entries presented in the server's order together with up-to-date cookies.

#### 27.4.10. Attribute Mask Change Notifications

```
struct notify_changeam4 {  
    uint32_t      ncam_order;  
    bitmap4       ncam_damask;  
    bitmap4       ncam_chmask;  
    bitmap4       ncam_flags;  
};
```

When this notification is sent, the associated data will be in the form of a notify\_changem4, as defined above.

This notification is sent whenever the server wishes to change the set of attributes for which updates are to be sent. This includes the case in which one or both the masks is set to indicate an empty attribute mask. This enables to respond to excessive attribute notification traffic without recalling the delegation.

The fields in the notification are used as follows;

- \* ncam\_order is used to protect against the potential effects of notification misordering. The responder need to compare the ncam\_order value received to the last such value received and only modify the attribute masks if the new value is greater than the last one received.
- \* ncam\_damask is a bit mask identifying the set of attributes of the delegated directory that will be included in subsequent NOTIFY4\_CHANGE\_DIR\_ATTR notifications.



- \* ncam\_chmask is a bit mask identifying the set of attributes for objects identified in entries within the delegated directory that will be included in subsequent NOTIFY4\_CHANGE\_CHILD\_ATTR notifications.
- \* ncam\_flags contains new values for flags originally returned as part of the response to the request to create a directory delegation but might have been changed because the attribute masks are being changed.

The values for the flags NOTIFY4\_PRAN\_CONTENT and NOTIFY4\_PRAN\_AUTH need to be used to determine whether prompt notification for various classes or attribute notifications is present since it is possible a change of masks might have caused the discontinuance of prompt notifications or their new presence when previously not in effect.

#### 27.4.11. Change of Authorization Notifications

```
const NCAU_OWNER           = 1;
const NCAU_GROUP           = 2;
const NCAU_OTHERS          = 4;

typedef uint32_t usemask4;

struct notify_upair {
    usemask4      nusp_ok;
    usemask4      nusp_acc;
};

struct notify_changeau4 {
    uint32_t      ncau_order;
    utf8str_mixed ncau_owner;
    utf8str_mixed ncau_group;
    notify_upair  ncau_lookup;
    notify_upair  ncau_readdir;
    usemask4      ncau_flush;
};
```

The notify\_upair structure used to encode ncau\_lookup and ncau\_readdir has a special arrangement including two usemask4s as described below:

- \* nusp\_ok defines the sets of users for which an authorization check is unnecessary so that, if the covered operation is to be performed locally, the client can be certain that the operation is authorized, without actually performing the checks,

When a local check is necessary, the result can be cached, subject to later flushing as directed by later notifications' `ncau_flush` values.

- \* `nsup_acc` defines the set of users for which an authorization requires that a (remote) `ACCESS` call be done. The `ACCESS` check cannot be replaced by a local authorization check since the authorization result is not motivating this need. Instead, there might be a need for special server-side processing mandated by `AUDIT` or `ALARM` ACEs.

When processing a request involving a user within one of the specified groups, the result cannot be cached, since the check needs to be done irrespective of any cached result.

When this notification is sent, the associated data will be in the form of a `notify_changeu4`, as defined above.

This notification is used to inform the client of necessary changes in the appropriate means of authorization of the local equivalents of `LOOKUP` and `READDIR` operations.

The fields in the notification are used as described below. Many of the fields are in the form of a `usetmask4` which defines the handling of a set of users by including or excluding the directory owner, set of users in the owning group but excluding the directory owner, and all other users, with one bit used for each of those sets.

- \* `ncau_order` is a numeric value used to avoid mistakes when notifications are processed in an unexpected order. The value incremented each time such a notification is sent for a given directory delegation and the client can check for ascending values as discussed below.

For the fields `ncau_owner`, `ncau_group`, `ncau_lookup`, and `ncau_readdir`, the specified changes are to be used to update the client's state only if the `ncau_order` is greater than the last one received.

The field `ncau_flush` is to be acted on unconditionally, regardless of the value of `ncau_order`. Such actions are not qualified by ordering since flushing a cache is an idempotent operation.

- \* `ncau_owner` and `ncau_group` provide the updated values of the directory owner and directory owning group to be used in classifying requests for authorization and in the caching of results from those authorization checks.

- \* `ncal_lookup` provides, for each of the three group of users specified in a `usetmask4`, whether requests to lookup a file by users in that group can be granted without an explicit authorization check or whether an `ACCESS` check is always needed and cannot be foreclosed by a client-side check.
- \* `ncal_readdir` provides, for each of the three group of users specified in a `usetmask4`, whether requests to read the directory by users in that group can be granted without an explicit authorization check or whether an `ACCESS` check is always needed and cannot be foreclosed by a client-side check.
- \* `ncal_flush` indicates, for each of the three group of users specified in a `usetmask4`, whether the cache of `ACCESS` check results for users of that class is to be flushed.

Such cache flushing is necessary when changes in the mode or ACL-related attributes make previous results unreliable and when changes in the owning user or group affect the categorization of users.

#### 27.4.12. Change of GETATTR Processing Notifications

```
enum ncga_astate4 {
    NCGAAS_NO           = 1,
    NCGAAS_ALLOK        = 2,
    NCGAAS_MOSTOK       = 3
};

enum ncga_fstate4 {
    NCGFS_FETCH        = 0,
    NCGFS_COND         = 1
};

enum ncga_ftype4 {
    NCGAT_RESET        = 0,
    NCGAT_SET          = 1,
    NCGAT_ADD          = 2
};

union ncga_finfo4 switch(ncga_fstate f) {
case NCGFS_COND:
    changeid4          chval;
    fattr4_fileid      idc;
default:
    fattr4_fileid      idd;
};

struct notify_changega4 {
    uint32_t           ncga_order;
    ncga_ftype4        ncga_ftype;
    ncga_astate        ncga_astate;
    ncga_finfo4        ncga_finfo;
};
```

When this notification is sent, the associated data will be in the form of a `notify_changega4`, as defined above.

The purpose of this notification is to inform the client of the need to make changes in the handling local equivalents of the `GETATTR` operation using cached data. The potential changes concern:

- \* The need for explicit authorization checks for the local equivalents of `GETATTR`.
- \* The need for explicit refetching of attributes, as opposed to using a previously cached value or making the fetch conditional on knowledge of a particular change id value for the object in question.

At any time, the handling is as directed by the client's current value of its GETATTR authorization state that is represented by one of the three values below.

- \* NCGAS\_NO indicates that explicit ACCESS checks are always necessary.
- \* NCGAS\_ALLOK indicates that explicit ACCESS checks are never necessary, since the operation is known to be authorized a priori.
- \* NCGAS\_SOMEOK that explicit ACCESS checks are necessary only for use of a specific set of files identified by fileid.

The set of files mentioned above includes the following additional information for each entry:

- \* The fileid of the designated file.
- \* An indication whether the current file attributes need always to be explicitly fetched or are to be fetched conditionally, based on the client's current knowledge of the object's change attribute.
- \* A change attribute to compare to the one in the currently cached attribute.

The approach allows multiple entry changes to be resolved using a single READDIR and could be compatible with later delegation changes that allowed entry changes to trigger a prompt notification of directory attribute changes (but probably only when the directory holds no multiply-linked files).

Transition between these states are effected depending on the value of the ncga\_type field of the notification, as described below.

- \* NCGAT\_RESET causes the state to be set to NCGAS\_NO and resets the list of exceptions to empty. In addition all entries will be treated as if a new attribute fetch was always required.
- \* NCGAT\_SET causes the state to be set to NCGAS\_ALLOK or NCGAS\_SOMEOK, depending on whether the list of fileids to be set as exceptions is empty or not. In either case, attribute fetch is specified within the new list while requiring a refetch is the default for fileids not mentioned.
- \* NCGAT\_ADD adds to the list fileids to be used as exceptions and set the state to NCGAS\_SOMEOK if it is currently NCGAS\_ALLOK. Fetching of entries whose fileids are already in the list are updated as specified in the new list.

The `notify_changega4` contains the following fields:

- \* `ncga_order` is a numeric value used to prevent problems when notifications are received by the client in an order different from the one in which they are sent.

Every notification carries the current value maintained on the server. The value is incremented for every notification of type `NCGAT_RESET` that is sent.

For notifications of type `NCGAT_RESET`, the notification is only acted upon if the order value sent is greater than the last one received of that type that is acted upon.

For notifications of other types, the notification is only acted upon if the order value sent is equal the last one received of type `NCGAT_RESET` that is acted upon.

- \* `ncga_ftype` defines the type of state transition, as described above.
- \* `ncga_astype` indicates the new authorization state as described above.
- \* `ncga_finf` contains set of fileids of object within the directory which are to be added to the list of objects for which the local equivalent of `GETATTR` requires explicit `ACCESS` checks. In addition, information relevant to the need to fetch attributes is included, as described above.

#### 27.5. Operation 7: `CB_PUSH_DELEG` - Offer Previously Requested Delegation to Client

##### 27.5.1. ARGUMENT

```
struct CB_PUSH_DELEG4args {
    nfs_fh4      cpda_fh;
    open_delegation4 cpda_delegation;
};
```

##### 27.5.2. RESULT

```
struct CB_PUSH_DELEG4res {
    nfsstat4 cpdr_status;
};
```

### 27.5.3. DESCRIPTION

CB\_PUSH\_DELEG is used by the server both to signal to the client that the delegation it wants (previously indicated via a want established from an OPEN or WANT\_DELEGATION operation) is available and to simultaneously offer the delegation to the client. The client has the choice of accepting the delegation by returning NFS4\_OK to the server, delaying the decision to accept the offered delegation by returning NFS4ERR\_DELAY, or permanently rejecting the offer of the delegation by returning NFS4ERR\_REJECT\_DELEG. When a delegation is rejected in this fashion, the want previously established is permanently deleted and the delegation is subject to acquisition by another client.

### 27.5.4. IMPLEMENTATION

If the client does return NFS4ERR\_DELAY and there is a conflicting delegation request, the server MAY process it at the expense of the client that returned NFS4ERR\_DELAY. The client's want will not be cancelled, but MAY be processed behind other delegation requests or registered wants.

When a client returns a status other than NFS4\_OK, NFS4ERR\_DELAY, or NFS4ERR\_REJECT\_DELAY, the want remains pending, although servers may decide to cancel the want by sending a CB\_WANTS\_CANCELLED.

## 27.6. Operation 8: CB\_RECALL\_ANY - Keep Any N Recallable Objects

### 27.6.1. ARGUMENT

```
const RCA4_TYPE_MASK_RDATA_DLG      = 0;
const RCA4_TYPE_MASK_WDATA_DLG      = 1;
const RCA4_TYPE_MASK_DIR_DLG        = 2;
const RCA4_TYPE_MASK_FILE_LAYOUT    = 3;
const RCA4_TYPE_MASK_BLK_LAYOUT     = 4;
const RCA4_TYPE_MASK_OBJ_LAYOUT_MIN = 8;
const RCA4_TYPE_MASK_OBJ_LAYOUT_MAX = 9;
const RCA4_TYPE_MASK_OTHER_LAYOUT_MIN = 12;
const RCA4_TYPE_MASK_OTHER_LAYOUT_MAX = 15;
```

```
struct CB_RECALL_ANY4args {
    uint32_t      craa_objects_to_keep;
    bitmap4       craa_type_mask;
};
```

### 27.6.2. RESULT

```
struct CB_RECALL_ANY4res {  
    nfsstat4      crar_status;  
};
```

### 27.6.3. DESCRIPTION

The server may decide that it cannot hold all of the state for recallable objects, such as delegations and layouts, without running out of resources. In such a case, while not optimal, the server is free to recall individual objects to reduce the load.

Because the general purpose of such recallable objects as delegations is to eliminate client interaction with the server, the server cannot interpret lack of recent use as indicating that the object is no longer useful. The absence of visible use is consistent with a delegation keeping potential operations from being sent to the server. In the case of layouts, while it is true that the usefulness of a layout is indicated by the use of the layout when storage devices receive I/O requests, because there is no mandate that a storage device indicate to the metadata server any past or present use of a layout, the metadata server is not likely to know which layouts are good candidates to recall in response to low resources.

In order to implement an effective reclaim scheme for such objects, the server's knowledge of available resources must be used to determine when objects must be recalled with the clients selecting the actual objects to be returned.

Server implementations may differ in their resource allocation requirements. For example, one server may share resources among all classes of recallable objects, whereas another may use separate resource pools for layouts and for delegations, or further separate resources by types of delegations.

When a given resource pool is over-utilized, the server can send a CB\_RECALL\_ANY to clients holding recallable objects of the types involved, allowing it to keep a certain number of such objects and return any excess. A mask specifies which types of objects are to be limited. The client chooses, based on its own knowledge of current usefulness, which of the objects in that class should be returned.

A number of bits are defined. For some of these, ranges are defined and it is up to the definition of the storage protocol to specify how these are to be used. There are ranges reserved for object-based storage protocols and for other experimental storage protocols. An RFC defining such a storage protocol needs to specify how particular bits within its range are to be used. For example, it may specify a mapping between attributes of the layout (read vs. write, size of



area) and the bit to be used, or it may define a field in the layout where the associated bit position is made available by the server to the client.

#### RCA4\_TYPE\_MASK\_RDATA\_DLG

The client is to return OPEN\_DELEGATE\_READ delegations on non-directory file objects.

#### RCA4\_TYPE\_MASK\_WDATA\_DLG

The client is to return OPEN\_DELEGATE\_WRITE delegations on regular file objects.

#### RCA4\_TYPE\_MASK\_DIR\_DLG

The client is to return directory delegations.

#### RCA4\_TYPE\_MASK\_FILE\_LAYOUT

The client is to return layouts of type LAYOUT4\_NFSV4\_1\_FILES.

#### RCA4\_TYPE\_MASK\_BLK\_LAYOUT

See [RFC5663] for a description.

#### RCA4\_TYPE\_MASK\_OBJ\_LAYOUT\_MIN to RCA4\_TYPE\_MASK\_OBJ\_LAYOUT\_MAX

See [RFC5664] for a description.

#### RCA4\_TYPE\_MASK\_OTHER\_LAYOUT\_MIN to RCA4\_TYPE\_MASK\_OTHER\_LAYOUT\_MAX

This range is reserved for telling the client to recall layouts of experimental or site-specific layout types (See Section 9.3.13).

When a bit is set in the type mask that corresponds to an undefined type of recallable object, NFS4ERR\_INVAL MUST be returned. When a bit is set that corresponds to a defined type of object but the client does not support an object of the type, NFS4ERR\_INVAL MUST NOT be returned. Future minor versions of NFSv4 may expand the set of valid type mask bits.

CB\_RECALL\_ANY specifies a count of objects that the client may keep as opposed to a count that the client must return. This is to avoid a potential race between a CB\_RECALL\_ANY that had a count of objects to free with a set of client-originated operations to return layouts or delegations. As a result of the race, the client and server would have differing ideas as to how many objects to return. Hence, the client could mistakenly free too many.

If resource demands prompt it, the server may send another CB\_RECALL\_ANY with a lower count, even if it has not yet received an acknowledgment from the client for a previous CB\_RECALL\_ANY with the same type mask. Although the possibility exists that these will be received by the client in an order different from the order in which

they were sent, any such permutation of the callback stream is harmless. It is the job of the client to bring down the size of the recallable object set in line with each `CB_RECALL_ANY` received, and until that obligation is met, it cannot be cancelled or modified by any subsequent `CB_RECALL_ANY` for the same type mask. Thus, if the server sends two `CB_RECALL_ANY`s, the effect will be the same as if the lower count was sent, whatever the order of recall receipt. Note that this means that a server may not cancel the effect of a `CB_RECALL_ANY` by sending another recall with a higher count. When a `CB_RECALL_ANY` is received and the count is already within the limit set or is above a limit that the client is working to get down to, that callback has no effect.

Servers are generally free to deny recallable objects when insufficient resources are available. Note that the effect of such a policy is implicitly to give precedence to existing objects relative to requested ones, with the result that resources might not be optimally used. To prevent this, servers are well advised to make the point at which they start sending `CB_RECALL_ANY` callbacks somewhat below that at which they cease to give out new delegations and layouts. This allows the client to purge its less-used objects whenever appropriate and so continue to have its subsequent requests given new resources freed up by object returns.

#### 27.6.4. IMPLEMENTATION

The client can choose to return any type of object specified by the mask. If a server wishes to limit the use of objects of a specific type, it should only specify that type in the mask it sends. Should the client fail to return requested objects, it is up to the server to handle this situation, typically by sending specific recalls (i.e., sending `CB_RECALL` operations) to properly limit resource usage. The server should give the client enough time to return objects before proceeding to specific recalls. This time should not be less than the lease period.

#### 27.7. Operation 9: `CB_RECALLABLE_OBJ_AVAIL` - Signal Resources for Recallable Objects

##### 27.7.1. ARGUMENT

```
typedef CB_RECALL_ANY4args CB_RECALLABLE_OBJ_AVAIL4args;
```

##### 27.7.2. RESULT

```
struct CB_RECALLABLE_OBJ_AVAIL4res {  
    nfsstat4      croa_status;  
};
```

### 27.7.3. DESCRIPTION

CB\_RECALLABLE\_OBJ\_AVAIL is used by the server to signal the client that the server has resources to grant recallable objects that might previously have been denied by OPEN, WANT\_DELEGATION, GET\_DIR\_DELEG, or LAYOUTGET.

The argument `craa_objects_to_keep` means the total number of recallable objects of the types indicated in the argument `type_mask` that the server believes it can allow the client to have, including the number of such objects the client already has. A client that tries to acquire more recallable objects than the server informs it can have runs the risk of having objects recalled.

The server is not obligated to reserve the difference between the number of the objects the client currently has and the value of `craa_objects_to_keep`, nor does delaying the reply to CB\_RECALLABLE\_OBJ\_AVAIL prevent the server from using the resources of the recallable objects for another purpose. Indeed, if a client responds slowly to CB\_RECALLABLE\_OBJ\_AVAIL, the server might interpret the client as having reduced capability to manage recallable objects, and so cancel or reduce any reservation it is maintaining on behalf of the client. Thus, if the client desires to acquire more recallable objects, it needs to reply quickly to CB\_RECALLABLE\_OBJ\_AVAIL, and then send the appropriate operations to acquire recallable objects.

## 27.8. Operation 10: CB\_RECALL\_SLOT - Change Flow Control Limits

### 27.8.1. ARGUMENT

```
struct CB_RECALL_SLOT4args {  
    slotid4      rsa_target_highest_slotid;  
};
```

### 27.8.2. RESULT

```
struct CB_RECALL_SLOT4res {  
    nfsstat4     rsr_status;  
};
```

### 27.8.3. DESCRIPTION

The CB\_RECALL\_SLOT operation requests the client to return session slots, and if applicable, transport credits (e.g., RDMA credits for connections associated with the operations channel) of the session's fore channel. CB\_RECALL\_SLOT specifies `rsa_target_highest_slotid`, the value of the target highest slot ID the server wants for the session. The client MUST then progress toward reducing the session's highest slot ID to the target value.

If the session has only non-RDMA connections associated with its operations channel, then the client need only wait for all outstanding requests with a slot ID > `rsa_target_highest_slotid` to complete, then send a single COMPOUND consisting of a single SEQUENCE operation, with the `sa_highestslot` field set to `rsa_target_highest_slotid`. If there are RDMA-based connections associated with operation channel, then the client needs to also send enough zero-length "RDMA Send" messages to take the total RDMA credit count to `rsa_target_highest_slotid + 1` or below.

### 27.8.4. IMPLEMENTATION

If the client fails to reduce highest slot it has on the fore channel to what the server requests, the server can force the issue by asserting flow control on the receive side of all connections bound to the fore channel, and then finish servicing all outstanding requests that are in slots greater than `rsa_target_highest_slotid`. Once that is done, the server can then open the flow control, and any time the client sends a new request on a slot greater than `rsa_target_highest_slotid`, the server can return NFS4ERR\_BADSLOT.

## 27.9. Operation 11: CB\_SEQUENCE - Supply Backchannel Sequencing and Control

### 27.9.1. ARGUMENT

```

struct referring_call4 {
    sequenceid4    rc_sequenceid;
    slotid4        rc_slotid;
};

struct referring_call_list4 {
    sessionid4     rcl_sessionid;
    referring_call4 rcl_referring_calls<>;
};

struct CB_SEQUENCE4args {
    sessionid4      csa_sessionid;
    sequenceid4     csa_sequenceid;
    slotid4         csa_slotid;
    slotid4         csa_highest_slotid;
    bool            csa_cachethis;
    referring_call_list4 csa_referring_call_lists<>;
};

```

#### 27.9.2. RESULT

```

struct CB_SEQUENCE4resok {
    sessionid4      csr_sessionid;
    sequenceid4     csr_sequenceid;
    slotid4         csr_slotid;
    slotid4         csr_highest_slotid;
    slotid4         csr_target_highest_slotid;
};

union CB_SEQUENCE4res switch (nfsstat4 csr_status) {
case NFS4_OK:
    CB_SEQUENCE4resok    csr_resok4;
default:
    void;
};

```

#### 27.9.3. DESCRIPTION

The CB\_SEQUENCE operation is used to manage operational accounting for the backchannel of the session on which a request is sent. The contents include the session ID to which this request belongs, the slot ID and sequence ID used by the server to implement session request control and exactly once semantics, and exchanged slot ID maxima that are used to adjust the size of the reply cache. In each CB\_COMPOUND request, CB\_SEQUENCE MUST appear once and MUST be the first operation. The error NFS4ERR\_SEQUENCE\_POS MUST be returned when CB\_SEQUENCE is found in any position in a CB\_COMPOUND beyond the first. If any other operation is in the first position of

CB\_COMPOUND, NFS4ERR\_OP\_NOT\_IN\_SESSION MUST be returned.

See Section 25.46.3 for a description of how slots are processed.

If `csa_cachethis` is TRUE, then the server is requesting that the client cache the reply in the callback reply cache. The client MUST cache the reply (See Section 7.6.1.3).

The `csa_referring_call_lists` array is the list of COMPOUND requests, identified by session ID, slot ID, and sequence ID. These are requests that the client previously sent to the server. These previous requests created state that some operation(s) in the same CB\_COMPOUND as the `csa_referring_call_lists` are identifying. A session ID is included because leased state is tied to a client ID, and a client ID can have multiple sessions. See Section 7.6.3.

The value of the `csa_sequenceid` argument relative to the cached sequence ID on the slot falls into one of three cases.

- \* If the difference between `csa_sequenceid` and the client's cached sequence ID at the slot ID is two (2) or more, or if `csa_sequenceid` is less than the cached sequence ID (accounting for wraparound of the unsigned sequence ID value), then the client MUST return NFS4ERR\_SEQ\_MISORDERED.
- \* If `csa_sequenceid` and the cached sequence ID are the same, this is a retry, and the client returns the CB\_COMPOUND request's cached reply.
- \* If `csa_sequenceid` is one greater (accounting for wraparound) than the cached sequence ID, then this is a new request, and the slot's sequence ID is incremented. The operations subsequent to CB\_SEQUENCE, if any, are processed. If there are no other operations, the only other effects are to cache the CB\_SEQUENCE reply in the slot, maintain the session's activity, and when the server receives the CB\_SEQUENCE reply, renew the lease of state related to the client ID.

If the server reuses a slot ID and sequence ID for a completely different request, the client MAY treat the request as if it is a retry of what it has already executed. The client MAY however detect the server's illegal reuse and return NFS4ERR\_SEQ\_FALSE\_RETRY.

If CB\_SEQUENCE returns an error, then the state of the slot (sequence ID, cached reply) MUST NOT change. See Section 7.6.1.3 for the conditions when the error NFS4ERR\_RETRY\_UNCACHED\_REP might be returned.

The client returns two "highest\_slotid" values: `csr_highest_slotid` and `csr_target_highest_slotid`. The former is the highest slot ID the client will accept in a future `CB_SEQUENCE` operation, and should not be less than the value of `csa_highest_slotid` (but see Section 7.6.1 for an exception). The latter is the highest slot ID the client would prefer the server use on a future `CB_SEQUENCE` operation.

#### 27.10. Operation 12: `CB_WANTS_CANCELLED` - Cancel Pending Delegation Wants

##### 27.10.1. ARGUMENT

```
struct CB_WANTS_CANCELLED4args {  
    bool cwca_contended_wants_cancelled;  
    bool cwca_resourced_wants_cancelled;  
};
```

##### 27.10.2. RESULT

```
struct CB_WANTS_CANCELLED4res {  
    nfsstat4      cwcr_status;  
};
```

##### 27.10.3. DESCRIPTION

The `CB_WANTS_CANCELLED` operation is used to notify the client that some or all of the wants it registered for recallable delegations and layouts have been cancelled.

If `cwca_contended_wants_cancelled` is `TRUE`, this indicates that the server will not be pushing to the client any delegations that become available after contention passes.

If `cwca_resourced_wants_cancelled` is `TRUE`, this indicates that the server will not notify the client when there are resources on the server to grant delegations or layouts.

After receiving a `CB_WANTS_CANCELLED` operation, the client is free to attempt to acquire the delegations or layouts it was waiting for, and possibly re-register wants.

#### 27.10.4. IMPLEMENTATION

When a client has an OPEN, WANT\_DELEGATION, or GET\_DIR\_DELEGATION request outstanding, when a CB\_WANTS\_CANCELLED is sent, the server may need to make clear to the client whether a promise to signal delegation availability happened before the CB\_WANTS\_CANCELLED and is thus covered by it, or after the CB\_WANTS\_CANCELLED in which case it was not covered by it. The server can make this distinction by putting the appropriate requests into the list of referring calls in the associated CB\_SEQUENCE.

#### 27.11. Operation 13: CB\_NOTIFY\_LOCK - Notify Client of Possible Lock Availability

##### 27.11.1. ARGUMENT

```
struct CB_NOTIFY_LOCK4args {
    nfs_fh4      cnla_fh;
    lock_owner4  cnla_lock_owner;
};
```

##### 27.11.2. RESULT

```
struct CB_NOTIFY_LOCK4res {
    nfsstat4      cnlr_status;
};
```

##### 27.11.3. DESCRIPTION

The server can use this operation to indicate that a byte-range lock for the given file and lock-owner, previously requested by the client via an unsuccessful LOCK operation, might be available.

This callback is meant to be used by servers to help reduce the latency of blocking locks in the case where they recognize that a client that has been polling for a blocking byte-range lock may now be able to acquire the lock. If the server supports this callback for a given file, it MUST set the OPEN4\_RESULT\_MAY\_NOTIFY\_LOCK flag when responding to successful opens for that file. This does not commit the server to the use of CB\_NOTIFY\_LOCK, but the client may use this as a hint to decide how frequently to poll for locks derived from that open.

If an OPEN operation results in an upgrade, in which the stateid returned has an "other" value matching that of a stateid already allocated, with a new "seqid" indicating a change in the lock being represented, then the value of the OPEN4\_RESULT\_MAY\_NOTIFY\_LOCK flag when responding to that new OPEN controls handling from that point



going forward. When parallel OPENS are done on the same file and open-owner, the ordering of the "seqid" fields of the returned stateids (subject to wraparound) are to be used to select the controlling value of the OPEN4\_RESULT\_MAY\_NOTIFY\_LOCK flag.

#### 27.11.4. IMPLEMENTATION

The server MUST NOT grant the byte-range lock to the client unless and until it receives a LOCK operation from the client. Similarly, the client receiving this callback cannot assume that it now has the lock or that a subsequent LOCK operation for the lock will be successful.

The server is not required to implement this callback, and even if it does, it is not required to use it in any particular case. Therefore, the client must still rely on polling for blocking locks, as described in Section 14.6.

Similarly, the client is not required to implement this callback, and even it does, is still free to ignore it. Therefore, the server MUST NOT assume that the client will act based on the callback.

#### 27.12. Operation 14: CB\_NOTIFY\_DEVICEID - Notify Client of Device ID Changes

##### 27.12.1. ARGUMENT

```

/*
 * Device notification types.
 */
enum notify_deviceid_type4 {
    NOTIFY_DEVICEID4_CHANGE = 1,
    NOTIFY_DEVICEID4_DELETE = 2
};

/* For NOTIFY4_DEVICEID4_DELETE */
struct notify_deviceid_delete4 {
    layouttype4    ndd_layouttype;
    deviceid4      ndd_deviceid;
};

/* For NOTIFY4_DEVICEID4_CHANGE */
struct notify_deviceid_change4 {
    layouttype4    ndc_layouttype;
    deviceid4      ndc_deviceid;
    bool           ndc_immediate;    /* Unused */
};

struct CB_NOTIFY_DEVICEID4args {
    notify4 cnda_changes<>;
};

```

#### 27.12.2. RESULT

```

struct CB_NOTIFY_DEVICEID4res {
    nfsstat4      cndr_status;
};

```

#### 27.12.3. DESCRIPTION

The CB\_NOTIFY\_DEVICEID operation is used by the server to send notifications to clients about changes to pNFS device IDs. The registration of device ID notifications is optional and is done via GETDEVICEINFO. These notifications are sent over the backchannel once the original request has been processed on the server. The server will send an array of notifications, cnda\_changes, as a list of pairs of bitmaps and values. See Section 9.3.7 for a description of how NFSv4.1 bitmaps work.

As with CB\_NOTIFY (Section 27.4.3), it is possible the server has more notifications than can fit in a CB\_COMPOUND, thus requiring multiple CB\_COMPOUNDS. Unlike CB\_NOTIFY, serialization is not an issue because unlike directory entries, device IDs cannot be re-used after being deleted (Section 18.4.9).

All device ID notifications contain a device ID and a layout type. The layout type is necessary because two different layout types can share the same device ID, and the common device ID can have completely different mappings for each layout type.

The server will send the following notifications:

#### NOTIFY\_DEVICEID4\_CHANGE

A previously provided device-ID-to-device-address mapping has changed and the client uses GETDEVICEINFO to obtain the updated mapping. The notification is encoded in a value of data type `notify_deviceid_change4`. This data type contains an unused boolean field, `ndc_immediate`, which provides no useful information to the client. The client is permitted to finish outstanding I/O that references the previously provided device-ID-to-device-address mapping. Before requesting new layouts, the client needs to replace the previously provided device-ID-to-device-address mapping using a GETDEVINFO operation. All outstanding layouts remain valid after a notification of type NOTIFY\_DEVICEID4\_CHANGE. If the device-ID-to-device-address mapping changed in an incompatible way, that would invalidate outstanding layouts, the server MUST recall all outstanding layouts and send a NOTIFY\_DEVICEID4\_DELETE notification instead

#### NOTIFY4\_DEVICEID\_DELETE

Deletes a device ID from the mappings. This notification MUST NOT be sent if the client has a layout that refers to the device ID. In other words, if the server is sending a delete device ID notification, one of the following is true for layouts associated with the layout type:

- \* The client never had a layout referring to that device ID.
- \* The client has returned all layouts referring to that device ID.
- \* The server has revoked all layouts referring to that device ID.

The notification is encoded in a value of data type `notify_deviceid_delete4`. After a server deletes a device ID, it MUST NOT reuse that device ID for the same layout type until the client ID is deleted.

### 27.13. Operation 10044: CB\_ILLEGAL - Illegal Callback Operation

#### 27.13.1. ARGUMENT

void;

#### 27.13.2. RESULT

```
/*
 * CB_ILLEGAL: Response for illegal operation numbers
 */
struct CB_ILLEGAL4res {
    nfsstat4      status;
};
```

#### 27.13.3. DESCRIPTION

This operation is a placeholder for encoding a result to handle the case of the server sending an operation code within CB\_COMPOUND that is not defined in the NFSv4.1 specification. See Section 26.2.3 for more details.

The status field of CB\_ILLEGAL4res MUST be set to NFS4ERR\_OP\_ILLEGAL.

#### 27.13.4. IMPLEMENTATION

A server will probably not send an operation with code OP\_CB\_ILLEGAL, but if it does, the response will be CB\_ILLEGAL4res just as it would be with any other invalid operation code. Note that if the client gets an illegal operation code that is not OP\_ILLEGAL, and if the client checks for legal operation codes during the XDR decode phase, then an instance of data type CB\_ILLEGAL4res will not be returned.

## 28. Security Considerations

The majority of the Security Considerations relevant to NFS Version 4.1 will appear in the corresponding section of the document devoted to the security of NFS Version 4 as a whole [I-D.dnoveck-nfsv4-security]. Some Security considerations relating to the use of pNFS and other NFSv4.1-specific features will be dealt with in subsections of this section.

### 28.1. Issues with Inherited Security Considerations Section

The following significant issues need to be addressed in a new Security Considerations section for NFSv4 in whatever document it appears:

- \* The absence of a threat analysis.
- \* The lack of attention to the security consequences of transmission of user data in the clear.
- \* The treatment of AUTH\_SYS as OPTIONAL without any discussion of the security consequences of using it.

It is anticipated that such a revised Security Considerations section will appear in the NFSV4-wide security document and that the corresponding section will deal with Security Considerations (including a threat analysis) for NFSv4.1-specific features such as parallel NFS.

## 28.2. Threat Analysis

Possible additional threats raised by new features in NFSv4.1 will be dealt as follows:

- \* There do not appear to be additional threats arising from the use of sessions per se. State protection, originally discussed, as part NFSv4.1, is now dealt with in NFSv4-wide security document, rather than in this one.

Threats related to the persistent storage of session state and locking state are dealt with in Section 28.2.1.

- \* Threats related to the use of pNFS will be dealt with in Section 28.2.2 and its subsections.

### 28.2.1. Threat Analysis for Use of Persistent Sessions and Locking State

Locking state can be transferred between two different client-server associations as a result of server restart. This raises the possibility that the transfer might be made inappropriately if a hostile client presents itself as the owner of existing persistently-preserved locking state created before the server restart.

In order to prevent any such misuse, servers that implement persistent sessions or locking state MUST do the following:

- \* Limit the persistent storage of state to situations in which the client peer owning that state is identifiable (e.g. by the use of client-peer identification together with use of RPC-with-tls).

- \* Persistently store the client identification together with the locking or reply state being maintained across potential server restart.
- \* Only restore persisted state when the successor client peer securely identifies itself with identification matching that stored when the state was created.
- \* Mark such persisted state as having been restored to prevent it being used by a second client peer instance.

It should be noted that there are similar situations in which state created by one client peer might be incorrectly accessed by another with current specifications taking a laxer approach as described in [I-D.dnoveck-nfsv4-security]. For various reasons, it was not possible for these older features to require the level of strictness applied above to persisted locking state:

#### 28.2.2.2. Threat Analysis for Use of pNFS

The threat analysis is divided based on layout type and coupling mode. Although most layout types only support a single coupling mode, the flexible files layout is designed to support multiple coupling modes with the result that its use with tight and loose coupling need to be dealt with separately.

- \* For layout types that use RPC for data access and rely on the support of a separate control protocol (i.e. The files layout type described in Section 20 and the flexible files layout described in [RFC8435] when used in tight coupling mode.), the material in Section 18.11.2 provides a general picture of the security issues while the corresponding threat analysis appears in Section 28.2.2.3.
- \* For layout types that use RPC for data access and have no separate control protocol (i.e. The flexible files layout described in [RFC8435] when used in loose coupling mode.), the material in Section 18.11.3 provides a general picture of the security issues while the corresponding threat analysis appears in Section 28.2.2.2.
- \* For layout types that do not use RPC for data access and have a separate control protocol (e.g. the blocks layout [RFC5663], the SCSI layout [RFC8154], and the objects layout [RFC5664]), the material in Section 18.11.1 provides a general picture of the security issues while the corresponding threat analysis appears in Section 28.2.2.1.

#### 28.2.2.1. Threat Analysis for pNFS Layout Types Involving non-RPC Data Access

Because data is accessed using non-RPC protocols, there might be no provision for authenticating or even identifying the requester asking for data to be read or written. In such situations, every client whose data access requests are executed has to be trusted to make such requests only in cases in which they would be authorized if pNFS were not used.

Clients that make IO requests in such environments have the ability to access or modify data when they are not authorized to do so. This provides a means whereby hostile clients can compromise data security.

Servers can address such threats by limiting use of non-RPC data access to clients who can be trusted to only make data access requests that would be authorized in the non-pNFS case. The client-peer authentication provided by RPC-with-TLS can be helpful in this regard while the use of such data access without such authentication gives rise likely compromise of necessary data security.

#### 28.2.2.2. Threat Analysis for Layout Types Using NFS versions as Data Access Protocols without Control Protocol Assistance

Although RPC is used for data access in this case, it is not used identify or authenticate the principal making the data access request. Instead, AUTH\_SYS is used and the requesting principal is specified by the layout. As a result, the situation, in terms of threats to data integrity, is similar to that described in Section 28.2.2.1.

[Author Aside]: TH Needs to review this section.

#### 28.2.2.3. Threat Analysis for pNFS Layout Types using NFSv4.1 as a Data Access Protocol Together with Control Protocol Assistance

Here we need to look at each of the possible pairs of communicating entities individually:

- \* For communication between the clients and either the metadata server to data storage devices (aka data servers), RPC should be used and the threat analysis in the NFSv4-wide security applies just as it does in the non-pNFS case. There will be no separate threat analysis in this document.

- \* For communication between the data storage devices and the metadata server or between two data storage devices, it is necessary that the communicating peers mutually authenticate and there needs to be a trust relationship between the peers. See Section 28.2.2.4 for an analysis of inter-component trust relationships for the various layout types.

#### 28.2.2.4. pNFS and Inter-component Trust Relationships.

When pNFS is not involved, there are only two actors involved in file access: the server and the client. There are two possibilities regarding trust relationships:

- \* With authenticated principals (e.g. RPCSEC\_GSS), there is no need for the server to trust the client or the user.

The client does need to trust the server to obey the protocol and not provide information about its requests to others

- \* When AUTH\_SYS is used, the server has to trust the client to correctly authenticate user principals. To deal with the possibility that clients might take advantage of this situation to cause the server to execute unauthenticated requests, RPC-with-tls can be used together with client peer authentication to limit the set of clients whose unauthenticated requests are accepted or to limit the set of principals that can be identified in this way. At a minimum, acceptance of requests identified as made by root would be limited.

When pNFS is involved, the situation is potentially more complicated in that there are three sorts of actors and six potential trust relationships.

- \* The interactions between the client and the metadata server are as described above.
- \* The interactions between client and data server will vary depending on the layout type.

For the cases described in Section 28.2.2.3, interactions between client and data server follows the same model as that between the client and metadata server. This case, when principals are authenticated on both the MDS and the data server (i.e., when RPC\_SECGSS is used) is the only one in which the client does not have to be specially trusted

Other mapping types will be discussed below.



- \* Trust relationships between metadata server and data servers need to exist, although, depending on the details of the mapping type, the client might not be aware of this distinction, and consider both servers together as a unified entity.

How this issue is addressed for various mapping types are discussed below.

Unlike the cases in which tight coupling is provided, for other mapping types we have no way of authenticating IO requesters and need to trust the client to determine when IO operations are authorized. The resulting situation is similar to that in which AUTH\_SYS is used. As is the case when AUTH\_SYS is used, the client peer needs to identify itself so that the server can use this identity to avoid a case in which hostile node represents itself as having the ability to individually authorize IO requests given a layout obtained when file was opened.

Because layouts are per-client, rather than per-principal, objects, IO authorization needs to be checked for each request, as would be the case if pNFS were not used. Of particular concern is the case in which the principal performing the IO is not the same as the one opening the file.

- \* In the case of such layout types as block or object, RPC is not used so there is no opportunity to identify the principal making the request to the data server.

In the case in which the principal making the IO request is not the opener, the client needs to use ACCESS to ascertain the authorization for the IO request.

- \* In the case of the flexible files layout in the loose coupling mode, the issues are similar, even though RPC is used. Although, the principal identification could be sent, this mapping type specifies that a different unrelated principal identified in the layout is to be passed. As a result, the principal issuing the request is not identified in the RPC request.

Just as with mapping types that do not use RPC, in the case in which the principal making the IO request is not the opener, the client needs to use ACCESS to ascertain the authorization for the IO request. The principal can be authenticated at this point if the metadata server is accessed using RPC\_SECGSS.

Issues regarding trust relationships between the metadata server and data servers are discussed below:

- \* In implementing the file layout and, one assumes, in the case of the flexible files layout in tight coupling mode, there is necessarily a (two-way) trust relationship between metadata server and data server. However, because the nature of their relationship requires adherence to an undocumented control protocol, those outside of the two servers are not in a position to verify or understand the mutual requirements of the two servers.

For all practical purposes, the client considers the two server interfaces as a unified whole and has a trust relationship with those two server interfaces considered together.

In particular, authorization decisions for IO requests are to have the same results whether the data server or metadata server and is the responsibility of the two servers to appropriately coordinate to ensure that, just as it needs to coordinate locking globally.

- \* For layout types such as block and object there is no special control protocol but the data access protocol is a subset of the data device's protocol while the control protocol has access to a different (usually larger) subset of that same data device protocol.

For the clients to be assured that their data is safe, there need to be restrictions on the use of the data device's protocol to prevent hostile clients getting access to data without authorization by the metadata sever.

- \* When the flexible files layout type is used in the loose coupling mode, the situation is similar to those in which block or object protocols are used. Just as in those cases, the data device's protocol is NFSv3 but the data access protocol uses a restricted subset with special constraints about how authorization is determined. The metadata server has access to essentially the full NFSv3 protocol.

Although clients are unlikely to be aware of the details, the safety of their data depends on limiting access by additional clients to data servers that are not restricted to the NFSv3 subset to be used by data servers. In order to effect the necessary limiting, the use of client peer authentication is needed to prevent use by hostile clients that are not prepared to implement the equivalent of the metadata server's authorization decisions. The potential damage is expanded if such hostile clients have access to the full NFSv3 protocol.

## 29. IANA Considerations

This section uses terms that are defined in [RFC8126].

### 29.1. IANA Actions

This update does not require any modification of, or additions to, registry entries or registry rules associated with NFSv4.1. However, since this document obsoletes RFC 8881, IANA is presumed to have updated all registry entries and registry rules references that point to RFC 8881 to point to this document instead.

Previous actions by IANA related to NFSv4.1 are listed in the remaining subsections of Section 29.

### 29.2. Named Attribute Definitions

IANA has created a registry called the "NFSv4 Named Attribute Definitions Registry". It has not been used and there is considerable doubt about whether it ever will be.

The NFSv4.1 protocol supports the association of a file with zero or more named attributes. The identifiers for these attributes are defined as string names. The protocol does not define the specific assignment of the namespace for these file attributes. The IANA registry was intended to promote interoperability where common interests exist but that approach now seems ill-advised for the following reasons:

- \* Since the feature creates a virtual directory, there is no more reason for users to register their names than there is for any other directory. The fact that the feature is referred to as "named attributes" doesn't change that.
- \* There is no point in registering this as part of the protocol, since, when it is used there are similar names in the local filesystem and for access by SMB that are outside the scope of what the IETF can and should control.

As a result, specification of a naming standard that applies to NFSv4 only when these named data streams are accessed using NFSv4 needs to be avoided.

While application developers are allowed to define and use attributes as needed, they were previously encouraged to register the attributes with IANA. However, as discussed above, the fact that the names apply to file systems other than those accessed by NFSv4 implies that it is not an appropriate way to promote interoperability.

The registry was specified as being maintained using the Specification Required policy as defined in Section 4.6 of [RFC8126].

The registry of named attributes was defined a list of assignments, each containing three fields for each assignment.

1. A US-ASCII string name that is the actual name of the attribute. This name must be unique. This string name can be 1 to 128 UTF-8 characters long.
2. A reference to the specification of the named attribute. The reference can consume up to 256 bytes (or more if IANA permits).
3. The point of contact of the registrant. The point of contact can consume up to 256 bytes (or more if IANA permits).

#### 29.2.1. Initial Registry

There is no initial registry.

#### 29.2.2. Updating Registrations

The registrant is always permitted to update the point of contact field. Any other change was expected to require Expert Review or IESG Approval.

#### 29.3. Device ID Notifications

IANA created a registry called the "NFSv4 Device ID Notifications Registry".

The potential exists for new notification types to be added to the CB\_NOTIFY\_DEVICEID operation (See Section 27.12). This can be done via changes to the operations that register notifications, or by adding new operations to NFSv4. This requires a new minor version of NFSv4, and requires a Standards Track document from the IETF. Another way to add a notification is to specify a new layout type (See Section 29.5).

Hence, all assignments to the registry are made on a Standards Action basis per Section 4.6 of [RFC8126], with Expert Review required.

The registry is a list of assignments, each containing five fields per assignment.

1. The name of the notification type. This name must have the prefix "NOTIFY\_DEVICEID4\_". This name must be unique.

2. The value of the notification. IANA will assign this number, and the request from the registrant will use TBD1 instead of an actual value. IANA MUST use a whole number that can be no higher than  $2^{32}-1$ , and should be the next available value. The value assigned must be unique. A Designated Expert must be used to ensure that when the name of the notification type and its value are added to the NFSv4.1 `notify_deviceid_type4` enumerated data type in the NFSv4.1 XDR description [RFC5662], the result continues to be a valid XDR description.
3. The Standards Track RFC(s) that describe the notification. If the RFC(s) have not yet been published, the registrant will use RFCTBD20, RFCTBD21, etc. instead of an actual RFC number.
4. How the RFC introduces the notification. This is indicated by a single US-ASCII value. If the value is N, it means a minor revision to the NFSv4 protocol. If the value is L, it means a new pNFS layout type. Other values can be used with IESG Approval.
5. The minor versions of NFSv4 that are allowed to use the notification. While these are numeric values, IANA will not allocate and assign them; the author of the relevant RFCs with IESG Approval assigns these numbers. Each time there is a new minor version of NFSv4 approved, a Designated Expert should review the registry to make recommended updates as needed.

### 29.3.1. Initial Registry

The initial registry is in Table 26. Note that the next available value is zero.

Notification Name	Value	RFC	How	Minor Versions
NOTIFY_DEVICEID4_CHANGE	1	RFC 8881	N	1
NOTIFY_DEVICEID4_DELETE	2	RFC 8881	N	1

Table 26: Initial Device ID Notification Assignments

### 29.3.2. Updating Registrations

The update of a registration will require IESG Approval on the advice of a Designated Expert.

#### 29.4. Object Recall Types

IANA created a registry called the "NFSv4 Recallable Object Types Registry".

The potential exists for new object types to be added to the CB\_RECALL\_ANY operation (see Section 27.6). This can be done via changes to the operations that add recallable types, or by adding new operations to NFSv4. This requires a new minor version of NFSv4, and requires a Standards Track document from IETF. Another way to add a new recallable object is to specify a new layout type (See Section 29.5).

All assignments to the registry are made on a Standards Action basis per Section 4.9 of [RFC8126], with Expert Review required.

Recallable object types are 32-bit unsigned numbers. There are no Reserved values. Values in the range 12 through 15, inclusive, are designated for Private Use.

The registry is a list of assignments, each containing five fields per assignment.

1. The name of the recallable object type. This name must have the prefix "RCA4\_TYPE\_MASK\_". The name must be unique.
2. The value of the recallable object type. IANA will assign this number, and the request from the registrant will use TBD1 instead of an actual value. IANA MUST use a whole number that can be no higher than  $2^{32}-1$ , and should be the next available value. The value must be unique. A Designated Expert must be used to ensure that when the name of the recallable type and its value are added to the NFSv4 XDR description [RFC5662], the result continues to be a valid XDR description.
3. The Standards Track RFC(s) that describe the recallable object type. If the RFC(s) have not yet been published, the registrant will use RFCTBD2, RFCTBD3, etc. instead of an actual RFC number.
4. How the RFC introduces the recallable object type. This is indicated by a single US-ASCII value. If the value is N, it means a minor revision to the NFSv4 protocol. If the value is L, it means a new pNFS layout type. Other values can be used with IESG Approval.
5. The minor versions of NFSv4 that are allowed to use the recallable object type. While these are numeric values, IANA will not allocate and assign them; the author of the relevant

RFCs with IESG Approval assigns these numbers. Each time there is a new minor version of NFSv4 approved, a Designated Expert should review the registry to make recommended updates as needed.

#### 29.4.1. Initial Registry

The initial registry is in Table 27. Note that the next available value is five.

Recallable Object Type Name	Value	RFC	How	Minor Versions
RCA4_TYPE_MASK_RDATA_DLG	0	RFC 8881	N	1
RCA4_TYPE_MASK_WDATA_DLG	1	RFC 8881	N	1
RCA4_TYPE_MASK_DIR_DLG	2	RFC 8881	N	1
RCA4_TYPE_MASK_FILE_LAYOUT	3	RFC 8881	N	1
RCA4_TYPE_MASK_BLK_LAYOUT	4	RFC 8881	L	1
RCA4_TYPE_MASK_OBJ_LAYOUT_MIN	8	RFC 8881	L	1
RCA4_TYPE_MASK_OBJ_LAYOUT_MAX	9	RFC 8881	L	1

Table 27: Initial Recallable Object Type Assignments

#### 29.4.2. Updating Registrations

The update of a registration will require IESG Approval on the advice of a Designated Expert.

#### 29.5. Layout Types

IANA created a registry called the "pNFS Layout Types Registry".

All assignments to the registry are made on a Standards Action basis, with Expert Review required.

Layout types are 32-bit numbers. The value zero is Reserved. Values in the range 0x80000000 to 0xFFFFFFFF inclusive are designated for Private Use. IANA will assign numbers from the range 0x00000001 to 0x7FFFFFFF inclusive.

The registry is a list of assignments, each containing five fields.

1. The name of the layout type. This name must have the prefix "LAYOUT4\_". The name must be unique.
2. The value of the layout type. IANA will assign this number, and the request from the registrant will use TBD1 instead of an actual value. The value assigned must be unique. A Designated Expert must be used to ensure that when the name of the layout type and its value are added to the NFSv4.1 layouttype4 enumerated data type in the NFSv4.1 XDR description [RFC5662], the result continues to be a valid XDR description.
3. The Standards Track RFC(s) that describe the notification. If the RFC(s) have not yet been published, the registrant will use RFCTBD2, RFCTBD3, etc. instead of an actual RFC number. Collectively, the RFC(s) must adhere to the requirements listed in Section 29.5.3 and 19.1
4. How the RFC introduces the layout type. This is indicated by a single US-ASCII value. If the value is N, it means a minor revision to the NFSv4 protocol. If the value is L, it means a new pNFS layout type. Other values can be used with IESG Approval.
5. The minor versions of NFSv4 that are allowed to use the notification. While these are numeric values, IANA will not allocate and assign them; the author of the relevant RFCs with IESG Approval assigns these numbers. Each time there is a new minor version of NFSv4 approved, a Designated Expert should review the registry to make recommended updates as needed.

#### 29.5.1. Initial Registry

The initial registry is in Table 28.



Layout Type Name	Value	RFC	How	Minor Versions
LAYOUT4_NFSV4_1_FILES	0x1	RFC 8881	N	1
LAYOUT4_OSD2_OBJECTS	0x2	RFC 5664	L	1
LAYOUT4_BLOCK_VOLUME	0x3	RFC 5663	L	1

Table 28: Initial Layout Type Assignments

### 29.5.2. Updating Registrations

The update of a registration will require IESG Approval on the advice of a Designated Expert.

### 29.5.3. IANA-Related Requirements for Layout Type Specifications

Specification of new layout types takes the form Standards-track RFCs. Relevant requirements for these specifications are defined in Section 19.

In addition, the specification needs to include an IANA considerations section, which will in turn include:

- \* A request to IANA for a new layout type per Section 29.5.
- \* A list of requests to IANA for any new recallable object types for CB\_RECALL\_ANY; each entry is to be presented in the form described in Section 29.4.
- \* A list of requests to IANA for any new notification values for CB\_NOTIFY\_DEVICEID; each entry is to be presented in the form described in Section 29.3.

### 29.6. Path Variable Definitions

This section deals with the IANA considerations associated with the variable substitution feature for location names as described in Section 17.17.3. As described there, variables subject to substitution consist of a domain name and a specific name within that domain, with the two separated by a colon. There are two sets of IANA considerations here:

1. The list of variable names.
2. For each variable name, the list of possible values.

Thus, there will be one registry for the list of variable names, and possibly one registry for listing the values of each variable name.

#### 29.6.1. Path Variables Registry

IANA created a registry called the "NFSv4 Path Variables Registry".

##### 29.6.1.1. Path Variable Values

Variable names are of the form "\${", followed by a domain name, followed by a colon (":"), followed by a domain-specific portion of the variable name, followed by "}". When the domain name is "ietf.org", all variables names must be registered with IANA on a Standards Action basis, with Expert Review required. Path variables with registered domain names neither part of nor equal to ietf.org are assigned on a Hierarchical Allocation basis (delegating to the domain owner) and thus of no concern to IANA, unless the domain owner chooses to register a variable name from his domain. If the domain owner chooses to do so, IANA will do so on a First Come First Serve basis. To accommodate registrants who do not have their own domain, IANA will accept requests to register variables with the prefix "\${FCFS.ietf.org:" on a First Come First Served basis. Assignments on a First Come First Basis do not require Expert Review, unless the registrant also wants IANA to establish a registry for the values of the registered variable.

The registry is a list of assignments, each containing three fields.

1. The name of the variable. The name of this variable must start with a "\${" followed by a registered domain name, followed by ":", or it must start with "\${FCFS.ietf.org". The name must be no more than 64 UTF-8 characters long. The name must be unique.
2. For assignments made on Standards Action basis, the Standards Track RFC(s) that describe the variable. If the RFC(s) have not yet been published, the registrant will use RFCTBD1, RFCTBD2, etc. instead of an actual RFC number. Note that the RFCs do not have to be a part of an NFS minor version. For assignments made on a First Come First Serve basis, an explanation (consuming no more than 1024 bytes, or more if IANA permits) of the purpose of the variable. A reference to the explanation can be substituted.
3. The point of contact, including an email address. The point of contact can consume up to 256 bytes (or more if IANA permits). For assignments made on a Standards Action basis, the point of contact is always IESG.

## 29.6.1.1.1. Initial Registry

The initial registry is in Table 29.

Variable Name	RFC	Point of Contact
<code>\${ietf.org:CPU_ARCH}</code>	RFC 8881	IESG
<code>\${ietf.org:OS_TYPE}</code>	RFC 8881	IESG
<code>\${ietf.org:OS_VERSION}</code>	RFC 8881	IESG

Table 29: Initial List of Path Variables

IANA has created registries for the values of the variable names `${ietf.org:CPU_ARCH}` and `${ietf.org:OS_TYPE}`. See Sections 29.6.2 and 29.6.3.

For the values of the variable `${ietf.org:OS_VERSION}`, no registry is needed as the specifics of the values of the variable will vary with the value of `${ietf.org:OS_TYPE}`. Thus, values for `${ietf.org:OS_VERSION}` are on a Hierarchical Allocation basis and are of no concern to IANA.

## 29.6.1.1.2. Updating Registrations

The update of an assignment made on a Standards Action basis will require IESG Approval on the advice of a Designated Expert.

The registrant can always update the point of contact of an assignment made on a First Come First Serve basis. Any other update will require Expert Review.

29.6.2. Values for the `${ietf.org:CPU_ARCH}` Variable

IANA created a registry called the "NFSv4 `${ietf.org:CPU_ARCH}` Value Registry".

Assignments to the registry are made on a First Come First Serve basis. The zero-length value of `${ietf.org:CPU_ARCH}` is Reserved. Values with a prefix of "PRIV" are designated for Private Use.

The registry is a list of assignments, each containing three fields.

1. A value of the `${ietf.org:CPU_ARCH}` variable. The value must be 1 to 32 UTF-8 characters long. The value must be unique.

2. An explanation (consuming no more than 1024 bytes, or more if IANA permits) of what CPU architecture the value denotes. A reference to the explanation can be substituted.
3. The point of contact, including an email address. The point of contact can consume up to 256 bytes (or more if IANA permits).

#### 29.6.2.1. Initial Registry

There is no initial registry.

#### 29.6.2.2. Updating Registrations

The registrant is free to update the assignment, i.e., change the explanation and/or point-of-contact fields.

#### 29.6.3. Values for the `${ietf.org:OS_TYPE}` Variable

IANA created a registry called the "NFSv4 `${ietf.org:OS_TYPE}` Value Registry".

Assignments to the registry are made on a First Come First Serve basis. The zero-length value of `${ietf.org:OS_TYPE}` is Reserved. Values with a prefix of "PRIV" are designated for Private Use.

The registry is a list of assignments, each containing three fields.

1. A value of the `${ietf.org:OS_TYPE}` variable. The value must be 1 to 32 UTF-8 characters long. The value must be unique.
2. An explanation (consuming no more than 1024 bytes, or more if IANA permits) of what CPU architecture the value denotes. A reference to the explanation can be substituted.
3. The point of contact, including an email address. The point of contact can consume up to 256 bytes (or more if IANA permits).

#### 29.6.3.1. Initial Registry

There is no initial registry.

#### 29.6.3.2. Updating Registrations

The registrant is free to update the assignment, i.e., change the explanation and/or point of contact fields.

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## Appendix A. Nature of the Changes Being Made for This Update

A large number of changes being made in this document are made to effect corrections to previous NFS Version 4.1 specifications. These include changes to address errata reports connected with those specifications, including some that were assigned the status REJECTED. In addition, similar changes are being made without explicit errata reports.

It is important to note that there are also a number of important organizational changes discussed below that will be made in this updated specification. As work on this document progresses, the status of those changes, together with other necessary changes, will be summarized in Appendix B.

\* The updated specification will depend on a number of NFSv4-wide documents, as described in Appendix A.1, rather than trying to deal with every aspect of the protocol description itself.

In the case of security, there will have to be decisions on how v4.1-specific security-related features will be addressed. See Appendix A.2 for details.

### A.1. Reliance on NFSv4-wide Documents

In many cases, matters previously described within the NFSv4.1 specification, will be addressed in separate NFSv4-wide documents.

- \* The process of protocol extension and creation of new minor versions is described in a separate NFSv4-wide document, [RFC8178], dealing with the issue for the NFSv4 protocols as a whole.
- \* Internationalization will be described in a separate document describing internationalization for all of the NFSv4 protocols, currently [I-D.ietf-nfsv4-internationalization]. The only v4.1-specific feature is the `fs_charset_cap` attribute, which is described in the current document rather than the internationalization document, although that document does discuss our choices in the matter.
- \* Security will also be described in a separate document applying to all minor versions. The handling is different because there is a wider range of security-relevant features within v4.1, despite the fact that the fundamental approach is the same for all minor versions. As a result, for some features, the security document will have the lead role while, for others, the v4.1 specification will be the main source of information about the feature, although the basic security functionality will be as defined by the NFSv4-wide security document.

### A.2. Adaptation of the NFSv4-wide Security Document to v4.1-specific Features

The v4.1-specific security-related features are dealt with as described below:

- \* Security issues regarding pNFS will be the responsibility of this v4.1 specification document. In doing this, we will rely, where we can, on the security facilities described in the NFSv4-wide security document.

This is necessary because some layout types will access data without using the RPC-based facilities that are discussed in the security document. In addition, the requirements for security-related coordination between data server and metadata server are best dealt with in this document, including cases in which RPC is used by both the data server and the metadata server, in which the necessary coordination requirements are defined by the layout type specification document.

- \* The description of the SECINFO\_NO\_NAME operation, will remain in the v4.1 specification, even though the description of SECINFO pseudo-flavors will be consolidated in the security. document.

This approach is necessary because the description of SECINFO pseudo-flavors needs to be augmented to allow negotiation of security-related transport characteristics in addition to auth-flavors, associated mechanisms and RPCSEC\_GSS services.

- \* The description of the MACH\_CRED and SSV features will remain in the v4.1 specification document and will only be mentioned in passing in the security document.

Instead, the focus with regard to state protection will be on client-peer authentication which applies to all minor versions.

### A.3. Changes to Effect Necessary Cleanup and Correction

The review of the existing specification text has discovered a set of areas that require clarification or correction, even though the problems had not been noticed as part of the pre-publication review of [RFC8881] and no errata reports have yet been filed.

In the following cases, it was necessary to make revisions to make the use of certain terms uniform throughout the document or to clarify the definitions which have come to disagree with the initial definitions.

- \* The treatment of the term "client owner" has been clarified to deal with the fact that previous specifications were inconsistent about whether the verifier was part of the client owner or added to it.

In this draft, a "client owner" always includes a verifier. When it is necessary to refer to the opaque string within it, the term "client owner id" is used.

These changes appear in Sections 2.5 and 5.5.

- \* The definition of "verifier" has needed to be substantially revised to reflect the fact that there are multiple verifiers within the protocol, each with its own use.

These changes appear in Section 2.5

- \* There has been a set of changes motivated by a need to clarify the circumstances under which delegation might be revoked.

This involved parallel changes in the description of leases where existing text was confusing because it was sometimes assumed that all locks were included rather than non-recallable ones, which obscures discussion of delegation/layout recall and revocation.

These changes appear in Sections 2.5 and 3.

A large set of changes were made to address issues within Section 7.6. These include:

- \* The requirement to wait forever for a response before reusing a slot has been modified to allow such waits to be terminated because of extraordinary circumstances such as termination of the task issuing the request.

That had to be changed because clients were unable to conform and because of the weakness of the proposed justification for the prohibition, i.e., that it resulted in a choice as to the next sequence value to be used.

The replacement makes clear why the sequence number is to be advanced, which is useful in reducing the probability of false retries.

- \* The prohibition on request retry was changed from a normative requirement to implementation guidance because it was clearly not a "fundamental requirement of the specification". Also the justification for a strict prohibition was undercut by work done in NFSv4.1 to implement Exactly-One Semantics, whose goal was to avoid negative consequences due to retries.

The replacement text clearly indicates why such retries are useless and best avoided, which is consistent with current practice. However it was necessary, in order to limit the occasions in which false retries could occur to use MUST NOT to forbid issuing of retries for abandoned requests once the slot had been used to send a later request.

- \* The discussion of false retries had to be extensively revised to make it clear that, while there were requirements to report certain false retries when detected, there were not corresponding requirements to check for this possibility. Instead situations in which such checking might be prudent are provided.

In the revised section, it is clear that false retries cannot occur if requests are never abandoned without a response and the protocol is implemented correctly. In addition, it is made clear how unlikely such false retries are if such request retries are constrained as required by the text related to valid reasons for request abandonment.

## Appendix B. Status of The Changes Being Made in this Update

Like all internet drafts, this document is a work in progress. In this particular case, that designation is particularly appropriate since there are specific changes that need to be made and either have not made or have been started but not completed. Information regarding changes made or to be made in this update is to be found in Appendices B.1 through B.4.

The current form in which the material is presented is designed for internal use within the working group, in order to help track the document's progress towards its goals.

Ultimately, the material regarding these changes will be re-organized in an eventual RFC.

### B.1. Changes Completed So Far in this Update

Work on the necessary changes discussed below has already been completed, although necessary review might not yet have occurred. At least for a while, changes made in later drafts of the working group document (i.e. those beyond -00) will not be reflected in this section and will be found within a subsection of Appendix C

The discussion of minor versioning has been updated to refer to [RFC8178], instead of the former approach which allowed each minor version to make its own versioning rules.

There are issues that have been addressed regarding how retried requests are to be terminated, including the fact the most common client handling of this situation violates a "MUST" in the existing specification. This took form of a revised Section 7.6.2, as discussed in Appendix D.2.1.

The document has been updated to eliminate the current (erroneous) treatment of internationalization, derived from earlier NFSv4.1 specifications [RFC5661], [RFC8881]. The section dealing with internationalization has been deleted, since it was never implemented. In its place, the specification has been modified to reference an external document which is to define the appropriate handling for internationalization for the NFSv4 protocols as a whole.

Currently, that document is the I-D draft-ietf-nfsv4-internationalization [I-D.ietf-nfsv4-internationalization]. In addition, the treatment of the `fs_charset_cap` attribute has been revised to reflect the analysis presented in the internationalization document.

Despite the completion of the internationalization work within this document, there remains work to do, most of which involves the completion of review for the NFSv4-wide internationalization document

- \* The new document was based on the treatment of internationalization within [RFC7530], which served as a useful starting point, since implementation of all NFSv4 minor version followed the same approach to internationalization issues. However the following issues still needed to be addressed:
- \* There was a need to update the treatment within RFC7530 to reflect IDNA changes made soon after the document was published.
- \* There was a need to deal better with client name-caching issues, especially in the context of case-insensitive file systems. Text has been written and submitted but review is needed.
- \* There was a need to address more fully the provision of representation-independent name lookup, which maps all canonically equivalent name strings in a directory to the same file.

However, these issues are being addressed in the context of the internationalization document, rather than the NFSv4.1 specification.

## B.2. Changes Made in this Update to Address NFSv4.1 Errata Reports

Work has been done to deal with errata reports associated previous NFSv4.1 specifications. These include a large set of errata reports associated with RFC5661 and a few associated with RFC8881. This work can be categorized as follows:

- \* The following errata reports associated with RFC5661 were dealt with in RFC8881, either because their substance related to issues to be dealt with in RFC8881 or because the simplicity of the needed change and its non-controversial nature made it simple to address the report as part of the RFC editing process for RFC8881: 2062, 2280, 2324, 2330, 2548, 3558, 5212.
- \* Work needed to be done to address many errata reports relevant to RFC 5661 that were not addressed in RFC8881, because the change was too large or too potentially too controversial to address in the context of RFC editing for RFC8881.



- \* There are a number of errata reports associated with RFC8881, that also had to be addressed.
- \* A few of the existing errata reports might have implications for RFC5662 and would need to be addressed by an eventual rfc5662bis RFC.

The errata reports that remain and that are being addressed in this document include reports currently assigned a range of statuses in the errata reporting system, including reports marked Verified, Hold for Document Update, and Rejected. These statuses are relevant to the processing of the associated errata but not in a way as direct as might be anticipated since errata reports marked Rejected might be addressed, as a result of a justified change in working group consensus.

- \* The following errata reports associated with RFC5661, have already been addressed in this document draft, in some cases by splitting out the associated change, if still necessary, into a related document: 2005, 2006, 2249, 2291, 2299, 2326, 2327, 2328, 2505, 2722, 3064, 3065, 3066, 3068, 3208, 3379, 3653, 3714, 3901, 4119, 4215, 4492, 4572, 4711, 4712, 4914, 5040, 5417, 5467, 5476, 6015, 6324.
- \* The following errata reports associated with RFC8881 already been incorporated into this document draft: 6308, 6337, 6865, 6611, 8705.
- \* The following errata reports associated with RFC5661 had not previously been addressed but will be resolved with publication of this draft as described in Appendix C.5: 2751, 3067, 4118, 5982.

The following issues need to be discussed with the errata report authors and the rest of the working group to enable the issues raised to be addressed in the resulting RFC:

- \* Errata reports 2751 and 3067 are related as both have to do with LAYOUTCOMMIT on the file layout type. As a result they are best discussed together.

The current status of 2751 is REJECTED which is justified given the scope of the proposed change. Nevertheless, it seems the working group needs to address this area, if not necessarily using the current proposed text for this report.

These reports need further working group discussion before the necessary changes are made in the document proper.

- \* Errata report 4118 has a current status of DEFER and for the most part appears unproblematic.

The proposed text uses the RFC2119-defined keyword "SHOULD" in a way that is not in accord with its definition and adds confusion to the proposal.

Once agreement is reached on the details of the replacement text, this issue should be easy to address.

- \* Errata report 5982 has current status of REJECTED. After further consideration of the issues, the proposer decided that the proposed replacement text addressed the wrong issue and so will be dropped.

Material related to the issues which this report was intended to address are being dealt with as described in Appendix D.2.1.

The only reports that still need to be addressed are 2751, 3067, and 4118.

### B.3. Changes Being Made Now in this Update

Work on the necessary changes discussed below has started but is not yet complete. This includes cases in which work to be completed is not within this document, but in a document referred to by this document. In such cases, matters formerly dealt within the NFSv4.1 specification, in the form of a single document, need to be addressed in a number of documents, each dealing with all NFSv4 minor versions together.

As noted previously, there are significant problems with the treatment of security within previous NFSv4.1 specifications [RFC3530], [RFC5661], and within other current NFSv4 specifications (e.g. [RFC7530], [RFC8881]). These are listed in Section 28.1. Work has started on these issues, although it is not as advanced as that for internationalization, since many important decisions need to be made. There is now a security I-D [I-D.dnoveck-nfsv4-security] which will serve as a guide to the decisions that will need to be made to guide the further work to arrive at a Proposed Standard discussing security for all the NFSv4 protocols, which rfc5661bis will refer to normatively.

Work is being done incorporate the material within [RFC8434], which establishes the requirements for parallel NFS (pNFS) layout types, which were not clearly specified in RFC 5661, leading to confusion.

There are a number of issues that are being addressed regarding the treatment of Directory Delegations. Although preliminary changes have been made, there might not be a clear consensus as to how existing issues should be addressed in this minor version. As a result further working group discussion will be needed. See Appendix D.2.2. for details.

Work has been done in Sections 2.7 and 5.3 to make the presentation more suitable to an environment in which RPC makes transport-level encryption and client-host authentication available. However, there is a need for some working group decisions to be made before completion of the transition to a security framework that fully embraces these new elements. In addition, the writing of a new Security Considerations section will require substantial progress on a standards-track security document for NFSv4 as a whole. Once that work is done, there will need to be a re-organization of those sections and their role will primarily be to refer to the standards-track security document.

In addition, work has been done to address security issues for NFSv4.1-specific features:

- \* Significant work has been done to clarify security implications of pNFS.

This work has primarily consisted of a major revision of Section 18.11 although there are significant updates to Sections 2.7 and 2.8.2.

It has been made clear that the only cases in which there are essentially no security consequences from the use of pNFS, are those in which RPC is used by the storage protocol, correcting text in previous specifications which gives a contrary impression.

The text has been revised to take account of the existence of services provided by rpc-tls including encryption and client host authentication.

There has been a re-organization of Section 18.11, including separate subsections dealing with non-RPC-based storage protocols and RPC-based storage protocols with either loose or tight coupling between storage server and metadata server.

- \* Significant work has been done to provide rpc-tls-based state protection which can be taken advantage even by clients who have not implemented SP4\_MACH\_CRED or SP4\_SSV or who are using AUTH\_SYS.

The Section 5.5.3 has been revised to allow, when SP4\_NONE is used, client host authentication to be used for state protection.

It is made clear in Section 7.8.3 that the use of SP4\_NONE, when host-client authentication is active, provides state protection against other clients rather than waiving state protection.

For many of the changes mentioned above, the definitive treatment will appear in the NFSv4-wide security document and there might also be a temporary references to the preliminary security I-D.

Further changes along these lines will most likely be necessary wherever in the document the SP4\_\* values are referred to.

There are a number of issues relating to the use of the key words defined in [RFC2119]. While the issues below could be treated individually and distributed among Sections B.1 through B.4, for now, we will treat them together.

- \* A shift has been made from only citing [RFC2119] to citing [RFC8174]. While it is sometimes said that, in the absence of RFC8174, "must" and "MUST" are to be considered synonymous, the working group has never interpreted RFC2119 in that way, although the clarification provided by RFC8174 was helpful. In light of this, it might be considered that all the necessary work has been done, apart from necessary review. However, given the working group discussion about this issue in connection with RFC8881, it appears that the working group will need to further discuss this issue soon after this document becomes a working group document. That would enable us to consider this aspect of the work complete.
- \* The use of the term "RECOMMENDED" to describe NFSv4 attributes which are not REQUIRED as been addressed by switching to the term "OPTIONAL", since "RECOMMENDED" is not in accord with [RFC2119]. This work is considered complete.
- \* There is ongoing work to deal with what appears to be a misclassification of protocol-defined attributes, making a number of attributes OPTIONAL, when the practical difficulties for clients in dealing with the absence of server support, makes this an inappropriate choice.

The security-related attributes owner, owner\_group, and mode have been made REQUIRED, both in this document and in [I-D.dnoveck-nfsv4-security].

The working group needs to review existing OPTIONAL attributes to see if similar changes need to be made for other attributes derived from NFSv3.

- \* For many of the previously existing uses of the terms "SHOULD" and "SHOULD NOT", it is not clear that the meaning is compatible with RFC2119. The difficulty is compounded by a long history of uncertainty as to the proper use of these terms, about which there may be disagreement within the working group. See Appendix D.1.1 for a detailed discussion of these issues.

It now appears that IESG is concerned about overuse of "SHOULD" and that this it could be left to justify such remaining use, while it seems that such justification does not exist.

For this reason the text now eliminate uses of "SHOULD" that cannot be clearly justified,

- \* There are some cases in which the terms "MUST" and "MUST NOT" have essentially been ignored by implementations or there are other reasons to believe that these terms may have been used inappropriately. See Appendix D.1.2 for a detailed discussion. Some work has been done toward addressing these issues but it is not complete, because further discussion is needed regarding changes to be made in Section 7.6.2.

These issues have been addressed by changes in Sections 5.7.1 and 5.7.2.

#### B.4. Changes That Will Need to be Made Later in this Update

Although this section is currently empty, one should be aware that

- \* There may be reasons to add new items discovered as the document progresses.
- \* There is still significant work listed in Appendix B.3.
- \* Much work already done still needs discussion and review.

#### Appendix C. Work Done in Various Drafts

The work in the subsections below cover changes made in various drafts of draft-ietf-nfsv4-rfc{5661,8881}bis and does not cover changes made in drafts of draft-dnoveck-nfsv4-rfc5661bis. As a result all such changes appear in Appendix C.1.

### C.1. Changes Made in 5661bis Draft -00

This draft made major organizational changes in the text inherited from [RFC8881] and started the work to clean up many of the troublesome issues discussed in Appendices D.1 and D.2.1.

The organizational changes included the following:

- \* Creating a new top-level section describing the reasons for this update.
- \* Moving most of the security-related material into its own NFSv4-wide document.
- \* Deleting the existing treatment of internationalization and referring the reader to the new NFSv4-wide internationalization document.
- \* Creating the initial versions of Appendices A, B, and D to track and explain changes needed and made.

### C.2. Changes Made in 5661bis Draft -01

Beyond limited editorial changes, this section lists the work done in draft -01.

The toc depth has been returned to the default value of three, with exclusions for subsections of operations and callbacks. The value of two left too many important third-level sections that did not appear the table of contents.

A large part of the changes consist of the changes described in in more detail in Appendix A.3.

- \* The re-organization of the description of client owner
- \* The revised explanation of the term "verifier".
- \* The enhanced description of delegation revocation.
- \* The set of changes made to address issues regarding EOS, retry and false retry, made to Section 7.6.

An interrelated set of changes were made in the pNFS area in order to clarify and re-organize the treatment of pNFS security, with some of it being the responsibility of the security document and to revise the material to clearly address the issues dealt with RFC8434 which is now obsoleted by this document. The following specific changes were made:

- \* Changes were made in terminology so that the general description of pNFS is consistent both with the files layout type and the layout types that appropriately deal with "storage devices."

The general term "data access protocol" refers both storage protocol and the use of file protocol to access a data server.

- \* The treatment of the term "control protocol" has been revised so as to avoid a contradiction between the presentation in RFC8881, which implied there was always a control protocol and RFC8435 which claimed it covered cases in which there was no control protocol.
- \* The section on pNFS security has been substantially revised for greater clarity and generality.
- \* A start has been made organizing the portions of the threat analysis that are to reside in rfc8881bis proper.

A new section for issues that need discussion was added as Appendix D.2.4. It deals with the following issues:

- \* A lack of clarity regarding possible persistence of lock state.
- \* A number of issues in the description of reply cache persistence that either make the feature difficult to implement or unnecessarily suggest that it needs to do things that are inherently difficult or impossible.

A new section (Appendix D.2.3) was added regarding issues raised by the discussion of memory-mapped files, now in Section 15.7. In addition, that section has been revised, in this draft to address the following issues:

- \* A neglect of the expected role of opening files for read, which would have caused delegation recall, rendering many of the issues worried about irrelevant.
- \* An unjustified expectation that, with mandatory locking in effect IO operations would result in obtaining locks making deadlock likely.

- \* An unjustified assumption that locking issues present in the last of mandatory locking apply as well in the case of advisory locking.

### C.3. Changes Made in 5661bis Draft -02

A number of changes related to the classification of attributes have been made:

- \* Attributed described as "Recommended", which previously has been described (incorrectly) as RECOMMENDED, were described as OPTIONAL, in accord with [RFC2119].
- \* The attributes Mode, Owner, Owner\_group, previously OPTIONAL, although referred to as "Recommended", have been made REQUIRED.

This change parallels similar changes in the NFSv4-wide security document [I-D.dnoveck-nfsv4-security].

### C.4. Changes Made in 5661bis Draft -03

A number of changes have been made to adapt to the splitting of ACL-related material from the security document and its presentation in a separate document devoted to ACLs [I-D.ietf-nfsv4-acls-update].

- \* Many reference to the security document have been updated to include the acls document as well [I-D.ietf-nfsv4-acls-update].
- \* Many reference to specific sections of the security document have been updated to reflect changes to that document.

Some of these have been modified to reference sections that are now in the acls document [I-D.ietf-nfsv4-acls-update].

In addition, a number of changes were made regarding the handling of various security-related attributes, introducing the topic with the understanding that the full treatment of the associated issues will be done within the NFSv4-wide security document [I-D.dnoveck-nfsv4-security].

- \* The paragraph regarding GETATTR and SETATTR of the name attribute directory has been rewritten to eliminate the dubious logic even though the protocol has not been changed, leaving a gap that still need to be addressed separately for POSIX authorization and ACLs.



- \* The section regarding the interpretation of owner and owner-group strings has been rewritten to introduce the possible choices, leaving the policy issues to the NFSv4-wide security document.

Further work was done on the issues discussed in Appendix D.2.1 including addressing issues originally intended to be dealt with a part of errata report 5982.

The description of the errata report status has been revised making it clear that only three reports still need to be addressed

We have created new per-draft Appendices C.1, C.2, and C.4 to keep track of specification changes.

There has been considerable work within Appendix D.2 including the following:

- \* Reorganization of the listing of the statuses of Appendix D.2 subsections according to the drafting of corresponding changes
- \* Considering Appendix D.2.1 Complete and ready for review.
- \* Creation of Appendices D.2.5 and D.2.6.

#### C.5. Changes Made in 5661bis Draft -04

Major revision have been made to Section 8 and Appendix D.2.4 in order to:

- \* update the description of persistence-related issues to reflect recent discoveries.
- \* provide an updated description of a persistence that has a good chance of being implemented and that could eliminate most grace period delays in the event of server restart.
- \* Clarify how a client becomes aware of persistence cross a server restart.

Considerable work was done in line with the suggestions in Appendix D.2.2. The following issues were addressed:

- \* The uncertainty about what assumptions could be made about the stability of cookie values and directory entry ordering by a directory delegation holder.

- \* The potential confusion about the effect of batching and delays of notifications needed to be addressed by making it clearer that these only applies to updates of attribute of file in the directory, rather than to the directory contents.

As a consequence of this work, the drafting associated with the various subsections of Appendix D.2 reached completion and it was necessary to revise Appendix D.2 proper to guide the necessary working group discussion of those changes.

There was considerable work clarifying the handling of CLAIM\_DELEG\_PREV. This includes:

- \* Distinguishing the special period allowed after client restart from the grace period used as part of server restart.
- \* Defining use of DELEG\_PREV claim types as reclaim-type operations while making it clear that they have no relation to the grace period or RECLAIM\_COMPLETE

Completed all necessary errata-based changes for this document by making the changes listed below:

- \* Errata report 2751, although rejected, has been incorporated into this draft.

The proposed changes were followed fairly closely, although the proposed new section has been moved from the pNFS chapter to the pNFS file chapter.

- \* Errata report 3067 has been incorporated into this draft.

The only divergence from the proposed text were the deletion of the word "deprecated" which was replaced by "are no longer used". Some incorrect uses of RFC2119-defined keywords were made lower-case.

- \* Errata report 4118 has been incorporated into this draft.

The only divergence from the proposed text were the deletion of the phrase "SHOULD be ignored" and its replacement by a statement that the field has no use. According to the author's reading of [RFC2119], "no harm, no 'SHOULD' applies here.

- \* Errata report 5982, which was rejected, has not been incorporated into this draft.

The change proposed, mentioning XID in connection with the false retry discussion turned out to be inadvisable and was dropped.

Despite that, other changes, discussed above, were made to satisfy the original motivation of this errata report, to make it clearer why extensive checking to detect false retry is not likely to be done, an issue which has needed to be addressed for a while.

Appendix D.2.7 was created to track discussion of errata that had been rejected.

#### C.6. Changes Made in 5661bis Draft -05

Changes to make it clearer when a change in the link count made by anyone other than the delegation holder cause a recall of a delegation. In particular, the case of changes done by NFSv4.0 clients needed clarification as did the rules for write delegations.

Added a discussion of the defects fixed as part of the rfc8881bis effort, focusing on the possibility of compatibility issues and the limited use of protocol extension, as provided for in Section 9 of [RFC8178]. This work included:

- \* Adding a summary of the work done to correct existing protocol defects in a new Section 1.4
- \* Extension of the notification enum to enable changes discussed below to enable implementation of directory delegations.
- \* Preliminary discussion of a new OPTIONAL attribute aclchoices to avoid having to wait for v4.2 to implement an aclfeatures attribute, as previously anticipated.

Made a series of changes to the discussion of directory delegations in Appendix D.2.2 and elsewhere. This work was prompted by suggestions from Rick Macklem and others.

- \* Additions to Appendix D.2.2 discussing reasons to provide a more flexible approach to the provision of position information within content update notifications.
- \* Adding of new Appendices D.2.2.2 and D.2.2.3 discussing how that flexibility might be provided
- \* Adding of a new Appendix D.2.2.5 together with corresponding work in Section 16.2

In addition to the above work within the Appendices, a number of changes were made to the specification proper to create new facilities to deal with issues discussed above and to incorporate Rick's suggestions to make the discussion directory delegations clearer. The changes included the following:

- \* Creating a new top-level section (now Section 16.2) explaining the directory delegation feature.
- \* A major revision of Section 25.39 to explain the use of the input and result notification bitmaps.
- \* A major rework/restructuring of Section 27.4 providing separate subsections for notification types.

Additional work was done in material first discussed in Appendix D.2.1 in order more clarity about the motivation for the changes and the connection to the probability of false retries. These changes were made in response to Olga Kornievskaja's review of version of this work appearing in the -04 rfc56661bis drafts. The changes include the following:

- \* Made a number of small changes to various subsection of Section 7.6. One of the most importance of these is connected to the use of retry without disconnection, formerly normatively prohibited but now a just a pointless, useless exercise. A normative prohibition was added about retries of requests that had been abandoned which was necessary to further limit the possibility of false retries.
- \* Major changes were made to Section 7.6.1.3.1 to make it clear that, while there were some requirements regarding the reporting of false retries that came to the replier's attention, there were no requirements regarding the work that replier's needed to do to make sure these would be found, if they occurred.

In addition, it was made clear why it was quite unlikely for these to occur, if the requirements laid out in Section 7.6 are followed.

- \* Description of the changes and their motivations were added to Appendix A.3.

#### C.7. Changes Made in 5661bis Draft -06

Significant additions were made to Section 1.3, adding new items to the existing list.

- \* In item 7, description of a new OPTIONAL attribute defining what extensions of the core UNIX ACL model are supported by the server.
- \* In item 8, discussion of a new ACE flag to support differences in handling of partial ACE satisfaction for draft-POSIX ACLs,
- \* In Item 9, discussion of a new ACE flag to support implementation of "default" ACLs as provided for by draft-POSIX ACLs.

#### C.8. Changes Made in 5661bis Draft -07

The following changes were made:

- \* A new item (#10) was added to the list in Section 1.3.  
  
This new item mentions the definition of GROUPNOTOWNER@ and OTHERS@ to be used to translate reverse-slope modes where DENY ACLs are not supported.
- \* Added definition of utf8pref to be used for components, since use of UTF8 is preferred but not required for these.

#### C.9. Changes Made in 5661bis Draft -08

The following changes were made in this draft:

- \* Moved Kerberos-specific material from 61bis to the security document.
- \* Moved SECINFO description back to this document.

#### C.10. Changes Made in 5661bis Draft -09

The following changes were made in this draft:

- \* Changes were made to clarify the role of the grace period as described in Section 13.4.2.

It appears that changes made in connection with the addition of RECLAIM\_COMPLETE seemed to make it necessary to delay all IO until completion of the grace period rather than the acquiring of new (i.e. not reclaimed) locks. This increased the negative effects of server recovery on clients, particularly if there were clients whose lock reclaim processing were delayed.

- \* Major additions were made to complete Section 28.2.

These included writing Sections 28.2.1, 28.2.2.3, 28.2.2.1, 28.2.2.2, and 28.2.2.4.

#### C.11. Changes Made in 5661bis Draft -10

The following changes were made in this draft:

- \* After-the-fact additions of "Changes made in Draft" section for -06, -07.
- \* Creation of a new section for RFC Editor Notes containing definitions of RFCs to replace the strings RFCTBD{10,20,21,22,30} in the final published document.

Use of these strings to simplify discussions of the relationships among documents that are part of the rfc5661bis effort.

- \* Create a new section 1.3 dealing with Compatibility Issues.
- \* Remove outdated material from the current Section 1.4 (previously 1.3).

List items 8-10 were removed because these had become unnecessary due to the change of approach to draft-POSIX ACLs. Instead of enhancing NFSv4 ACLs to support their semantics, it is now intended that support will be available via an NFSv4.2 extension.

- \* Work was done to clarify the description regarding returning NFS4ERR\_INVALID when inappropriate attributes are specified for (N)VERIFY.

#### C.12. Changes Made in 5661bis Draft -11

A number of changes were necessary to clarify/correct issues with the text in [RFC8881] in response to Working Group questions. These questions did not arise as part of the rfc8881bis effort, but suggest areas where the current specification is inadequate.

- \* Deal more thoroughly with situation regarding the interactions of delegation and REMOVE where the suggested recall can be safely dispensed with.
- \* Clarify/correct the semantics of PRESERVED\_UNLINKED to be POSIX-compatible, avoiding a gratuitous "MUST NOT".
- \* Restructure/revise the Implementation description sections for the REMOVE and RENAME operations.

This involved a reorganization of REMOVE to treat possible rejections first and a clear separation of the various REMOVE-induced file system changes.

RENAME was clarified to refer to REMOVE regarding the handling of files removed because they were renamed-over. This addresses issues regarding delegation non-recall and handling of deletion upon last close.

In addition, Appendix D.2.8 was created to provide helpful background for the review of the above changes.

#### C.13. Changes Made in 8881bis Draft -00

This document has been renamed as rfc8881bis since RFC8881 is obsoleted by it, despite the fact that the issues it addresses were introduced in RFC5661.

In addition, a number of small clarification/corrections were made in this draft.

- \* It needed to be made clearer how unsupported attributes are dealt with when requested by READDIR. Unfortunately, [RFC8881] says very little about the attribute fetch feature, leading to confusion about the case in which the set of supported attributes changes as result of crossing a mount point.
- \* In the discussion of pNFS, the text is not clear as to the client's and server's responsibilities with regard to shared state and did not provide normative guidance, substituting implementation advice in its place.
- \* Some work was done regarding pNFS terminology, following Tom Haynes' suggestions.
- \* A discussion of expected work with regard to [RFC8434] has been added to Appendix D.2.10

#### C.14. Changes Made in 8881bis Draft -01

The primary change made in this draft involves the creation of Section 19 as a replacement for [RFC8434] and the redesignation of this draft as obsoleting that RFC. Associated with these changes are the following:

- \* Corrections within Section 20 to make it a valid layout type specification according to Section 19.

- \* Changes to existing layout type specifications discussed in Section 19.2 and summarized in Section 19.2.5

Additional changes include the following:

- \* Updates within the Appendices B.1 through B.4 to reflect status changes made in recent drafts.

Quite important was moving [RFC8434]-related work and material related to directory delegations from Appendix B.4 to B.3.

Also important was moving the work regarding termination of RPC requests from Appendix B.4 to B.1.

- \* Appendix B.4 was rewritten to reflect the fact that it is, for now, empty.
- \* A lot of material relating to the revisions of Sections 18 through 18 was discussed in Appendix D.2.

Discussion of the issues raised in Appendices D.2.10 through D.2.13 is important to validate the changes made in this area.

- \* Some work as done to clarify and reorganize the description of directory delegations and associated notifications.

In additions, a few substantive changes were made regarding the presentation of attribute changes, with special handling for some child attributes that can never be changed and directory attributes inherently tied to associated content modifications.

To support the latter, the structure of a notify4 was explained as similar to a fattr4, and the document was explicit about combination of notifications that happen at the same time being presented within a single notify4.

#### C.15. Changes Made in 8881bis Draft -02

A number of changes were made to reflect a basic problem in the discussion of open-upgrade that dates back, at least to [RFC3010]:

- \* Significant revisions were made to Section 14.9 to correct confusion about IO authorization that would arise if open upgrade were done when the two principals doing the successive OPENS are different



- \* Changes were made to Section 20.14.1 to deal with the special challenges that these erroneous open-upgrades pose for the file layout type.

In addition, there were minor changes to correct omissions found when defining new layout types for an NFSv4.2 extension:

- \* Changes to Section 19.1.2 to include `lrf_body` among the nominally opaque fields that might have a layout-type-specific overlay
- \* Changes to Section 20.7 to explicitly indicate that there is no layout-type-specific overlay for `lrf_body` within the files layout type.

#### C.16. Changes Made in 8881bis Draft -03

A major focus number of changes were made in this draft related to the ACCESS operation. These were made necessary because of a failure to keep ACCESS fully aligned with the finer-grained authorization model introduced by the NFSv4 ACL model. The changes listed below include some which cannot be completed immediately for a number of reasons including the current uncertainty about the semantics of the ACE append bit, and the need to decide whether desirable additions are reasonably defined as "defects" which might be addressed in a bis document or need to wait to be added in an extension to a later minor version.

- \* Making it clear that for files, the semantics of EXTEND bit needs to match that of the ACE mask bit for append.

The need for this match is made clear in this document but the substance of those semantics will need to be addressed later in co-ordination with work on the ACL document and the other issues discussed in Appendix D.2.16.

- \* Clarifying the discussion of the addition of directory entries, taking into account the fact the NFSv4 ACL model has separate ACE mask bits controlling the additions of subdirectories and other file system objects.
- \* Possible access bits covering the authorization for the creation of (and possibly the scanning or modification of) the named attribute directory associated with a filesystem object.

This issues need to be coordinated with further work on the ACLs document, which needs change in the current description of related ACE bits appearing in the -00 draft of [I-D.ietf-nfsv4-acls-update].

In addition, the following changes were necessary:

- \* Updates of references to reflect the promotion and consequent renaming of the acl document.
- \* A number of clarifications were necessary in Sections 19.9.6 and 15.9.7 to clarify the needs that led to authorization support extensions for directory delegations and the possible alternatives to their use. The result was to restructure the material into Sections 16.2.6 through 16.2.9.

During the restructuring, it was discovered that some substantive changes were needed. These reflected the ability, previously discouraged, for the client to make authorization decisions on its own and the need to deal with AUDIT and ALARM ACEs that makes that impossible.

These changes resulted in flags to be added to the response to a request to create a directory delegation and significant functional additions to the handling of local GETATTR with respect to both authorization and caching.

#### C.17. Changes Made in 8881bis Draft -04

These include a number of related changes made to deal issues related, direct or indirectly, to the handling of large ACLs.

- \* Creation of a new Section 7.6.5 that discusses the constraints of suitable values of `ca_maxrequestsize`, `ca_maxresponsesize`, and `ca_maxresponsesize-cached`.
- \* A number of changes to make clear that while EOS is needed for modifying requests, caching of the reply is only needed for non-idempotent requests.

This previously was not as clear as it should have been/

- \* Associated additions to Acknowledgments.

In addition, text was added to address the issue discussed in errata report 8705, in which [RFC8881] identified the last legal operation code incorrectly.

## C.18. Changes Made in 8881bis Draft -05

Significant changes were made with regard to the discussion of corrections of protocol defects in bis documents. These concerned cases in which the size of the extension is important in deciding whether the extension is appropriately added to a bis document rather than being addressed in an extension to a late minor version

Work was done to make clear order-agnostic content notifications do not need to be sent to client responsible for initiating content changes.

A number of changes was made to further clarify Named Attributes, their previous uses, and discuss their possible connections with future work related to extended attributes:

- \* A major revision to Section 2.9.3 to clarify the relationship between named attributes and extended attributes.
- \* Substantial revisions to Section 29.2 to explain why the creation of this registry was ill-advised and to make it clearer why it has never been used and is unlikely to be used in the future.
- \* Creation of a new Appendix D.2.17 to deal with issues related to our future handling of extended attributes taking into account the desire for a unified approach to the problem for both Windows and POSIX, the difficulties of accommodating both in the same feature, and the need for compatibility with previous implementations using the ADS-based approach.

As our unfortunate experience with ACLs suggests, the need to accommodate both traditions must be approached with caution so that inconvenient semantic differences are not ignored, despite the important advantages of a common approach.

## C.19. Changes Made in 8881bis Draft -06

A number of changes were made related to attribute caching prompted by discussion of potential new caching-related advice, to be provided in later extensions:

- \* Restructuring of the section dealing with client-side caching into two sections, dealing with file and attribute caching (still Section 15 and with name and directory caching (a new Section 16.

This change was not made to correct a problem with [RFC8881]. It addresses the fact I should have done this when incorporating discussion of directory delegations and notification and that recent developments have made that poor choice insupportable going forward.

This division provides the opportunity, in Section 15, to discuss the important connections between file and attribute caching and their common interaction with file delegation and share reservations.

Similarly, this division provides, in Section 16, the opportunity to discuss directory caching and its differences from other forms of caching with regard to the issue of cache coherence.

- \* Creation of a new Section 15.3.3, acknowledging the previously ignored difficulties in sharing files, including those being written, without effective cache-coherence features.
- \* A major revision of Section 15.6 discussed below.
- \* Revisions to Section 25.30.3 to make clearer the necessary connection between changing size and opens for WRITE.
- \* Creation of new Appendices D.2.18, D.2.19, and D.2.20 to provide for necessary working group discussions of the suitability and adequacy of the changes made to the presentation of data and attribute caching, their relationship to ongoing work to address these issues, and the possibility of further work to address the issues that Tom Haynes has brought to the Working Group's attention.

Note that, in this context, the approach Tom has proposed, that of providing for the elimination of caching is considered one way of avoiding cache incoherence and is therefore dealt with in Appendix D.2.19.

Work on the revised Section 15.6 has included the following:

- \* The reorganization of Section 15.6 into multiple subsections including Sections 15.6.1 through 15.6.3 dealing with, respectively, coherence between a change-requesting client and the server, coherence among clients, and coherence between the requesting client and remote applications.

- \* Effecting a distancing from previous ways of treating write-behind caching, which tended, unfortunately to encourage its use. The new treatment stresses that it may be used while giving appropriate non-normative guidance about troubles it might cause in some environments.
- \* Revising the unrealistic desire to cache all attributes which was probably unrealistic when written and certainly unrealistic now.
- \* Converting the "rules" for sharing attributes to "guidelines" since these have never been understood well enough to form verifiable rules and we cannot promise that they work but know only they have been used successfully in the past.

#### C.20. Changes Made in 8881bis Draft -07

The following changes have been made:

- \* Continuation of the caching changes to deal more forthrightly with "close-to-open consistency".

These included creation of a new Section 15.3.2 together with minor changes to Sections 15.3.1 and 15.3.3

- \* Update of the errors tables to reflect new acl-related errors introduced in REFTBD22.

- \* A thorough review of instances of existing uses of "SHOULD" and revision of those that are not consistent with the definition of this term in [RFC2119].

See Appendix D.1.1 for more detail regarding the need to address these issues together with discussion of the question of why these terms have been misused in the past.

These were all done in a single draft to simplify the review process by giving the working group the opportunity to focus on this single issue without needing to separately relitigate each of many needed changes.

- \* Clarification was added to the definition of the `fs_layout_type` attribute to make it clear that "supporting" means that a server needs to be able to serve as the metadata server for the file system.

For file-oriented layout types, many servers cooperate to support the layout type, making this clarification desirable.

## Appendix D. Issues Requiring Further Discussion

This Appendix discusses issues that the working group needs to discuss before making decisions regarding potentially necessary specification changes. Despite the need for working group decisions on certain policy matters, some of the specific examples cited have already been addressed by revised text within the draft specification proper.

### D.1. Appropriate Uses of RFC2119 Keywords

Although, as stated in Section 1.1, this document intends to use these keywords as described in RFC2119, there are a number of issues that have resulted due to uses of these keywords in RFC5661 and RFC8881 that may not be clearly in accord with these definitions, possibly requiring some corrective action, once the working group has reached a consensus regarding the appropriate path forward.

- \* Because of a lack of clarity within RFC2119, there is considerable uncertainty about appropriate situations in which to use "SHOULD" and "SHOULD NOT", resulting in a number of cases in which they are inappropriately used or in which it is unclear whether particular uses are appropriate.

The working group needs to discuss these examples (see Appendix D.1.1) so that these terms can be used consistently in this specification and within the boundaries established by RFC2119, even if those boundaries have some level of uncertainty surrounding them.

The use of the related term "RECOMMENDED" in connection with file attributes is not included in the above discussion, since it is already clearly understood that this use is incorrect.

- \* Although RFC2119 is appropriately clear, there are a number of cases in which uses of "MUST" and "MUST NOT" are problematic, since they are used in RFCs 5661 and 8881 in ways not in accord with their definition while the existence of clients and servers that ignore such statements gives one reason to doubt whether these are truly required for successful interoperation.

The working group needs to discuss these examples (see Appendix D.1.2) so that such uses are corrected and to reduce the probability of similar occurrences in the future.

- \* Even apart from the definitions of these keywords, there is the further statement in RFC2119 that these terms are to be used "sparingly". Given the size of the v4.1 specification, it is desirable that all contributors adopt a common approach to issues about where these terms are appropriately used.

The working group needs to discuss the issues described in Appendix D.1.3) so that the new specification has a consistent approach to these matters.

#### D.1.1. Appropriate Use of "SHOULD" and "SHOULD NOT"

RFC2119 defines "SHOULD" as follows, with the definition of "SHOULD NOT" paralleling it.

This word, or the adjective "RECOMMENDED", means that there may exist valid reasons in particular circumstances to ignore a particular item, but the full implications must be understood and carefully weighed before choosing a different course.

This definition makes it clear how "SHOULD" differs from "MUST" but the specific difference with "MAY", while these terms are clearly intended to be distinct, is left unclear. Since it would not normally be expected for the other peer to be able to judge the validity of the reasons chosen by the SHOULD-using peer (or even whether the full implications of the choice made have been understood and carefully weighed by the peer's implementer), the other peer is in the same position as it would have been if "MAY" had been used. It needs to be prepared for the "SHOULD" to be followed or not followed, as the SHOULD-directed peer chooses.

Although one gets the sense that not following "SHOULD" or "SHOULD NOT" is in some way disapproved of, since one often does not have specific "valid reasons" to either follow or not follow a "MAY". However, this leaves a great deal of uncertainty remaining as to when "SHOULD" is justified, especially given the indication within RFC2119 that these terms are to be used "sparingly".

Given the existing definition and the widespread use of this keyword in situations in which valid reasons to bypass the recommendation are hard to determine, recent drafts (See Appendix C.8) have eliminated this term except where valid reasons to bypass the recommendation can be determined:

- \* One valid reason to bypass an otherwise controlling recommendation is the need to accommodate previously valid uses which are now viewed to be mistaken. In such cases, reliance on previous specification is a valid reason to continue to do so, rather than

retroactively declaring an existing implementation non-compliant, while the harm resulting from bypassing the recommendations is made clear in explaining the need to eliminate the allowance previously made, except for existing implementations.

- \* For very large portion of previous uses of "SHOULD" it does not appear that such valid reasons exist and that "should" would have been more appropriate. It appears that the advice to use these terms "sparingly" was ignored and authors felt that "SHOULD" was to be used based on a feeling that it made the document more IETF-like, or by IESG member comments suggesting that a greater use of such terms was somehow desirable.

A number of uses of "SHOULD" show clear evidence that the definitions within [RFC2119] were ignored by authors unwilling to use the word "should". This includes "SHOULD try" and multiple occurrences of "SHOULD have" where there is no true recommendation.

- \* A common pattern was to appropriately use normative text in describing the server's obligation, while describing what the client needed to do to use the feature effectively using "SHOULD" with no clarity as to what might be reasons to bypass the recommendation.

Even though "should" would be appropriate in these cases and is normative as that word is generally used, within IETF such statements are often felt not to be helpful, as they are not sufficiently directive in that they can be characterized as advice, which is ignorable and thus not "normative". In such situations, "SHOULD" is used without clear reasons to bypass the recommendation, even though it is just as ignorable as "should" would be.

Given the difficulty of revising this long-existing terminology, it seems that use of "should" is the best available alternative to address these issues.

In light of the changes discussed above, the working group needs to be given an appropriate time to review them and resolve any disagreements regarding this approach to the terminology choices. After that, this item will cease to be a matter for active discussion with the motivation for the changes recorded elsewhere,



#### D.1.2. Uses of "MUST" and "MUST NOT" that are Problematic

While the definitions of "MUST" and "MUST NOT" are quite clear, there are still instances within the existing specifications in which it is not clear that particular uses are appropriate or in which common client and servers do not follow the offered direction while interoperating successfully.

Some interesting examples from RFC8881 [RFC8881] follow. Note that, unlike the case in Appendix D.1.1 which looked at each instance of the target terms in a given section of the document, here we only look at a subset of uses which appear, in some way, spurious or otherwise questionable.

- \* There are reasons to question the use of "MUST" in the following statement appearing in Section 5.7.1 of RFC8881:

Where an NFSv4.1 implementation supports operation over the IP network protocol, any transport used between NFS and IP MUST be among the IETF-approved congestion control transport protocols.

This statement would make invalid the use of NFSv4.1 using RPC-over-RDMA when the RDMA connection is implemented using RoCE while allowing it for Infiniband and iWARP.

Although the peer might depend on operating together with an implementation having adequate congestion control, there is no basis for requiring that specific protocols (i.e. SCTP and TCP) be used, particularly since RFC2119 indicates that these keywords not be used "to try to impose a particular method on implementers where the method is not required for interoperability".

Regardless of one's judgment of the propriety of using "MUST" in this context, the working group needs to discuss and decide, by consensus, how to address the issue of RoCE use in supporting NFSv4.1 using RPC-over-RDMA.

- \* There are reasons to question the use of "MUST NOT" in the following statement appearing in Section 5.7.2 of this document and the similar statement appearing in Section 7.6.2.

A requester MUST NOT retry a request unless the connection the request was sent over was lost before the reply received.

Given that the text states that this is to "reduce congestion", it is hard to see how the mandated behavior is an "absolute requirement of the specification."

The following statement appearing in Section 7.6.2, phased as implementation advice, provides a positive explanation of the motivation, without making the use of "MUST" or similar terms, resulting in a shift between a normative introduction and the implementation advice providing the underlying substance:

Note that it is not fatal for a requester to retry without a disconnect between the request and retry. However, the retry does consume resources, especially with RDMA, where each request, retry or not, consumes a credit. Retries for no reason, especially retries sent shortly after the previous attempt, are a poor use of network bandwidth and defeat the purpose of a transport's inherent congestion control system.

Not mentioned in this section but one possible motivation for such a restriction is the potential need to simplify the work discussed in Section 7.6.1.3 particularly the possible need of the server to checksum data to be written to detect false retry, possibly undercutting the performance benefits of RDMA, as discussed in Appendix D.2.1.

If the issues relating to limiting the work necessary to detect false retries is not an appropriate basis for this prohibition, it seems better to avoid a shift between a normative introduction and later implementation advice by saying something like the following:

Given that NFSv4.1 uses transport providing reliable delivery, there is little point retrying a request, except in cases in which the occurrence of a connection disconnect leaves the requester uncertain as to whether the initial request was successfully delivered. The session-based reply cache allows the replier to deal correctly with retries after reconnect whether the initial request was delivered and executed or not.

- \* The following statement, appearing in Section 5.7.2 RFC8881, leaves one uncertain about whether the use of "MUST NOT" is justified, since it gives no clear explanation of why the prohibited behavior is troublesome.

A replier MUST NOT silently drop a request, even if the request is a retry. (The silent drop behavior of RPCSEC\_GSS [4] does not apply because this behavior happens at the RPCSEC\_GSS layer, a lower layer in the request processing.) Instead, the replier SHOULD return an appropriate error (See Section 3.9.6.1), or it MAY disconnect the connection.

This uncertainty is exacerbated by the introduction which states, incorrectly, that this is "to reduce congestion" and that it is paired in a bulleted list with the previous statement using "MUST NOT" where its use is also problematic.

It should be considered whether the explanation would be clearer if the focus is on the responsibilities of the replier in the session model, rather than on one particular case of the replier ignoring those responsibilities. One possible approach:

The replier MUST attempt to obtain and send a reply each compound request received. This applies with equal force to the case in which the request is a retry, with the instructions in Section 7.6.1.3 followed in generating the reply which the replier needs to send to the requester in all cases, except where a disconnection event makes this impossible

- \* The following statement, appearing in Section 5.7.2 of RFC8881, requires further analysis since the justification provided for the prohibition merely cites a possible difficulty, without consideration of whether this difficulty could be resolved without this prohibition.

A requester MUST wait for a reply to a request before using the slot for another request. If it does not wait for a reply, then the requester does not know what sequence ID to use for the slot on its next request. For example, suppose a requester sends a request with sequence ID 1, and does not wait for the response. The next time it uses the slot, it sends the new request with sequence ID 2. If the replier has not seen the request with sequence ID 1, then the replier is not expecting sequence ID 2, and rejects the requester's new request with NFS4ERR\_SEQ\_MISORDERED (as the result from SEQUENCE or CB\_SEQUENCE).

Beyond the problem with the justification provided, is the fact, that many clients, including those most commonly used, essentially ignore the "MUST NOT", yet successfully interoperate with most servers. This essentially makes the "MUST NOT" untenable.

There are two problems with the current justification:

- Only a single value is assumed as the sequence value to be chosen for the next request (i.e. two) while the possibility of the alternative choice (i.e. one) is not addressed at all.
- The occurrence of an error is treated as disposing of the matter, without consideration of potential recovery approaches.

It appears likely that whichever value is used as the next sequence, the resulting error is not fatal, making use of "MUST NOT" inappropriate. Possible replacement approaches will be discussed in Appendix D.2.1 and explored in a modified Section 7.6

#### D.1.3. Issues Regarding Use of RFC2119 Keywords "Sparingly"

RFC2119 contains the following statement:

Imperatives of the type defined in this memo must be used with care and sparingly. In particular, they MUST only be used where it is actually required for interoperation or to limit behavior which has potential for causing harm (e.g., limiting retransmissions) For example, they must not be used to try to impose a particular method on implementers where the method is not required for interoperability.

The following issues make this statement difficult to interpret.

- \* In fact, none of the terms defined in this RFC is an "imperative". They range among adjectives, participles, and modal auxiliaries, making it hard to determine which terms are being referred to.
- \* The terms "MUST", "MUST NOT", "SHALL", and "REQUIRED" might be thought of loosely as imperatives, since they are directing implementers to do something or to not do something.
- \* Although the terms "SHOULD", "SHOULD NOT", and "RECOMMENDED" do not have the sense of imperatives, they might be thought of as fundamentally carrying an imperative message, albeit one with a rather unclear provision for the recognition of exceptions.
- \* The terms "MAY" and "OPTIONAL", cannot reasonably be considered imperatives. Furthermore, the final sentences of the paragraph do not really make sense when applied to uses of these terms.
- \* Although the paragraph would normally be read assuming that the subject of the first sentence (i.e. "imperatives of the type defined in this memo") and "they" as used in the final two sentences, designate the same group of terms, that may not be possible since "MAY" and "OPTIONAL", do not make sense in the final sentences while it is hard to believe that the author really meant that these terms did not need to be used with care. The case of "sparingly" is not as clear cut but it is hard to conclude that only the first two classes of terms need to be used sparingly.

If these terms are to be used "sparingly", whether terms like "MAY" are included or not, a meaningful distinction must be made between things that are "an absolute requirement of the protocol" and the far more numerous set of things that simply describe how the protocol works. While it is required for interoperability that the client and server agree on the XDR for operations and results, and the actions to be performed for each operation, it is not clear how one could decide which of those interoperability requirements is "an absolute requirement of the protocol" meriting use the word "MUST", since deciding that they all do would not use these terms "sparingly" and is likely to result in an unreadable specification as well.

At times in the past there has been inconclusive working group discussion of the possible use the word "MUST" in connection with the need to return certain errors. While it was clear that the need for interoperability meant that this was a requirement within the definition of "MUST", there was concern about what that did to the style of explanation since returning errors, like setting appropriate operation parameters and results and performing the requested operations, are simply "the way the protocol works", and not an "absolute requirement of the specification" assuming those can be distinguished from ordinary requirements of implementing the protocol. The possible need to use such terms "sparingly" adds additional weight to this concern.

In any case, there seems to be a need for the working group to discuss and come to some consensus regarding the routine use of the word "MUST" even when the situation is not one which the question to be addressed is whether the definition of the word is adhered to, as discussed in Appendix D.1.2.

#### D.1.4. Going Forward Regarding Use of RFC2119 Keywords

Although many of the specific issue discussed here have been addressed, the working group needs further discussion in order to arrive at a consensus regarding policies to be followed on these issues in general.

#### D.2. Issues Regarding Proposed and Actual Changes

The subsections within this appendix each concern a set of changes that have been made to address various issues in the existing specification for NFSv4.1 [RFC8881] which are discussed in Appendices D.2.1 through D.2.8. Some, although not all of these, relate to matters raised in Appendix D.1.

Each of these require further working group discussion, although The nature of the discussion may vary, based on the nature of work to address the cited issues and the possibility that further related work might be required. The author assumes that, all these cases, leaving the material in the form it had in [RFC8881] would not be acceptable.

- \* For sub-sections in which the changes already been made in the current draft, if correct, fully address the issues now known, The necessary discussion will involve discussion of the proposed changes, in the form suggested in this draft in an attempt to move to a working group consensus on their correctness, adequacy, and clarity.

These sub-sections includes Appendices D.2.1, D.2.3, D.2.4, D.2.5, D.2.7 and D.2.8 changes have already been drafted and appear in the current document draft.

- \* For some other sub-sections, the author is unsure whether the changes made so far, even if correctly done, fully address the underlying issue. In these cases, the working group discussion of the proposed changes will need to be combined with a discussion of whether further changes are necessary and possible in the context of the current document.

These sub-sections includes Appendices D.2.2 and D.2.6,

#### D.2.1. Changes Regarding Request Aborts, Retries, and the Session Model

This issues discussed in this section have been dealt with by a set of changes to the document proper in Sections 7.6 and 7.6.1.3.1. These changes have been made in multiple drafts as described below and now need working group review to make sure that the changes made are adequate to deal with the problems in the existing text in [RFC8881]

- \* Previously, the question had been whether and how to deal with the misuse of RFC2119-defined keywords and the problem that the ability to terminate RPC requests, required by many client implementations, was inconsistent with exactly-once semantics, since zero times is not "exactly once".
- \* Now that there is an alternate approach available for consideration, the relevant question is whether that approach is adequate and how it might need to be changed. For more detail on the changes made, see Appendix A.3.

We deal here with three related issues that are connected, in some way, with the new session feature and the associated reply cache logic;

1. The potential need, as might be inferred from the discussion in Section 7.6.1.3 of [RFC8881], to checksum request data, particularly data to be written in order to eliminate the possibility of not actually acting on a request which is a "false retry", potentially resulting in data corruption.
2. The prohibition, discussed in Appendix D.1.2 on reissuing a request without the occurrence of a disconnection of the connection on which the request as issued.
3. The prohibition, discussed in Appendix D.1.2 on ceasing to wait for a response without actually receiving the response.

These issues have been addressed together since considering them in the context of the design of the sessions feature sheds light on the troublesome issues mentioned above. Specifically, we are looking at the possibility that we have arrived at a suitable framework to discuss these issues so that:

- \* This framework provides a more realistic and convincing explanation for any necessary prohibitions and/or recommendations.
- \* This framework allows such prohibitions to be safely downgraded to recommendations or implementation advice
- \* This framework might encourage client implementations to implement EOS without allowing the possibility of false retries, making it advisable for server implementations to avoid extensive recording of request contents or the checksumming of requests in order to prevent the undetected occurrence of false retries.

The issues and the changes made to address them in rfc5661bis drafts can be summarized as follows:

- \* The use of "MUST" in Section 7.6 and force clients to wait for responses for all requests needed to be adjusted because many clients were unable to comply when tasks were terminated, rendering the requirement useless.

In draft-01, the text was modified to make it clear that, realistically, there were situations in which the client could not wait forever and that real goal of the EOS logic was at-most-once semantics. In addition, the discussion now covers the possibility of getting SEQ\_MISORDERED in this situation.

- \* The use of "MUST NOT" in Section 5.7.2 to prohibit retransmissions needed to be revised since sending such retries while undesirable does not cause the kinds of harm that use of the RFC2119-defined term implies exist.

In draft-00, we eliminated the prohibition regarding retry and replaced it with implementation advice indicating why there is little reason for retries without clear motivation and explaining the unfortunate consequence of such retries.

- \* The text in Section 7.6.1.3.1 needed revision to make it clear when checks for false retries were either required or desirable and to make clearer how they could arise given the implementation of Exactly-once semantics.

In draft-01, added an initial paragraph indicated reasons that checks for false retry result from implementation problems and that there are practical limits as to when they will be done.

In draft-03, extensive changes were made to explain the situations in which false retry was a real concern and suggesting approaches to checking that can realistically be implemented. This includes addressing issues that were intended to be addressed by errata report 5982.

#### D.2.2. Issues Regarding Directory Delegation that Need to be Resolved

Directory delegations and notifications were added to NFSv4.1 but have never been implemented using the specification in [RFC8881]. During working group discussions of NFSv4 performance issues with regard to directory handling and later discussions with those working on implementations, it was discovered that there are a number of issues with regard to handling of directory delegations that need to be addressed, including cases where the design is adequate but substantial clarification is needed.:

- \* It was assumed that clients holding a delegation would maintain locally an image of the server directory, that needed to match that of the server with regard to directory entry order and the values of directory cookies.

As implementation efforts proceeded, it became apparent that this assumption was unduly limiting and needed to be addressed. In addition, this approach made it unduly difficult to use the feature for file systems in which sets of distinct characters are treated as equivalent (i.e those supporting normalization-related processing or case- insensitivity).



As a result, it was decided that allowance be made for clients using a different approach to caching, as described in Appendix D.2.2.2

- \* In addition, for clients that were prepared to maintain a local directory image, there were important gaps in the explanation that resulted in a lack of clarity regarding the ability of the notification scheme to allow the client image of the directory to be kept in sync with that of the server. It is not made as explicit as it might be that server support for continuation of directory delegations requires that the information provided in directory notifications is adequate to provide to the client the information needed to appropriately update the client's image of the directory to so that it can serve in place of a READDIR to the server. This includes the ability to maintain a directory ordering matching that on the server and READDIR cookies that match those held on the server. In situations in which the changes to the directory are of such a nature that this sort of update cannot be done based on a directory notification, the directory delegation needs to be recalled and returned. With the clarified, the value of directory delegations in avoiding the need to refetch large directories because of a small number of directory changes, would be more obvious. See Appendix D.2.2.1 for some suggestions in that regard.

While it is unusual for CREATES, RENAMES, and REMOVES to cause wholesale changes in the directory entry ordering or READDIR cookie values, there has previously been no way for the client to be sure that no such changes are being made, even when no other client is changing the directory. As a result, many clients are accustomed to refetch directories when they are changed, despite the consequent negative effect on performance.

As a result, it has seemed to many that there is little value in implementing directory delegations and notifications, leaving concerns about directory performance unaddressed.

- \* In addition, there is uncertainty as to whether a client making a change to a directory will receive timely notice of the details of the changes that will be made to the modified directory. The means by which notification is typically provided, using an asynchronous callback with provision for notification delay and batching of notifications was primarily directed to cases in which another client is making the modifications and the client receiving the notifications needs to sheltered from excessive notifications. This requires two issues to be addressed.

- \* It needs to be made clearer that the batching and delay of notifications apply only to the notifications of directory attribute changes and not to those notifying the client of changes in directory contents.

Although this might ultimately require a major rework of the text of Section 16.2, some useful suggestions can be found in Appendix D.2.2.4

- \* Since the notification model for changes in directory contents is an asynchronous one, it needs to be made clearer how clients making changes to directory contents can use these notifications can avoid refetching directory contents.

Some suggestions regarding useful change in this area can be found in Appendix D.2.2.1

- \* Although improvements have been made to deal with support for local equivalents of GETATTR using cached data provided using directory delegations, further work might need to be done as discussed in Appendix D.2.2.6.

Changes might be possible to respond to difficulties arising from the lack of prompt notifications of attribute changes for directory entries.

#### D.2.2.1. Clarifying the Role of Directory Delegations and Notifications in Avoiding the Need to Refetch Directory Contents

The first step is to clearly define the problem that content notifications address. This could be addressed by adding the following new paragraph to the end of Section 16.1.2:

Even when a client is certain that no other client is modifying a cached directory, there still might be a need to refetch the directory contents to satisfy additional READDIR requests. This is because there is no way to determine the modified ordering of the directory or the associated cookie entries using the knowledge of the directory changes requested and the corresponding responses. For a discussion of how directory delegation together with directory content updates can avoid the need to refetch directory contents, see Section 16.2.11.

In addition the following paragraphs need to be added at an appropriate place within Section 16.2.

When a directory delegation is held and notifications of changes of directory content updates are provided, the need to refetch the directory contents to satisfy READDIR requests can be avoided. This is of considerable benefit when the directories are large.

Although the directory content updates provided are asynchronous, they are not batched or delayed for considerable periods of time. Because of this, clients can keep track of the set of pending updates expected to avoid refetching directories when a content update is likely to enable the client to avoid the extra work.

#### D.2.2.2. Dealing with Various Client Caching Schemes

As a result of discussions with those involved in working on directory delegation implementations, it was discovered that:

- \* There are a significant set of clients that want to be aware of directory content changes but have no need for the position information currently provided since they either do not try to avoid repeated READDIRs by means of caching previous ones or synthesize READDIR results from cached contents without depending on the server's choice of directory entry order or directory entry cookies.

While the client is free to ignore position information provided, the effort to provide it where it is not needed might be a significant barrier to implementation.

- \* Some servers could produce useful position information with less difficulty if the directory cookies were defined so as to be acceptable to clients who do want to replicate the server's directory entry order and cookie values.
- \* Also relevant are those clients who do wish to provide cached READDIR results conforming to the server's order (not formally required but shuffling these might not be accepted by some users).

Addressing these issues requires providing the server more flexibility in the form of notifications used to inform clients of directory content changes while respecting client needs, which might be different for different clients. One way of doing so which has already been explored would allow three forms of content update notifications, such as are listed below:

- (A): Provision of position information in the way defined in [RFC8881].

- (B): Provision of position information in a simplified fashion using directory cookies only and avoiding the need to provide the names of nearby entries.

This approach is valuable for use in the large class of servers for which the entry cookie value is monotonically increasing as successive directory entries are transmitted as part of READDIR.

- (C): Omission of position information i content update notifications.

This approach is only useful to clients that do not try to maintain the server's directory order but are only interested maintaining the set of entries, independent of their order.

Selection of the proper format for any given update requires a new mechanism for the client server to arrive at the chosen format based on client needs and server capabilities. This mechanism is described in Appendix D.2.2.3. Because existing operations were specified without any provision for such selection, certain desirable options will only be available once a v4.2-based extension is available. However, as discussed below, there will be opportunities to avoid the Procrustean approach currently described in [RFC8881] in which only (A) is allowed.

In addition to the matters discussed above, the issues raised in Appendices D.2.2.5 D.2.2.7 need to be addressed.

#### D.2.2.3. Version Control and Extension

We need an mechanism that is acceptable in the v4.1 context to allow some limited extension of the approach to directory delegation specified in [RFC8881]. This needed to:

- \* Provide a way of including more substantial authorization support using additional notifications.
- \* Allow selection of the form of position information in content update notifications and to decide on the necessity of certain recalls based on the following factors:
  - Certain clients are concerned about the order of directory while other might not care.
  - There are clients concerned about entry order who want to have knowledge about server directory cookies, while there might be other that do not.

- There are servers that maintain directory cookies so they are always monotonically increasing with directory position while there others where that connection cannot be relied upon.
- There might be servers and clients written based on the approach of [RFC8881] rather than the more flexible one described in this document.

We need to deal with this possibility even though there are good reason to believe that no such implementations exist. Until [RFC8881] is obsoleted, which could take a while, we have no way of tracking ongoing development activities.

Dealing with all of the above, in a general way, requires a general extension mechanism, the overall structure of which will be discussed below, while the details will be decided in later extension document.

As part of this effort, we will need a way to provide greater flexibility, if possible, in an NFSv4.1 context, while interoperating correctly with client and server's using the [RFC8881] approach. This requires some way of distinguishing implementations without excessive XDR additions.

Regarding the selection of extension mechanism, it appears that the best approach is to generalize the use of the existing argument and result notification bitmaps. This requires only very limited XDR changes that follow the approach laid out in Section 9 of [RFC8178]. Bits can be defined without corresponding notifications and allowing other interface changes to be inferred from the presence or absence of particular bits.

For distinguishing new and old implementations, it seems the best approach is to rely on the notification bitmaps. Inclusion of an extension could signal client awareness of the extension with the appearance of the bit in the response signaling the server's knowledge of the extensions.

#### D.2.2.4. Clearly Separating Directory Content Updates from Other Notifications

Although an extensive set of small changes to clearly split content notifications from attribute notifications is probably necessary, to begin the clarification of this issue, Section 11.15 needs a clarifying final paragraph, to read as follows:

Note that these attributes apply only to directory attribute notifications and not to directory content updates, which, although asynchronous are not subject to batching or explicit delays.

Also important are the following replacement paragraphs for the first two paragraphs of the IMPLEMENTATION section of GET\_DIR\_DELEGATION, Section 25.39.4.

Directory delegations provide the benefit of improving cache consistency of namespace information. This is done through synchronous callbacks. A server must support synchronous callbacks in order to support directory delegations. In addition to that asynchronous notifications provide a way to provide a client a way to maintain an accurate client-side image of a changing directory (through client content updates) and to reduce network traffic as well (through both sorts of notifications)

Directory update notifications are specified in terms of potential changes to the directory. A client can ask to be notified of events by setting one or more bits in `gdda_notification_types`. The client can ask for notifications on addition of entries to a directory (by setting the `NOTIFY4_ADD_ENTRY` in `gdda_notification_types`), notifications on entry removal (`NOTIFY4_REMOVE_ENTRY`), and renames (`NOTIFY4_RENAME_ENTRY`). Cookie verifier changes, although not directory content updates, can be obtained by setting `NOTIFY4_CHANGE_COOKIE_VERIFIER` `gdda_notification_types` field. All of these notifications are asynchronous but they are not, like directory attribute notifications, subject to batching or time-based delays.

Directory attribute changes are requested using the notification type `NOTIFY4_CHANGE_DIR_ATTRIBUTES`, with the specific attributes that require notifications specified by `gddr_dir_attributes`. Like the child attribute notifications discussed below, these notifications are subject to matching and time-based delay in order to limit network traffic.

#### D.2.2.5. Directory Delegations and Permissions

It appears that Section 16.2 was written without sufficient attention to authorization issues that arise when LOOKUP, READDIR, GETATTR, and ACCESS operations are satisfied from cached data.

As a result, significant work will need to be done in related subsections to address that gap. This will inevitably, have to involve consideration of the increased difficulty of dealing with situations in the presence of ACLs. Given the current uncertain

state of ACLs, this will require, at least for NFSv4.1, steps prohibit or give permission to the server to prohibit use of directory delegations in situations in which their existence might compromise needed authorization restrictions.

#### D.2.2.6. Possible improvements in Support for Local GETATTR

Given that the protocol already has taken significant steps to deal with the following issues, the rest of this section will focus on the difficulties that result from the lack of prompt notification of directory entry attribute changes, the difficulties in providing such support and how limited server support for prompt notification might allow significantly more efficient client implementations.

- \* Handling AUDIT and ALARM ACE entries.

The server can scan the directory for such ACEs and inform the client of their non-existence using the CHANGE\_GA notification.

- \* Handling ACLs denying ACE4\_READ\_ATTRIBUTES (expected to be unusual).

The server can scan the directory for such ACLs and inform the client of their non-existence using the CHANGE\_GA notification.

The same scan, we could be done in the background, can address both issues.

Turning back to the possibility of prompt notification, we need to consider:

- \* The possibility that such notifications might slow down attribute changes making it a net performance loss despite the benefits of local GETATTR.
- \* The difficulty of implementing such notification for multiply-linked file given that filesystems that can find all directories referencing a file are quite rare with those that can do so quickly rarer still.

Given the above issues we should consider the possibility of specially managed notifications that are sometimes provided promptly, but have a useful fallback path allowing efficient local GETATTR in unusual situations. The following list is worth considering:

- \* In each case in which a requested attribute (e.g. change) can no longer be provided, the client is guaranteed prompt notification of that fact through use of existing facilities within a CHANGE\_GA notification.
- \* This guarantee is presumed negated whenever the directory entry in question is multiply linked or becomes so.

The existence of files with high link counts can be addressed in the same asynchronous scan used to find troublesome ACLs.

Transfer of a multiply-linked file into an existing directory for which a delegation is held, whether as the result of a cross-directory RENAME or a LINK can trigger the CHANGE\_GA since the target directories identified and the delegation-related information can be found from that.

- \* Where significant number of attribute changes, the server is free to cancel the promise of prompt notification.

Low-frequency asynchronous scans would be free to re-establish the promise using CHANGE\_GA.

#### D.2.2.7. Going Forward Regarding Directory Delegations

The material in this section should provide a suitable basis for working group discussion, in the hope that it will enable those changes to be moved into the specification proper in a later draft revision.

Once that work is done, the working group will be able to decide;

- \* Whether, with implementation of directory delegations, NFSv4.1 still has a directory performance problem that needs to be addressed
- \* Whether there is a need for extensions to improve directory delegations (synchronous notifications).
- \* Whether additional directory performance features are worth pursuing.

#### D.2.3. Changes Regarding Memory Mapping

It has been necessary to make major changes to material currently dealt with in Section 10.7 of RFC8881. The replacement is Section 15.7. Extensive changes have been necessary for the following reasons:



- \* The previous text consistently ignores the need for those reading and writing files to open them.

As a result, many the problems the previous section was concerned with, regarding a concurrently held write delegation only apply in the unusual case of files being read using special stateids.

- \* There had been assumption that CB\_GETATTR would always be used when attributes are interrogated, ignoring the possibility of the delegation being recalled.

This ignored the fact that CB\_GETATTR is an OPTIONAL feature and that there is no requirement for clients implementing to use it for access and modified time.

- \* In citing byte-range locking, there was no consideration of the fact that none of the cited issues poses any difficulty in the case if advisory byte-range lock.
- \* The treatment of mandatory byte-range locking assume, incorrectly that it requires as part a each IO, that a lock be obtained to enable that operation.

In fact, mandatory byte-range locking only requires that no inconsistent lock be held by another process performing IO.

As a result most of the potential issues cited do not exist for NFSv4.1 and if they did, they would apply to local IO as well, making this sort of locking untenable.

#### D.2.4. Issues Regarding Handling of Persistence

There are a number of important issues relating to persistence that need working group discussion and corresponding specification changes. These involve both reply cache persistence and the potential persistence of locking state to allow lock reclaim to be avoided.

The treatment of issues related to the persistence of protocol data shared by the client and server needs substantial remedial work, as described below. Without such remedial work, we would be stuck with a confusing description of a hypothetical feature that has never been implemented and has no prospect of being implemented, whose description is unusable due to confusion about the handling of locking state persistence. The issues to be addressed include the following;

- \* The excessive demands as to request atomicity and continuation across server restart, leading to a feature which cannot be implemented, as described in Appendix D.2.4.2.
- \* Ambiguity about the possibility of transparent state recovery and the means by which the client might be informed of its existence, as described in Appendix D.2.4.1.
- \* The confusion about the role of the clientid in connection with session recovery as described in Appendix D.2.4.3.

An alternative approach to these issues is presented in Appendix D.2.4.4. and will be the basis for a revised Section 8.

#### D.2.4.1. Ambiguity Regarding Locking State Persistence

As to the potential persistence of locking state, current specifications are unclear, mostly due to different handling of the issue in different sections and undue focus on what the server will provide with no attention to the question of how the client finds out about persistence or the lack thereof and deals appropriately with the situation.

First of all, the section entitled "Client Identifiers and Client Owners" (Section 2.4 in [RFC8881] and 5.5 in this document) gives the impression that, in the event of a server restart, the client will inevitably find out, by getting an NFS4ERR\_BAD\_CLIENTID error that locking state has been lost. While there is no explicit statement to this effect, the presentation of expected sequences of events (there are separate discussions of this for the cases of persistent and non-persistent sessions) leads one to suppose alternatives are not anticipated.

On the other hand, the section entitled "Loss of Session" (Section 2.10.13.1.4 in [RFC8881] and 7.13.1.4 in this document) strongly suggests that the NFS4ERR\_BAD\_CLIENTID is not inevitable, opening the way for locking state to be persisted across a server reboot, even though there is no explicit statement allowing servers to do so.

Given this divergence, it makes sense to determine which of these approaches is correct and make explicit descriptions of recovery make clear how clients are to deal with servers that do maintain state across reboot and avoid reclaim just as they do in the event of migration. A part of that discussion will concern potential compatibility issues which are not troublesome if clients do follow the approach laid out in the loss-of-session section

Adding to the existing confusion are occasional references to the possibility of certain forms of state persistence, with no discussion of how the client might find out about this potentially persistent state. For example, Section 25.43.3 contains the following paragraph:

If the metadata server is in a grace period, and does not persist layouts and device ID to device address mappings, then it MUST return NFS4ERR\_GRACE (See Section 13.4.2.1).

While this strongly implies that metadata servers could persist layouts across server failure, given the existing confusion it is hard to see how clients could effectively use this functionality or why server might provide it other than by providing a general lock persistence.

#### D.2.4.2. Implementability of Persistent Reply Cache as Currently Described

Another important topic of discussion concerns a number of statements in the section entitled "Persistence" (Section 2.10.6.5 in [RFC8881] and 8 in this document, which make implementation of persistent reply caches significantly harder than it needs to be or give the reader the impression that it is nearly unimplementable. This might have led to lack of implementation effort as part of a vicious spiral, that might result in the loss of this helpful feature, that needs implementation to take advantage the availability of lower-latency persistent storage. The following issues need to be addressed:

- \* One concern is that the statement "The execution of the sequence of operations (starting with SEQUENCE) and placement of its results in the persistent cache MUST be atomic" might convince the reader that the execution of each COMPOUND needs to be atomic as well, making conformance difficult and would seriously undercut any attempt to provide file system parallelism.

This might not have been the author's intention, even though it is the most natural reading of the sentence in question

There are necessary atomicity guarantees required but they have to be more limited and explicit to make implementation possible.

- \* Even more troubling is the issue raised in the statement "A server could fail and restart in the middle of a COMPOUND procedure that contains one or more non-idempotent or idempotent-but-modifying operations". The text goes on to say, as mildly as possible, "This creates an even higher challenge for atomic execution and placement of results in the reply cache.", but the indicating that

this is a greater challenge is likely to convince most reader that the feature is essentially unimplementable. The rest of the paragraph gives no reason to expect something workable except in special environments in which implementing this feature is the only goal.

Fortunately, the essential unimplementability derives, not from the feature but from the assumption that COMPOUNDS be executed atomically across a server restart rather than being terminated as part of the termination of the previous server instance, which this paragraph assume will never happen.

There needs to be a way to terminate CONPOUNDS still active at the time of server reboot, if there is no way to forbid execution of such troublesome COMPOUNDS.

- \* The final paragraph does nothing to correct this impression of unimplementability.

First, it says the following which would be unexceptionable for features for which there is at least one way to implement them: "While the description of the implementation for atomic execution of the request and caching of the reply is beyond the scope of this document".

Following this, it drives the final nail into the coffin of this feature by saying "An example implementation for NFSv2 [45] is described in [46]". The important fact here is that NFSv2 does not have COMPOUND, allowing the troublesome atomicity and cross-server-instance request continuity issues dealt with in the first two paragraphs to be bypassed, making this citation in this context inapposite.

#### D.2.4.3. Confusion about Clientid Role in Persistent Reply Cache

The existing discussion of reply cache persistence describes two possible variants that primarily differ as to the persistent storage of clientid-related information, with each variant deficient in some important way. This leads us to conclude that confusion about the role of the clientid and clientid-scoped state information and its persistent storage was not taken into account when this arrangement was arrived at. For example:

- \* While the persistent recording of some clientid-related information is presented as part of the second option given, there is no mention whatsoever of clientid-scoped locking state and its persistent storage.

Given the lack of explicit discussion of these matters, it is hard to tell whether it was expected that the server would or could persist this state.

- \* If clientid-related information is not saved (i.e. the first option), then the persistent reply cache does provide EOS across the server failure but does not allow the existing session to be used for new requests.

While this is a valid use case for clients worried about the possibility of EOS problems across server failures, the fact that there is no locking state persistence means that server failure will disrupt operation by requiring a grace period before, for example, opening a file. However, giving this limited benefit, it is troubling that the existing spec, due to confusion about clientid state, requires persistent recording of idempotent non-modifying operations (e.g. READs) with no real benefit because taking advantage of that benefit would require issuing new requests and checking their sequence ids against a persistent sorted sequence id.

#### D.2.4.4. Replacement Approach to Persistence in the Case of Server Failure

Regardless of the original intent with regard how to how these two aspects of data persistence were to be tied together, it seems that these need to be defined as independent features. Even if one could determine that these were intended to be tied together, which seems unlikely it is not possible to tie these two aspects of data persistence, each with its own scope, together at this point. Making that choice now would undercut the adaptation of NFSv4 to more available low-latency persistent storage in a number of ways:

- \* Since the amount of locking state is not bounded, there would need to be accommodations to the situation in which a surfeit of locking state makes use of persistent storage impossible. If these two sets of state were tied together, such situations would unnecessarily interfere with the a need to provide EOS semantics across server failure.
- \* In the context of clustered servers, the loci for the update of session-related and clientid-related data might be different, especially where clientid trunking is used. In such situations, there will inevitably be occasions where only one of these two forms of persistence is implemented.

- \* Given that the flow of locking operations is often a small part of the total and likely to be below the total of non-idempotent and modifying operations well, for many server implementors it would have a higher priority for implementation and use, As a result, it is reasonable to expect servers that implement it without implementing reply cache persistence even after the issues discussed in Appendix D.2.4.2 are successfully addressed

If persistence of locking state is to be made available as its own feature, allowing clientids to persist across server failure, then it is necessary to decide how to deal appropriately with the existing two options for session-based state persistence, once the issue of clientid-based state persistence is put aside.

- \* Persistence of the reply cache (only) will still be a viable useful option, not providing session continuation across server failure.

In defining this as a possible choice, it needs to be stressed that servers aiming to provide this functionality do not need to persistently store changes in sequence ids that would not be perceivable by a reconnecting client.

- \* Full session persistence would remain an option, even though there would be no need to persist data beyond the reply cache, current sequence id array and clientid.

Actual use of a persistent session would require persistence of the associated clientid-based locking state information. However, the server would not commit itself to maintain this information at session creation time and would find out if full session continuation was available, at the point at which the first new requests were issued

#### D.2.5. Changes in Attribute Categorization

The following changes in the categorization of attributes have been completed in 5661bis draft -03 but need to be further discussed by the working group. That discussion might involve other documents in addition to this one.

- \* The detailed description of authorization-related attributes has been moved to the documents [I-D.dnoveck-nfsv4-security] and [I-D.ietf-nfsv4-acls-update].

In line with this shift, the attribute categorizations have been made the province of [I-D.dnoveck-nfsv4-security] with that controlling in the event of any conflict.

- \* The attributes mode, owner, and owner\_group have been made REQUIRED rather than RECOMMENDED (with the meaning OPTIONAL)
- \* The attributes acl, sacl, and dacl have been described as "Experimental" in NFSv4.1, since, unlike other OPTIONAL attributes, the existing specifications do not describe the attribute sufficiently to allow interoperable client and server implementations to be developed.

#### D.2.6. Changes in Treatment of Attributes for Named Attribute Directories

Previous specification were self-contradictory in that:

- \* There were statements that made SETATTR and GETATTR on named attribute directories were undefined operations.  
  
The explanation offered for this exclusion did not make sense. For an explanation of why this text was eventually removed See Author Aside #66a in Section 5.3.5 of [I-D.dnoveck-nfsv4-security].
- \* There were other statements saying that named attribute directories could have attribute and stating that they were to include all the REQUIRED ones.

Changes have been made to eliminate this contradiction, as described below:

- \* The statement regarding SETATTR and GETATTR on named attribute directories being undefined operations was retained although the text "explaining" this exclusion was deleted.  
  
Since this material had been moved to [I-D.dnoveck-nfsv4-security], the replacement appears in Section 5.3,5 of that document.
- \* The statements about the ability to have (non-named) attributes for the named attribute directory in Section 11.7 have been deleted.

Discussion of these changes needs to continue to address the following issues:

- \* Even though the contradiction has been resolved, it is not certain why the exclusion is justified, given the inadequacy of the existing explanation.

Providing the ability to access and modify the attributes associated with named attribute directories might address some of the authorization issues discussed below, but could be expected to add additional complexity.

Leaving this as it is in the current set of documents would avoid additional complexity but still make it possible to reference named attribute directory attributes as part of dealing with authorization of operation involving the named attribute directory.

- \* There is no existing discussion of POSIX-based authentication of operations involving the named attribute directory, leaving a gap that needs to be filled.

Using the mode, owner and owner\_group attributes of the base object in place of those for the named attribute directory runs into troublesome issues since the X bit, controlling exec privileges for a (non-directory) base file controls lookup for the named attribute directory.

Any necessary changes will be made as part of Consensus Item #66 in [I-D.dnoveck-nfsv4-security].

- \* There exist ACE mask bits devoted to control of named attribute directories but it is clear that some changes need to be made.

As currently defined, these bits only control access to and creation of a named attribute directory, while allowing creation of new named attributes without authorization controls. Cleaning this up will be easier when we know how implementations behave but so far, none have been found.

Any necessary changes will be made as part of Consensus Item #100 in [I-D.ietf-nfsv4-acls-update].

#### D.2.7. Changes Made as a Result of REJECTED Errata Reports

Changes were made in response to the errata reports listed below, each of which was assigned a REJECTED status. The author, based on his own sense of the working group's need and wants, has addressed those errata reports. Even though there is no reason to suppose these reports do not need to be addressed, their rejection needs to be addressed by establishing that there is a working group consensus to make the change.



- \* Errata report 2722, reported by Ricardo Labiaga, was an editorial change that was rejected for reasons that are not clear. In addition, it is also unclear why this report was retracted.

In any case, the author addressed the troublesome area in way he feels satisfactory without necessarily following the text in the retracted report. The rejection, whatever its motivation implies we need a working group consensus as to the acceptability of the change made.

- \* Errata report 2751, reported by Ricardo Labiaga, was a technical change that was rejected because it proposed a substantive change in the handling of LAYOUTCOMMIT. Despite this justified rejection, it appears that implementation adopted the suggested approach, making it necessary that we, even at this late date, to adjust the specification so that implementation and specification no longer differ.

This specification draft has adopted the proposed changes without major changes except in one respect: The proposed new section, slated to be part of pNFS chapter will be done as part of the chapter devoted to the pNFS files layout. In addition, the presentation of changes in the form of text replacement complicated the process by requiring decompilation of the changes into xml. While the author did as well as he could, the complexity of the process calls for extra review.

In any case, further review of these change is necessary to make sure that the resultant text has working group consensus.

- \* Errata report 5982, reported by David Noveck was a technical change that was rejected. As things turned out the proposed text was misguided and was dropped.

Although other changes were made with same ultimate motivation and do need review, no special review is needed based on the rejection of the errata report.

#### D.2.8. Changes Made to Address Problems in Description of REMOVE/RENAME

In addition to the restructurings/clarifications discussed in Appendix C.12, there are a number of substantive changes for which a consensus needs to be arrived at. These include changes in which the change is arguably substantive because ambiguities in the existing text makes it hard to determine whether the existing text is or is not compatible with the new treatment.

- \* A change of approach to the PRSERVE\_UNLINKED flag, making it only apply to the OPEN which returned it. This avoids dealing with the possibility of it being returned differently for successive opens and makes it clear why it does not apply to NFSv3 or NFSv4.0 OPENS.
- \* Providing rules allowing the recall of delegations help by a client doing a REMOVE can be dispensed with, depending on the nature of the server's restrictions regarding REMOVE of open files.

This includes a requirement for the client holding the delegation and doing a REMOVE to make the server aware of any OPENS denying READ or WRITE.

There are also a set of potential changes that might be made once it is clear whether or not compatibility issues would prevent any substantive change.

- \* Making it explicit that the restrictions regarding removal of a renamed-over file are identical to those removed using REMOVE.
- \* Applying similar logic regarding the non-recall of delegations for OPENS of the file being removed.

It is important to note that the need for a write-delegation-holding client to make the server aware of OPENS denying write before doing a removal operation potentially raises compatibility issues. However, the practical problems arising are likely to be small since:

- \* A large portion of existing clients do not issue OPENS denying read or write.
- \* For those that do issue them, the complexity cost of doing so locally is likely to inhibit broad use of this technique. If a client were to do this locally, it would be taking on the work of the server in dealing with multiple processes doing such OPENS and corresponding CLOSEs, and detecting conflicts. This is unlikely to be worth doing since the benefit is likely to be quite small.

#### D.2.9. Changes Made to Address Lack of Clarity Regarding "The Forgetful Model"

Section 12.5.3 of [RFC8881] (and Section 18.7.3 of the previous draft this document) contains the following statement;

The client must fully process the operations before the "seqid" can be used. For LAYOUTGET results, if the client is not using the forgetful model (Section 12.5.5.1/>), it MUST first update its record of what ranges of the file's layout it has before using the seqid.

This statement raises the following issues/questions:

- \* There is no explanation of what "using the forgetful model" means and how it might differ from a presumably more persistent model.
- \* There is no explanation of how a client might choose a particular model and how a server might be affected by the client's choice of model.
- \* The sentence fragment "if the client is not using the forgetful model (Section 12.5.5.1), it MUST" is hard to understand since "MUST", according to [RFC2119] is used to introduce "a fundamental requirement of the specification". Given that definition, it is hard to understand how or why a client not following this requirement might be excused due to following this model.
- \* The section referenced in the troublesome text (i.e., Section 12.5.5.1), does not define "forgetful model"
- \* The title of the referenced section, referring to "Recall Robustness" seems to make its use where recalls are not involved, confusing.

In order to address the existing confusion and provide suitable normative text regarding the layout sequencing requirements and server and client requirements regarding obligations to retain layout information (independent of a presumed model choice), we are doing the following:

- \* Rename the sections so it is clear that it is about possible protocol requirements (and non-requirements).
- \* Re-organizing the section so that we are clear that what was previously referred to as an "assumption" which was not always the case is now explicitly a non-requirement.
- \* Put all the normative requirements together near the start of the section.
- \* Retain the useful implementation suggestion but make it clearer that this material is non-normative.

- \* Divide the previous ancillary section regarding recall robustness into two sections, one about retention requirements and another about recall/return interaction.

#### D.2.10. Need for Replacement of RFC8434

Discussion of recent drafts has made it clear that further work is needed to provide an up-to-date replacement for [RFC8434]. For that reason, draft-00 of rfc8881bis is not marked as obsoleting RFC8434, while it is intended to renew that designation in a later draft.

Overall, it seems best if most of the material in this RFC is placed in a new section after the pNFS Section and before the pNFS files Section.

The new placement, the passage of time and existence of new layout types leads to a number of issues that will need to be addressed when this new section is added.

- \* The use of RFC2119-defined keywords within [RFC8434] needs further discussion because these keywords are defined as applying to implementations rather than to possible future documents.

The Working Group faced a similar issue in the drafting of [RFC8178]. During discussion of that document some group members argued that each working group was free to specify things as it might choose and it was inappropriate for one Working Group to foreclose what a future working group might do.

As a result, we adopted the approach documented in Section 2.1 of [RFC8178], which can be adopted in a replacement for [RFC8434]

- \* The nature of this respecification means that we can no longer assume that the new section updates rfc8881bis, since it is part of rfc8881bis.

As a result there is no longer a meaningful distinction between Sections 3.1 and 3.2 of [RFC8434]

Although Section 4.2 and 4.3 still make sense, section 4.1 makes more sense in the pNFS files section. With regard to flexible files, it should only need a brief reference to [RFC8435].

Remaining issues that still need to be addressed are discussed Appendix D.2.11.

## D.2.11. Remaining Issues for Discussion After Replacement of RFC8434

The following work done as part of replacing [RFC8434] needs Working Group review:

- \* Making substantial revisions to Section 18 to take advantage of the more substantial role of layout types introduced by [RFC8434]

This enabled deletion of the confusing practice of referring to specific layout types in this section, which led to an unclear expansion path and an undue attention to unusual features of those layout types, with the placement of the file layout specification being misunderstood as applying to pNFS as a whole.

- \* Creating a new top-level Section 19.

While the material was derived from [RFC8434], its new placement (between Sections 18 20) clarified the relationship between individual layout type and pNFS as a whole.

- \* Revising and retitling Section 29.5.3 to focus on IANA-related requirements only.
- \* The addition of subsections 20.4, 20.26, and 20.27 to Section 20.

The following choices to change/restructure things as part of this replacement will need to be discussed:

- \* The change from using RFC2119-defined keywords to describe the obligations of layout type specifications to describing these, as [RFC8178] did, as normative despite the lack of these keywords.
- \* The reorganization of these requirements from one section (now Section 29.5.3 in [RFC5661] to two sections (19 and 29.5.3) in the bis document.
- \* There are good reasons to change the prohibition of the use of all-zero and all-one stateids that appears using the phrase "MUST NOT".

I can see no good reason for this prohibition, let alone any reason to consider this restriction "a fundamental requirement of the specification. On the other hand, it is difficult to change at this point given that this assumption may have affected implementations in ways that are difficult to change now.

\* The treatment of authorization-checking needs a major rework because of the following issues in the way this is dealt with in Section 19.9.2.3 of [RFC8881]:

- This treatment essentially ignores the access checking done by OPEN with troublesome negative effects on performance and requiring a functional divergence from the on-metadata-server approach that many versions of NFS have needed to deal with to prevent the authorization checking done by OPEN from being invalidated post facto.
- It treats ACLs differently from other authorization-related attributes, adding complexity and possible confusion..

It contradicts what should be the main focus of the requirements, making handling on data servers and the metadata server as nearly identical as possible.

The new treatment in Sections 20.15 and 20.14.2 takes a different approach:

- It limits the need for post-OPEN access checking to case in which the requester is different from the opener, as provided for by POSIX authorization semantics

Requiring this checking be redone on all IOs is a performance problem and functional problem if it applies to all IO and a troublesome difference if it applies only to data-server-directed IO.

- The needed sharing is not limited to ACLs, but applies to all changes of authorization-related attributes.
- Rather than assuming attribute propagation is always needed, the new treatment, integrated to discussion of sharing of locking information, allowing various form of propagation or remote interrogation as chosen by the implementer.
- Because of the previous uncertainty regarding authorization and the likely inability to suitably restrict the semantics of NFSv4 ACLs in the near future, the identity behavior is made the primary requirement, potentially overriding what is said in Section 20.15. See Section 20.14.2

- \* In a number of areas within existing NFSv4.1 specifications of pNFS, there are situations where explicit use of per-layout-type approach would be helpful but, because the text was written before the concepts behind RFC8434 were understood are dealt with in unsatisfactory ways:

- the differences are addressed by referring to the needs of specific layout types, presented as examples.

This results in uncertain handling regarding layout types not yet defined.

- The need for different sorts of handling is addressed as if it were implementation advice, even where there is no way that a particular layout type could do anything other than what is presented.
- The implementation is allowed to choose its handling (through use of "MAY") in circumstances in which this is inappropriate since the choice is better tied to the layout type used.

For the following areas, we have addressed additional case in which choices need to be made explicitly layout-type-specific. These are discussed in Appendices D.2.12 and D.2.13.

Aa a result of the need for extensive changes in these areas, Working Group review of the changes discussed above is necessary to make sure we have not inadvertently restricted needed implementation freedom or invalidated an implementations.

#### D.2.12. Layout-Type-Specificity for LAYOUTCOMMIT

The handling of attribute coordination/synchronization needed to be addressed by making more of the handling of LAYOUTCOMMIT layout-type-specific, as discussed below:

- \* As far as arriving at a global modified-time/change attribute, it is hard to arrive at a justification for the variation/complexity allowed in how the MDS's value is to be arrived at. One approach explicitly allowed the MDS to use the time of the LAYOUTCOMMIT as the new modified time and proceeding along similar lines for the change attribute. This avoids a lot of unnecessary complexity in the treatment within [RFC8881] regarding possible clock synchronization issues, discussion of possible estimates and how they might be used or "sanity-checked".

While a more radical simplification of this might be possible, we have made the matter layout-type-specific so that, if there is a justification for this complex treatment, it needs to be made in the context of a layout type that needs it.

- \* The change attribute is not explicitly deal with although a global value needs to be arrived at. The approach taken needs to be tied together with the layout-type-specific approach to setting of the modified time.
- \* As far arriving at a global size attribute at the MDS, the issues are similar and there seems no case for variation beyond making the choice layout-type-specific.
- \* Although there are indications within [RFC8881] that LAYOUTCOMMIT is not always required where one might expect it to be, the discussion is separated from other aspects of the description of LAYOUTCOMMIT as part of a general description of a layout type (files). Description of any situation in which the requirements are different is now a necessary as part of the layout type specifications for all layout types.
- \* Similarly, there had been within Section 18.9 material that suggested that LAYOUTCOMMIT would always be necessary, The section has now been rewritten to make it clear that this requirement is always layout-type-specific.

#### D.2.13. Layout-Type-Specificity for Layout Termination

The handling of layout recall, return revocation, and the discarding of layout-related information needs to be layout-type-specific for a number of reasons. In the new treatment, they are addressed together in Section 19.1.9 as discussed below:

- \* In the case of layout return, any restrictions other than requiring waiting for the completion of pending IOs, if they exist, need to be clearly explained in the layout type specification. These include the potential need to make the data storage device aware of the change in layout status.
- \* The case of layout recall is similar in that restrictions, if they exist, need to be described by the layout type specification. In cases in which recall needs to be done for some layout types (e.g. REMOVE), the requirements need to be clearly stated with a explanation as to how layout types that do not have such requirements prevent erroneous access to unused storage.



- \* While layout revocation is always possible due to an unacceptably delayed layout recall or a failure-like situation such as an expired lease, some layout types might allow it in other situations. These need to be described clearly.

Where these situations are described using the term "conflicting layouts", the meaning of this term needs to be explained clearly. This is particularly important because [RFC8881] defined this in a way that does not apply to the files layout when server multi-pathing was in effect. Also, it needs to be made clear whether non-layout-based IO operations can have the same effect.

Descriptions of how these situations are to be dealt with need to explain how IO operations using the revoked layouts are prevented and existing one drained. In this context, it needs to be clear that referring to "fencing" is stating a requirement which needs to be supplemented by an explanation of how the requirement is dealt with in all situations. For example, informing a data server of a layout termination needs to be supplemented by consideration of the situation in which the data server cannot be contacted.

- \* As discussed in Section 18.7.3.1, clients and metadata servers may discard layout information without worry about the inconsistency between the MDA and clients thereby created. However, there can be restrictions on such discarding based on the layout type. These need to be clearly described.

#### D.2.14. Possible Generalization Data Server Failure Handling

As discussed in Section 20.23.2, it is possible to generalize recovery to allow server implementations more freedom to maintain lock continuity across data server failures. This might become necessary if servers exist that manifest these behaviors.

#### D.2.15. Discussion Needed re Status of RFC5663

In light of the discussion in Section 19.2.1, there are reasons to either update or obsolete [RFC5663]. These problems arise because we now have a situation in which we have two different block layout types with corresponding specification documents:

- \* [RFC5663] was published with some troublesome problems in its handling of locking semantics and authorization. In addition, it indicated a need for fencing but remained unclear on how that need was to be dealt with.

- \* Later [RFC8154] was published to provide a clear description of how fencing was to be implemented. This document substantially improved the handling of locking semantics and authorization but these changes, unrelated to the fencing clarification which motivated the new document, were not backported to [RFC5663].

As a result, we now have two closely related documents describing block layout types. Given the effort to maintain these two, we need to decide:

- \* Whether we need to maintain a separate non-scsi block layout type.
- \* If we do want that, what is the best available way to provide for that?

Normally, to deal with such situations, we put the maintenance effort into the base specification and make the extension (in this case the SCSI handling of fencing\_ a small update but that option is no longer available to us.

The working group has the following options to address this situation and might have others.

- \* Update [RFC5663] in the same way that [RFC5664] is being update and mark it as updated by the bis document wen approved.

This leaves us with the job of maintaining these two for an indefinite period.

- \* Consider [RFC5663] as superseded by [RFC8154] and mark it as obsoleted by the bis document when approved.
- \* Write a statement within this document explaining support for the BLOCKS Layout type stating that is to follow [RFC8154] except that the scsi-related facilities are not necessarily available.

Mark [RFC5663] as obsoleted by the bis document when approved

#### D.2.16. Discussion of Possible Additional Access Flags

Discussion is needed to clarify the needs for further control bits for the ACCESS operation and the appropriate relationship between these bits and the corresponding ACE mask bits. Although there gas some additions to these bits to reflect the finer-grained permission model introduced by NFSv4 ACLs, we need to understand whether these need to be brought into closer alignment, and, if we do not do so, how to deal with the differences we leave.

In earlier specifications in which the client's was discouraged from making its own authorization decisions, it might have seemed essential for ACCESS to be made of equal or finer granularity than the ACE mask, this is no longer the case, leaving us to make a number of specific decisions listed below:

- \* Whether, given that we have separate ACE bits controlling adding of files and directories, ACCESS need to follow suite.
- \* Whether we need ACCESS bit corresponding to the ACE mask bits controlling access to and modification of the named attributes directory

If we choose not to do the above, the description of ACCESS needs to explain why not and how to determine authorization in these situations.

Another possibility is the definition of a new operation similar to ACCESS that returns an ACE mask word.

#### D.2.17. Discussion of Named Attributes and Their Possible Relationship with Extended Attributes

The following elements of the handling of extended attributes within NFSv4 so far makes it necessary for the Working Group to discuss and decide on an appropriate way forward for extended attributes:

- \* NFSv4.2 contains an XATTR extension oriented solely toward Linux extended attributes, which, while outside the scope of POSIX, are commonly available in POSIX-oriented clients and servers.
- \* The named attribute feature has been used to provide support for Windows extended attributes, despite the troubling performance issues, and the fact that named attribute functionality is far beyond what is necessary to support extended attributes.

This use presents the Working Group with an important compatibility issue that would need to be addressed by any new approach to support of Windows extended attributes.

- \* Although the significance of this fact is unclear, it appears that there has been, during the evolution of NFSv4, a "deprecation" of the use of the use of multiple data stream within Windows in favor of a model closer to that of Linux XATTRs although the exact differences remain unclear.

In any case, the working Group now has three possible choices to consider, identified below by "A" through "C". In analyzing these choices, we will need to refer to similar issues that arose in the handling of ACLs. While the desire to avoid a repeat of that unfortunate situation is part of our motivation, there are important differences between the two cases. Nevertheless, the similarities make it useful to use what we have learned from the ACL case in this analogous situation regarding extended attributes.

The three choice we have to select from are as follows:

A: Maintain the status quo in NFsv4.2 going forward.

Important advantages of this approach are that there is no work to do and no compatibility issue to address.

One important disadvantage is that we wind up with two different features to address extended attributes in Linux and Windows. Beyond the unsatisfactoriness of this approach from an architectural point of view, it seems that these two feature will contribute to the further separate evolution of Windows-oriented and Linux-oriented variants of NFsv4, undercutting the important value provided by unifying these in a single protocol.

While it might be reasonably assumed that this would make it difficult to transfer files including, between servers supporting the two distinct modes of extended attributes, this is only portly true. While Windows XATTRs would not be representable on server filesystems without support for multiple data streams (i.e., almost all Linux filesystems), the reverse is not the case since the Linux XATTR extension could be easily supported.

B: Create a separate XATTR extension oriented toward the extended attributes implemented in Windows systems, similar to wat has already been done in v4.2 for XATTRs in Linux systems.

While this approach shares the architectural issues discussed above for "A", the file transfer issue discussed there is fully addressed. The situation is similar to the handling of ACL transfers in the proposed draft-POSIX ACL extensions in which a server can support storage of both forms while making a separate choice as to which attribute(s) to use for authorization. The XATTR case is simpler because the filesystem does not actually use the extended attributes.

C: Extend the XATTR feature described in [RFC8276] so that it supports Window Extended attributes as well.

This approach, while architecturally more satisfying, requires extra work to deal with compatibility for servers using the named attribute approach. Proposals to adopt this approach need to be evaluated to ensure that they deal satisfactorily with filesystems that have stored extended attributes in this form and clients built to store them in this form.

The most important potential disadvantage of this approach is our current uncertainty as to the semantic requirements for Windows extended attributes and the possibility, as occurred with ACLs that neither the Windows nor the POSIX-oriented approach is a struct subset of the other. Although it is easy and helpful to assume that this is the case, if it is not true in practice, we would be putting ourselves in a situation similar to that which occurred in the ACLs case, which we have no desire to repeat.

The decision as to the appropriate approach will occur with the analysis of proposed extensions. This is better than trying to select "A", "B", or "C" in a vacuum, and then look for a proposal in the selected class. Fortunately, we have time to make this choice, as we can continue with "A", until a B-style or C-style proposal is prepared, adopted, and published.

#### D.2.18. Discussion of Neglected Caching Needs

This section has two major functions:

- \* To provide a basis for a discussion of the adequacy of the changes made to the treatment of caching made within Section 15 as a way of contributing to addressing the troublesome issues that Tom Haynes' work with workloads including access to files with multiple writers as exposed.

This discussion will be limited to work that could be done in context of the NFSv4.1 respecification effort.

Given that anticipated extensions will probably make reference to the same issue WG discussion (and presumably assent) to these changes will provide preparation for discussion of the needed extensions even though the extensions will be published before completing the rfc8881bis effort.

- \* To introduce possible ways of dealing, in later minor versions, of a range of possible ways to address the difficulties in handling cases in which a file is accessed by multiple clients including at least one writer.

We need to address situations in which a cache coherence mechanism is needed but might not, in practical terms be implementable. Complicating the situation is the fact that the protocol's method of preventing unwanted inter-client interactions, share reservations is not implemented in many client environments and there is no easy way to change this.

We will consider ways of directly avoiding cache incoherence including by eliminating caching as discussed in Appendix D.2.19. Other alternatives in which we try to avoid difficult sharing that provides no value by targeted semantic changes, or provide coordination to deal with sharing more effectively will be explored in Appendix D.2.20.

As we begin this discussion, I need to make clear how the issues we need to deal with, and how that view differs from that of others who might well agree on the means to fix them, at least in part. I think it is important to be clear about differences as we work to address the pressing issues that have been identified for sharing in situations in which cache incoherence is a troubling problem, despite the usefulness of caching for many other workloads.

In my view, the existence of these problems, decades after the basic decisions were made as for NFSv4 points to a troubling pattern of neglect in which the need to deal appropriately, with shared access to files that are not read-only was simply ignored. This happened because the only facility available to control such sharing, share reservations, was not widely implemented in many important environments. This happened because the needed feature was outside the POSIX file access framework and no alternatives were provided or even thought about. It is now time to correct this situation, which will take considerable time. Luckily, Tom has provided an expeditious means of surviving this delay, even though it might not be the right long-term solution from an architectural point of view.

With this in mind we need to proceed as follows:

- \* Verify that we have done all that we can reasonably to address the problem as part of the respecification effort, even if it is of no practical value. However, even if that is the case, we need to be aware that there is a problem to be addressed, even if we have a likely partial solution that can be made available in the immediate NFSv4.2 time-frame.
- \* Discuss the strengths and weakness of possible mechanisms of caching elimination, as discussed in Appendix D.2.19.

These need to address the immediate needs for an expeditious solution together with the longer term needs for a finer-grained approach.

- \* Discuss other alternative semantics to address sharing issues as discussed in Appendix D.2.20.

#### D.2.19. Discussion Regarding Ways of Avoiding Cache Incoherence

As noted above, an important way of avoiding the problems associated with cache incoherency is simply to avoid caching when such problems are likely to arise. Because caching is so important in other environments, it is desirable to limit this avoidance to situations in which it is actually needed.

Given this situation the work already done in specifying the "uncacheable" attributes needs to be made realizable as quickly despite possible concerns, which I share, about its coarse-grained nature. As we note the progress that has been made so far, the important advance is the recognition, in later drafts, of the advisory nature of the attribute,, although the prominent use of the word "uncacheable" sometimes makes it hard to see this properly. Despite this progress, there are a number of remaining issues that need to be addressed before a more current version is submitted (and adopted if necessary) and we prepare to move toward a successful WGLC.

- \* Although there has been considerable discussion of the relationship between advising against caching of files and of directories, it seems that additional discussion is needed in this area.

Although the approach taken in discussion is that these are unrelated is probably correct, the draft needs to make this clearer and we need an explanation of what problems might cause the server (or one of the clients ?) to advise clients to avoid directory caching.

Although the author of these drafts is concerned about the possibilities of inconsistencies and believes that caching elimination for files and data needs to be done together, the lack of connection between the file-related caching advice and the directory-related caching advice seems to be inconsistent with that.

Apart from issues regarding the number of attributes or documents, it appears that there is no obvious motivation for elimination of directory caching, as there is for file and attribute caching.

If the justification for directory caching elimination is, as one suspect it is, access-based enumeration, the document needs to say so. In my view, however, it would be better to enhance NFS to support ABE in an NFSv4 extension rather than letting each server do this in its own way.

- \* Although the switch to a model of advising clients whether to cache has been successful, there are some unresolved issues left from the original model in which much of the guidance was stated with normative text or an approximation thereof.

Part of this is that the word "uncacheable" seems to imply a mandate rather than helpful advice. More important is a lack of clarity about the basis for this advice.

This problem is compounded by the specification of this attribute as writable since the server is in a better position to decide on such advice on the basis of sharing patterns.

- \* If, as might be reasonably expected, the advice is conditional on the sharing pattern, the issue of notification promptness needs to be addressed.
- \* The issue of advice granularity needs to be addressed. Although making this per-file is satisfactory, there are suggestions, given the description, that something more coarse-grained is likely to be implemented. If so, much of the description about the basis for the server's advice to avoid caching is hard to understand since the sharing patterns for each file are likely to be different.

Assuming we can quickly achieve a useful coarse-grained solution, we need to prepare for a more fine-grained approach that can, unlike the prototypes already done, be useable in all environments because it does not suppress caching when there is no need to do so

One way of providing a more fine-grained approach to the problem of inter-client file sharing involves the extension of file delegations in a manner similar to the way they are used in directory delegation. In this case, sharing would result in a notification of the new sharing pattern rather than revocation of the delegation. Such notifications would allow the client to avoid revocation while updating its handling of the open file to respond to current sharing conditions:

- \* When the file transitions to a state in which it is shared, with one or more writers, to cease data caching.



- \* When the file transitions to a state in which any sharing is read-only. to enable use of data and attribute caching.
- \* n When the file transitions to a state in which it is shared, with more than one writer, to apply any special coordination measures to deal with that situation, in addition to refraining from use of data and attribute caching.

Within such a framework, the initial sharing state could be provided as part of the OPEN response or as an initial pro forma notification of the initial sharing state.

#### D.2.20. Discussion Regarding Possible Semantic Changes to Address Cache Incoherence

It appears we have a range of issues that arise from the fact share reservations are the architecturally correct solution to such unneeded sharing but that we have nearly zero chance of having them adopted widely in a reasonable time frame. As a result, we will look at ways to provide facility to avoid unwanted use of incoherent caching that do not involve changes to the signature of the POSIX OPEN operation.

Given the current state of affairs, it appears that a large part of the cases involving troublesome inter-client interactions are inadvertent because client applications have no means to prevent them.

Given this troubling situation, the approaches discussed so far, which suppresses caching in order to prevent cache incoherence, appears to be the correct solution to a problem that would be unnecessary, if the sharing is actually not needed.

In this section we will look at alternatives to unwanted sharing that can be used as defaults if the client-side API cannot be adapted to prevent unneeded sharing. We will deal separately with:

- \* Undesirable sharing between a single writer and one or more readers.

It seems that a considerable portion of such interactions are inadvertent and only happen because share reservations are unavailable or not used.

While caching suppression is desirable here, it does not address all the problems that result from such inadvertent sharing. For detail regarding how to address this see the discussion below regarding the single-writer case.

- \* Undesirable sharing among multiple writers and possibly some readers as well.

Where the sharing is inadvertent, the same techniques used for the single-writer case are useful as well.

For other cases, we separately address the interacting writer and non-interacting writer cases below.

Although the desire to avoid changes to OPEN leads us to focus on providing default mechanism that could be selected by things like mount options, there will be cases in which needs for more extensive co-ordination can be addressed by API additions that do not modify OPEN. Because of the difficulty of changing OPEN, we will focus on function that could be invoked after completion of the OPEN operation.

There are a number of changes that are likely to be needed for all of the cases discussed below:

- \* Creation of a default-DENY flag to be used when no denial mode is specified.

This is particularly useful for operating environments in which the user is unable to specify a particular deny mode. However, it will also be useful if the default depends, as it should, on actual sharing including also the possibility of caching elimination or the change deferral mechanism mentioned below.

- \* Providing for a sort of delegation-like object such as that described in Appendix D.2.19 that provides sharing-mode updated without necessarily implying revocation.

- \* Interfaces to file system facilities to delay visibility of changes until close or other synchronizing event.

The single-writer-case can be addressed as follows by allowing the open to proceed (unlike the case with share reservations) and allowing the client to respond to the existence of sharing, whether it discovers the sharing, whether this happens as a result of the OPEN response, or via a later notification when the sharing starts. In either case, the client will have the following options with a client-based default selected if he does not make a choice;

- \* Elimination of caching to avoid reliance on caching that does not provide coherence.
- \* Delay of changes by the writer until there are no readers.

In this case existing readers will see the old copy, most likely stored persistently as a spares file with later modification of the existing file happening by modification of the files block mapping when all the readers are gone.

- \* The current default of continuing to cache. This is a poor default but it should be selectable by those who want it.

The case of multiple writers requires distinguishing between interacting and non-interacting writers. Where both are possible, use of sharing based defaults is inadequate since the application would need to signal its need for the less-used handling (probably the interacting-write case) explicitly. If this is done after the OPEN, when the existence of multiple-writer sharing is established, implementers can avoid the difficult work of providing difficult-to-justify extensions to the OPEN API.

Given the above:

- \* It appears that noninteracting writers can be dealt with by multiple instances of the single-writer approach although it is likely that a different default will be chosen for this case.

Such cases are likely to involve uses of distinct portions of the file by different clients, as would be likely for local files modified by multiple processes.

- \* The interacting-write case would require a range of co-ordination facilities that are not provided but could be added without needing changes in OPEN. The client should have the option of obtaining access to these coordination facilities whenever it is found out that the file is being shared by multiple writers.

These facilities are likely to require some form of mandatory byte-range locks adapted to provide inter-client coordination facilities to support multi-client write-sharing of open files.

The development of further extensions might be necessary if uses of this form becomes more common. Likely developments include the provisions of atomic block swaps.

Although NFSv4 currently has mandatory byte-range locking as an OPTIONAL feature but there are a number of potential weaknesses that will need to be resolved:

- \* The ability to select mandatory or advisory byte-range locking as opposed to the current server-chosen one-size-fits-all approach.

- \* More detailed discussion of deadlock detection and recovery therefrom.
- \* Some consideration of the relationship between byte-range locking and caching and how the guarantee provided by the former might allow routine use of the latter.

Writers should probable have the ability to atomically swap a set of blocks existing blocks and replacements. Compare and swap of such block sets should be considered as well. Although byte-ranges might seem more natural, it is probably not worth the additional implementation complexity.

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#### RFC Editor Notes

[RFC Editor: please remove this section prior to publishing this document as an RFC]

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