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Module-Lattice Key Exchange in SSH  
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## Abstract

This document defines pure post-quantum key exchange methods based on Module-lattice post-quantum key encapsulation schemes for use in the SSH Transport Layer Protocol.

## About This Document

This note is to be removed before publishing as an RFC.

The latest revision of this draft can be found at [https://alharrison.github.io/ssh\\_mlkem\\_draft/draft-harrison-mlkem-ssh.html](https://alharrison.github.io/ssh_mlkem_draft/draft-harrison-mlkem-ssh.html). Status information for this document may be found at <https://datatracker.ietf.org/doc/draft-harrison-sshm-mlkem/>.

Source for this draft and an issue tracker can be found at [https://github.com/alharrison/ssh\\_mlkem\\_draft](https://github.com/alharrison/ssh_mlkem_draft).

## Status of This Memo

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## 1. Introduction

Secure Shell (SSH) [RFC4251] is a secure remote login protocol. The key exchange protocol described in [RFC4253] supports an extensible set of methods. The security of traditional key exchange methods used in Secure Shell (SSH) [RFC4251] relies on the algorithms being too computationally complex to be broken. The development of quantum computers poses a threat to the complexity of these algorithms. Given sufficiently powerful quantum computers, these traditional algorithms would be vulnerable to attack. Additionally, the threat of "harvest-now-decrypt-later" attacks could create a risk in the current landscape before sufficiently powerful quantum computers are available. In this attack, the data would be collected and decrypted by these quantum computers at a later date.

This document addresses the problem by proposing the use of post-quantum key encapsulation mechanisms (KEMs) to extend the SSH [RFC4253] key exchange. [I-D.draft-ietf-sshm-mlkem-hybrid-kex] introduces ML-KEM in PQ/T Hybrid mode [draft-ietf-pquip-pqt-hybrid-terminology] which combines the shared secrets established by an ECDH and a ML-KEM key exchange. This document uses ML-KEM in a single-algorithm scheme without combining it with a traditional ECDH exchange.

In the context of the [NIST\_PQ], key exchange algorithms are formulated as key encapsulation mechanisms (KEMs), which consist of three algorithms:

- \* 'KeyGen() -> (pk, sk)': A probabilistic key generation algorithm, which generates a public key 'pk' and a secret key 'sk'.
- \* 'Encaps(pk) -> (ct, ss)': A probabilistic encapsulation algorithm, which takes as input a public key 'pk' and outputs a ciphertext 'ct' and shared secret 'ss'.
- \* 'Decaps(sk, ct) -> ss': A decapsulation algorithm, which takes as input a secret key 'sk' and ciphertext 'ct' and outputs a shared secret 'ss', or in some cases a distinguished error value.

The main security property for KEMs is indistinguishability under adaptive chosen ciphertext attack (IND-CCA2), which means that shared secret values should be indistinguishable from random strings even given the ability to have arbitrary ciphertexts decapsulated. IND-CCA2 corresponds to security against an active attacker, and the public key / secret key pair can be treated as a long-term key or reused. A weaker security notion is indistinguishability under chosen plaintext attack (IND-CPA), which means that the shared secret values should be indistinguishable from random strings given a copy of the public key. IND-CPA roughly corresponds to security against a passive attacker, and sometimes corresponds to one-time key exchange.

The post-quantum KEM discussed in this document is ML-KEM which is based on CRYSTALS-KYBER. [FIPS203] standardized the ML-KEM scheme in 2024 with three parameter variants, ML-KEM-512, ML-KEM-768, ML-KEM-1024. ML-KEM is a NIST approved mechanism that is believed to be secure against an attacker with a quantum computer.

## 2. Conventions and Definitions

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

### 2.1. ML-KEM Key Exchange Message Numbers

When using ML-KEM as the Key Exchange Method, the following private namespace message numbers are defined in this document: #define SSH\_MSG\_KEX\_KEM\_INIT 30 #define SSH\_MSG\_KEX\_KEM\_REPLY 31

## 3. Key Exchange Method: ML-KEM

The client sends SSH\_MSG\_KEX\_KEM\_INIT. With this, the client sends C\_INIT which is the ephemeral client ML-KEM public key, C\_PK. C\_PK represents the 'pk' output of the post-quantum KEM's 'KeyGen' at the client.

The server sends SSH\_MSG\_KEX\_KEM\_REPLY. With this, the server sends S\_REPLY which is the concatenation of S\_CT. S\_CT is the ciphertext 'ct' output of the 'Encaps' algorithm generated by the server which encapsulates a secret to the client public key C\_PK. Before producing S\_CT, the server MUST perform the encapsulation key checks defined in Section 6.2 of [FIPS203], and abort using a disconnect message (SSH\_MSG\_DISCONNECT) with a SSH\_DISCONNECT\_KEY\_EXCHANGE\_FAILED as the reason, if they fail.

C\_PK and S\_CT are used to establish the shared secret, K\_PQ. K\_PQ is the post-quantum shared secret decapsulated from S\_CT. Before decapsulating, the client MUST check if the ciphertext S\_CT length matches the selected ML-KEM variant. The client MUST abort using a disconnect message (SSH\_MSG\_DISCONNECT) with a SSH\_DISCONNECT\_KEY\_EXCHANGE\_FAILED as the reason if the S\_CT length does not match the ML-KEM variant or decapsulation fails for any other reason.

The derivation of encryption keys is done from the shared secret K\_PQ according to Section 7.2 in [RFC4253] with a modification on the exchange hash H. The hash H is the result of computing the HASH, where HASH is the hash algorithm specified in the named key exchange method name, over the concatenation of the following

string V\_C, client identification string (CR and LF excluded)  
string V\_S, server identification string (CR and LF excluded)  
string I\_C, payload of the client's SSH\_MSG\_KEXINIT  
string I\_S, payload of the server's SSH\_MSG\_KEXINIT  
string K\_S, server's public host key  
string C\_INIT, client message octet string  
string S\_REPLY, server message octet string  
string K\_PQ, SSH ML-KEM shared secret

### 3.1. ML-KEM Key Exchange Method Names

The ML-KEM key exchange method names defined in this document (to be used in SSH\_MSG\_KEXINIT [RFC4253]) are

mlkem512-sha256  
mlkem768-sha256  
mlkem1024-sha384

Below we define

#### 3.1.1. mlkem512-sha256

mlkem512-sha256 defines the ml-kem-512 C\_PK public key and ciphertext S\_CT from the client and server respectively which are encoded as octet strings. The K\_PQ shared secret is decapsulated from the ciphertext S\_CT using the client post-quantum KEM private key as defined in [FIPS203]. K\_PQ is encoded as mpint [RFC4251].

The HASH function used in this key exchange [RFC4253] is SHA-256 nist-sha2 [RFC6234]

#### 3.1.2. mlkem768-sha256

mlkem768-sha256 defines the ml-kem-768 C\_PK public key and ciphertext S\_CT from the client and server respectively which are encoded as octet strings. The K\_PQ shared secret is decapsulated from the ciphertext S\_CT using the client post-quantum KEM private key as defined in [FIPS203]. K\_PQ is encoded as mpint [RFC4251].

The HASH function used in this key exchange [RFC4253] is SHA-256 nist-sha2 [RFC6234]

### 3.1.3. mlkem1024-sha384

mlkem1024-sha384 defines the ml-kem-1024 C\_PK public key and ciphertext S\_CT from the client and server respectively which are encoded as octet strings. The K\_PQ shared secret is decapsulated from the ciphertext S\_CT using the client post-quantum KEM private key as defined in [FIPS203]. K\_PQ is encoded as mpint [RFC4251].

The HASH function used in this key exchange [RFC4253] is SHA-384 nist-sha2 [RFC6234]

## 4. Security Considerations

The security of ML-KEM is based on the presumed difficulty of solving the Module Learning With Errors (MLWE) problem, based on the computational problems in module lattices. [FIPS203]

## 5. IANA Considerations

IANA is requested to register new method names "mlkem512-sha256", "mlkem768-sha256", "mlkem1024-sha384", and to be registered by IANA in the "Key Exchange Method Names" registry for SSH [IANA-SSH] with a reference field to this RFC and the "OK to implement" field of "MAY".

## 6. References

### 6.1. Normative References

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Open Quantum Safe has an existing implementation of ML-KEM based key encapsulation methods in all three parameter variants. Their fork of OpenSSH (OQS-SSH) contains an implementation using these algorithms for SSH key exchange algorithms. The authors would like thank Open Quantum Safe for their example implementations of postquantum algorithms.

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